

# Database Systems

## Lecture #05

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# Objectives



- ◆ To learn the relational data model
  - Concepts of a relational model
  - Notations of a relational model
  - Constraints of a relational model

- ◆ Relational Model Concepts
- ◆ Characteristics of Relations
- ◆ Relational Model Constraints
  - Key Constraints
  - Relational Database Schema
  - Integrity Constraints

# Relational Model Concepts

## ◆ Database

- Collection of *relations* (or *tables*)

## ◆ Relation

- Set of *tuples* (or *rows*)

# Relational Model Concepts

## ◆ Tuple

- A fact that typically corresponds to a real-world *entity or relationship*
- Consists of a set of *attributes*

## ◆ Attribute

- Defines a property of each object corresponding to a *tuple*

# Relational Model Concepts



## STUDENT

| Name  | Student_number | Class | Major |
|-------|----------------|-------|-------|
| Smith | 17             | 1     | CS    |
| Brown | 8              | 2     | CS    |

## COURSE

| Course_name               | Course_number | Credit_hours | Department |
|---------------------------|---------------|--------------|------------|
| Intro to Computer Science | CS1310        | 4            | CS         |
| Data Structures           | CS3320        | 4            | CS         |
| Discrete Mathematics      | MATH2410      | 3            | MATH       |
| Database                  | CS3380        | 3            | CS         |

## PREREQUISITE

| Course_number | Prerequisite_number |
|---------------|---------------------|
| CS3380        | CS3320              |
| CS3380        | MATH2410            |
| CS3320        | CS1310              |

## SECTION

| Section_identifier | Course_number | Semester | Year | Instructor |
|--------------------|---------------|----------|------|------------|
| 85                 | MATH2410      | Fall     | 07   | King       |
| 92                 | CS1310        | Fall     | 07   | Anderson   |
| 102                | CS3320        | Spring   | 08   | Knuth      |
| 112                | MATH2410      | Fall     | 08   | Chang      |
| 119                | CS1310        | Fall     | 08   | Anderson   |
| 135                | CS3380        | Fall     | 08   | Stone      |

## GRADE\_REPORT

| Student_number | Section_identifier | Grade |
|----------------|--------------------|-------|
| 17             | 112                | B     |
| 17             | 119                | C     |
| 8              | 85                 | A     |
| 8              | 92                 | A     |
| 8              | 102                | B     |
| 8              | 135                | A     |

# Relational Model Concepts



Diagram illustrating the components of a Relational Model:

- Relation Name:** POINTS TO **STUDENT**
- Attributes:** POINTS TO the column headers: **Name**, **Ssn**, **Home\_phone**, **Address**, **Office\_phone**, **Age**, and **Gpa**
- Tuples:** POINTS TO the rows of data in the table

| Name           | Ssn         | Home_phone    | Address              | Office_phone  | Age | Gpa  |
|----------------|-------------|---------------|----------------------|---------------|-----|------|
| Benjamin Bayer | 305-61-2435 | (817)373-1616 | 2918 Bluebonnet Lane | NULL          | 19  | 3.21 |
| Chung-cha Kim  | 381-62-1245 | (817)375-4409 | 125 Kirby Road       | NULL          | 18  | 2.89 |
| Dick Davidson  | 422-11-2320 | NULL          | 3452 Elgin Road      | (817)749-1253 | 25  | 3.53 |
| Rohan Panchal  | 489-22-1100 | (817)376-9821 | 265 Lark Lane        | (817)749-6492 | 28  | 3.93 |
| Barbara Benson | 533-69-1238 | (817)839-8461 | 7384 Fontana Lane    | NULL          | 19  | 3.25 |

# Relational Model Concepts

## ◆ Domain $D$

- A set of *atomic* values
- Example: domain of Age attribute
  - set of values between 16 ~ 80



# Relational Model Concepts

## ◆ Attribute $A_i$

- *Name of a role* played by some domain  $D$  in the relation schema  $R$
- $\text{dom}(A_i)$ : Domain of attribute  $A_i$  in relation schema  $R$
- Example: Employee\_ages attribute in relation schema  $R$ 
  - Has domain  $\text{dom}(\text{Employee\_ages})(=D)$

