

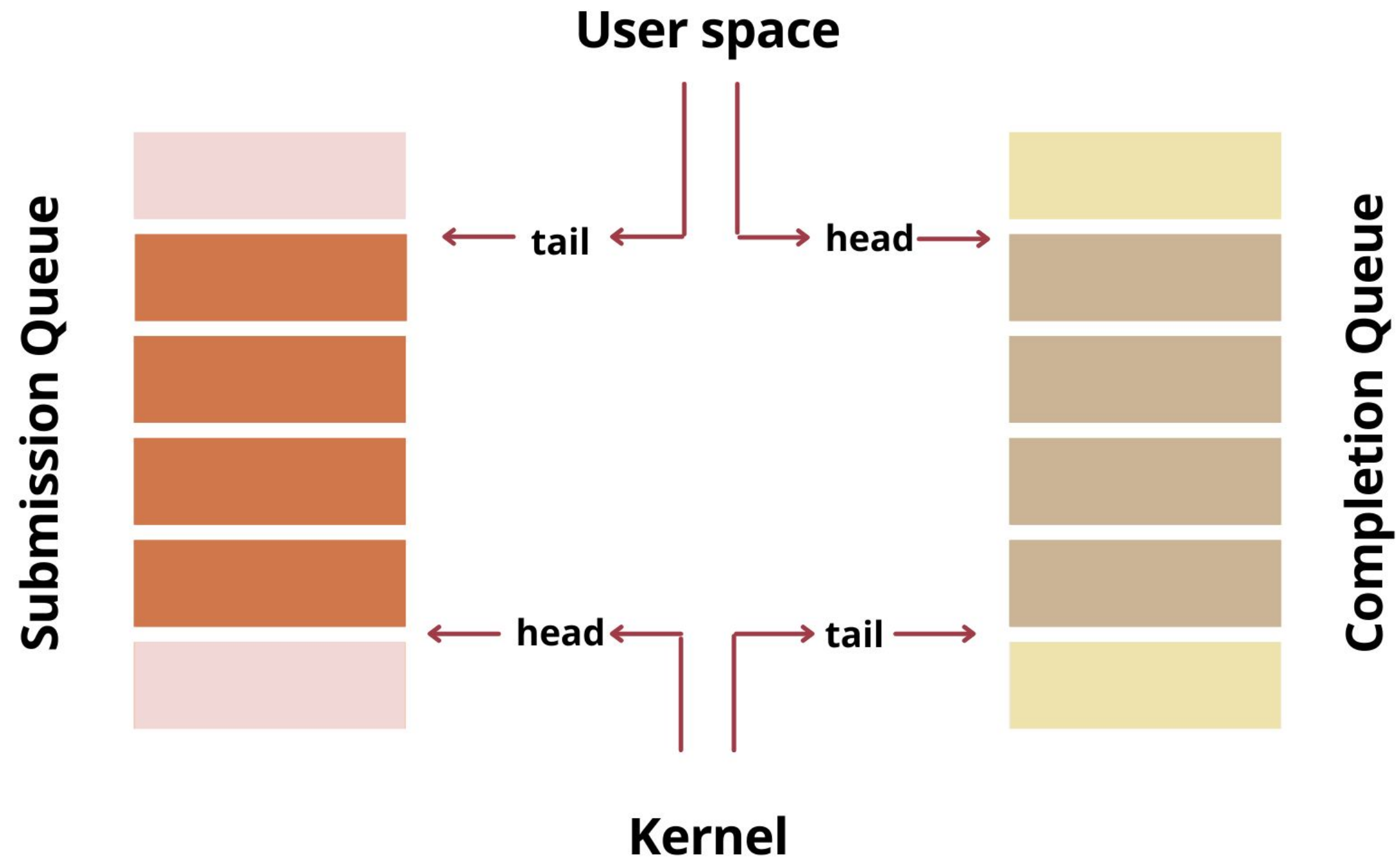
# io\_uring: BPF controlled I/O

Linux Plumbers 2021

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# io\_uring: introduction



# Lots of operations

```
enum {  
    IORING_OP_NOP,  
    IORING_OP_READV,  
    IORING_OP_WRITEV,  
    IORING_OP_FSYNC,  
    IORING_OP_READ_FIXED,  
    IORING_OP_WRITE_FIXED,  
    IORING_OP_POLL_ADD,  
    IORING_OP_POLL_REMOVE,  
    IORING_OP_SYNC_FILE_RANGE,  
    IORING_OP_SENDMSG,  
    IORING_OP_RECVMSG,  
    IORING_OP_TIMEOUT,  
    IORING_OP_TIMEOUT_REMOVE,  
    IORING_OP_ACCEPT,  
    IORING_OP_ASYNC_CANCEL,  
    IORING_OP_LINK_TIMEOUT,  
    IORING_OP_CONNECT,  
    IORING_OP_FALLOCATE,  
    ...  
};
```

```
...  
    IORING_OP_OPENAT,  
    IORING_OP_CLOSE,  
    IORING_OP_FILES_UPDATE,  
    IORING_OP_STATX,  
    IORING_OP_READ,  
    IORING_OP_WRITE,  
    IORING_OP_FADVISE,  
    IORING_OP_MADVISE,  
    IORING_OP_SEND,  
    IORING_OP_RECV,  
    IORING_OP_OPENAT2,  
    IORING_OP_EPOLL_CTL,  
    IORING_OP_SPLICE,  
    IORING_OP_PROVIDE_BUFFERS,  
    IORING_OP_REMOVE_BUFFERS,  
    IORING_OP_TEE,  
    IORING_OP_SHUTDOWN,  
    IORING_OP_RENAMEAT,  
    IORING_OP_UNLINKAT,  
    IORING_OP_MKDIRAT,  
    IORING_OP_SYMLINKAT,  
