

# Machine Language

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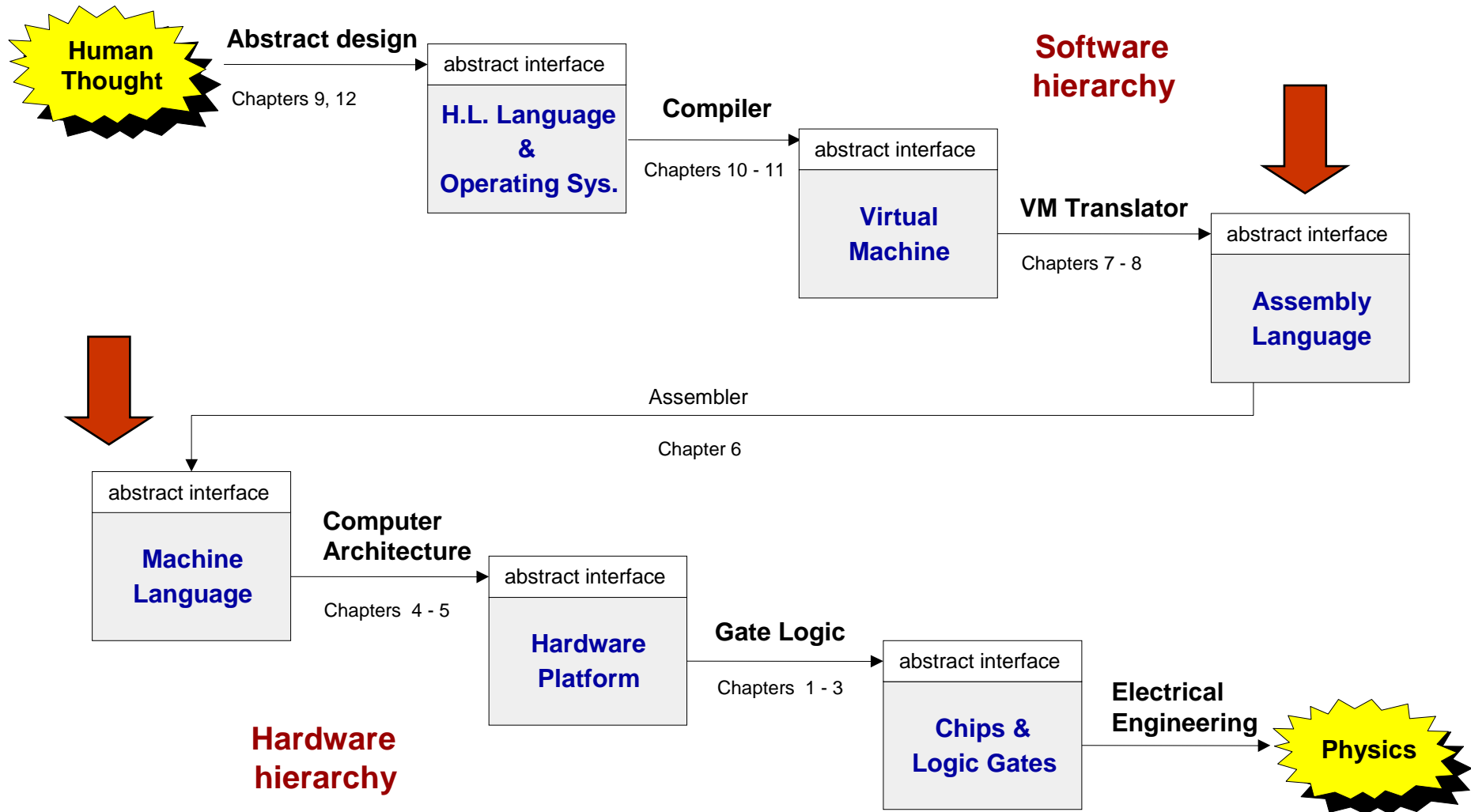
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# Where we are at:



# Machine language is “the soul of the machine”

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## Duality:

- Machine language (= instruction set) can be viewed as an abstract (programmer-oriented) description of the hardware platform
- The hardware can be viewed as a physical means for realizing an abstract machine language

## Another duality:

- Binary version
- Symbolic version

## Loose definition:

- Machine language = an agreed upon formalism for manipulating a *memory* using a *processor* and a set of *registers*
- Same spirit but different syntax across different hardware platforms.

