# **Embedded Operating Systems**

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#### **Outline**

- 1. Introduction to Embedded Operating Systems
- 2. Examples of Embedded Operating Systems
- 3. Multitasking without OS Support

### **Definition of Embedded OS**

- ☐ OSes used in embedded systems
  - Past embedded systems did not use operating systems, but implemented in firmware
  - Modern embedded systems are adopting operating systems
- ☐ Main requirements on embedded OS
  - Small size requirement for limited resources
  - Real-time services
  - Predictability

## Why OS in Embedded Systems?

#### ☐ Technical aspects

- Multitasking
- Networking and communications
- Storage and memory management
- Protection and security
- GUI

#### ☐ Economic aspects

- Time-to-market
- Low NRE (non-recurring engineering) cost
- Maintenance cost
  - 40% 80% in the software lifecycle

## **Typical Embedded Operating Systems**

- ☐ RTOS VxWorks, QNX, pSOS, VRTX, Nucleus, WinCE ☐ Mobile OS iOS, Android, Tizen, WinCE, Symbian, PalmOS □ Open general purpose OS Linux, Solaris 10, FreeBSD, eCOS ☐ DSP OS ARC/MQX, DSP/BIOS □ Others
  - TinyOS (smart sensor), Nachos (education), Xinu (education),
  - IOS (Internetwork OS, Cisco routers)
  - ThreadX (pripherals)

### **Commercial RTOSes**

#### □ Commercial RTOSes

QNX, VxWorks, pSOS, VRTX, Nucleus

#### ☐ Features

- Multithreading, real-time synchronization and POSIX APIs
- Lightweight kernel
  - Even scaled down to fit in the ROM of the system
- Modularized structure
  - Minimum kernel + service components
- Responsiveness
  - Minimal interrupt and switching latency

### RTOS does not mean "fast"

- RTOS adds runtime overheads
- Slightly slower performance than OSless systems

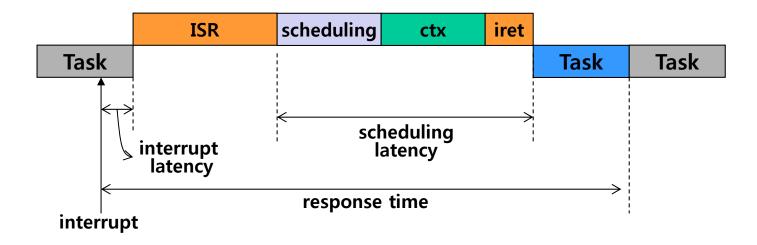
### RTOS does mean "deterministic"

- Guarantees "worst case" times
- Better performance than GPOS

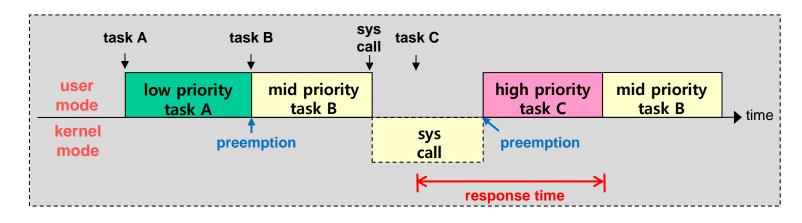
# **Timing Guarantees from RTOS**

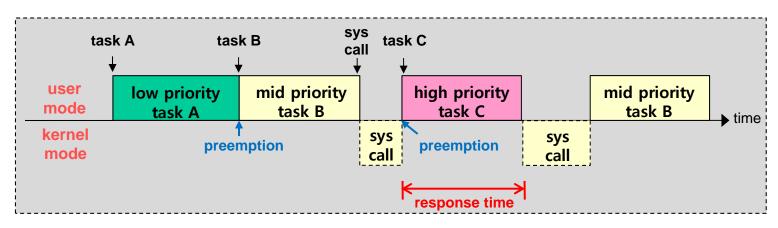
#### ☐ Key requirements

- Fully preemptive priority-based scheduling
  - High priority tasks preempt low ones
- Bounded latency
  - eg.) interrupt off time  $< \alpha$  and scheduling latency  $< \beta$



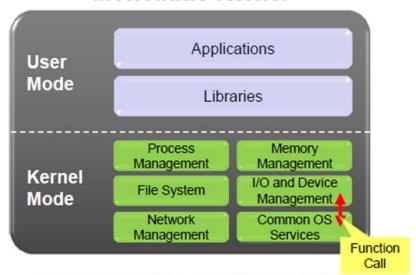
### Partially Preemptive vs. Fully Preemptive





