

# LINFO1104 – LSINC1104

## Concepts, paradigms, and semantics of programming languages

### Lecture 1

#### Basic concepts, functions, and invariants

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## Context of this course

- LEPL1402 – LSINC1402: First semester programming course
  - Introduction to programming based on Java and IntelliJ
  - Java and object-oriented programming
  - Algorithms, data structures, and invariants
  - Complexity and software engineering
  - Introduction to functions and concurrency in Java
- **LINFO1104 – LSINC1104**: Second semester course
  - Programming paradigms based on Oz multiparadigm language
  - Formal semantics including lambda calculus
  - Procedural abstraction (higher-order programming)
  - Data abstraction (objects and abstract data types)
  - Symbolic programming
  - Concurrent programming
  - Multi-agent programming and Erlang



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# Programming paradigms



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*Hundreds of programming languages are in use...*



python™



Scala



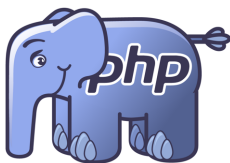
Java



Erlang

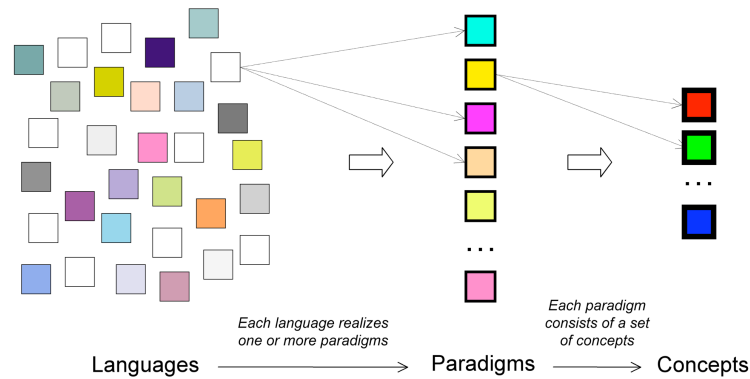


OCaml



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## So many, how can we understand them all?



- Key insight: languages are based on paradigms, and there are many fewer paradigms than languages
- We can **understand many languages by learning few paradigms!**

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## What is a paradigm?

- A **programming paradigm** is an approach to programming a computer based on a coherent set of principles or a mathematical theory
- A program is written to solve problems
  - Any realistic program needs to solve different kinds of problems
  - Each kind of problem needs its own paradigm
  - So we need **multiple** paradigms and we need to **combine** them in the same program

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## How can we study *multiple* paradigms?



- How can we study multiple paradigms without studying multiple languages (since most languages only support one, or sometimes two paradigms)?
- Each language has its own syntax, its own semantics, its own system, and its own quirks
  - Picking many languages, like Java, Scala, Erlang, Scheme, and Haskell, and structuring our course around them would be complicated
- Our pragmatic solution: **two languages**, Oz (a multiparadigm research language) and Erlang (a multi-agent industrial language)
  - Our textbook, *Concepts, Techniques, and Models of Computer Programming*, uses Oz to cover many paradigms



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## How can we *combine* paradigms in a program?



- Each paradigm is a **different way of thinking**
  - How can we combine different ways of thinking in one program?
- We can do it using the concept of a **kernel language**
  - Each paradigm has a simple core language, its kernel language, that contains its essential concepts
    - Every practical language, even complicated ones, can be translated easily into its kernel language
  - Even very different paradigms have kernel languages that have much in common; often there is only one concept difference
- We start with a simple kernel language that underlies our first paradigm, **functional programming**
  - We then **add concepts one by one** to give the other paradigms
  - Scientific method: understand a system in terms of its parts



