```
Christian Zinck
Problem 45
10/9/18
```

C.

```
A.  <e, \mathsf{rho}_0, \mathsf{sigma}_0>\downarrow <\mathsf{v}, \mathsf{sigma}_1> \\ \hline <\mathsf{VAL}(x,e), \mathsf{rho}_0, \mathsf{sigma}_0>\downarrow <\mathsf{rho}_0\{x\to 1\}, \mathsf{sigma}_1\{1\to v\}> \\ B. \\ ;; In uscheme this code creates a variable x with value #f, sets it to #t, and ;; the if statement evaluates to uscheme. In the new semantics, this code creates a ;; a variable x with value #f, creates a new and separate variable x with value #t, and the ;; if statement evaluates to new. <math display="block"> (\mathsf{val} \ x \ \#f) \\  (\mathsf{val} \ x \ \#f) \\  (\mathsf{val} \ x \ \#t) \\  (\mathsf{if} \ x \ `\mathsf{uscheme} \ `\mathsf{new})
```

I prefer the new design because a val binding should create a new variable and assign it the given value regardless of whether or not the variable already exists and what value the variable holds if it does exist.