Prog. Lang. Design and Implementation	Homework 2
ENGR 3599 Spring 2014 (Pucella)	Due Date: Friday, Feb 7

This homework is meant to be done individually. You may discuss problems with fellow students, but all work must be entirely your own, and should not be from any other course, present, past, or future. If you use a solution from another source, please cite it. If you consult fellow students, please indicate their names.

Submission information:

- For this problem set, solutions should be put in a file called homework2.sml.
- Your file homework2.sml should be emailed as an attachment before the beginning of class the day the homework is due at the following address:

homeworks.pldi.sp14@gmail.com

• Your file homework2.sml should begin with a block comment that lists your name, your email address, and any remarks that you wish to make about the homework.

Notes:

- All code should compile in the latest release version of Standard ML of New Jersey.
- There is a file homework2.sml available from the web site containing code you can use, including code described in this write-up.
- All questions are compulsory.
- In general, you should feel free to define any helper function you need when implementing a given function. In many cases, that is the cleanest way to go.

This homework uses a variant of the intermediate representation we saw in class that supports single-argument functions and in which primitive operations are built into the representation. (Yes, there is a reason why I restrict to single-argument functions; see Question 1.)

```
datatype value = VInt of int
               | VBool of bool
                    (* the new stuff *)
               | VPair of value * value
               | VList of value list
               | VFun of function
and expr = EVal of value
         | EAdd of expr * expr
         | ESub of expr * expr
         | EMul of expr * expr
         | ENeg of expr
         | EEq of expr * expr
         | EIf of expr * expr * expr
         | ELet of string * expr * expr
         | Eldent of string
         | ECall of string * expr
             (* the new stuff *)
         | EPair of expr * expr
         | EFirst of expr
         | ESecond of expr
         | ESlet of (string * expr) list * expr
         | ECons of expr * expr
         | EIsEmpty of expr
         | EHead of expr
         | ETail of expr
         | ECallE of expr * expr
and function = FDef of string * expr
```

A couple of things to note: we have both new value forms and new expression forms to create and work with those new values. The questions on this homework are about implementing the functionality of those expressions. Because of these additions, the datatypes are now mutually recursive, since they refer to each other in a cycle. (Can you spot the cycle?)

First, look and internalize the provided code—it is just an interpreter structured like every other interpreter weve seen. There are a few places where functions call a function unimplemented that simply raises an exception. Those are the places where I want you to do something, by replacing that function call by something that does the right thing.

Question 1 (Pairs).

This question concerns the new value form

```
| VPair of value * value
```

and the new expression forms

```
| EPair of expr * expr
| EFirst of expr
| ESecond of expr
```

This lets us work with pairs in the object language.

A value VPair (v_1, v_2) is a pair of value v_1 and value v_2 . Values v_1 and v_2 can be arbitrary values, and need not be of the same type. For example,

is a perfectly fine pair whose first component is integer 1 and second component is itself a pair of Boolean true and integer 3.

How do we create pairs in the object language? We can create them directly by using EVal (VPair (v_1, v_2)), or we can use expression form EPair (e_1, e_2) which should evaluate to the pair value made up of the values to which e_1 and e_2 evaluate, respectively. Thus,

```
EPair (EAdd (EVal (VInt 1),
EVal (VInt 2)),
EVal (VInt 4))
```

should evaluate to value VPair (VInt 3, VInt 4).

We only have two operations on pairs, EFirst and ESecond, that are used to extract the first and second components of a pair. Thus,

```
EFirst (EPair (EVal (VInt 1), EVal (VInt 2)))
```

should evaluate to VInt 1, and

```
ESecond (EPair (EVal (VInt 1), EVal (VInt 2)))
```

should evaluate to VInt 2. It is an evaluation error to try to extract the first or second component of a value that is not a pair.

Once we have pairs, as mentioned in class, we can "fake" a two-arguments function by passing the arguments in a pair, a three-arguments function by passing the arguments in a pair whose first component is the first argument to the function and the second component is a pair of the second and third arguments; similarly for functions with more arguments.

