

What is Recursion?

A function calling itself to solve a problem

Solving a problem using sub-problems of the same type is called Recursion

$$\begin{aligned} \text{sum}(N) &= \text{Sum of first } N \text{ natural numbers} \\ &= 1 + 2 + 3 + \dots + N-2 + N-1 + N \end{aligned}$$

$$\text{sum}(N) = \text{sum}(N-1) + N$$

sum(4)

$$10 \rightarrow 4 + \text{sum}(3)$$

$$6 \rightarrow 3 + \text{sum}(2)$$

$$3 \rightarrow 2 + \text{sum}(1) \checkmark$$

$$1 \rightarrow 1 + \text{sum}(0)$$

$$\rightarrow 0 + \text{sum}(-1)$$

### 3 steps of Recursion

Assumption: Define what your recursive function has to do and just assume it works.

Eg:  $\text{sum}(N)$  returns sum of  $N$  natural nos

Main logic: Using the function defined in step 1, break the bigger problem into smaller problems

$$\text{sum}(N) = \text{sum}(N-1) + N$$

Base condition: When do you want to stop.

$\text{if}(N == 0) \text{return } 0$

```
int sum(int N) {
```

```
    if(N == 0) return 0;
```

```
    return sum(N-1) + N;
```

```
}
```

Question: find factorial

$$N! = 1 \cdot 2 \cdot 3 \cdot 4 \cdot \dots \cdot (N-1) \cdot N$$

$$5! = 1 \cdot 2 \cdot 3 \cdot 4 \cdot 5 = 120$$

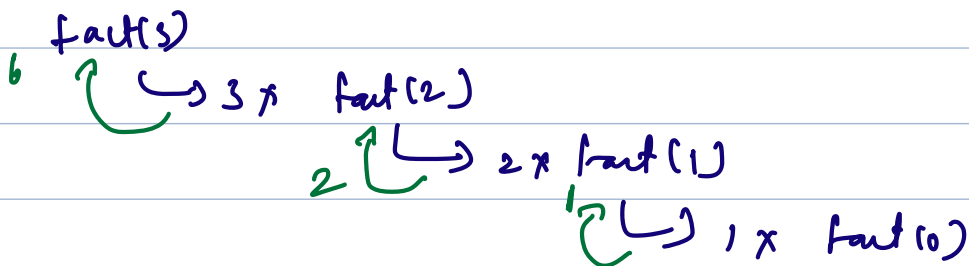
Assumption: `fact(N)` returns  $N!$

Main Logic:

$$\text{fact}(N) = \text{fact}(N-1) \times N$$

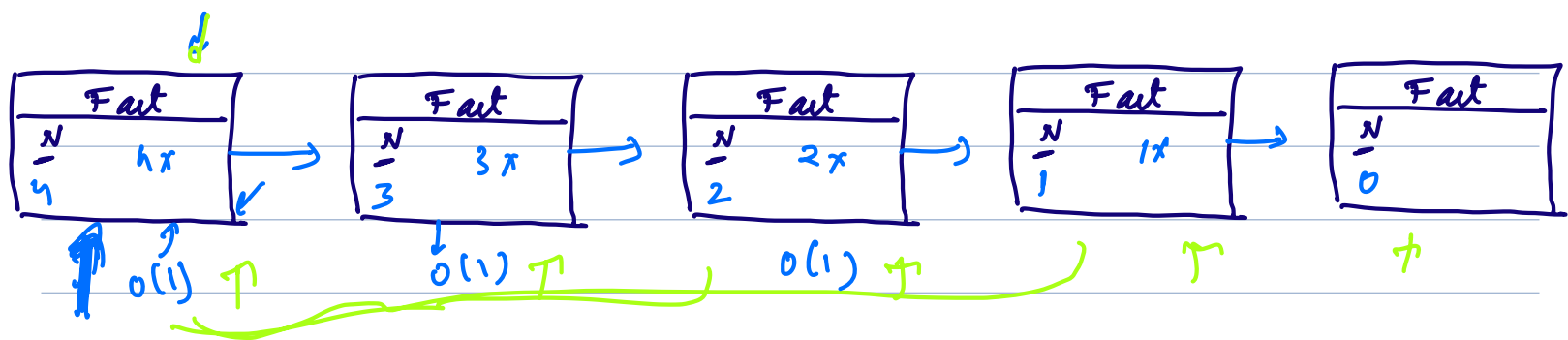
Base Condition

$$0! = 1$$



```
int fact(int N) {
    if (N == 0) return 1;
    return N * fact(N-1);
}
```

}



- Whenever a function is called, some memory is allocated on top of call stack
- After function completely execution, it is popped from the call stack

