# COMP30023 Project 2 Remote Procedure Call

Out date: 28 April 2023

Due date: No later than 5pm Monday 22 May, 2023 AEST

Weight: 15% of the final mark

### **Project Overview**

Remote Procedure Call (RPC) is a crucial technology in distributed computing that enables software applications to communicate with each other seamlessly over a network. It provides a way for a client to call a function on a remote server as if it were a local function call. This abstraction allows developers to build distributed systems and applications that span In this project, you will be it seamlessly between multiple computers. This

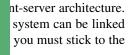
will still apply. Your RPC system must be the one you have been p using existing RPC librari ng principles of RPC

inux cloud VM, like n RPC code, without

## **RPC System Arc**

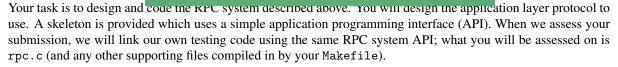
Your task is to design and The RPC system will be to either a client or a serv proposed API carefully.

For testing purposes, you



ur VM).

### **Project Details**



Note that implementing the API will require you to use sockets. This uses material covered in the lectures after the project is released. There is plenty to do on the project before you need to use sockets, so please do not say that you cannot start because you have not yet learned about sockets.

#### 3.1 API

The basic API implemented by rpc.c consists of the following data structures and functions.

#### 3.1.1 Data structures

The API will send and receive data structures of the form:

```
typedef struct {
    int data1;
    size_t data2_len;
    void *data2;
} rpc_data;
```

where data1 is an integer to be passed to the other side, and data2 is a block of bytes (of length data2\_len) to be sent. The purpose of data1 is to allow simple functions that only pass an integer to avoid memory management issues, by setting data2\_len=0 and data2=NULL. Your protocol can limit data1 to being no more than 64 bits.

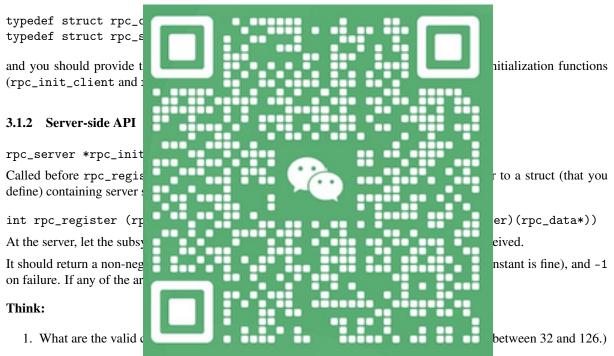
Note that size\_t depends on the architecture, and the sender and receiver can have different architectures. Think how this will affect your protocol.

The handler that implements the actual remote procedure will have the signature:

```
rpc_data *procedure (rpc_data *d);
```

That is, it takes a pointer to an rpc\_data object and returns a pointer to another rpc\_data object. This function will dynamically allocate memory with malloc for both the rpc\_data structure and its data2 field. It is the responsibility of the RPC system to free those after use.

The state of the client and server will be in data structures that you define. These are declared in rpc.h as:



- 2. How does the server know the length of name? (We will not test names longer than 1000 bytes.)
- 3. Should there be a minimum length for name? (We will not test for empty names.)

If there is already a function registered with name name, then the old function should be forgotten and the new one should take its place.

To get full marks, you should be able to register at least 10 functions. You can still get most of the marks as long as you can register one function, so implement that first.

```
void rpc_serve_all (rpc_server *srv)
```

This function will wait for incoming requests for any of the registered functions, or rpc\_find, on the port specified in rpc\_init\_server of any interface. If it is a function call request, it will call the requested function, send a reply to the caller, and resume waiting for new requests. If it is rpc\_find, it will reply to the caller saying whether the name was found or not, or possibly an error code.

This function will not usually return.

It should only return if srv is NULL or you're handling SIGINT (not a requirement).

#### 3.1.3 Client-side API

rpc\_client \*rpc\_init\_client (const char \*addr, int port)

Called before rpc\_find or rpc\_call. Use this for whatever you need. The string addr and integer port are the text-based IP address and numeric port number passed in on the command line.

The function should return a non-NULL pointer to a struct (that you define) containing client state information on success and NULL on failure.

void rpc\_close\_client (rpc\_client \*cl)

Called after the final rpc\_call or rpc\_find (i.e., any use of the RPC system by the client).