# Binary and symbolic notation

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1010 0001 0010 1011

ADD R1, R2, R3



Jacquard loom  
(1801)

## Evolution:

- Physical coding
- Symbolic documentation
- Symbolic coding
- Translation and execution
- Requires a *translator*.



Augusta Ada King,  
Countess of Lovelace  
(1815-1852)

# Lecture plan

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- Machine languages at a glance
- The Hack machine language:
  - Symbolic version
  - Binary version
- Perspective

(The assembler will be covered in lecture 6).

# Instructions in a typical machine language

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```
// In what follows R1,R2,R3 are registers, PC is program counter,  
// and addr is a value.
```

```
ADD R1,R2,R3      //  $R1 \leftarrow R2 + R3$ 
```

```
ADDI R1,R2,addr   //  $R1 \leftarrow R2 + \text{addr}$ 
```

```
AND R1,R1,R2      //  $R1 \leftarrow \text{And}(R1,R2)$  (bit-wise)
```

```
JMP addr          //  $PC \leftarrow \text{addr}$ 
```

```
JEQ R1,R1,addr    // IF  $R1 = R2$  THEN  $PC \leftarrow \text{addr}$  ELSE  $PC++$ 
```

```
LOAD R1, addr     //  $R1 \leftarrow \text{RAM}[\text{addr}]$  Where v is an address
```

```
STORE R1, addr    //  $\text{RAM}[\text{addr}] \leftarrow R1$  where v is an address
```

```
NOOP              // Do nothings
```

```
// Plus several more commands that are essentially versions  
// or extensions of the above commands.
```

# The Hack computer

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The 16-bit Hack computer consists of the following elements:

Data memory: `RAM` - a series of 16-bit words

Instruction memory: `ROM` - a series of 16-bit words

Registers: `D`, `A`, `M`, where `M` stands for `RAM[A]`

Processing: `ALU`, capable of computing various functions

Program counter: `PC`, holding an address

Control: The `ROM` is loaded with a sequence of 16-bit instructions, one per memory location, beginning at address 0. The next instruction is always fetched from `ROM[PC]`

Instruction set: Two instructions: `A`-instruction, `C`-instruction.

# A-instruction

---

```
@value      // A ← value
```

Where *value* is either a number or a symbol referring to some number.

## Used for:

- Entering a constant value  
( A = value)

## Coding example:

```
@17      // A = 17  
D = A     // D = 17
```

- Selecting a RAM location  
( register = RAM[A] )

```
@17      // A = 17  
D = M     // D = RAM[17]
```

- Selecting a ROM location  
( fetch ROM[A] )