# Relational Model Concepts

## ◆ Relation schema $R$

- Made up of a relation name  $R$  and a list of attributes,  $A_1, A_2, \dots, A_n$
- Denoted by  $R(A_1, A_2, \dots, A_n)$
- Example
  - STUDENT(Name, SSN, HomePhone, Address, OfficePhone, Age, GPA)

# Relational Model Concepts

- ◆ Degree (or arity) of a relation
  - Number of attributes  $n$  of its relation schema
  - Example: degree of STUDENT relation
    - 7 (seven attributes)

# Relational Model Concepts

## ◆ Tuple $t$ of relation $R$

- Also called  *$n$ -tuple*
- Ordered list of  $n$  values  $t = \langle v_1, v_2, \dots, v_n \rangle$
- Each value  $v_i$ ,  $1 \leq i \leq n$ , is
  - an element of  $\text{dom}(A_i)$  or
  - a special *NULL* value

# Relational Model Concepts

## ◆ Relation instance $r(R)$

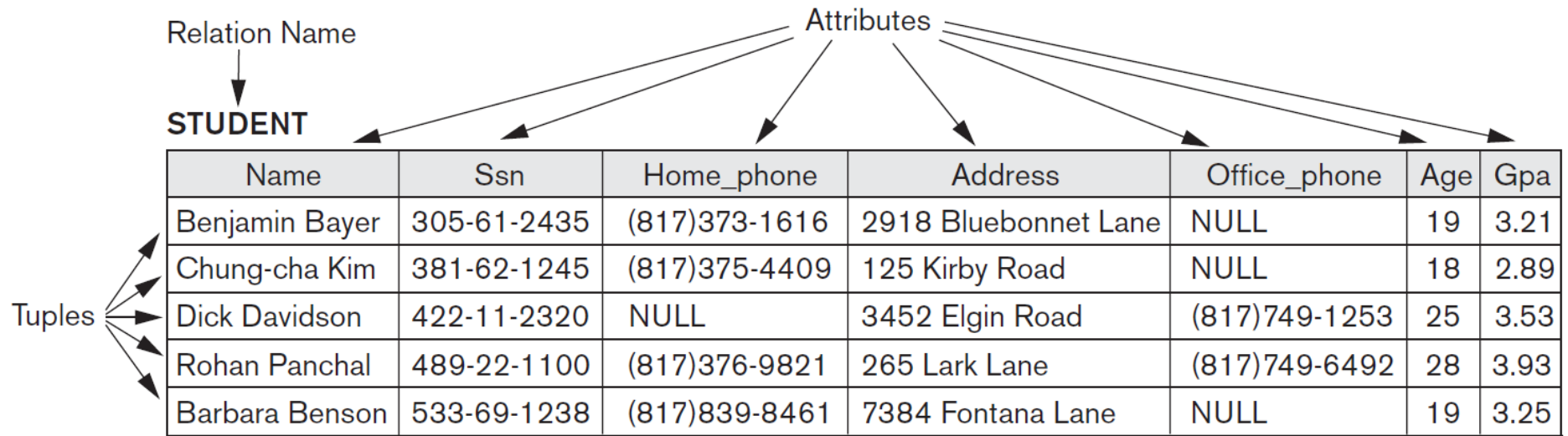
### ● Set of *tuples*

- $r(R) = \{t_1, t_2, \dots, t_m\}$

### ● *Subset* of the *Cartesian product* of the domains that define R:

- $r(R) \subseteq (\text{dom}(A_1) \times \text{dom}(A_2) \times \dots \times \text{dom}(A_n))$

# Relational Model Concepts



# Characteristics of Relations

- ◆ Ordering of *tuples* in a relation
  - Relation defined as a *set* of tuples
  - Tuples have *no order* among them
    - Although tuples can be represented as a table

# Characteristics of Relations

- ◆ Ordering of *attributes* within a tuple
  - *n-tuple* defined as a *ordered list of n values*
  - Order of attributes and values is not that important
    - As long as the correspondence between attributes and values maintained



# Characteristics of Relations

## ◆ Values in tuples

- Each value in a tuple is *atomic*
  - Composite and multivalued attributes not allowed
  - *First normal form* assumption
- NULLs
  - Represent the values of attributes that may be *unknown* or may be *not applicable* to a tuple

