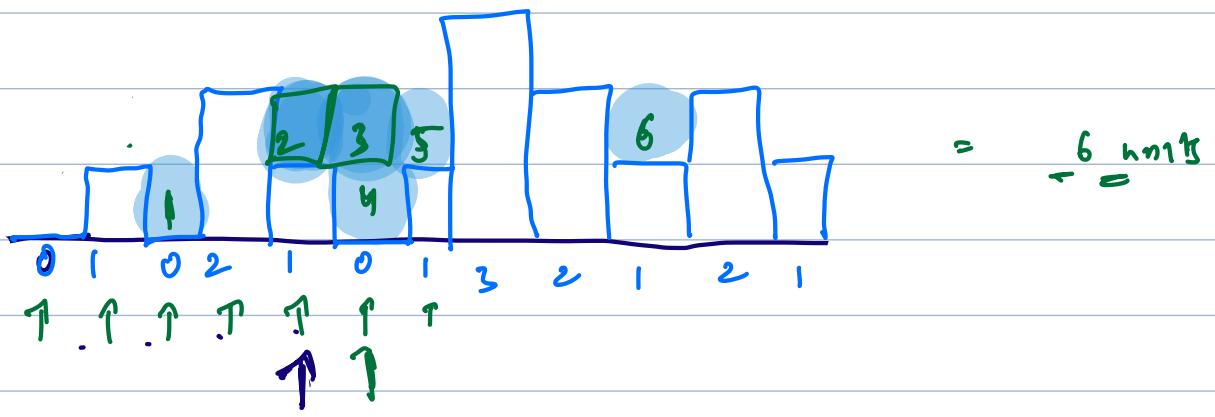


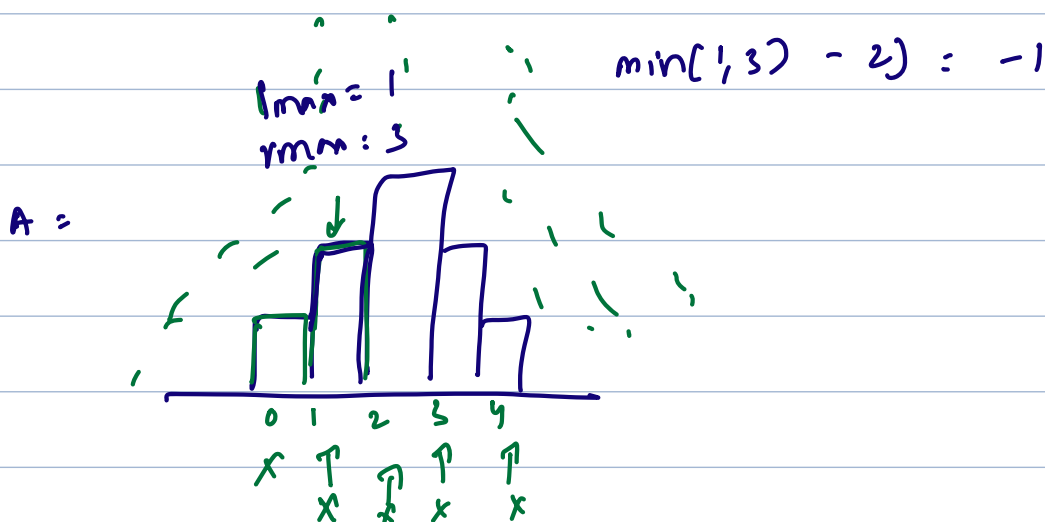
## Question: Rain Water Trapping

Given an array of size  $N$  representing height of buildings, compute how much water is trapped after it rains