    IORING_OP_LINKAT,  
};
```

# Features

- **SQPOLL** for syscall-less submission
- **IOPOLL** for beating performance records
- Registered resources with fast updates
  - **IORING\_REGISTER\_FILES**: optimised file refcounting
  - **IORING\_REGISTER\_BUFFERS**: eliminates page refcounting, no page table walking, etc.
  - dynamic fast updates: no more full io\_uring quiesce
- **IOSQE\_IO\_LINK**: request links for execution ordering
- **IORING\_FEAT\_FAST\_POLL**: automatic poll fallback, no need for epoll
- **IO-WQ**: internal thread pool, when nothing else works
- multi-shot requests, e.g. poll generating multiple CQEs
- executors (**IO-WQ**, **SQPOLL**) sharing
- *and more ...*

# Execution flow

**First try nowait:** IOCB\_NOWAIT, LOOKUP\_CACHED, etc.

- might just complete, e.g. if data is already there
- O\_DIRECT goes async, -EIOCBQUEUED
- added to a waitqueue, e.g. poll requests

**Try async buffered read,** see FMODE\_BUF\_RASYNC

**Internally try polling** if supported, see IORING\_FEAT\_FAST\_POLL

- once fires, goto nowait attempts again

Any other way to go genuinely async; will be more in the future

Fall back to a **thread pool**, slower but often necessary

# Misconception debunking

io\_uring is not "just a worker pool"

