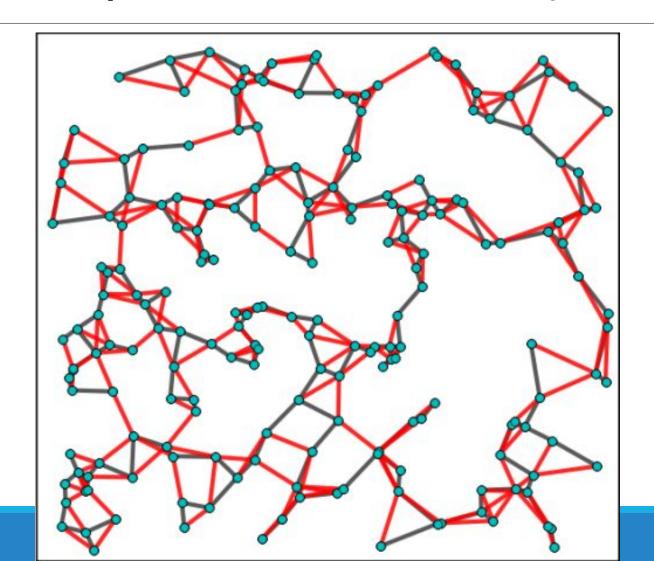
Grafos 2

Temario

- >Árbol de Expansión Mínimo
- ➤ Algoritmo de Kruskal
- ➤ Algoritmo de Prim
- ➤ Algoritmo de Dijkstra



Entrada:

➤ Un grafo no dirigido (conectado) G = (V,E) con una función de pesos

 \rightarrow W: E \rightarrow R

Salida:

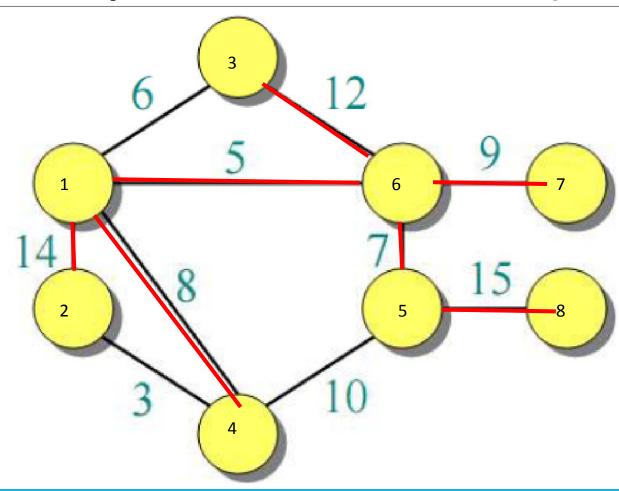
Un árbol de expansión T (árbol que conecta todos los nodos con un mínimo peso)

$$w(T) = \sum_{(u,v)\in T} w(u,v)$$

Ejemplo:

Entrada

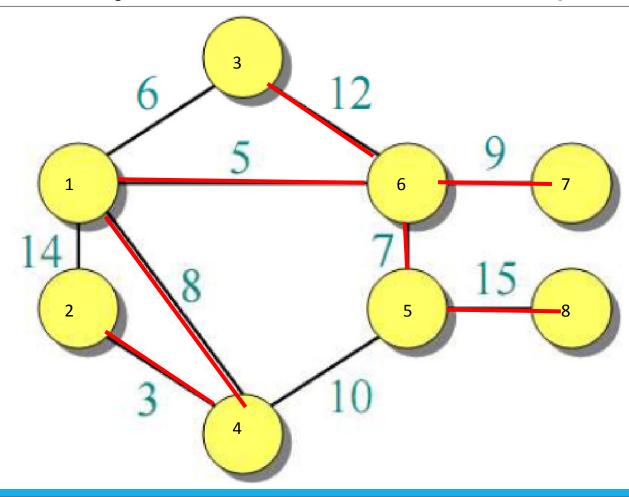
MST = 70



Ejemplo:

Entrada

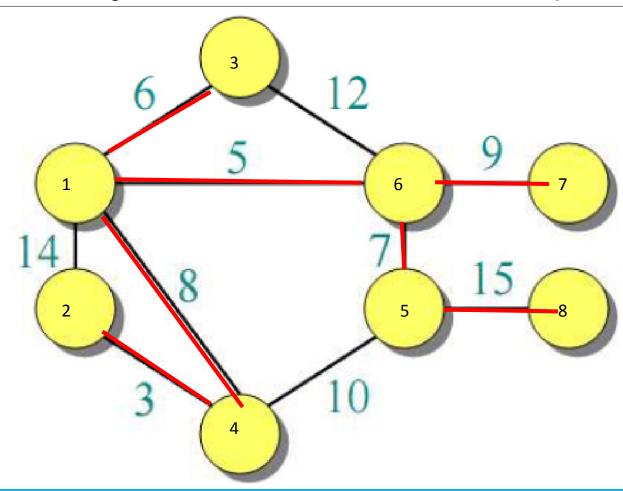
MST = 59



Ejemplo:

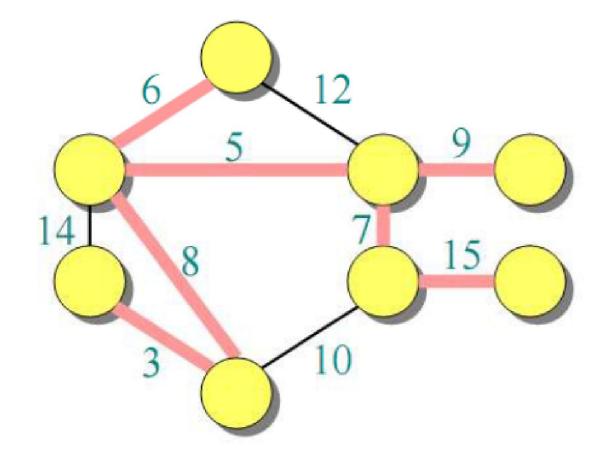
Entrada

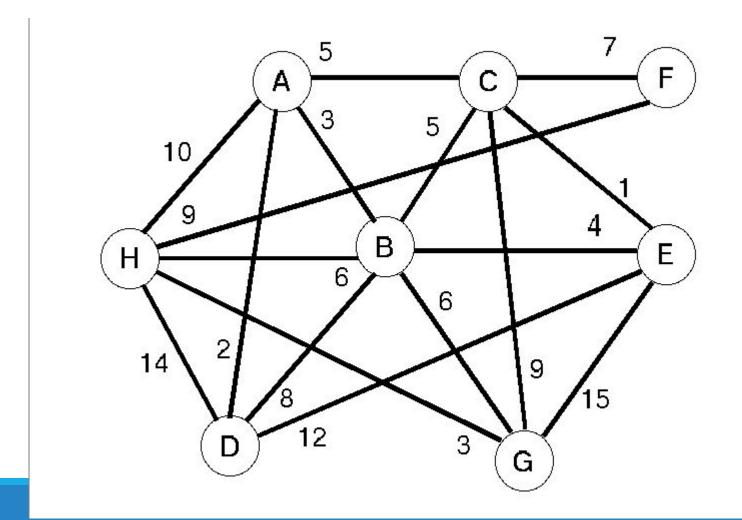
MST = 53



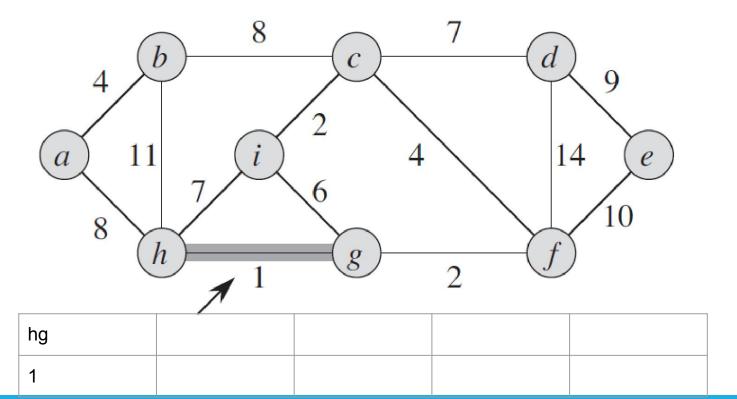
Ejemplo:

Salida

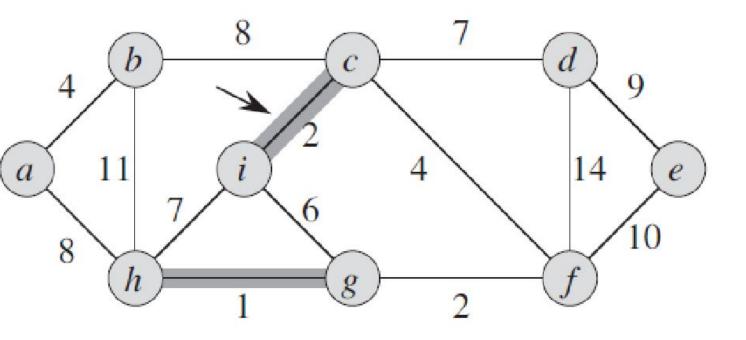




Ejemplo

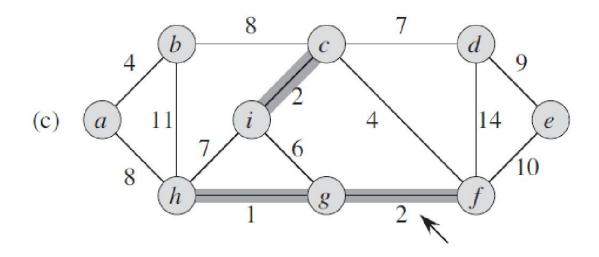


hg	1
gf	2
ic	2
cf	4
ab	4
ig	6
hi	7
cd	7
bc	8
ah	8
de	9
fe	10
bh	11
df	14



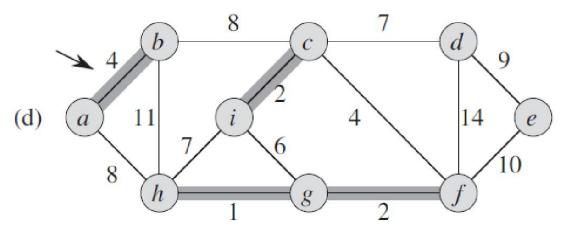
hg	ic		
1	2		

<mark>hg</mark>	1
ic	2
gf	2
cf	4
ab	4
ig	6
hi	7
cd	7
bc	8
ah	8
de	9
fe	10
bh	11
df	14



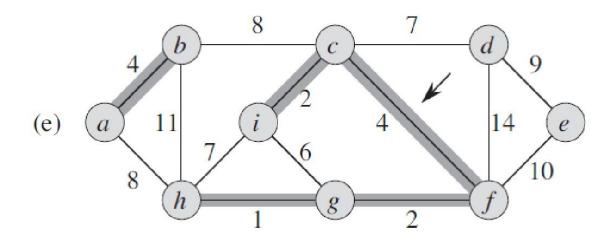
hg	ic	gf	
1	2	2	

<mark>hg</mark>	1
ic	2
gf	2
ab	4
cf	4
ig	6
hi	7
cd	7
bc	8
ah	8
de	9
fe	10
bh	11
df	14



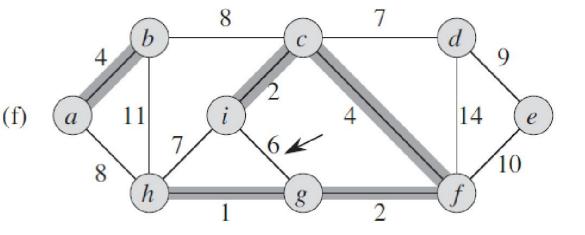
hg	ic	gf	ab	
1	2	2	4	

hg	1
ic	2
gf	2
ab_	4
cf	4
ig	6
hi	7
cd	7
bc	8
ah	8
de	9
fe	10
bh	11
df	14



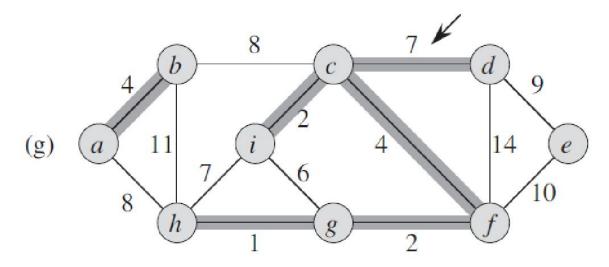
hg	ic	gf	ab	cf
1	2	2	4	4

hg	1
ic	2
gf	2
ab	4
cf	4
ig	6
hi	7
cd	7
bc	8
ah	8
de	9
fe	10
bh	11
df	14



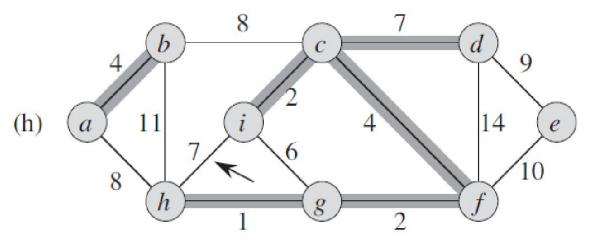
hg	ic	gf	ab	cf
1	2	2	4	4

hg	1
ic	2
gf	2
ab	4
cf	4
ig	6
cd	7
hi	7
ah	8
bc	8
de	9
fe	10
bh	11
df	14



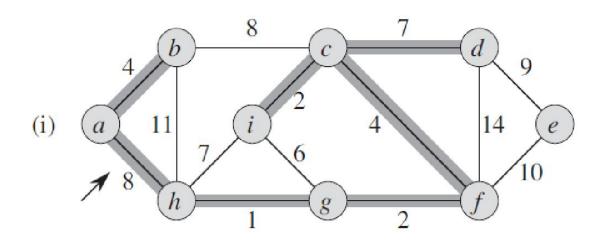
hg	ic	gf	ab	cf	cd
1	2	2	4	4	7

hg	1
ic	2
gf	2
<mark>ab</mark>	4
<u>cf</u>	4
ig	6
cd	7
hi	7
ah	8
bc	8
de	9
fe	10
bh	11
df	14



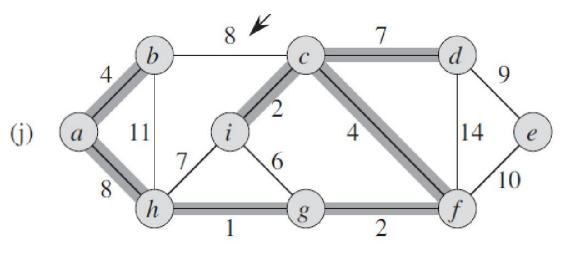
hg	ic	gf	ab	cf	cd
1	2	2	4	4	7

hg	1
ic	2
gf	2
ab	4
<mark>cf</mark>	4
ig	6
cd	7
hi	7
ah	8
bc	8
de	9
fe	10
bh	11
df	14



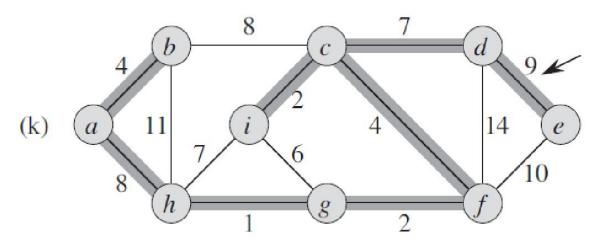
hg	ic	gf	ab	cf	cd	ah
1	2	2	4	4	7	8

hg	1
ic	2
gf	2
ab	4
cf	4
ig	6
cd	7
hi	7
<mark>ah</mark>	8
bc	8
de	9
fe	10
bh	11
df	14