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## Summary of the approach



- **Hundreds of languages** are used in practice: we cannot study them all in one course or in one lifetime
  - Solution: **focus on paradigms**, since each language is based on a paradigm and there are many fewer paradigms than languages
- **One language per paradigm is too much** to study in a course, since each language is already complicated by itself
  - Solution: **use one research language**, Oz, that can express many paradigms (plus Erlang for multi-agent programming!)
- **Realistic programs need to combine paradigms**, but how can we do it since each paradigm is a different way of thinking?
  - Solution: **define paradigms using kernel languages**, since different paradigms have kernel languages that are almost the same
  - Kernel languages allow us to define many paradigms by focusing on their differences, which is much more economical in time and effort

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## Let's get started



- You should already know an object-oriented language!
  - **Object-oriented programming**, as used in Java, is clearly an important paradigm (as seen in LEPL1402)
  - But what about the other paradigms?
- Isn't object-oriented programming by far the most important and useful paradigm?
  - Actually, no, it's not!
  - **Many other paradigms are extremely useful**, often more so than OOP! For example, to make robust and efficient distributed programs on the Internet, OOP does not solve the right problems. **Multi-agent programming** is much better for that, which is why we will give an introduction to Erlang.
  - LINFO1104 covers **five paradigms** that solve many problems

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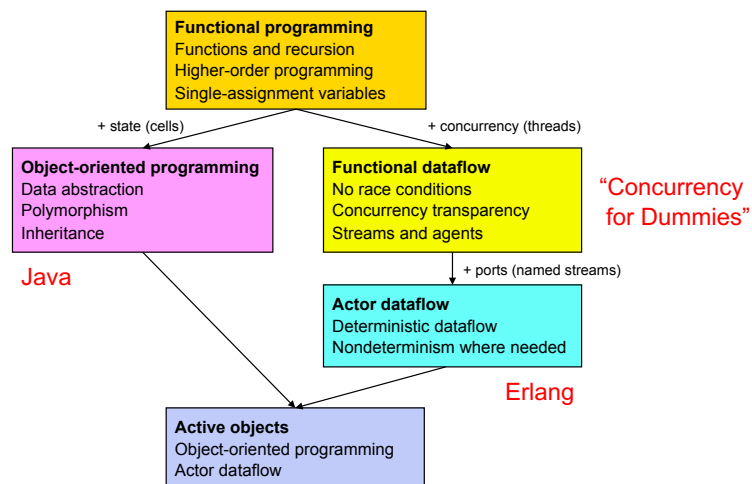
## Five paradigms



- LINFO1104/LSINC1104 covers five paradigms:
  - Functional programming
  - Object-oriented programming (used in Java)
  - Functional dataflow programming
  - Actor dataflow programming (multi-agent) (used in Erlang)
  - Active objects
- These are probably the most important programming paradigms for general use
  - But there are many other paradigms, made for other problems: LINFO1104/LSINC1104 gives you a good foundation for studying them later if you wish

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## Five paradigms



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## Our first paradigm



- Functional programming
  - It is the simplest paradigm
  - It is the foundation of all other paradigms
  - It is a form of *declarative programming*: say what, not how
- Functional programming is the right way to start
  - It is our first introduction to **programming concepts**
  - It is our first introduction to a **kernel language**
  - We use it to explain **invariant programming** (how to do loops)
  - We use it to explain **symbolic programming**
  - We use it to explain **higher-order programming**
  - We give a **formal semantics** based on the kernel language
- Everything we learn in functional programming stays true later!

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## Practical organization



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## Course organization LINFO1104



- Weekly lecture
  - Tuesday 14h00-16h00 (A.02 SCES)
- Weekly lab session
  - Tuesday 16h15-18h15 (MERC 04)
  - Wednesday 8h30-10h30 (MERC 02)
  - Wednesday 10h45-12h45 (MERC 02, sauf en S1 MERC 01)
  - Wednesday 14h00-16h00 (PCUR 01, sauf en S1 BARB 13)
  - Monday 14h00-16h00 (MERC 04 à partir de S2)
- Midterm exam (around week 7)
  - Optional, counts for 5 points on final grade (max. with final exam Q1)
- Course project (second half of semester)
  - Practical programming project in groups of two students
  - Mandatory, counts for 5 points on final grade
- Final exam (15 points; 5 points correspond to midterm)
- Bonus (1 point extra if you attend all lab sessions)