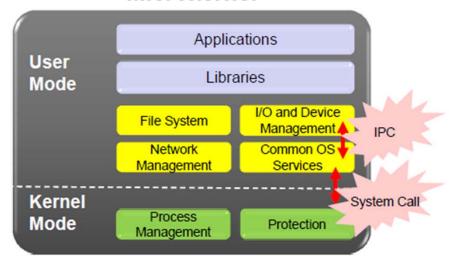
### Monolithic Kernel vs. Microkernel

#### Monolithic Kernel



- Every kernel functionality runs in kernel mode
- Better Performance
- E.g., UNIX, Linux, MS Windows 9x

#### Microkernel



- Most kernel functionality runs in user mode except only the elementary functions
- Better extensibility and modularity
- E.g., MINIX, L4, QNX

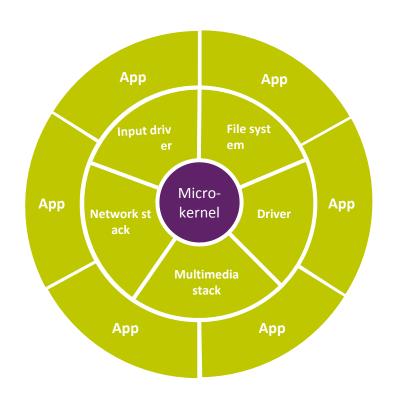
### **QNX Neutrino RTOS**

#### ☐ Microkernel structure

Moves as much from the kernel into "user" space

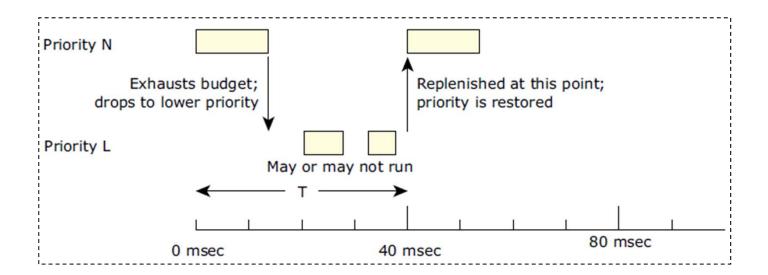
#### □ Rich features

- Momentics C/C++ IDEs
- Photon microGUI
- HMIs such as HTML5 and Qt
- Multimedia and acoustics support
- Mobile connectivity framework and navigation



## **QNX Support for Sporadic Scheduling**

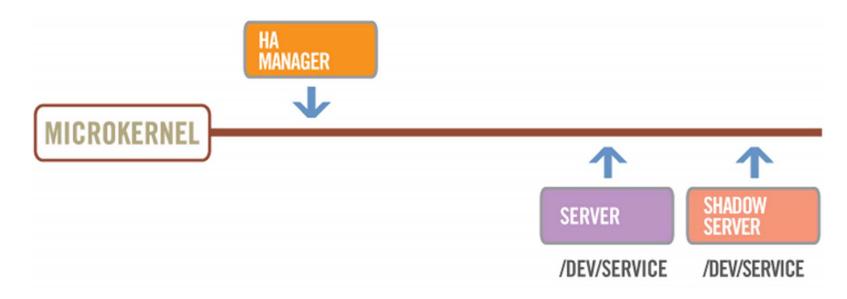
- ☐ Scheduling of real-time periodic tasks
  - Initial budget: the amount of time a task is allowed to run
  - Replenishment period: the period of task execution



# **QNX High Availability Manager**

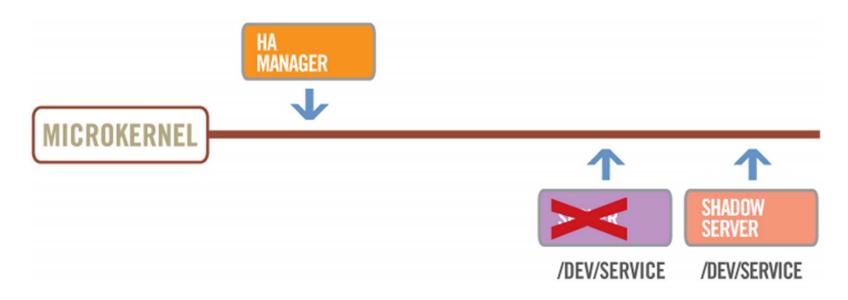
- ☐ HAM (High Availability Manager)
  - Monitor tasks and services (smart watchdog)
  - Perform multistage recovery
- ☐ HAM Hierarchy
  - Entities
    - Units of observation and monitoring
  - Conditions
    - Represent entity states
  - Actions
    - Executed when the appropriate conditions are true with respect to a specific entity