Use this for whatever you need; it should at least free(cl).

If it is (mistakenly) called on a client that has already been closed, or cl == NULL, it should return without error. (Think: How can you tell if it has already been closed?)

rpc\_handle \*rpc\_find (rpc\_client \*cl, const char \*name)

At the client, tell the subsystem what details are required to place a call. The return value is a handle (not handler) for the remote procedure, which is passed to the following function.

If name is not registered, it should return NULL. If any of the arguments are NULL then NULL should be returned. If the find operation fails, it a) rpc\_data \*rpc\_call This function causes the s If the call fails, it returns ULL. If this returns a non-NULL value, then it s ucture and its data2 field. The client will free The skeleton gives an exa 3.1.4 Shared API void \*rpc\_data\_free Frees the memory allocate Note that there is a referen 3.2 Planning task When you are designing the wer for this particular project, and separately for

Put the answers in a *plain text* file answers.txt, beginning with your name, login ID and student ID. These, together with protocol description below, are worth **1 mark**.

- 1. Should the server accept calls from everyone, or just a subset of users?
- 2. Should authentication etc. be provided by the RPC framework, or by the functions that use the RPC framework?
- 3. What transport layer protocol should be used? What are the trade-offs?
- 4. In which function(s) should the socket(s) be created?
- 5. Should rpc\_client and rpc\_server be allocated dynamically or statically? What are the implications for the client and server code?
- 6. What happens if one host uses big-endian byte order and the other uses little-endian? How does that relate to "network byte order"?

#### 3.3 Protocol

You should design and document a simple application layer protocol for this RPC system. (It can be very simple.) Describe it in the file answers.txt, in enough detail that someone else could implement the protocol. Together with the questions above, this is worth **1 mark**.

Note that size\_t is system-dependent. You will need to choose a system-independent way to encode the size of the data block in the packet. You can use a universal encoding, like Elias gamma coding (see Wikipedia) which can specify arbitrarily long strings, or you can use a fixed sized field, which is simpler to decode but limits the size of the string. *Explain your choice in your protocol description*.

In all test cases, len < 100 000.

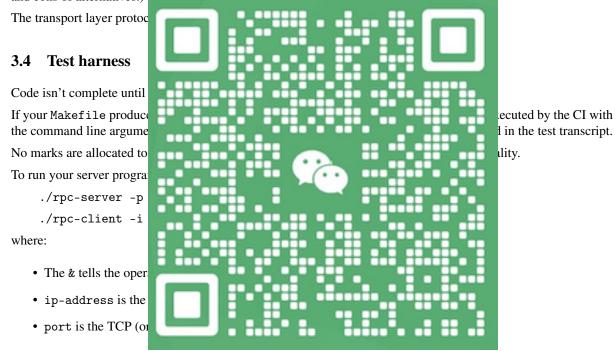
If data2\_len is too large to be encoded in your packet format, the relevant function in rpc.c should print "Overlength error" to stderr and return an error.

The protocol should specify error responses for routine failures, such as a request for a procedure that does not exist.

The protocol should work correctly even if there are IP layer packet loss and duplication.

The protocol should handle the fact that IP packets have a maximum allowed size.

Decide what transport layer protocol to use. (You will almost certainly choose TCP, but briefly mention the pros and cons of alternatives.)



The server is expected to listen for incoming connections on the port passed via command line arguments, on any of the hosts IPv6 network addresses.

#### 4 Stretch goal: non-blocking performance

Many RPC operations, such as database look-ups, take a substantial time to complete. Instead of making each request wait for all those before it to complete, it is possible to continue to accept and start processing new requests while previous requests are executing, with each call returning as soon as it is finished. The simplest way to do this is with multiple threads: each time a request (or connection) is received, a new thread is spawned to execute the procedure, and is destroyed once the result has been sent back to the caller. Alternatives include creating new processes with fork(2), or using select(2).

