dented line and showing that Q can be proved on a following line.

Conditional proof: prove $P \Rightarrow Q$ by assuming P on an in- Indirect proof: prove P by assuming $\neg P$ on an indented line and showing that both Z and $\neg Z$ can be proved on following lines.

Equivalences.

$$A \wedge true \equiv A$$

$$A \Rightarrow B \quad \equiv \quad \neg A \vee B$$

$$A \wedge (B \vee C) \equiv (A \wedge B) \vee (A \wedge C)$$

$$A \wedge false \equiv false$$

 $A \wedge A \equiv A$

$$A \Rightarrow B \equiv \neg B \Rightarrow \neg A$$

$$A \lor (B \lor C) \equiv (A \lor B) \lor (A \lor C)$$
$$A \lor (B \land C) \equiv (A \lor B) \land (A \lor C)$$

$$A \wedge \neg A \equiv false$$

$$\neg (A \Rightarrow B) \equiv A \land \neg B$$

 $A \land (A \lor B) \equiv A$

$$\neg \neg A \equiv A$$

$$A \Rightarrow true = true$$

$$A \wedge (A \vee D) = A$$

$$A \lor A \equiv A$$

$$A\Rightarrow true \equiv true$$

$$A \lor (A \land B) \equiv A$$

$$A \lor \neg A \equiv true$$

$$A \Rightarrow false \equiv \neg A$$

$$A \wedge (\neg A \vee B) \equiv A \wedge B$$

$$A \lor true \equiv true$$

$$true \Rightarrow A \equiv A$$

 $false \Rightarrow A \equiv true$

$$A \lor (\neg A \land B) \equiv A \lor B$$

 $\neg (A \land B) \equiv \neg A \lor \neg B$

$$A \lor false \equiv A$$

$$A \Rightarrow A \equiv true$$

$$\neg (A \land B) \equiv \neg A \lor \neg B$$
$$\neg (A \lor B) \equiv \neg A \land \neg B$$

• Modus Ponens

$$\frac{A \Rightarrow B, A}{B}$$

$$\frac{A \wedge B}{A}$$

$$\frac{A \Rightarrow B, B \Rightarrow C}{A \Rightarrow C}$$

• Modus Tollens

$$\frac{A \Rightarrow B, \neg B}{\neg A}$$

• Addition

$$\frac{A}{A \vee B}$$

• Constructive dilemma

$$\frac{A \lor B, A \Rightarrow C, B \Rightarrow D}{C \lor D}$$

Conjunction

$$\frac{A,B}{A \wedge B}$$

• Disjunctive syllogism

$$\frac{A \vee B, \neg A}{B}$$

• Destructive dilemma

$$\frac{\neg C \lor \neg D, A \Rightarrow C, B \Rightarrow D}{\neg A \lor \neg B}$$

• Assignment Axiom

$$\{Q(x/t)\} \text{ x := t } \{Q\}$$

• If-then

$$\frac{\{P \wedge C\} \ \mathtt{S} \ \{Q\}, P \wedge \neg C \Rightarrow Q}{\{P\} \ \mathtt{if} \ \mathtt{C} \ \mathtt{then} \ \mathtt{S} \ \{Q\}}$$

• Composition

$$\frac{\{P\} \; \mathtt{S1} \; \{R\}, \{R\} \; \mathtt{S2} \; \{Q\}}{\{P\} \; \mathtt{S1}; \mathtt{S2} \; \{Q\}}$$

• If-then-else

$$\frac{\{P \wedge C\} \text{ S1 } \{Q\}, \{P \wedge \neg C\} \text{ S2 } \{Q\}}{\{P\} \text{ if C then S1 else S2 } \{Q\}}$$

Consequence

$$\begin{split} \frac{P \Rightarrow R, \{R\} \text{ S } \{Q\}}{\{P\} \text{ S } \{Q\}} \\ \frac{\{P\} \text{ S } \{T\}, T \Rightarrow Q}{\{P\} \text{ S } \{Q\}} \end{split}$$

• While

$$\frac{\{P \wedge C\} \; \mathrm{S} \; \{P\}}{\{P\} \; \mathrm{while} \; \mathrm{C} \; \mathrm{do} \; \mathrm{S} \; \{P \wedge \neg C\}}$$

• Await rule

$$\frac{\{P \wedge B\} \; \mathrm{S} \; \{Q\}}{\{P\} \; \langle \mathrm{await} \; \; (\mathrm{B}) \; \; \mathrm{S}; \rangle \; \{Q\}}$$

• Co rule

$$\frac{\{P_i\} \text{ Si } \{Q_i\} \text{ are interference free}}{\{P_1 \wedge \ldots \wedge P_n\} \text{ co S1; // }\ldots \text{ // Sn; oc } \{Q_1 \wedge \ldots \wedge Q_n\}}$$