# Relational Model Notation

## ◆ $t[A_i]$

- refers to the value  $v_i$  in  $t$  for attribute  $A_i$

## ◆ $t[A_U, A_W, \dots, A_Z]$

- refers to the sub-tuple of values  $\langle v_U, v_W, \dots, v_Z \rangle$  from  $t$  corresponding to the attributes specified in the list

# Relational Model Constraints

## ◆ Constraints

- *Restrictions* on the actual values in a relation schema
- Every instance in the relation should follow them
- Derived from the *rules* in the mini-world

# Relational Model Constraints

- ◆ Important constraints in relational model
  - Key constraint
  - Entity integrity constraint
  - Referential integrity constraint

# Key Constraints

- ◆ *Superkey* of relational schema  $R$ 
  - A set of attributes SK in  $R$
  - No two distinct tuples in any relation instance  $r(R)$  of  $R$  can have the same value for SK
  - For any  $t_1$  and  $t_2$  ( $t_1 \neq t_2$ ) in  $r(R)$ 
    - $t_1[SK] \neq t_2[SK]$
  - Called *uniqueness constraint*

# Key Constraints

## ◆ *Candidate key* $K$

- Superkey of  $R$
- *Minimal* superkey
  - Removing any attribute  $A$  from  $K$
  - Leaves a set of attributes  $K$  that *is not a superkey* of  $R$  any more
- A relation schema may have more than one key
  - Candidate

# Key Constraints



## CAR

| <u>License_number</u> | Engine_serial_number | Make       | Model   | Year |
|-----------------------|----------------------|------------|---------|------|
| Texas ABC-739         | A69352               | Ford       | Mustang | 02   |
| Florida TVP-347       | B43696               | Oldsmobile | Cutlass | 05   |
| New York MPO-22       | X83554               | Oldsmobile | Delta   | 01   |
| California 432-TFY    | C43742               | Mercedes   | 190-D   | 99   |
| California RSK-629    | Y82935               | Toyota     | Camry   | 04   |
| Texas RSK-629         | U028365              | Jaguar     | XJS     | 04   |

# Key Constraints



- ◆ Example: STUDENT relational schema
  - {SSN, Name, Age}: is a superkey / *not a candidate key*
  - {SSN}: is a superkey / is a candidate key



# Key Constraints

- ◆ Primary key (PK) of the relation
  - Designated among *candidate keys*
    - By a database designer
  - Underlined attribute
  - Example: DEPARTMENT relational schema
    - DepartmentID, DepartmentName

# Key Constraints



## ◆ Key constraint

- Specifies a relation must have candidate key attributes

# Relational Database Schema

- ◆ Relational database schema
  - Set of *relation schemas* belonging to the same database
    - $S = \{R_1, R_2, \dots, R_m\}$
  - Contains a set of *integrity constraints* IC

# Relational Database Schema

## EMPLOYEE

|       |       |       |            |       |         |     |        |           |     |
|-------|-------|-------|------------|-------|---------|-----|--------|-----------|-----|
| Fname | Minit | Lname | <u>Ssn</u> | Bdate | Address | Sex | Salary | Super_ssn | Dno |
|-------|-------|-------|------------|-------|---------|-----|--------|-----------|-----|

## DEPARTMENT

|       |                |         |                |
|-------|----------------|---------|----------------|
| Dname | <u>Dnumber</u> | Mgr_ssn | Mgr_start_date |
|-------|----------------|---------|----------------|

## DEPT\_LOCATIONS

|                |                  |
|----------------|------------------|
| <u>Dnumber</u> | <u>Dlocation</u> |
|----------------|------------------|

## PROJECT

|       |                |           |      |
|-------|----------------|-----------|------|
| Pname | <u>Pnumber</u> | Plocation | Dnum |
|-------|----------------|-----------|------|

## WORKS\_ON

|             |            |       |
|-------------|------------|-------|
| <u>Essn</u> | <u>Pno</u> | Hours |
|-------------|------------|-------|

## DEPENDENT

|             |                       |     |       |              |
|-------------|-----------------------|-----|-------|--------------|
| <u>Essn</u> | <u>Dependent_name</u> | Sex | Bdate | Relationship |
|-------------|-----------------------|-----|-------|--------------|