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## Software support



- LINFO1104 Moodle
  - All practical information is given here
  - Register yourself!
- Mozart Programming System
  - For practical exercises; [www.mozart2.org](http://www.mozart2.org)
  - Download and install Mozart 2.0.1
- Erlang OTP
  - For multi-agent programming; [www.erlang.org](http://www.erlang.org)
  - Download and install OTP 25.2 (in second half of course)

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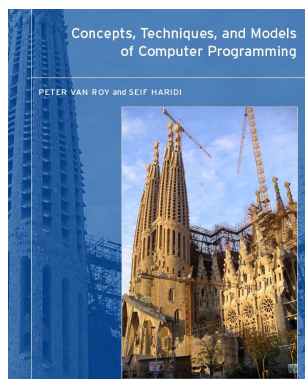
## Educational team



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## Course textbook and slides



- “Concepts, Techniques, and Models of Computer Programming” by Peter Van Roy and Seif Haridi, MIT Press
- Slides for each week’s lecture
  - The slides contain most of the course material, but please look at the book if anything is unclear or if you want more examples
  - Some material will be given that is not in the book, in particular regarding semantics (lambda calculus), for that part you should refer to the slides
  - It is up to you to read the book if anything is unclear! (or ask your tutor, a teaching assistant, or me)

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# Basic concepts



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## Mozart Programming System: interactive Emacs interface



**declare**

$X = 1234 * 5678$

{Browse X}

- Select a region in the Emacs buffer
- Feed the region to the system
  - The text is compiled and executed
- Interactive system can be used as a powerful calculator →



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## Creating variables

### declare

X = 1234 \* 5678

{Browse X}

- **Declare** (create) a variable **designated** by X
- **Assign** to the variable the value 7006652
  - Result of the calculation 1234\*5678
- **Call** the procedure Browse with the argument designated by X
  - Opens a window that displays 7006652

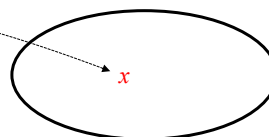
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## Identifiers and variables

*Program text*

```
declare X  
X=11*11  
{Browse X}
```

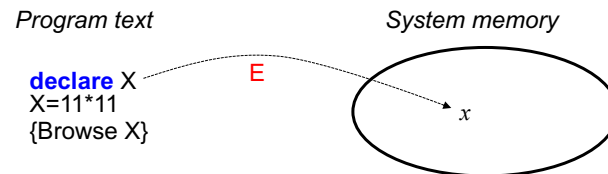
*System memory*



- There are two concepts hiding in plain view here
  - **Identifier X** : what you type (character sequence starting with uppercase)  
Var, A, X123, FirstCapitalBank
  - **Variable x** : what is in memory (used to store the value)
- Variables are short-cuts for values (= constants)
  - Can only be assigned to one value (like [mathematical variables](#))
  - Multiple assignment is another concept! We will see it later in the course.
  - The type of the variable is only known when it is assigned ([dynamic typing](#))

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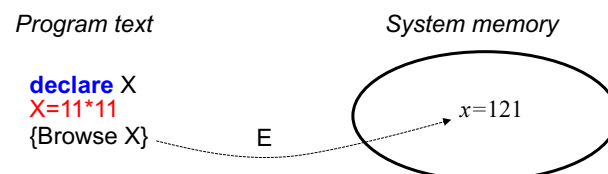
# Environment



- **declare** is an interactive instruction
  - Creates a new variable in memory
  - Links the identifier and its corresponding variable
- Third concept: **environment**  $E = \{X \rightarrow x\}$ 
  - A function that takes an identifier and returns a variable:  $E(X) = x$
  - Links identifiers and their corresponding variables (with the values they are bound to)

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# Assignment



- The assignment instruction  $X=121$  binds the variable  $x$  to the value 121

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## Single assignment



- A variable can only be bound to **one value**
  - It is called a **single-assignment variable**
  - Why? Because we are doing functional programming!
- Incompatible assignment: **signals an error**  
X = 122
- Compatible assignment: **accepted**  
X = 121

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## Why single assignment?