- worker threads is a slower path

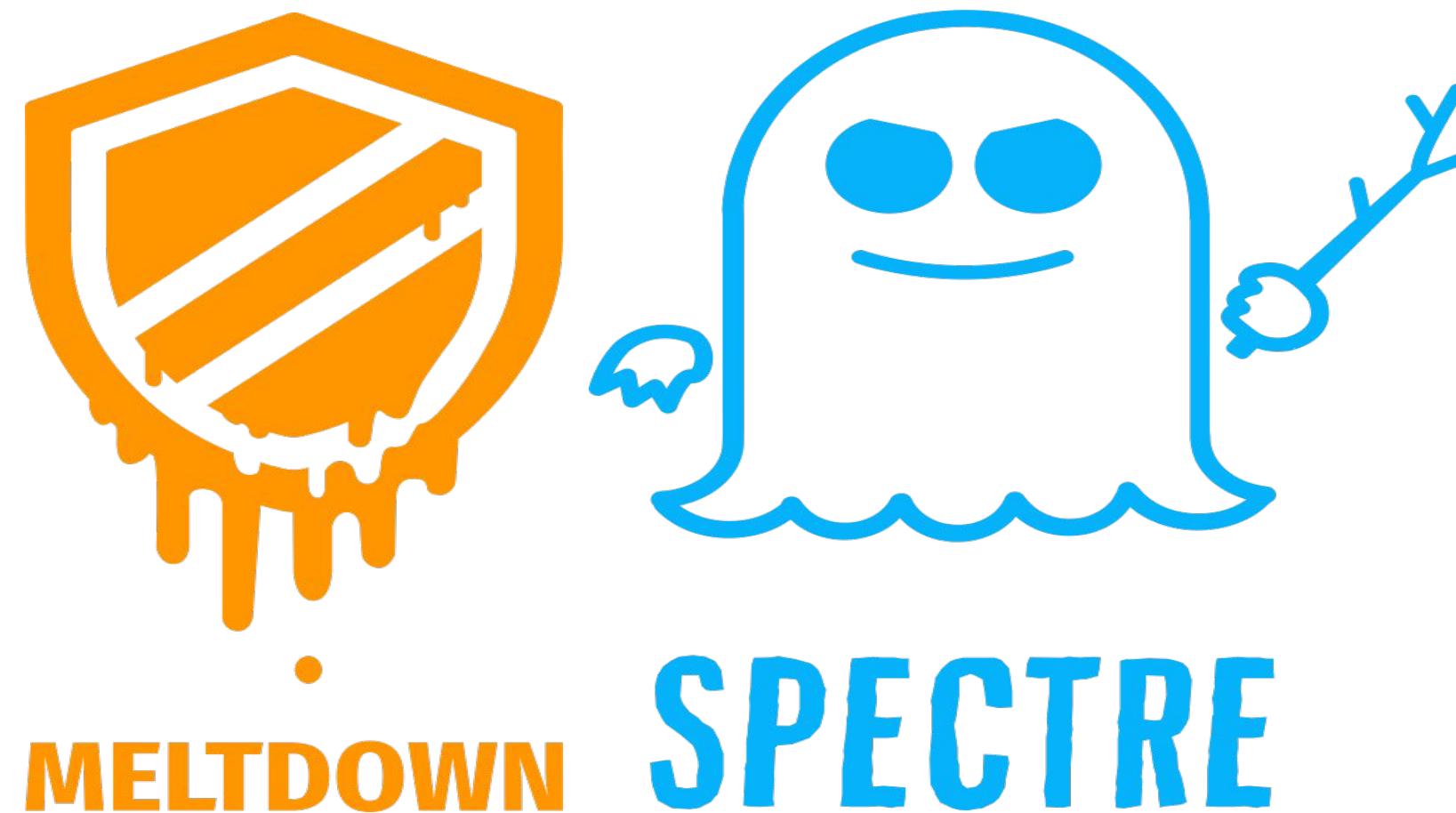
io\_uring is not I/O Completion Ports (IOCP)

- ... Microsoft is now developing a io\_uring for Windows

io\_uring is not only about syscall elimination/reduction

- provides asynchrony
- easy parallelism
- provides a state to base optimisations on, e.g. registered files

# The problem



By Natascha Eibl - <https://meltdownattack.com/>, CC0,  
<https://commons.wikimedia.org/w/index.php?curid=65233480>  
<https://commons.wikimedia.org/w/index.php?curid=65235937>

# syscall overhead

Vulnerability mitigations are **expensive**, and so are syscalls

- cost varies with CPU and enabled mitigations

Overhead for syscalls in a tight loop with little work can take **20-50%**  
(apparently, tested CPU is the worst case)

```
# copy by 4KB at a time
# cp_4kb ./file /dev/zero
```

```
29.47%  busybox  [kernel.vmlinux]  [k]  syscall_exit_to_user_mode
12.68%  busybox  [kernel.vmlinux]  [k]  entry_SYSCALL_64
12.49%  busybox  [kernel.vmlinux]  [k]  syscall_return_via_sysret
0.52%   busybox  [kernel.vmlinux]  [k]  do_syscall_64
...
```



```
# mitigations enabled
# nop requests, batch 32
# fio/t/io_uring -d32 -s32 -c32 -N1

16.41% io_uring [kernel.vmlinux] [k] io_submit_sqe
14.78% io_uring [kernel.vmlinux] [k] syscall_exit_to_user_mode
10.70% io_uring [kernel.vmlinux] [k] __io_submit_flush_completions
10.17% io_uring [kernel.vmlinux] [k] io_submit_sqes
9.78% io_uring [kernel.vmlinux] [k] io_issue_sqe
7.61% io_uring [kernel.vmlinux] [k] __io_queue_sqe
7.28% io_uring [kernel.vmlinux] [k] io_req_free_batch
5.07% io_uring [kernel.vmlinux] [k] entry_SYSCALL_64
4.79% io_uring [kernel.vmlinux] [k] syscall_return_via_sysret
4.29% io_uring io_uring [.] submitter_fn
2.75% io_uring [kernel.vmlinux] [k] io_alloc_req
...
```

```
# mitigations enabled
# Null block device, "realistic batching" 4 requests at a time
# modprobe null_blk no_sched=1 irqmode=1 completion_nsec=0 submit_queues=16
# fio/t/io_uring -d4 -s4 -c4 -p1 -B1 -F1 -b512 /dev/nullb0
```

