Drools Usage Manual

version 2.0-beta-12

The Drools Project

drools.org

Contents

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Chapter 1

Functional Overview

1.1 Rules

1.1.1 Declarative Form

Drools directly supports *declarative* rules, as opposed to *procedural* logic. Declarative rules typically take the form as follows:

> **if** *condition*
> **then** *consequence*

By "declarative', it is meant that the rules *declare*, by way of the *condition*, what should occur but do not specify the procedure for actually testing the conditions. For example, a procedural method for ensuring you have an umbrella if it is raining would be:

> Step outside and determine if it is raining. If it is raining, then go to the closet and get an umbrella.

In a declarative form, the above could be represented by two rules:

- **if** *it is raining*
  then *you need an umbrella*

- **if** *you need an umbrella*
  then *get one from the closet*

Given declarative rules, the knowledge that it is raining could produce two courses of action:

1. You already have an umbrella, perhaps because you always carry one, in which case, you're ready.

2. You don't have an umbrella, so you go get one from the closet.

1.1.2   Applicable Context

The set of rules to be considered at any point of time depend upon the current context. As a human, you have certain rules to think about when you dine at a fine restaurant, which are exclusive to the set of rules you consider when spending a sunny day at the swimming pool. Even so, there are some other overriding rules that may be important regardless of the context, such as laws against homocide.

So, in a given context, different sets of rules are pertinent. A single set of rules may be user in one context, while a different set is used in a different context. Within a rule-engine, available sets of rules are called *rule sets*. The set or sets of rules currently applicable given the context is called a *rule base*.

1.2   Knowledge

You become aware of knowledge, in the form of facts, over time. Likewise, over time, facts may change or cease to be true, in which case, they facts as you know them must be altered or purged from your memory. Likewise, within a rule-engine, there are operations for becoming aware of a fact, purging a fact, or modifying a known fact. These operations are as follow:

- **Assert** Add a fact to what is known.
  The weather is rainy. Betty is in the room.

- **Retract** Remove a fact from what is known.
  Removing the fact *Betty is in the room* when Betty leave the room.

- **Modify** Alter a fact from what is know.
  Changing the knowledge about the weather from *The weather is rainy* to *The weather is sunny* when the rain stops and the clouds go away.

Within a rule-engine, the collected knowledge is called the *working memory*. Knowledge is asserted, retracted and modified within the working memory and the rules are evaluated to determine what actions, if any, should be taken.

1.3   Why a rule-engine?

While the logic expressed in a rule can and has often been writen within the code a system, a rule-engine offers many benefits. Instead of locking the logic up in code written by developers, the logic can be moved out-board external to the actual application. In this way it is possible for non-developers to change the logic without having to rebuild the system. Additionally, by codifying all of the system rules in a central location, they are no longer scattered throughout the application. This allows for easier validation of the system's requirements and analysis of the logic of the system.

Additionally, a rule-engine such as Drools is built upon an intelligent algorithm that allows for the evaluation of many rules against many facts in an effecient manner. In a procedural system, a change in a single fact might require double-checking *every* rule to determine if any action needs to be taken. A rule-engine which uses the *Rete* algorithm is optimized to minimize the amount of processing effort that is required to evaluate the rules that may have been affected by a change in knowledge.

1.4   A note about "business rules"

A higher-level form of rule is the *business rule*. Business rules do not necessarily follow the *if-then* form, but may be specified in different formats that are not as closely linked to the underlying rule-engine implementation. Business rules tend to use the *must* or *must not* form to expression constraints or inferences.

- An order **must not** be billed before it ships.

- An applicant for store credit **must** be 18 years of age.

Drools does not directly support this level of business rules, but other projects built upon Drools[1] may easily support such notation.

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[1]The **Fluxtapose** suite of tools from The Werken Company is one such product that supports business rules.

Chapter 2

Architecture

2.1 Rules, rule-sets, and rule-bases

Within Drools, the concepts of rules, rule-sets and rule-bases are directly modelled by the classes `Rule`, `RuleSet` and `RuleBase`. A `Rule` may be a member of multiple `RuleSet`s, and multiple `RuleSet`s may be active within a given `RuleBase`.

