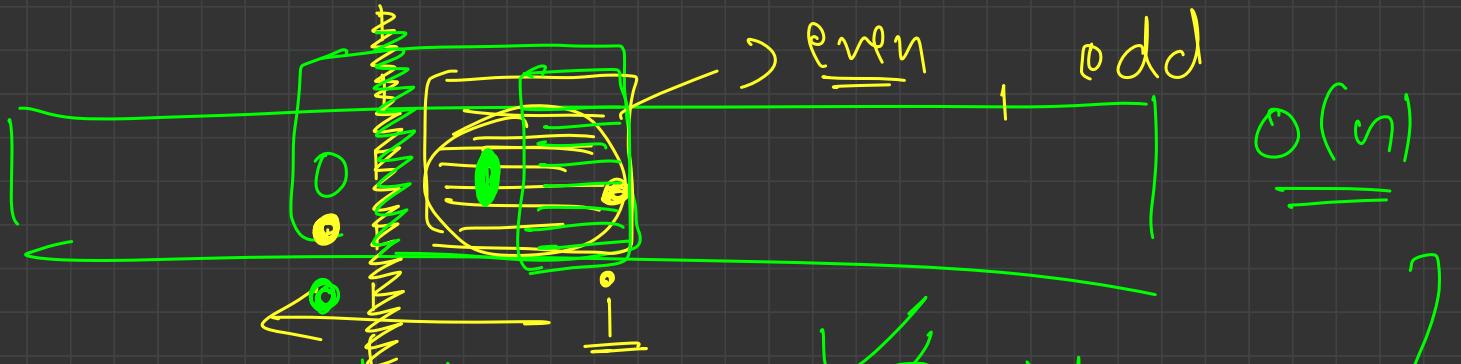


Dynamic Programming 2.3

- Priyansh Agarwal

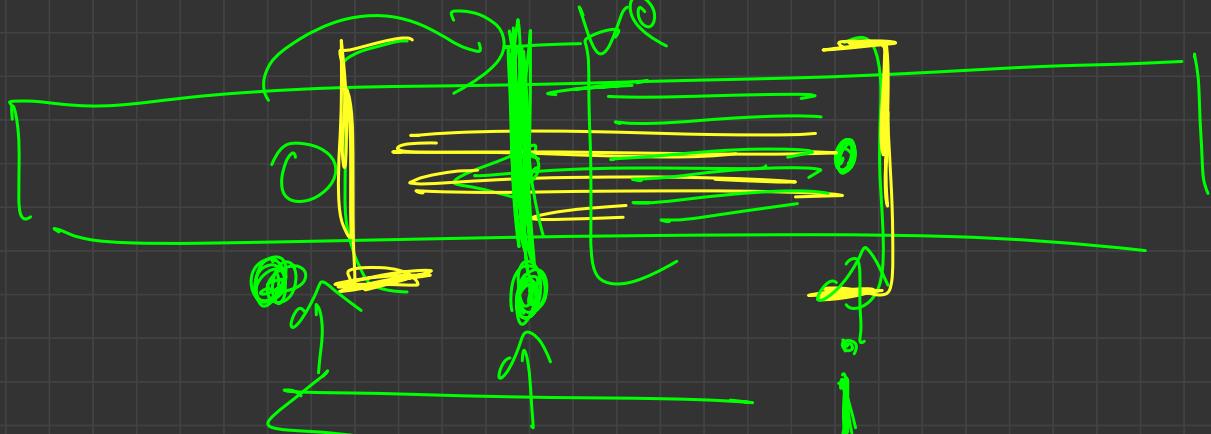


\rightarrow has product $=$ Yes / No

$2 \rightarrow$ has product \rightarrow +ve

$3 \rightarrow$ has product \rightarrow -ve

