

Sorting

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Notes

Java

```
static void change(int a) {
    a = 50;
}

public static void main(String args[]) {
    ✓ int a = 10;
    ✓ change(a);
    ✓ System.out.println(a);
}
```

Python

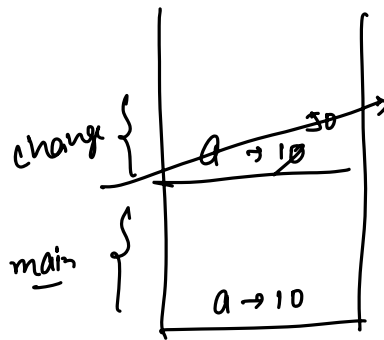
```
def change(a):
    a = 50

def main():
    a = 10
    change(a)
    print(a)

if __name__ == "__main__":
    main()
```

o/p. ~ 10

Stack



Java

```
static int[] fun(int[] a) {
    a = new int[2];
    a[0] = 50; a[1] = 60;
    return a;
}

public static void main(String args[]) {
    int[] a = {10, 20, 30};
    a = fun(a);
    System.out.println(a[0]);
}
```

Python

```
def fun(a):
    a = [0, 0]
    a[0] = 50
    a[1] = 60
    return a

def main():
    a = [10, 20, 30]
    a = fun(a)
    print(a[0])

if __name__ == "__main__":
    main()
```

o/p. ~ 50



Sorting

↳ arrangement of data in any particular order on the basis of some parameter.

arr[] → [2, 3, 9, 12, 17, 19]

parameter → value

arr[] → [19, 6, 5, 2, -1, -15]

parameter → value

arr[] → [1, 13, 9, 6, 12]

parameter → factors count.

factors →
↓ ↓ ↓ ↓ ↓
1 2 3 4 6

Why sorting?

- to make search faster.
- analyze & represent the data

Java. → Arrays.sort(arr);

C++ → sort(arr.begin(), arr.end())

Python → arr.sort()

} check the exact syntax.



T.C → $O(N \log N)$

S.C → depends on inbuilt sorting algo.

**Question** (Elements Removal)

Given N elements, at every step remove an array element.

Cost to remove an element = Sum of array of elements present in an array

Find minimum cost to remove all elements.

NOTE : First add the cost of removal and then remove it.

arr - [2, 1, 4]

remove 2 \Rightarrow 7

remove 1 \Rightarrow 5

remove 4 \Rightarrow 4

total cost \Rightarrow 16

Arr (1 \rightarrow [4, 1])

remove 6 \rightarrow 11

remove 4 \rightarrow 5

remove 1 \rightarrow 1

total cost \rightarrow 17

[2, 1, 4]

remove 4 \rightarrow 7

remove 2 \rightarrow 3

remove 1 \rightarrow 1

total cost \Rightarrow 11

Arr (1 \rightarrow [2, 1, -3])

remove 5 \rightarrow 6

remove 3 \rightarrow 1

remove 1 \rightarrow -2

remove -3 \rightarrow -3

total cost \rightarrow 2

arr[7] \rightarrow [~~a~~ ~~b~~ ~~c~~ d]

Remove $a \rightarrow a + b + c + d$

Remove $b \rightarrow b + c + d$

Remove $c \rightarrow$ $c + d$

Remove $d \rightarrow$ d

total cost $\rightarrow a + 2b + 3c + 4d$

$$\underline{a > b > c > d.}$$

to get the minimum cost \Rightarrow largest element should appear minimum no. of times.

Idea. \rightarrow Sort the arr. in decreasing order. Eg \rightarrow $[6, 4, 1]$
0 1 2

#code. →

sort(arr, desc.)

