Trinity College Dublin
Coláiste na Tríonóide, Baile Átha Cliath
The University of Dublin

2.1 – Microprocessor Basics

CSU11021 – Introduction to Computing I

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Simple Model of a Microprocessor System

A **processing unit** (or **processor** or **CPU**) which performs operations on data

Memory, which stores:

Data: representing text, images, videos, sensor readings, π, audio, etc. ...

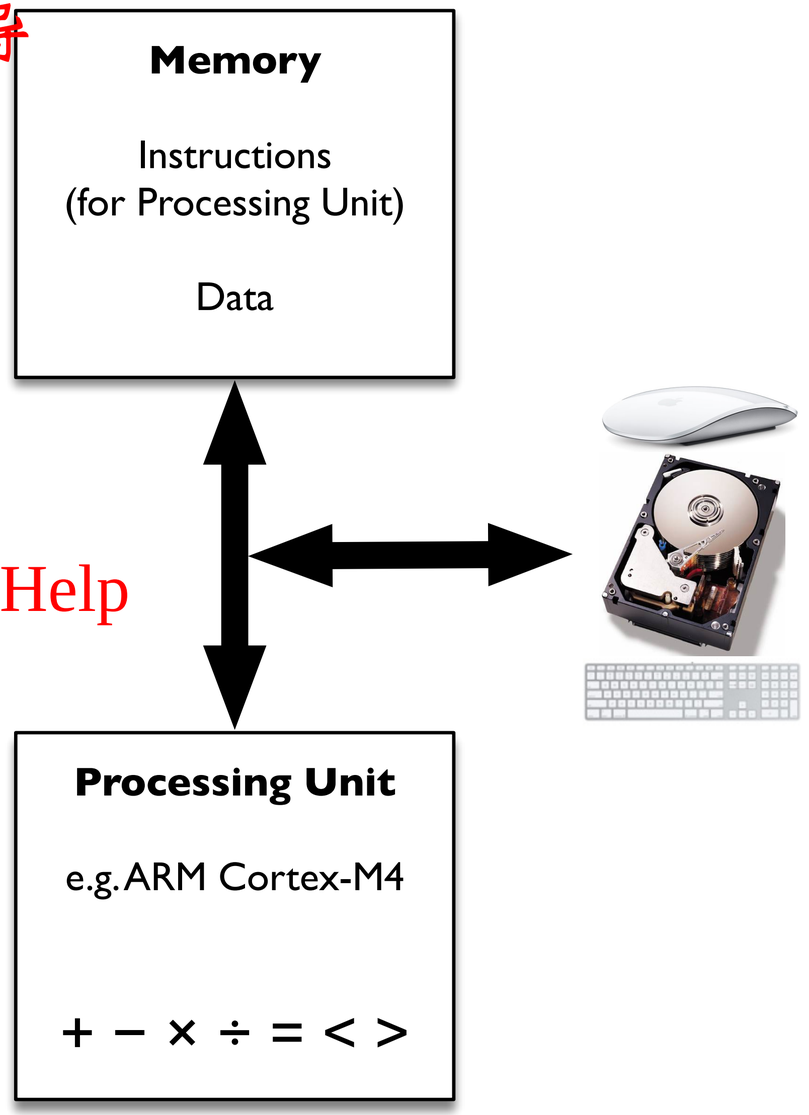
Instructions: Programs are composed of sequences of instructions that control the actions of the processing unit

Instructions typically describe very simple operations, e.g.

Add two values together

Move a value from one place to another

Compare two values



Simple Model of a Microprocessor System

Memory can be viewed as a series of equally-sized "**locations**", each of which stores a small piece of information

Each location has a unique "**address**"

