This is a review of context-free grammars from the lecture on Thursday; in each example, the grammar itself is on the left; the explanation for each non-terminal is on the right.

• Properly nested strings of parentheses.

$$S \to \epsilon \mid S(S)$$
 properly nested parentheses

Here is a different grammar for the same language:

$$S \to \epsilon \mid (S) \mid SS$$
 properly nested parentheses

• $\{0^m1^n \mid m \neq n\}$. This is the set of all binary strings composed of some number of 0s followed by a different number of 1s.

$$S \to A \mid B \\ A \to 0A \mid 0C \\ B \to B1 \mid C1 \\ C \to \epsilon \mid 0C1$$

$$\{0^{m}1^{n} \mid m \neq n\} \\ \{0^{m}1^{n} \mid m < n\} \\ \{0^{m}1^{n} \mid m = n\}$$

- 1 $\{0^{2n}1^n \mid n \ge 0\}$ https://tutorcs.com
- 2 $\{0^m1^n \mid m \neq 2n\}$ (Hint: If $m \neq 2n$, the with m 22 transfer mesorevious grammar, but pay attention to parity. This language contains the string 01.)
- 3 $\{0,1\}^* \setminus \{0^{2n}1^n \mid n \ge 0\}$ (Hint: Extend the previous grammar. What is missing?)

Work on these later:

- 4 $\{w \in \{0,1\}^* \mid \#(0,w) = 2 \cdot \#(1,w)\}$ Binary strings where the number of 0s is exactly twice the number of 1s.
- [Anti-hint: The language $\{ww \mid w \in \{0,1\}^*\}$ is **not** context-free. Thus, the complement of a context-free language is not necessarily context-free!]