$$\cos t = 0;$$

```
for ( i = 0; i < N; i++) {
```

for (i = 0; i < n; i++)
{
 cost += (arr[i] * (i+1));
}

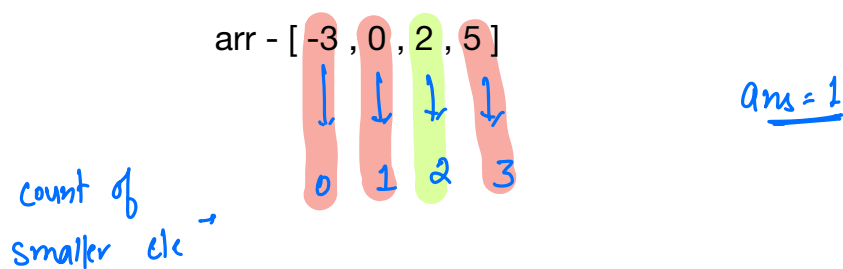
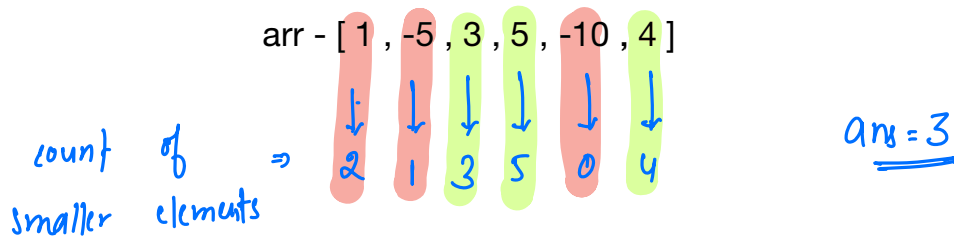
```
return cost;
```

$\left[\begin{array}{l} \text{T.C} \rightarrow O(N \log N) \\ \text{S.C} \rightarrow \text{depends on inbuilt sorting algo} \end{array} \right]$

**Question** (Noble Integers) { Distinct data }

Given N array elements, calculate number of noble integers.

An element ele in $arr[]$ is said to be noble if { count of smaller elements = ele itself }



Bf idea. -
 For every element, iterate and find the count of elements smaller than current element.

```

ans = 0
for (i = 0; i < n; i++) {
    count-smaller = 0;
    for (j = 0; j < n; j++) {
        if (arr[j] < arr[i]) { count-smaller++; }
    }
    if (count-smaller == arr[i]) { ans++; }
}
return ans;
  
```

T.C $\rightarrow O(N^2)$
 S.C $\rightarrow O(1)$



optimisation

if array is sorted in ascending order.

Count of smaller elements for $arr[i] \rightarrow \underline{i}$

#code:-

```
sort(arr, asc.)
```

```
ans = 0;
```

```
for (i = 0; i < n; i++) {
```

```
    if (arr[i] == i) { ans++; }
```

```
}
```

```
return ans;
```

T.C $\Rightarrow O(N \log N)$
S.C \Rightarrow depends on sorting algo



Question (Noble Integers) : { Data can repeat }

arr = [-10, 1, 1, 3, 100]

0 1 2 3 4

ans = 3

arr = [-10, 1, 1, 2, 4, 4, 4, 8, 10]

0 1 2 3 4 5 6 7 8

ans = 5

arr = [-3, 0, 2, 2, 5, 5, 5, 5, 8, 8, 10, 10, 10, 14]

0 1 2 3 4 5 6 7 8 9 10 11 12 13

ans = 7

count-smaller = 0 1 2 4 8 10 13

</> Code

```
Arrays.sort(arr, asc.);
```

```
count-smaller = 0, ans = 0;
```

```
if (arr[0] == 0) { ans++; }
```

```
for (i = 1; i < n; i++) {
```

```
    if (arr[i] != arr[i-1]) {
```

```
        count-smaller = i;
```

```
        if (arr[i] == count-smaller) { ans++; }
```

```
}
```

```
return ans;
```

[T.C → $O(n \log n)$
S.C → depends on sorting algo]



Selection Sort

idea : Select the minimum element and send that elements to correct position by swapping.

arr[] = ¹5, ²6, ³4, ⁴1, ⁵2

_{0 1 2 3 4}

min-ele = 1

min-idx = 3

↓

²1, ³6, ⁴4, ⁵5, ⁶2

_{0 1 2 3 4}

min-ele = 2

min-idx = 4

↓

[¹1, ²2] [³4, ⁴5, ⁵6]

_{0 1 2 3 4}

min-ele = 4

min-idx = 2

↓

[¹1, ²2, ³4] [⁴5, ⁵6]

_{0 1 2 3 4}

min-ele = 5

min-idx = 3

↓

[¹1, ²2, ³4, ⁴5] [⁵6]

_{0 1 2 3 4}



</> Code

```
for( i = 0; i < N-1 ; i++) {
```

```
    min-element = arr[i];
```

```
    min-idx = i;
```

```
    for( j = i+1; j < N; j++) {
```

```
        if( arr[j] < min-element) {
```

```
            min-element = arr[j];
```

```
            min-idx = j;
```

```
        }
```

```
    swap( arr[i] with arr[min-idx]);
```

