SCALER

Array - prefix sum

TABLE OF CONTENTS

Practical Scenrio

You want to develop a feature for a banking app that allow users to quickly accus their stock portfolio's profit or loss over specific periods.

arr → profit or loss from a particular stck over a period of days.

arr = [-5, 10, 20, 40, 50, -10, 80, -90, -20, -10]
indices: 0 1 2 3 4 5 6 7 8 9

start	end	ans
0	9	85
1	4	120
0	0	-5
2	7	90

< *Question* > : Given an array of N integers and Q queries. For each query calculate the sum of elements in the range - [L , R]

Note : L and R are indices such that L ≤ R.

$1 \le N, Q \le 10^5$

arr[10] → [-3 6 2 4 5 2 8 -9 3 1]

0 1 2 3 4 5 6 7 8 9

Queries

L	R	
4	8	→ 9
3	7	→ 10
1	3	→ 12
0	4	→ 14
7	7	→ -9

Brute force Approach →

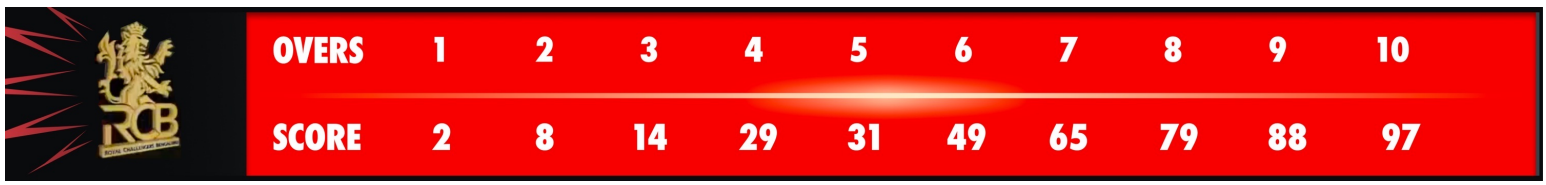
For every query,
iterate from L to R and find the sum.

```
for (i = 0; i < Q ; i++) {
    L → Q[i][0]
    R → Q[i][1]
    sum = 0;
    for( j = L; j ≤ R ; j++){
        sum += arr[j];
    }
    print(sum);
}
```

T.C → O(Q*N)
S.C → O(1)

→ T.L.E

- **Given Royal Challengers Bengaluru's cricket scores for first 10 overs of batting.**



OVERS	1	2	3	4	5	6	7	8	9	10
SCORE	2	8	14	29	31	49	65	79	88	97

⇓

Cumilative Sum / prefix Sum

- Runs scored in 7th over = 65 - 49 = 16
- Runs scored from 6 - 10th over = 97 - 31 = 66
- Runs scored in 10th over = 97 - 88 = 9
- Runs scored from 3 - 6th over = 49 - 8 = 41
- Runs scored from 4 - 9th over = 88 - 14 = 74
- Runs scored in lth - rth over = runs[r] - runs[l-1]

How to create psum()

arr[10] → [-3 6 2 4 5 2 8 -9 3 1]

0 1 2 3 4 5 6 7 8 9

psum[10] = [-3, 3, 5, 9, 14, 16, 24, 15, 18, 19]

0 1 2 3 4 5 6 7 8 9

Queries:

4	8
3	7
1	3
0	4
-1	7

arr[6] → [10 32 6 12 20 1]

0 1 2 3 4 5

psum[10] = [10, 42, 48, 60, 80, 81]

0 → psum[8] - psum[4-1] ⇒ 9

1 → psum[7] - psum[3-1] ⇒ 10

2 → psum[3] - psum[1-1] ⇒ 12

3 → psum[4] ⇒ 14

4 → psum[7] - psum[7-1] ⇒ -9

code to create psum[]

```
int psum[N];
psum[0] = arr[0];
for( i = 1; i < N; i++){
    psum[i] = psum[i-1] + arr[i];
}
```

</> Code

// 1. Create psum[]

```
int psum[N];
psum[0] = arr[0];
for( i = 1; i < N; i++){
    psum[i] = psum[i-1] + arr[i];
}
```

// 2. Answer the queries

$1 \le N, Q \le 10^5$

```
for( i = 0; i < Q; i++){
    L -> Q[i][0];
    R -> Q[i][1];

    if( L == 0) {
        print( psum[R];
    }
    else {
        print( psum[R] - psum[L-1]);
    }
}
```

T.C → O(N+Q)

