

# CS & IT ENGINEERING

Discrete Maths  
Graph Theory



Lecture No. 09



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## TOPICS TO BE COVERED

01 Analysis In Connectivity

02 Various definition in Connectivity

03 Edge Connectivity

04 Vertex Connectivity

05 Largest inequality theorem

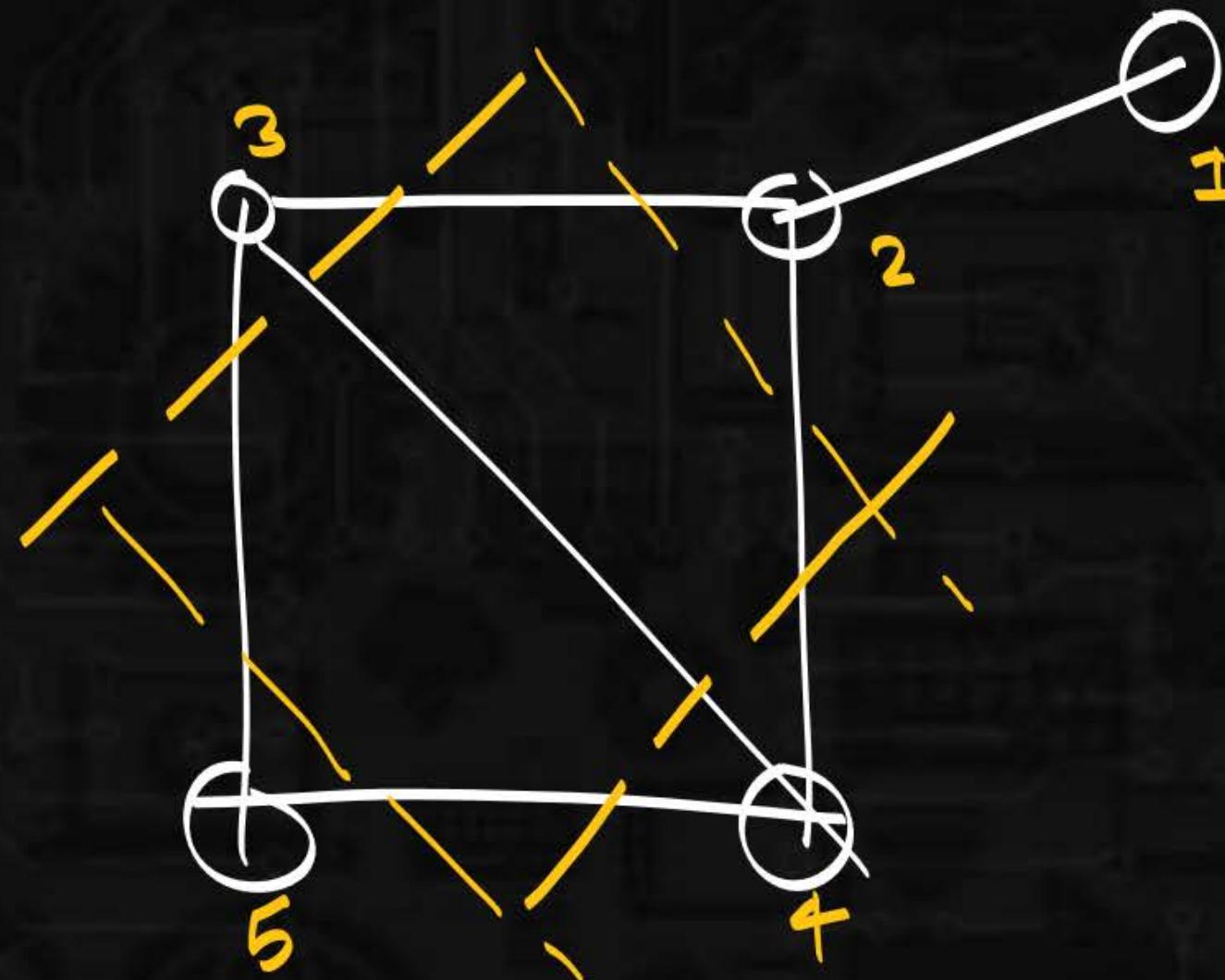
# Connectivity in Graphs

cut edge  
↳ 1 edge.

\* cut set  
↳ set of edges.  
1 → cut edge.  
2  
3  
4

edge connectivity  
↳ min + set of edges.

# Connectivity in Graphs



edge can  
 $\Delta(G) = 1$

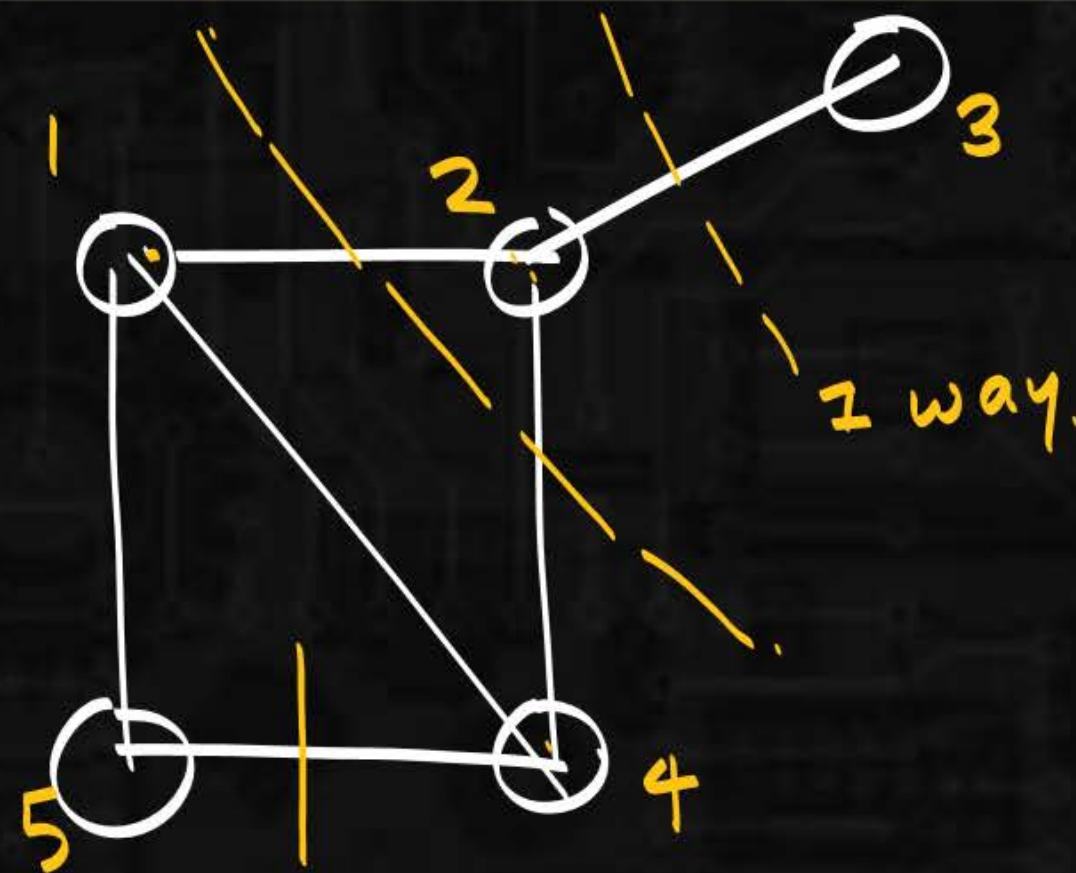
**Cut set**  
Size 1  $\rightarrow \{1, 2\}$  cut edge.

Scan.  
min.