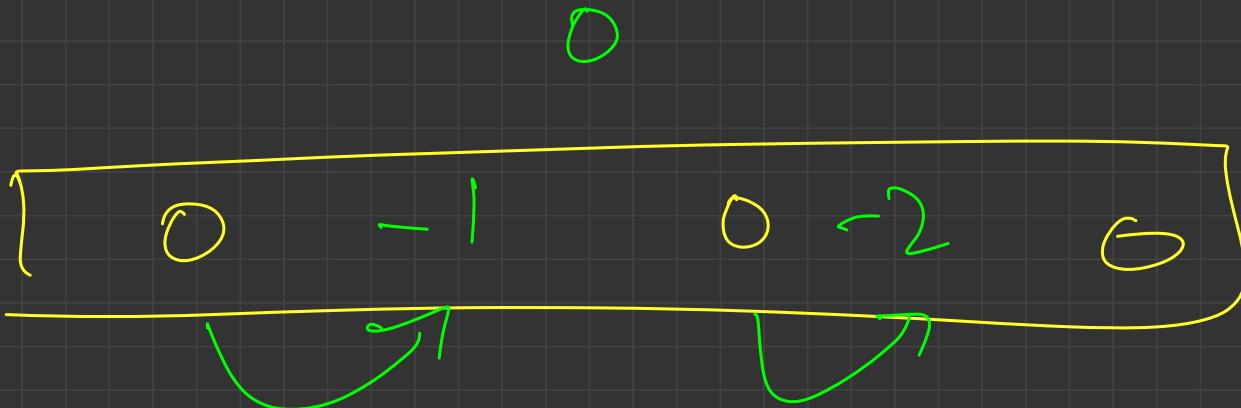
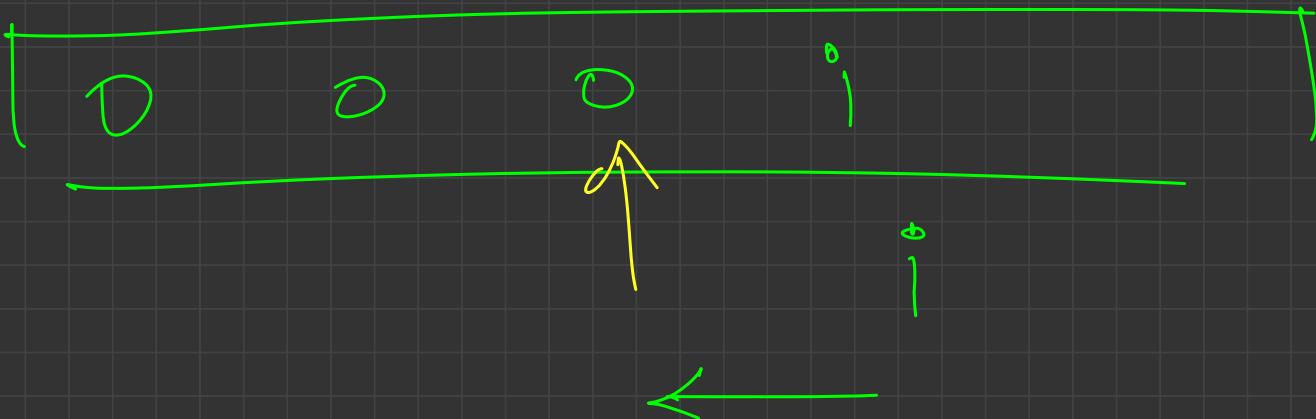
$-1 - (- - 1 - 1)$



Even $\rightarrow -\sqrt{e}$

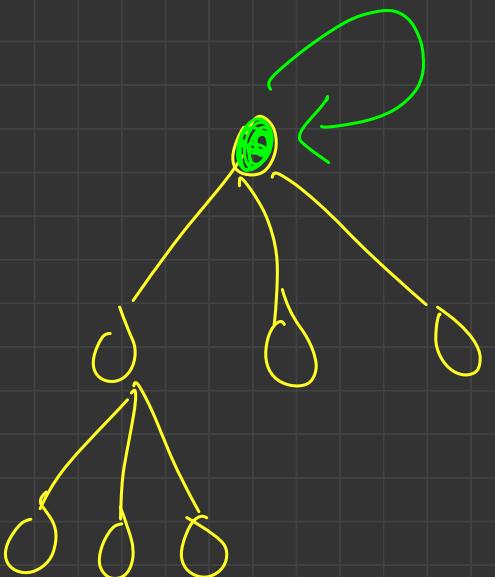
$O(n)$ \rightarrow O
 $O(n)$ \rightarrow even fold - we \rightarrow odd
 $O(a)$ \rightarrow \sqrt{e}

