Project 2: The Raft Consensus Algorithm

Please read <u>all</u> pages before writing code.

1 Overview

Important Dates:

Project release: Thursday, October 26, 2023

Checkpoint due: Monday, November 6, 2023 at 11:59pm Full project due: Thursday, November 16, 2023 at 11:59pm Submission limits: 15 Gradescope submissions per checkpoint

In this project, you'll implement the Raft consensus algorithm, a replicated state machine protocol. You will want to start early. For more information regarding what portion of the project is expected to be completed for the checkpoint and the final test please refer to section and algorithm.

The starter code for this project is hosted as a read-only repository on GitHub. For instructions on how to build, run, test, and submit your server implementation, see the README.md file in file triples strottly the triples of the

Wte Contracts:/@stublem17640/p2.git

You must work on this project **individually**. You will have 15 submissions for each due date. The same policy of P1 will apply here as well - at most 3 late days per due date. No submissions will be accepted 3 days after each deadline.

About Gradescope:

Your Gradescope submission will output a message showing how many submissions you have made and how many you have left. Gradescope will allow you to submit beyond this limit, but we will be checking manually. Gradescope will also allow you to select a submission to grade. Only your selected submission within the first 15 submissions counts. If your selected submission is not within the first 15 submissions, your last submission of the first 15 submissions will be graded. We won't accept requests to update scores if you miscount the number of your submissions.

Please remove all your print statements before making the submission. The autograder may not work properly with print statements.

There will be no manual style grading for the checkpoint. However, 4 points are allocated for manual style grading on the final submission. Specifically, we are looking for good function headers, comments for variables, and no debugging print statements or chunks of dead code. For more details on the rubric, please refer to the Style Guide. You're allowed to use mutexes for this project (see Project Requirements).

2 Raft

A replicated service (e.g., key/value database) achieves fault tolerance by storing copies of its data on multiple replicas. Replication allows the service to continue operating even if some of its replicas experience failures (crashes or a broken/flaky network). The challenge is that failures may cause the replicas to hold differing copies of the data.

One protocol to ensure all of these copies of the data are consistent across all non-faulty replicas is Raft. Raft implements a replicated state machine by sequencing client requests into a log, and ensuring that the replicas agree on the contents and ordering of the log entries. Tack replica asynchronously applies the client requests in the older they appear in the replica's log of the service's state. If a replica fails and fater recovers, Raft takes care of bringing the log of the recovered replica up to date. Raft will continue to operate as long as at least a quorum of replicas is alive and able to communicate. If a quorum is not available, Raft will sportaking plages for the surplus soon as a quorum becomes available.

In this project, you will implement Raft in Go. Your Raft module will implement a replicated log using the Raft protectly using RPC to communicate between replicas. Your implementation should support an indefinite sequence of numbered commands (log entries). Each log entry is comprised of a client command and an index number. After a log entry is committed, Raft will "apply" the log entry by sending the committed log entry to the application that is using your Raft module.

Note: Only RPC may be used for interaction between different Raft instances. For example, different peers in your Raft implementation are not allowed to share Go variables or access shared files/sockets. Communicating between 2 replicas using anything except the approved rpc package will result in losing 50% of the points on this project.

You'll implement a part of the Raft protocol described in the extended paper. You do not need to implement persistence, cluster membership changes (Section 6) or log compaction / snapshotting (Section 7).

You should consult the extended Raft paper. You may also find it useful to look at this illustrated guide to Raft. For a broader perspective, have a look at Paxos, Chubby, Paxos Made Live, Spanner, Zookeeper, Harp, Viewstamped Replication, and Bolosky et al.

Tip 1: Start early. Although you can implement Raft in a relatively small number of lines of code (our reference solution is ≈ 700 lines), getting it to work correctly will be challenging. Both the algorithm and the code are tricky, and there are many corner cases to consider. When one of the tests fails, it may take a bit of puzzling to understand in what scenario your solution isn't correct, and how to fix your solution.

Tip 2: Read and understand the extended Raft paper before you start. Your implementation should follow the paper's description closely, particularly Figure 2, since that's what the tests expect. Figure 2 is reproduced on the last page of this handout.

Tip 3: Your first implementation may not be clean enough to easily reason about its correctness. Give yourself enough time to rewrite your implementation at least once.

3 The code

Implement Raft by adding code to raft/raft.go. In that file, you'll find a bit of skeleton code and examples of how to send and receive RPCs.