- Single assignment is part of functional programming
  - It means that variables are *mathematical variables*, like in an equation
- Programming with mathematical variables is easy
  - It may seem strange for Java programmers, but it is actually quite easy
- Advantages
  - Programs are much **easier to prove, analyze, test, and debug**
- Disadvantages
  - It's a different way of thinking and takes some getting used to
  - We will see later that the main disadvantage is that **interaction with the real world is harder** (e.g., with human users that connect to the system and interact). This is a bit subtle so we do not explain it now.

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## The functional style can be done in all languages

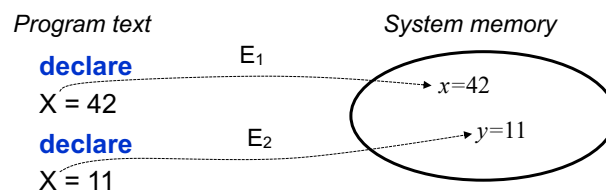


- “A program that works today will work tomorrow”
  - Functions and variables don’t change
  - All changes are in the arguments, not in the functions
- This programming style is encouraged in all languages (incl. Java)
  - “Stateless server” for a client/server application
  - “Stateless component” for a service application
- Learning functional programming helps us think in this style
  - As well as helping us understand all programming languages, today’s and tomorrow’s



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## Redeclaring an identifier



- An identifier can be **redeclared**
  - The same identifier refers to a different value
  - There is no conflict with single assignment. Each occurrence of X corresponds to a different variable.
- The interactive environment always has the last declaration
  - **declare** keeps the same correspondance until redeclared (if ever)
  - In this example X will refer to 11

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## Scope of an identifier occurrence



```
local
  X
in
  X = 42 {Browse X}
  local
    X
  in
    X = 11 {Browse X}
  end
  {Browse X}
end
```

- The instruction `local X in <stmt> end` declares X between `in` and `end`
- The **scope** of an identifier occurrence is that part of the program text for which the occurrence corresponds to the same variable declaration
- The scope can be determined by **inspecting the program text**; no execution is needed. This is called **lexical scoping** or **static scoping**.
- Why is there no conflict between X=42 and X=11, even though variables are single assignment?
- What will the third Browse display?

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## Tips on Oz syntax



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## Tips on Oz syntax



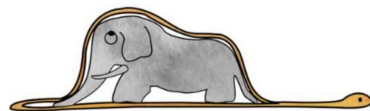
- You can see that Oz syntax is not like Java syntax
- The most popular syntax in mainstream languages (Java and C++) is « C-like », where identifiers are statically typed « int i; » and start with lowercase, and code blocks are delimited by braces { ... }
  - Oz syntax is definitely not C-like!
- Oz syntax is designed for multiparadigm programming
  - Oz syntax is inspired by many languages: Prolog (logic programming), Scheme and ML (functional programming), C++ and Smalltalk (object-oriented programming), and so forth
- (Later on we will see Erlang, with even another syntax)

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## Syntax is just the surface



"My drawing was not a picture of a hat.  
It was a picture of a boa constrictor digesting an elephant."



*"It is only with the heart that one can see  
rightly; what is essential is invisible to the eye."*

*"On ne voit bien qu'avec le cœur. L'essentiel  
est invisible pour les yeux."*

– *The Little Prince, Antoine de Saint-Exupéry*

- Programmers are too much in love with their favorite syntax
  - Syntax is just the surface. Programmers must learn to look beyond syntax to the true reality, which is semantics.
  - In this course we will mainly use Oz syntax, but the semantics you learn applies to all the languages you will use in your whole career
  - Programmers should be polyglot and move easily from one syntax to another
    - I hope that this course helps you achieve that ability!