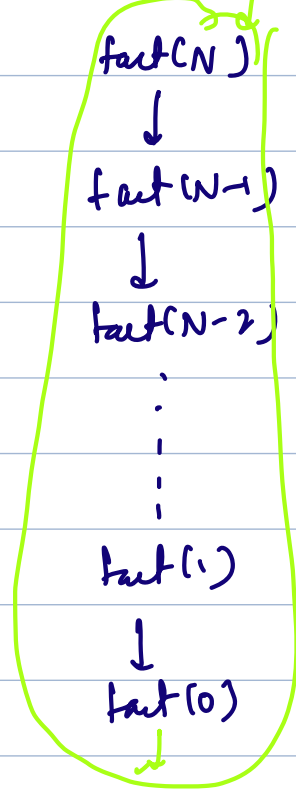
```
int fact(int N) {
```

```
1. if (N == 0) return 1; →  $O(1)$ 
```

```
2. return N * fact(N-1); →  $O(1)$ 
```

```
}
```

|                    |   |
|--------------------|---|
| <del>fact(0)</del> |   |
| <del>fact(1)</del> | $1 \times 1 = 1$                          |
| <del>fact(2)</del> | $2 \times \underline{1} = 2$              |
| <del>fact(3)</del> | $3 \times \underline{2} = 6$              |
| <del>fact(4)</del> | $4 \times \underline{6} = \underline{24}$ |



$$[0, N] = N+1$$

T.C:  $O(\underbrace{\# \text{ Function calls}}_{(N+1)} \times \text{T.C per call}) \quad \checkmark$

$$(N+1) \times O(1) \Rightarrow O(N)$$

S.C:  $O(\# \text{ map } \uparrow \text{ function calls } \times \text{S.C per active in stack function call})$

$$(N+1) \times O(1) = \underline{O(N)}$$

→ Every recursive solution can also be implemented through iterative solutions

Question: Print 1 to N in increasing order

Assumption: printInc(N) prints 1 to N in 7 order

Main Logic

```
printInc(N-1);  
print(N);
```

Base Case

```
if (N == 1) {  
    print(1);  
    return;  
}
```

OR

```
if (N == 0) {  
    return;  
}
```

```
void printInc(int N) {  
    if (N == 0) return;  
    printInc(N-1);  
    print(N);  
}
```

```
printInc(N-1);  
print(N);
```

// 1. 2. 3. ... N-1

y

## Quesn:

### Problem

**Whirlpool** wants to design a timer for their **washing machines**. This feature is a simple countdown timer. When a user sets a time, for example, 10 minutes, the washing machine needs to show each minute passing, counting down until it reaches 0.

Your task is to write a program that takes an integer **A** (the time in minutes set by the user) and then prints out each minute as it counts down to 0. The requirement is that after a user set a timer for the washing machine for some time say **A**, the washing machine should display each minute after that decremented one by one till the time becomes 0.

### Simplified Problem statement

Given **N**, print all numbers from **1 to N** in decreasing order.

5:      5      4      3      2      1

```
void printDec (int N) {  
    if (N == 0) return;
```

```
    print(N);  
    printDec (N-1);  
}
```

y

T.C:  $N \times O(1) = O(N)$

S.C:  $N \times O(1) = O(N)$

8:23

Question: Find  $N^{\text{th}}$  number in fibonacci series

|   |   |   |   |   |   |   |    |    |    |    |     |
|---|---|---|---|---|---|---|----|----|----|----|-----|
| 0 | 1 | 1 | 2 | 3 | 5 | 8 | 13 | 21 | 34 | 55 |     |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7  | 8  | 9  | 10 | ... |
| - | - | ↑ | ↑ | ↑ | ↑ | ↑ | ↑  |    |    |    |     |

$$\text{fib}(N) = \text{fib}(N-1) + \text{fib}(N-2)$$

Assumption:  $\text{fib}(N)$  returns  $N^{\text{th}}$  fibonacci series

Main Logic

$$\text{fib}(N) = \text{fib}(N-1) + \text{fib}(N-2);$$

Base Case:

if  $(N == 0)$  return 0;

if  $(N == 1)$  return 1;



Note: In most of the cases, the # base cases = # distinct recursive sub-problems

```
int fib (int N) {
```