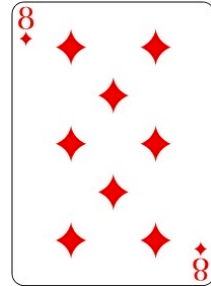
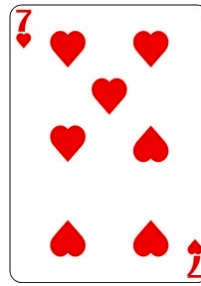
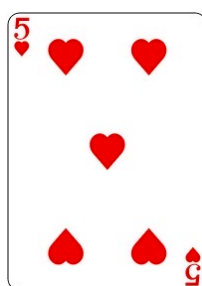
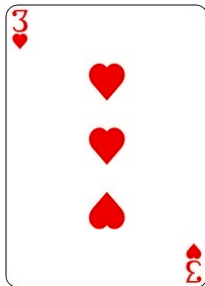
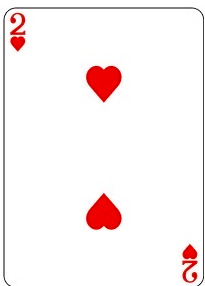
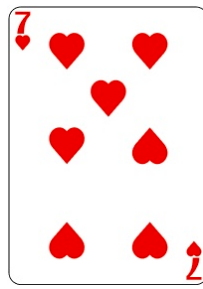
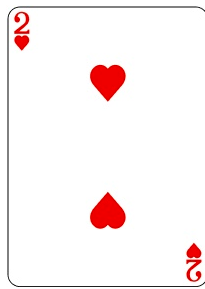
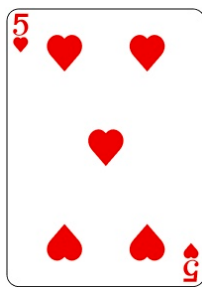
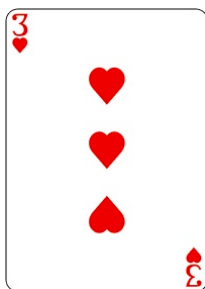
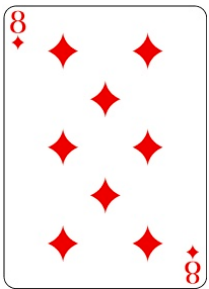
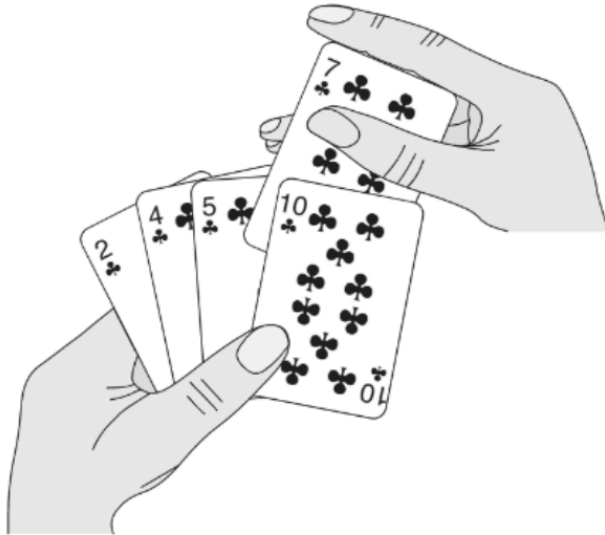
```
}
```

selecting the
min element
from idx = i to N-1

$T.C \rightarrow O(N^2)$
 $S.C \rightarrow O(1)$



Insertion Sort (Arrangement of playing cards)





arr \rightarrow $\begin{bmatrix} 3 & 8 \\ 8 & 3 & 5 & 2 & 7 \end{bmatrix}$ $\text{idx} \rightarrow 1$

0 1 2 3 4

↓

$\begin{bmatrix} 3 & 5 & 8 \\ 8 & 3 & 5 & 2 & 7 \end{bmatrix}$ $\text{idx} \rightarrow 2$

0 1 2 3 4

↓

$\begin{bmatrix} 2 & 2 & 3 & 2 & 5 & 8 \\ 8 & 3 & 5 & 2 & 7 \end{bmatrix}$ $\text{idx} \rightarrow 3$

0 1 2 3 4

↓

$\begin{bmatrix} 2 & 3 & 5 & 7 & 8 \\ 8 & 3 & 5 & 2 & 7 \end{bmatrix}$ $\text{idx} \rightarrow 4$

0 1 2 3 4

#code:-

```
for ( i = 1; i < N; i++) {
    for ( j = i - 1; j >= 0; j--) {
        if ( arr[j] > arr[j+1]) {
            swap ( arr[j] with arr[j+1]);
        }
        else {
            break;
        }
    }
}
```

$\begin{bmatrix} T.C \rightarrow O(N^2) \\ S.C \rightarrow O(1) \end{bmatrix}$

	3	25	77	12
1	8	7	12	3
0	1	2	3	4