S.C → O(N)

Modification of same array into psum

arr[10] → [-3 6 2 4 5 2 8 -9 3 1]

0 1 2 3 4 5 6 7 8 9

arr[10] → [-3, 3, 5, 9, 14, 16, 24, 15, 18, 19]

```
for( i = 1; i < N; i++){
{    arr[i] = arr[i] + arr[i-1];
}
```

S.C → O(1)

< *Question* > : Given an arr[N] and Q queries with start(s) and end(e) index. For every query print sum of all even indexed elements from s to e.

arr[ ] →

2	3	1	6	4	5
0	1	2	3	4	5

Queries

s	e	
1	3	→ 1
2	5	→ 5
0	4	→ 7
3	3	→ 0

[BF Idea] - For every query, iterate from s to e.

[T.C → O(N x Q)
S.C → O(1)]

[Idea] - **Using prefix sum** of even. indexed elements only.

arr[ ] →

2	3	1	6	4	5
0	1	2	3	4	5

psum[10] →

2	2	3	3	7	7
0	1	2	3	4	5

Queries

s	e	
1	3	→ 1
2	5	→ 5
0	4	→ 7
3	3	→ 0

psum[e] - psum[s-1]

< / > Code

```
psum[N];
psum[0] = arr[0];
for( i=1; i<N; i++){
    if( i%2 == 0){
        psum[i] = psum[i-1] + arr[i];
    }
    else{
        psum[i] = psum[i-1];
    }
}
// → todo. (similar to previous ques)
```

N → 100

Q → 10^5

T.C → O(N+Q)

S.C → O(N)

modify array

↓

O(1)

< Question > : For all the queries, find the sum of odd indexed elements from s to e.

0 1 2 3 4 5

⇒ todo.

Special Index

< *Question* > : Given an arr[N], count the number of special indices in the array.

Special Index : Index after removing which,

Sum of even indexed elements = sum of odd indexed elements.

$1 \le N \le 10^5$

arr - [4, 3, 2, 7, 6, -2]
0 1 2 3 4 5

ans = 2.

remove idx = 0 → [4 3 2 7 6 -2] → [3 2 7 6 -2]
0 1 2 3 4 5 → 0 1 2 3 4 ✓

remove idx = 1 → [4 3 2 7 6 -2] → [4 2 7 6 -2]
0 1 2 3 4 5 → 0 1 2 3 4 ✗

remove idx = 2 → [4 3 2 7 6 -2] → [4 3 7 6 -2]
0 1 2 3 4 5 → 0 1 2 3 4 ✓

remove idx = 3 → [4 3 2 7 6 -2] → [4 3 2 6 -2]
0 1 2 3 4 5 → 0 1 2 3 4 ✗

remove idx = 4 → [4 3 2 7 6 -2] → [4 3 2 7 -2]
0 1 2 3 4 5 → 0 1 2 3 4 ✗

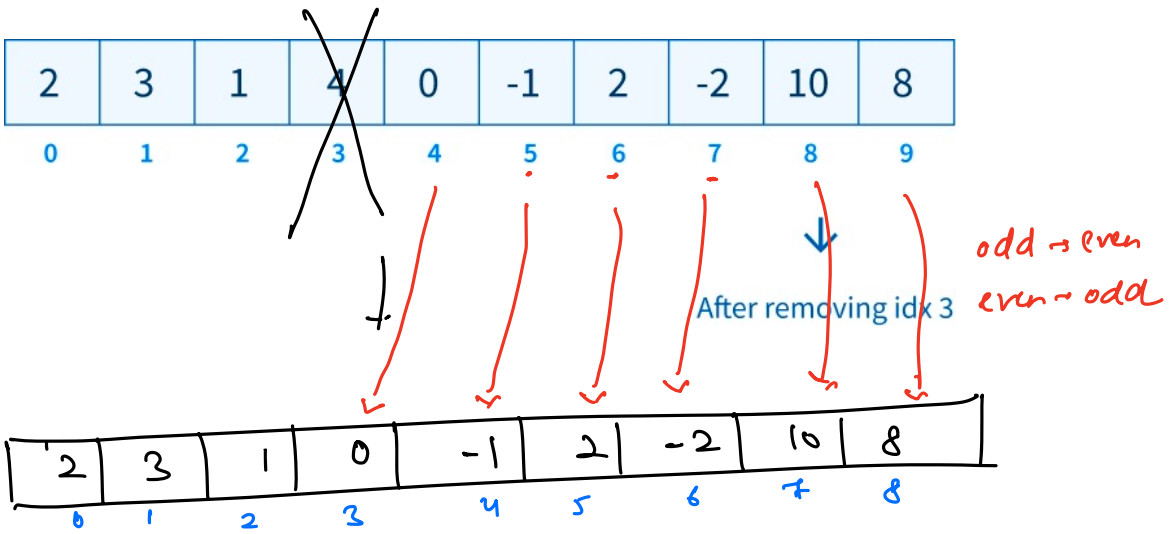
remove idx = 5 → [4 3 2 7 6 -2] → [4 3 2 7 6]
0 1 2 3 4 5 → 0 1 2 3 4 ✗