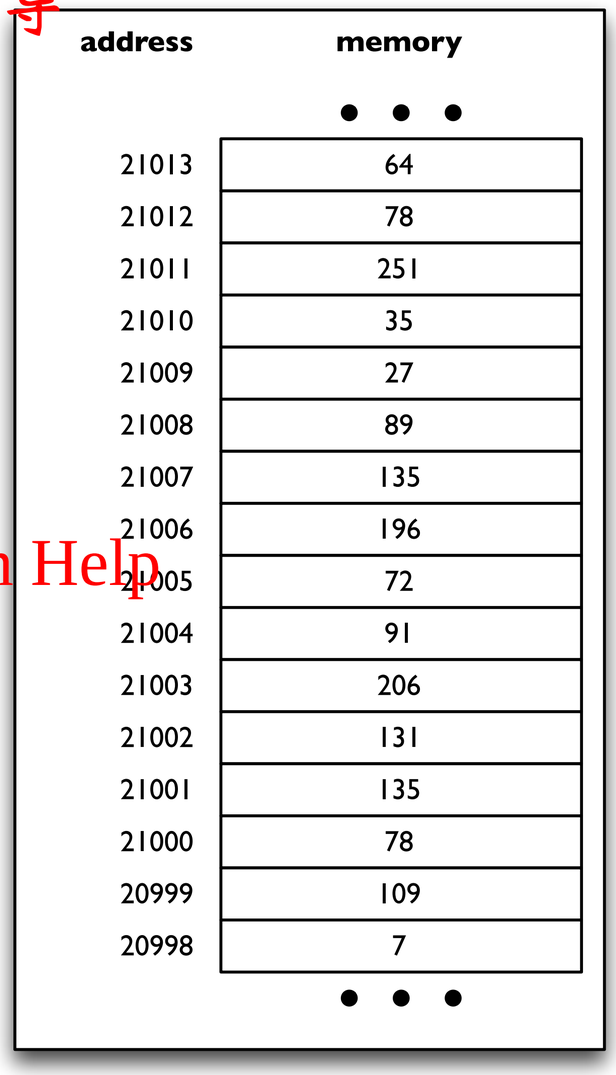
The information at each location may be

data, e.g. the value 91 or

an **instruction** that tells the processor how to manipulate the data

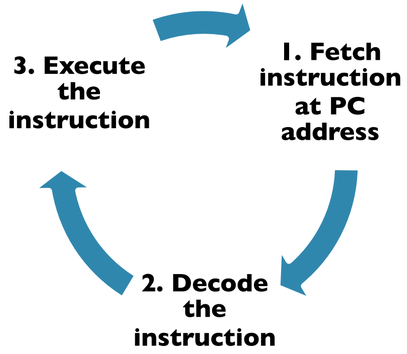
But instructions are also encoded as values!!

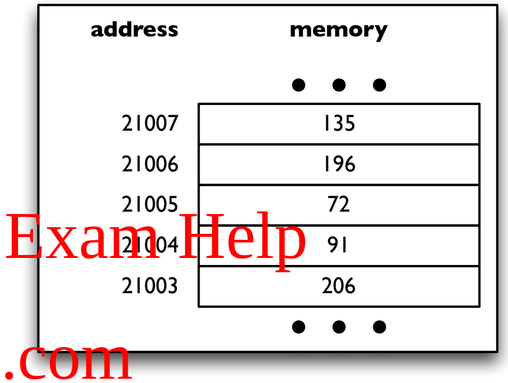
e.g. the value 91 might be a code used to tell the processing unit to add two other values together



Program Execution

When the computer is turned on, the processing unit begins executing the instruction in the me[illegible] the address stored in the Program Counter or PC





After fetching an instruction, the value of the Program Counter is changed to the address of the next instruction in the program

Processing unit keeps doing this until the computer is turned off

Program Execution

This simple model of a microprocessor system (computer) is the model used by computers familiar to us (PCs, games consoles, mobile phones, engine control units, ...)

Behaviour is predictable (deterministic)

> *"On two occasions I have been asked, — 'Pray, Mr. Babbage, if you put into the machine wrong figures, will the right answers come out?' ... I am not able rightly to apprehend the kind of confusion of ideas that could provoke such a question."*
>
> Charles Babbage, Passages from the Life of a Philosopher (1864), ch. 5 "Difference Engine No. 1"

The "power" of computers arises because they perform a lot of simple operations very quickly

The complexity of computers arises because useful programs are composed of many thousands or millions of simple instructions

Possibly executing in parallel on more than one processor/computer!

Registers


```
MOV  total, a
ADD  total, total, b
ADD  total, total, c
ADD  total, total, d
```