$O(n)$



① Representing Non-Integer DP
parameters

② Cycling of states



Representing non-integer parameters

- How will you store the dp states if instead of integer parameters you had a string or a vector or a map or any complex data type?
- Use a map instead of an array.
- Tradeoff `map<pair<int, string>>` DP or `vector<map<string>>` DP

$d f(i) \{ j \} (k)$ \rightarrow (i, j, k) are all
the integers

$d f("Priyanshi") (10) \{ 2 \}$

$\delta g^{(1)}, \delta g^{(2)} \dots \delta g^{(n)}$

$\delta g^{("mikanu")}$ $\delta g^{("TLE")}$

Memoization :

→ states

already calculated

Array

to be calculated

Calculate & then store very fast

$\sum \underline{\underline{O(1)}}$

to access it

$\underline{\underline{O(1)}}$

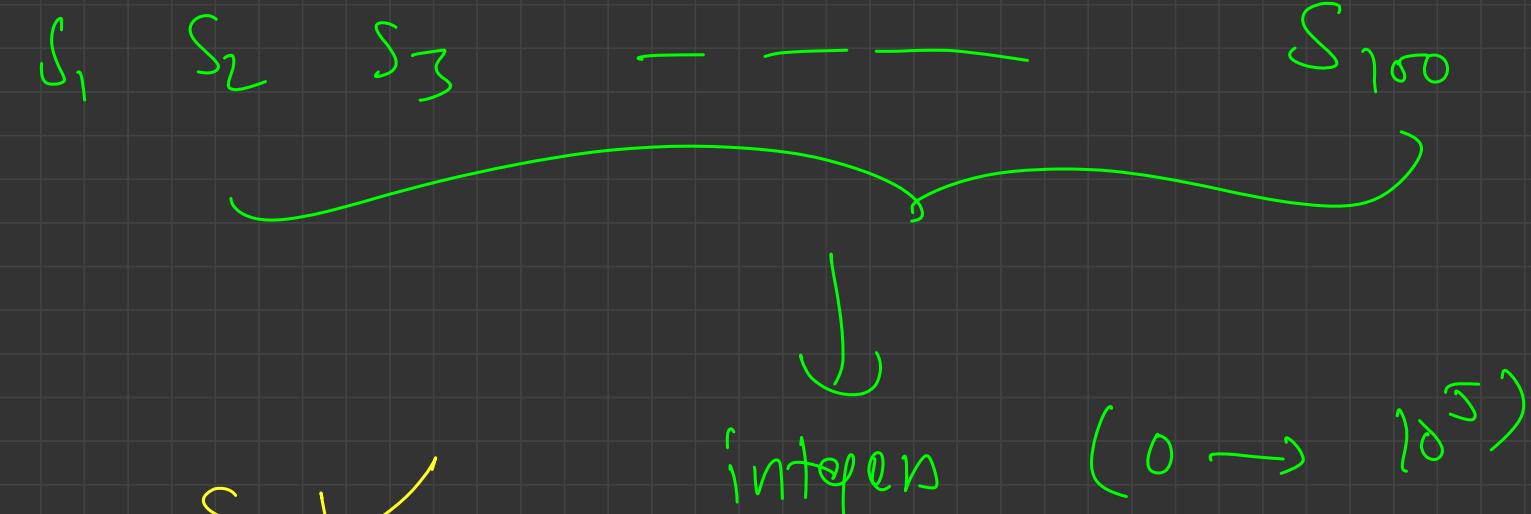


$d\varphi(\gamma, d\varphi^2 \gamma) - \underline{\underline{d\varphi^n}}$

Access time $\rightarrow \mathcal{O}(1)$

Update | Store time $\rightarrow \mathcal{O}(1)$ $\leq \underline{\underline{10}}$

$d\varphi^{(1)}(\text{tryout}) \rightarrow$ hash ("tryout")
into an integer
 $\leq w^5$