<b>9.01%</b>	io_uring	[kernel.vmlinux]	[k] <b>syscall_exit_to_user_mode</b>
4.87%	io_uring	[kernel.vmlinux]	[k] blkdev_direct_IO
<b>3.27%</b>	io_uring	[kernel.vmlinux]	[k] <b>entry_SYSCALL_64</b>
<b>2.92%</b>	io_uring	[kernel.vmlinux]	[k] <b>syscall_return_via_sysret</b>
2.89%	io_uring	[kernel.vmlinux]	[k] kmem_cache_free
2.74%	io_uring	[null_blk]	[k] null_queue_rq
2.68%	io_uring	io_uring	[.] submitter_fn
2.31%	io_uring	[kernel.vmlinux]	[k] blkdev_bio_end_io
2.27%	io_uring	[kernel.vmlinux]	[k] io_issue_sqe
2.19%	io_uring	[kernel.vmlinux]	[k] io_do_iopoll
2.12%	io_uring	[kernel.vmlinux]	[k] kmem_cache_alloc
...			

Sweet spot for optimisation. How about **SQPOLL**?

- still needs userspace to process completions
- takes a CPU core; high CPU consumption
- cache bouncing

**BPF** is there to help! Can also help latency

# Requirements

**Flexibility:** what capabilities BPF has to have?

- submitting new requests
- accessing CQEs, multiple if needed
- poking into userspace memory

**Low overhead**

- Traditionally we've optimised batched submission more
- BPF is expected to have a lower batch ratio

## Idea 1: let's add a callback to each request and run it on completion

- needs hooks in generic paths, non-zero cost
- limits control over execution context
- can't do waiting and other async stuff
- BPF needs context, would need allocation
- looks horrible ...

```
struct io_uring_sqe {  
    ...  
    u32 callback_id;  
};  
  
int io_init_req(struct io_uring_sqe *sqe)  
{  
    if (sqe->callback_id)  
        req->bpf_cb = get_bpf(sqe->callback_id);  
    ...  
}  
  
void io_req_complete(struct io_kiocb *req, long res)  
{  
    if (req->callback)  
        req->bpf_cb(req, res);  
    ...  
}
```

## New io\_uring request type: **IORING\_OP\_BPF**

No extra per request overhead, everything is enclosed in opcode handlers.

And we can use generic io\_uring infrastructure:

- locking and better control of execution context
- completion and other batching
- space in the internal request struct, i.e. struct io\_kiocb
- can be linked to other requests
- possible to execute multiple times, i.e. keeping a BPF request alive