```
@17      // A = 17  
JMP      // fetch the instruction  
          // stored in ROM[17]
```

Later



## Coding examples (programming practice)

---

Write the Hack instructions that implement the following tasks:

- ❑ Set A to 17
- ❑ Set D to A-1
- ❑ Set both A and D to A + 1
- ❑ Set D to 19
- ❑ Set both A and D to A + D
- ❑ Set RAM[5034] to D - 1
- ❑ Set RAM[53] to 171
- ❑ Add 1 to RAM[7], and store the result in D.

Hack commands:

@value // set A to value

dest = x op y

op is + or -

x is A, D, or M

y is A, D, M or 1

(op y) is optional

dest is D, M, MD, A, AM, AD, AMD, or null

## Coding examples (cont.)

Write the Hack instructions that implement the following tasks:

- `sum = 0`
- `j = j + 1`
- `q = sum + 12 - j`
- `arr[7] = 0`

Etc.

Hack commands:

`@value // set A to value`

`dest = x op y`

`op` is `+` or `-`

`x` is `A`, `D`, or `M`

`y` is `A`, `D`, `M` or `1`

`(op y)` is optional

`dest` is `D`, `M`, `MD`, `A`, `AM`, `AD`, `AMD`, or null

Symbol table:

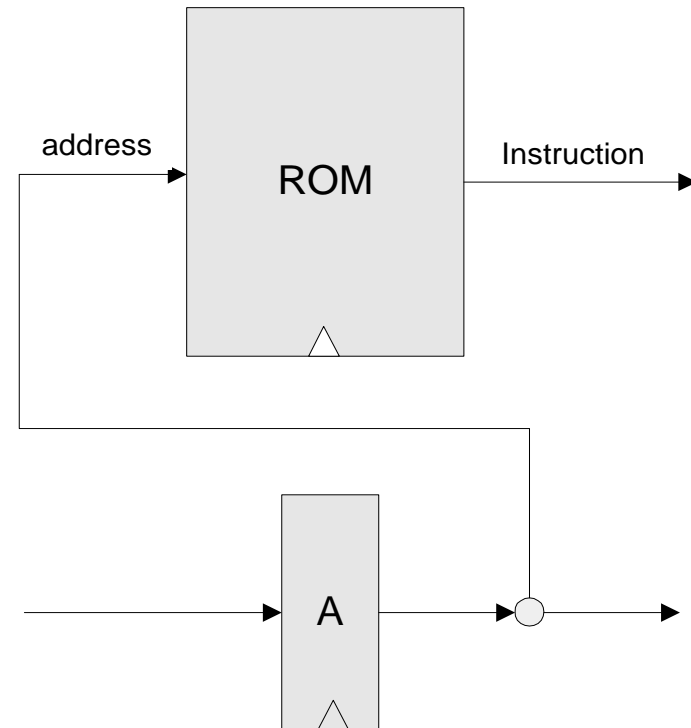
<code>j</code>	17
<code>sum</code>	22
<code>q</code>	21
<code>arr</code>	16

(All symbols and values in are arbitrary examples)

# Control (first approximation)

---

- ROM = instruction memory
- Program = sequence of 16-bit numbers, starting at ROM[0]
- Current instruction = ROM[A]
- To select instruction  $n$  from the ROM, we set A to  $n$ , using the instruction @ $n$



(The actual architecture is slightly different, as we'll see in the next lecture)

## Coding examples (practice)

Write the Hack instructions that implement the following tasks:

- ❑ GOTO 50
- ❑ IF D = 0 GOTO 112
- ❑ IF D < 9 GOTO 507
- ❑ IF RAM[12] > 0 GOTO 50
- ❑ IF sum > 0 GOTO END
- ❑ IF axis] <= 0 GOTO NEXT.

### Hack commands:

@value // set A to value

dest = comp ; jump // “dest = “ is optional

// Where:

comp = 0 , 1 , -1 , D , A , !D , !A , -D , -A , D+1 ,  
A+1 , D-1 , A-1 , D+A , D-A , A-D , D&A ,  
D | A , M , !M , -M , M+1 , M-1 , D+M , D-M ,  
M-D , D&M , D | M

dest = M , D , MD , A , AM , AD , AMD , or null

jump = JGT , JEQ , JGE , JLT , JNE , JLE , JMP , or null

All conditional jumps refer to the current value of D.

### Symbol table:

sum	200
x	4000
i	151
END	50
NEXT	120

(All symbols and values in are arbitrary examples)

# C-instruction syntax (final version)

---

```
dest = comp ; jump           // comp is mandatory
                               // dest and jump are optional
```

Where:

**comp** is one of:

```
0, 1, -1, D, A, !D, !A, -D, -A, D+1, A+1, D-1, A-1, D+A, D-A, A-D, D&A, D | A,
M,      !M,      -M,      M+1,      M-1, D+M, D-M, M-D, D&M, D | M
```

**dest** is one of:

```
null, M, D, MD, A, AM, AD, AMD
```

**jump** is one of:

```
null, JGT, JEQ, JGE, JLT, JNE, JLE, JMP
```

# IF logic – Hack style

---

High level:

```
if condition {  
    code segment 1}  
else {  
    code segment 2}  
// next instruction
```

Hack:

```
D ← not condition)  
@IF_TRUE  
D;JEQ  
code segment 2  
@END  
0;JMP  
(IF_TRUE)  
code segment 1  
(END)  
// next instruction
```

- To prevent conflicting use of the *A* register, in well-written Hack programs a *C*-instruction that includes a jump directive should not contain a reference to *M*, and vice versa.

# WHILE logic – Hack style

---

High level:

```
while condition {  
    code segment 1  
}  
// next instruction
```

Hack:

```
(LOOP)  
    D ← not condition)  
    @END  
    D;JEQ  
    code segment 1  
    @LOOP  
    0;JMP  
  
(END)  
  
// next instruction
```

# Complete program example

C:

```
// Adds 1+...+100.  
into i = 1;  
into sum = 0;  
while (i <= 100){  
    sum += i;  
    i++;  
}
```

Demo  
CPU emulator

Hack:

```
// Adds 1+...+100.  
@i      // i refers to some memo. location  
M=1     // i=1  
@sum    // sum refers to some memo. location  
M=0     // sum=0  
(LOOP)  
@i  
D=M     // D = i  
@100  
D=D-A   // D = i - 100  
@END  
D;JGT   // If (i-100) > 0 got END  
@i  
D=M     // D = i  
@sum  
M=D+M   // sum += i  
@i  
M=M+1   // i++  
@LOOP  
0;JMP   // Got LOOP  
(END)  
@END  
0;JMP   // Infinite loop
```



# Lecture plan

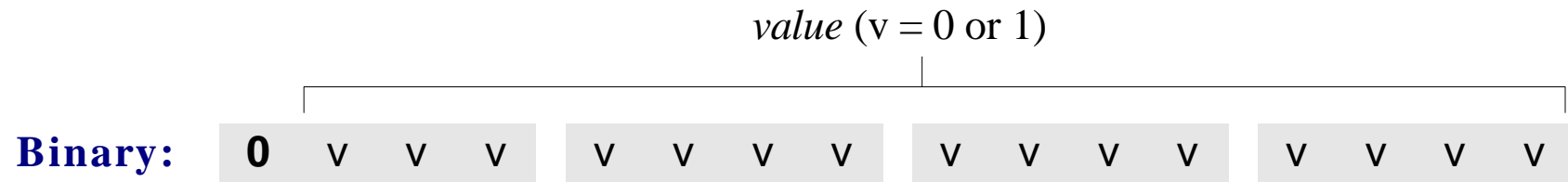
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- Symbolic machine language
- Binary machine language

# A-instruction

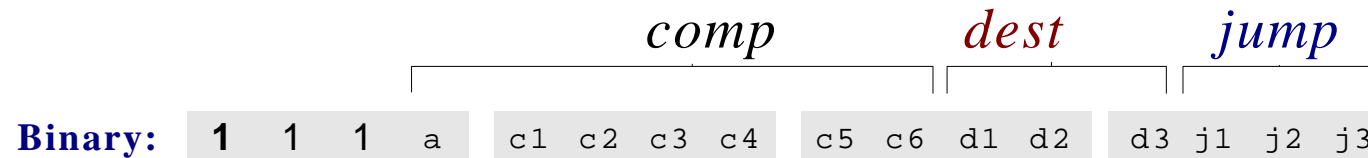
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**Symbolic:**    *@value*        // Where *value* is either a non-negative decimal number  
                                 // or a symbol referring to such number.



# C-instruction

**Symbolic:** *dest=comp;jump* // Either the *dest* or *jump* fields may be empty.



(when a=0) <i>comp</i>	c1	c2	c3	c4	c5	c6	(when a=1) <i>comp</i>	d1	d2	d3	Mnemonic	Destination (where to store the computed value)
0	1	0	1	0	1	0		0	0	0	null	The value is not stored anywhere
1	1	1	1	1	1	1		0	0	1	M	Memory[A] (memory register addressed by A)
-1	1	1	1	0	1	0		0	1	0	D	D register
D	0	0	1	1	0	0		0	1	1	MD	Memory[A] and D register
A	1	1	0	0	0	0	M	1	0	0	A	A register
!D	0	0	1	1	0	1		1	0	1	AM	A register and Memory[A]
!A	1	1	0	0	0	1	!M	1	1	0	AD	A register and D register
-D	0	0	1	1	1	1		1	1	1	AMD	A register, Memory[A], and D register
-A	1	1	0	0	1	1	-M					
D+1	0	1	1	1	1	1			j1	j2	j3	
A+1	1	1	0	1	1	1	M+1	(out < 0)	(out = 0)	(out > 0)	Mnemonic	Effect
D-1	0	0	1	1	1	0		0	0	0	null	No jump
A-1	1	1	0	0	1	0	M-1	0	0	1	JGT	If out > 0 jump
D+A	0	0	0	0	1	0	D+M	0	1	0	JEQ	If out = 0 jump
D-A	0	1	0	0	1	1	D-M	0	1	1	JGE	If out ≥ 0 jump
A-D	0	0	0	1	1	1	M-D	1	0	0	JLT	If out < 0 jump
D&A	0	0	0	0	0	0	D&M	1	0	1	JNE	If out ≠ 0 jump
D A	0	1	0	1	0	1	D M	1	1	0	JLE	If out ≤ 0 jump
								1	1	1	JMP	Jump

## Symbols (user-defined)

- Label symbols: User-defined symbols, used to label destinations of got commands. Declared by the pseudo command `(xxx)`. This directive defines the symbol `xxx` to refer to the instruction memory location holding the next command in the program
- Variable symbols: Any user-defined symbol `xxx` appearing in an assembly program that is not defined elsewhere using the "`(xxx)`" directive is treated as a variable, and is assigned a unique memory address by the assembler, starting at RAM address 16
- By convention, label symbols are upper-case and variable symbols are lower-case.