2.1.1 `Rule`, `Condition` **and** `Consequence`

A single `Rule` may have one-or-more `Condition`s associated with it. Each condition must be met before the rule is considered to be activated. Once activated the rule's `Consequence` is a candidate for being fired. In pattern parlance, the `Condition` class is simply a *predicate object* which evaluates itself against the known facts to return a boolean value of either *true* or *false*. The `Consequence` class is likewise simply a *functor* which objectifies a function and performs an arbitrary task when executed.
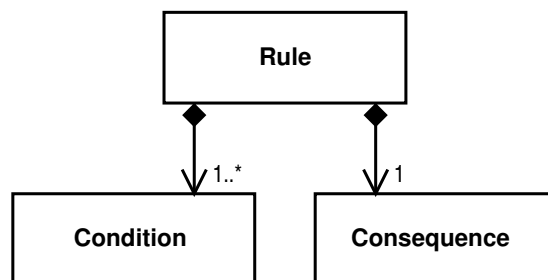
Figure 2.1: Object model for `Rule`, `Condition` and `Consequence`
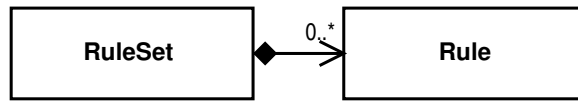
Figure 2.2: Object model for `RuleSet` and `Rule`
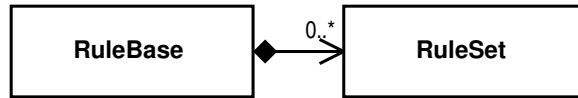
Figure 2.3: Object model for `RuleBase` and `RuleSet`

2.1.2 `RuleSet`

A `RuleSet` is simply a collection of `Rule`s. It serves only to associate a group of rules with one another so that they may be worked with as a set.

2.1.3 `RuleBase`

A `RuleBase` is an *active* collection of `RuleSet`s. A `RuleBase` contains rules that are all considered to be in effect for a given set of knowledge. Multiple `RuleSet`s may be a part of a given `RuleBase`.

2.2 Knowledge

The set of knowledge that is examined is modelled by the class `WorkingMemory`, which is backed by a particular `RuleBase`. It is through the `WorkingMemory` that knowledge is *asserted*, *retracted* and *modified*. Each `WorkingMemory` is backed by exactly one `RuleBase` which determines which rules are evaluated as knowledge is manipulated.
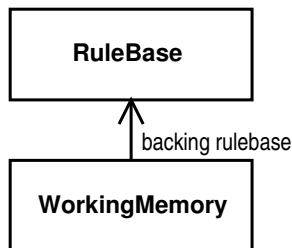
Figure 2.4: Object model for `WorkingMemory` and `RuleBase`

2.3   Complete Model
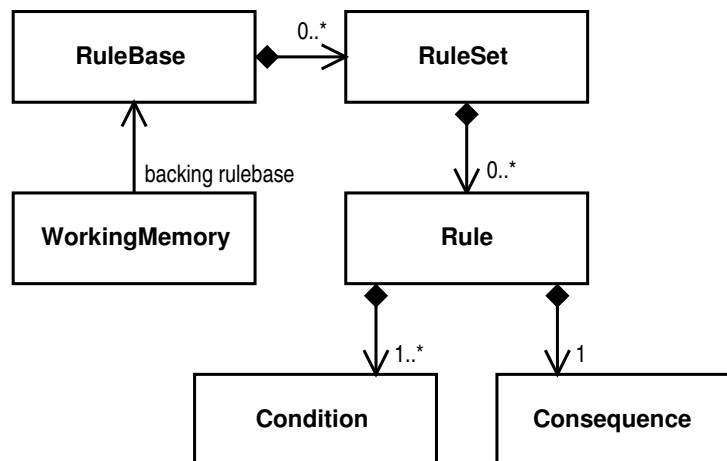
Figure 2.5: Complete object model

Chapter 3

Drools Rule Language

3.1 Introduction

The Drools Rule Language (DRL) is an XML-based language that uses an extensible tag mechanism. The available tags can be extended through the use of *semantic modules*. By itself the DRL is not fully usable. At least one semantic module must be selected and used. What follows is a reference for the available core DRL tags.

3.2 Core DRL namespace

The core DRL tags belong to the `http://drools.org/rules` namespace. In order to use the tags, the root tag must include an `xmlns` declaration binding the DRL namespace to either a prefix (Figure 3.1) or as the default namespace (Figure 3.2). Common practice is to bind the DRL namespace URI to the default namespace upon the root `<rule-set>` tag of a document.

```
<drl:rule-set xmlns:drl="http://drools.org/rules">
```

Figure 3.1: Binding the DRL namespace URI to a prefix

```
<rule-set xmlns="http://drools.org/rules">
```

Figure 3.2: Binding the DRL namespace URI as the default namespace

3.3 Core DRL tags

3.3.1 `<rule-set>`

The root tag of a DRL document is the `<rule-set>` tag, which serves to collect a group of rules. A rule-set must have a `name` attribute defining its name.