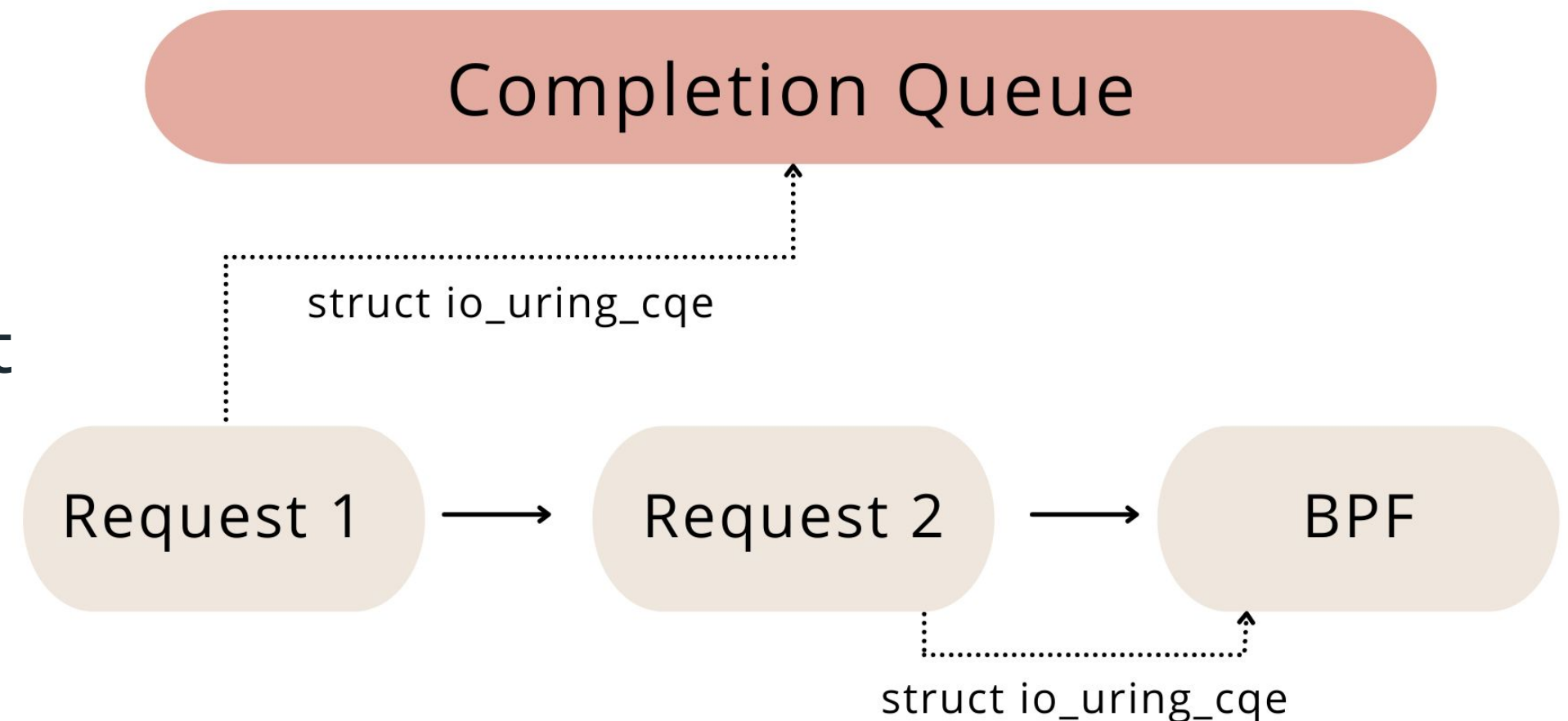
**The downside** is that extra requests are not free, there is a cost to that, but we can work with it.

# Feeding BPF completions

BPF needs feedback from other requests.