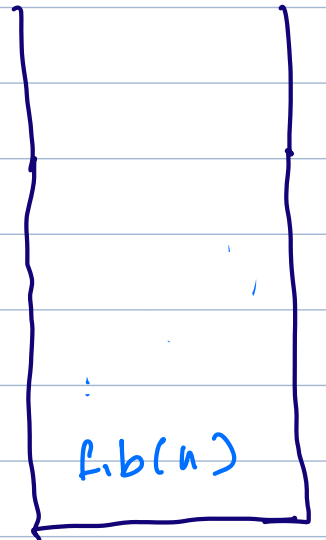
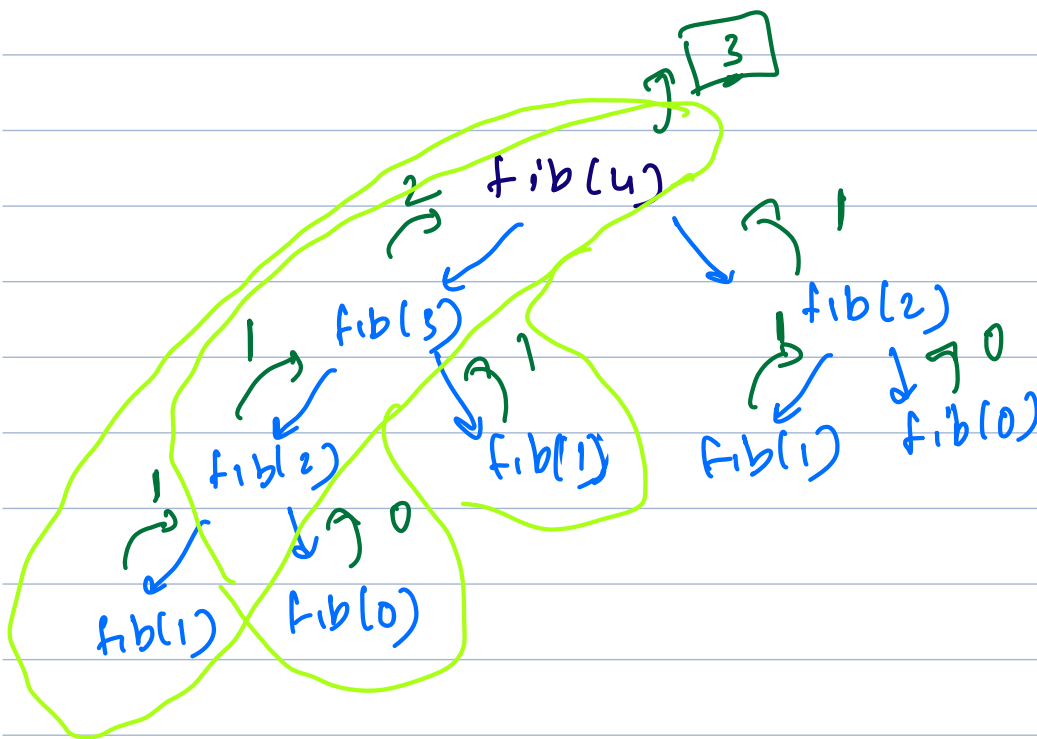
T.C:  $O(1)$

```
if (N == 0 || N == 1) return N;
```

S.C:  $O(1)$

```
return fib(N-1) + fib(N-2);
```

}



call stack

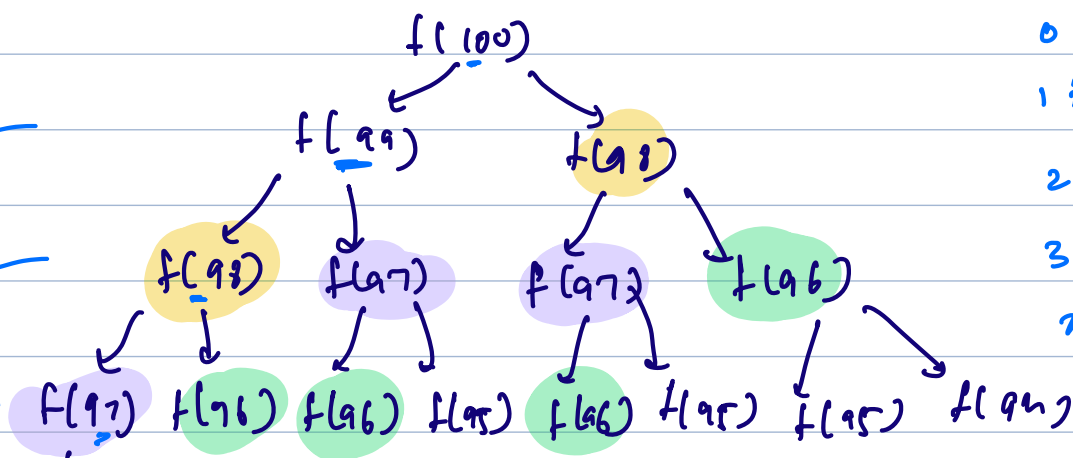
S.C:  $N \times O(1) \Rightarrow O(N)$

T.C:

21 ~~~~~

$$\frac{2^2}{2} \quad \sim$$

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$$\frac{99}{2} \Rightarrow f(1)$$


$$x+1 = 100$$

 $[0, N-1]$ 

# Function calls :  $2^0 + 2^1 + 2^2 + \dots + 2^{N-1}$

$$Sum_K = \frac{a(r^K - 1)}{r - 1}$$

$$a = 1$$

$$\gamma = 2$$

$$K \approx N$$

$$\therefore \frac{1(2^N - 1)}{2 - 1} = 2^N - 1$$

T.C:  $(2^N - 1) \times O(1) = \underline{O(2^N)}$

↓ DP

$O(N)$   $\leftarrow N = 10^6$

$$2^{20} \approx 10^6$$

$$2^{23} \approx \underline{10^7}$$

```
for (i = 1; i ≤ N; i++) {
```

```
    print("Hi");
```

```
    ↪ func()
```

// N<sup>2</sup>

}

O(N<sup>3</sup>)

```
sum(N) {
```

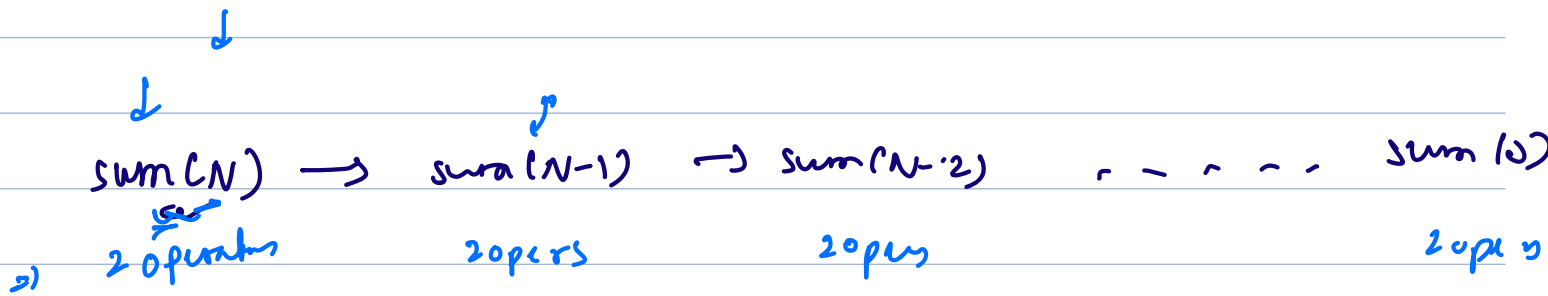
```
    return (sum(N-1)) + N;
```

}

sum(N-1) → sum(N-2) → sum(N-3) . . . sum(0) //

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[N-1]

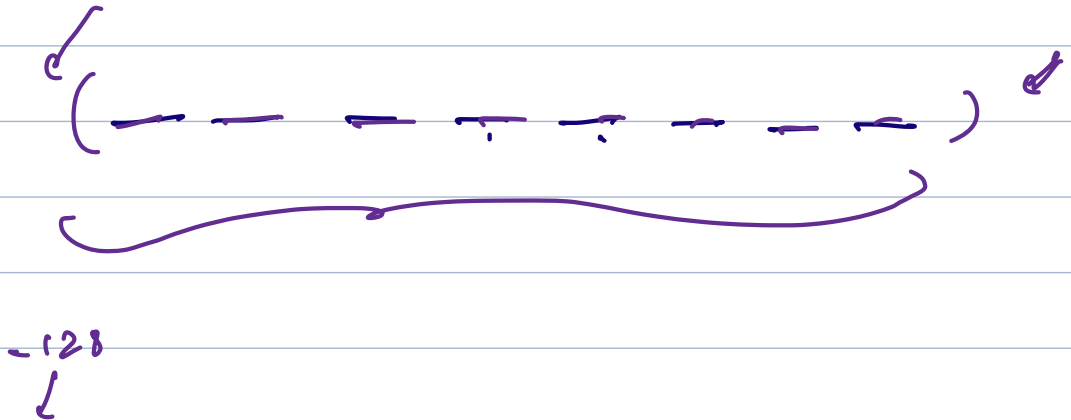


$$(N+1) \times 2 \text{ ops/iter} = 2N + 2$$

$$\Downarrow$$

$$O(N)$$

$$\approx 10^6 - 10^8$$



$a = 10^5$   
 $b = 10^6$   
 $a \times b$

$\int \text{int } c = a \times b;$