### **HAM Example**



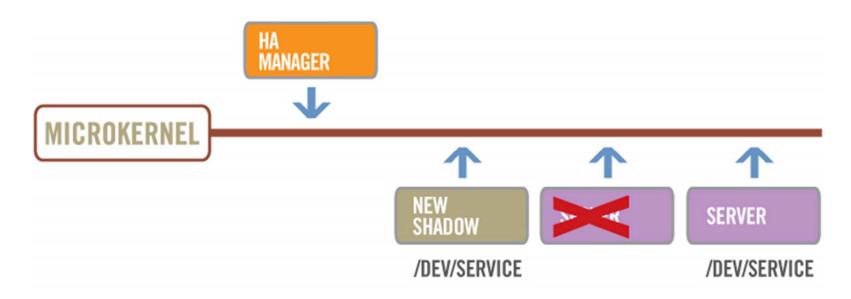
☐ A second "shadow" server runs

### **HAM Example**



- ☐ A second "shadow" server runs
- ☐ If a fault occurs, new clients use the shadow server

## **HAM Example**



☐ Start a new "shadow" server

### **QNX Certification**

☐ Several certified variants for use in products with high criticality and low tolerance for failure

Operating System	Certification/Compliance
QNX OS for Automotive	ISO 26262 ASIL D
QNX OS for Medical	IEC 62304
QNX OS for Safety and Se curity	Common Criteria ISO/IEC 15408 Evaluation Assurance Level (EAL) 4+ and IEC 61508 Safety Integrity Level 3 (SIL 3)
QNX OS for Security	Common Criteria ISO/IEC 15408 Evaluation Assurance Level ( EAL) 4+
QNX OS for Safety	IEC 61508 Safety Integrity Level 3 (SIL 3)

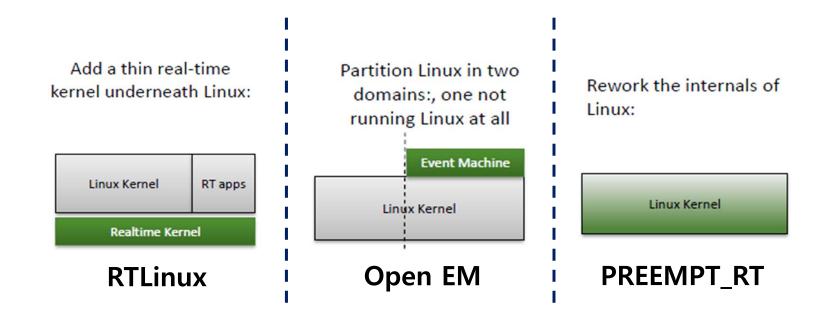
#### **Linux for Real-Time?**

- ☐ Linux is a general purpose OS
  - Aimed at high performance and stability
- ☐ Old Linux did not support
  - Full preemption and bounded latency
- ☐ Current Linux supports
  - ✓ Deterministic scheduler (SCHED\_FIFO and EDF)
  - ✓ Priority inheritance mutexes
  - ✓ Lockable memory (disabling demand paging)
  - ✓ High resolution timers
  - ? Bounded latency inside the kernel
  - ? Boot-up time

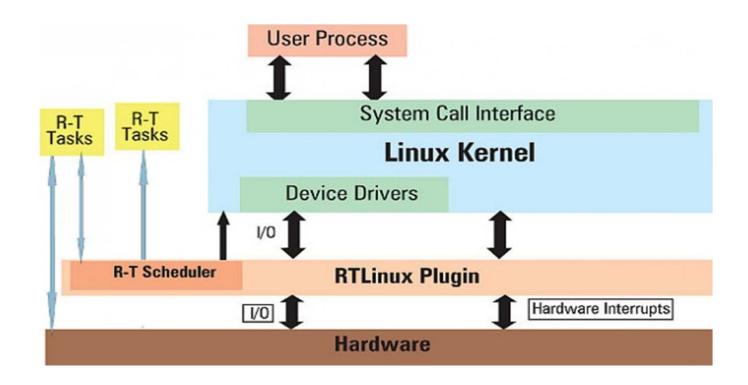
### **Linux Variants for RT**

#### ☐ Extensions for real-time systems

- RTLinux
- Open Event Machine
- PREEMP\_RT patch



### **RTLinux**



- Developed by Victor Yodaiken, Michael Barabanov, Cort Dougan as a commercial product at FSMLabs Wind River Systems acquired FSMLabs in February 2007, but ended it in 2011

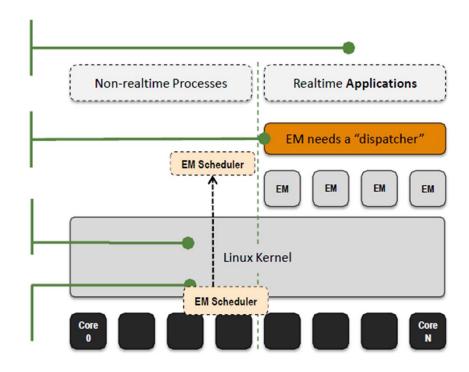
### **Open Event Machine**

EM partitions the system into one RT domain and one non-RT domain

EM uses an event-based model, a run-to-completion model, for individual "context-less" work packages (NO threading or OS model)

