

Chapter 3 Context-free Grammars & Parsing

한양대학교 컴퓨터공학부 컴파일러 2014년 2학기



Introduction





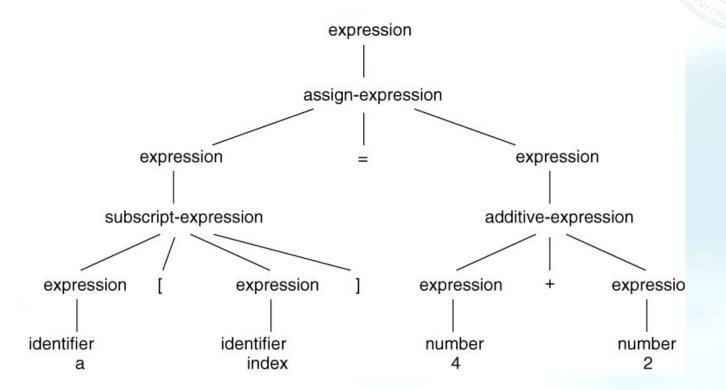
- Parsing
 - task of determining the syntax, or structure, of a program
 - also called syntax analysis
- Syntax of a programming language
 - usually given by the grammar rules of a context-free grammar
- major difference between regular expressions and the rules of a context-free grammar
 - recursion

The parsing process





sequences of tokens -> parse tree or syntax tree

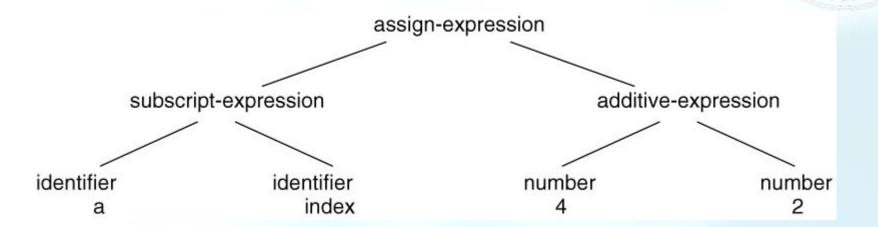


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The parsing process







Context-free grammars



- Representing lexical structures
 - Regular expressions
- Representing syntactic structures
 - Context-free grammars



A context-free grammar is a set of grammar rules.

Context-free grammars





- A context-free grammar rule in BNF (Backus Naur form) consists of a string of symbols.
 - \bullet exp \rightarrow exp op exp | (exp) | number
 - The first symbol is a name for a structure.
 - The second symbol is " \rightarrow "
 - which is followed by a string of symbols, each of which is either a token, a name, or |.

Context-free grammars





Other form of grammar rules

- \bullet exp \rightarrow exp op exp | (exp) | number
- <exp>::= <exp> <op> <exp> | (<exp>) | NUMBER
- Equivalence of grammar rules
 - \bullet exp \rightarrow exp op exp | (exp) | number
 - $exp \rightarrow exp$ op exp $exp \rightarrow (exp)$ $exp \rightarrow number$

Grammar and derivations





Grammar: set of grammar rules

start
$$exp \rightarrow exp \ op \ exp \ | \ (exp \) \ | \ number$$
symbol $exp \rightarrow + | - | *$

- A derivation is a sequence of replacements of structure names by choices on the right-hand sides of grammar rules.
 - (34 3) * 42
 - \bullet exp => exp op exp
 - ⇒ exp op **number**
 - *⇒ exp * number*
 - \Rightarrow (exp) * number
 - \Rightarrow (exp op exp) * number
 - ⇒ (exp op number) * number
 - ⇒ (exp number) * number
 - \Rightarrow (number number)* number

sentential forms

sentence

The language defined by a grammar





- The language defined by a grammar is the set of all strings of token symbols obtained by derivation.
 - $L(G) = \{ s \mid exp => * s \}$
- Nonterminals and terminals
 - \bullet exp \rightarrow exp op exp | (exp) | number
 - \circ op \rightarrow + | | *