	`<rule-set>`
Attribute	**Description**
`name`	Name of the rule-set.

Tag	**Description**
`<rule>`	One or more rule members of the set.

```
<rule-set xmlns="http://drools.org/rules"
          name="My First Rule Set">
    ...
</rule-set>
```

3.3.2 `<rule>`

A `<rule>` tag defines a single rule within the rule-set. A rule must have at least one parameter[1]. A rule contains conditions, extractors and exactly one consequence.

	`<rule>`
Attribute	**Description**
`name`	Name of the rule.
`salience`	Optional rule salience/priority.

Tag	**Description**
`<parameter>`	Parameter declaration.
`<declaration>`	Additional local declaration.
`<ns:condition>`	Semantic condition.
`<ns:extractor>`	Semantic extractor.
`<ns:consequence>`	Semantic consequence.
`<duration>`	Truthness duration.

---

[1]Semantic modules may add implicit parameters to rules, so it's not strictly true that each `<rule>` must have a `<parameter>` or a `<consequence>`

```
<rule-set xmlns="http://drools.org/rules"
          name="My First Rule Set">

    <rule name="My First Rule"
          salience="10">
        ...
    </rule>

</rule-set>
```

3.3.3 `<parameter>`

A `<parameter>` defines an input parameter for a rule. Each input parameter will be supplied an object from the working memory in order to attempt satisfying the conditions. Each `<parameter>` is required to have an `identifier` that is unique within the rule and a semantic object type as its child element.

	`<parameter>`
Attribute	**Description**
`identifier`	Parameter identifier.

Tag	**Description**
`<ns:objectType>`	Object type of the parameter.

```
<rule name="My First Rule">

    <parameter identifier="factOne">
        ...
    </parameter>
    <parameter identifier="factTwo">
        ...
    </parameter>

</rule>
```

3.3.4 `<declaration>`

A `<declaration>` defines an addition *local* declared fact which may be the target of an *extractor*. Its `identifier` must be unique across all `<parameter>` and `<declaration>` tags. As with a `<parameter>`, it must have a semantic object type as its child element.

	`<declaration>`
Attribute	**Description**
`identifier`	Declaration identifier.

Tag	**Description**
`<ns:objectType>`	Object type of the local declaration.

```
<rule name="My First Rule">

    <declaration identifier="factOne">
        ...
    </declaration>
    <declaration identifier="factTwo">
        ...
    </declaration>

</rule>
```

3.3.5 `<duration>`

The `<duration>` tag is optionally used to specify the truthness duration of the
rule. The truthness duration determines how long the rule's conditions must be
continually true in order to activate the rule.

	`<duration>`
Attribute	**Description**
`weeks`	Number of weeks.
`days`	Number of days.
`hours`	Number of hours.
`minutes`	Number of minutes.
`seconds`	Number of seconds.

```
<rule name="My First Rule">

    <duration hours="2"
              minutes="30"/>

</rule>
```

Chapter 4

Java Semantic Module

Chapter 5

Python Semantic Module

Chapter 6

XML Semantic Module

Chapter 7

Client API

The Drools client API is simply a fairly lightweight interface for starting a session, manipulating knowledge and executing rules.

7.1 Retrieving a `RuleBase`

Before any activites may be performed, the client application must obtain a `RuleBase` through either building one using the administrative API or by performing a lookup within a JNDI directory or other managed object data source. Building a `RuleBase` is discussed in the next chapter.

7.1.1 Looking up a `RuleBase` from JNDI

A `RuleBase` may be a managed object which is deployed by an administrator within a JNDI directory. This method of deployment allows an application to be isolated from changes in the rules. The rules can be easily altered and redeployed, enabling an on-the-fly modification of the behaviour of an application. See Figure 7.1 for an example of accessing a `RuleBase` deployed within a JNDI data source.

7.1.2 Deserializing a `RuleBase` from a file

Fully assembled `RuleBase` objects may be serialized to a file or other bytestream. This allows for the rules to be assembled before deployment and enables easy

```
1   InitialContext ctx = ...
```

Figure 7.1: Loading a `RuleBase` from a JNDI data source

```
1   File ser = new File(...);
```

Figure 7.2: Deserializing a `RuleBase` from a file

```
1   RuleBase ruleBase = obtainRuleBase();
2
3   WorkingMemory workingMemory = ruleBase.newWorkingMemory();
```

Figure 7.3: Creation of a `WorkingMemory` from a `RuleBase`

loading. See Figure 7.2 for an example of loading a serialized `RuleBase` from a file.