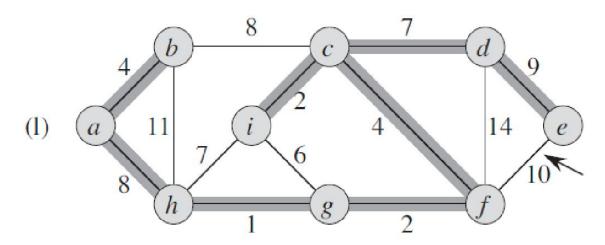
hg	ic	gf	ab	cf	cd	ah
1	2	2	4	4	7	8

hg	1
ic	2
gf	2
ab	4
cf	4
ig	6
cd	7
hi	7
<mark>ah</mark>	8
bc	8
de	9
fe	10
bh	11
df	14



hg	ic	gf	ab	cf	cd	ah	de
1	2	2	4	4	7	8	9

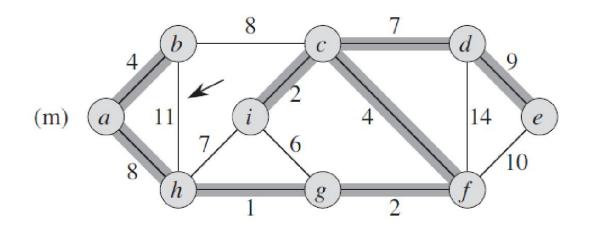
hg	1
ic	2
gf	2
ab	4
cf	4
ig	6
cd	7
hi	7
ah	8
bc	8
de	9
fe	10
bh	11
df	14

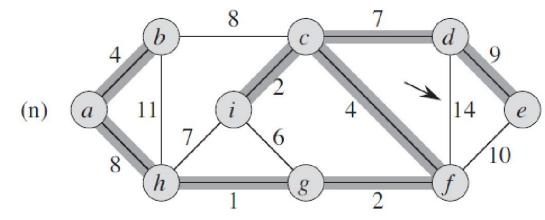


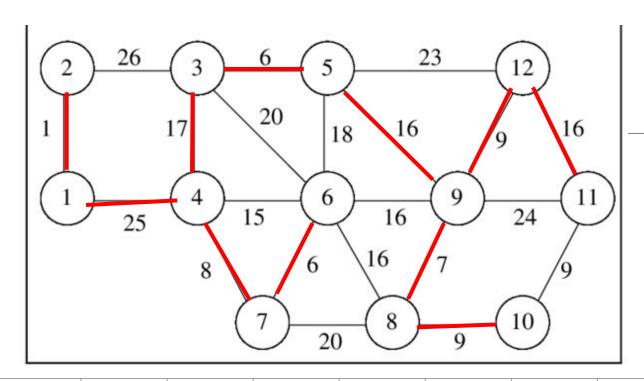
hg	ic	gf	ab	cf	cd	ah	de	Total
1	2	2	4	4	7	8	9	37

hg	1
ic	2
gf	2
ab	4
cf	4
ig	6
cd	7
hi	7
ah	8
bc	8
de	9
fe	10
bh	11
df	14

Ejemplo





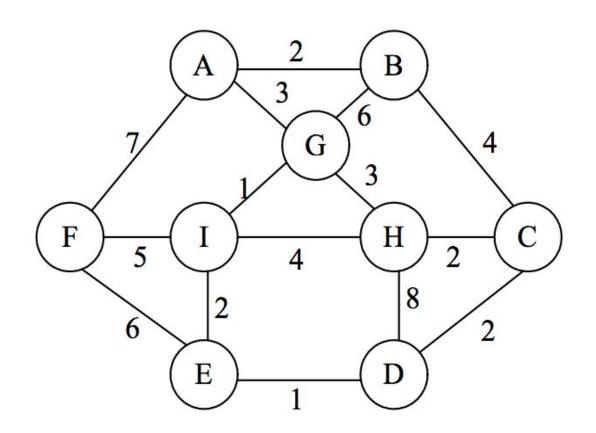


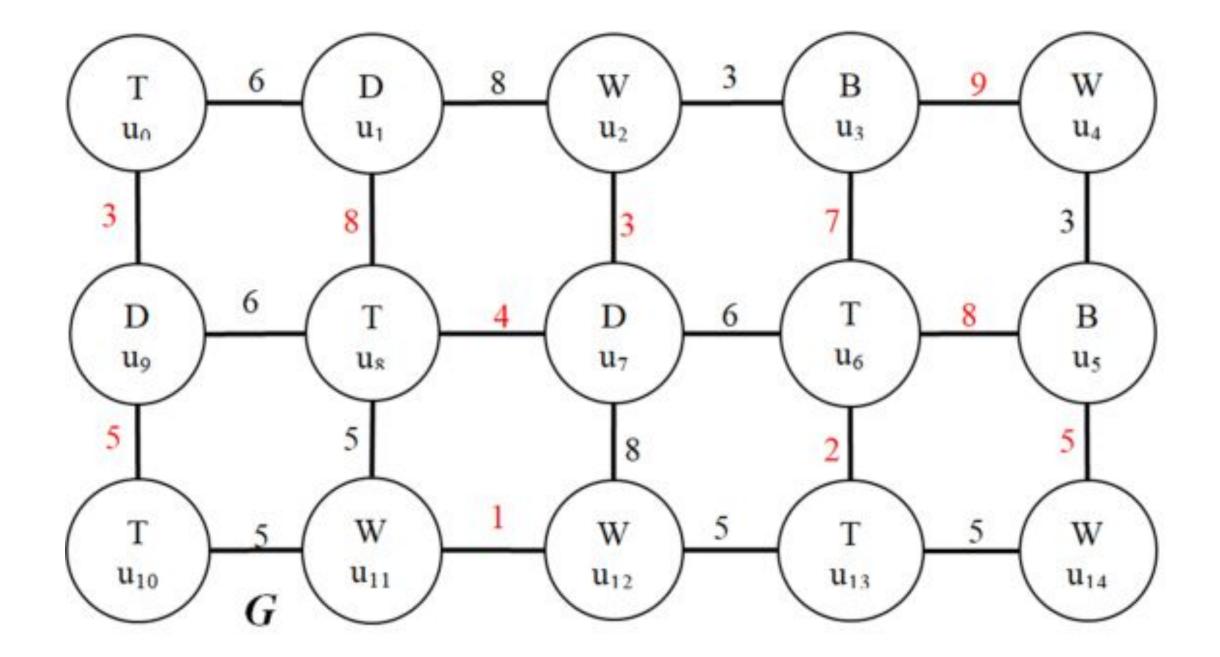
1,2	3,5	6,7	9,8	4,7	12,9	11,10	8,10
1	6	6	7	8	9	9	9

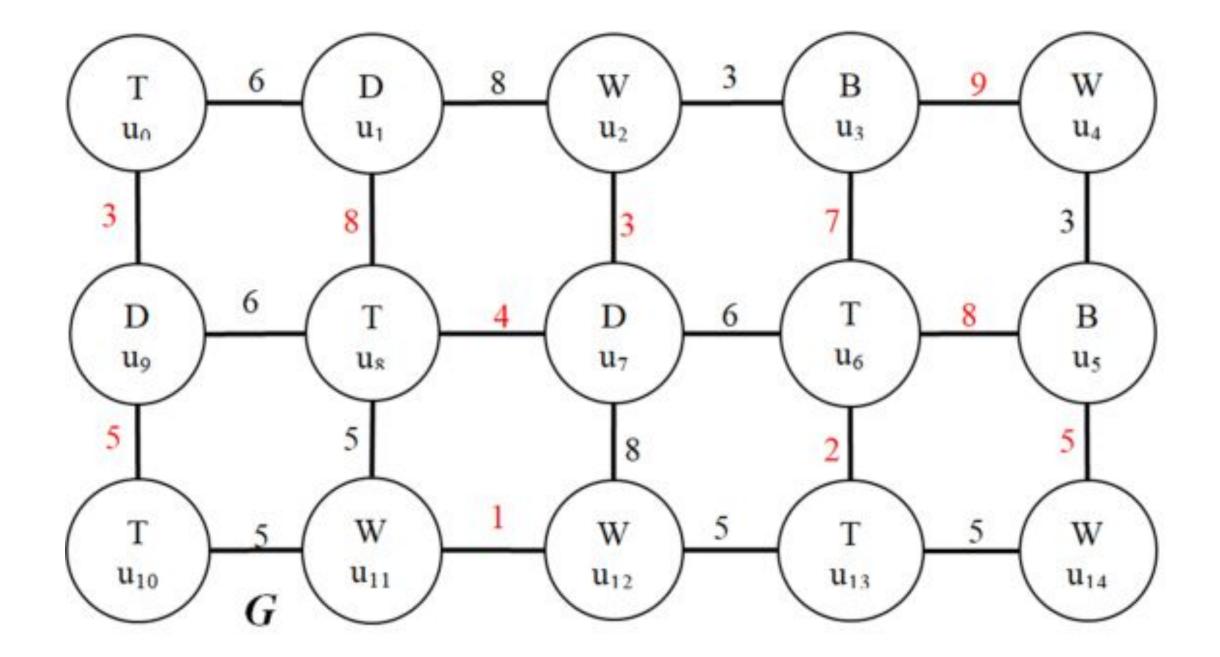
6,9	5,9	12,11	3,4	SUMA TOTAL
16	16	16	17	120