Non-essential Linux processes and interrupts are migrated away from the EM cores

EM does not need a special interrupt handling model. Needs a "scheduler" in either Linux partition OR in HW



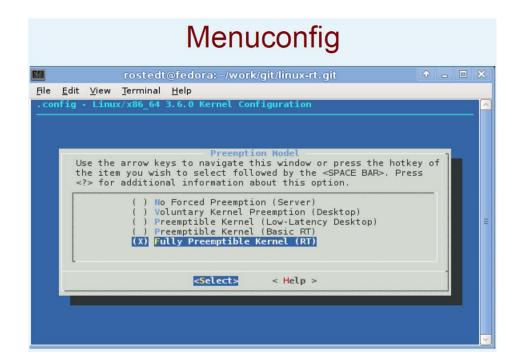
Open EM has been chosen as a starting point for the OpenDataPlaneTM (ODP) initiative by Linaro, the not-for-profit engineering organization developing open source software for the ARM® architecture

### PREEMPT\_RT

#### ☐ PREEMPT\_RT has merged into the mainline kernel

- Since August 2006 by Ingo Molnar et al.
- https://www.kernel.org/pub/linux/kernel/projects/rt/





### PREEMPT\_RT

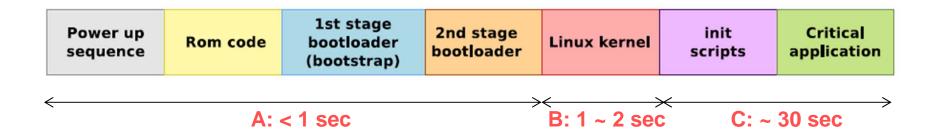
#### ☐ Features

- Preemptible critical sections
- Preemptible interrupt handlers
- Priority inheritance for in-kernel spinlocks and semaphores
- Deferred operations
- Some latency-reduction measures
- ☐ Work is still going on

### **Linux Boot-up Time**

#### ☐ Generic boot sequence

- A. Firmware initialization phase
- B. Kernel initialization phase
- C. User Space initialization phase



#### ☐ Typical boot time

- < 60 seconds: desktops, set top boxes, GPS devices, ...</p>
- < 30 seconds: smartphones</p>
- < 5 seconds: smart TVs</p>

### **Boot-up Time Reduction**

- ☐ General techniques for boot-up time reduction
  - Optimize each job of the boot sequence
  - See backup slides for details
- □ Snapshot-based techniques
  - Save RAM content to nonvolatile storage before power-off
  - Upon power-on, restore RAM state
- ☐ Limitations of snapshot approach
  - Stability problem
  - Retake snapshot every time power is turned off
  - Take snapshot only at "stable" state
  - Android adds another SW layer, increasing snapshot size

#### **GPL License Issue**

- ☐ GPL guarantees end users the freedoms to use, study, share (copy), and modify the software
  - Written by Richard Stallman of the Free Software Foundation (FSF) for the GNU project
- □ Restriction
  - Modified or derived code must be released under GPL
- ☐ LGPL (Lesser GPL)
  - Library linking is allowed without releasing the code
  - Android uses Bionic libc derived from BSD libc to isolate apps from GPL and LGPL

# **Multitasking without OS Support**

# Four Approaches to Multitasking

- ☐ Four approaches
  - Polling-based approach (PA)
  - Event-based approach (EA)
  - Schedule-based approach (SA)
  - Thread-based approach (TA)
- ☐ The first three can be implemented without operating system's support
- ☐ The above approaches can be combined
  - PA + EA, PA + SA, EA + SA, PA + TA ...

# **Polling-based Approach**

### **□** Extremely simple

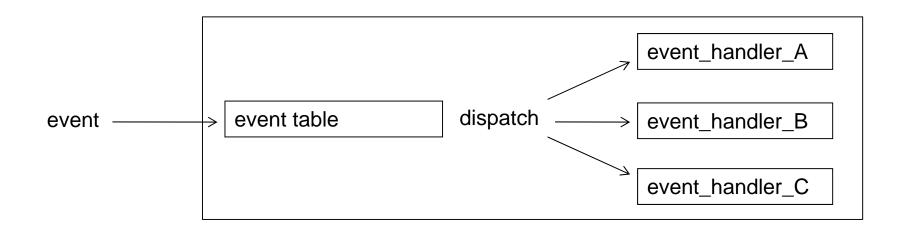
No preemption and no context switch

#### **Problems with PA**

- ☐ PA does not support multi-rate systems
  - Voice recorder samples microphone at 20kHz, samples switches at 15Hz, and updates display at 4Hz
- ☐ Polling frequency is limited by the time required to execute the loop
  - We may obtain more performance by testing more often
    - A/B/A/C/A/B/A/C ...
- ☐ Waiting time may be very long
  - In the worst case, we need to wait for all services to be completed (no preemption)
- ☐ The architecture is fragile
  - Adding a new service will affect timing of all other services
  - Changing rates is tedious and difficult