7.2   Creating a `WorkingMemory`

A `WorkingMemory` object represents an isolated rule sessio n where knowledge is manipulated and evaluated to determine which rules should fire. Each `WorkingMemory` is independent from all other instances. Each `WorkingMemory` is backed by the `RuleBase` that created it (Figure 7.3).

7.3   Working with knowledge

The three knowledge manipulation operations that can be performed are *assertion*, *retraction*, and *modification*. The `WorkingMemory` class models these three operations through three methods:

- `assertObject(...)` Adds a fact to the working knowledge.

- `retractObject(...)` Removes a fact from the working knowledge.

- `modifyObject(...)` Modifies a fact within the working knowledge.

7.3.1   Adding knowledge

To add knowledge to a `WorkingMemory`, a fact, in the form of any Java object is asserted into it (Figure 7.4 on the next page). A `FactHandle` is returned from each assertion and is used when performing further manipulation of the asserted fact. Once asserted, the fact may be used to satisfy some portion of any rule's activation condition.

```
1   WorkingMemory memory = ruleBase.newWorkingMemory();
2
3   FactHandle handle1 = memory.assertObject( factOne );
4   FactHandle handle2 = memory.assertObject( factTwo );
5   FactHandle handle3 = memory.assertObject( factThree );
```

Figure 7.4: Aasserting knowledge into a `WorkingMemory`

```
1   WorkingMemory memory = ruleBase.newWorkingMemory();
2
3   FactHandle handle = memory.assertObject( factOne );
4
5   memory.retractObject( handle );
```

Figure 7.5: Retracting knowledge from a `WorkingMemory`

7.3.2 Removing knowledge

To remove knowledge from a `WorkingMemory`, a `FactHandle` obtained from a previous invocation of `assertObject(...)` is used to reference the fact to be removed (Figure 7.5). Once retracted, the fact can not be used to satisfy any portion of any rule's activation condition.

7.3.3 Modifying knowledge

Over time, knowledge may change. Instead of always retracting old facts and asserting new ones, it is sometimes advisable to view a change in knowledge as the *modification of a single fact*. For example, the fact about the current state of the weather may change as the weather changes. Instead of retracting a "rainy" fact and asserting a "sunny" fact, it is possible to simply change the value of the general "weather" fact (Figure 7.6 on the next page). As with retraction, modification of knowledge depends on a `FactHandle` obtained from a previous `assertObject(...)` call (Figure 7.6 on the following page).

The most important reason to perform fact modification instead of an assertion and a retraction is that rules with truth durations will automatically deactivate if a matching fact is retracted. If a fact is modified but still matches a rule with a truth duration, the clock on the duration timer is *not* reset.

7.4 Firing rules

As knowledge is manipulated within a `WorkingMemory`, no rules are actually triggered to be fired. Instead, rules are *activated*, which makes them candidates for being fired. Rule activations are placed upon the *agenda*.

```
1   WorkingMemory memory = ruleBase.newWorkingMemory();
2
3   FactHandle weatherHandle = memory.assertObject( new Weather( "rainy" ) );
4
5   memory.modifyObject( weatherHandle,
6                        new Weather( "sunny" ) );
```

Figure 7.6: Modifying knowledge within a `WorkingMemory`

```
1   WorkingMemory memory = ruleBase.newWorkingMemory();
2
3   FactHandle weatherHandle = memory.assertObject( new Weather( "rainy" ) );
4
5   memory.fireAllRules();
```

Figure 7.7: Firing rules

When an application determines that the current state of knowledge is consistent and should be acted upon, the application may cause the activations on the agenda to begin firing (Figure 7.7). The `fireAllRules(...)` method blocks until the agenda has been completely cleared of activations.

Note that as an activation fires, its consequence may perform knowledge manipulation by assertion new facts, retracting existing facts, or modifying existing facts. All of these activities may cause activations to be added to or removed from the agenda.

Chapter 8

Administrative API

The administrative functions of drools include the assembly and deployment of rules. These activities may certainly be performed by the same application that acts as the client, The API has been designed to allow for the separation between the client and administrative responsibilities.

8.1   Building rules

Programmatically building rules using the Java API is covered in Chapter 10 on page 24. The next section discusses loading rules that were built using the Drools Rule Language, which is described in depth in Chapter 3 on page 7.