- Example 3.1
 - $\bullet E \rightarrow (E) \mid a$
- Example 3.2
 - $\bullet E \rightarrow (E)$
- Example 3.3

$$\bullet E \rightarrow E + a \mid a$$







- Example 3.4
 - statement → if-stmt | **other**
 - if-stmt → if (exp) statement | if (exp) statement else statement
 - \bullet exp \rightarrow 0 | 1

other

if (0) other

if (1) other else other

if (0) if (0) other

if (0) if (1) other else other

if (1) other else if (0) other else other





Repetition

- **o** a+
 - $A \rightarrow Aa \mid a$
 - \bullet $A \Rightarrow Aa \Rightarrow Aaa \Rightarrow Aaaa \Rightarrow aaaa$ (left recursive)
 - $A \rightarrow aA \mid a$
 - \bullet $A \Rightarrow aA \Rightarrow aaA \Rightarrow aaaA \Rightarrow aaaa$ (right recursive)

•
$$A \rightarrow A\alpha | \beta$$

• βα*

•
$$A \rightarrow \alpha A | \beta$$

$$\circ \alpha * \beta$$





- Repetition
 - a*
 - $\bullet A \rightarrow Aa \mid \varepsilon$
 - $A \rightarrow aA \mid \varepsilon$
- Example 3.5
 - $\bullet A \rightarrow (A)A \mid \varepsilon$







- Example 3.4
 - statement → if-stmt | other
 - if-stmt → if (exp) statement | if (exp) statement else statement
 - \bullet exp \rightarrow 0 | 1

Example 3.6

- statement → if-stmt | other
- if-stmt \rightarrow if (exp) statement else-part
- else-part \rightarrow else statement | ε
- \bullet exp \rightarrow 0 | 1





Example 3.7

- o stmt-sequence → stmt; stmt-sequence | stmt
- \circ stmt \rightarrow s

(add ε)

- stmt-sequence \rightarrow stmt; stmt-sequence | ε
- $stmt \rightarrow s$

$$\{\varepsilon, s;, s;s;, s;s;s;, \dots\}$$

(no; in the end)

- stmt-sequence \rightarrow nonempty-stmt-sequence $\mid \varepsilon \mid$
- nonempty-stmt-sequence → stmt; nonempty-stmt-sequence | stmt
- \circ stmt \rightarrow s

$$\{\varepsilon, s, s; s, s; s; s, \dots\}$$



Parse trees

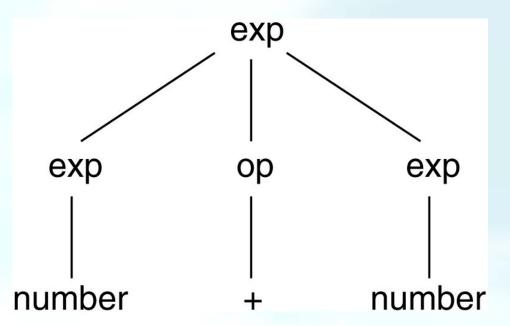




- Grammar
 - \bullet exp \rightarrow exp op exp | (exp) | number

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- \circ op \rightarrow + | | *
- A derivation for *number + number*
 - exp => *exp op exp*
 - \Rightarrow **number** op exp
 - \Rightarrow number + exp
 - ⇒ number + number



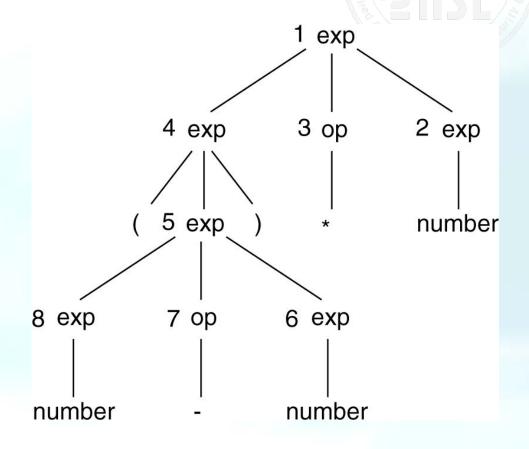
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Parse trees