1,2	1	<mark>5,6</mark>
<mark>3,5</mark>	6	<mark>3,6</mark>
<mark>6,7</mark>	6	<mark>7,8</mark>
9,8	7	5,12
<mark>4,7</mark>	8	9,11
12,9	9	<mark>1,4</mark>
11,10	9	<mark>2,3</mark>
8,10	9	
<mark>4,6</mark>	15	
<mark>6,9</mark>	16	
6,8	16	
<mark>5,9</mark>	16	
12,11	16	
3,4	17	

<mark>5,6</mark>	18
3,6	20
7,8	20
5,12	23
9,11	24
<mark>1,4</mark>	25
2,3	26





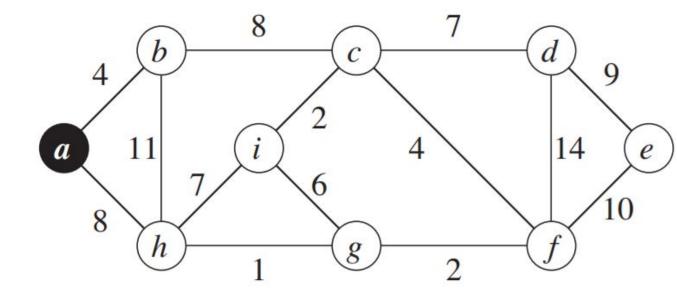


```
MST-PRIM(G, w, r)
    for each u \in G.V
        u.key = \infty
         u.\pi = NIL
   r.key = 0
   Q = G.V
    while Q \neq \emptyset
         u = \text{EXTRACT-MIN}(Q)
 8
         for each v \in G.Adj[u]
 9
              if v \in Q and w(u, v) < v. key
10
                  \nu.\pi = u
                  v.key = w(u, v)
```

nodo	а	b	С	d	е	f	g	h	i
key	0	∞	∞	∞	∞	∞	∞	∞	∞
padre	-	-	-	-	-	-	-	-	-

MST-PRIM(G, w, r)

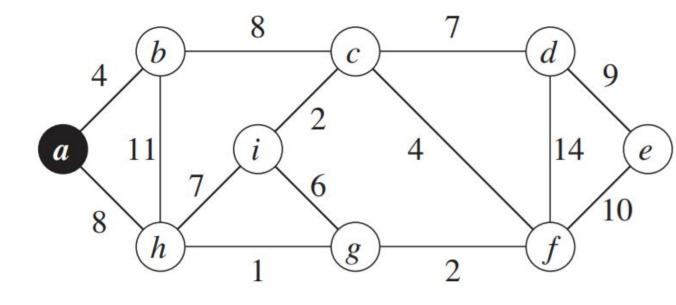
```
for each u \in G.V
         u.key = \infty
        u.\pi = NIL
    r.key = 0
    Q = G.V
    while Q \neq \emptyset
         u = \text{EXTRACT-MIN}(Q)
 8
         for each v \in G.Adj[u]
              if v \in Q and w(u, v) < v.key
10
                   \nu.\pi = u
11
                  v.key = w(u, v)
```



nodo	а	b	С	d	е	f	g	h	i
key	0	∞/4	∞	∞	∞	∞	∞	∞/8	∞
padre	-	-/a	-	-	-	-	-	-/a	-

MST-PRIM(G, w, r)

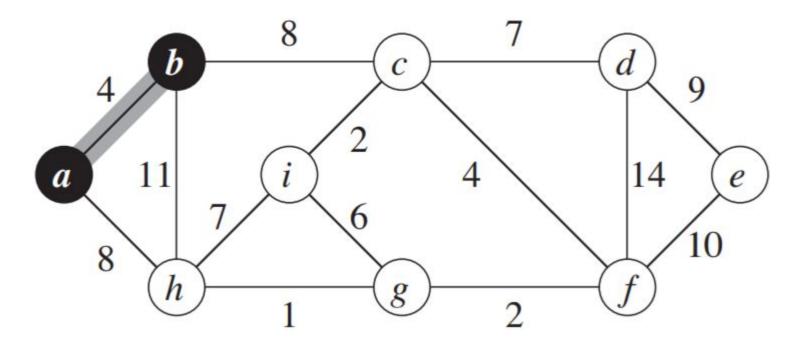
```
for each u \in G.V
         u.key = \infty
         u.\pi = NIL
    r.key = 0
    Q = G.V
    while Q \neq \emptyset
         u = \text{EXTRACT-MIN}(Q)
 8
         for each v \in G.Adj[u]
              if v \in Q and w(u, v) < v.key
10
                   \nu.\pi = u
11
                  v.key = w(u, v)
```



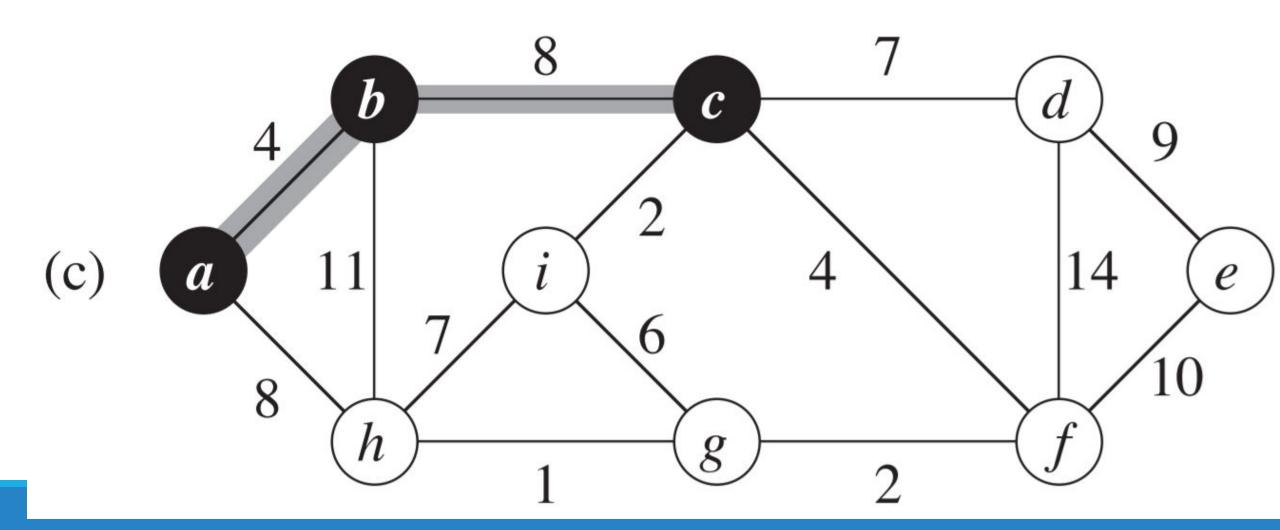
nodo	а	b	С	d	е	f	g	h	i
key	0	∞/4	∞/8	∞	∞	∞	∞	∞/8	∞
padre	-	-/a	-/b	-	-	-	-	-/a	-

MST-PRIM(G, w, r)