# Relational Database Schema

## ◆ Relational database instance

- A set of relation instances within the same database
  - $\{r_1, r_2, \dots, r_m\}$
- $r_i$  is an instance of  $R_i$  which satisfies the integrity constraints specified in IC

# Relational Database Schema



## EMPLOYEE

| Fname    | Minit | Lname   | <u>Ssn</u> | Bdate      | Address                  | Sex | Salary | Super_ssn | Dno |
|----------|-------|---------|------------|------------|--------------------------|-----|--------|-----------|-----|
| John     | B     | Smith   | 123456789  | 1965-01-09 | 731 Fondren, Houston, TX | M   | 30000  | 333445555 | 5   |
| Franklin | T     | Wong    | 333445555  | 1955-12-08 | 638 Voss, Houston, TX    | M   | 40000  | 888665555 | 5   |
| Alicia   | J     | Zelaya  | 999887777  | 1968-01-19 | 3321 Castle, Spring, TX  | F   | 25000  | 987654321 | 4   |
| Jennifer | S     | Wallace | 987654321  | 1941-06-20 | 291 Berry, Bellaire, TX  | F   | 43000  | 888665555 | 4   |
| Ramesh   | K     | Narayan | 666884444  | 1962-09-15 | 975 Fire Oak, Humble, TX | M   | 38000  | 333445555 | 5   |
| Joyce    | A     | English | 453453453  | 1972-07-31 | 5631 Rice, Houston, TX   | F   | 25000  | 333445555 | 5   |
| Ahmad    | V     | Jabbar  | 987987987  | 1969-03-29 | 980 Dallas, Houston, TX  | M   | 25000  | 987654321 | 4   |
| James    | E     | Borg    | 888665555  | 1937-11-10 | 450 Stone, Houston, TX   | M   | 55000  | NULL      | 1   |

## DEPARTMENT

| Dname          | <u>Dnumber</u> | Mgr_ssn   | Mgr_start_date |
|----------------|----------------|-----------|----------------|
| Research       | 5              | 333445555 | 1988-05-22     |
| Administration | 4              | 987654321 | 1995-01-01     |
| Headquarters   | 1              | 888665555 | 1981-06-19     |

## DEPT\_LOCATIONS

| <u>Dnumber</u> | <u>Dlocation</u> |
|----------------|------------------|
| 1              | Houston          |
| 4              | Stafford         |
| 5              | Bellaire         |
| 5              | Sugarland        |
| 5              | Houston          |

# Relational Database Schema



**WORKS\_ON**

| <u>Essn</u> | <u>Pno</u> | Hours |
|-------------|------------|-------|
| 123456789   | 1          | 32.5  |
| 123456789   | 2          | 7.5   |
| 666884444   | 3          | 40.0  |
| 453453453   | 1          | 20.0  |
| 453453453   | 2          | 20.0  |
| 333445555   | 2          | 10.0  |
| 333445555   | 3          | 10.0  |
| 333445555   | 10         | 10.0  |
| 333445555   | 20         | 10.0  |
| 999887777   | 30         | 30.0  |
| 999887777   | 10         | 10.0  |
| 987987987   | 10         | 35.0  |
| 987987987   | 30         | 5.0   |
| 987654321   | 30         | 20.0  |
| 987654321   | 20         | 15.0  |
| 888665555   | 20         | NULL  |

**PROJECT**

| Pname           | <u>Pnumber</u> | Plocation | Dnum |
|-----------------|----------------|-----------|------|
| ProductX        | 1              | Bellaire  | 5    |
| ProductY        | 2              | Sugarland | 5    |
| ProductZ        | 3              | Houston   | 5    |
| Computerization | 10             | Stafford  | 4    |
| Reorganization  | 20             | Houston   | 1    |
| Newbenefits     | 30             | Stafford  | 4    |