Your implementation must support the following interface, which the tester will use. You'll find more details in comments in raft.go.

```
// rf = NewPeerhttps://tutorcs.com
//
     Create a new Raft peer.
//
  rf.PutCommand Interface } (index term, Start agreement on a new log entry
//
//
// rf.GetState() (me, term, isLeader)
     Ask a Raft peer for "me" (see line raft.go:75),
//
//
     its current term, and whether it thinks it is a leader
//
// ApplyCommand
//
    Each time a new entry is committed to the log, each Raft peer
//
     should send an ApplyCommand message to the service (e.g. tester) on the
//
     same server, via the applyCh channel passed to NewPeer()
```

A service calls NewPeer(peers, me, ···) to create a Raft peer. The peers argument is an array of established RPC connections, one to each Raft peer (including this one). The me argument is the index of this peer in the peers array. PutCommand(command) asks Raft to start the processing to append the command to the replicated log. PutCommand() should

return immediately, without waiting for this process to complete. The service expects your implementation to send an ApplyCommand for each new committed log entry to the applyCh argument to NewPeer().

Your Raft peers should exchange RPCs using the rpc Go package that we provide to you (https://github.com/15-440/p2/src/github.com/cmu440/rpc). It is modeled after Go's rpc library, but internally uses Go channels rather than sockets. raft.go contains some example code that sends an RPC (sendRequestVote()) and that handles an incoming RPC (RequestVote()). The reason you must use rpc instead of Go's RPC package is that the tester tells rpc to delay RPCs, re-order them, and delete them to simulate challenging network conditions under which your code should work correctly. Any modifications you make to rpc will be discarded before grading.

Note: The rpc package only provides a subset of the functionality of Go's RPC system. For instance, asynchronous RPC calls are not provided by rpc.

4 Checkpoint Assignment Project Exam Help

Implement leader election and heartbeats (AppendEntries RPCs with no log entries). The goal for the checkpoint is for a single leader to be elected, for the leader to remain the leader if there are no failures, and for a new leader to take over if the old leader fails or if packets to/from the old leader are lost. Run go test -race -run 2A to test your checkpoint code.

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Be sure you pass the checkpoint tests before submitting. Note that the checkpoint tests test the basic operation of leader election. The final tests will test leader election under more challenging settings and may expose bugs in your leader election code which the checkpoint tests miss.

4.2 General Guidelines

Add any state you need to the Raft struct in raft.go. You'll also need to define a struct to hold information about each log entry. Your code should follow Figure 2 in the paper as closely as possible.

Fill in the RequestVoteArgs and RequestVoteReply structs. Modify NewPeer() to create a background goroutine that will kick off leader election periodically by sending out RequestVote RPCs when it hasn't heard from another peer for a while. This way a peer

will learn who the leader is, if there is already a leader, or become the leader itself. Implement the RequestVote() RPC handler so that servers will vote for one another.

To implement heartbeats, define an AppendEntries RPC struct (though you may not need all the arguments yet), and have the leader send them out periodically. Write an AppendEntries RPC handler method that resets the election timeout so that other servers don't step forward as leaders when one has already been elected.

Make sure the election timeouts in different peers don't always fire at the same time, or else all peers will vote only for themselves, and no one will become the leader.

The tester requires that the leader send heartbeat RPCs no more than ten times per second.

The tester requires your Raft to elect a new leader within five seconds of the failure of the old leader (if a majority of peers can still communicate). Remember, however, that leader election may require multiple rounds in case of a split vote (which can happen if packets are lost or if candidates unluckily choose the same random backoff times). You must pick election timeouts (and thus heartbeat intervals) that are short enough that it's very likely that an election will complete in less than five seconds even if it requires multiple rounds.

The papers Sexton In mentions declined in the transaction to RW in Diseconds. DO NOT USE THEIR RECOMMENDED TIMEOUT INTERVALS. Those values only make sense if you use a heartbeat interval significantly smaller than 150 ms (the lower bound), which is not the case in this project. Instead, here are some general guidelines on how to choose intervals in this project.

- 1. You should not send teartbeat RPCs more than 10 times per second. CSTUTOTCS
- 2. Election timeout should be multiple times greater than your heartbeat interval.
- 3. Election timeout should be small enough to allow re-election of a leader in under 5 seconds.
- 4. Make sure the election timeouts on different peers don't always fire at the same time, as explained earlier.
- 5. We recommend you run tests > 10 times locally to make sure you do not have concurrency issues before you submit to Gradescope.

4.3 Notes and Hints

There are some details and hints we want to emphasize here to help you pass our tests:

- If an RPC arg or arg struct field is unexpectedly nil on the receiver side:
 - Check that all struct and sub-struct field names are capitalized. Go RPC only serializes struct fields with capitalized names. Sub-structures must also have capitalized field names (e.g., fields of log records in an array). Forgetting to capitalize field names sent by RPC is the single most frequent source of bugs while using RPCs.
 - Check that the args struct and its sub-structs are not completely empty. (If the struct or any sub-structs are completely empty, it will cause an encoding error and give nil for the whole args value.) Also, it appears that sending an empty slice will cause that slice to decode as nil instead of an empty slice. So either add a check for the (slice == nil) case (treated as an empty slice), avoid empty slices, or use an array instead of a slice.