Make the first number the subtotal

Add the second number to the subtotal

rd number to the subtotal

rth number to the subtotal

Recall our simple program from Lecture #1

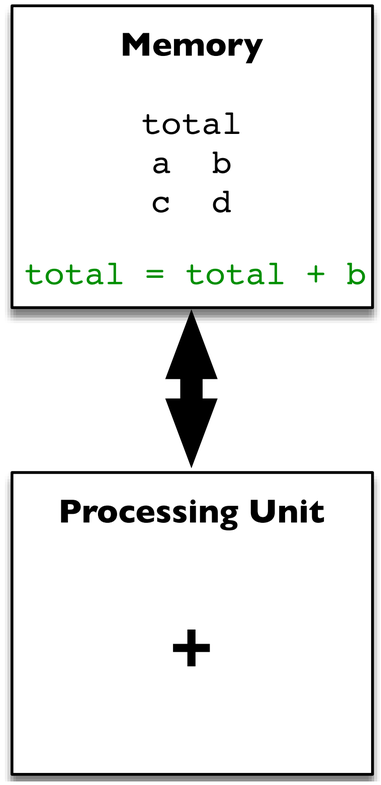
Add four numbers together

total = a + b + c + d

total, a, b, c, and d are stored in memory

operations (move and add) are performed in the processing unit

Question: How many memory ↔ processing unit transfers?



Registers

total = total + b

Load instruction from Memory to Processing Unit

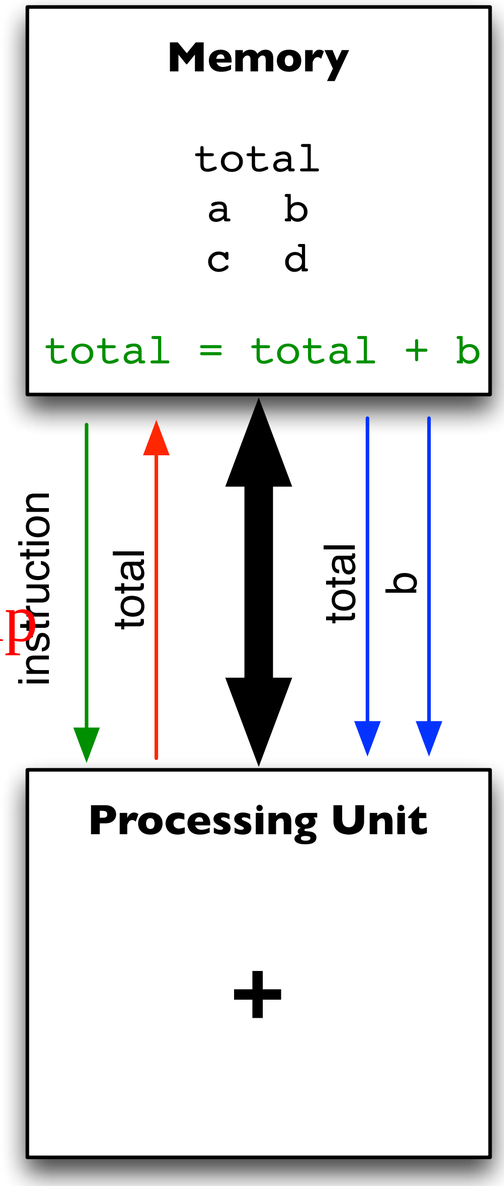
Load total from Memory to Processing Unit

Load b from Memory to Processing Unit

Compute total + b

Store total from Processing Unit to Memory

Accessing memory is slow relative to the speed at which the processor can execute instructions



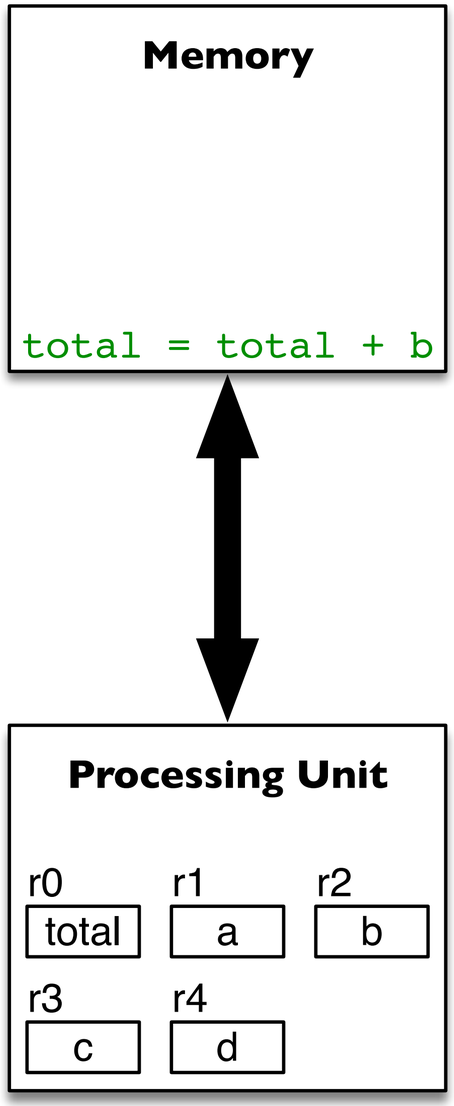
Registers

Processors use small fast internal storage to temporarily store values called registers

ARM microprocessors have 16 registers

- Labelled R0, R1, ..., R15
- R15 is special – the Program Counter
- R13 and R14 are also special (you should avoid using them for now)

Question: How many memory ↔ Processing Unit transfers if total, a, b, c, and d are stored in registers?