8.2   Loading rules

Rule-sets can be loaded using the `RuleSetReader` from a variety of sources including `URL`, `InputStream` and `Reader` objects. A `RuleSetReader` builds a `RuleSet` from the contents of a DRL document (Chapter 3 on page 7).

The `RuleSetReader` by default may use any and all semantic modules available on the classpath (see Chapter 9 on page 20). It may be optionally parameterized with a `SAXReader` and a `SemanticsRepository`. By default, the `RuleSetReader` uses the Java API for XML Parsing (JAXP) for reading the XML documents. See Figure 8.1 for an example of loading a `RuleSet` from a `URL`.

```
1   RuleSetReader reader     = new RuleSetReader();
2   URL           ruleSetUrl = new URL( "http://myco.com/theRuleSet.xml" );
3   RuleSet       ruleSet    = reader.read( ruleSetUrl );
```

Figure 8.1: Loading a `RuleSet` from a `URL`

8.3   **Building a** `RuleBase`

8.4   **Deploying a** `RuleBase` **to a JNDI data store**

8.5   **Serializing a** `RuleBase` **to a file**

Chapter 9

Semantic Module Framework

The Semantic Module Framework (SMF) provides an extension point to the Drools Rule Language syntax (see Chapter 3 on page 7). By creating semantic modules, domain-specific rule types, object-types, conditions, extractors and consequences can be added to the core DRL language.

9.1 Semantic components

9.1.1 `RuleType`

9.1.2 `ObjectType`

Implementations of `ObjectType` allow rules to be defined in terms of semantic types. Instead of requiring all rules to be defined in terms of Java classes, rules may be defined in terms of higher semantics. For example, while all XML documents may be instances of `org.w3c.dom.Document` each document may have a different semantic type based upon the name and namespace of the root tag. Given an `Object`, an `ObjectType` implementation must simply determine if it matches its semantic type (Figure 9.1).

```
1   package org.drools.spi;
2
3   public interface ObjectType
4       extends SemanticComponent
5   {
6       boolean matches(Object object);
7   }
```

Figure 9.1: `ObjectType` interface

```
1  package org.drools.smf;
2
3  public interface ConfigurableObjectType
4      extends ObjectType
5  {
6      void configure(Configuration config) throws ConfigurationException;
7  }
```

Figure 9.2: `ConfigurableObjectType` interface

`ConfigurableObjectType`

Configuration information may be passed to an `ObjectType` if the implementation is marked with the `ConfigurableObjectType` interface. A `configure(...)` method is added by the `ConfigurableObjectType` interface. through which additional information may be passed (Figure 9.2).

9.1.3 `Condition`

`ConfigurableCondition`

9.1.4 `Extractor`

`ConfigurableExtractor`

9.1.5 `Condition`

`ConfigurableCondition`

9.1.6 `Consequence`

`ConfigurableConsequence`

9.2 **The** `Configuration` **structure**

For the `Configurable...` form of the semantic components, configuration information is communicated through a tree of `Configuration` objects. Each `Configuration` object acts as a node in the tree, and may contain the following data:

- **Attributes.** Zero or more name/value pairs of strings (Figure 9.3 on the following page).

- **Text.** A single string text value (Figure 9.4 on the next page).

- **Child `Configuration` nodes.** Zero or more named child `Configuration` nodes (Figure 9.5 on the following page).

```
1  String[] attrNames = config.getAttributeNames();
2  String   someValue = config.getAttribute( someName );
```

Figure 9.3: Attribute-related operations of `Configuration`

```
1  String nodeText = config.getText();
```

Figure 9.4: Text-related operation of `Configuration`

The tree of `Configuration` nodes may be thought of as a simplified version of an XML structure. For configurable semantic components used through the DRL (Chapter 3 on page 7), the root `Configuration` is based upon the component's own tag, and children tags are represetned by children `Configuration` nodes.

9.3   Semanitc Module Descriptor

9.4   `SemanticsRepository`

A `SemanticsRepository` manages a set of `SemanticModule` objects and allows each to be looked-up by its URI (Figure 9.6 on the following page). Primary a `SemanticsRepository` is used by a `RuleSetReader` (Section 8.2 on page 18) in order to extend the core DRL (Chapter 3 on page 7) syntax.

9.5   The `DefaultSemanticsRepository` **helper**

The `DefaultSemanticsRepository` helper class is useful in that it contains all conforming semantic modules available on the classpath. Each module that has a module descriptor (Section 9.3) located as `META-INF/drools-semantics.xml` will be automatically discovered by the `DefaultSemanticsRepository` upon first use.