### **Event-based Approach**

- □ A typical approach for interrupt handling
  - Possible to execute periodic tasks by using timer events
  - Response time is short

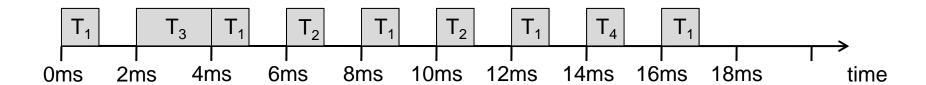


#### **Problems with EA**

- ☐ Waiting time may be long
  - We need to wait for the current job to be completed (no preemption)
- ☐ It is hard to control the execution rate of each task
  - It may depend upon the frequency of event occurring
- ☐ Hard to apply a priority scheme
  - But some hardware supports priorities

## **Schedule-based Approach**

- □ Basically meant for real-time periodic task scheduling
  - Off-line schedule creation
  - On-line task dispatching according to the schedule
  - Uses periodic timer interrupts (time-triggered)
- ☐ Four task scheduling: T<sub>i</sub>(period, execution time) =calendar scheduling
  - $T_1(4, 1), T_2(10, 1), T_3(20, 2), T_4(20, 1)$



### Some Issues with SA

- ☐ Advantages
  - Deterministic behavior
  - Amenable to timing analysis
  - Jitter control
- ☐ We need an algorithm for off-line schedule creation
  - RM (rate monotonic)
    - Tasks with small periods get high priorities
  - EDF (earliest deadline first)
    - Tasks with early deadlines get high priorities
- ☐ We need WCET (worst-case execution time) for each task
  - Schedule must be based on WCETs
  - If not, tasks may overrun

## Cyclic Executive Example

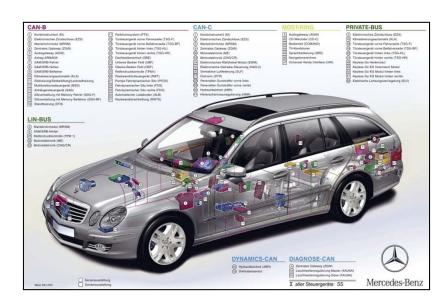
```
#define MAX 10
int schedule[MAX] = \{1,3,1,2,1,2,1,4,1,0\}; <- task
int tick = 0:
/* timer interrupt every 2ms */
void Timer_ISR()
   switch (schedule[tick]) {
          case 0: break;
          case 1: execute_task_1(); break;
          case 2: execute_task_2(); break;
          case 3: execute_task_3(); break;
          case 4: execute_task_4(); break;
   tick = (tick + 1)\%MAX;
```

- ☐ Good for managing the inherent concurrency of an application
- □ Easy to apply a priority scheme or real-time scheduling policy
- ☐ Response time can be short
  - Preemption is supported
  - Blocking can be allowed in a thread, but there is no blocking for the entire application

# AUTomotive Open System ARchitecture

#### **Software Crisis?**

- ☐ Complexity: a lot of ECUs, buses, and gateways
- ☐ Productivity: poor reusability and composability
- ☐ Diversity: diverse OEMs and their suppliers



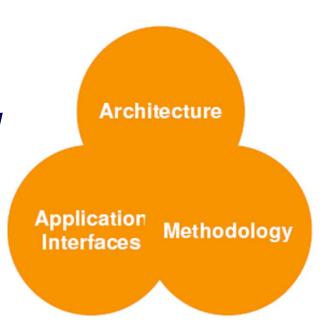
- 55 ECUs
- 7 Buses of 4 types
- 4 Gateways
- More than 10 million LOC

