

# More polymorphism!

Oslo Haskell, April 2015

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## Introduction

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- Two main kinds of polymorphism used in Haskell.
- Parametric & ad-hoc polymorphism.
- Generic ADTs & functions “solve” parametric polymorphism.
- Typeclasses are a form of constrained polymorphism which in Haskell “solve” ad-hoc polymorphism.

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- Parametric polymorphism: Operate on values independently of their type.
- `length`  $:: [\alpha] \rightarrow \text{Int}$
- `length []`  $= 0$
- `length (_:xs)`  $= 1 + \text{length } xs$
- `length` only cares about the shape! There are plenty of functions like this in Prelude – try looking through it & identifying some of them!

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- Ad-hoc polymorphism: Operate on values of different types.
- The perhaps most common need for ad-hoc polymorphism: avoiding the need for separate `addInt` & `addFloat` functions.
- Solved by typeclasses in Haskell.

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- Typeclasses is a form of constrained polymorphism (AKA bounded qualification).
- `class Num  $\alpha$  where`  
    `(+), (-), (*) :: Num  $\alpha$  =>  $\alpha \rightarrow \alpha \rightarrow \alpha$`
- The functions are constrained to work on any type  $\alpha$  which has a `Num` instance.

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Bit bigger example of where ad-hoc polymorphism is useful...

- `class Visible  $\alpha$  where`  
    `render ::  $\alpha$  → Picture`
- `instance Visible Board where`  
    `render (Board bs) = Pictures (map render bs)`
- `instance Visible Brick where`  
    `render b = Color (mixColors 1.0 (health b) (colour b) white`  
        `$ uncurry Translate (centre b)`  
        `$ rectangleSolid (width b) (height b)`

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## Semigroups & Monoids



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## Semigroups & Monoids

- Two extremely simple typeclasses.
- Semigroups are a set of stuff with a binary operation that combines the stuff.
- A Monoid is a Semigroup with an element for which their binary operation is id.

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## Semigroups

- class Semigroup  $\alpha$  where  
     $(\langle \rangle) :: \alpha \rightarrow \alpha \rightarrow \alpha$
- Associative binary operation.
- Examples: addition & multiplication of numbers,  
    conjunction & disjunctions of booleans

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## Monoids

- Simple API. Just the identity, and the binary operation for combining things.
- `class Monoid  $\alpha$  where`
  - `mempty ::  $\alpha$`  -- Identity
  - `mappend ::  $\alpha \rightarrow \alpha \rightarrow \alpha$`  -- Binary operation

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## Monoids

- Some really simple examples:
- Numbers under addition:  $42 + 0 = 42$
- Numbers under multiplication:  $42 * 1 = 42$
- $[4] ++ [2]$  is the same as `mappend [4] [2]`
- $[42] ++ []$  is the same as `mappend [42] mempty`

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## Monoids

- Let's say you have some ByteStrings which you want to turn into Text.
- $f :: (T.Text \rightarrow T.Text) \rightarrow T.Text \rightarrow T.Text \rightarrow T.Text$   
 $f\ g\ x\ y = g\ x\ `T.append` g\ y$  -- annoying to “port”!
- $f :: Monoid\ m \Rightarrow (m \rightarrow m) \rightarrow m \rightarrow m \rightarrow m$   
 $f\ g\ x\ y = g\ x\ <>\ g\ y$  -- no “porting” necessary!

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**$\ast \rightarrow \ast$  polymorphism**

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## \*->\* polymorphism

- So far we've seen \* polymorphism.
- By \*, we mean “term level” values.
- We've seen functions that operate on data of different types. **length** doesn't care if you have a list of integers or a list of wibbles.

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## \*->\* polymorphism

- Now we're going to take a look at things that don't even care if it's a list of things!
- Lists are after all just a shape. What if you have a tree? Sometimes you don't care if you have a list of wibbles or a tree of wibbles, or indeed a wobble of wibbles!



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## $*$ -> $*$ polymorphism

- But first – what does  $*$  and  $*$ -> $*$  really mean??  $*$  is a concrete type, whilst  $*$ -> $*$  is incomplete.
- Here are some example of things of kind  $*$ :
- $42 \quad \quad \quad :: \text{Int} \quad \quad \quad :: *$
- $[42] \quad \quad \quad :: [\text{Int}] \quad \quad \quad :: *$
- $\text{Just } [42] :: \text{Maybe } [\text{Int}] :: *$

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## $\ast \rightarrow \ast$ polymorphism

- Now let's see some things that are kind  $\ast \rightarrow \ast$ :
- `Maybe`  $:: \ast \rightarrow \ast$
- `[]`  $:: \ast \rightarrow \ast$
- Aha... Interesting... Hmm... So how do we use this?
- Allow me to demonstrate.

