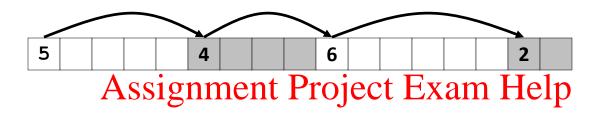
Keeping Track of Free Blocks

■ Method 1: *Implicit list* using length—links all blocks



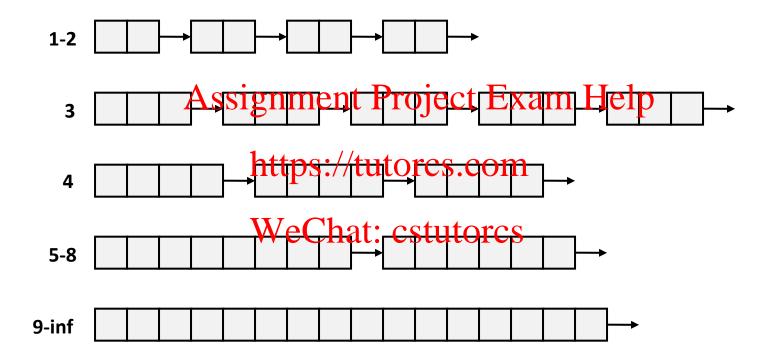
■ Method 2: Explicitlist among the free blocks using pointers



- Method 3: Segregated free list
 - Different free lists for different size classes
- Method 4: *Blocks sorted by size*
 - Can use a balanced tree (e.g. Red-Black tree) with pointers within each free block, and the length used as a key

Segregated List (Seglist) Allocators

Each size class of blocks has its own free list



- Often have separate classes for each small size
- For larger sizes: One class for each two-power size

Seglist Allocator

Given an array of free lists, each one for some size class

To allocate a block of size n:

- Search appropriate free list for block of size m > n Assignment Project Exam Help If an appropriate block is found:
- - Split block and place fragment on appropriate list (optional)
- If no block is found, try next larger class
- Repeat until blocky for cstutores

If no block is found:

- Request additional heap memory from OS (using **sbrk ()**)
- Allocate block of *n* bytes from this new memory
- Place remainder as a single free block in largest size class.

Seglist Allocator (cont.)

- To free a block:
 - Coalesce and place on appropriate list (optional)
- Advantages of seglishallooptersect Exam Help
 - Higher throughput
 - log time for bttps://twtorzesclasses
 - Better memory utilization
 - First-fit search of segregated free list approximates a best-fit search of entire heap.
 - Extreme case: Giving each block its own size class is equivalent to best-fit.