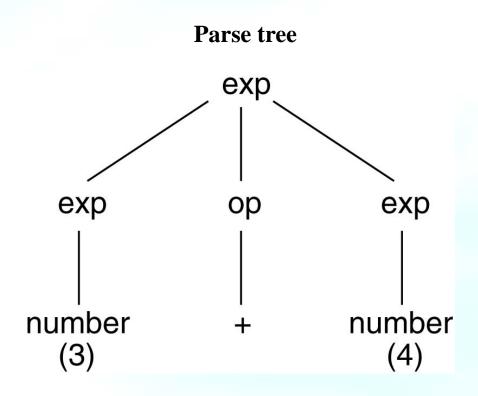
- Grammar
 - $exp \rightarrow exp \ op \ exp \ | \ (exp) \ | \ number$
 - $op \to + |-|*$
- A parse tree for **(34-3)*42**

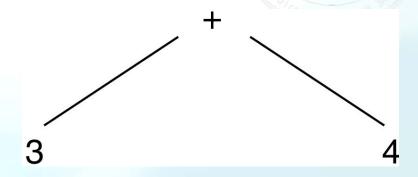






3+4

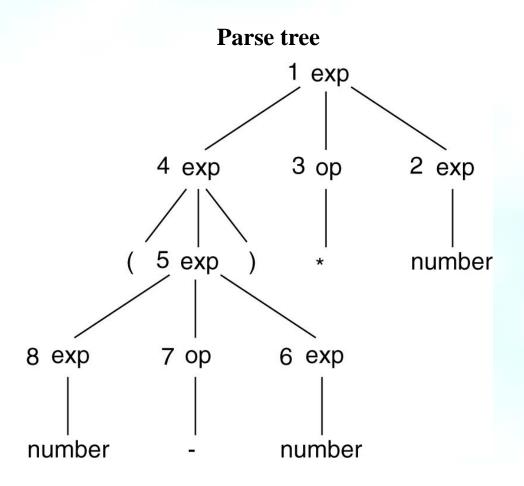


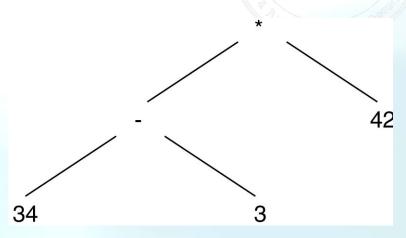






• (34-3)*42

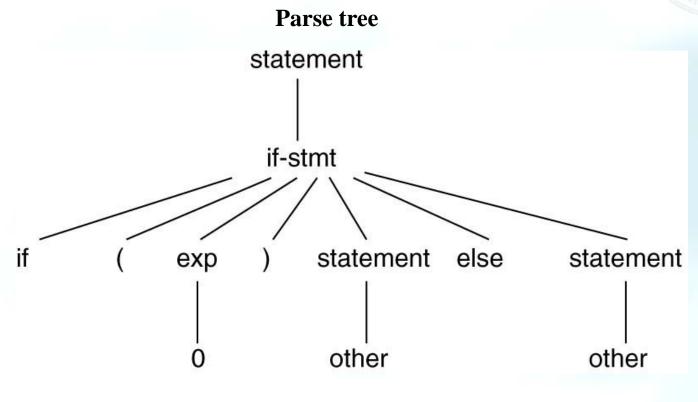








- if (0) other else other
 - statement → if-stmt | other
 - if-stmt \rightarrow if (exp) statement | if (exp) statement else statement
 - \bullet exp \rightarrow 0 | 1

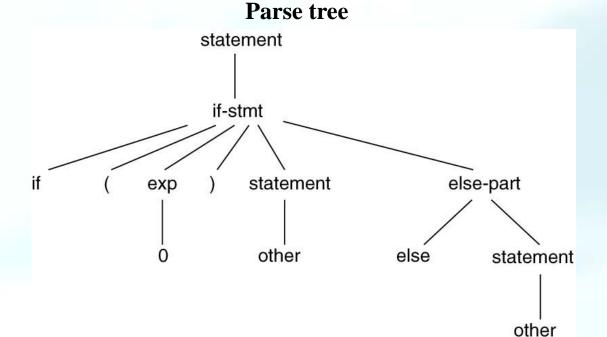






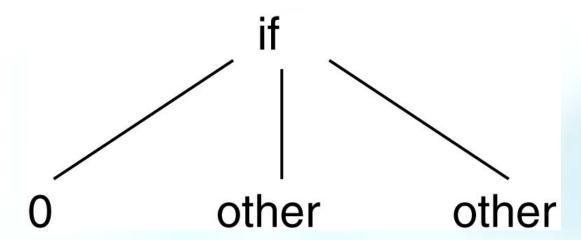
- if (0) other else other
 - statement → if-stmt / other
 - if-stmt \rightarrow if (exp) statement else-part
 - else-part \rightarrow else statement | ε
 - \bullet exp \rightarrow 0 | 1