$S_1 \rightarrow \text{hash}(S_1) \rightarrow 100$

$S_{78} \rightarrow \text{hash}(S_{78}) \rightarrow 100$

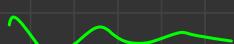
Hash a string and convert to
integers and then use
normal dp



Access $\rightarrow O(1)$
Update $\rightarrow \underline{O(1)}$

No ~~XX~~

map < key, value >



Access time } $\geq \log n$
insert | update } =

integer, string, vector, map, set

1. choose value

Hash map

Hash map

Tree map

Hash map

Keys are in a random order

Update

Insert

Delete

Search

Avg time $\rightarrow \Theta(1)$

Worst time $\rightarrow \Theta(n)$

~~Θ(n^2)~~

Normal Map | Tree Map

Keys are sorted

Avg $\rightarrow \Theta(\log n)$

Worst $\rightarrow \Theta(\log n)$

$\delta_f \{ \| T(\epsilon) \| \} \rightarrow \text{int-d } \delta(f)$

$\overbrace{\quad\quad\quad}$ $\overbrace{\quad\quad\quad}$ \times

$\text{map} < \text{string}, \text{int} > \underline{\delta_f}$

$\overbrace{\quad\quad\quad}$

df | string
vector
map
set
double

map ✓

dp [integer] (string)

$dp(2)["TLE"]$

\equiv

(1) $dp[\{ 2, "TLE" \}]$

\equiv

(map < pair < int, string >, int > dp) =

(2)

$dp(2)["TLE"]$

vector < map < string, int > > dp

of $f_1 f_2$ (dry run)

10^2 10^2

10^4

(1)

$\log(10^8)$

Map < pair < pair < int, int >, string >, int > df

(2)