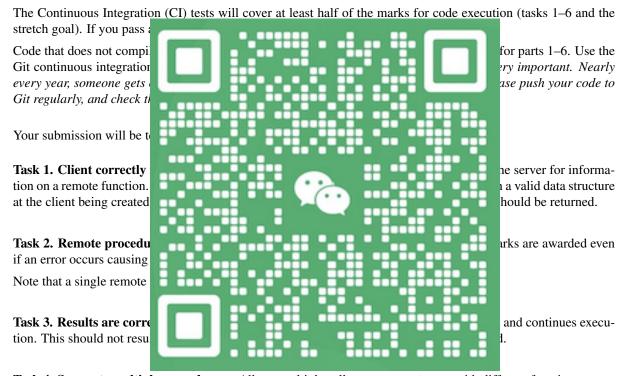
If you implement non-blocking, include the line "#define NONBLOCKING" in your code. Otherwise, this functionality will not be tested/marked.

#### 5 Marking Criteria

The marks are broken down as follows.

The column "In CI" specifies how many of the marks allocated for this are tested by the continuous integration system available before submission. If you pass these in CI, you can be confident of getting those marks. There are also tests that will only be run on your final submission, which give the remaining marks.

Task # and description	Marks	In CI
Client correctly finds module on server	2	1
2. Remote procedure is called correctly	2	1
3. Results are correctly returned to client	2	1
4. Supports multiple procedures	2	1
5. Portability and safety	2	1
6. Build quality	1	1
7. Quality of software practices	2	0
8. Planning task and protocol description, in answers.txt	1	0
9. Stretch goal. Non-blocking operation works	1	0.5



**Task 4. Supports multiple procedures** Allows multiple calls to rpc\_register with different function names. The protocol must indicate which function is being called by rpc\_call, and NULL if an unregistered function is passed to rpc\_find or an invalid handle is passed to rpc\_call.

**Task 5. Portability and safety** The server component of the RPC system return errors on failure and does not crash after listen(2), e.g., due to malformed input, unexpected termination of connection etc.

The client component of the RPC system return errors on failure, e.g. cannot connect to server, if the server shuts down in the middle of an operation etc.

If you implement timeouts (not a requirement), the timeout duration must be longer than 30 seconds.

Data is sent in network byte order, regardless of whether the client and server are big-endian or little-endian.

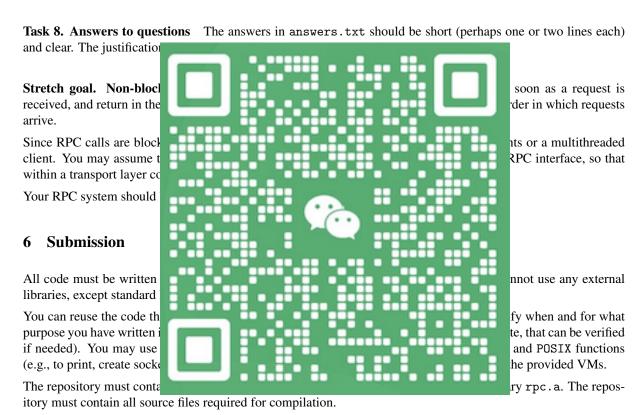
(Note that the testing system will check this by snooping on packets. Encryption or compression will break this. If you use either of those, please comment it clearly in your code so that it can be tested by hand.)

Task 6. Build quality Running make clean && make -B should produce an object file rpc.o or static library rpc.a, which contains everything needed for the RPC system, and will be linked to our test client and server code. Include in your Makefile (in 1 line) a LDFLAGS variable describing any linker flag(s) that your submission requires.

If this fails for any reason, you will be told the reason, and be allowed to resubmit (with the usual late penalty). If it still fails, you will get 0 for Tasks 1–5 and the stretch goal. Test this by committing regularly, and checking the CI feedback. (If you need help, ask on the forum.)

A 0.5 mark penalty will be applied if compiling using "-Wall" yields a warning or if your final commit contains any executable, .o, or .a files.

**Task 7. Quality of software practices** Factors considered include **quality of code**, based on the choice of variable names, comments, formatting (e.g. consistent indentation and spacing), structure (e.g. abstraction, modularity), use of global variables, proper management of memory including not leaking memory, and **proper use of version control**, based on the regularity of commit and push events, their content and associated commit messages. Profanity or abuse in the commit messages will also result in mark deductions; everyone gets frustrated, but commits must remain professional.