```
// Recto program
@R0
D=M
@INFINITE_LOOP
D;JLE
@counter
M=D
@SCREEN
D=A
@addr
M=D
(L00P)
@addr
A=M
M=-1
@addr
D=M
@32
D=D+A
@addr
M=D
@counter
MD=M-1
@LOOP
D;JGT
(INFINITE_LOOP)
@INFINITE_LOOP
0;JMP
```

## Symbols (pre-defined)

- Virtual registers: **R0**, ..., **R15** are predefined to be 0, ..., 15
- I/O pointers: The symbols **SCREEN** and **KBD** are predefined to be 16384 and 24576, respectively (base addresses of the *screen* and *keyboard* memory maps)
- Predefined pointers: the symbols **SP**, **LCL**, **ARG**, **THIS**, and **THAT** are predefined to be 0 to 4, respectively.

```
// Recto program
@R0
D=M
@INFINITE_LOOP
D;JLE
@counter
M=D
@SCREEN
D=A
@addr
M=D
(LOOP)
@addr
A=M
M=-1
@addr
D=M
@32
D=D+A
@addr
M=D
@counter
MD=M-1
@LOOP
D;JGT
(INFINITE_LOOP)
@INFINITE_LOOP
0;JMP
```

# Perspective

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- Hack is a simple machine language
- User friendly syntax:  $D=D+A$  instead of `ADD D,D,A`
- Hack is a " $\frac{1}{2}$ -address machine"
- A Macro-language can be easily developed
- A Hack assembler is needed and will be discussed and developed later in the course.

## End-note: a macro machine language (can be implemented rather easily)

---

### Assignment:

1. `x = constant` (e.g. `x = 17`)
2. `x = y`
3. `x = 0` , `x = 1`, `x = -1`

### Arithmetic / logical:

4. `x = y op z`  
where `y`, `z` are variables or constants and  
`op` is some ALU operation like `+`, `-`, `and`, `or`, etc.

### Control:

5. `GOTO s`
6. `IF condo THEN GOTO s`  
where `condo` is an expression `(x op y) {=|<|>|...} {0|1}`  
e.g. `IF x+17>0 got loop`

### White space or comments:

7. White space: ignore
8. `//` comment to the end of the line: ignore.