

Give regular expressions for each of the following languages over the alphabet  $\{0, 1\}$ .

- 1** All strings containing the substring 000.

**Solution:**  $(0 + 1)^*000(0 + 1)^*$

- 2** All strings *not* containing the substring 000.

**Solution:**  $(1 + 01 + 001)^*(\varepsilon + 0 + 00)$

**Solution:**  $(\varepsilon + 0 + 00)(1(\varepsilon + 0 + 00))^*$

- 3** All strings in which every run of 0s has length at least 3.

**Solution:**  $(1 + 0000^*)^*$

**Solution:**  $(\varepsilon + 1)((\varepsilon + 0000^*)1)^*(\varepsilon + 0000^*)$

- 4** All strings in which 1 does not appear after a substring 000.

**Solution:**  $(1 + 01 + 001)^*0^*$

- 5** All strings containing at least three 0s.

**Solution:**  $(0 + 1)^*0(0 + 1)^*0(0 + 1)^*0(0 + 1)^*$

**Solution:**  $0^*01^*01^*0(0 + 1)^*$  or  $(0 + 1)^*001^*01^*01^*$

- 6** Every string except 000. (**Hint:** Don't try to be clever.)

**Solution:** Every string  $w \neq 000$  satisfies one of three conditions: Either  $|w| < 3$ , or  $|w| = 3$  and  $w \neq 000$ , or  $|w| > 3$ . The first two cases include only a finite number of strings, so we just list them explicitly. The last case includes *all* strings of length at least 4.

$$\begin{aligned} & \varepsilon + 0 + 1 + 00 + 01 + 10 + 11 \\ & + 001 + 010 + 100 + 000 + 001 + 010 + 100 + 110 + 111 \\ & + (1 + 0)(1 + 0)(1 + 0)(1 + 0)(1 + 0)^* \end{aligned}$$

**Solution:**  $\varepsilon + 0 + 00 + (1 + 01 + 001 + 000(1 + 0))(1 + 0)^*$

- 7** All strings  $w$  such that *in every prefix of  $w$* , the number of 0s and 1s differ by at most 1.

**Solution:** Equivalently, strings that alternate between 0s and 1s:  $(01 + 10)^*(\varepsilon + 0 + 1)$

- 8** (**Hard.**) All strings containing at least two 0s and at least one 1.

**Solution:** There are three possibilities for how such a string can begin:

- Start with 00, then any number of 0s, then 1, then anything.
- Start with 01, then any number of 1s, then 0, then anything.
- Start with 1, then a substring with exactly two 0s, then anything.

All together:  $000^*1(0 + 1)^* + 011^*0(0 + 1)^* + 11^*01^*0(0 + 1)^*$

Or equivalently:  $(000^*1 + 011^*0 + 11^*01^*0)(0 + 1)^*$

## Solution:

There are three possibilities for how the three required symbols are ordered:

- Contains a 1 before two 0s:  $(0+1)^*1(0+1)^*0(0+1)^*0(0+1)^*$
- Contains a 1 between two 0s:  $(0+1)^*0(0+1)^*1(0+1)^*0(0+1)^*$
- Contains a 1 after two 0s:  $(0+1)^*0(0+1)^*0(0+1)^*1(0+1)^*$

So putting these cases together, we get the following:

$$\begin{aligned} & (0+1)^*1(0+1)^*0(0+1)^*0(0+1)^* \\ & + (0+1)^*0(0+1)^*1(0+1)^*0(0+1)^* \\ & + (0+1)^*0(0+1)^*0(0+1)^*1(0+1)^* \end{aligned}$$

**Solution:**  $(0+1)^*(101^*0 + 010^*011^*0 + 01^*01)(0+1)^*$

**9 (Hard.)** All strings  $w$  such that *in every prefix of  $w$* , the number of 0s and 1s differ by at most 2.

**Solution:**  $(0(01)^*1 + 1(10)^*0)^* \cdot (\varepsilon + 0(01)^*(0 + \varepsilon) + 1(10)^*(1 + \varepsilon))$

**10 (Really hard.)** All strings in which the substring 000 appears an even number of times.

(For example, 0001000 and 0000 are in this language, but 00000 is not.)

**Solution:** Every string in  $\{0,1\}^*$  alternates between (possibly empty) blocks of 0s and individual 1s; that is,  $\{0,1\}^* = (0^*1)^*0^*$ . Trivially, every 000 substring is contained in some block of 0s. Our strategy is to consider which blocks of 0s contain an even or odd number of 000 substrings.

Let  $X$  denote the set of all strings in  $0^*$  with an even number of 000 substrings. We easily observe that  $X = \{0^n \mid n = 1 \text{ or } n \text{ is even}\} = 0 + (00)^*$ .

Let  $Y$  denote the set of all strings in  $0^*$  with an *odd* number of 000 substrings. We easily observe that  $Y = \{0^n \mid n > 1 \text{ and } n \text{ is odd}\} = 000(00)^*$ .

We immediately have  $0^* = X + Y$  and therefore  $\{0,1\}^* = ((X+Y)1)^*(X+Y)$ .

Finally, let  $L$  denote the set of all strings in  $\{0,1\}^*$  with an even number of 000 substrings. A string  $w \in \{0,1\}^*$  is in  $L$  if and only if an odd number of blocks of 0s in  $w$  are in  $Y$ ; the remaining blocks of 0s are all in  $X$ .

$$L = ((X1)^*Y1 \cdot (X1)^*Y1)^*(X1)^*X$$

Plugging in the expressions for  $X$  and  $Y$  gives us the following regular expression for  $L$ :

$$\left( ((0 + (00)^*)1)^* \cdot 000(00)^*1 \cdot ((0 + (00)^*)1)^* \cdot 000(00)^*1 \right)^* \cdot ((0 + (00)^*)1)^* \cdot (0 + (00)^*)$$

Whew!