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## Why is Oz syntax different?



- It is different because Oz supports many programming paradigms
  - The syntax is carefully designed so that the paradigms don't interfere with each other
  - It's possible to program in just one paradigm. It's also possible to program in several paradigms that are cleanly separated in the program text.
- So it is important not to get confused by the differences between Oz syntax and other syntaxes you may know
- We show the main differences so that you will not be hindered by them

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## Four main differences in Oz syntax



- Identifiers in Oz always start with an **uppercase letter**
  - Examples: X, Y, Z, Min, Max, Sum, IntToFloat.
  - Why? Because lowercase is used for symbolic constants (atoms).
- Procedure and function calls in Oz are surrounded by **braces** { ... }
  - Examples: {Max 1 2}, {SumDigits 999}, {Fold L F U}.
  - Why? Because parentheses are used for record data structures.
- **Local identifiers** are introduced by **local ... end**
  - Example: **local** X **in** X=10+20 {Browse X} **end**.
  - Why? Because all compound instructions in Oz start with a keyword (such as « **local** ») and terminate with **end**.
- Variables in Oz are **single assignment**
  - Examples: **local** X Y **in** X=10 Y=X+20 {Browse Y} **end**.
  - Why? Because the first paradigm is functional programming. Multiple assignment is a concept that we will introduce later.

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## Oz syntax in the programming exercises



- Most programming bugs, at least early on, are due to syntax errors
  - Most common error: lowercase letter to start an identifier
- Please take into account the four main differences. Once you have assimilated them, reading and writing Oz will become straightforward.
- And now let's introduce functions

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## Functional programming



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## Functions



- We would like to execute the same code many times, each time with different values for some of the identifiers
  - To avoid repeating the same program code, we can **create a function**
- A function defines program code to execute
  - A function is just another kind of value in memory, like a number (as we will see later)
  - Variables can be bound to functions just as easily as to numbers
- The function Sqr returns the square of its input:

```
declare  
fun {Sqr X} X*X end
```

- The **fun** keyword identifies the function. The identifier Sqr refers to a variable that is bound to the function.

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## Numbers



- There are two kinds of numbers in Oz
  - **Exact numbers**: integers
  - **Approximate numbers**: floating point
- Integers are exact (arbitrary precision)
- Floats are approximations to real numbers (around 15 digits precision, using 64-bit representation)
- In Oz there is **never any automatic conversion** from exact to approximate and vice versa
  - To convert, we use functions IntToFloat or FloatToInt
  - Design principle: **don't mix incompatible concepts**

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## SumDigits3



- Function SumDigits3 calculates the sum of digits of a three-digit positive integer:

```
declare
fun {SumDigits3 N}
  (N mod 10) + ((N div 10) mod 10) +
  ((N div 100) mod 10)
end
```

- **mod** and **div** are integer functions
- / (division) is a float function
- + (addition) is a function on both floats and integers

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## SumDigits6



- Sum of digits of a six-digit positive integer

```
fun {SumDigits6 N}
  {SumDigits3 (N div 1000)} +
  {SumDigits3 (N mod 1000)}
end
```

- This is an example of **function composition**: defining a function in terms of other functions
  - This is a **key ability for building large systems**: we can build them in **layers**, where each layer is built by a different person
  - This is the first step towards **data abstraction**

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## SumDigitsR (first try)



- Sum of digits of any positive integer
- We use **recursion**: the function calls itself with a smaller argument

```
fun {SumDigitsR N}  
  (N mod 10) + {SumDigitsR (N div 10)}  
end
```

- This function calls itself with a smaller value
  - What is the problem with this definition?

