

Boolean Arithmetic

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This presentation contains lecture materials that accompany the textbook “The Elements of Computing Systems” by Noam Nisan & Shimon Schocken, MIT Press, 2005.

We provide both PPT and PDF versions.

The book web site, www.idc.ac.il/tecs , features 13 such presentations, one for each book chapter. Each presentation is designed to support about 3 hours of classroom or self-study instruction.

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Counting systems

quantity	decimal	binary	3-bit register
	0	0	000
*	1	1	001
**	2	10	010
***	3	11	011
****	4	100	100
*****	5	101	101
*****	6	110	110
*****	7	111	111
*****	8	1000	overflow
*****	9	1001	overflow
*****	10	1010	overflow

Rationale

$$(9038)_{ten} = 9 \cdot 10^3 + 0 \cdot 10^2 + 3 \cdot 10^1 + 8 \cdot 10^0 = 9038$$

$$(10011)_{two} = 1 \cdot 2^4 + 0 \cdot 2^3 + 0 \cdot 2^2 + 1 \cdot 2^1 + 1 \cdot 2^0 = 19$$

$$(x_n x_{n-1} \dots x_0)_b = \sum_{i=0}^n x_i \cdot b^i$$

Binary addition

- Assuming a 4-bit system:

$$\begin{array}{r} 0 \ 0 \ 0 \ 1 \\ \hline 1 \ 0 \ 0 \ 1 \\ 0 \ 1 \ 0 \ 1 \\ \hline 0 \ 1 \ 1 \ 1 \ 0 \end{array}$$

no overflow

$$\begin{array}{r} 1 \ 1 \ 1 \ 1 \\ \hline 1 \ 0 \ 1 \ 1 \\ 0 \ 1 \ 1 \ 1 \\ \hline 1 \ 0 \ 0 \ 1 \ 0 \end{array}$$

overflow

- Algorithm: exactly the same as in decimal addition
- Overflow (MSB carry) has to be dealt with.

Representing negative numbers (4-bit system)

0	0000		
1	0001	1111	-1
2	0010	1110	-2
3	0011	1101	-3
4	0100	1100	-4
5	0101	1011	-5
6	0110	1010	-6
7	0111	1001	-7
		1000	-8

- The codes of all positive numbers begin with a "0"
- The codes of all negative numbers begin with a "1"
- To convert a number:
leave all trailing 0's and first 1 intact,
and flip all the remaining bits

Example: $2 - 5 = 2 + (-5) =$ 0 0 1 0

 + 1 0 1 1

 1 1 0 1 = -3

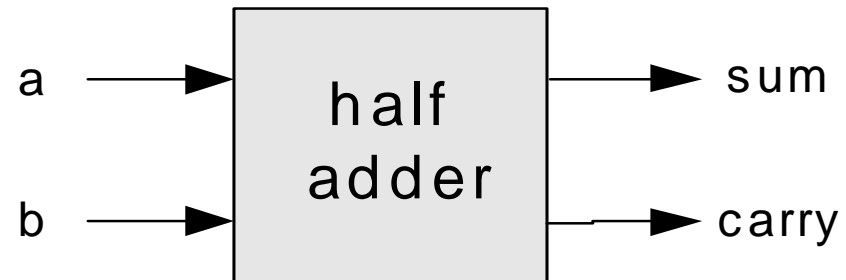
Building an Adder chip



- Adder: a chip designed to add two integers
- Proposed implementation:
 - Half adder: designed to add 2 bits
 - Full adder: designed to add 3 bits
 - Adder: designed to add two n -bit numbers.

Half adder (designed to add 2 bits)

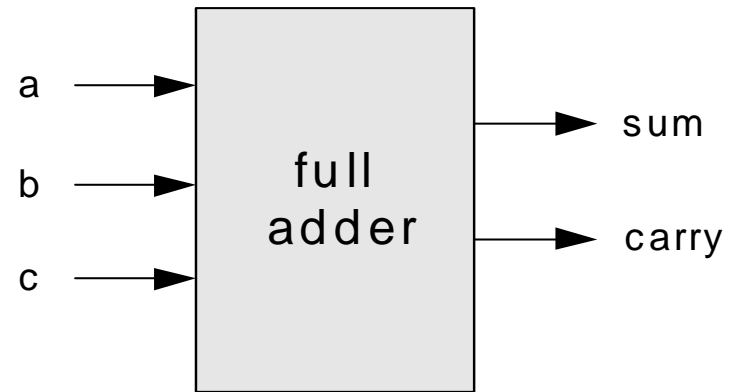
a	b	carry	sum
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0



- Implementation: based on two gates that you've seen before.

Full adder (designed to add 3 bits)

a	b	c	carry	sum
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1



- Implementation: can be based on half-adder gates.