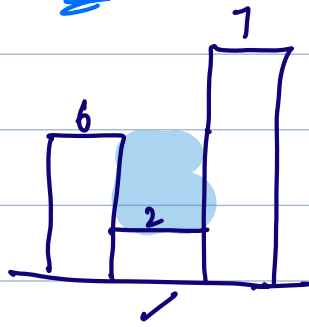
A: 0 1 0 2 1 0 1 3 2 1 2 1



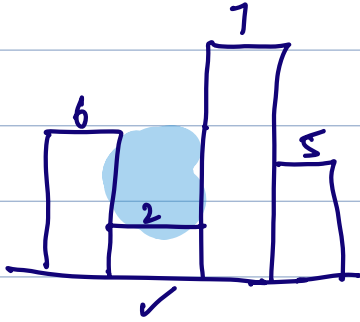
$$W_i = \min(\text{left-max}, \text{right-max}) - \text{building}$$



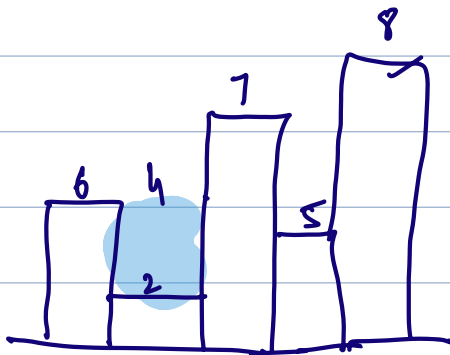
# Observations



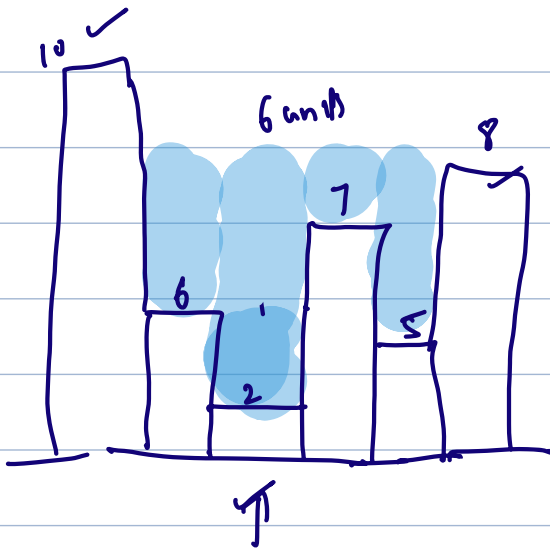
$$\Rightarrow \min(7, 6) - 2 = 4$$



$$\Rightarrow 4$$



$$\Rightarrow 4 \text{ units}$$



$$\min(10, 7) - 2 = 7 - 2 = 5$$

$$\text{water}[i] = \min(\text{leftMax}, \text{rightMax}) - \text{height}$$

## Brute Force

```
int totalWater = 0;
```

```
for (i = 0; i < N; i++) {
```

```
int lmax = max(0, i-1) // O(N)
```

```
int rmax = max(i+1, N-1) // O(N)
```

$$\text{water} = \min(lmax, rmax) - A[i];$$

```
if (water > 0) {
```

total water  $\neq$  water;

T.C:  $O(N^2)$

S.C: 0(1)

### Optimized Approach

	0	1	2	3	4	5	6	7	8	9	10
A :	4	2	5	7	4	2	3	6	8	2	3
map :	0	4	4	5	7	7	7	7	7	8	8
map :	8	8	8	8	8	8	8	8	3	3	0

$lmax[i] = \text{Max}(0, \dots, i-1)$

$lmax[i] = \max(lmax[i-1], A[i-1])$

$lmax[0] = 0;$

for( $i=1; i < N; i++$ ) {

$lmax[i] = \max(lmax[i-1], A[i-1]);$

}

T.C:  $O(N)$

$rmax[N-1] = 0;$

for( $i=N-2; i \geq 0; i--$ ) {

$rmax[i] = \max(rmax[i+1], A[i+1]);$

}

int totalWater = 0;

for( $i=0; i < N; i++$ ) {

left  $\rightarrow$  right

water =  $\min(lmax[i], rmax[i]) - A[i];$

if (water > 0) {

totalWater += water;

}

}

T.C:  $O(N)$

S.C:  $O(N)$

$\hookrightarrow$  2 arrays

optimization: We can only use 1 array by carrying forward  $lmax$ .