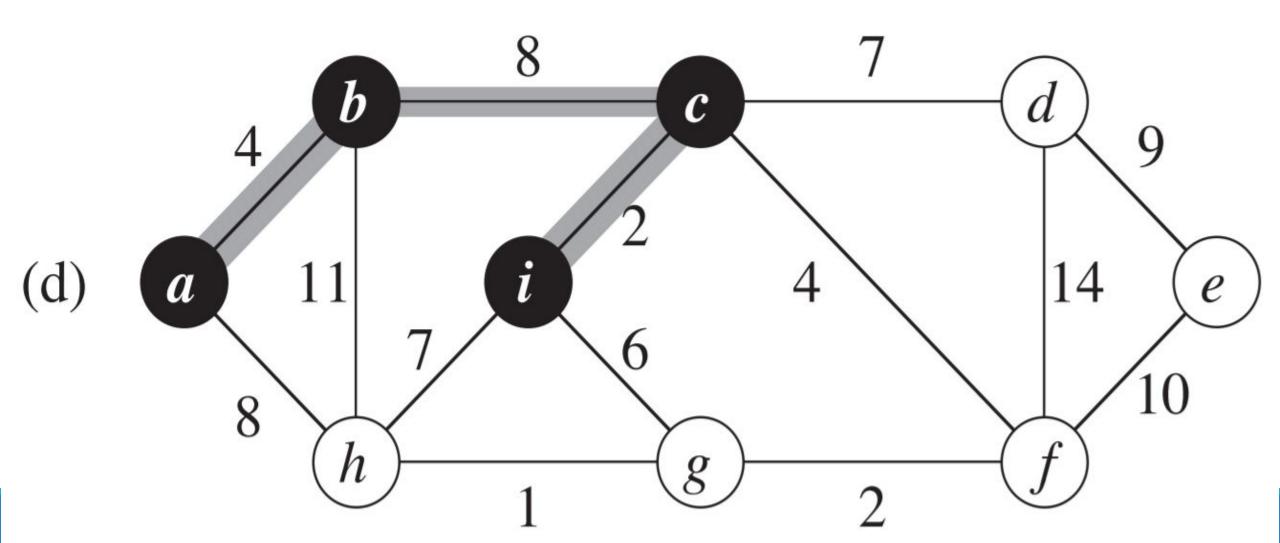
```
for each u \in G.V
         u.key = \infty
         u.\pi = NIL
    r.key = 0
     Q = G.V
    while Q \neq \emptyset
         u = \text{EXTRACT-MIN}(Q)
 8
         for each v \in G.Adj[u]
              if v \in Q and w(u, v) < v. key
10
                   \nu.\pi = u
11
                  v.key = w(u, v)
```



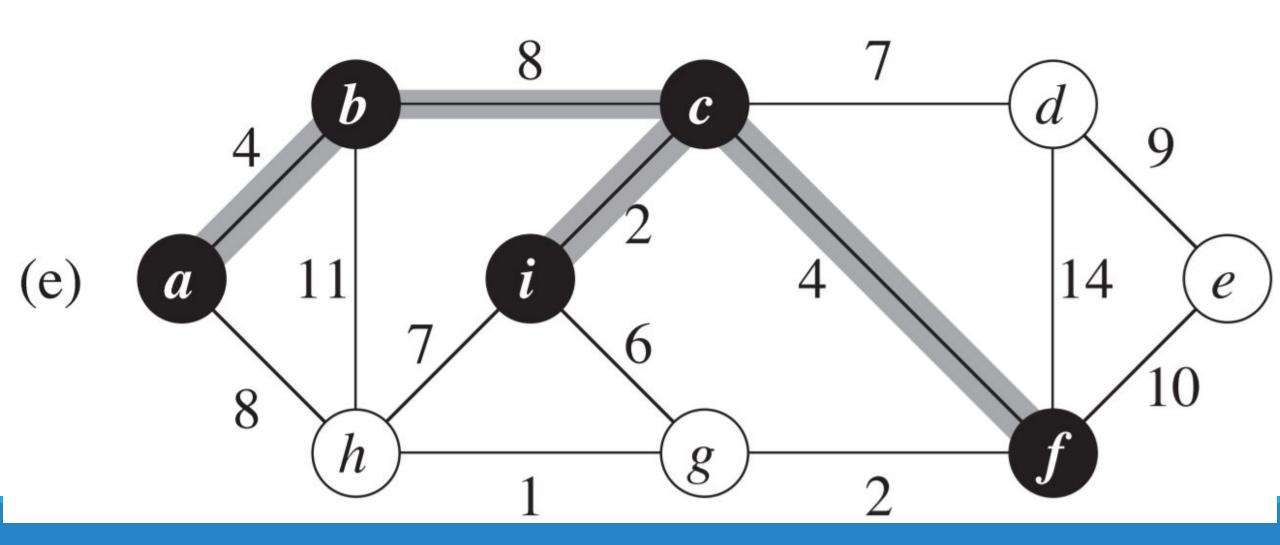
nodo	а	b	С	d	е	f	g	h	i
key	0	∞/4	∞/8	∞/7	∞	∞/4	∞	∞/8	∞/2
padre	-	-/a	-/b	-/c	-	-/c	-	-/a	-/c



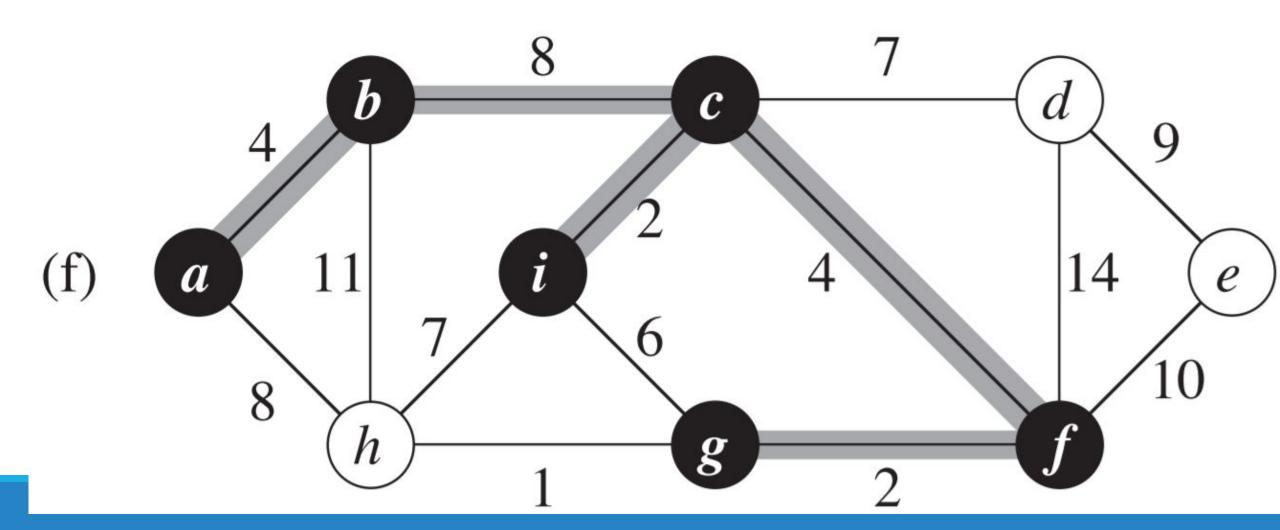
nodo	а	b	С	d	е	f	g	h	i
key	0	∞/4	∞/8	∞/7	∞	∞/4	∞/6	∞/8/7	∞/2
padre	-	-/a	-/b	-/c	-	-/c	-/i	-/a/i	-/c



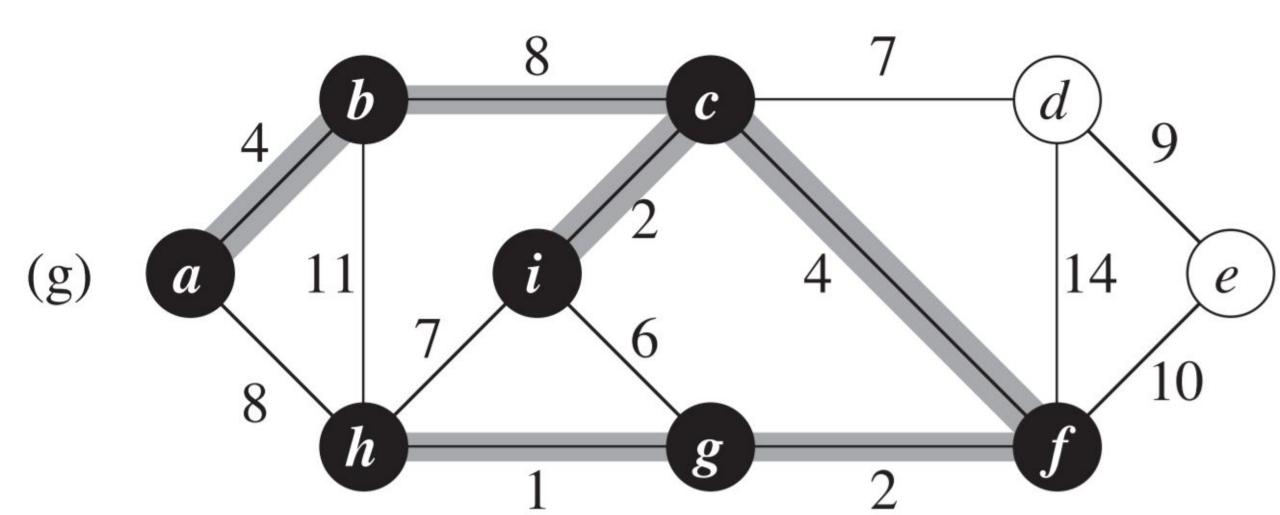
nodo	а	b	С	d	е	f	g	h	i
key	0	∞/4	∞/8	∞/7	∞/10	∞/4	∞/6/2	∞/8/7	∞/2
padre	-	-/a	-/b	-/c	-/f	-/c	-/i/f	-/a/i	-/c