# **AUTomotive Open System ARchitecture**

**AUTOSAR** -

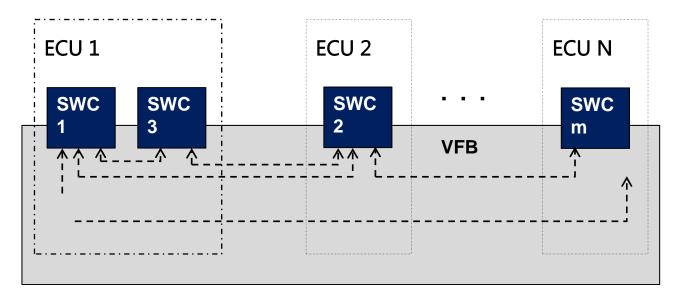
S/W convention

- ☐ Standardized system S/W architecture
  - OSEK kernel
  - OS services and H/W abstractions
- □ Component-based application S/W
  - Standard component interfaces
  - Communication middleware
- Model-based development
  - Authoring tools and XML formats
  - Automatic code generation



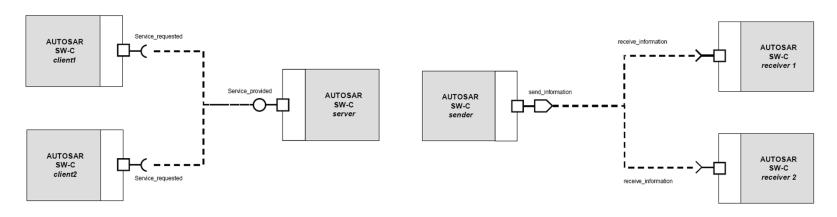
## **SW Components and Virtual Functional Bus**

- **☐ SWC (Software Components)** 
  - Standardized interfaces
  - Reusability and composability
- ☐ VFB (Virtual Functional Bus)
  - Separation between software components and infrastructure



### **SWC** and VFB

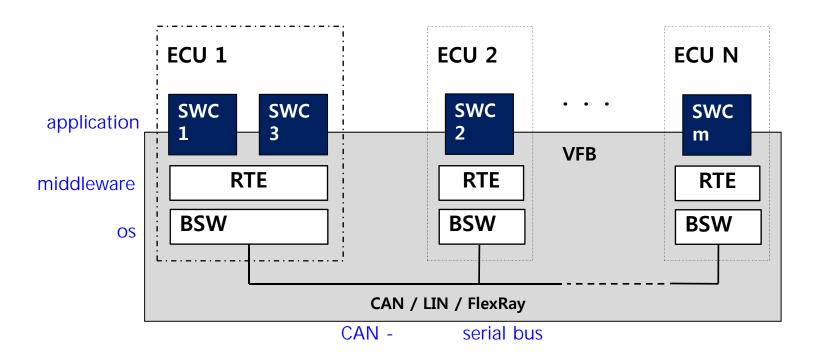
- ☐ Two types of communication
  - **Client-Server Communication**
  - Sender-Receiver Communication



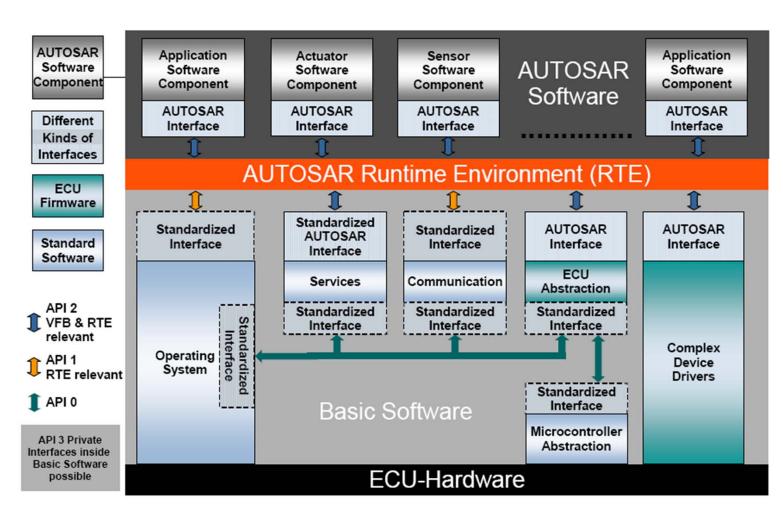
Client-Server Communication Sender-Receiver Communication

### **Runtime Environment and Basic SW**

- □ RTE (Runtime Environment)
  - Middleware for inter- and intra-ECU communication
- ☐ BSW (Basic Software)
  - Traditional OS services