Being a help class that is initialized once, it follows the singleton pattern.

```
1  Configuration[] allChildren    = config.getChildren();
2  Configuration   someFirstChild = config.getChild( someName );
3  Configuraiton[] someChildren   = config.getChildren( someName );
```

Figure 9.5: Child-related operations of `Configuration`

```
1  SemanticsRepository repo       = locateSemanticsRepository();
2  SemanticModule[]    modules    = repo.getSemanticModules();
3  SemanticModule      someModule = repo.lookupSemanticModule( someUri );
```

Figure 9.6: Usage of the `SemanticsRepository`

```
1  SemanticsRepository repo    = DefaultSemanticsRepository.getInstance();
2  SemanticModule[]    modules = repo.getSemanticModules();
```

Figure 9.7: Retrieving and using the `DefaultSemanticsRepository` helper

To use the `DefaultSemanticsRepository`, the `getInstance()` method will retrieve the singleton instance (Figure 9.7).

Chapter 10

Rule Assembly API

Chapter 11

Examples

Chapter 12

Algorithms

12.1 Efficient Matching

While it may be simple to create a rules engine that allows specification of business logic in a format that is comfortable to business analysts, the matching of the rules may still be problematic without a good algorithm.

The rules engine must be made aware of its environment, typically through a process called *fact assertion*. Fact assertion consists of the program asserting facts into a rules session, or *working memory*.

Whenever a fact is asserted, retracted or modified within the working memory, many rules may become candidates for firing, or may have become invalidated. A simplistic approach is to reevaluate all rules against the entirety of the working memory. This method is guaranteed to be correct but will also certainly be sub-optimal. Any individual fact modification only affects a small number of conditions in a small number of rules.

Variations of the Rete algorithm allow the rules engine to maintain a memory of the results of partial rule matches across time. Reevaluation of each condition is no longer necessary, as the engine knows which conditions might possibly change for each fact, and only those must be reevaluated.

12.2 Rete

Charles Forgy created the original Rete algorithm [**?**] around 1982 as part of his DARPA-funded research. Compared to many previous production-matching algorithms, Rete was very advanced. Even today, there have been few improvements to it in the general case[1]. Variations on Rete, such as TREAT [**?**], may have different performance characteristics depending on the environment. Some perform better with large rule sets but small numbers of objects, while other

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[1]Both ILOG and Haley claim to have optimized Rete algorithms, but details are not currently public.

perform well for steady-state environments, but react poorly to numerous successive changes in the data.

A *Rete network* is a graph through which data flows. Originally, data was specified using Cambridge-prefix tuples since Lisp-like languages were in style for logic programming.[2] The tuples were used to express attributes about objects. For example, tuples may be used to express a person's name and her pets. The tuples are dropped into the Rete network, and those that reach the far end cause the firing of a rule. The original production-matching was based upon matches against tuple patterns.

The Rete network is comprise of two types of nodes:

- **1-input/1-output nodes**
  The *1/1* nodes are constrictive nodes that only allow matching tuples to flow through. Any tuples that do not match are discarded by the node.

- **2-input/1-output nodes**
  The *2/1* nodes simply connect the output arcs from two other nodes (either *1/1* nodes or *2/1* nodes) merging tuples from both the left and right incoming arcs into a single tuple on the outgoing arc. Maintains a memory of tuples for matching against future facts.

A forest of *1/1* nodes acts as the entry-point into the entire Rete network for any incoming tuple. The network-entry nodes filter tuples purely by their type. Tuples about dogs and tuples about cats may each have a different type and may be differentiated from each other by the *1/1* network-entry nodes.

Each condition of a rule is merely a pattern for a particular tuple type. The condition describes the attributes that a tuple must have and acts as a filter. Each condition is transformed into a *1/1* node that only allows tuples matching the specified attributes to pass. An attribute value may be specified as a variable and implies that the variable must hold the same value in all occurrences. The *1/1* filter nodes are attached to the network downstream from the *1/1* entry-node that differentiates their tuple type.

Consider a condition such as "For any person who has a dog that has the same name as that person's sister's cat, then..." This could be expressed with the condition patterns of:

```
(1)  ( person   name=person?   sister=sister?  )
(2)  ( person   name=person?   dog=petName?  )
(3)  ( person   name=sister?   cat=petName?  )
```

Condition *#1* models the sister relationship so that the rule only applies to two people who are sisters. The `person?` and `sister?` tokens are variables that must be consistent across any set of tuples that match this rule.