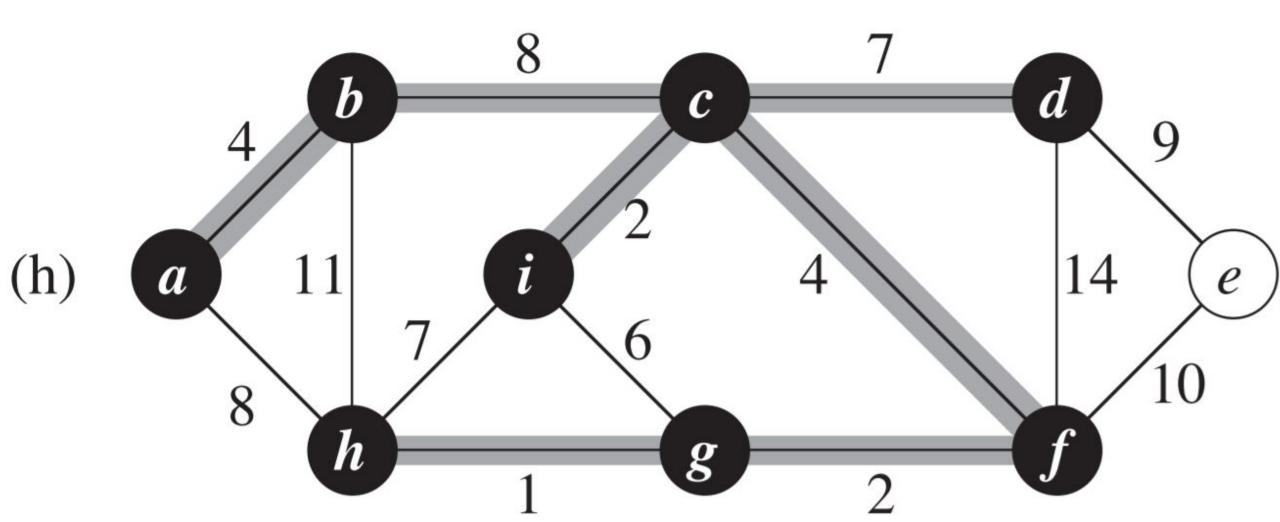
nodo	а	b	С	d	е	f	g	h	i
key	0	∞/4	∞/8	∞/7	∞/10	∞/4	∞/6/2	∞/8/7/1	∞/2
padre	-	-/a	-/b	-/c	-/f	-/c	-/i/f	-/a/i/g	-/c



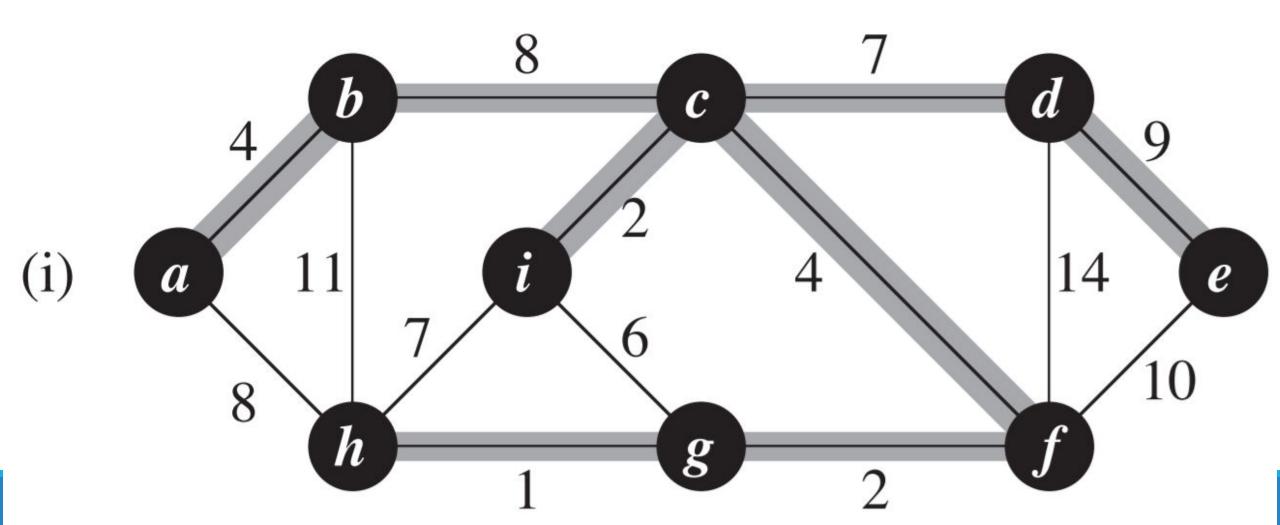
nodo	а	b	С	d	е	f	g	h	i
key	0	∞/4	∞/8	∞/7	∞/10	∞/4	∞/6/2	∞/8/7/1	∞/2
padre	-	-/a	-/b	-/c	-/f	-/c	-/i/f	-/a/i/g	-/c

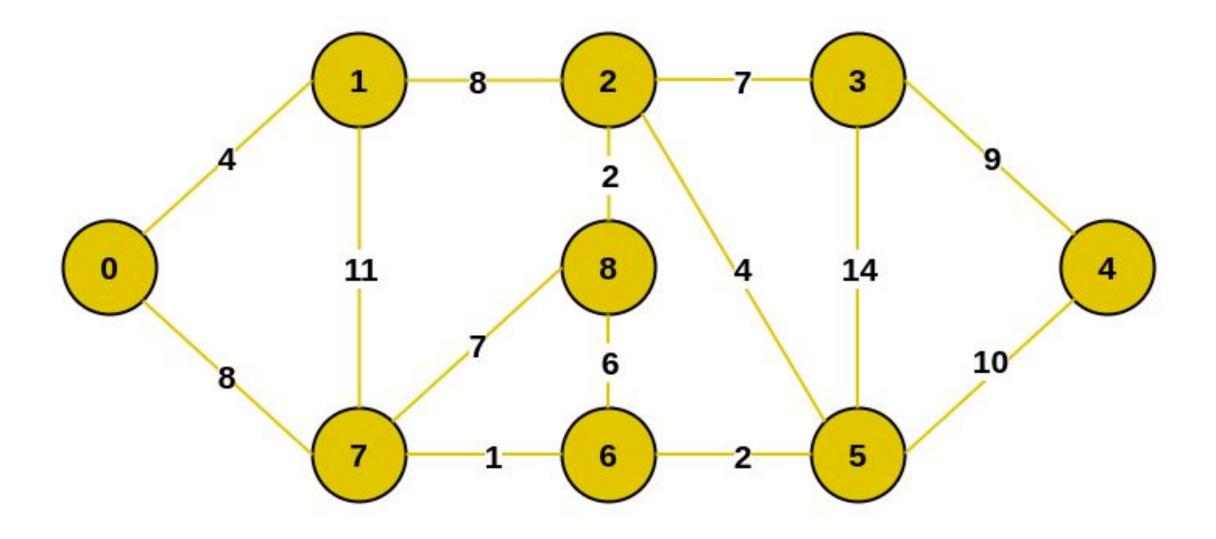


nodo	а	b	С	d	е	f	g	h	i
key	0	∞/4	∞/8	∞/7	∞/10/9	∞/4	∞/6/2	∞/8/7/1	∞/2
padre	-	-/a	-/b	-/c	-/f/d	-/c	-/i/f	-/a/i/g	-/c