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## SumDigitsR (correct)



- Sum of digits of any positive integer

```
fun {SumDigitsR N}  
  if (N==0) then 0  
  else  
    (N mod 10) + {SumDigitsR (N div 10)}  
  end  
end
```

- This introduces the conditional (**if**) statement
- This is a correct example of **function recursion**: defining a function that calls itself
  - This is a **key ability for building complex algorithms**: we divide a complex problem into simpler subproblems (divide and conquer)

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## We can now do functional programming



- We're done! We can now do functional programming.
  - Calculating with numbers, functions, and conditionals can do everything!
- This language is **Turing complete**
  - We can compute anything that any computer can compute
- We can extend it with **powerful concepts and techniques**
  - Invariant programming, symbolic programming, higher-order programming, data abstraction (including objects), concurrent programming
  - How is this better, since the original language is already Turing complete?
- Is there an even smaller language for functional programming?
  - Yes, the **lambda calculus ( $\lambda$  calculus)** is the smallest!
  - We will see the lambda calculus later in the course
  - **All programming languages are based on the lambda calculus**



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## Invariant programming



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## Invariant programming: loops!



- A **loop** is a part of a program that is repeated until a condition is satisfied
  - Loops are an important technique in all paradigms
  - In functional programming, loops are done with recursion
- We give a general technique, **invariant programming**, to program correct and efficient loops
  - Loops are often hard to get exactly right, and invariant programming is an excellent way to make them right
  - Invariant programming works for **all paradigms**, including functional programming and Java's object-oriented programming
- Each loop has one invariant
  - An invariant is a formula that is true at the beginning of each loop

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## Doing loops with recursion...



- Let's take another look at SumDigitsR:

```
fun {SumDigitsR N}
  if (N==0) then 0
  else (N mod 10) + {SumDigitsR (N div 10)} end
end
```
- The recursive call and the condition together act like a **loop**: a calculation that is repeated to achieve a result
  - Each execution of the function body is one iteration of the loop
- We see that recursion can make a loop
  - Let's show a simpler example to understand better how it works!

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## Naïve factorial function



- We can define factorial mathematically:
  - $0! = 1$
  - $n! = n \times (n-1)!$  when  $n > 0$
- We can define it as a program:

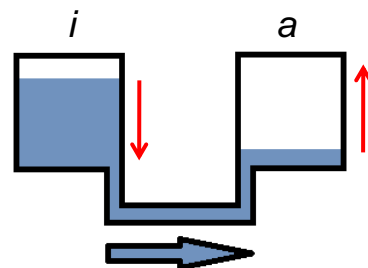
```
fun {Fact1 N}  
  if N==0 then 1  
  else N*{Fact1 N-1} end  
end
```
- This looks ok, right? It's actually very bad!

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## Better factorial function



- There is a much better way to define factorial!
- We start with an **invariant**, which is a logical formula that splits the work into two parts
  - We want to calculate  $n!$
  - Split it into two parts:  $n! = i! \times a$
- We do **communicating vases**
  - $n$  is constant
  - $i$  and  $a$  are varying
  - Start with  $i=n$  and  $a=1$
  - Decrease  $i$  and increase  $a$ , while keeping the formula true
  - When  $i=0$  then  $a$  is the answer



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## Better factorial function



- Using the invariant gives another factorial function
  - The varying parts ( $i$  and  $a$ ) are arguments of Fact2
  - The constant part ( $n$ ) is the first call to Fact2

```
fun {Fact2 I A}
  if I==0 then A
  else {Fact2 I-1 I*A} end
end
{Browse {Fact2 N 1}}
```

- This one is much better than Fact1
  - Let's see why by looking at the execution of both

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## Comparing Fact1 and Fact2



- $10 * \{ \text{Fact1 } 10-1 \} \Rightarrow$   
 $10 * (9 * \{ \text{Fact } 9-1 \}) \Rightarrow$   
 $10 * (9 * (8 * \{ \text{Fact } 8-1 \})) \Rightarrow$

Each line does one computation step.  
The expression gets bigger and bigger!