**DEPENDENT**

| <u>Essn</u> | <u>Dependent_name</u> | Sex | Bdate      | Relationship |
|-------------|-----------------------|-----|------------|--------------|
| 333445555   | Alice                 | F   | 1986-04-05 | Daughter     |
| 333445555   | Theodore              | M   | 1983-10-25 | Son          |
| 333445555   | Joy                   | F   | 1958-05-03 | Spouse       |
| 987654321   | Abner                 | M   | 1942-02-28 | Spouse       |
| 123456789   | Michael               | M   | 1988-01-04 | Son          |
| 123456789   | Alice                 | F   | 1988-12-30 | Daughter     |
| 123456789   | Elizabeth             | F   | 1967-05-05 | Spouse       |

# Entity Integrity Constraints

- ◆ Constraint for the attributes in PK of each relation schema  $R$
- ◆ No primary key value can be NULL
- ◆ For every tuple  $t$  in  $r(R)$ 
  - $t[\text{PK}] \neq \text{NULL}$
- ◆ The primary key value is used to identify individual tuples in a relation



# Referential Integrity Constraints

- ◆ Constraints specified *between two relations*
  - Key constraints and entity integrity constraints are specified on *an individual relation*
- ◆ Used to maintain the *consistency* among tuples in the two relations

# Referential Integrity Constraints

- ◆ A tuple in one relation should refer to an *existing tuple* in that relation
- ◆ Example: DNO attribute in EMPLOYEE should refer to *existing value* for DNumber attribute in DEPARTMENT

# Referential Integrity Constraints

## ◆ Foreign key (FK)

- A set of attributes FK that satisfies the following rules for relation schema  $R_1$  and  $R_2$ :
  - The attributes in FK in  $R_1$  have the same domain(s) as the primary key attributes PK in  $R_2$
  - FK in a tuple  $t_1$  of the current instance  $r_1(R_1)$  either has *a value of PK* for some tuple  $t_2$  in the current instance  $r_2(R_2)$  or is *NULL*

# Referential Integrity Constraints

## ◆ Foreign key (FK) (cont'd)

- $t_1[\text{FK}] = t_2[\text{PK}]$  (or NULL)
- ' $t_1$  *refers to*  $t_2$  '

# Referential Integrity Constraints

- ◆ Referential integrity constraint
  - Constraints specifying the foreign key
  - Represented as an *arrow from FK of  $R_1$  to PK of  $R_2$*

# Referential Integrity Constraints

## EMPLOYEE

| Fname | Minit | Lname | <u>Ssn</u> | Bdate | Address | Sex | Salary | Super_ssn | Dno |
|-------|-------|-------|------------|-------|---------|-----|--------|-----------|-----|
|-------|-------|-------|------------|-------|---------|-----|--------|-----------|-----|

## DEPARTMENT

| Dname | <u>Dnumber</u> | Mgr_ssn | Mgr_start_date |
|-------|----------------|---------|----------------|
|-------|----------------|---------|----------------|

## DEPT\_LOCATIONS

| <u>Dnumber</u> | <u>Dlocation</u> |
|----------------|------------------|
|----------------|------------------|