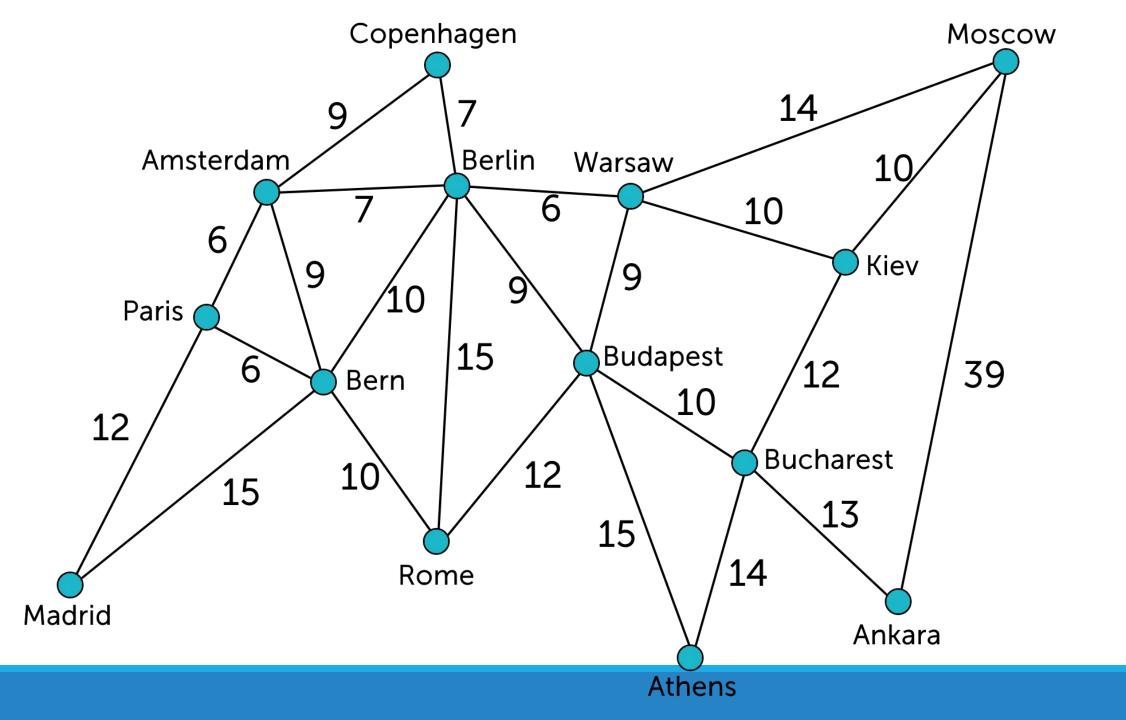


nodo	а	b	С	d	е	f	g	h	i
key	0	∞/4	∞/8	∞/7	∞/10/9	∞/4	∞/6/2	∞/8/7/1	∞/2
padre	-	-/a	-/b	-/c	-/f/d	-/c	-/i/f	-/a/i/g	-/c





Example of a Graph



Camino más corto

VARIANTES

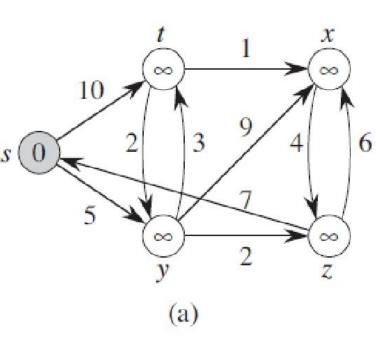
- Problema del camino más corto desde un nodo origen: dado un grafo G=(V,E), se debe encontrar el camino más corto desde un vértice de origen dado s ∈ V a cada vértice v ∈ V
- Problema del camino más corto hacia un nodo destino: Encontrar el camino más corto desde un vértice t hacia cada vértice v

VARIANTES

- •Problema del camino más corto entre un par de nodos: Encontrar le camino más corto desde un vértice t hacia un vértice dado v
- •Problema del camino más corto entre todos los pares de nodos: Encontrar el camino más corto para cada par de nodos u y v del grafo

- ▶ DIJKSTRA(G, w, s)
 - ► I INITIALIZE-SINGLE-SOURCE(G, s)
 - \triangleright 2 S \leftarrow Ø
 - ▶ 3 Q \leftarrow V[G]
 - ▶ 4 while Q ≠ Ø
 - ▶ 5 **do** $u \leftarrow \text{EXTRACT-MIN}(Q)$
 - 6 S ← S ∪{u}
 - ▶ 7 **for** each vertex $v \in Adj[u]$
 - \blacktriangleright 8 **do** RELAX(u, v, w)

- ► INITIALIZE-SINGLE-SOURCE(G, s)
 - ▶ I **for** each vertex $v \in V[G]$
 - ▶ 2 **do** $d[v] \leftarrow \infty$
 - ▶ 3 $\pi[v] \leftarrow NIL$
 - ▶ $4 d[s] \leftarrow 0$
- ▶ RELAX(*u*, *v*, *w*)
 - $ightharpoonup 1 ext{ if } d[v] > d[u] + w(u, v)$
 - ▶ 2 then $d[v] \leftarrow d[u] + w(u, v)$
 - ▶ 3 $\pi[v] \leftarrow u$



nodo	s	t	x	у	z
key	0	inf	inf	inf	nf
padre	-	-	-	-	-

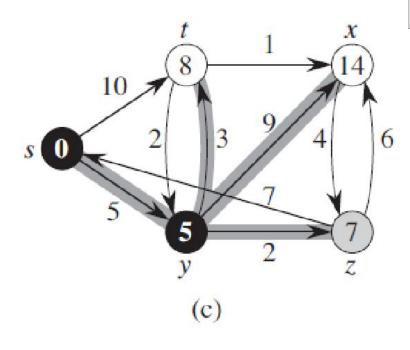
- ▶ RELAX(*u*, *v*, *w*)
 - Iif d[v] > d[u] + w(u, v)
 - ▶ 2 then $d[v] \leftarrow d[u] + w(u, v)$
 - ▶ 3 $\pi[v] \leftarrow u$

$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	6
(b)	

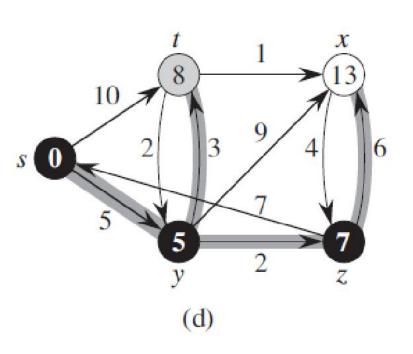
nodo	S	t	x	у	z
key	0	inf/10	nf	inf/5	nf
padre	-	-/s	-	-/s	-

- ▶ RELAX(*u*, *v*, *w*)
 - Iif d[v] > d[u] + w(u, v)
 - ▶ 2 then $d[v] \leftarrow d[u] + w(u, v)$
 - ▶ 3 $\pi[v] \leftarrow u$

nodo	s	t	x	у	z
key	0	inf/10/8	inf/14	inf/5	inf/7
padre	-	-/s/y	-/y	-/s	-/y

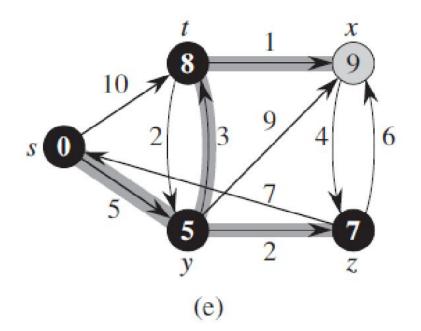


- ▶ RELAX(*u*, *v*, *w*)
 - \vdash I if d[v] > d[u] + w(u, v)
 - ▶ 2 then $d[v] \leftarrow d[u] + w(u, v)$
 - ▶ 3 $\pi[v] \leftarrow u$