• if (0) other else other

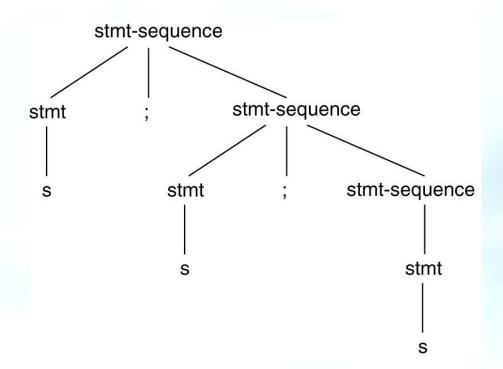






- S; S; S
 - o stmt-sequence → stmt; stmt-sequence | stmt
 - \circ stmt \rightarrow s

Parse tree



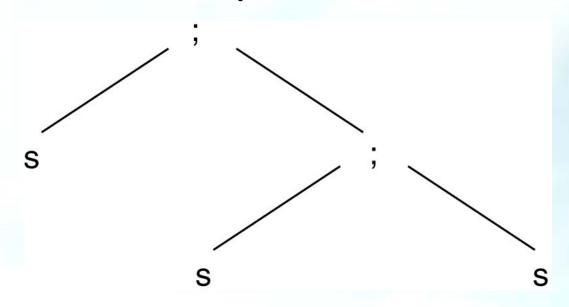






- S; S; S
 - stmt-sequence → stmt; stmt-sequence | stmt
 - \circ stmt \rightarrow s



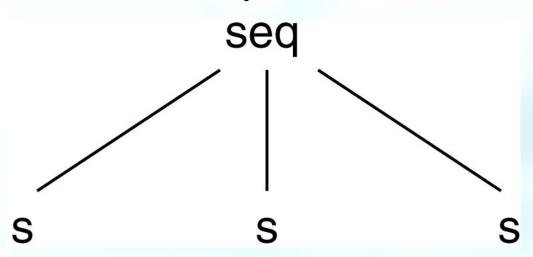




2nst

- S; S; S
 - o stmt-sequence → stmt; stmt-sequence | stmt
 - $stmt \rightarrow s$



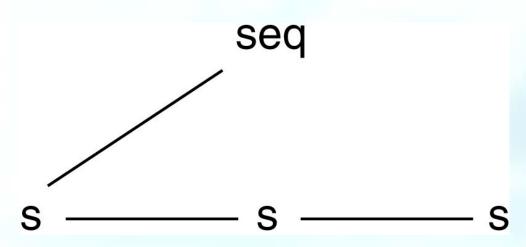






- S; S; S
 - stmt-sequence → stmt; stmt-sequence | stmt
 - $stmt \rightarrow s$









- S; S; S
 - stmt-sequence → stmt; stmt-sequence | stmt
 - $stmt \rightarrow s$

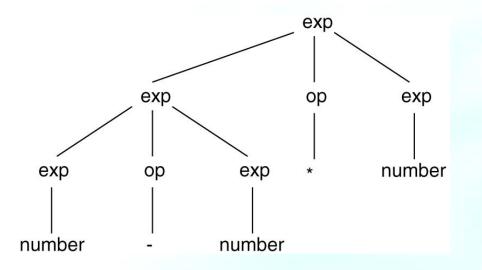


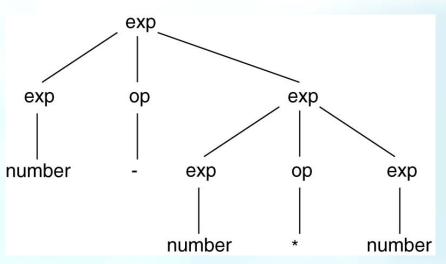




- number number * number
 - \bullet exp \rightarrow exp op exp | (exp) | number
 - \circ op \rightarrow + | | *