**Question:** Given a square matrix  $N \times N$ , print boundary elements in clockwise

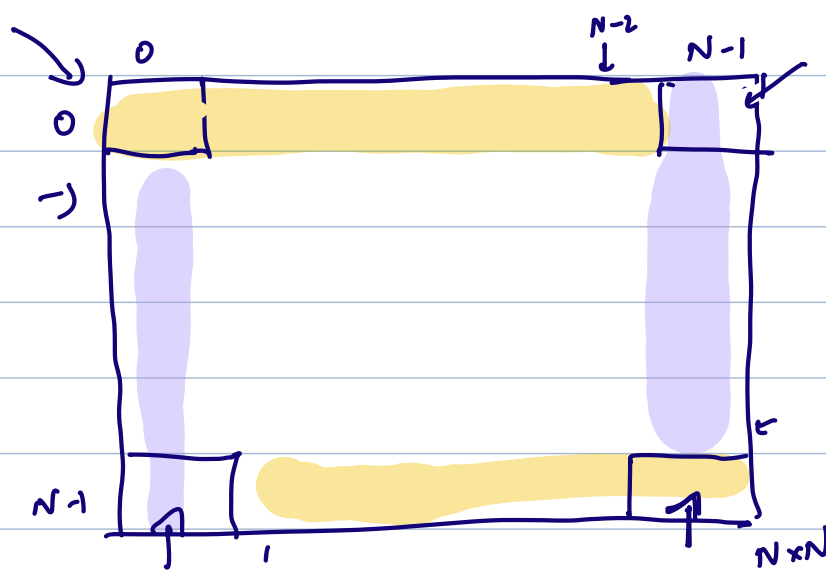
	0	1	2	3	4
0	1	2	3	4	5
1	6	7	8	9	10
2	11	12	13	14	15
3	16	17	18	19	20
4	21	22	23	24	25

5x5

1 2 3 4 5 10 15 20  
25 24 23 22 21  
16 11 6

Boundaries:

- 1)  $0^{th}$  row
- 2)  $(N-1)^{th}$  col
- 3)  $(N-1)^{th}$  row in reverse
- 4)  $0^{th}$  col in reverse



$$[0,0] \rightarrow [0,N-2] : N-1$$

$$[0,N-1] \rightarrow [N-2,N-1] : N-1$$

$$[N-1,N-1] \rightarrow [N-1,1] : N-1$$

$$[N-1,0] \rightarrow [1,0] : N-1$$

```
print Boundaries ( mat[][7] ) {
```

```
    int row = 0, col = 0 ;
```

```
    for ( k = 0; k < N-1; k++ ) {
```

```
        print ( A[row][col] );
```

```
        col++;
```

```
    }
```

```
    // row = 0    col = N-1
```

```
    for ( k = 0; k < N-1; k++ ) {
```

```
        print ( A[row][col] );
```

```
        row++;
```

```
    }
```

```
    // row = N-1, col = N-1
```

```
    for ( k = 0; k < N-1; k++ ) {
```

```
        print ( A[row][col] );
```

```
        col--;
```

```
    }
```

```
    // row = N-1 col = 0
```

```
    for ( k = 0; k < N-1; k++ ) {
```

```
        print ( A[row][col] );
```

```
        row--;
```

```
    }
```

```
    // row = 0, col = 0
```

```
}
```