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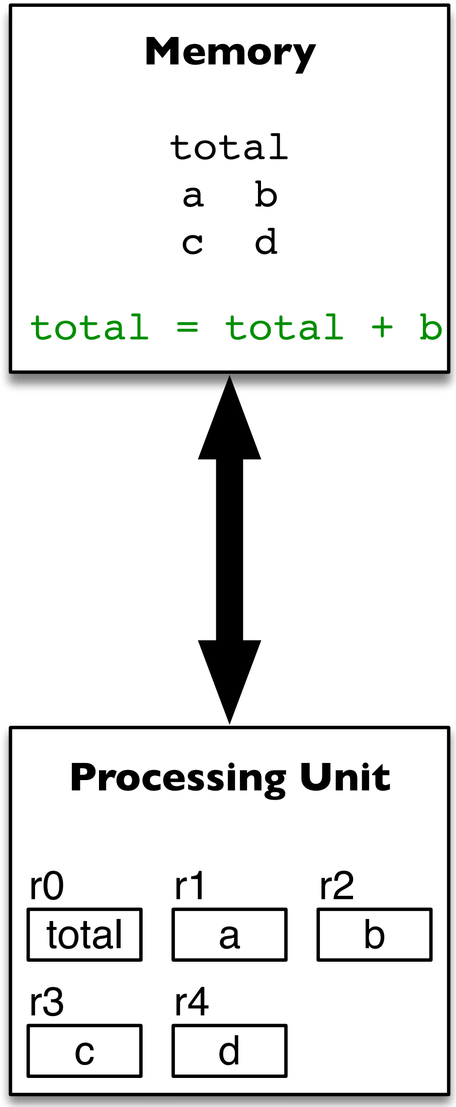

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2.2 – Machine Code and Assembly Language

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Machine Code

A program (any program, originally written using any language) is composed of a sequence of **machine code instructions** that are stored in memory

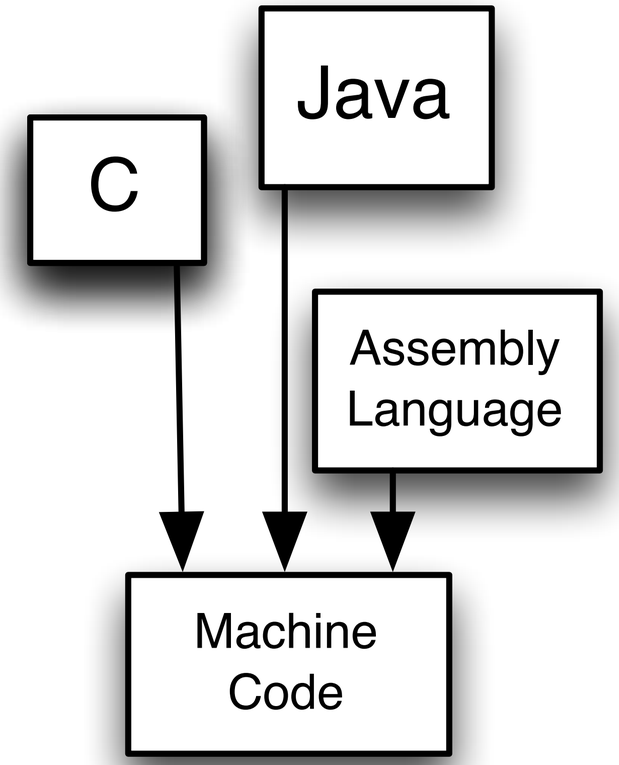
Instructions determine the operations performed by the processor (e.g. add, move, multiply, subtract, compare, ...)

