Číslo cvičení: 3 Jméno: Marek Bryša

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1. (a) 
$$x^6 - 1 = (1 - x + x^2)(1 + x + x^3 + x^4) \implies h(x) = x^2 - x + 1$$

(b)

- 2. (a) Not linear  $\implies$  not cyclic, not equivalent to a cyclic code.
  - (b) Not a cyclic code, not equivalent to a cyclic code.
  - (c) Is a cyclic code.
  - (d) Not a cyclic code, not equivalent to a cyclic code.

3.

4. MDS  $\iff$   $M=q^{n-d+1}$ . Ham(r,2) has a  $(r\times(2^r-1))=((n-k)\times n)$  parity check matrix  $\implies$   $n=2^r-1, n-k=r, d=3$ .  $M=q^k=2^{n-r}$ .

$$2^{2^{r}-1-3+1} = 2^{2^{r}-1-r}$$
$$2^{r} - 3 = 2^{r} - 1 - r$$
$$r = 2$$

5. 
$$(x^7 - 1)/(x^3 + x + 1) = x^4 + x^2 + x + 1 = h(x)$$
.

$$H = \begin{pmatrix} 1 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 1 \end{pmatrix}, \bar{h}(x) = x^4 + x^3 + x^2 + 1$$

 $6. \ \ x^6-1=(x-1)(x+1)(x^2-x+1)(x^2+x+1). \ \ \text{There are } 2^4=16 \ \text{such codes. Their generator polynomials are: } 1,(x-1),(x+1),(x^2-x+1),(x^2+x+1),(x^2-1),(x^3-2x^2+2x-1),(x^3-1),(x^3+1),(x^3+2x^2+2x+1),(x^4+x^2+1),(x^4+x^3-x+1),(x^4-x^3+x-1),(x^5-x^4+x^3-x^2+x-1),(x^5+x^4+x^3+x^2+x+1),(x^6-1)=0$ 

| Polynomial                        | Matrix |       |   |   |   |   |   |
|-----------------------------------|--------|-------|---|---|---|---|---|
| 1                                 |        | $I_6$ |   |   |   |   |   |
| x+1                               |        | 1     | 1 | 0 | 0 | 0 | 0 |
|                                   |        | 0     | 1 | 1 | 0 | 0 | 0 |
|                                   |        | 0     | 0 | 1 | 1 | 0 | 0 |
|                                   |        | 0     | 0 | 0 | 1 | 1 | 0 |
|                                   |        | 0     | 0 | 0 | 0 | 1 | 1 |
| $x^2 + 1$                         | ĪĪ     | 1     | 0 | 1 | 0 | 0 | 0 |
|                                   |        | 0     | 1 | 0 | 1 | 0 | 0 |
|                                   |        | 0     | 0 | 1 | 0 | 1 | 0 |
|                                   |        | 0     | 0 | 0 | 1 | 0 | 1 |
| $x^{3} + 1$                       | ĪĪ     | 1     | 0 | 0 | 1 | 0 | 0 |
|                                   |        | 0     | 1 | 0 | 0 | 1 | 0 |
|                                   |        | 0     | 0 | 1 | 0 | 0 | 1 |
| $x^4 + x^2 + 1$                   |        | 1     | 0 | 1 | 0 | 1 | 0 |
|                                   |        | 0     | 1 | 0 | 1 | 0 | 1 |
| $(x^5 + x^4 + x^3 + x^2 + x + 1)$ |        | 1     | 1 | 1 | 1 | 1 | 1 |

- 7. (a) Let  $C_1$  be a repetition code of length n,  $C_2$  a no-parity code of length  $n \implies \neg C_1 = C_1, C_1 \cap C_2 = C_1 = C_3$ .
  - (b)  $C_3$  does not exist. There must be at least one non-zero codeword in  $C_1$ . If  $C_3$  were to be cyclic, that non-zero bit would have to be shifted to all positions including the last one.
  - (c) Let  $C_1$  and  $C_2$  be the same as in (a)  $\implies C_1 \cup C_2 = C_2 = C_3$ .