For example, here is function exp from lecture, that expects two arguments a and n and computes a^n , modified to take its arguments in a pair:

(it relies on function pred given in lecture, and also provided in the sample code.)

(a) Code a function applyPair with type

```
value -> value -> value
```

where applyPair v1 v2 returns a pair value with v1 and v2 as components.

```
- applyPair (VInt 1) (VInt 2);
val it = VPair (VInt 1,VInt 2) : value
- applyPair (VInt 1) (VBool true);
val it = VPair (VInt 1,VBool true) : value
- applyPair (VInt 1) (VPair (VInt 1, VInt 2));
val it = VPair (VInt 1,VPair (VInt #,VInt #)) : value
```

(b) Code a function applyFirst with type

```
value -> value
```

where applyFirst v returns the first component of v if it is a pair, and returns an error otherwise.¹

¹Use function evalError provided in the code to create the error.

```
- applyFirst (VPair (VInt 1, VInt 2));
val it = VInt 1 : value
- applyFirst (VPair (VBool true, VBool false));
val it = VBool true : value
- applyFirst (VPair (VPair (VInt 1, VInt 2), VPair (VInt 3, VInt 4)));
val it = VPair (VInt 1, VInt 2) : value
- applyFirst (VInt 1);
uncaught exception Fail [Fail: Eval Error @ applyFirst]
  raised at: solution2.sml:43.28-43.54
- applyFirst (VBool true);
uncaught exception Fail [Fail: Eval Error @ applyFirst]
  raised at: solution2.sml:43.28-43.54
```

(c) Code a function applySecond with type

value -> value

where applyFirst v returns the second component of v if it is a pair, and returns an error otherwise.

```
- applySecond (VPair (VInt 1, VInt 2));
val it = VInt 2 : value
- applySecond (VPair (VBool true, VBool false));
val it = VBool false : value
- applySecond (VPair (VPair (VInt 1, VInt 2), VPair (VInt 3, VInt 4)));
val it = VPair (VInt 3, VInt 4) : value
- applySecond (VInt 1);

uncaught exception Fail [Fail: Eval Error @ applySecond]
  raised at: solution2.sml:43.28-43.54
- applySecond (VBool true);

uncaught exception Fail [Fail: Eval Error @ applySecond]
  raised at: solution2.sml:43.28-43.54
```

- (d) Complete the cases EPair, EFirst, and ESecond of substitution function subst.
- (e) Complete the cases EPair, EFirst, and ESecond of evaluation function eval. You will probably want to use the functions in (a)-(c).

```
- eval [] (EFirst (EPair (EVal (VInt 1),
                          EVal (VInt 2))));
val it = VInt 1 : value
- eval [] (ESecond (EPair (EVal (VInt 1),
                           EVal (VInt 2))));
val it = VInt 2 : value
- eval [] (EFirst (ELet ("x", EVal (VInt 1),
                         EPair (EIdent "x", EVal (VInt 2))));
val it = VInt 1 : value
- exp;
val it = ("exp",FDef ("args",ELet (#,#,#))) : string * function
- pred;
val it = ("pred",FDef ("n",ESub (#,#))) : string * function
- eval [exp, pred] (ECall ("exp", EPair (EVal (VInt 4),
                                         EVal (VInt 9)));
val it = VInt 262144 : value
```

Question 2 (Simultaneous Bindings).

This question concerns the new expression form

```
| ESlet of (string * expr) list * expr
```

This lets us define *simultaneous bindings*. In some kind of proto-syntax, a simultaneous binding is of the form:

```
slet x = \dots

y = \dots

z = \dots

in ...
```

(where of course the ... stand for actual expressions).

A simultaneous binding first evaluates all the expressions to be bound to the identifiers, and *then* binds the resulting values to those identifiers. This means, in particular, that the following code (still in our proto-syntax) would evaluate to (20,10):

Spend a few seconds to see why that is different from the following code, where the bindings for \mathbf{x} and \mathbf{y} are *not* simultaneous:

```
let x = 10
let y = 20
let x = y
let y = x
in (x,y)
```

This last expression evaluates to (20,20).

