

1. void fun(int N){
$$\rightarrow$$
 Space Complexity
int x = N;
int y = 2*x;
long z = x+y;
}

Space Complexity

 $x = 4B$
 $y = 4B \implies 2 \text{ space} = 16B$

2 = 8B

2 = 8B

2 = 8B

2 = 8B

3 O(1)

2. func(int N){ // 4 bytes
 int am[10]; //40 bytes
 int x; // 4 bytes

$$y = 4B$$
 $\Rightarrow \le space$

n = 4B

int x = N;
int y = x*x;
long z = y+y;
int[] arr = new int[N];
long[][] b = new long[N][N];

$$b = 8 \times N \times N B = 8 \times N^2 B$$

$$\leq 5 \text{ page} = 8 \times N^2 + 4 \times N + 16 \implies OCN^2$$

}

3.

void fun(int N){

→ We need to calculate how much int maxArr(int arr[], int N){ int ans = arr[0]; EXTRA space for(i=0; i<N; i++){ - Hence the space taken by input ans = Max(ans, arr[i]); parometer voriables is not accounted for

parameter variables is most account for
$$ons = 4R \implies O(1)$$

Rotate An Array

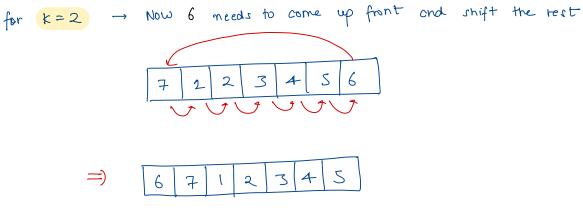
Cliven an array of size N & an integer K. Rotate arrill by K.

$$\underline{1 \le N \le 10^6}$$

$$\underline{0 \le K \le 10^9}$$

N = 7

rest of the array meeds to this one space towards the right
$$\frac{1}{2} \frac{3}{4} \frac{4}{5} \frac{6}{6} \frac{7}{4}$$



You get the idea, so in the Brute Force approach,

let's devise on algorithm which achieves the some

Step 1 -> Sove the value of last element

Step 2 -> Shift the values of the initial element 1 place to right

Step 3 -> Sot the value of the first element with the saved

value of the last element

int length = arr.length;

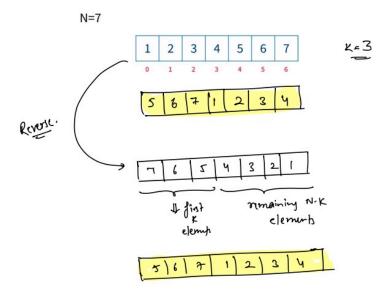
// * Rotating array 'K' times for (int i = 0; i < K; i++) {

int last = arr[length - 1]; for (int $j = length - 2; j >= 0; j--) {$ $arr[j + 1] = arr[j]; \longrightarrow Storing$

Soving current index value ahead arr[0] = last;

into the index ofter

Optimisation



- 6 K= K1.N
- 1) kererse the whole array
- @ Rivine fist k elements.
- 3 Reverse remaining al-k elements