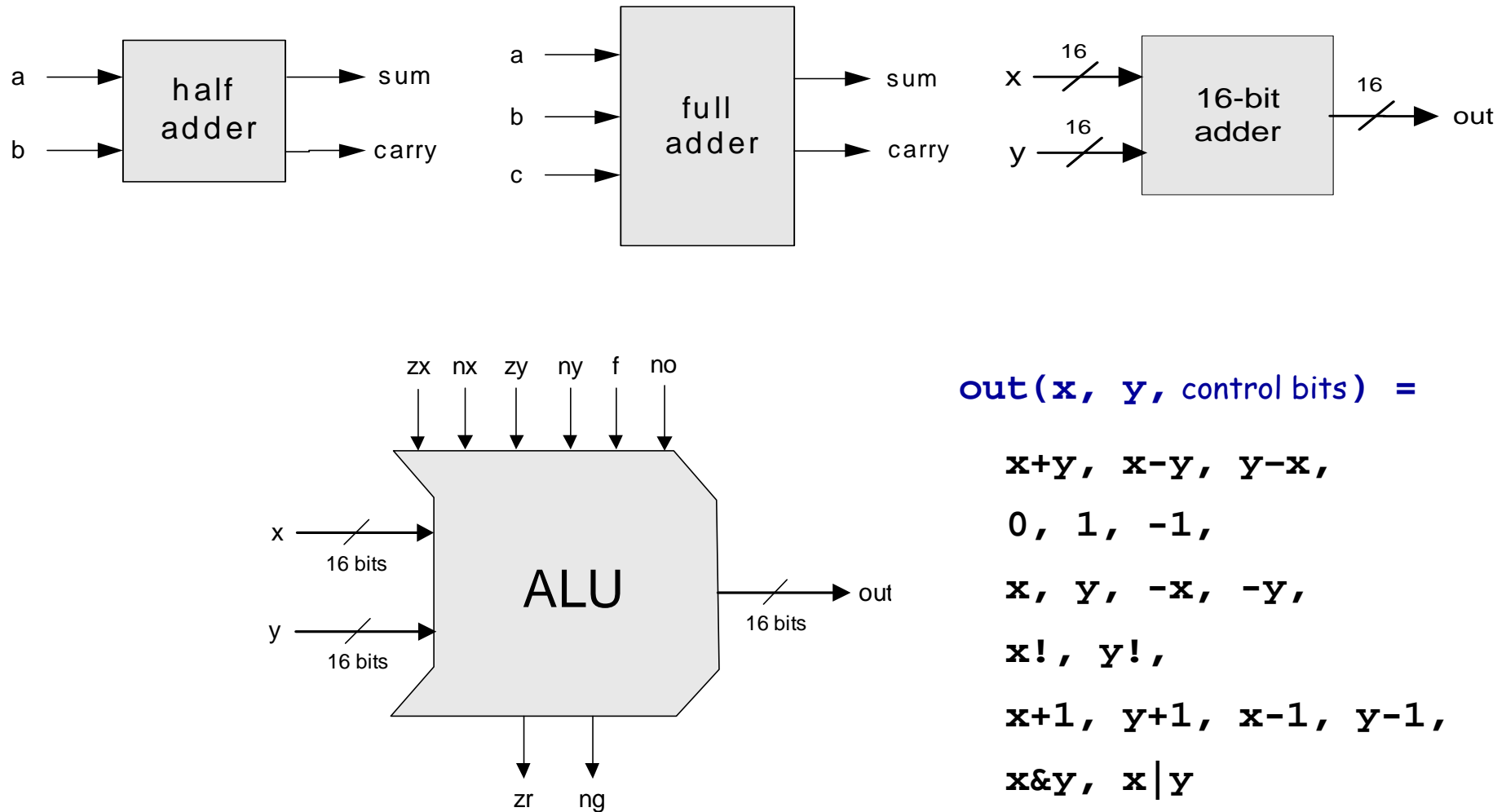
n -bit Adder



$$\begin{array}{rcccccl} \dots & 1 & 0 & 1 & 1 & a \\ & & & & & + \\ \dots & 0 & 0 & 1 & 0 & b \\ \hline \dots & 1 & 1 & 0 & 1 & \text{out} \end{array}$$

- Implementation: array of full-adder gates.

The ALU (of the Hack platform)