Size 2  $\rightarrow \{3, 2, 2+1\}$   
 $\rightarrow \{3, 5, 5+1\}$

Size 3  $\rightarrow \{ \}$

# Connectivity in Graphs



proper subset

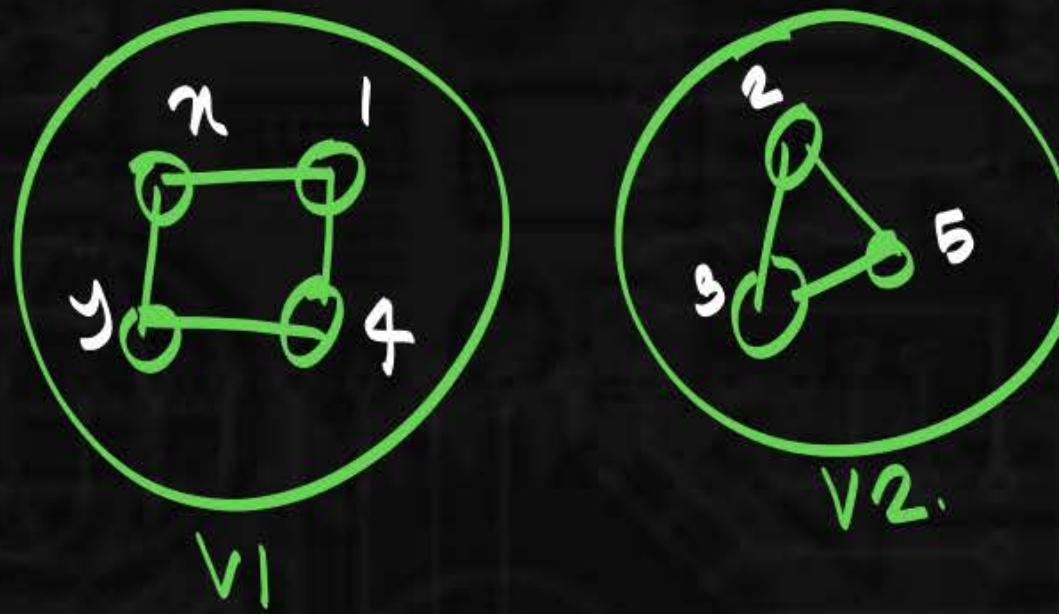
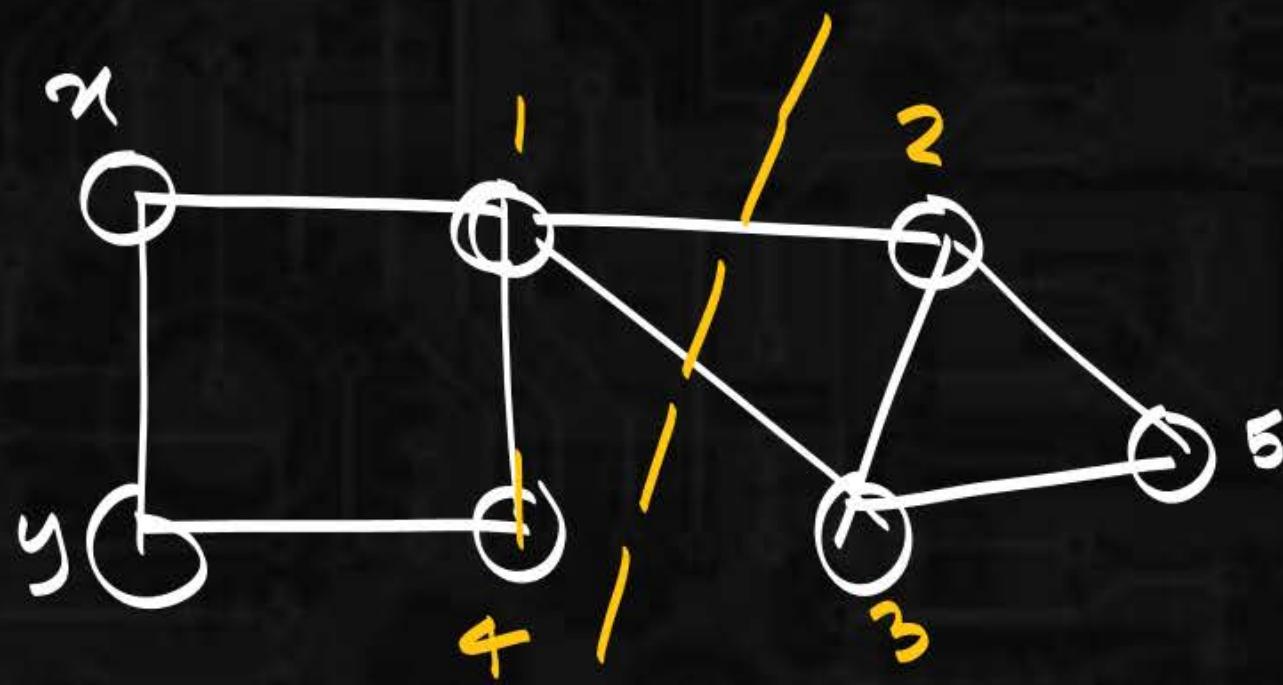
$\{1, 2, 2 \leftarrow\}$  cutset

$\{1, 2, 2 \leftarrow, 4, 5\} \rightarrow$

cutset

not cutset

# Connectivity in Graphs



{ {1, 2, 13} } → cut-set ✓  
{ {2, 13, 14} } not cut-set ?  
Cutset

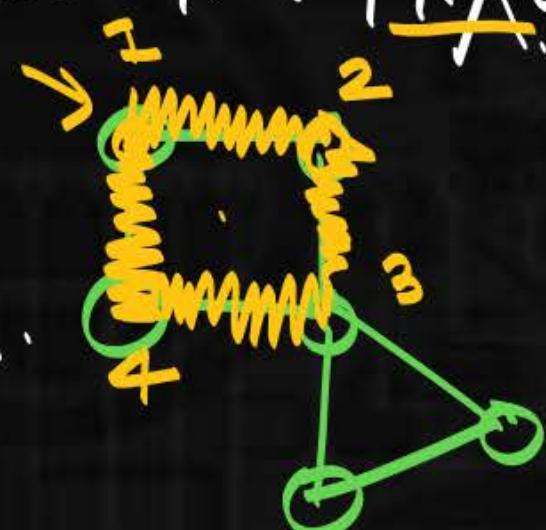
# Connectivity in Graphs

Euler Graph :

Trail: alternating sequences of vertices & edges  $R \cdot V | R \times E$

① - 2 - 3 - 4 - ①

Closed Trail: Trail + starting = ending vertex.

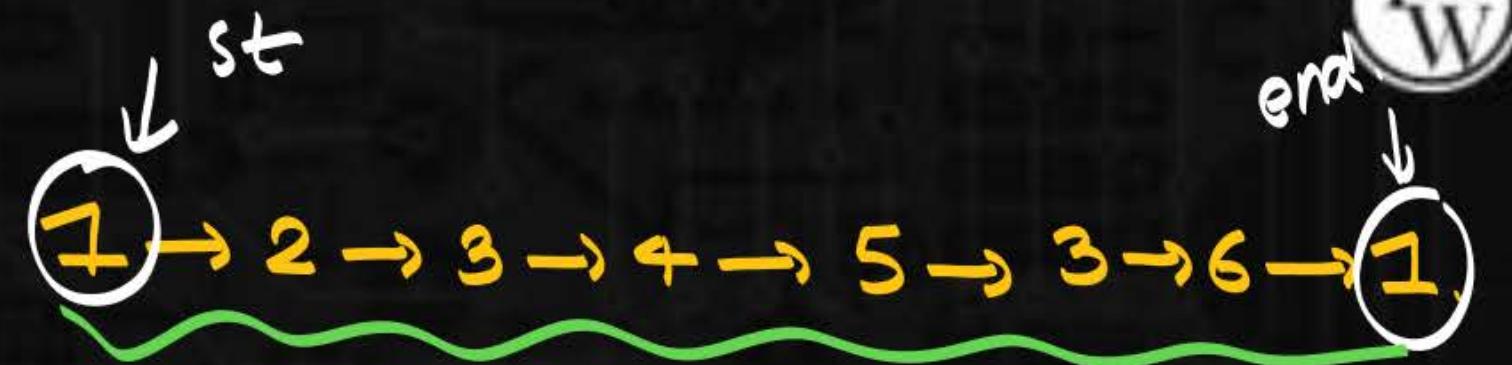


Euler circuit : Closed Trail + cover all edges exactly once.

# Connectivity in Graphs

Euler ckt.

{ Closed Trail.  
{ all edges.



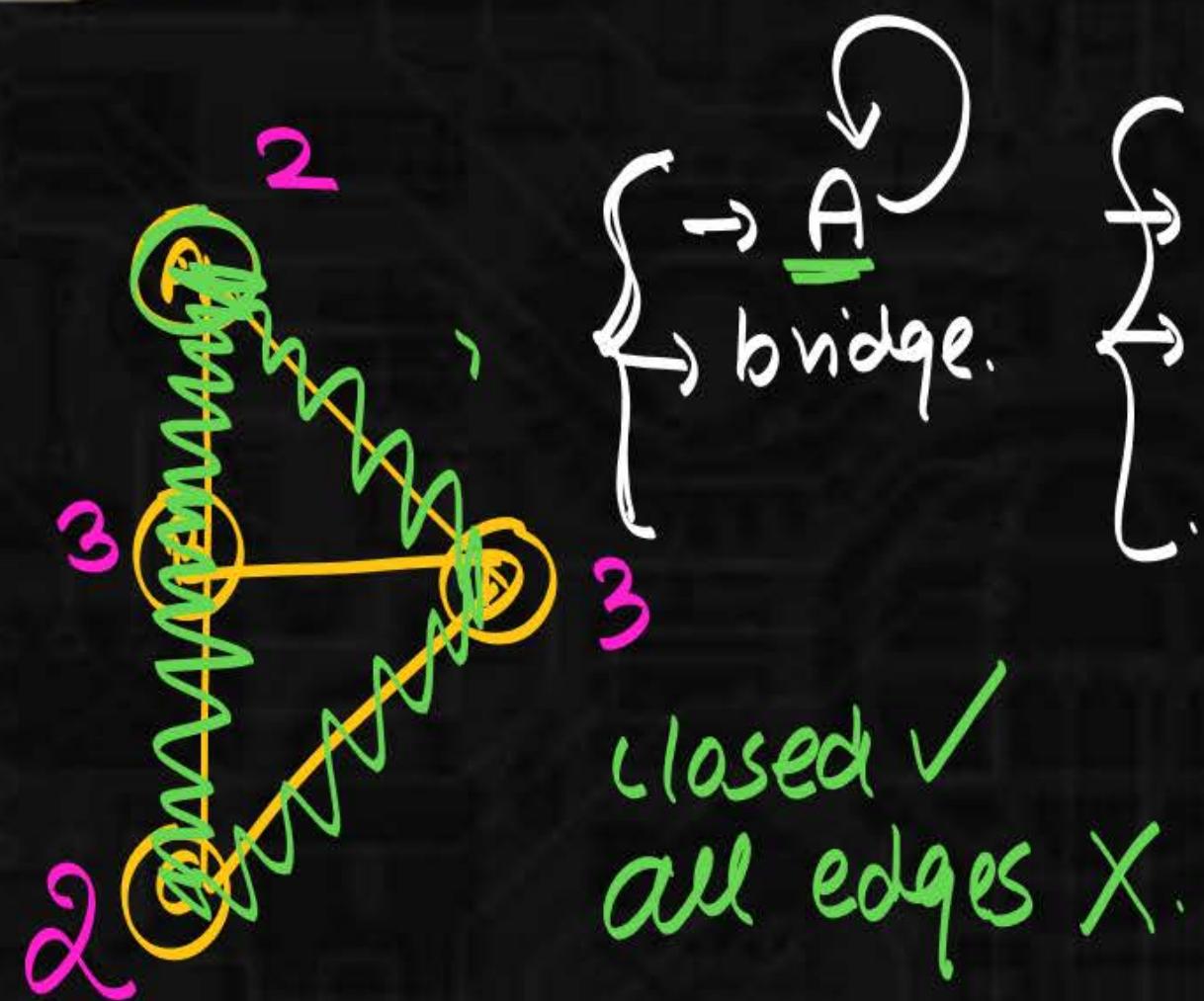
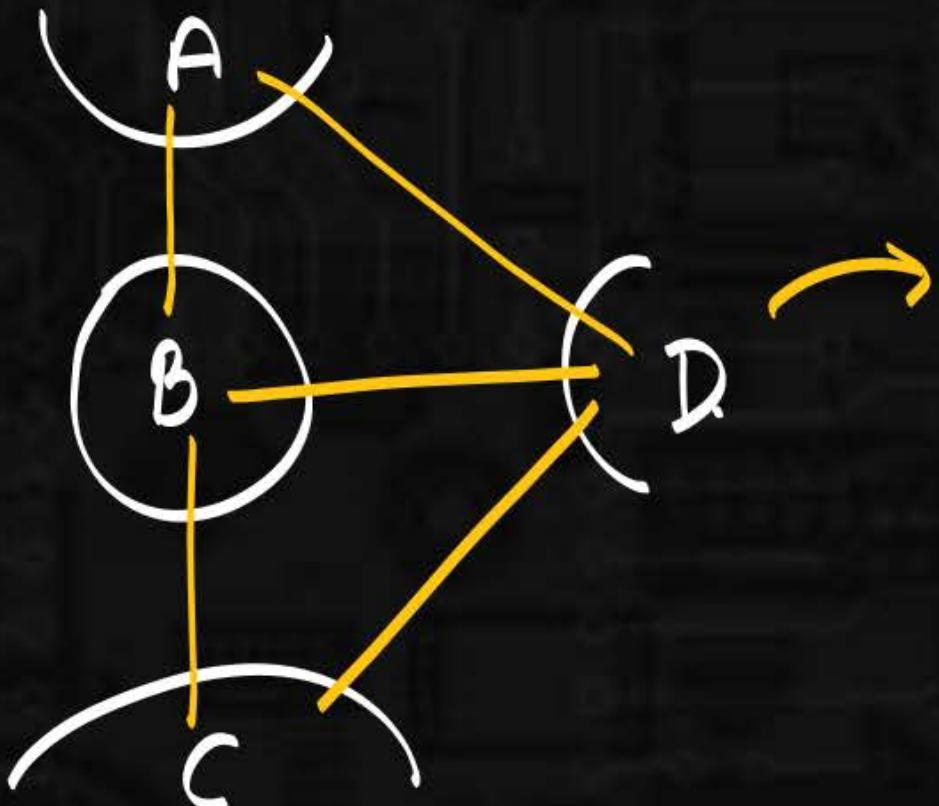
Repstn  $\rightarrow$  Euler ckt