## PROJECT

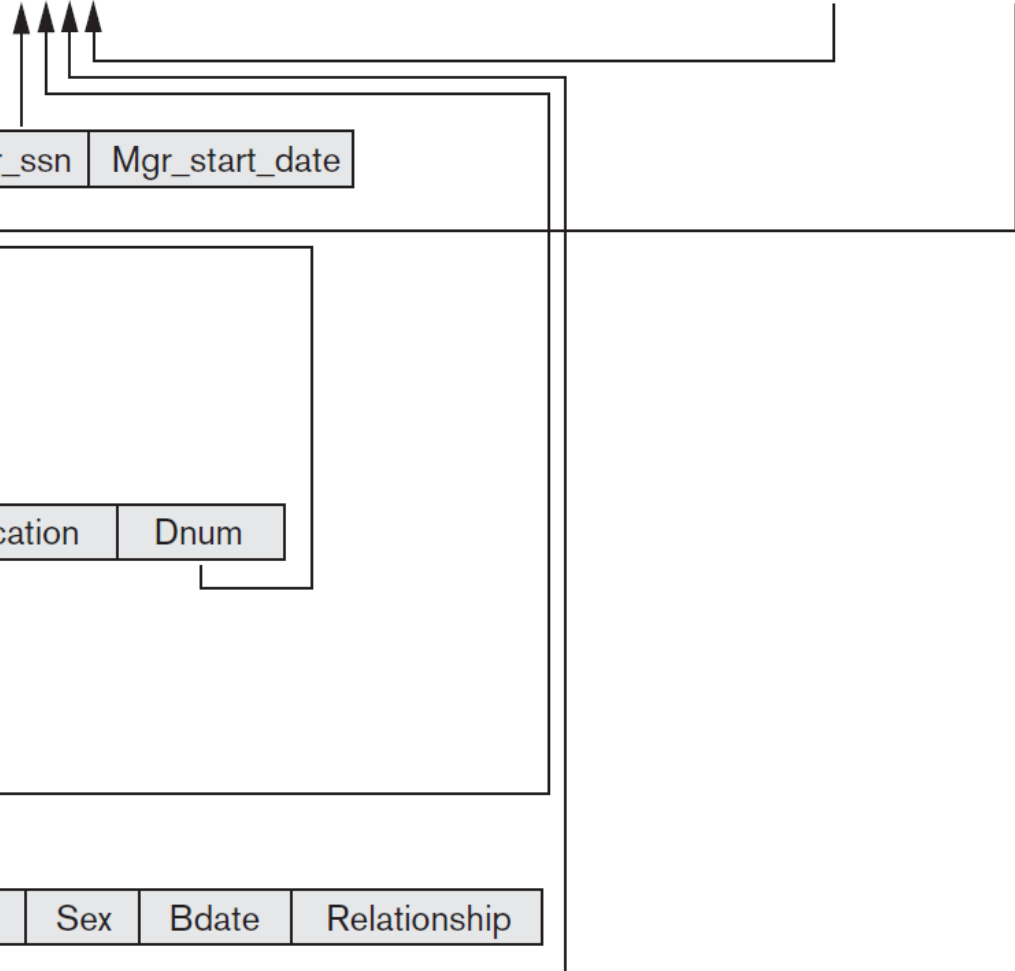
| Pname | <u>Pnumber</u> | Plocation | Dnum |
|-------|----------------|-----------|------|
|-------|----------------|-----------|------|

## WORKS\_ON

| <u>Essn</u> | <u>Pno</u> | Hours |
|-------------|------------|-------|
|-------------|------------|-------|

## DEPENDENT

| <u>Essn</u> | <u>Dependent_name</u> | Sex | Bdate | Relationship |
|-------------|-----------------------|-----|-------|--------------|
|-------------|-----------------------|-----|-------|--------------|



## ◆ Relational Model Concepts

- Relations, tuples, attributes, domains, relational schemas, and relational instances

## ◆ Characteristics of Relations

## ◆ Relational Model Constraints

- Key Constraints
- Integrity constraints
  - Entity integrity constraints
  - Referential integrity constraints

# References



1. Codd, Edgar F. "A relational model of data for large shared data banks." *Communications of the ACM* 13.6 (1970): 377-387.
2. Codd, Edgar F. "A data base sublanguage founded on the relational calculus." *Proceedings of the 1971 ACM SIGFIDET (now SIGMOD) Workshop on Data Description, Access and Control*. ACM, 1971.
3. Codd, Edgar F. *Relational completeness of data base sublanguages*. IBM Corporation, 1972.
4. Codd, Edgar F. "Recent Investigations in Relational Data Base Systems." *IFIP congress*. Vol. 74. 1974.
5. Codd, Edgar F. "Extending the database relational model to capture more meaning." *ACM Transactions on Database Systems (TODS)* 4.4 (1979): 397-434.



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9. Todd, Stephen JP. "The Peterlee relational test vehicle—a system overview." *IBM Systems Journal* 15.4 (1976): 285-308.

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10. Schmid, Hans Albrecht, and J. Richard Swenson. "On the semantics of the relational data model." *Proceedings of the 1975 ACM SIGMOD international conference on Management of data*. ACM, 1975.
11. Chen, Peter Pin-Shan. "The entity-relationship model—toward a unified view of data." *ACM Transactions on Database Systems (TODS)* 1.1 (1976): 9-36.
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13. Atzeni, Paolo, and Valeria De Antonellis. *Relational database theory*. Benjamin-Cummings Publishing Co., Inc., 1993.

Have a nice day!