

Introduction:

We've recently seen a few things about grammars, parsing, derivation and how to decide and design different elements of languages. Demonstrate your knowledge and facility with these mechanisms below.

Methodology:

Submit a PDF with your answers to the below on Sakai on or before 2355 23 September.

Note: in all of the below capital letters are non-terminals, non-grammar symbols are terminals.

0. Given the following grammar, G:

$$S \rightarrow A \mid S \# S \mid S @ S$$
$$A \rightarrow C \mid C A$$
$$C \rightarrow a \mid b \mid c$$

For each production below, state whether it is:

In the language G decides

- if so, prove it with a derivation and parse tree
- if not, explain why

If it is ambiguous in G

- if so, prove it is with another valid parse tree in G

Productions:

I. aa ## bb

II. a @ b # c

III. ab

1. Rewrite grammar G so that:

- no string in it is ambiguous
- # has higher precedence than @
- # and @ are both left-associative
(i.e. a # b # c should mean (a # b) # c)

2. Given grammar H below:

(in Extended BNF, terminals in bold and underlined):

<Statement> → <Assignment>
<Statement> → <While>
<Assignment> → <Var> ≡ <Value> [, <Value>] ;
<While> → while(<Value>) { {<Statement>} }
<Value> → <Var> | <Number>
<Var> → a | b | c
<Number> → 0 | 1

For each production below:

- determine if it is in the language H
- explain why or why not

Productions:

- I. a = 0,b;
- II. a = b,c,1;
- III. while(a){b = 0; while (b) { }}
- IV. a=1; while(a) {a = 1; while (a=0;}

3. Write a grammar for BNF in BNF
(not EBNF, e.g. no curly braces)

Evaluation:

All on time submissions will be graded out of 100 points, 25 points per question.