**The first idea:** just use links and pass a CQE of the previous request to BPF!

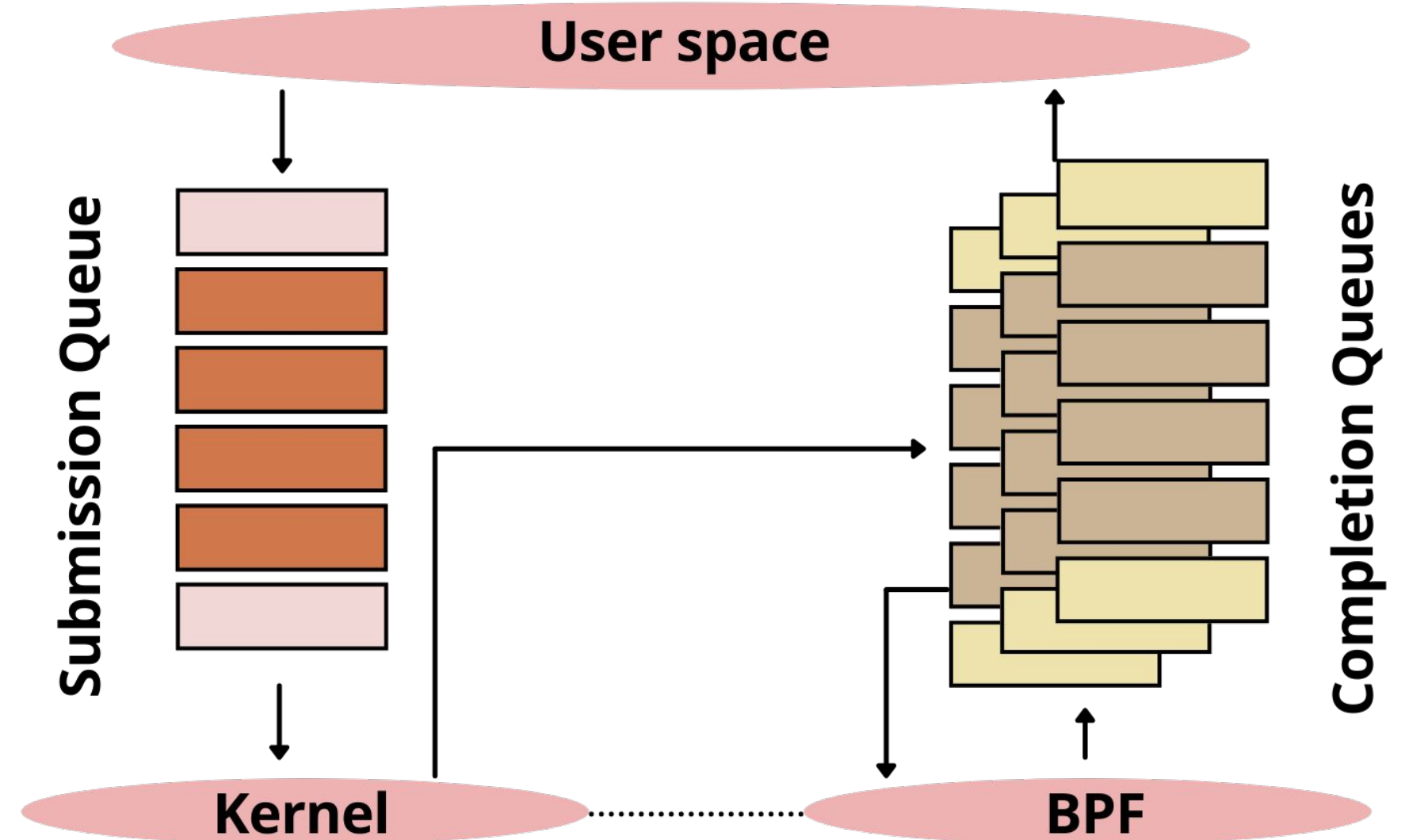
- ugly again
- bound to linking by design
- no way to pass multiple CQEs
- extra overhead for non-BPF code



# Multiple CQs

Introduce multiple CQs:

- `sqe->cq_idx`, each request specifies to which CQ its completion goes
- BPF can emit and consume CQEs to / from any CQ
- Can wait
- Synchronisation is up to the userspace / BPF