## **Layered Architecture**



microprocessor
- cpu
microcontroller ~ SOC
- cpu+ram+

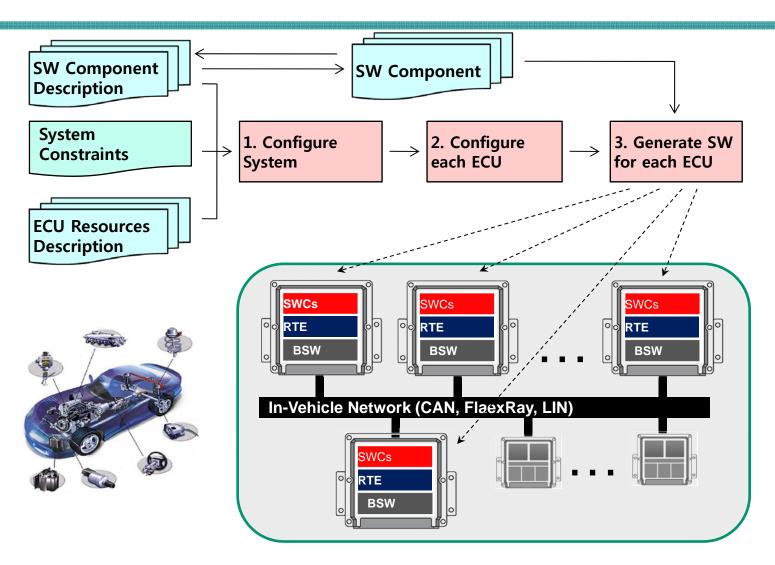
1. vector tool component <-BSW RTE 2. eb tool

s/w

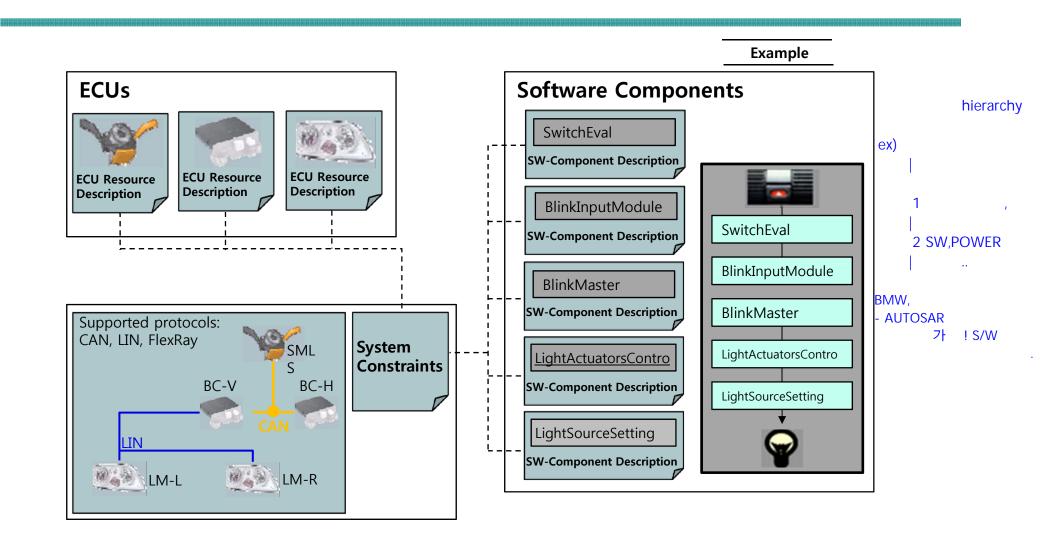
가 s/w

=> BSW/RTE가

compile binrary executable



# **Input Descriptions**



# **Input Descriptions**

#### **SW Component Description**

- General characteristics (name, manufacturer, etc.)
- Communication properties:
  - p\_ports
  - r\_ports
  - interfaces
- Inner structure (composition)
  - sub-components
  - connections
- Required HW resources:
  - processing time
  - scheduling
  - memory (size, type, etc.)

#### **ECU Resource Description**

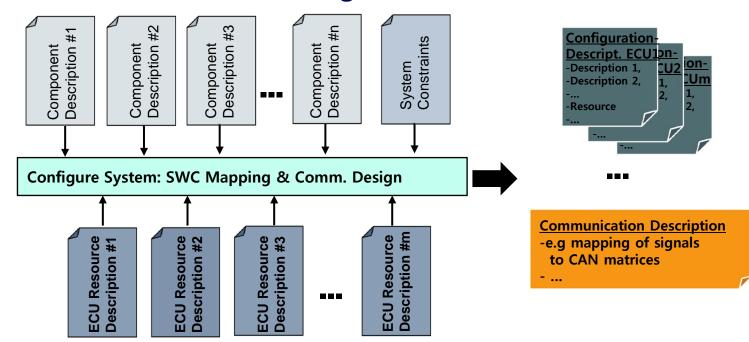
- General characteristics (name, manufacturer, etc.)
- Temperature (own, environment, cooling/heating)
- Available signal processing methods
- Available programming capabilities
- Available HW:
  - uC, architecture
  - Memory
  - Interfaces (CAN, LIN, MOST, FlexRay)
  - Periphery (sensor/actuator)
- SW below RTE for microcontroller
- Signal path from Pin to ECU-abstraction