Quiz

1.

Sum of odd indexed elements → !!

2.



$S_{o[0-9]} = S_{o[0-2]} + S_{e[4-9]}$

$(2) + (12) = 15$

$S_{e[0-9]}$ after removing 3rd index → $S_{e[0-2]} + S_{o[4-9]}$

→ 8.

[Idea - 2]

arr = [2, 3, 1, 4, 0, -1, 2, -2, 10, 8]
0 1 2 3 4 5 6 7 8 9

Delete 4th index

↳ [2, 3, 1, 4, -1, 2, -2, 10, 8]
0 1 2 3 4 5 6 7 8

$S_o[0-9]$ after deleting 4th index → $S_o[0-3] + S_e[5-9]$

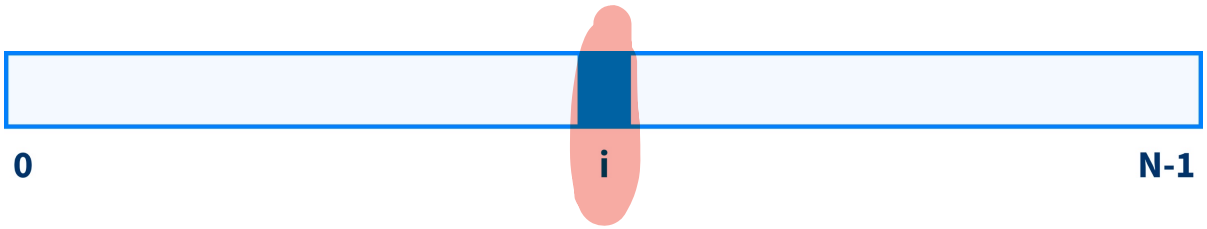
→ 19.

$S_e[0-9]$ after deleting 4th index → $S_e[0-3] + S_o[5-9]$

→ 8

pf Even[ ]
pf odd[ ]

Assumption, we already have pfeven[ ] and pfodd[ ].



sum of even indexed elements after removing ith indexed element →

$$Se[0, i-1] + So[i+1, N-1]$$

$$\downarrow \qquad \downarrow$$

$$pfEven[i-1] + pfodd[N-1] - pfodd[i]$$

sum of odd indexed elements after removing ith indexed element →

$$So[0, i-1] + Se[i+1, N-1]$$

$$\downarrow \qquad \downarrow$$

$$pfodd[i-1] + pfEven[N-1] - pfEven[i]$$

< / > Code

```
// Create pfEven[] and pfOdd[]

pfEven[N], pfOdd[N];

pfEven[0] = arr[0];

for(i=1; i<N; i++){
    if(i%2==0){
        pfEven[i] = pfEven[i-1] + arr[i];
    }
    else{
        pfEven[i] = pfEven[i-1];
    }
}

pfOdd[0]=0;

for(i=1; i<N; i++){
    if(i%2==1){
        pfOdd[i] = pfOdd[i-1] + arr[i];
    }
    else{
        pfOdd[i] = pfOdd[i-1];
    }
}
```

```
int count = 0;

if( pfodd[N-1] - pfodd[0] == pfEven[N-1] - pfEven[0]){
        count++;
}

for( i=1; i<N; i++){
        Se = pfEven[i-1] + pfodd[N-1] - pfodd[i];
        So = pfodd[i-1] + pfEven[N-1] - pfEven[i];
        if( Se == So) { count++ }
}

return count;
```

T.C → O(N)
S.C → O(N)

→ Carry forward

→ Basic of Subarrays.

→ optimisation → Sum of elements is req. again & again.

arr = [7, 2, 8, 3, 8, 6, 8]

↓

1 + 6 + 0 + 5 + 0 + 2 + 0 = 14

① find max elem

② keep on taking diff. with max elem.