$$G = (V, E)$$

Euler Graph.

# Connectivity in Graphs

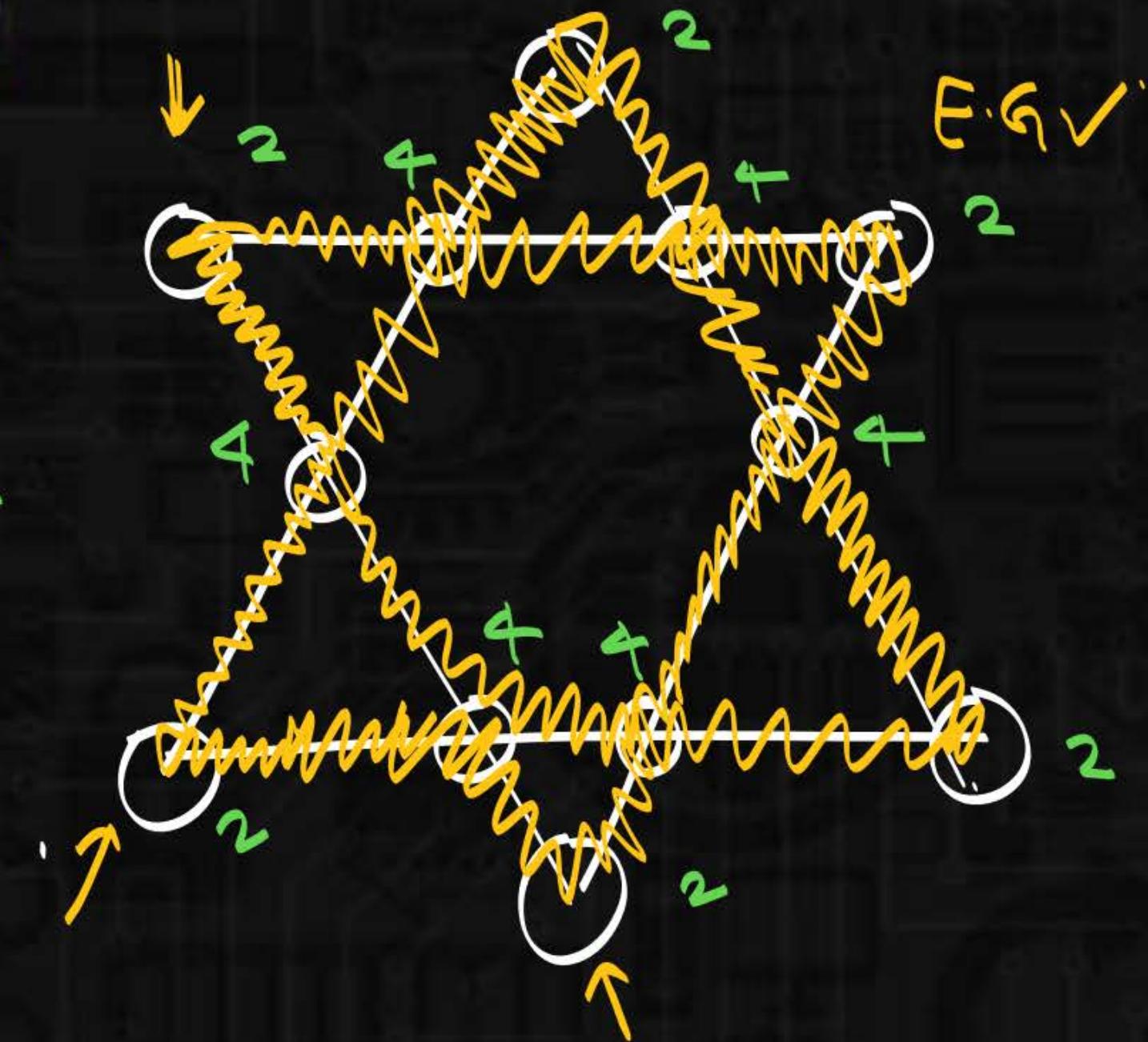
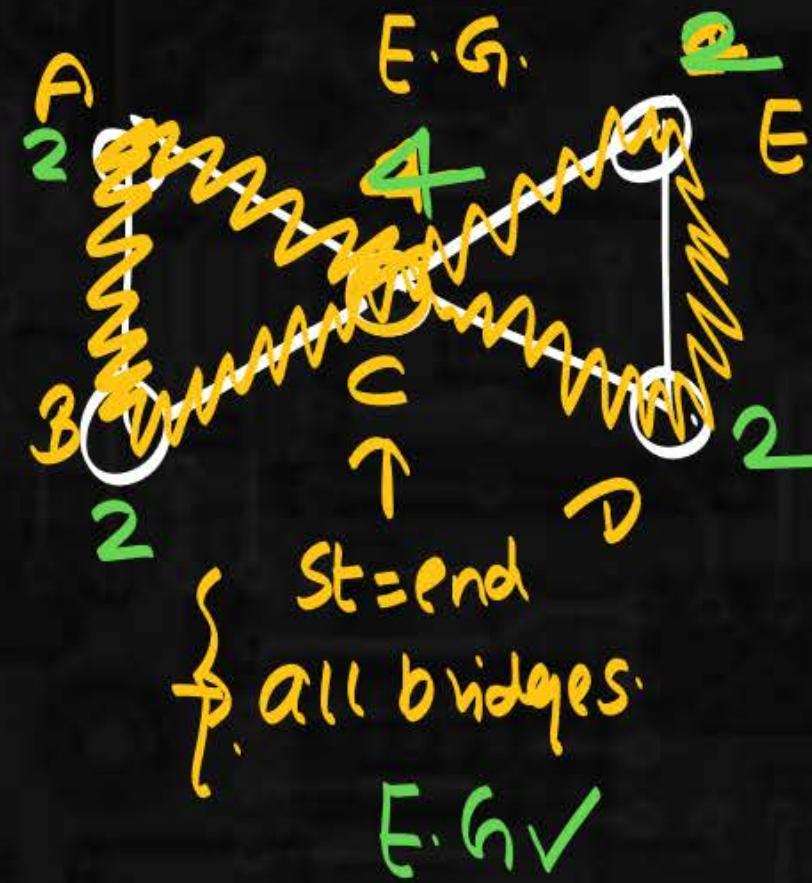
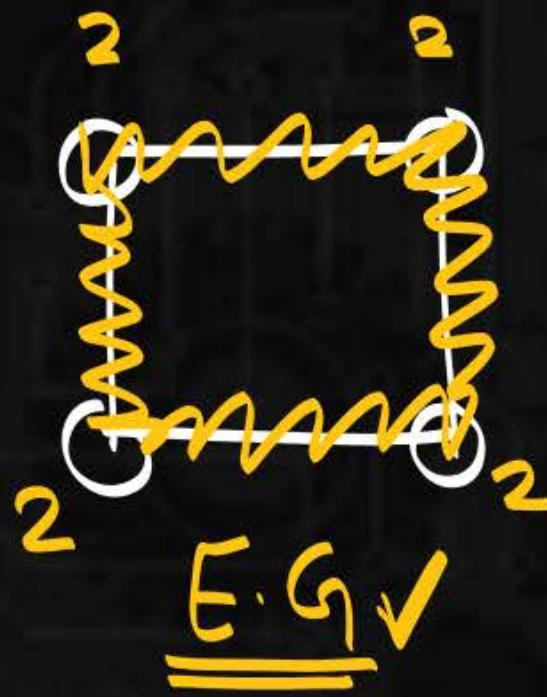
1736:



closed ✓  
all edges X.