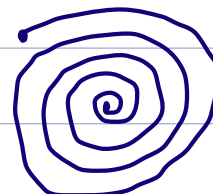
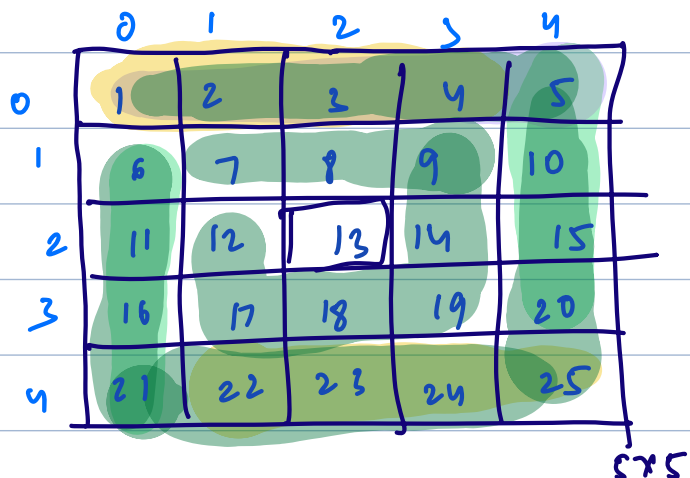
# Question: Lawn Mowing Challenge

## Scenario

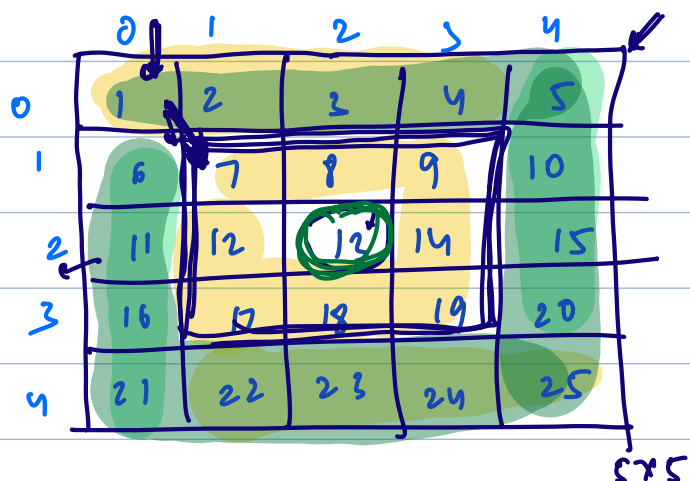
You need to program an automated grass-mowing robot for "GreenTech Robotics" to navigate a square lawn ( $N \times N$ ) represented as a grid. The lawn is designed as a **square grid**, where each cell in the grid shows the grass height in that area...

## Problem Statement

Your challenge is to find out the heights of grass patches that the robot needs to cut (in the order they are cut), with the robot's path following a **spiral pattern** from the **outer** edge towards the **center** of the lawn.



1 2 3 4 5 10 15 20 25 24 23 22  
21 16 11 6 7 8 9 14 19 18 17 12  
13



Step 1: (0,0)

$N = 5 \times 5$

Step 2: (1,1)

$N = 3$

Step 3: (2,2)

$N = 1$



Spiral Order ( mat[][7] ) {

```
    int row = 0, col = 0 ;  
    while ( N > 1 ) {  
        for ( k = 0; k < N-1; k++ ) {  
            print ( A[row][col] );  
            col++ ;  
        }
```

```
        // row = 0      col = N-1  
        for ( k = 0; k < N-1; k++ ) {  
            print ( A[row][col] );  
            row++ ;  
        }
```

```
        // row = N-1 ,   col = N-1  
        for ( k = 0; k < N-1; k++ ) {  
            print ( A[row][col] );  
            col-- ;  
        }
```

```
        for ( k = 0; k < N-1; k++ ) {  
            print ( A[row][col] );  
            row-- ;  
        }
```

```
        row ++ ;   col ++ ;  
        N = N - 2 ;
```

```
    }  
    // Edge case  
    if ( N == 1 ) print ( A[row][col] );  
}
```