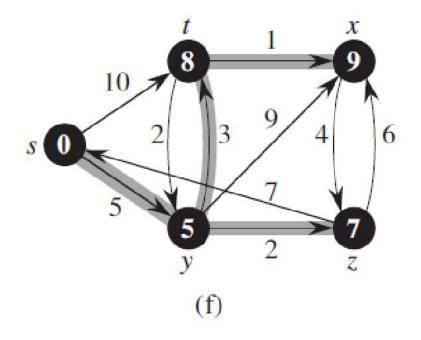


nodo	s	t	x	у	Z
key	0	inf/10/8	inf/14/13	inf/5	inf/7
padre	-	-/s/y	-/y/z	-/s	-/y

- ▶ RELAX(*u*, *v*, *w*)
 - I if d[v] > d[u] + w(u, v)
 - ▶ 2 then $d[v] \leftarrow d[u] + w(u, v)$
 - ▶ 3 $\pi[v] \leftarrow u$

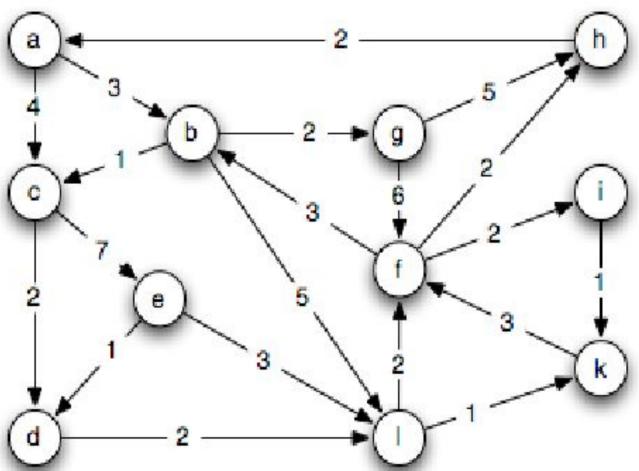


nodo	s	t	X	у	Z
key	0	inf/10/8	inf/14/13/9	inf/5	inf/7
padre	-	-/s/y	-/y/z/t	-/s	-/y

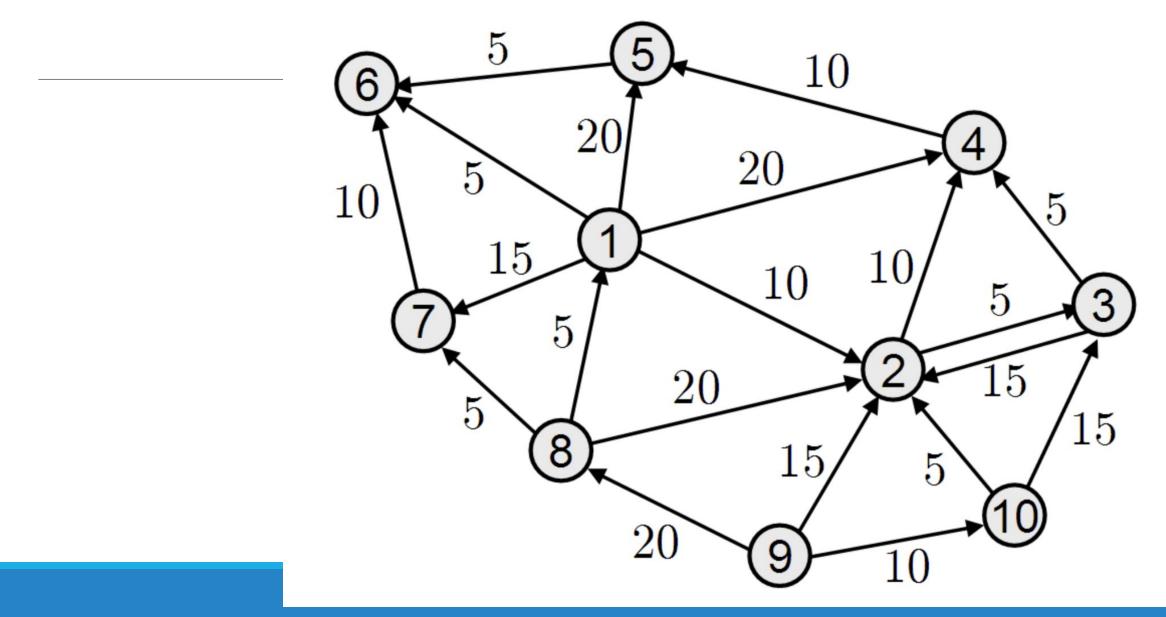


nodo	S	t	х	у	Z
key	0	inf/10/8	inf/14/13/9	inf/5	inf/7
padre	-	-/s/y	-/y/z/t	-/s	-/y

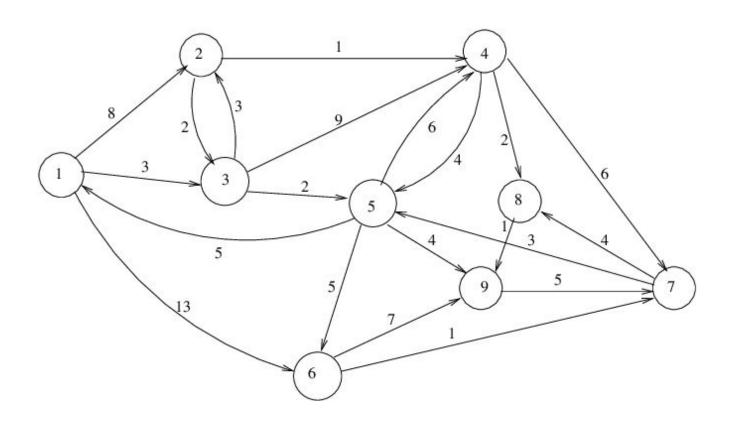
1a, 1g, 1b



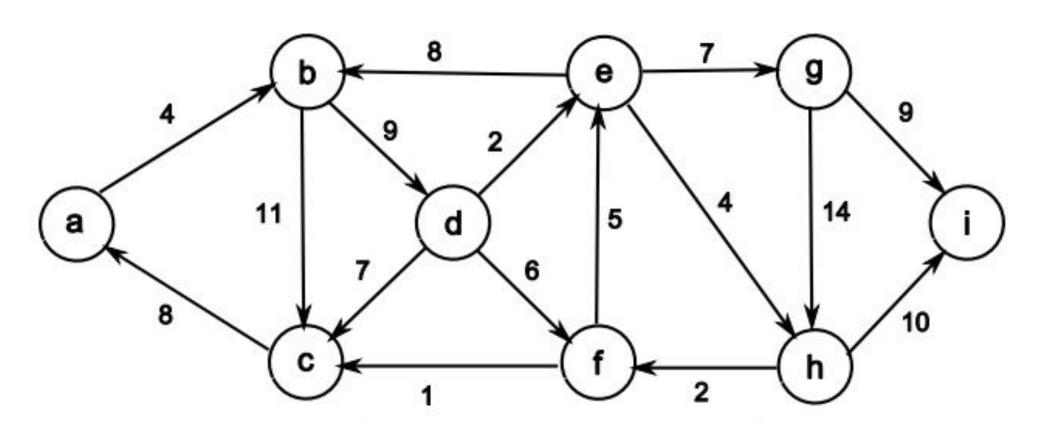
2-1,2-8,2-3

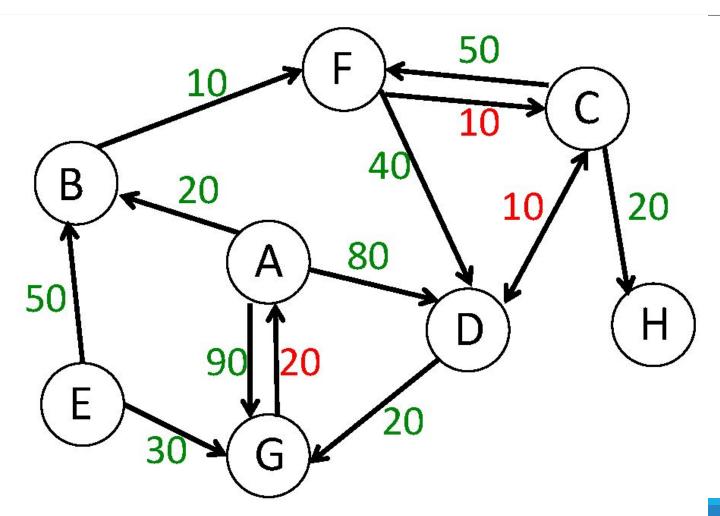


3-1, 3-5



4a, 4e





6a,