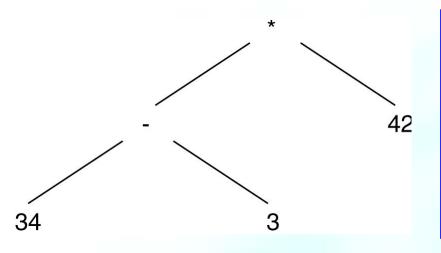


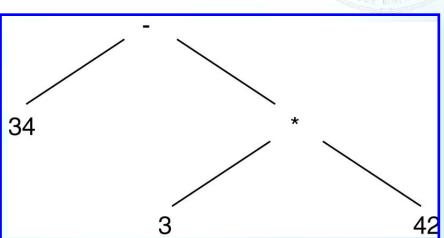






- 34 3 * 42 (precedence)
 - \bullet exp \rightarrow exp op exp | (exp) | number
 - \circ op \rightarrow + |-|*

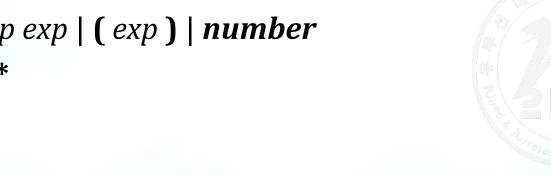


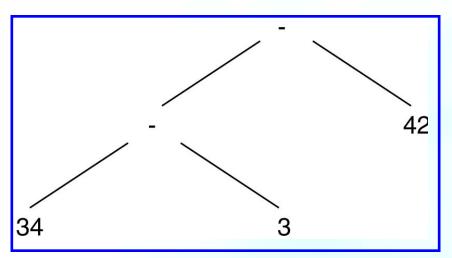


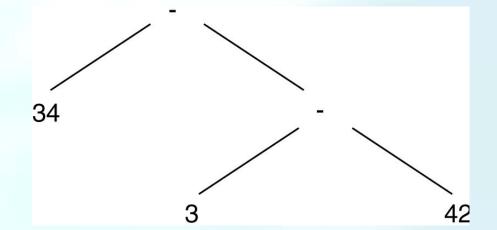




- 34 3 42 (Associativity)
 - \bullet exp \rightarrow exp op exp | (exp) | number
 - \circ op \rightarrow + | | *











Precedence cascade

- exp → exp addop exp | term
- addop \rightarrow + \mid -
- term → term mulop term | factor
- $mulop \rightarrow *$
- $factor \rightarrow (exp) / number$

Associativity (left)

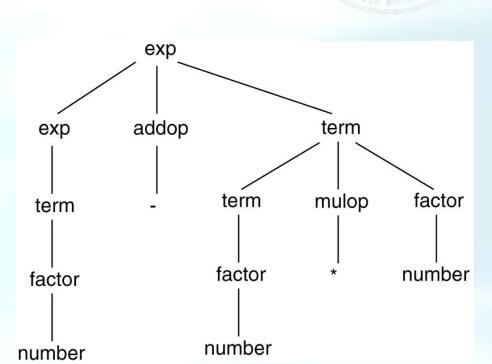
- exp → exp addop term | term
- addop \rightarrow + \mid -
- term → term mulop factor | factor
- $mulop \rightarrow *$
- $factor \rightarrow (exp) / number$







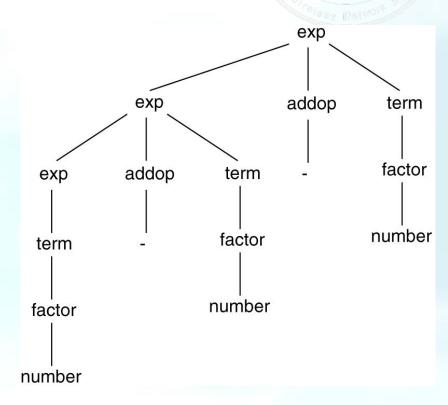
- 34 3 * 42
 - exp → exp addop term | term
 - $addop \rightarrow + | -$
 - term → term mulop factor | factor
 - $mulop \rightarrow *$
 - $factor \rightarrow (exp) / number$