vector < vector < map < string, int >>> df

) $\log(10^4)$

$\text{map} < \underline{\text{pair<int, string>}}, \text{int} > \text{df}$



total

combinations

20 iterations

$$\rightarrow 10^6$$

total keys in map

$$\rightarrow 10^6 \rightarrow$$

$\text{df}(\text{int})(\text{string})$

=

$$10^3$$

$$10^3$$

integer

string

$$\log(10^6) \rightarrow \underline{\underline{20}}$$

vector < map < string, int > df
20³
integer

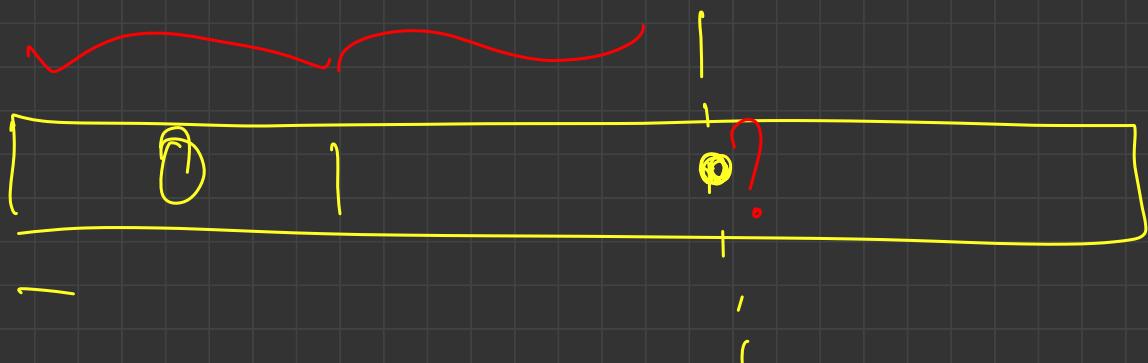
df(2){^{if T(E)}}
1 = 10
lo opertion

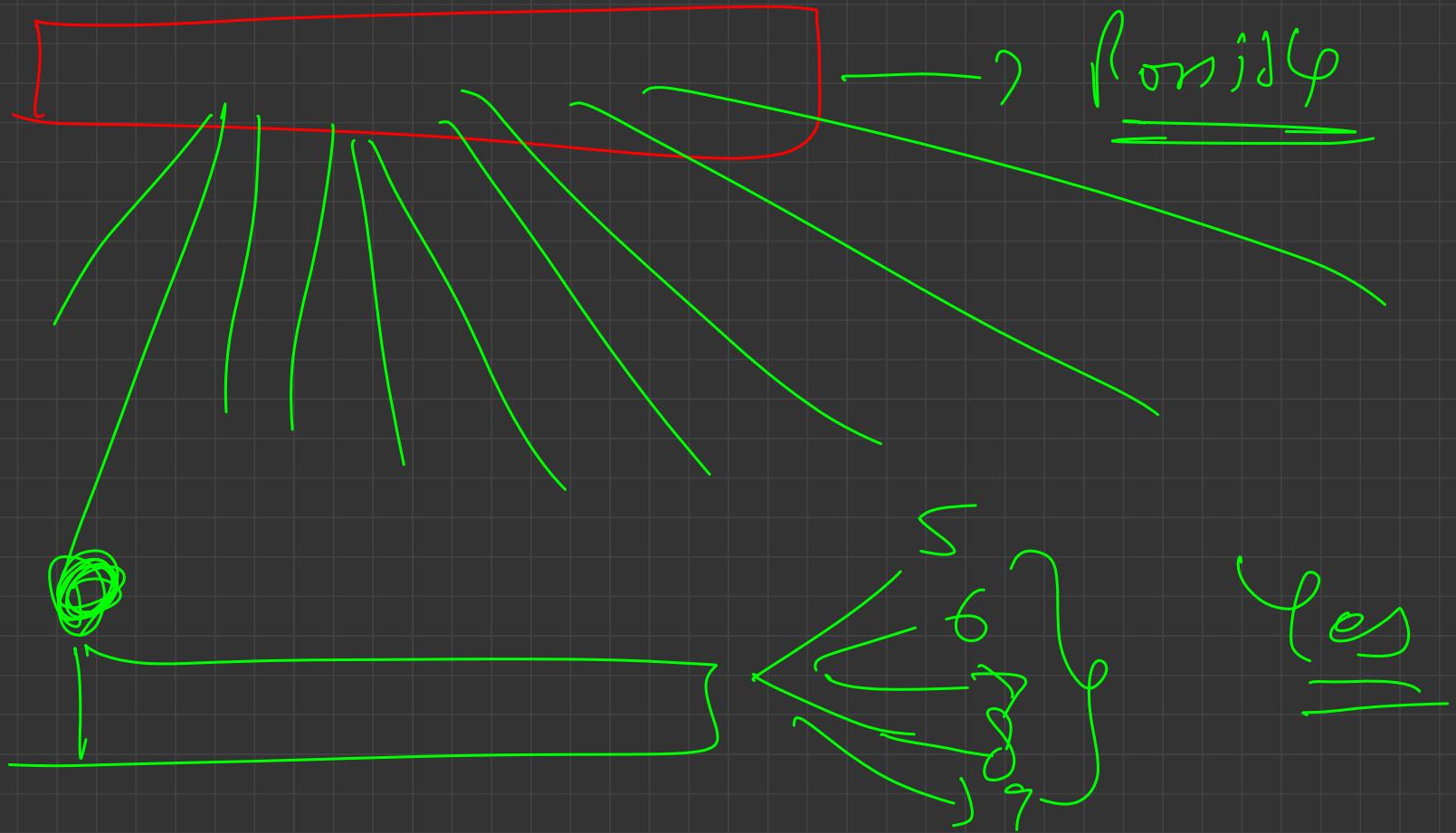
~~O(1)~~ → map < string, int >
11 T(E)
10³ string
O(log(10³)) 21.0

