

# Divide and Conquer - Complete Notes (Zero to Everything)

## 1. Intuition / Idea

Think of **breaking the problem into smaller pieces**, solving each independently, and **combining** the result.

Just like Ramayan mein **Vanar Sena** alag-alag dikh gayi thi Lanka tak jaane ke liye, par goal ek hi tha! 😊

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## 2. Steps / Strategy

*Remember this 3-step:*

### 1. Divide

Break the main problem into **smaller subproblems**.

### 2. Conquer

Solve each subproblem **recursively**.

### 3. Combine

Merge or combine solutions of subproblems to get the **final answer**.

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## 3. Time Complexity Format

Let's say problem size is  $n$

You divide it into  $k$  subproblems, each of size  $n/b$ , and combining takes  $O(n^d)$

Then:

$$T(n) = k * T(n/b) + O(n^d)$$

Use **Master Theorem** to solve this recurrence (explained below 📌).

1  
2  
3  
4

## 4. Master Theorem (Time Analysis Shortcut)

For recurrence:

$$T(n) = aT(n/b) + O(n^d)$$

We compare  $n^d$  vs  $n^{\log_b(a)}$

Now 3 cases:

Case	Condition	Time Complexity
1	$d < \log_b(a)$	$O(n^{\log_b(a)})$
2	$d == \log_b(a)$	$O(n^d \cdot \log n)$
3	$d > \log_b(a)$	$O(n^d)$

💡 “Smaller dominates  $\rightarrow \log$ , Equal  $\rightarrow \log n$ , Larger dominates  $\rightarrow \text{polynomial}$ ”



## 5. Famous Examples of Divide and Conquer

Problem	Divide Step	Combine Step	Time Complexity
✂ Merge Sort	Split array in half	Merge two sorted arrays	$O(n \log n)$
🔍 Binary Search	Halve the array	No combine	$O(\log n)$
★ Quick Sort	Partition the array	Nothing	$O(n \log n)$ avg
🔴 Maximum Subarray (Kadane Alt)	Divide into left & right halves	Combine crossing subarray	$O(n \log n)$
🟢 Matrix Multiplication	Split matrix into 4 submatrices	Add resulting products	$O(n^3)$ or better
⚡ Strassen's Algorithm	7 multiplications instead of 8	Matrix operations	$O(n^{2.81})$
📍 Closest Pair of Points	Divide point set	Find closest pair across mid	$O(n \log n)$

Problem	Divide Step	Combine Step	Time Complexity
📁 Karatsuba Multiplication	Split numbers	Combine result with shifts	$O(n^{\log_2 3}) \approx O(n^{1.58})$

## 🧠 6. Recursion Tree Insight

Divide and conquer problems can be visualized as **recursion trees**, where:

- Each level has **more subproblems**,
- Combine step may take **linear time** per level.

🕒 Depth of tree =  $\log(n)$  if divided by 2 every time.

⌘ Total work = work per level \* number of levels.

## 👤 7. Java Example: Merge Sort

```
public class MergeSort {  
    public static void mergeSort(int[] arr, int left, int right) {  
        if (left < right) {  
            int mid = left + (right - left) / 2;  
  
            // Divide  
            mergeSort(arr, left, mid);  
            mergeSort(arr, mid + 1, right);  
  
            // Conquer + Combine  
            merge(arr, left, mid, right);  
        }  
    }  
  
    public static void merge(int[] arr, int left, int mid, int right) {  
        int[] leftArr = Arrays.copyOfRange(arr, left, mid + 1);  
        int[] rightArr = Arrays.copyOfRange(arr, mid + 1, right + 1);  
  
        int i = 0, j = 0, k = left;
```

```

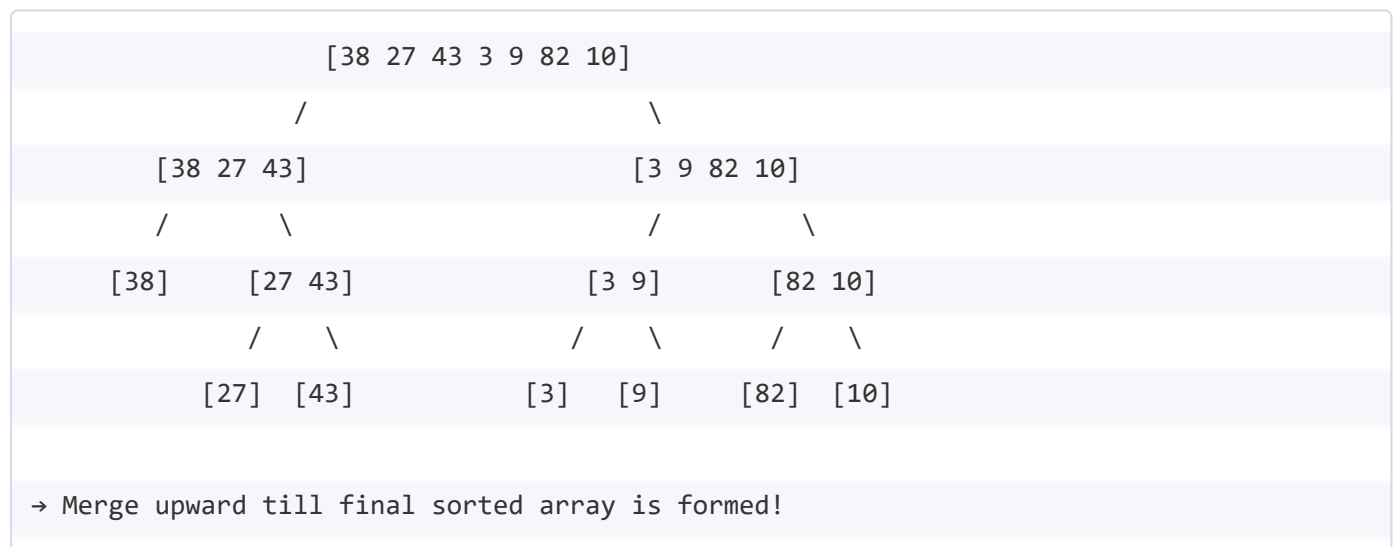
while (i < leftArr.length && j < rightArr.length) {
    if (leftArr[i] <= rightArr[j]) {
        arr[k++] = leftArr[i++];
    } else {
        arr[k++] = rightArr[j++];
    }
}

while (i < leftArr.length) arr[k++] = leftArr[i++];
while (j < rightArr.length) arr[k++] = rightArr[j++];
}
}

```

## 🎨 8. Visualization Example – Merge Sort

Let's sort: [38, 27, 43, 3, 9, 82, 10]



## ⚖️ 9. Pros & Cons

Pros ✅	Cons ❌
Breaks problem into manageable pieces	Recursion → Stack overhead
Often optimal in time complexity	Harder to debug sometimes
Ideal for parallelization 💻	May need extra space (e.g., Merge Sort)



## 10. When to Use Divide and Conquer

✓ When:

- Problem can naturally be **split into subproblems**
- Subproblems are **independent**
- You can **combine easily**

✗ Avoid When:

- Subproblems are tightly connected
- Combine step is very complex



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## Bonus: Strassen's Matrix Multiplication

- Standard:  $O(n^3)$
  - Strassen's: Reduces multiplications from 8  $\rightarrow$  7
  - Time:  $O(n^{\log_2 7}) \approx O(n^{2.81})$
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