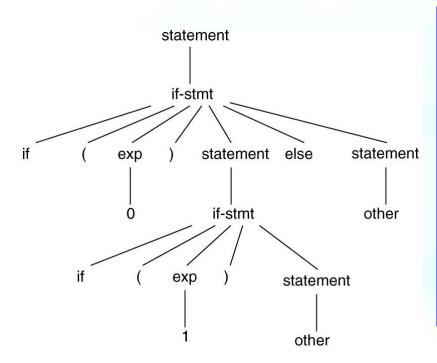
- 34 3 42
 - exp → exp addop term | term
 - $addop \rightarrow + | -$
 - term → term mulop factor | factor
 - $mulop \rightarrow *$
 - $factor \rightarrow (exp) / number$

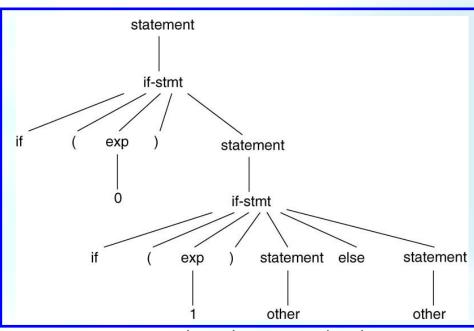






- if (0) if (1) other else other (Dangling else)
 - statement → if-stmt / other
 - if-stmt \rightarrow if (exp) statement | if (exp) statement else statement
 - \bullet exp \rightarrow 0 | 1





most closely nested rule

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- A new grammar without the dangling else problem.
 - statement → matched-stmt | unmatched-stmt
 - matched-stmt → if (exp) matched-stmt else matched-stmt | other

 - $exp \rightarrow 0 \mid 1$

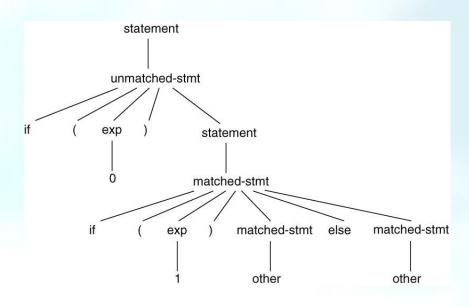




- if (0) if (1) other else other
 - statement → matched-stmt | unmatched-stmt
 - matched-stmt → if (exp) matched-stmt else matched-stmt | other
 - unmatched- $stmt \rightarrow if (exp) statement$

| **if (** exp **)** matched-stmt **else** unmatched-stmt

 \bullet exp \rightarrow 0 | 1



Ambiguity





- Other methods to solving dangling else problem
 - Disambiguating rules
 - Bracketing keywords (p. 122)
 - if-stmt → if condition then statement-sequence end if
 | if condition then statement-sequence else statement-sequence end if
- Inessential ambiguity (p. 122)
 - Example
 - stmt-sequence → stmt-sequence; stmt-sequence | stmt
 - $stmt \rightarrow s$
 - But still obtain unique syntax trees

EBNF





- {}
 - Repetition
 - $A \rightarrow A\alpha \mid \beta$ (left recursive) : $A \rightarrow \beta\{\alpha\}$
 - \circ *A* → α*A* | β (right recursive) : *A* → {α}β
 - o p. 124
- []
 - Optional
 - *if-stmt* → **if** (*exp*) *statement* [**else** *statement*]
- ()
 - Choice







- Representing EBNF
 - Nonterminals
 - Square or rectangle boxes
 - Terminals
 - round or oval boxes