A single instruction is composed of

- an **operator** (the operation to perform, e.g. ×, ÷, −, +)
- zero, one or more **operands** (e.g. R1, R2, R3, R4)

Each instruction and its operands are encoded using a numerical value

- e.g. 3766550530 is the machine code that causes the processor to add the values in R1 and R2 and store the result in R3 (`ADD R3, R1, R2`)



Assembly Language

Writing programs using machine code is possible …

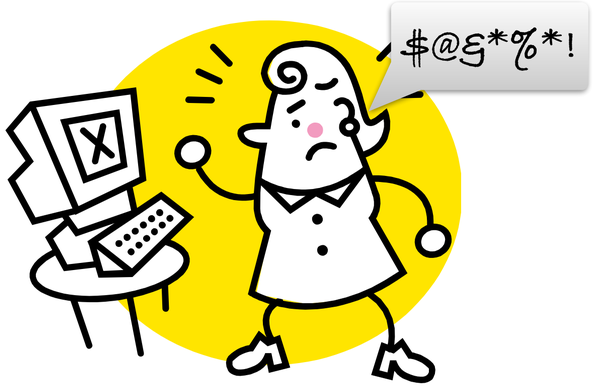
… but not practical

Instead, we will write programs using assembly language

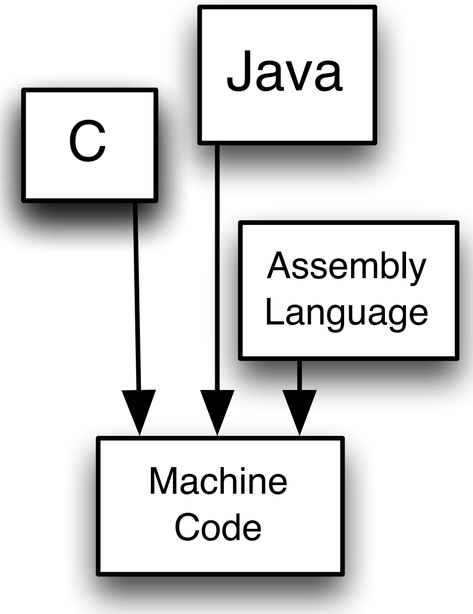
- Instructions are expressed using mnemonics
- e.g. the word “ADD” instead of the machine code 3766550530
- e.g. the expression “R2” to refer to register number two

Assembly language must still be translated into machine code

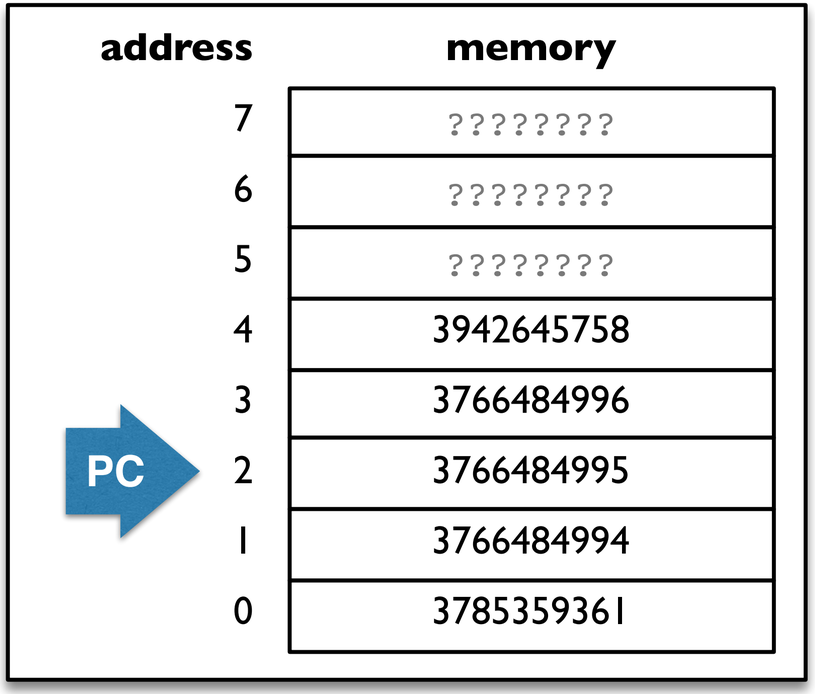
- This is done using a program called an assembler
- Machine code produced by the assembler is stored in memory and executed by the processor (the **fetch – decode – execute** cycle)