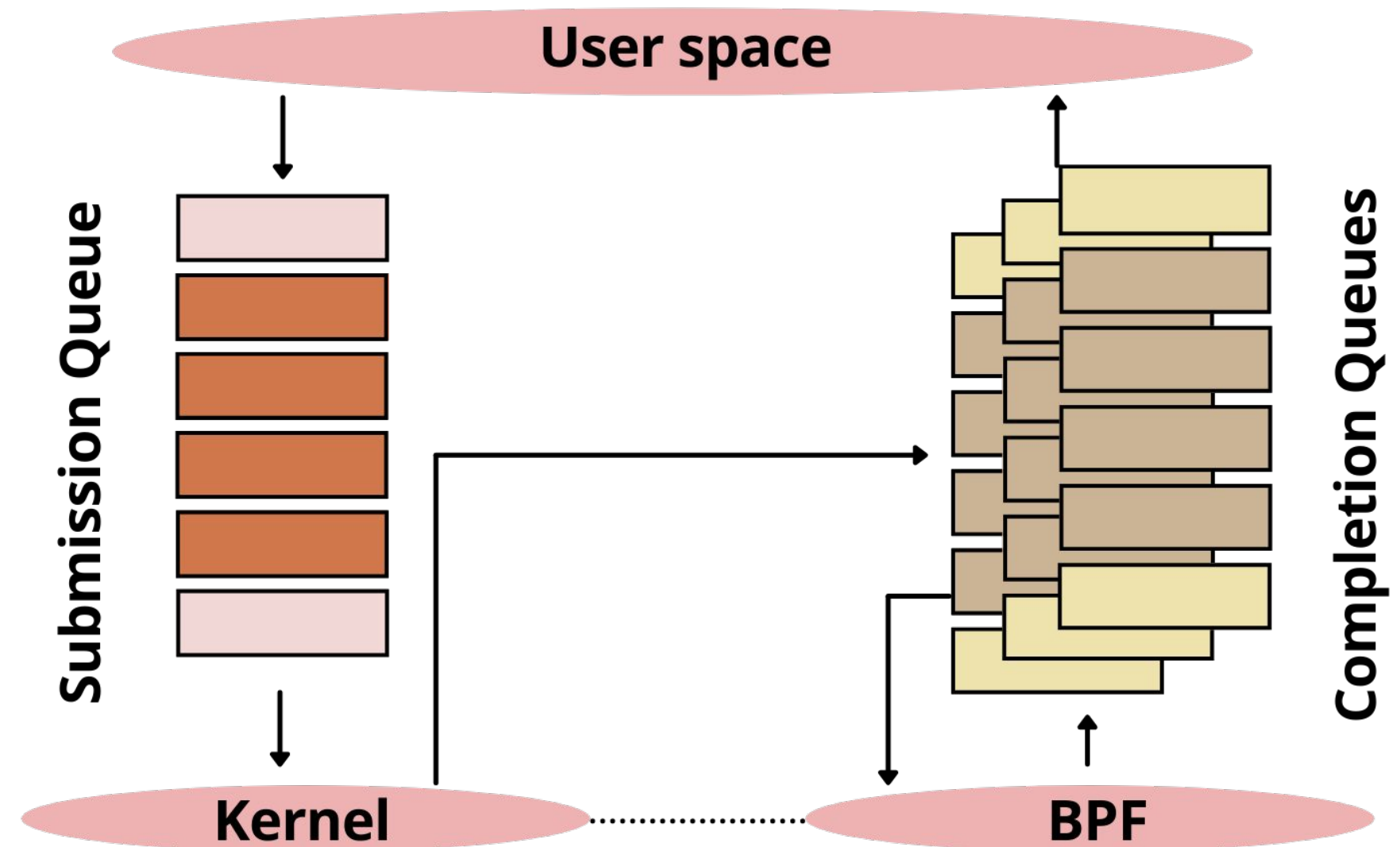


## Pros:

- Can pass multiple CQEs
- CQs can be waited on (including by BPF)
- Extra way of communication:  
posting to a CQ

## Example:

Each BPF request has its own CQ. It keeps a number of operations in-flight and posts to the main CQ when it's done with the job.



What about poking into the normal userspace memory?

BPF subsystem already has an answer: **sleepable BPF programs**

It does what it sounds like, allows BPF programs to sleep.

- reading userspace memory is already there
- writing is trivial to add
- a big deal for io\_uring as submission might need to sleep
- `bpf_copy_[from, to]_user()` + io\_uring performance is yet to be measured

There are also BPF maps / arrays and other infrastructure provided by BPF

- not everything is supported with sleepable programs, may get lifted (if not already)

# Overhead

There can be  $O(N)$  BPF requests, important to keep overhead low

A lot of work has been done! Highlights:

- persistent submission state, request caching
- infrastructure around task\_work and execution batching
- task\_struct referencing and other overhead amortisation
- removing request refcounting
- completion batching
- native io-wq workers (planned to use)
- upcoming IOSQE\_CQE\_SKIP\_SUCCESS
- just cutting the number of instructions required per request ...

QD1 should be in a good shape as well ...

... apart from syscalling and `__do_sys_io_uring_enter`

# API: program registration

API is not set in stone yet, can and will change

```
enum {  
    ...  
    IORING_REGISTER_BPF,  
    IORING_UNREGISTER_BPF,  
};  
  
int bpf_prog_fds[NR_PROGS] = {...};  
// BPF registration can be made optional  
ret = __sys_io_uring_register(ring->ring_fd, IORING_REGISTER_BPF, bpf_prog_fds, NR_PROGS);  
  
// unregister programs, inflicts full quiesce  
ret = __sys_io_uring_register(ring->ring_fd, IORING_UNREGISTER_BPF, 0, 0);  
// or cleaned up automatically on ring exit
```

# API: BPF request

```
enum {  
    ...  
    IORING_OP_BPF,  
};  
  
struct io_uring_sqe *sqe = ...;  
memset(sqe, 0, sizeof(sqe));  
sqe->opcode = IORING_OP_BPF;  
sqe->off = bpf_program_idx;  
// generic, for all request types  
sqe->user_data = (u64)data_ptr; // returned back in CQE. Also, BPF has access to its user_data  
sqe->cq_idx = completion_queue_idx; // CQ index to post CQE to  
sqe->flags = sqe_flags; // combination of IOSQE_*, as usual
```