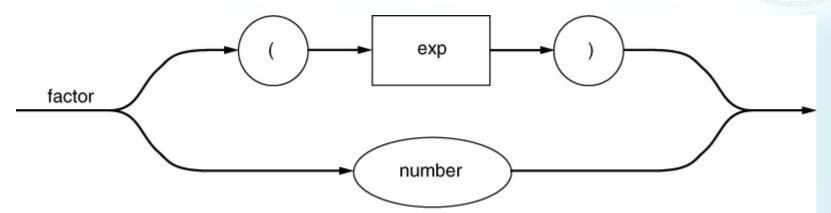






• $factor \rightarrow (exp) \mid number$

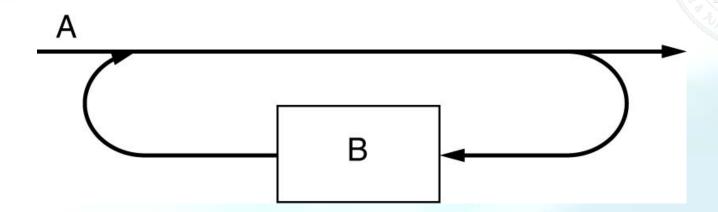








 $\bullet A \rightarrow \{B\}$



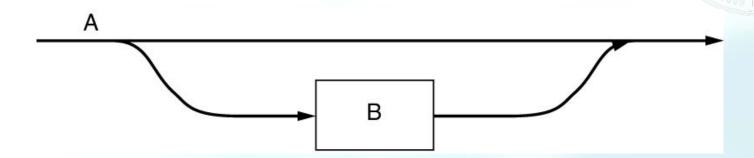






 \bullet $A \rightarrow [B]$







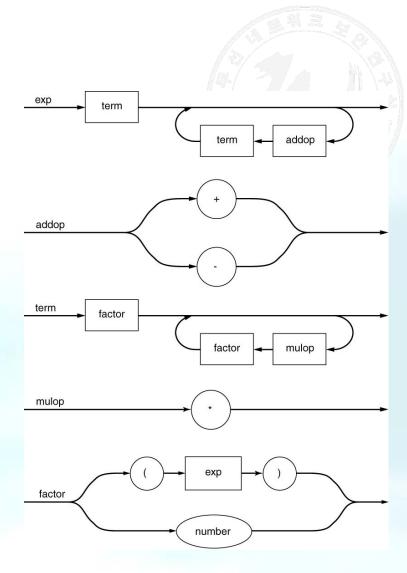


BNF

- exp → exp addop term | term
- addop \rightarrow + / -
- term → term mulop factor | factor
- $mulop \rightarrow *$
- $factor \rightarrow (exp) / number$

EBNF

- exp → term {addop term}
- $addop \rightarrow + | -$
- term → factor {mulop factor}
- $mulop \rightarrow *$
- $factor \rightarrow (exp) / number$







BNF

- statement → if-stmt / other
- if-stmt → if (exp) statement/ if (exp) statement else statement
- \bullet exp \rightarrow 0 | 1



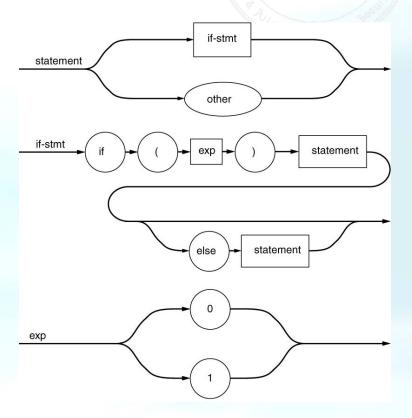
EBNF

- statement → if-stmt / other
- *if-stmt* → **if** (*exp*) *statement* [**else** *statement*]
- \bullet exp \rightarrow 0 | 1





- EBNF
 - statement → if-stmt | other
 - *if-stmt* → **if** (*exp*) *statement* [**else** *statement*]
 - $exp \rightarrow 0 \mid 1$



Syntax of the TINY language





```
program \rightarrow stmt-sequence
stmt-sequence \rightarrow stmt-sequence; statement \mid statement
statement → if-stmt | repeat-stmt | assign-stmt | read-stmt | write-stmt
if-stmt \rightarrow if exp then stmt-sequence end
          / if exp then stmt-sequence else stmt-sequence end
repeat-stmt -> repeat stmt-sequence until exp
assign-stmt → identifier := exp
read-stmt → read identifier
write-stmt \rightarrow write exp
exp \rightarrow simple-exp comparison-op simple-exp | simple-exp
comparison-op \rightarrow < | =
simple-exp \rightarrow simple-exp addop term | term
addop \rightarrow + \mid -
term → term mulop factor | factor
mulop \rightarrow * | /
factor \rightarrow (exp) / number / identifier
```