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## Functors, Applicatives, & Monads

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## Functor

- Arguably the most ubiquitous typeclass in Haskell.
- Represents a computational context & a way to operate on whatever data is in this context.
- Basically used for any type that can be “mapped over”.

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## Functor

- The API is very simple!
- `class Functor  $\varphi$  where`  
    `fmap :: ( $\alpha \rightarrow \beta$ )  $\rightarrow \varphi \alpha \rightarrow \varphi \beta$`
- Just a way to apply  $\alpha \rightarrow \beta$  on  $\varphi \alpha$  and get  $\varphi \beta$ .

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## Functor

- instance Functor [] where
  - fmap \_ [] = []
  - fmap f (x:xs) = f x : fmap f xs
  - Yes, this is just Prelude.map!
- instance Functor Maybe where
  - fmap \_ Nothing = Nothing
  - fmap f (Just a) = Just (f a)

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## Functor

### Lists

```
→ fmap (+1) [1, 2, 3] -- [2, 3, 4]
→ fmap (+1) []        -- []
```

### Maybes

```
→ fmap (+1) (Just 1) -- Just 2
→ fmap (+1) Nothing -- Nothing
```

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## Applicative Functor

- A bit more specialised than Functor.
- Useful for certain effective computations.
- Lets us apply a function in a computational context to values in the same context.



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## Applicative Functor

- Another simple API!
- `class Functor  $\varphi \Rightarrow$  Applicative  $\varphi$  where`
- `pure`     $:: \alpha \rightarrow \varphi \alpha$
- `(<*>)`    $:: \varphi (\alpha \rightarrow \beta) \rightarrow \varphi \alpha \rightarrow \varphi \beta$
- Apply  `$\varphi (\alpha \rightarrow \beta)$`  on  `$\varphi \alpha$`  and get  `$\varphi \beta$` .

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## Applicative Functor

→ So for lists, the instance looks like this:

→ instance Applicative [] where

pure = (:[ ])

f:fs <\*> xs = fmap f xs ++ (fs <\*> xs)

[ ] <\*> \_ = [ ]

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## Applicative Functor

- `pure 42 :: [Int]` `-- [42]`
- `pure (*2) <*> pure 21` `-- [42]`
- `pure (*) <*> [2, 20] <*> pure 21` `-- [42, 420]`
- `[(*2), (*20)] <*> pure 21` `-- [42, 420]`
- `[(*2), (*20)] <*> [21, 42]` `-- [42, 84, 420, 840]`
- `pure (*) <*> [2, 20] <*> [2, 20]` `-- [42, 84, 420, 840]`

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## Applicative Functor

- Combining the Applicative API with the Functor API, we get what is called “applicative style” programming.
- $\langle \$ \rangle = \text{fmap}$  -- Applicative style operator  $\text{fmap}$
- $f \langle \$ \rangle a \langle * \rangle b$

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## Applicative Functor

→ (\*) <\$> [2, 20] <\*> pure 21 -- [42, 420]

→ (\*) <\$> [2, 20] <\*> [21, 42] -- [42, 84, 420, 840]

→ And so on.

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## Monad

- A bit more specialised than `Applicative`.
- Useful for computational contexts which may be composed sequentially.
- Lets us apply a function that takes a regular value, and returns a value in a context, to values in the same context.

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## Monad

- The API is, of course, very simple!
  - `class Applicative  $\mu \Rightarrow$  Monad  $\mu$  where`
- |                           |  |
|---------------------------|--|
| <code>return</code>       | <code>:: <math>\alpha \rightarrow \mu \alpha</math></code>   |
| <code>(&gt;&gt;=)</code>  | <code>:: <math>\mu \alpha \rightarrow (\alpha \rightarrow \mu \beta) \rightarrow \mu \beta</math></code> |
| <code>(&gt;&gt;)</code>   | <code>:: <math>\mu \alpha \rightarrow \mu \beta \rightarrow \mu \beta</math></code>                      |
| <code>m &gt;&gt; n</code> | <code>= m &gt;&gt;= <math>\lambda \_ \rightarrow n</math></code>   |