A simultaneous binding

slet
$$x = 10$$

 $y = 20$
in e

is expressed in our internal language as

```
ESlet ([("x", EVal (VInt 10)), ("y", EVal (VInt 20))],
```

Thus, you see that ESlet takes a list of bindings (each binding a pair of an identifier name as a string and an expression to evaluate and bind to the identifier) and an expression representing the body of the slet.

(a) Complete the case ESlet of substitution function subst.

Substituting into an ESlet requires some thought. Here is the pseudo-code for substituting into an ESlet. It is very similar to substituting into an ELet, except that we need to substitute into all the expressions in the bindings.

```
To substitute id with e in ESlet (bdgs,body):

substitute id with e in every expression in bindings bdgs

if id is not one of the identifiers bound in bdgs:

substitute id with e in body
```

(b) Complete the case ESlet of evaluation function eval.

```
- eval [] (ESlet ([("x", EVal (VInt 1)),
                   ("y", EVal (VInt 2))],
                  EAdd (EIdent "x", EIdent "y")));
val it = VInt 3 : value
- eval [] (ESlet ([("x", EVal (VInt 20)),
                   ("y", EVal (VInt 10))],
                  ESlet ([("x", EAdd (EIdent "x", EIdent "y")),
                           ("y", ESub (EIdent "x", EIdent "y"))],
                         EPair (EIdent "x", EIdent "y"))));
val it = VPair (VInt 30, VInt 10) : value
- swap;
val it = ("swap", FDef ("args", ELet (#, #, #))) : string * function
- eval [swap] (ECall ("swap", EVal (VPair (VInt 10, VInt 20))));
val it = VPair (VInt 20, VInt 10) : value
- eval [] (ESlet ([("x", EVal (VInt 1)),
                   ("y", Eldent "x")],
                  EAdd (EIdent "x", EIdent "y")));
uncaught exception Fail [Fail: Eval Error @ eval/Eldent - x]
  raised at: solution2.sml:43.28-43.54
```

Question 3 (Lists).

This question concerns the new value form

```
| VList of value list
```

and the new expression forms

```
| ECons of expr * expr
| EIsEmpty of expr
| EHead of expr
| ETail of expr
```

This lets us work with lists in the object language.

A value $VList\ L$ is a list L of values. Those values in L can be arbitrary, and need not be of the same type. For example,

```
VList [VInt 1,
VBool true,
VList [],
VList [VInt 2]]
```

is a perfectly fine list of four elements, the last two of which are themselves lists.

How do we create lists in the object language? We can create them directly by using EVal (VList L), or we can use expression form ECons (e_1 , e_2) which roughly corresponds to the :: operation in Standard ML: it creates a new list in which the result of evaluating e_1 is the first element, and the rest of the elements are from the list obtained by evaluating e_2 . Thus,

should evaluate to value VList [VInt 3, VInt 1, VInt 2]

There are three operations on lists. Operation EIsEmpty checks if a list is empty. It takes an expression, and if the result of evaluating the expression is an empty list, it returns true, otherwise it returns false (as values in the object language, obviously). Thus,

```
EIsEmpty (ECons (EVal (VInt 1), EVal (VList [])))
```

should evaluate to VBool false.

Operations EHead and ETail extract out the first element of a list and the list made up of all but the first element, respectively. Thus,

```
should evaluate to VInt 1, while

EHead (EVal (VList [VInt 1, VInt 2, VInt 3]))

should evaluate to VList [VInt 2, VInt 3].

Here is a sample recursive function to compute the length of a list:

("length",
FDef ("xs",
EIf (EIsEmpty (EIdent "xs"),
EVal (VInt 0),
EAdd (EVal (VInt 1),
ECall ("length", ETail (EIdent "xs"))))))
```

EHead (EVal (VList [VInt 1, VInt 2, VInt 3]))

Here is another that appends two lists, using pairs from Question 1 to pass the two lists as arguments:

(a) Code a function applyCons with type

```
value -> value -> value
```

where applyCons v1 v2 returns a list value with v1 as first element of the list followed by the elements of v2. It should return an error if the second argument is not a list.

```
- applyCons (VInt 1) (VList [VInt 2, VInt 3, VInt 4]);
val it = VList [VInt 1,VInt 2,VInt 3,VInt 4] : value
- applyCons (VBool true) (VList [VInt 1, VBool false, VInt 2]);
val it = VList [VBool true,VInt 1,VBool false,VInt 2] : value
- applyCons (VList []) (VList []);
val it = VList [VList []] : value
- applyCons (VList [VInt 1, VInt 2]) (VList [VInt 3, VInt 4]);
val it = VList [VList [VInt #,VInt #],VInt 3,VInt 4] : value
- applyCons (VInt 1) (VInt 2);
uncaught exception Fail [Fail: Eval Error @ applyCons]
raised at: solution2.sml:43.28-43.54
```

(b) Code a function applyIsEmpty with type

```
value -> value
```

where applyIsEmpty v returns VBool true if v is an empty list, and VBool false otherwise.

```
- applyIsEmpty (VList []);
val it = VBool true : value
- applyIsEmpty (VList [VInt 1]);
val it = VBool false : value
- applyIsEmpty (VList [VInt 1, VInt 2]);
val it = VBool false : value
- applyIsEmpty (VInt 1);
val it = VBool false : value
- applyIsEmpty (VBool true);
val it = VBool false : value
```

(c) Code a function applyHead with type

```
value -> value
```

where applyHead v returns the first element of v if it is a non-empty list, and returns an error otherwise.

```
- applyHead (VList [VInt 1, VInt 2, VInt 3]);
val it = VInt 1 : value
- applyHead (VList [VBool true, VBool false]);
val it = VBool true : value
- applyHead (VList []);
```

```
uncaught exception Fail [Fail: Eval Error @ applyHead]
  raised at: solution2.sml:43.28-43.54
- applyHead (VInt 1);
uncaught exception Fail [Fail: Eval Error @ applyHead]
  raised at: solution2.sml:43.28-43.54
```

(d) Code a function applyTail with type

```
value -> value
```

where applyTail v returns the list made up of all the elements of v except the first if v is a non-empty list, and returns an error otherwise.

```
- applyTail (VList [VInt 1, VInt 2, VInt 3]);
val it = VList [VInt 2,VInt 3] : value
- applyTail (VList [VInt 1]);
val it = VList [] : value
- applyTail (VList [VBool true, VBool false]);
val it = VList [VBool false] : value
- applyTail (VList []);
uncaught exception Fail [Fail: Eval Error @ applyTail]
  raised at: solution2.sml:43.28-43.54
- applyTail (VInt 1);
uncaught exception Fail [Fail: Eval Error @ applyTail]
  raised at: solution2.sml:43.28-43.54
```

- (e) Complete the cases ECons, EIsEmpty, EHead, and ETail of substitution function subst.
- (f) Complete the cases ECons, EIsEmpty, EHead, and ETail of evaluation function eval. You will probably want to use the functions in (a)-(d).

```
val it = VList [VInt 1,VInt 2] : value
- eval [] (EHead (ECons (EVal (VInt 1),
                         EVal (VList [VInt 2])));
val it = VInt 1 : value
- eval [] (ETail (ECons (EVal (VInt 1),
                         EVal (VList [VInt 2])));
val it = VList [VInt 2] : value
- eval [] (EIsEmpty (ECons (EVal (VInt 1),
                            EVal (VList [VInt 2])));
val it = VBool false : value
- eval [] (EIsEmpty (ETail (ECons (EVal (VInt 1),
                                   EVal (VList []))));
val it = VBool true : value
- eval [] (ELet ("x", EVal (VInt 1),
                 ELet ("ys", EVal (VList [VInt 2, VInt 3]),
                       ECons (EIdent "x", EIdent "ys"))));
val it = VList [VInt 1, VInt 2, VInt 3] : value
- length;
val it = ("length",FDef ("xs",EIf (#,#,#))) : string * function
- eval [length] (ECall ("length",
                        EVal (VList [VInt 10, VInt 20, VInt 30]));
val it = VInt 3 : value
- append;
val it = ("append",FDef ("args",ELet (#,#,#))) : string * function
- eval [append] (ECall ("append",
                        EVal (VPair (VList [VInt 1, VInt 2, VInt 3],
                                     VList [VBool true, VBool false]))));
val it = VList [VInt 1, VInt 2, VInt 3, VBool true, VBool false] : value
```

Question 4 (Functions as Arguments).