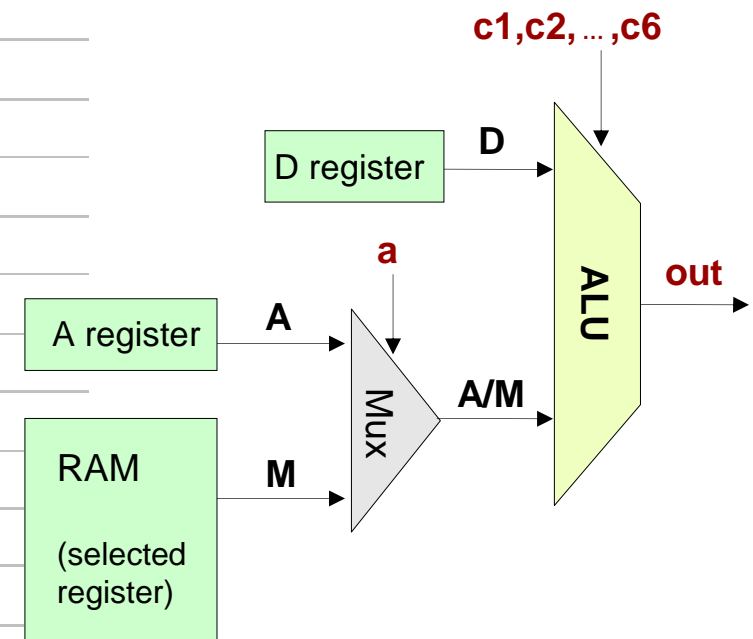
ALU logic (Hack platform)

These bits instruct how to pre-set the x input		These bits instruct how to pre-set the y input		This bit selects between + / And	This bit inst. how to post-set out	Resulting ALU output
zx	nx	zy	ny	f	no	out=
if zx then x=0	if nx then x=!x	if zy then y=0	if ny then y=!y	if f then out=x+y else out=x And y	if no then out=!out	f (X, y) =
1	0	1	0	1	0	0
1	1	1	1	1	1	1
1	1	1	0	1	0	-1
0	0	1	1	0	0	x
1						y
0						!x
1						!y
0						-x
1						-y
0	1	1	1	1	1	x+1
1	1	0	1	1	1	y+1
0	0	1	1	1	0	x-1
1	1	0	0	1	0	y-1
0	0	0	0	1	0	x+y
0	1	0	0	1	1	x-y
0	0	0	1	1	1	y-x
0	0	0	0	0	0	x&y
0	1	0	1	0	1	x y

Implementation: build a logic gate architecture that "reads" each control bit and does what the table specifies: if zx=1 then set x to 0, etc.

The ALU in the CPU context (Hack platform)

out (when a=0)	c1	c2	c3	c4	c5	c6	out (when a=1)
0	1	0	1	0	1	0	
1	1	1	1	1	1	1	
-1	1	1	1	0	1	0	
D	0	0	1	1	0	0	
A	1	1	0	0	0	0	M
!D	0	0	1	1	0	1	
!A	1	1	0	0	0	1	!M
-D	0	0	1	1	1	1	
-A	1	1	0	0	1	1	-M
D+1	0	1	1	1	1	1	
A+1	1	1	0	1	1	1	M+1
D-1	0	0	1	1	1	0	
A-1	1	1	0	0	1	0	M-1
D+A	0	0	0	0	1	0	D+M
D-A	0	1	0	0	1	1	D-M
A-D	0	0	0	1	1	1	M-D
D&A	0	0	0	0	0	0	D&M
D A	0	1	0	1	0	1	D M



Perspective

- Combinational logic
- Our adder design is very basic: no parallelism
- It pays to optimize adders
- Our ALU is also very basic: no multiplication / division
- Where is the seat of advanced math operations?
a typical hardware/software tradeoff.

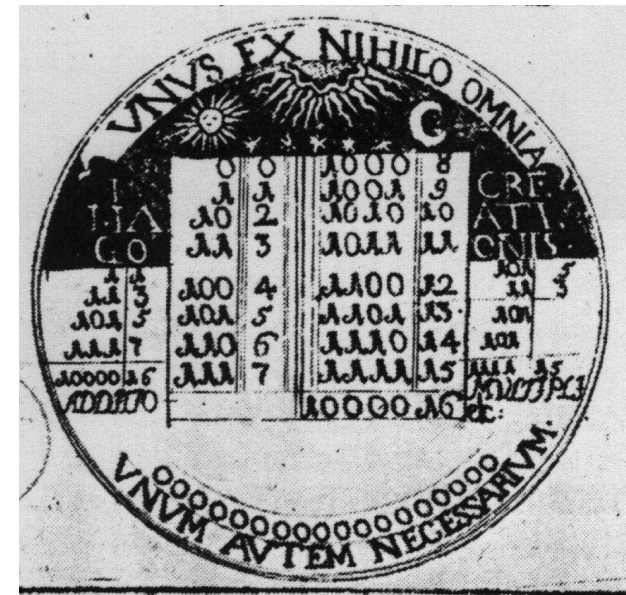
Historical note: Leibnitz (1646-1716)



- "The binary system may be used in place of the decimal system; express all numbers by unity and by nothing"
- 1679: built a mechanical calculator (+, -, *, /)



- CHALLENGE: "All who are occupied with the reading or writing of scientific literature have assuredly very often felt the want of a common scientific language, and regretted the great loss of time and trouble caused by the multiplicity of languages employed in scientific literature:
- SOLUTION: "*Characteristica Universalis*": a universal, formal, language of reasoning
- The dream's end: Turing and Goedel in 1930's.



Leibniz's medallion
for the Duke of Brunswick