Euler Graph  
Closed Trail.  
all edges.

# Connectivity in Graphs



# Connectivity in Graphs

non trivial connected graph is Euler Graph iff degrees of all vertices are even.



Degree = 0  
even degree.

Trivial  $\rightarrow$  Default



Initial Graph:

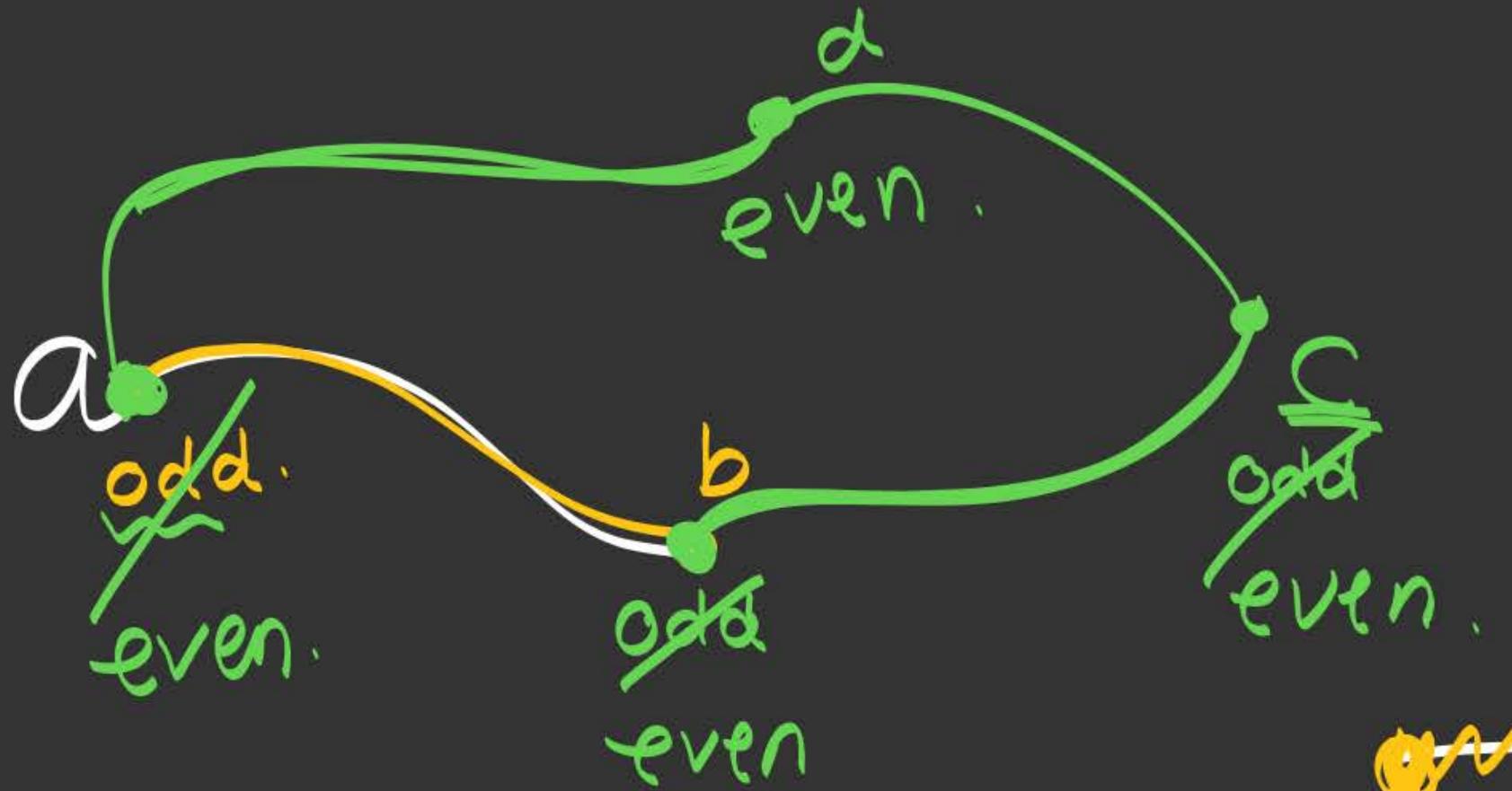
$$G = (V, E)$$

$\downarrow$        $\downarrow$   
set      set

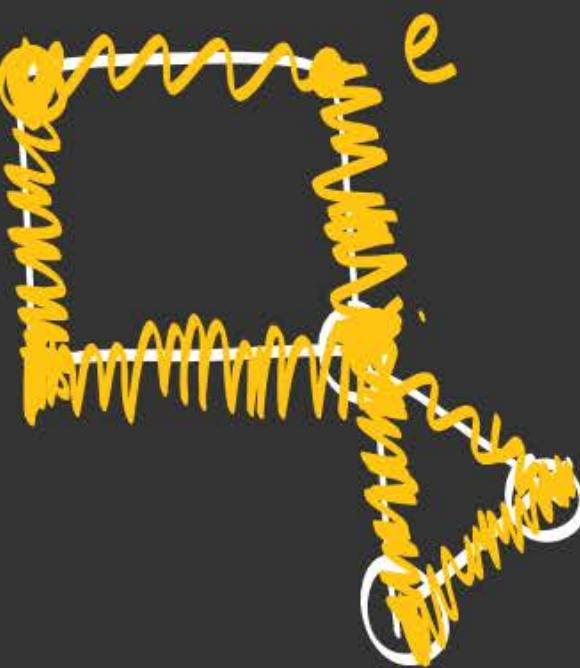
$$\left\{ \begin{array}{l} V \neq \emptyset \\ E \neq \emptyset \end{array} \right. \rightarrow$$

A diagram showing a single black circle representing a node. Inside the circle is a small yellow dot representing the vertex itself. There are no edges or self-loops present.

Initial Graph:



{ closed ✓  
all edges ✓  
Stop  $\rightarrow a$ .



## Connectivity in Graphs

Graph is Euler Graph iff degrees of all vertices are even.