...  
 $10 * (9 * (8 * (7 * (6 * (5 * (... (1 * \{ \text{Fact } 0 \}) ...)) \Rightarrow$   
 $10 * (9 * (8 * (7 * (6 * (5 * (... (1 * 1) ...)) \Rightarrow$   
 ...  
 3628800

- $\{ \text{Fact2 } 10-1 \ 10*1 \} \Rightarrow$   
 $\{ \text{Fact2 } 9-1 \ 9*10 \} \Rightarrow$   
 $\{ \text{Fact2 } 8-1 \ 8*90 \} \Rightarrow$   
 ...  
 $\{ \text{Fact2 } 1-1 \ 1*3628800 \}$

?

It seems that Fact2 is better than Fact1, because the expression does not grow. How can we make this intuition precise? We need to introduce a **formal semantics** of the execution.  
 → we will see it in lectures 2 & 3

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## Fact2 is tail-recursive



- **Tail recursion** is when the recursive call is the last operation in the function body
- $N * \{Fact1\ N-1\}$  % No tail recursion  
 ↑  
 After Fact1 is done, **we must come back** for the multiply.  
 Where is the multiplication stored? On a stack!
- $\{Fact2\ I-1\ I*A\}$  % Tail recursion  
 The recursive call **does not come back**!  
 All calculations are done *before* Fact2 is called.  
 No stack is needed (memory usage is constant).

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## Sum of digits using invariant programming



- Each recursive call handles one digit
- So we divide the initial number  $n$  into its digits:
  - $n = (d_{k-1}d_{k-2}\cdots d_2d_1d_0)$  (where  $d_i$  is a digit)
- Let's call the sum of digits function  $s(n)$
- Then we can split the work in two parts:
  - $s(n) = \underbrace{s(d_{k-1}d_{k-2}\cdots d_i)}_{s_i} + \underbrace{(d_{i-1} + d_{i-2} + \cdots + d_0)}_a$
- $s_i$  is the work not yet done and  $a$  is the work already done
- To keep the formula true, we set  $i' = i+1$  and  $a' = a + d_i$
- When  $i=k$  then  $s_k=s(0)=0$  and therefore  $a$  is the answer

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## Example execution



- Example with  $n=314159$ :

$$s(n) = s(d_{k-1}d_{k-2}\cdots d_i) + (d_{i-1} + d_{i-2} + \cdots + d_0)$$

- $s(314159) = s(314159) + 0$
- $s(314159) = s(31415) + 9$
- $s(314159) = s(3141) + 14$
- $s(314159) = s(314) + 15$
- $s(314159) = s(31) + 19$
- $s(314159) = s(3) + 20$
- $s(314159) = s(0) + 23 = 0 + 23 = 23$

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## Final SumDigits2 program



- $S = (d_{k-1}d_{k-2}\cdots d_i)$   
 $A = (d_{i-1} + d_{i-2} + \cdots + d_0)$

```
fun {SumDigits2 S A}
  if S==0 then A
  else
    {SumDigits2 (S div 10) A+(S mod 10)}
  end
end
{Browse {SumDigits2 314159 0}}
```

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## Invariant programming is best



- We have now programmed two problems
  - Factorial
  - Sum of digits
- For each problem we have defined two functions
  - First version based on a simple mathematical definition
  - Second version designed with **invariant programming**
- The second version has three interesting properties
  - It uses **constant memory space**, unlike the first version
  - It has two arguments: the growing one is called an **accumulator**
  - The recursive call is the last operation: it is **tail-recursive**



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## Tail recursion = while loop



- A while loop in functional programming:

```
fun {While S}
  if {IsDone S} then S
  else {While {Transform S}} end /* tail recursion */
end
```
- A while loop in imperative programming:  
(i.e., in languages with multiple assignment like Java and C++)

```
state whileLoop(state s) {
  while (IsDone(s))
    s=transform(s); /* assignment */
  return s;
}
```
- In *both* cases, **invariant programming is the right way to do loops**

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# Summary



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## Summary



- Overview of course
- Practical organization
- Basic concepts
- Tips on Oz syntax
- Functional programming
- Invariant programming
- Next week:  
symbolic programming

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