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## Monad

→ **instance Monad Maybe where**

**return = Just**

**Just x >>= f = f x**

**Nothing >>= \_ = Nothing**



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## Monad

- `return 42 :: Maybe Int -- Just 42`
- `doubleIfPos x`
  - `| x > 0 = Just (x * 2)`
  - `| otherwise = Nothing`
- `Just 21 >>= doubleIfPos -- Just 42`
- `Nothing >>= doubleIfPos -- Nothing`

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## do-notation

- Haskell is the world's finest imperative programming language.
- do-notation is a convenient and nice way to program with monads.
- Monads are programmable semicolons!

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## do-notation

→  $a \gg= \lambda x \rightarrow b\ x \gg= \lambda y \rightarrow c\ y$  -- tedious to chain!

→ do

$x \leftarrow a$

$y \leftarrow b\ x$

$c\ y$

-- Ahh... much nicer!

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## do-notation

→ Side-by-side:

```
a    >>= λx →  
b x  >>= λy →  
c y
```

```
do  x ← a  
    y ← b x  
    c y
```

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## do-notation

→ Bit more practical example:

→ do

```
x ← Just 4      -- x is now 4
y ← Just 2      -- y is now 2
let z = x+y      -- z = 6
return (z*(z+1)) -- Just 42
```

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## do-notation

- Currently only for Monad.
- Applicative-do: coming soon to a GHC near you!

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## Monad comprehensions

- Another cute syntax sugar for monads.
- A lot like set comprehension from maths, if you're into that kind of stuff.
- Usually used for lists (called list comprehensions).

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## Monad comprehensions

- `[z*(z+1) | x ← Just 4, y ← Just 2, let z = x+y] -- Just 42`
- `[x+y | x ← [1, 2], y ← [10, 100]] -- [11,101,12,102]`



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## Monad comprehensions

- We can now make the list applicative a bit nicer:  
  
→  $(f:fs) \langle^* \rangle xs = \text{fmap } f \text{ } xs ++ (fs \langle^* \rangle xs)$   
→  $fs \langle^* \rangle xs = [f \ x \mid f \leftarrow fs, x \leftarrow xs]$

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## Foldable & Traversable

# More polymorphism!

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## Foldable

- Functor is for things which may be mapped over, Foldable is for things which may be folded up.
- Folding is also known as reducing, injecting, and various other things in other languages.
- But WTF is folding anyway? Allow me to explain...

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## Foldable

→ Remember **length**?

→ **length**  $:: [\alpha] \rightarrow \text{Int}$

→ **length** [] = 0

→ **length** (\_:xs) = 1 + **length** xs

→ This is actually a variation of a super common pattern!

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## Foldable

Consider these:

- `sum`  $:: [\text{Int}] \rightarrow \text{Int}$
- `sum []`  $= 0$
- `sum (x:xs)`  $= x + \text{sum } xs$

- `reverse`  $:: [\alpha] \rightarrow [\alpha]$
- `reverse []`  $= []$
- `reverse (x:xs)`  $= \text{reverse } xs ++ [x]$

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## Foldable

- $(++)$   $:: [\alpha] \rightarrow [\alpha] \rightarrow [\alpha]$
- $[] ++ ys = ys$
- $(x:xs) ++ ys = x : xs ++ ys$
- $filter$   $:: (\alpha \rightarrow \text{Bool}) \rightarrow [\alpha] \rightarrow [\alpha]$
- $filter\ p\ [] = []$
- $filter\ p\ (x:xs)$ 
  - |  $p\ x$   $= x : filter\ p\ xs$
  - | otherwise  $= filter\ p\ xs$

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## Foldable

And what about these chaps?

- `map`  $:: (\alpha \rightarrow \beta) \rightarrow [\alpha] \rightarrow [\beta]$
- `map _ []`  $= []$
- `map f (x:xs)`  $= f\ x : \text{map } f\ xs$
- `id`  $:: [\alpha] \rightarrow [\alpha]$
- `id x`  $= x$

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## Foldable

→ These are all just variations on a theme encapsulated by what we call a fold.