T.C:  $O(N^2)$

S.C:  $O(1)$

8:17

Permutations:

$\Rightarrow$   $\begin{bmatrix} 1 & 2 & 3 \\ 2 & 3 & 1 \end{bmatrix}$   $\begin{bmatrix} 3 & 2 & 1 \\ 3 & 1 & 2 \end{bmatrix}$   $\begin{bmatrix} 1 & 3 & 2 \\ 2 & 1 & 3 \end{bmatrix}$

$\Rightarrow [1 \ 2 \ 3 \ 4] \rightarrow [1 \ 2 \ 4 \ 3] \rightarrow [1 \ 3 \ 2 \ 4] \rightarrow$   
 $[1 \ 3 \ 4 \ 2] \rightarrow [1 \ 4 \ 2 \ 3] \rightarrow [1 \ 4 \ 3 \ 2]$   
 $\rightarrow [2 \ 1 \ 3 \ 4] \rightarrow [2 \ 1 \ 4 \ 3] \rightarrow [2 \ 3 \ 1 \ 4]$   
 $\rightarrow [2 \ 3 \ 4 \ 1] \rightarrow [2 \ 4 \ 1 \ 3] \rightarrow \dots$

## Problem Statement

Implement the next permutation, which rearranges numbers into the numerically next greater permutation of numbers for a given array A of size N.

If such arrangement is not possible, it must be rearranged as the lowest possible order, i.e., sorted in ascending order.

### NOTE:

The replacement must be in-place, do not allocate extra memory.

DO NOT USE LIBRARY FUNCTION

$[1, 2, 3] \rightarrow [1, 3, 2]$

$[3, 2, 1] \rightarrow [1, 2, 3]$

$[1, 2, 6, 3, 4, 5] \rightarrow 1, 2, 6, 3, 5, 4$

$1, 2, 4, 3, 5, 6 \rightarrow 1, 2, 4, 3, 6, 5$

⇓

Ex:  $\begin{matrix} 7 & 8 & 4 & 1 & 2 & 3 & 5 & 6 & 7 \\ 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & \end{matrix}$

3 → 4

$7, 8, 3, 6, 5, 4, 2, 1$

6, 5, 4, 3, 2, 1

A :

4 3 1 5 8 7 6 2

0 1 2 3 4 5 6 7

A' :

4 3 1 5 2 6 7 8

i = -1

i  
↓  
2 1

8  
7  
6  
5

```
vector<int> Solution::nextPermutation(vector<int> &A) {
```

```
    int len = A.size();
```

```
    int i, j;
```

```
    for(i = len - 2; i >= 0; i--)
```

```
        if(A[i] < A[i + 1])
            break;
```

```
    // the array is in descending order
```

```
    if (i == -1) {
```

```
        reverse(A.begin(), A.end());
        return A;
```

// Find element just greater than A[i];

```
    for(j = len - 1; j > i; j--)
```

```
        if(A[j] > A[i])
            break;
```

```
    // swap with the smallest number in the suffix
```

```
    swap(A[i], A[j]);
```

```
    // reversing the suffix
```

```
    reverse(A.begin() + i + 1, A.end());
```

```
    return A;
```

```
}
```

O(N)

→ O(N)

→ O(N)

O(N)

T.C: O(N)

S.C: O(1)

i = -1 i i i

↓ ↓ ↓ ↓

6 5 3 2

$\downarrow$       $\downarrow$     $\downarrow$   
 1   2   9   8   7  
     $\downarrow$

$A[i] < A[i+1]$

break;

}

$i = i - 1$     $i$     $i$     $i$   
 $\downarrow$       $\downarrow$     $\downarrow$     $\downarrow$   
      9   8   7   6

$i$       $i$       $i$       $i$   
 $\downarrow$       $\downarrow$       $\downarrow$       $\downarrow$   
 7   4   6   5   3   2  
     $i$       $i$       $i$       $i$