Conditions *#2* and *#3* serve two roles. The `dog` and `cat` attributes share the same `petName?` variable and serve to identify two people who have a cat

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[2]As it is for many artificial intelligence projects.

and a dog with the same name. They each contain a `name` attribute with either the variable `person?` or `sister?` which ties the last two conditions back to the first two.
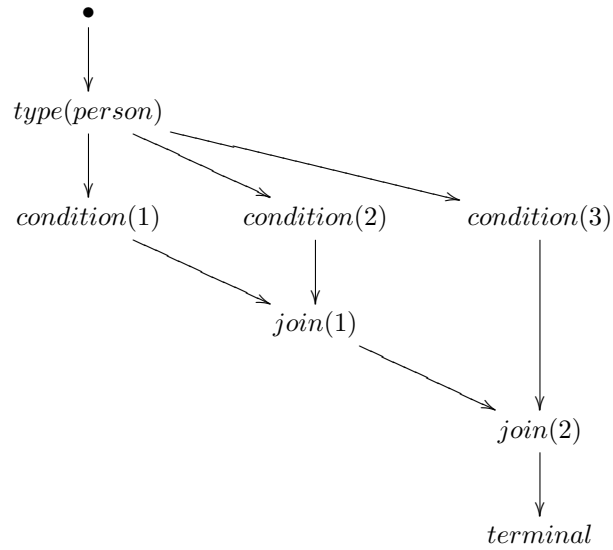

Figure 12.1: Rete network

type	person	sister	cat	dog
tuple set # 1				
person	rebecca	jeannie	zoomie	*null*
person	jeannie	rebecca	*null*	zoomie
tuple set # 2				
person	rebecca	jeannie	zoomie	*null*
person	jeannie	rebecca	*null*	toby

Figure 12.2: Example tuple sets

If two sets of tuples (see Figure 12.2) were asserted against the rule, *tuple set #1* would cause a firing of the rule, where *tuple set #2* would not. In both cases, the two tuples would pass node *condition(1)*, as the nodes simply associate the `person?` and `sister?` variables with the appropriate values from each tuple.

The *join(1)* node would allow both tuples to merge and propagate past it in both the first and second case. Additionally, for both cases, the *rebecca* tuple would pass node *condition(2)* and the *jeannie* tuple would pass node *condition(3)*.

The *join(2)* node is where the two cases differ. In the first case, nodes *condition(2)* and *condition(3)* have each associated the value of "ugly" to the `petName?` variable. In the second case, the two nodes has assigned different values to the variable. The *join(2)* node only allows those tuples that have consistent associations with all variables to pass.

12.3   Rete-OO

The Rete algorithm works wonderfully in language systems such as Lisp where pertinent attributes about objects are directly asserted to the rules engine. In an object-oriented language, such as C++ or Java, and entire graph of objects can be reachable from a single named root object. Expressing highly complex relationships between entities using Cambridge-prefix notation may require many separate assertions. In an OO language, the single root object is all that should be asserted, since attributes and relationships can be *extracted* using normal language constructs.

Bob McWhirter of The Werken Company adapted Forgy's original Rete algorithm to object-oriented constructs, creating the Rete-OO algorithm. As with Rete, there are *1/1* nodes and *2/1* nodes. Unlike Rete, there are nodes that exist simply to extract reachable attributes and add columns to passing tuples. Rete always constructs the condition *1/1* nodes toward the root of the tree leaving the bottom portion to be comprised of purely aggregating *2/1 join* nodes. Rete-OO must interleave both *1/1* and *2/1* nodes.

The same example as in Section 12.2, the conditions could be expressed in terms of object-oriented language boolean and assignment expressions. The choice of Java as the expression language is purely arbitrary.

```
(0)   Person personOne, personTwo
(1)   personOne.hasSister(personTwo)
(3)   petName = personOne.getCat().getName()
(3)   petName = personTwo.getDog().getName()
```

Rete-OO adds the concept of *root object declaration*, where the root objects of the condition are declared with a name and type. The object's type maps directly to the tuple type in Rete. The root object name has no direct mapping in Rete and causes the addition of a *parameter node* in Rete-OO. Boolean

expressions in Rete-OO conditions are equivalent to Rete's condition patterns against attributes. The assignment expressions map to place-holder variables in Forgy's algorithm.

The types of nodes used in Rete-OO graph construction are listed here. Those that are new or different from Rete are denoted with a '*'.

- **Object type**
  Object type nodes differentiate objects by filtering on their defined type.