This question concerns the new value form

```
| VFun of function
```

and the new expression form

```
| ECallE of expr * expr
```

Together, these provide a way to call functions passed as arguments, the way we can do it in Standard ML (and other languages).

A value $VFun\ f$ represents a function. Not a function name—a bona fide function. For example,

```
VFun (FDef ("n", EAdd (EIdent "n", EVal (VInt 1))))
```

is the function that takes an argument n and adds 1 to it.

How do we create such function values in the object language? We can create them directly using $VFun\ f$ as above, or we can look them up in the function environment. The idea is that an idenfier (an expression EIdent) that has not been substituted for evaluates down to a function definition, if that identifier's name is a function name in the function environment. For example, if we have a function environment

```
[ ("pred", FDef ("n", ESub (EIdent "n", EVal (VInt 1)))) ]
```

then the expression

```
Eldent "pred"
```

should evaluate to VFun (FDef ("n", ESub (EIdent "n", EVal (VInt 1)))).

How do we use such functions? We use expression ECallE, which intuitively works just like ECall, except that instead of determining which function to call by looking up the name of the function in the function environment, it simply calls the function to which its first argument evaluates. Thus, for example, and in the context of the sample function environment above containing pred,

```
ECallE (EIdent "pred",
EVal (VInt 10))
```

should evaluate to VInt 9.

More interestingly, this new expression form allows us to write functions that expect functions as arguments and apply them to other values. For instance, the following function twice, described in lecture, takes a function f as argument and a value v and returns f (f v):

Note that it uses a pair to accept two arguments, as seen in Question 1.

The following function takes a function and a list as arguments and produces a new list made up of the results of applying the function to every element of the list.

- (a) Complete the case ECallE of substitution function subst.
- (b) Complete the case ECallE and EIdent of evaluation function eval.

```
- succ;
val it = ("succ",FDef ("n",EAdd (#,#))) : string * function
- eval [succ] (EIdent "succ");
val it = VFun (FDef ("n",EAdd (#,#))) : value
- eval [] (EIdent "nonexistent");
```

```
uncaught exception Fail [Fail: Eval Error @ eval/Eldent - nonexistent]
 raised at: solution2.sml:43.28-43.54
- eval [succ] (ECallE (EIdent "succ", EVal (VInt 10)));
val it = VInt 11 : value
- eval [succ] (ECallE (EIdent "succ",
                       ECallE (EIdent "succ", EVal (VInt 10)));
val it = VInt 12 : value
- eval [succ] (ELet ("f", EIdent "succ",
                     ECallE (EIdent "f", EVal (VInt 15))));
val it = VInt 16 : value
- twice;
val it = ("twice",FDef ("args",ELet (#,#,#))) : string * function
- eval [twice, succ] (ECall ("twice", EPair (EIdent "succ",
                                             EVal (VInt 25))));
val it = VInt 27 : value
- pred;
val it = ("pred",FDef ("n",ESub (#,#))) : string * function
- eval [twice, pred] (ECall ("twice", EPair (EIdent "pred",
                                             EVal (VInt 25))));
val it = VInt 23 : value
- mapf;
val it = ("mapf",FDef ("args",ELet (#,#,#))) : string * function
- eval [mapf,succ]
     (ECall ("mapf",
             EPair (EIdent "succ",
                    EVal (VList [VInt 10, VInt 20, VInt 30])));
val it = VList [VInt 11, VInt 21, VInt 31] : value
- eval [mapf]
     (ECall ("mapf",
             EPair (EVal (VFun (FDef ("n",
                                      EMul (EIdent "n", EVal (VInt 2)))),
                    EVal (VList [VInt 10, VInt 20, VInt 30])));
val it = VList [VInt 20, VInt 40, VInt 60] : value
```