+

Connected

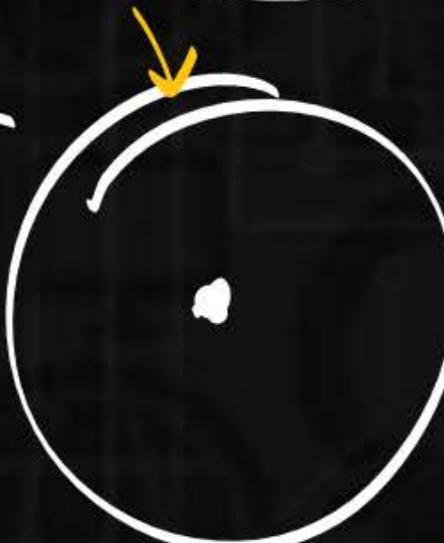
+

non trivial

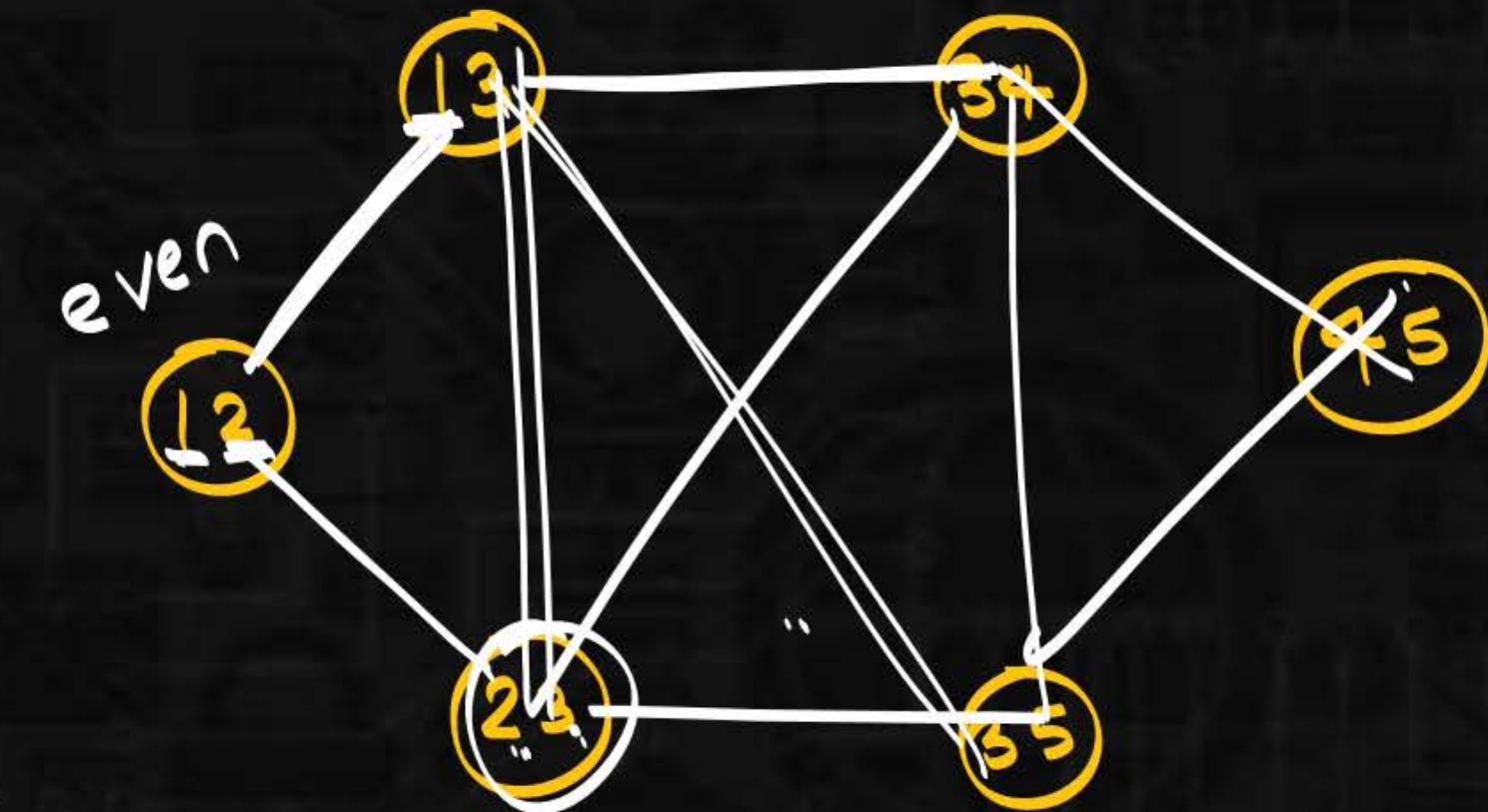
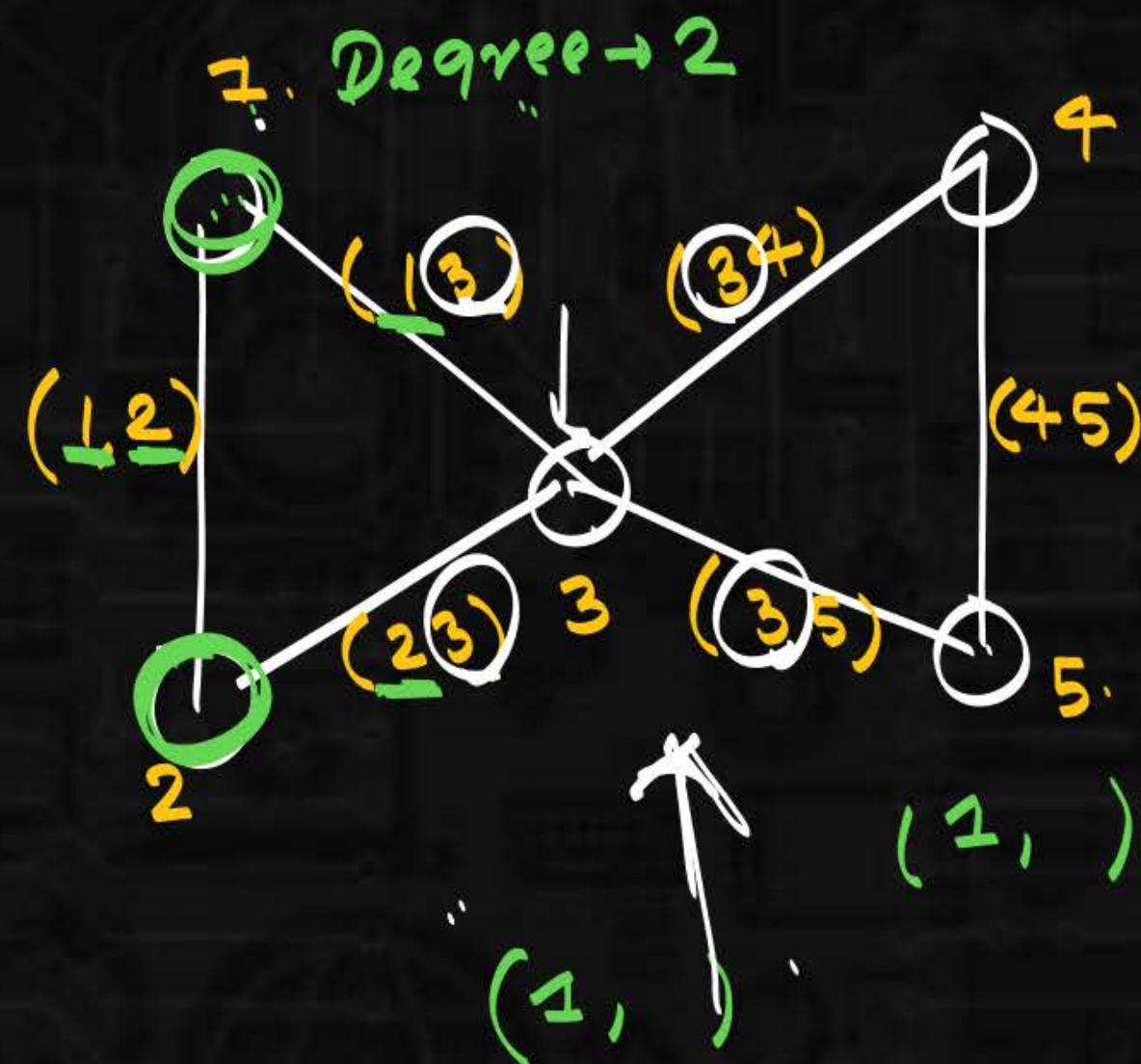
Doubt 1:



Doubt 2:

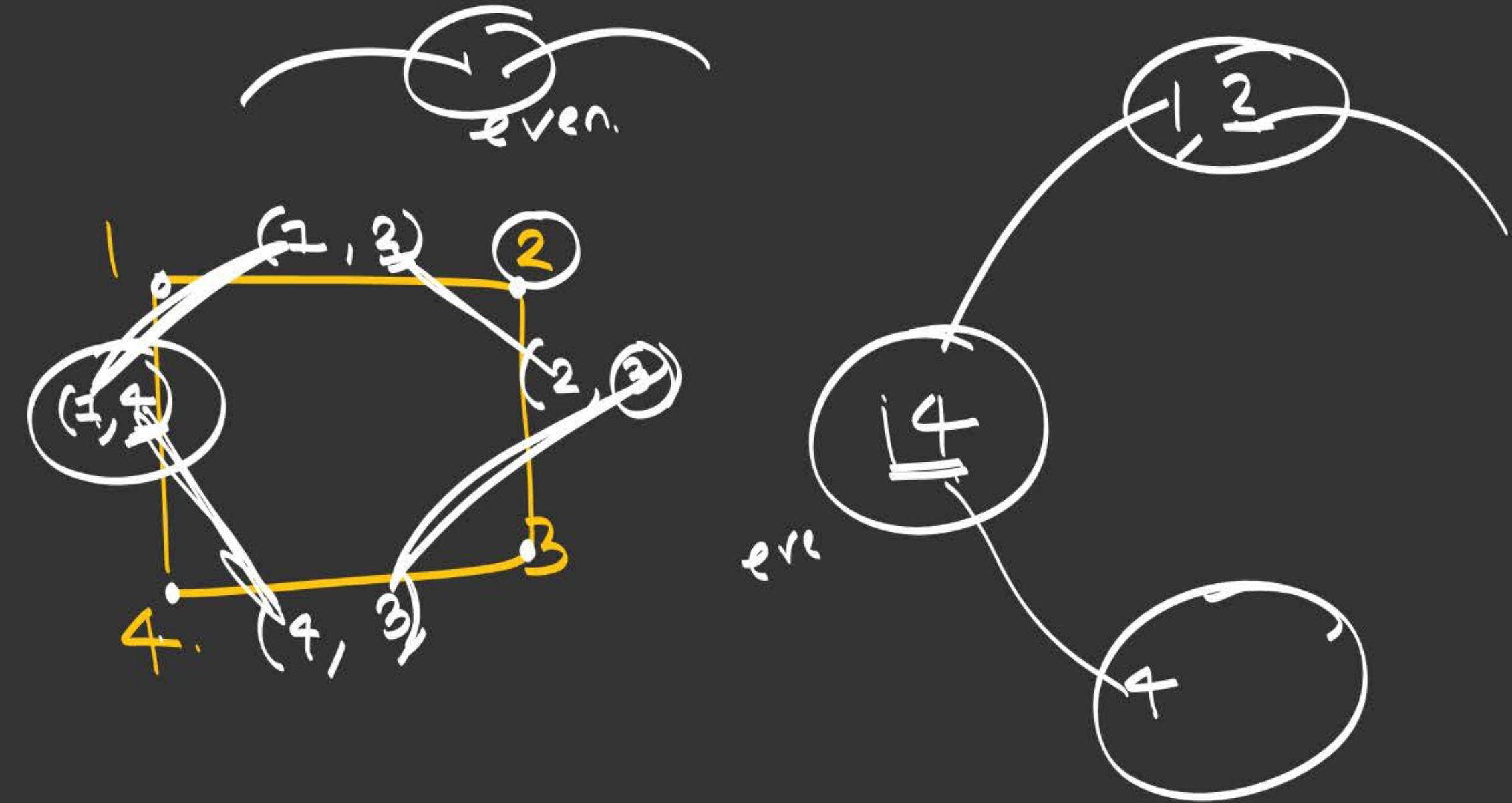
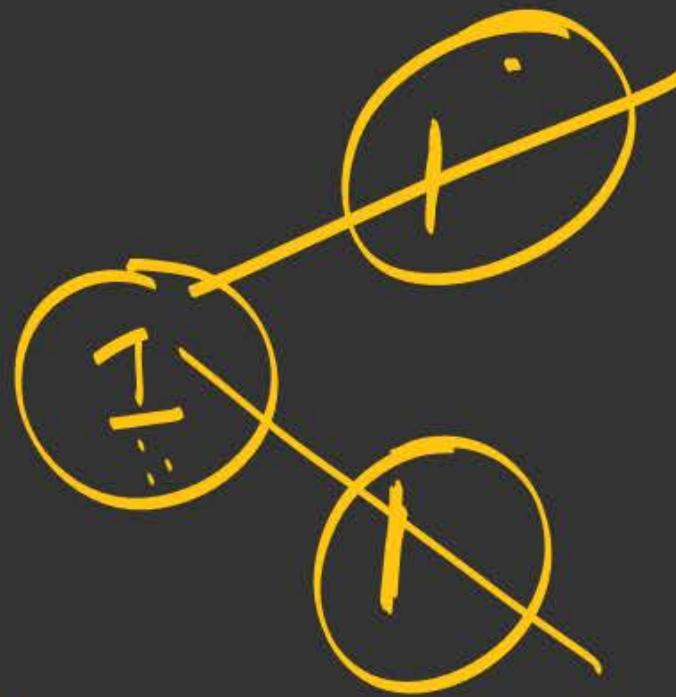