Problem: Link

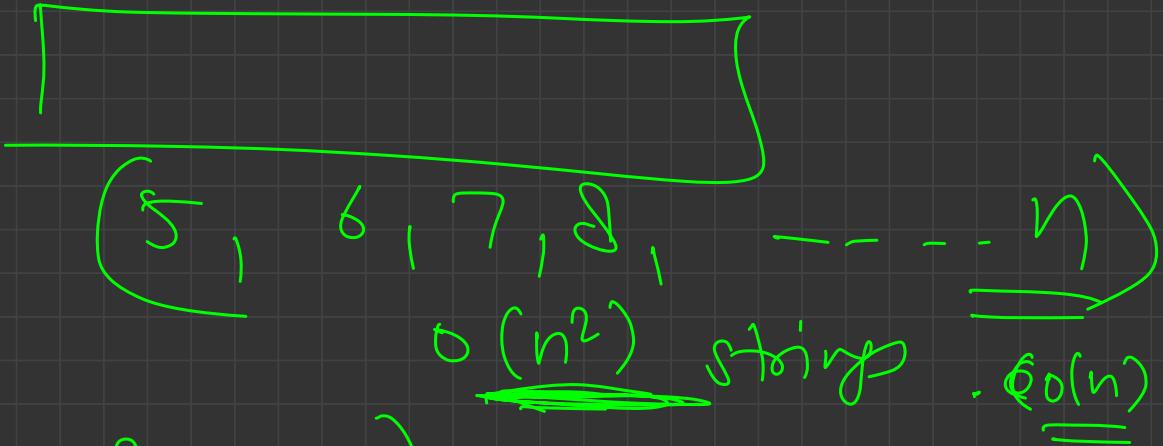
- State:
 -
- Transition:
 -
- Base Case:
 -
- Final Subproblem:
 -

String of length \underline{N}





$N \rightarrow$



$O(2^n \cdot n^2 \cdot n)$



generation

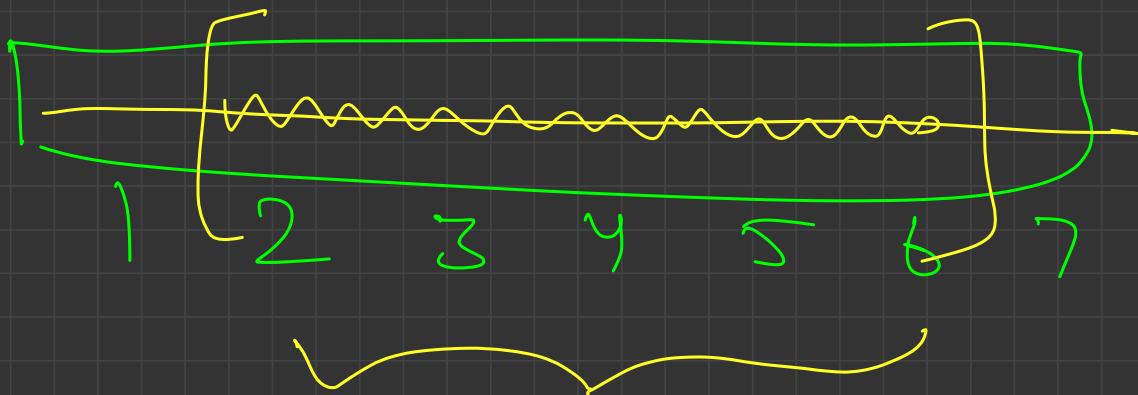


substring

$N \rightarrow 15$

falindromes
(check)

Palindrome of length \rightarrow 7



① if a string doesn't have a
palindrome of length 5
then it won't have a palindrome

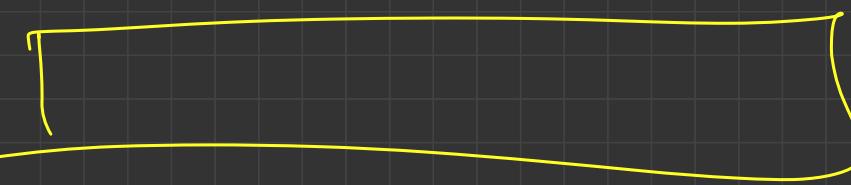
of length 7, 9, 11, 13 ...

② Similarly for length 6

for a string:

$S_{(6,7,8,7,6)}$
---. 5

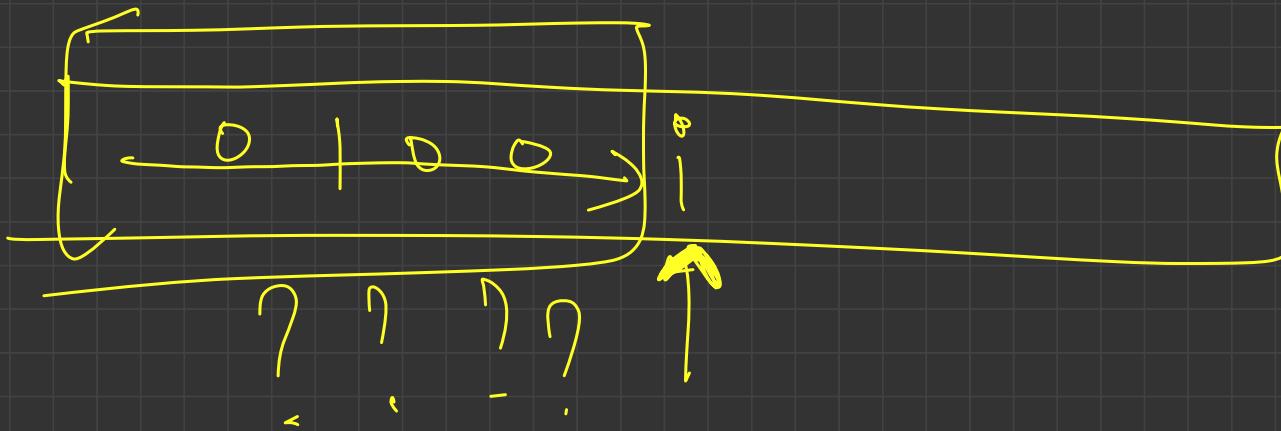
$S_{(6, \cancel{7}, 8)}$

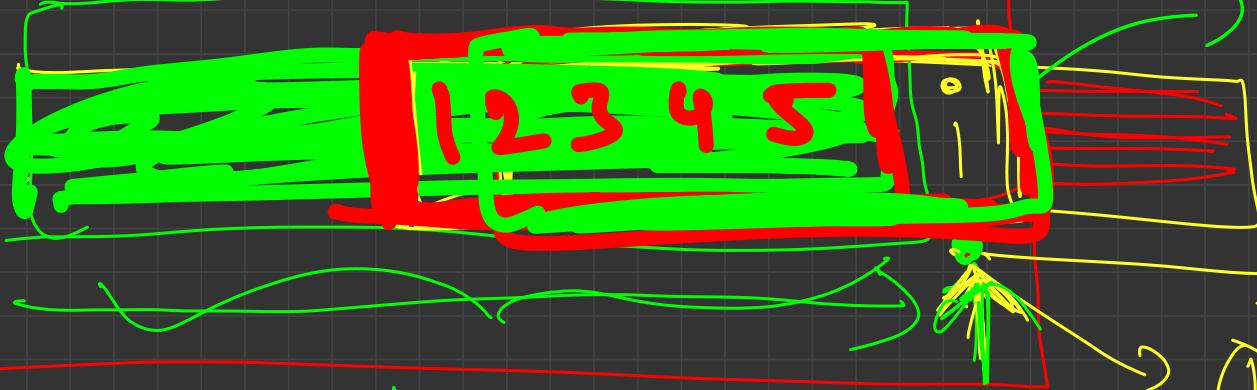
$$2^n \cdot n^3$$


$$\underbrace{5}_1 = \underbrace{6}_{\sigma(n)}$$

$$\underbrace{\sigma(n)}_{\Omega(n)} \cdot \underbrace{\sigma(n)}_{\Omega(n)}$$

$$2^n \cdot n^2$$



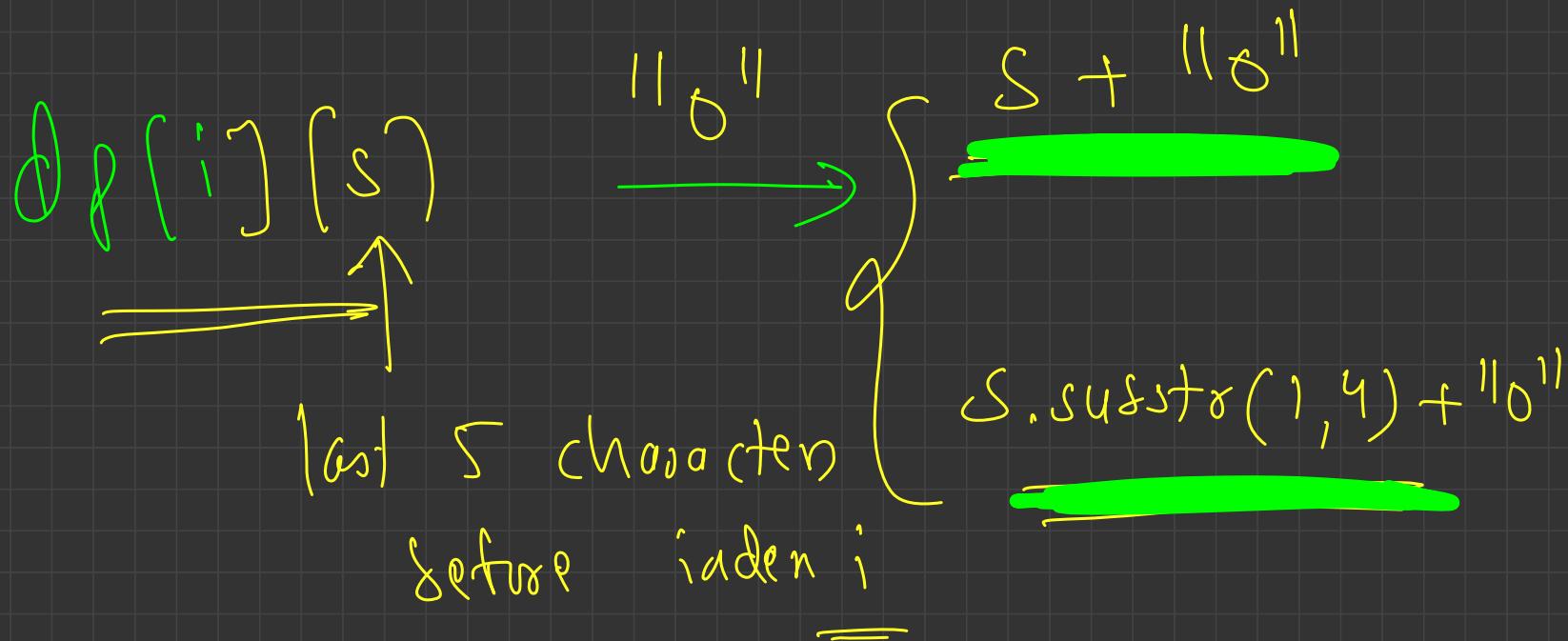