Syntax of the TINY language





A context-free grammar for TINY (Fig. 3.6)

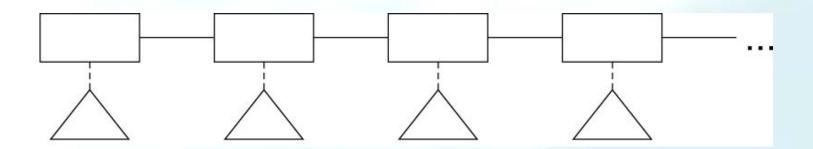
- 1: program is a statement sequence.
- 2: A statement sequence is a list of statement separated by ;.
- 3: There are five kinds of statements.
- 4-8: if, repeat, assign, read, and write statements.
- 9-15: expressions
- Precedence: *, / > +,- > <, =</pre>





- statement sequence
 - o stmt-sequence → stmt-sequence ; statement | statement

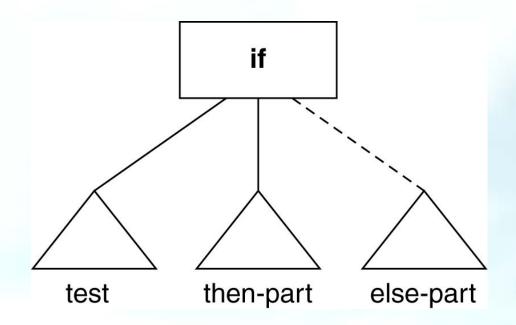






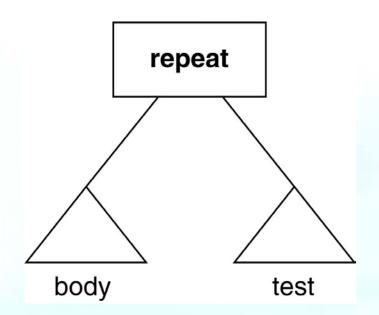


- if statement
 - if-stmt → if exp then stmt-sequence end
 | if exp then stmt-sequence else stmt-sequence end



- - Z_nSt

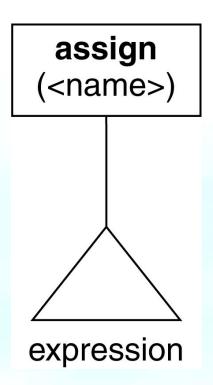
- repeat statement
 - repeat-stmt → repeat stmt-sequence until exp







- assign statement
 - assign-stmt → identifier := exp

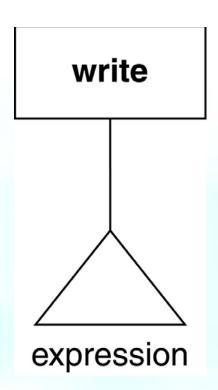








- write statement
 - write-stmt → write exp





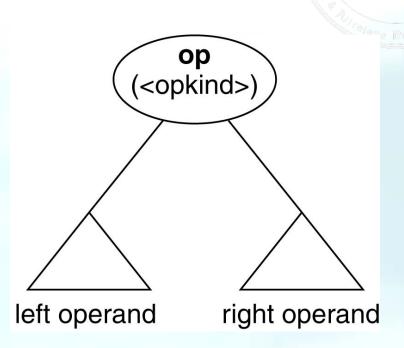




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operation

- comparison-op \rightarrow < | =
- addop → + | -
- $mulop \rightarrow * | /$







- other statements?
 - read-stmt → read identifier
 - \bullet addop \rightarrow + \mid -
 - *mulop* → * | /
 - No syntax trees for grammar rules with only nonterminals on the right hand.





Sample program in TINY language

```
{ Sample program
  in TINY language -
  computes factorial
read x; { input an integer }
if 0 < x then { don't compute if x <= 0 }
fact := 1;
 repeat
 fact := fact * x;
  x := x - 1
 until x = 0;
 write fact { output factorial of x }
end
```







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