# Connectivity in Graphs



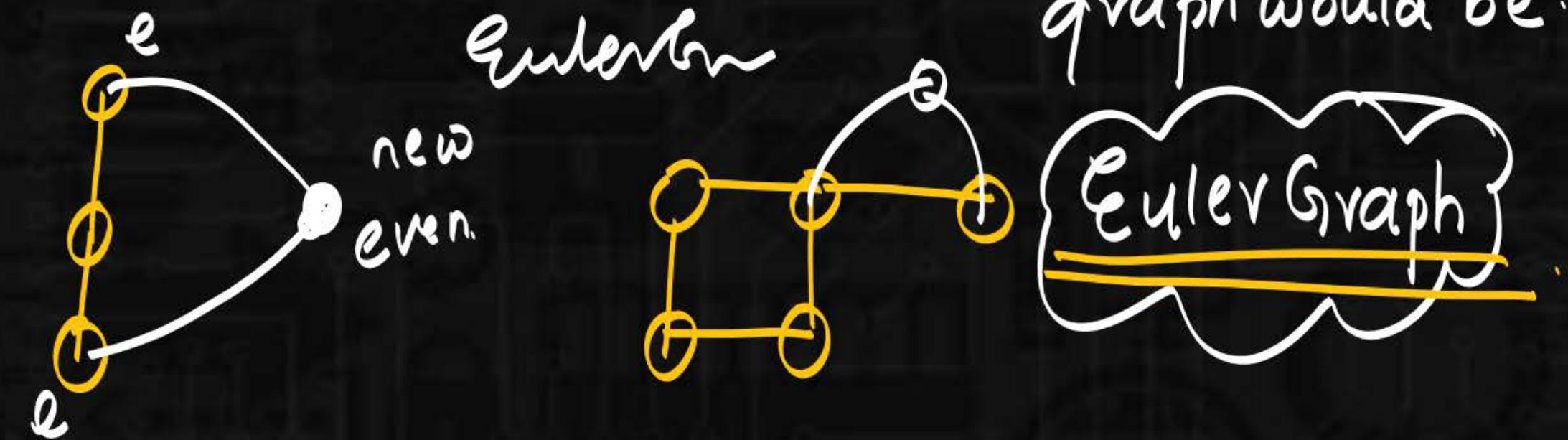
Line Graph of Euler Graph will always be  
Euler Graph.

Degrees → 2  
2 paths



# Connectivity in Graphs

Thm 2:  
Consider a connected graph having some odd degree vertices  
add new vertex make it connect to all odd degree  
(GATE) vertices, then the resultant  
graph would be.



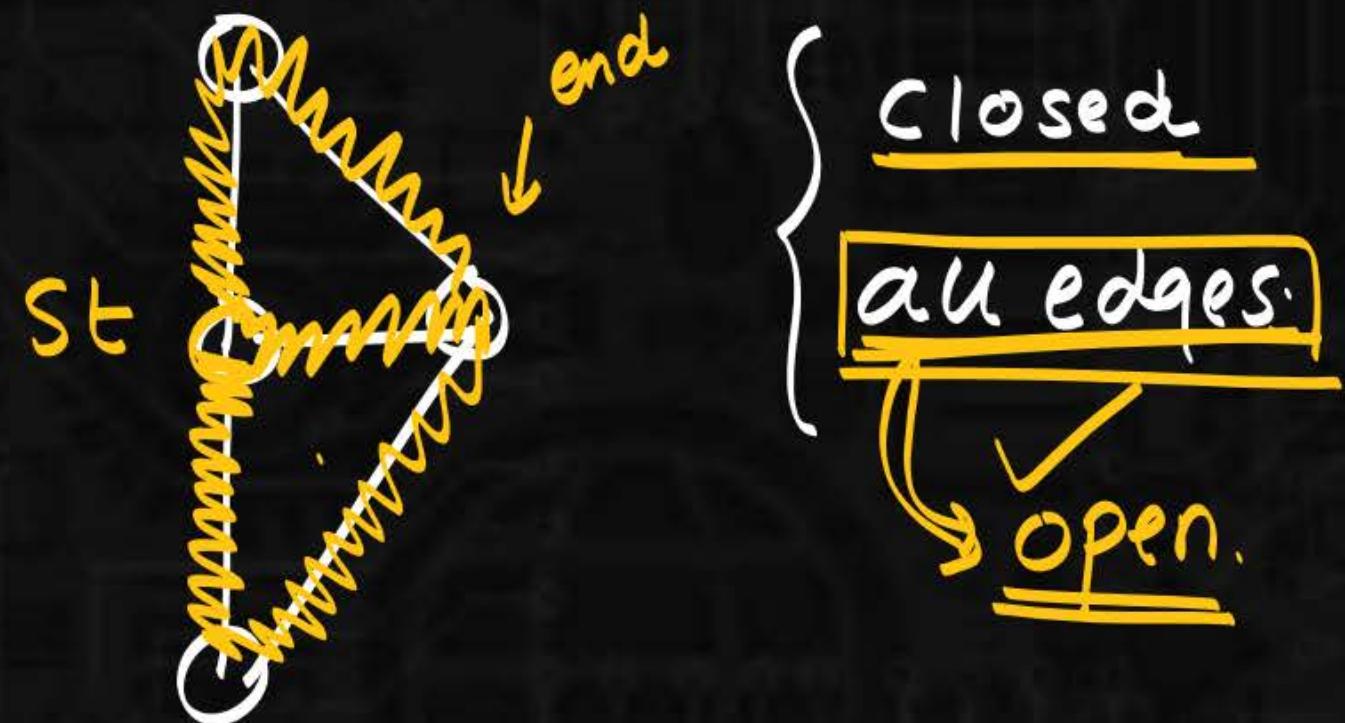
# Connectivity in Graphs

Trail:

open Trail: Trail + st  $\notin$  ending vertex

Euler line: open Trail + covering all edges exactly once.

Unicursal line  
Euler Path



# Connectivity in Graphs



Eg x.

All edges ✓  
+  
open.

# Connectivity in Graphs

covering all edges.

Closed .

Euler ckt

Open .

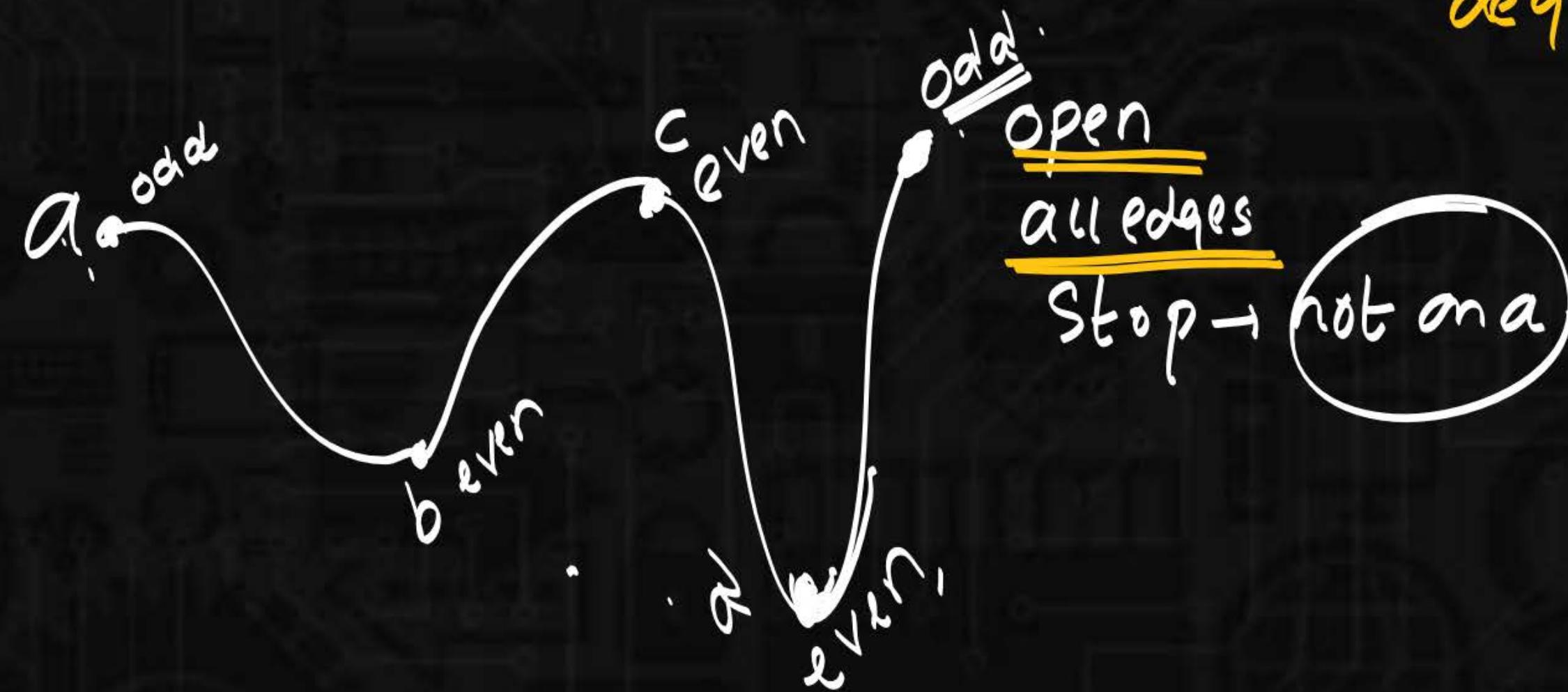
Eulerline / EulerPath :

## Connectivity in Graphs

Note: Start  $\rightarrow$  odd degree  
End  $\rightarrow$  odd degree

Theorem:

Graph contains Euler line iff it contains exactly 2 odd degree vertices.