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- Unlike in P1, you may use buffered channels of any size in this project. Additionally, you are also permitted to use mutexes in this project.
- You'll need to write code that takes actions periodically or after delays in time. The easiest way to do this is to create a goroutine with a loop that calls time. Sleep().
- If your code was to ble parting the Settliffed fie Sper's Figure 2 again; the full logic for leader election is spread over multiple parts of the figure.
- A good way to debug your code is to insert debug logs when a peer sends or receives a message, redirect it to a file (go test -race > out.txt) and examine the file to trace the execution of your system.
- You should check your code with go test -race, and fix any races it reports.
- You should try your code with varying numbers of CPU cores (go test -race -cpu=N)

4.4 Debugging

In the starter code, we provide you a simple logging framework to allow you to maintain separate debug logs for for each peer, with the option to output to files or stdout. You

are welcome to use, modify, remove, or ignore this logging code – however, we expect you to have clear, understandable debug log files before asking questions at office hour or on Edstem. The easiest way to do this is to use our included debug logging framework.

These debug logs can easily be disabled by setting

kEnableDebugLogs = false

at the top of raft.go. Be sure to disable your debug logs before submitting to Gradescope. You can output your debug log output to stdout by setting

kLogToStdout = true

at the top of raft.go. Each log line will be prefixed by the peer's unique name.

If you set kLogToStdout to false, each peer's log will instead be output into a separate .txt file in the directory described by kLogOutputDir. By default, we enable debug logs and output the transfer project Exam Help

5 Final Test https://tutorcs.com

We want Raft to keep a consistent, replicated log of operations. A call to PutCommand() at the leader starts the process of adding a new operation to the log; the leader sends the new operation to the track servers using appending the RPCs.

5.1 Task

Implement the leader and follower code to append new log entries. This will involve implementing PutCommand(), completing the AppendEntries RPC structs, sending them, fleshing out the AppendEntry RPC handler, and advancing the commitIndex at the leader. Your first goal should be to pass the TestBasicAgree2B() test (in raft_test.go). Once you have that working, you should get all the final tests to pass (go test -race -run 2B).

5.2 General Guidelines

While the Raft leader is the only server that initiates appends of new entries to the log, all the replicas need to independently give each newly committed log entry to their local service

replica (via their applyCh passed to NewPeer()). You should try to keep the goroutines that implement the Raft protocol as separate as possible from the code that sends committed log entries on the applyCh (e.g., by using a separate goroutine for delivering committed messages). If you don't separate these activities cleanly, then it is easy to create deadlocks. Without a clean separation, a common deadlock scenario is as follows: an RPC handler sends on the applyCh, but it blocks because no goroutine is reading from the channel (e.g., perhaps because it called PutCommand()). Now, the RPC handler is blocked while holding the mutex on the Raft structure. The reading goroutine is also blocked on the mutex because PutCommand() needs to acquire it. Furthermore, no other RPC handler that needs the lock on the Raft structure can run.

You will need to implement the election restriction (section 5.4.1 in the paper).

6 Hand In

Reminder: Please disable or remove all debug prints regardless of whether you are using our logging framework or not before immitting to Gradescope. This liefs aloid inadvertent failures, mess autograder outputs and siyle point deductions.

For both the checkpoint and the final submission, create handin.zip using the following command under the p2/directory and then upload it to Gradescope. sh make_submit.nltps.//tutorcs.com

7 Testing and Scalingt: cstutores

We will use Gradescope to automatically grade your implementation for correctness, and manual grading for style. In addition to tests we provide you in raft_test.go, we will run additional, more extensive tests on Gradescope for both your checkpoint and the final submissions. We will not be able to provide you any details (except for the stdout you'll see on Gradescope) about any of these new tests. We will run each test multiple times – you should pass every invocation of a test to pass that test.

You are encouraged to write new tests on your own. You are, however, not required to do so and will not be graded on any new tests you write. Here are some tips to do this. Refer to raft_test.go as you read these tips.

• See TestInitialElection2A: you can use cfg.checkOneLeader() to check for a leader's election and to get the current leader's ID.

• See TestFailAgree2B: you can use cfg.one(value, num_servers) to start an agreement.

• See TestFailNoAgree2B: you can use cfg.disconnect(server_id) and cfg.connect(server_id) to disconnect and connect servers. You can also directly call PutCommand() on one of the Raft peers by using cfg.rafts.