# API: BPF definitions

```
enum { // Return values for io_uring BPF programs
    IORING_BPF_OK = 0, // complete request
    IORING_BPF_WAIT,   // wait on CQ for completions
};

struct io_uring_bpf_ctx { // BPF io_uring context
    __u64    user_data;    // sqe->user_data specified at submission
    __u32    wait_nr;      // number of requests to wait for
    __u32    wait_idx;     // CQ index to wait on
};

// Returns the number of submitted requests or a negative error if failed.
long (*bpf_io_uring_submit)(void *ctx, void *sqe, __u32 size);
// Returns 0 on success, -ENOMEM if the CQE has been dropped.
long (*bpf_io_uring_emit_cqe)(void *ctx, __u32 cq_idx, __u64 user_data, __s32 res, __u32 cflags);
// Returns 0 on success, -ENOENT if there are no CQEs in the CQ.
long (*bpf_io_uring_reap_cqe)(void *ctx, __u32 cq_idx, struct io_uring_cqe *cqe, __u32 size);
```

# API: libbpf example

```
SEC("iouring") // io_uring BPF program
int bpf_program_name(struct io_uring_bpf_ctx *ctx) {
    struct io_uring_cqe cqe;
    ret = bpf_io_uring_reap_cqe(ctx, cq_idx, &cqe, sizeof(cqe));

    struct io_uring_sqe sqe;
    io_uring_prep_nop(&sqe); // helper copy-pasted from liburing
    sqe.user_data = 42;
    ret = bpf_io_uring_submit(ctx, &sqe, sizeof(sqe));

    u64 data, *uptr = (u64 *)ctx->user_data;
    bpf_copy_from_user(&data, sizeof(data), uptr);

    if (exit) return IORING_BPF_OK;
    ctx->wait_idx = cq_idx_to_wait;
    ctx->wait_nr = nr_cqes_to_wait;
    return IORING_BPF_WAIT; // wait for @nr_cqes_to_wait CQEs in @cq_idx_to_wait CQ
}
```

# API: ideas?

- make BPF registration optional
- extra data to pass in SQE, e.g. maps or shared memory
- more convenient `bpf_copy_[from,to]_user()`, e.g. plain pointers
- other synchronisation, e.g. futex
- batched version of `bpf_io_uring_submit()`
- anything missing?



# Testing

Not yet conclusive. Test case:

- Copy a file by 4KB at a time into `/dev/zero`, buffered and fully cached

Mitigations	Test case	Time (ms)
ON	<code>read(2)/write(2)</code>	1350
ON	<code>io_uring</code> , simple QD=1	1630
any	<code>io_uring</code> + BPF	810
OFF	<code>read(2)/write(2)</code>	550

However, let's take another CPU:

Mitigations	Test case	Time (ms)
ON	<code>read(2)/write(2)</code>	1320
any	<code>io_uring</code> + BPF	1250

# Applicability

Applicability: shouldn't be of interest if batching is naturally “high enough”.

High queue depth is not always possible and/or desirable.

- batching hurts latency
- may care about ordering, e.g. TCP sockets.
- slow devices and memory/responsiveness restrictions

Use cases to try:

- databases / engines, caching systems
- Intelligent file-file splicing, e.g. based on data
- broadcast / collect
- mentioned that may be of use to QUIC
- explore applicability to FUSE
- ideas are welcome

# Next steps

Need to explore more test cases ...  
... and more “interesting” tests.

Each new case requires some tuning and optimisation.  
Upside: also usually benefits non-BPF io\_uring.

Have to solve some slight performance regressions from multi-CQ

- good chance extra CQs will only be visible to BPF

TODO: selftests, bpf\_link, API changes

# Resources

## Kernel

<https://github.com/isilence/linux.git> bpf\_v3

**Liburing**, see <liburing>/examples/bpf/\*

<https://github.com/isilence/liburing.git> bpf\_v3

## io\_uring mailing list

io-uring@vger.kernel.org

## io\_uring guide

[https://kernel.dk/io\\_uring.pdf](https://kernel.dk/io_uring.pdf)

**benchmark**, <fio>/t/io\_uring.c

git://git.kernel.dk/fio