# Connectivity in Graphs



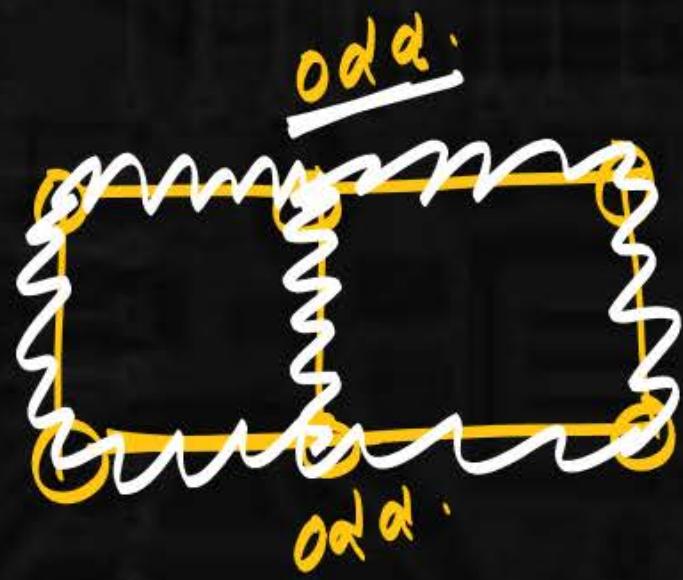
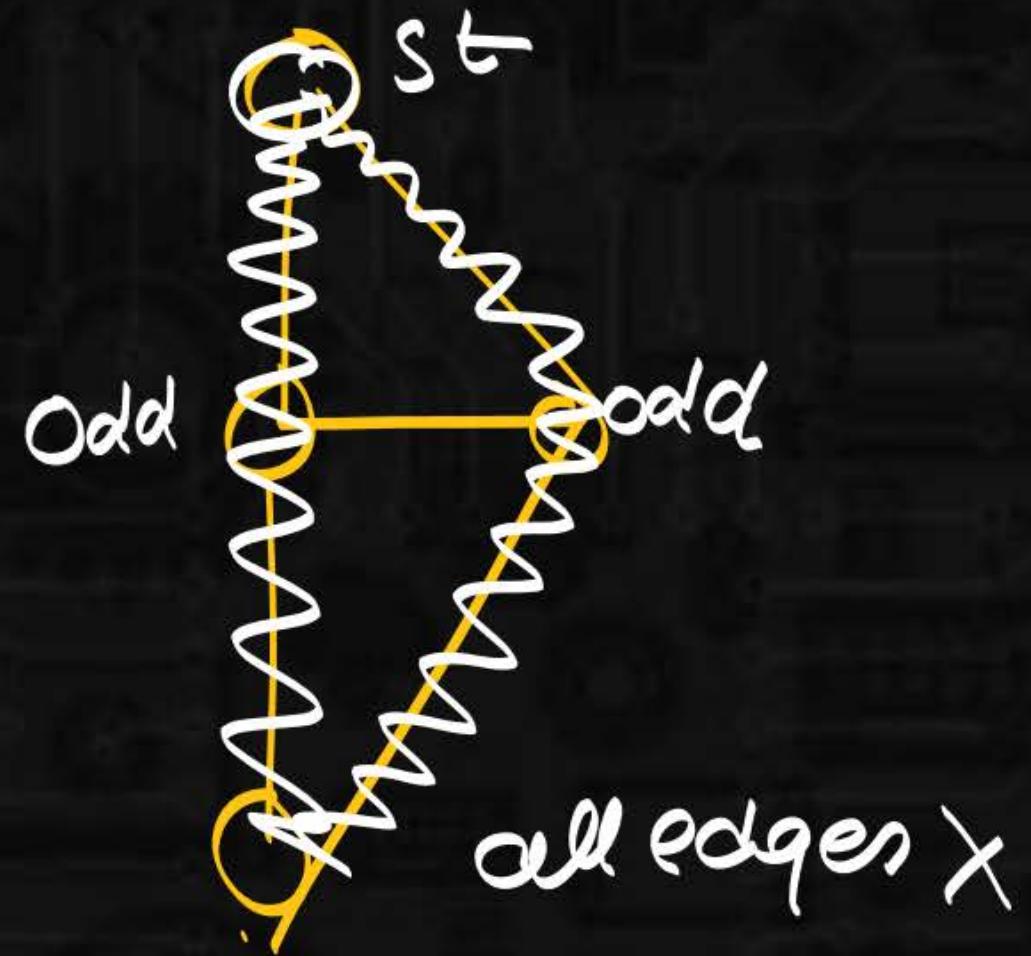
St  $\rightarrow$  a.

Stop  $\rightarrow$  b

all edges.



# Connectivity in Graphs



# Connectivity in Graphs

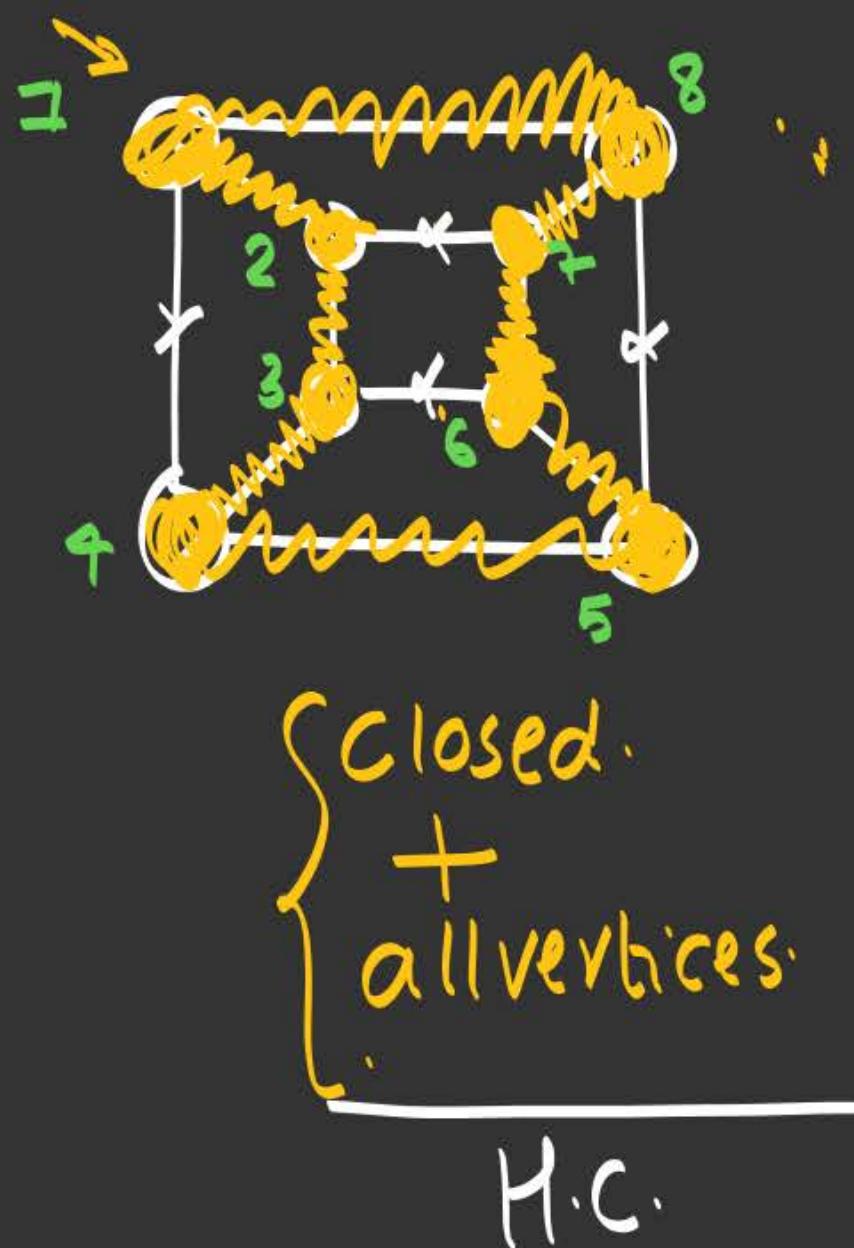
Path :

Closed Path : Path +  $s = t$  ending vertex.

Hamiltonian ckt : Closed Path + cover all vertices exactly once.

→ Starting ✓  
✓  
✓

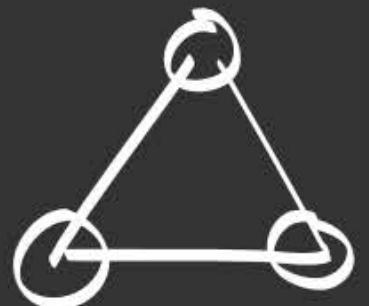
{ closed.  
+  
all vertices.  
H.C.