The checkpoint (Part A) is worth 45 points. The following table shows the point distribution for the final version of the project:

Part A	45 points
Part B	145 points
Manual Grading of Part B	4 points
go fmt of Part B	1 point
Total	195 points

Note that the tests for Part A are rup and graded for points for both the checkpoint and the final version by the pldjees diadescape assignment. However, there is no manual style grading for the checkpoint.

8 Project Requirements torcs.com

As you write code for this croject, also keep in mind the following requirements: CSTUTORS

- You must work on this project individually. You are free to discuss high-level design issues with other people in the class, but every aspect of your implementation must be entirely your work.
- You must format your code using go fmt and must follow Go's standard naming conventions. See the Formatting and Names sections of Effective Go for details.
- You may use any of the synchronization primitives in Go's sync package for this project.
- For the tester to function correctly, please use the provided rpc package instead of Go's native one.

State

Persistent state on all servers:

(Updated on stable storage before responding to RPCs)

currentTerm latest term server has seen (initialized to 0

on first boot, increases monotonically)

votedFor candidateId that received vote in current

term (or null if none)

log entries; each entry contains command

for state machine, and term when entry was received by leader (first index is 1)

Volatile state on all servers:

commitIndex index of highest log entry known to be

committed (initialized to 0, increases

monotonically)

lastApplied index of highest log entry applied to state

machine (initialized to 0, increases

monotonically)

Volatile state on leaders:

(Reinitialized after election)

nextIndex[] for each server, index of the next log entry

A to send to that server (initialized to lender Alast Ise index 1-1)

matchIndex[]

for each server, index of highest log entry

known to be replicated on server (initialized to 0, increases monotonically)

Appendintrix RSC//[U[O

Invoked by leader to replicate log entries (§5.3); also used as heartbeat (§5.2).

Arguments: term

leader's telm e Chat: C

leaderId so follower can redirect clients index of log entry immediately preceding

new ones

prevLogTerm term of prevLogIndex entry

entries[] log entries to store (empty for heartbeat;

may send more than one for efficiency)

leaderCommit leader's commitIndex

Results:

term currentTerm, for leader to update itself success true if follower contained entry matching

prevLogIndex and prevLogTerm

Receiver implementation:

- 1. Reply false if term < currentTerm (§5.1)
- Reply false if log doesn't contain an entry at prevLogIndex whose term matches prevLogTerm (§5.3)
- 3. If an existing entry conflicts with a new one (same index but different terms), delete the existing entry and all that follow it (§5.3)
- 4. Append any new entries not already in the log
- If leaderCommit > commitIndex, set commitIndex = min(leaderCommit, index of last new entry)

RequestVote RPC

Invoked by candidates to gather votes (§5.2).

Arguments:

term candidate's term candidateId candidate requesting vote

lastLogIndex index of candidate's last log entry (§5.4) term of candidate's last log entry (§5.4)

Results:

term currentTerm, for candidate to update itself true means candidate received vote

Receiver implementation:

- 1. Reply false if term < currentTerm (§5.1)
- 2. If votedFor is null or candidateId, and candidate's log is at least as up-to-date as receiver's log, grant vote (§5.2, §5.4)

Rules for Servers

All Servers:

• If commitIndex > lastApplied: increment lastApplied, apply log[lastApplied] to state machine (§5.3)

If RPG request or response on aims throat > currentTerm:

Followers (§5.2):

- · Respond to RPCs from candidates and leaders
- If election timeout elapses without receiving AppendEntries IPC from unrent leader or granting vote to candidate:
- convert to candidate

Candidates (§5.2):

• On conversion to candidate, start election:

On reme Courrent Term Vote for self

- Reset election timer
- · Send RequestVote RPCs to all other servers
- · If votes received from majority of servers: become leader
- If AppendEntries RPC received from new leader: convert to follower
- · If election timeout elapses: start new election

Leaders:

- Upon election: send initial empty AppendEntries RPCs (heartbeat) to each server; repeat during idle periods to prevent election timeouts (§5.2)
- If command received from client: append entry to local log, respond after entry applied to state machine (§5.3)
- If last log index ≥ nextIndex for a follower: send
 AppendEntries RPC with log entries starting at nextIndex
- If successful: update nextIndex and matchIndex for follower (§5.3)
- If AppendEntries fails because of log inconsistency: decrement nextIndex and retry (§5.3)
- If there exists an N such that N > commitIndex, a majority
 of matchIndex[i] ≥ N, and log[N].term == currentTerm:
 set commitIndex = N (§5.3, §5.4).

Figure 2: A condensed summary of the Raft consensus algorithm (excluding membership changes and log compaction). The server behavior in the upper-left box is described as a set of rules that trigger independently and repeatedly. Section numbers such as §5.2 indicate where particular features are discussed. A formal specification [31] describes the algorithm more precisely.