

CS 563  
Assignment Project Exam Help  
Concurrent Programming  
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Lecture 10: Synchronization, Atomic  
Actions, and Await Statements

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# Research Project Timeline

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- ❖ 3 / 1: Reading list 1 due
- ❖ 3 / 22: Reading list 2 due **Assignment Project Exam Help**  
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- ❖ **3 / 23: Midterm** **WeChat: cstutorcs**
- ❖ 3 / 30-4 / 20: Research paper presentations (2 students / lecture time)
- ❖ 4 / 5: Research project proposal due
- ❖ 4 / 22-4 / 29: Research project presentations (4 students / lecture time)
- ❖ 5 / 8: Research project due (paper and source code)



# Example

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`y = 0;      z = 0;`

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P1:      `x = y + z;`      P2:      `y = 1;`  
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- ❖ What are the final values of x, y, and z?
- ❖ Answer depends on execution order and what is atomic



# Concepts

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- ❖ **State:** values of variables at a point in time  
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- ❖ **Atomic action:** indivisible state change  
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  - ❖ hardware: load, store, add, ...
  - ❖ software: critical sections (later...)
- ❖ **History:** a trace of ONE execution; an interleaving of atomic actions
- ❖ **Property:** an attribute of ALL histories of a program, e.g., correctness



# Histories

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- ❖ How many histories are there in a program with  $n$  processes and  $m$  atomic actions per process?

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$$(n \cdot m)! / (m!)^n$$

- ❖ If  $n = 2$  and  $m=4$ ...



# Synchronization

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❖ Synchronization <https://tutorcs.com>

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- ❖ restricts the number of histories
- ❖ Example: given array  $a[1:n]$  of positive integers
  - ❖ find the maximum value  $m$  by examining all elements in parallel



# Finding Max of Array

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- ❖ Goal -- expressed as a predicate

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$$(\forall j : 1 \leq j \leq n : m \geq a[j]) \wedge$$
$$(\exists j : 1 \leq j \leq n : m = a[j])$$

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- ❖ Sequential program

```
int m = 0
for [i = 0 to n-1] {
    if a[i] > m
        m = a[i];
}
```



# Concurrent Program

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## ❖ Program outline

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```
int m = 0
```

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```
co [i = 0 to n-1] {
```

body

```
}
```



# Program 1

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- ❖ No synchronization (hence no constraints)

```
if a[i] > m {  
  m = a[i];  
}
```

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- ❖ What is true before and after this statement?
- ❖ What is wrong with this program?
  - ❖ interference
- ❖ What's the most that we can say?
  - ❖ m is SOME a[i]



# Program 2

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- ✧ Make it a single atomic action

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```
< if (a[i] > m) {  
    m = a[i]  
} >
```

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- ✧ The program is now correct, but...



# Program 3

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- ❖ Use smaller atomic actions

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`if (a[i] > m) {`

`< m = a[i] >`

`}`



# Program 4

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- ✧ Combine 2 and 3: Double checking

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```
if (a[i] > m) {  
    if (a[i] > m) {  
        m = a[i]  
    }  
}
```



# Points to Remember

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- ❖ Synchronization is needed when we don't have independence
- ❖ Angle brackets specify atomic actions
- ❖ Double-checking technique -- especially when it is possible that the first check is false



# Synchronization

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- ✧ Prevents undesirable interleavings by  
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- ✧ Combining fine-grained atomic actions into coarse-grained atomic actions
- ✧ **Mutual exclusion**
- ✧ Delaying process execution until some predicate is satisfied
- ✧ **Condition synchronization**



# Atomic Actions

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- ❖ Fine grained -- reads and writes of simple variables (loads / stores)
- ❖ Coarse grained -- programmed using  $\langle \dots \rangle$  or real code (critical section)



# Example: Fine-grained Atomicity

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```
int y = 0, z = 0;
```

```
co x = y + z;
```

```
// y = 1; z = 2;
```

```
co;
```

co statement:  
a simple way  
to represent  
concurrency

- ❖ What is the final value of x
- ❖ How to achieve expression atomicity?



# Definitions

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- ❖ A **critical reference** in an expression is a reference to a variable that is changed by other processes
- ❖ An assignment  $x=e$  satisfies the At-Most-Once Property if one of the following is true:
  - ❖  $e$  contains at most one critical reference and  $x$  is not referenced by another process
  - ❖  $e$  contains no critical references, in which case  $x$  may be read by other processes

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# Appearance of Atomicity

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- ❖ If an assignment meets the requirements of the At-Most-Once property
  - ❖ execution of the assignment statement will appear to be atomic



# Examples

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```
int x = 0, y = 0;
```

```
co x = x + 1; // y = y + 1; oc;
```

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```
int x = 0, y = 0;
```

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```
co x = y + 1; // y = y + 1; oc;
```

```
int x = 0, y = 0;
```

```
co x = y + 1; // y = x + 1; oc;
```



# Examples

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int y = 0, z = 0;

CO x = y + z;

// y = 1; z = 2;

CO;



# Coarse-grained Atomicity

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- ✧ Atomicity may be required when At-Most-Once Property is not held

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- ✧ Need mechanism for constructing coarse-grained atomic actions
  - ✧ Sequence of fine-grained atomic actions
  - ✧ Appearance



# Specifying Atomic Actions

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- ❖ `< S; >`

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- ❖ execute statement list S indivisibly (mutual exclusion)

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- ❖ `< await(B); >`

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- ❖ wait for B to be true (conditional synchronization)

- ❖ `< await(B) S; >`

- ❖ wait for B to be true, then execute S ALL AS A SINGLE ATOMIC ACTION



# Example of use of await

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`< s = s+1; >` <https://tutorcs.com>

`< await (s>0) s = s-1; >` WeChat: cstutorcs



# At-Most-Once Property

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✧  $\langle S; \rangle == S;$

✧ if either:

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✧ S is a single assignment and meets At-Most-Once Property requirement, OR

✧ S is implemented by a single indivisible machine instruction

✧  $\langle \text{await } (B); \rangle == \text{while } (\text{not } B)$

✧ if B meets At-Most-Once Property requirement (i.e., contains only one critical reference)



# Producer/Consumer Synchronization

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- ❖ Problem:
  - ❖ copy  $a[n]$  in Producer to  $b[n]$  in Consumer using a single shared buffer
  - ❖ How? Use synchronization to alternate access to the buffer



# Program

---

```
int buf, p = 0, c = 0;
```

```
process Producer {  
    int a[n];  
    while (p < n) {  
        < await (p == c); >  
        buf = a[p];  
        p=p+1;  
    }  
}
```

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```
process Consumer {  
    int b[n];  
    while (c < n) {  
        < await (p > c); >  
        b[c] = buf;  
        c=c+1;  
    }  
}
```