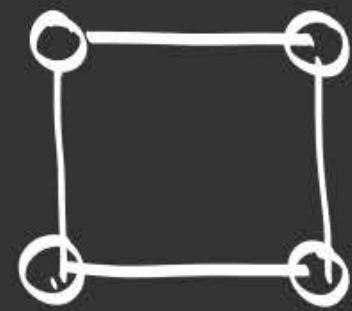
Graph contains Hamiltonian  
Ckt then Graph is called.  
Hamiltonian Graph.

H.C.  
1 - 2 - 3 - 4 - 5 - 6 - 7 - 8 - 1

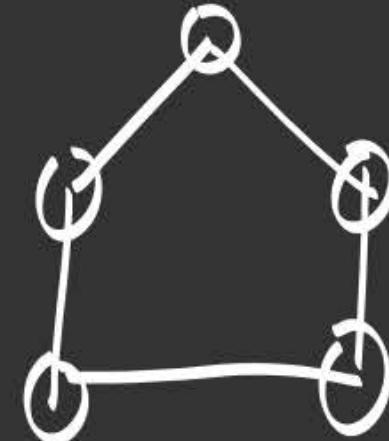
$$G = (V, E)$$



H.C ✓

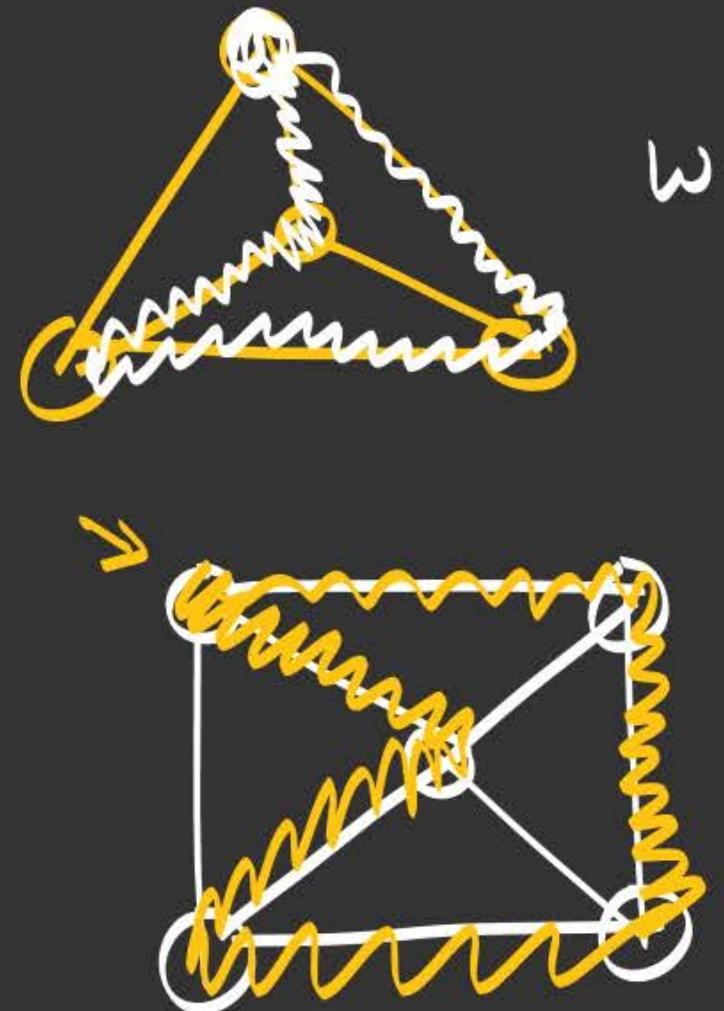


H.C ✓



H.C ✓

all cn are H.G.



$\omega \leftarrow H.G.$

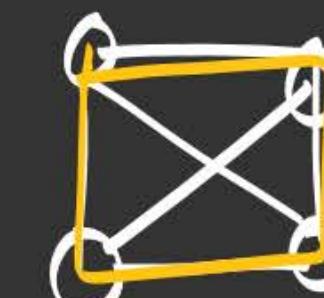
all  $\omega_n \rightarrow H.G.$



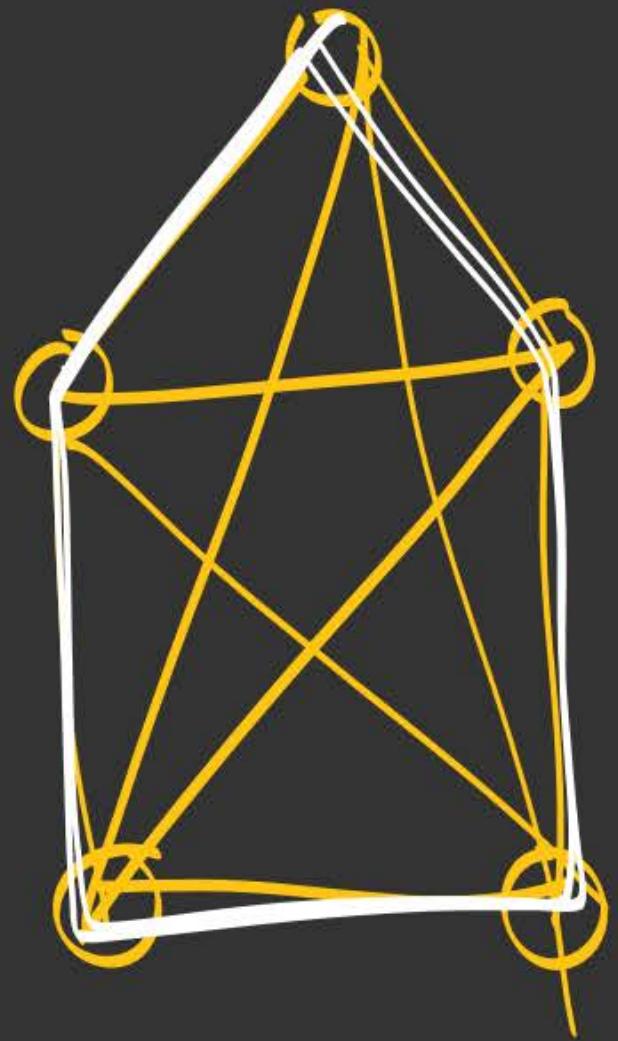
$\kappa_1$   
X.  
 $\kappa_2$ .



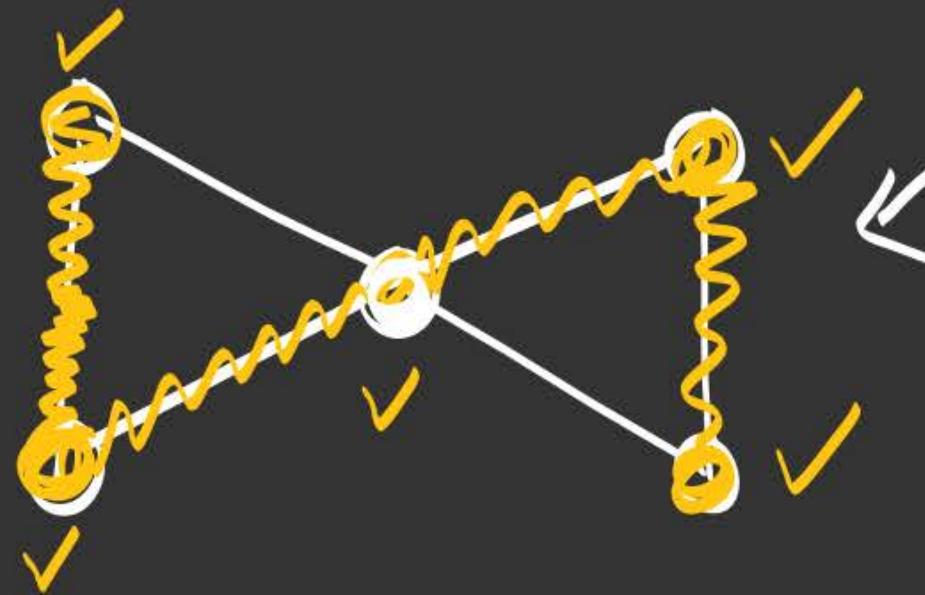
$\kappa_3$  ✓



H.G ✓



all  $K_n \rightarrow H \cdot G$  ( $n \geq 3$ )



Hamiltonian Path. :

open Path + covers all vertices  
exactly once.

not Hamiltonian Graph.

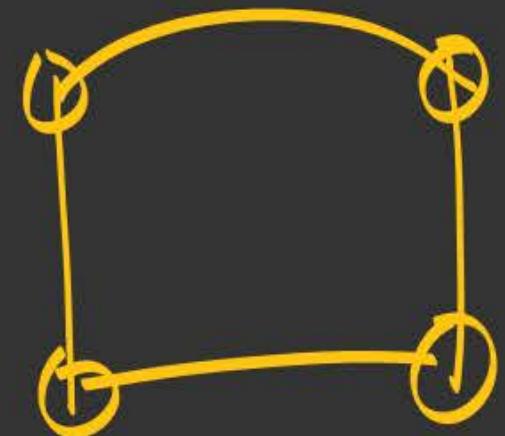
{ all vertices ✓  
closed ✗



all vertices ✓

Closed.

Every Hamiltonian ckt  
contains H-path. ✓



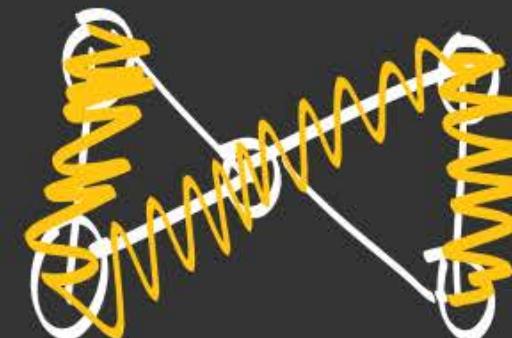
all vertices

open

→ Every Hamiltonian path

contains Hamiltonian ckt.

(false)



# Covering