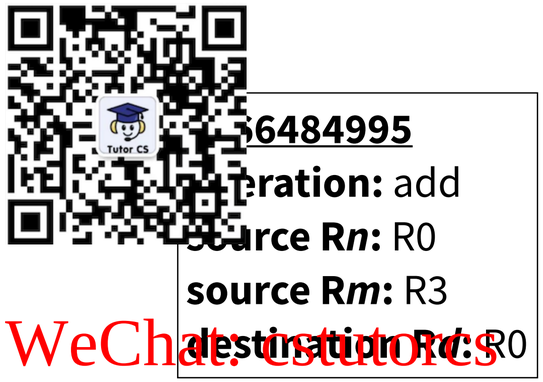


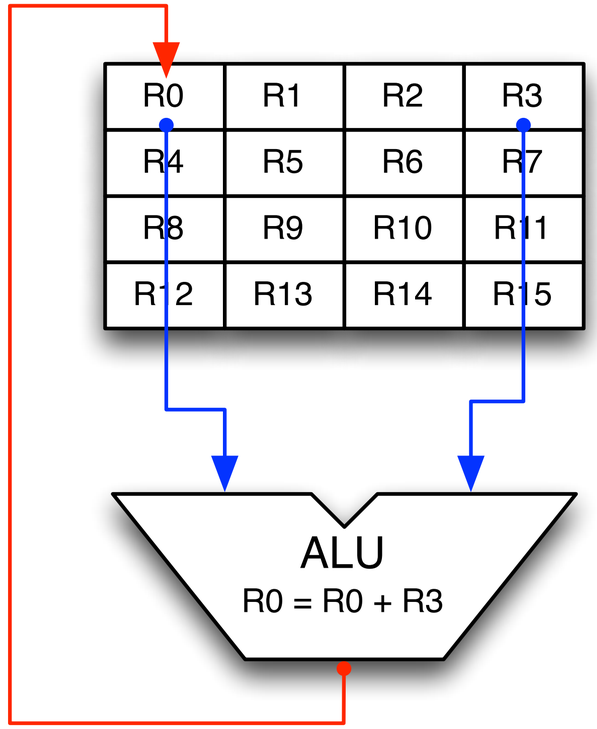
This is you writing a machine code program!!



Instruction Execution







Fetch next instruction from memory at the address contained in the Program Counter (PC ≡ R15)

PC advances (increments) to next instruction

Machine Code instruction is decoded to determine operation and source / destination operands

Instruction is executed. (In this example the **ALU** adds the values in R0 and R3, storing the result back in R0.)

General form of an ARM Assembly Language instruction:

```
OP      dest, source_1, source_2
```

e.g. Add the values in R2 and R3, store the result in R1


```
ADD     R1, R2, R3      @ R1 = R2 + R3
```


Some ARM instructions have just one input operand


```
MOV     R4, R5          @ R4 = R5
```

Immediate Operands

Register Operands

```
ADD     R1, R2
MOV     R2
```

Register Operand

Often we want to use constant values, instead of registers (variables)

e.g. move the value 0 (zero) into register R3

```
MOV     R0, #0
```

e.g. set R1 = R2 + 1

```
ADD     R1, R2, #1
```

Immediate Operands

Simple Arithmetic

Write an ARM Assembly Language program to compute $4x^2+3x$ if x is stored in reg R1. Store the result in reg R0.

```
@ Write a program to compute 4x^2+3x
MUL   R0, R1, R1        @ result = x * x
MOV   R2, #4            @ tmp = 4
MUL   R0, R2, R0        @ result = 4 * x * x

MOV   R2, #3            @ tmp = 3
MUL   R2, R1, R2        @ tmp = x * tmp

ADD   R0, R0, R2        @ result = result + tmp
```

Cannot use MUL to multiply by a constant value
(a.k.a. immediate operand)

R1 is unmodified by our solution … which may be something we want … or maybe we don't care

LoaD Register (LDR)

`MOV Rx, #y` can only be used to load certain small values into Registers (Which values and how small? Long story for another day …)

Loa**D** **R**egister instruction can be used to load any* value into a register

```
LDR     R4, =45673857    ; tmp = 45673857
```


Note use of **=x** syntax instead of **#x** with LDR instruction


```
LDR     R2, =3           ; tmp = 3
MUL     R2, R1, R2       ; tmp = x * tmp
```


Note use of **=** again in operand **=3**

If in doubt, use `LDR Rx, =y`.

* well, not quite any

Provide meaningful comments and assume someone else will be reading your code


```
MUL     R2, R1, R2      @ r2 = r1 * r2
```

✗


```
MUL     R2, R1, R2      @ tmp = x * tmp
```

✔

Break your programs into small blocks separated by white space

While starting out, keep programs simple

Pay attention to initial values in registers (and memory)