Place the Makefile at the root of your repository, and ensure that running make places the requested files there too. If the GitLab CI does not find these files, the marking system will not either.

Ensure that your RPC system does not write to stdout. However, the client and server can.

The server component of the RPC system should not shut down by itself. SIGINT (like CTRL-C) will be used to terminate the server between test cases. You may notice that a port and interface which has been bound to a socket sometimes cannot be reused until after a timeout. To make your testing and our marking easier, please override this behaviour by placing the following lines before the bind() call:

```
int enable = 1;
if (setsockopt(sockfd, SOL_SOCKET, SO_REUSEADDR, &enable, sizeof(int)) < 0) {
    perror("setsockopt");
    exit(EXIT_FAILURE);
}</pre>
```

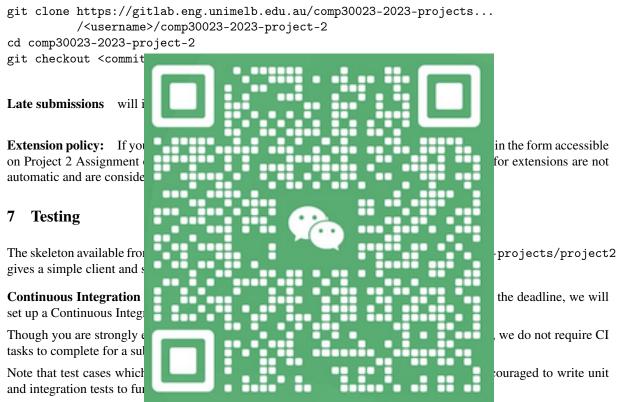
Make sure that all source code is committed and pushed to git. **Executable files** (that is, all files with the executable bit which are in your repository) **will be removed** before marking. Hence, ensure that none of your source files have the executable flag set. (You can verify this by cloning your repo onto your VM, and using ls -l; only directories should have "x" flags.)

For your own protection, it is advisable to commit your code to git at least once per day. Be sure to push after you commit. The git history may be considered for matters such as special consideration, extensions and potential plagiarism. Your commit messages should be a short-hand chronicle of your implementation progress and will be used for evaluation in the Quality of Software Practices criterion.

You must **submit the full 40-digit SHA1 hash** of your chosen commit to the **Project 2 Assignment** on LMS. You must **also push your submission** to the repository named comp30023-2023-project-2 in the subgroup with your username of the group comp30023-2023-projects on gitlab.eng.unimelb.edu.au.

You will be allowed to update your chosen commit. However, only the last commit hash submitted to LMS before the deadline will be marked without late penalty.

You should ensure that the commit which you submitted is accessible from a fresh clone of your repository. For example (below, the . . . are added for clarity to break the line):



The requisite .gitlab-ci.yml file will be provided and placed in your repository, but is also available from the project2 repository linked above.

Please, please use this CI feature. Almost all failed projects come from not fixing bugs that are reported by CI.

#### 8 Getting help

Please see Project 2 Module on LMS.

### 9 Collaboration and Plagiarism

You may discuss this project abstractly with your classmates but what gets typed into your program must be individual work, not copied from anyone else. Do **not** share your code and do **not** ask others to give you their programs. The best way to help your friends in this regard is to say a very firm "**no**" if they ask to see your program,

point out that your "no", and their acceptance of that decision, are the only way to preserve your friendship. See https://academicintegrity.unimelb.edu.au for more information.

Note also that solicitation of solutions via posts to online forums, whether or not there is payment involved, is also Academic Misconduct. You should not post your code to any public location (e.g., GitHub) until final subject marks are released.

If you use a **small** amount of code not written by you, you must attribute that code to the source you got it from (e.g., a book or Stack Exchange).

Do **not** post your code on the subject's discussion board Ed.

**Plagiarism policy:** You are reminded that all submitted project work in this subject is to be your own individual work. Automated similarity checking software will be used to compare submissions. It is University policy that cheating by students in any form is not permitted, and that work submitted for assessment purposes must be the independent work of the student concerned.

Using git properly is an important step in the verification of authorship. We should see the stages of your code being written, not just the finished product.

AI software such as ChatGPT can generate code, but it will not earn you marks. You are allowed to use AI tools, but if you do then you must strictly adhere to the following rules.