→ `foldr`  $:: (\alpha \rightarrow \beta \rightarrow \beta) \rightarrow \beta \rightarrow [\alpha] \rightarrow \beta$   
→ `foldr _ z []`  $= z$   
→ `foldr f z (x:xs)`  $= f\ x\ (foldr\ f\ z\ xs)$



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## Foldable

- `length` = `foldr (const (1 +)) 0`
- `sum` = `foldr (+) 0`
- `reverse` = `foldr (flip (++) . return) []`
- `xs ++ ys` = `foldr (:) ys xs`
- `filter p` = `foldr (\x xs → if p x then x : xs else xs) []`
- `map f` = `foldr ((:) . f) []`
- `id` = `foldr (:) []`

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## Foldable

- So that's how you fold lists.
- But, surely!, you won't be surprised to learn that lists aren't all that special.
- We can fold anything that has a Foldable instance!

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## Foldable

- The API looks a bit more complicated this time around.
- Don't be scared though! To make a Foldable, you only need to implement one method! You get the others for free.
- Your pick between **foldMap** and **foldr**.

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## Foldable

→ class Foldable  $\tau$  where

fold :: Monoid  $\mu \Rightarrow \tau \mu \rightarrow \mu$

foldMap :: Monoid  $\mu \Rightarrow (\alpha \rightarrow \mu) \rightarrow \tau \alpha \rightarrow \mu$

foldr ::  $(\alpha \rightarrow \beta \rightarrow \beta) \rightarrow \beta \rightarrow \tau \alpha \rightarrow \beta$

foldl ::  $(\alpha \rightarrow \beta \rightarrow \alpha) \rightarrow \alpha \rightarrow \tau \beta \rightarrow \alpha$

foldr1 ::  $(\alpha \rightarrow \alpha \rightarrow \alpha) \rightarrow \tau \alpha \rightarrow \alpha$

foldl1 ::  $(\alpha \rightarrow \alpha \rightarrow \alpha) \rightarrow \tau \alpha \rightarrow \alpha$

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## Foldable

- instance Foldable Maybe where
  - foldr \_ z Nothing = z
  - foldr f z (Just x) = f x z
- instance Foldable ((,) a) where
  - foldr f z (\_, y) = f y z

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## Traversable

- Traversable represents data structures which can be traversed while preserving the shape.
- A Foldable Functor that lets us commute two functors.

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## Traversable

- Like Foldable, what initially looks to be a semi-complicated API.
- Like Foldable, only one function needs to be implemented. Pick between **sequence** or **traverse**.

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## Traversable

→ `class (Functor  $\tau$ , Foldable  $\tau$ ) => Traversable  $\tau$  where`

<code>traverse</code>	<code>:: Applicative <math>\varphi</math></code>	<code>=&gt; (<math>\alpha \rightarrow \varphi \beta</math>) <math>\rightarrow \tau \alpha \rightarrow \varphi (\tau \beta)</math></code>
<code>sequenceA</code>	<code>:: Applicative <math>\varphi</math></code>	<code>=&gt; <math>\tau (\varphi \alpha) \rightarrow \varphi (\tau \alpha)</math></code>
<code>mapM</code>	<code>:: Monad <math>\mu</math></code>	<code>=&gt; (<math>\alpha \rightarrow \mu \beta</math>) <math>\rightarrow \tau \alpha \rightarrow \mu (\tau \beta)</math></code>
<code>sequence</code>	<code>:: Monad <math>\mu</math></code>	<code>=&gt; <math>\tau (\mu \alpha) \rightarrow \mu (\tau \alpha)</math></code>



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## Traversable

- `instance Traversable Maybe where`  
    `traverse _ Nothing = pure Nothing`  
    `traverse f (Just x) = Just <$> f x`
- `instance Traversable ((,) a) where`  
    `traverse f (x, y) = (,) x <$> f y`

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## Traversable

- `sequence [Just 42] -- Just [42]`
- `sequence Just [42] -- [Just 42]`

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## Traversable

- `doubleIfPos <$> [-5, 21]` -- `[Nothing, Just 42]`
- `doubleIfPos <$> [21, 42]` -- `[Just 42, Just 84]`
- `doubleIfPos `traverse` [-5, 21]` -- `Nothing`
- `doubleIfPos `traverse` [21, 42]` -- `Just [42, 84]`
- `sequence $ doubleIfPos <$> [-5, 21]` -- `Nothing`
- `sequence $ doubleIfPos <$> [21, 42]` -- `Just [42, 84]`