#### **System Description**

- Network topology
  - bus systems: CAN, LIN, FlexRay
  - connected ECUs, Gateways
  - power supply, system activation
- Communication (for each channel)
  - K-matrix
  - gateway table
- Mapping / Clustering of SW components

# **System Configuration**

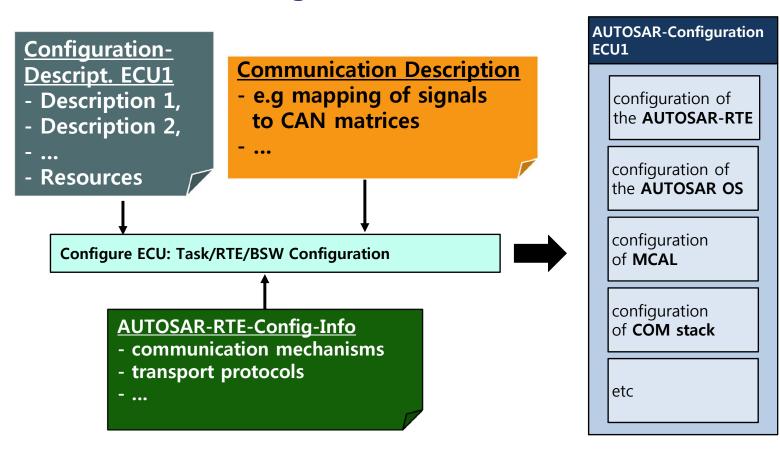
- ☐ SWC mapping
  - Maps the software components to the ECUs
- □ Communication design
  - Completely describes the frames running on the networks and the contents and timing of those frames



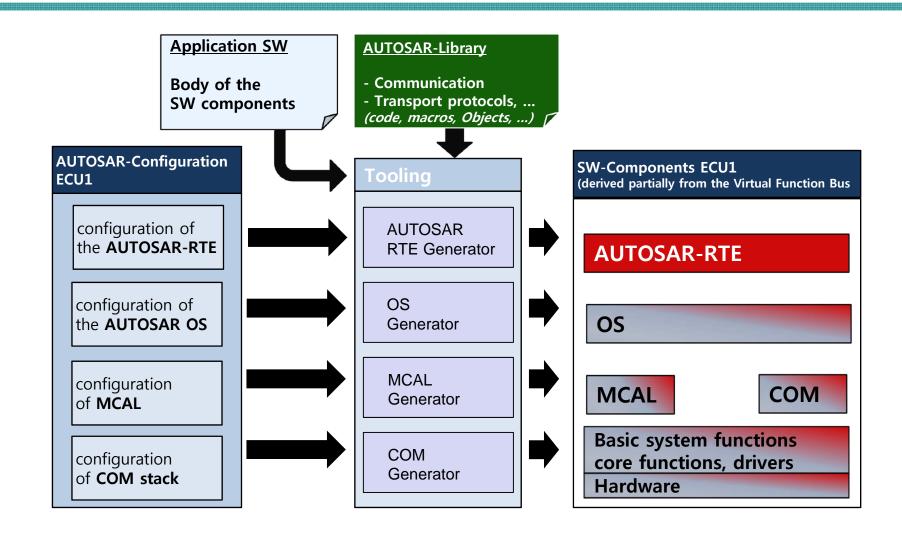
# **ECU Configuration**

component - Runnable <- os(scheduler)가 thread

- □ Runnable-to-task mapping and task scheduling
- ☐ RTE and BSW configuration



#### **Software Generation**



## **Workflow and Tools**

Work	Tools	
1. Control Function Design	<ul> <li>MATLAB, Simulink, Stateflow (Mathworks)</li> <li>→ MIL (Model-in-the-Loop) simulation support</li> </ul>	
2. Code Generation	<ul><li>Real-Time Workshop Embedded Coder (Mathworks)</li><li>TargetLink (dSPACE)</li></ul>	matlab code tran
3. System Architecture Design	<ul> <li>SystemDesk (dSPACE)         → SIL (SW-in-the-Loop) simulation support</li> <li>DaVinci Developer (Vector Informatik)</li> <li>RTA-RTE (ETAS)</li> <li>Volcano Vehicle System Architect (Mentor Graphics)</li> </ul>	·
4. BSW Configuration and Generation	<ul><li>EB tresos Studio and AutoCore (Electrobit)</li><li>MICROSAR (Vector Informatik)</li></ul>	
5. Experimenting, Testing and Debugging	<ul> <li>ControlDesk (dSPACE)</li> <li>NUnit (Vector Informatik)</li> <li>DaVinci Component Tester (Vector Informatik)</li> <li>EB tresos Inspector (Electrobit)</li> </ul>	