- **Parameter***
  Parameter nodes create a tuple with a single entry binding the object to the name.

- **Condition**
  Condition nodes simply tests a tuple against an a boolean expression.

- **Extraction***
  Extraction nodes extract new attributes, create new columns on tuples, and store the results.

- **Join**
  Join nodes connect the output arcs from two other nodes and allows consistent tuples to be merged and passed through.

- **Terminal**
  Terminal nodes fire to indicate a successful match for the rule.

The resulting Rete-OO graph is constructed in a different manner than the equivalent Rete graph, due to the addition and rearrangement of some nodes.
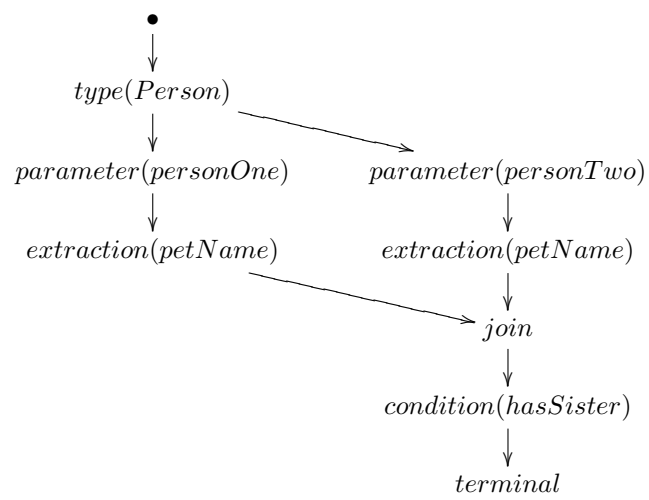
Figure 12.3: Rete-OO network

Appendix A

Project Information

A.1 Web Site

All development resources related to Drools are hosted by **The Codehaus**, the open-source arm of The Werken Company. Drools maintains a website at:

```
http://drools.org/
```

A.2 Mailing Lists

The drools project maintains two mailing lists. The first, known as `drools-interest` is for general discussion by users and developers of drools. The second list is `drools-cvs` which simply tracks changes made to the source-code through the CVS repository. For information about subscribing to each list or access to the list archives:

```
http://lists.codehaus.org/listinfo/drools-interest
http://lists.codehaus.org/listinfo/drools-cvs
```

A.3 Source Repository

The drools project maintains a revision control repository using CVS. To check-out the latest sources, you must issue two CVS commands. The first is used to login. When presented with a prompt for a password, simply press *ENTER*.

```
cvs -d:pserver:anonymous@cvs.codehaus.org:/scm/cvspublic login
cvs -d:pserver:anonymous@cvs.codehaus.org:/scm/cvspublic co drools
```

A.4   Internet Relay Chat

There is a dedicated channel on The Werken Company's IRC server for drools:

address	`irc.codehaus.org`
port	`6667`
channel	`#drools`
url	`irc://irc.codehaus.org:6667/drools`

A.5   Bug, Issue & Feature Tracking

For bug, issue and feature tracking, the Drools project uses the Jira project management system provided by The Codehaus.

    http://jira.codehaus.org/

A.6   Project Team

A.6.1   Bob McWhirter

Bob McWhirter originally founded the Drools project in 2000 and developed the Rete-OO algorithm used by the engine. Bob is also the founder of The Werken Company and the chief architect behind the commercial **Fluxtapose** suite of tools which build upon Drools to provide a complete solution for implementing business rules.

A.6.2   Thomas Diesler

Thomas Diesler researched and supplied the JSR-94 Rule-Engine API bindings for Drools.

A.6.3   Roger F. Gay

Roger F. Gay devised the XML Schemas for the core DRL syntax and each semantic module.

A.6.4   Contributors

Others have contributed ideas, patches and testing assistance over the years:

- Dave Cramer *(eBox)*
- Martin Hald
- Matt Ho
- Pete Kazmier *(iBasis)*
- Christiaan ten Klooster
- James Roome
- Bart Selders *(iBanx)*
- James Strachan *(CoreDevelopers Network)*
- Tom Vasak

Appendix B

Licensing

B.1   Drools license

B.1.1   The license

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B.1.2   Summary

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Index

This manual was produced without the aid of Microsoft products. It was authored on a Linux laptop using XEmacs and LaTeX. Drafts were previewed using `xdvi`. Diagrams were prepared using `dia`, exported in EPS format and converted for inclusion in the PDF document with `epstopdf`. The index was prepared using the `makeindex` package.