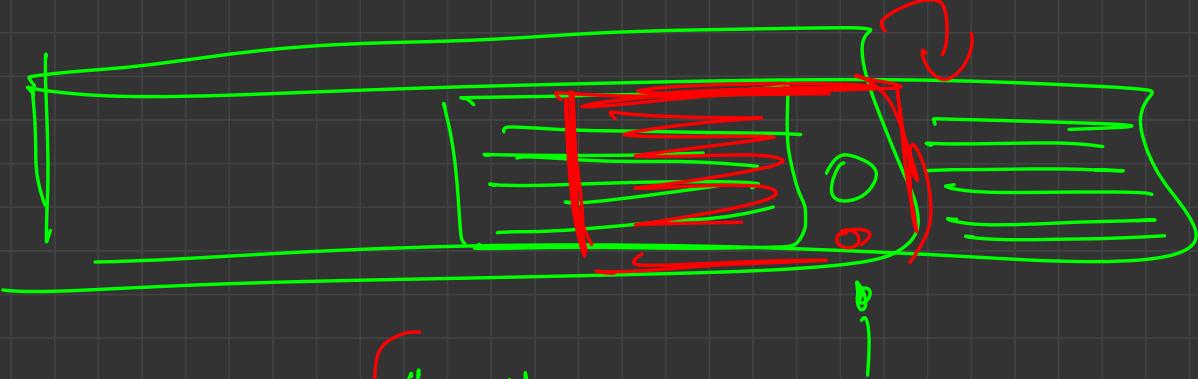
there are no substrings of
length 5 or greater which are
palindromes

$dp(\text{index})["\text{last } 5 \text{ characters}"]$

→ I have filled up the string till $(\text{index} - 1)$ such that the last 5 characters are → is there a way to replace question marks from (index) to $(n - 1)$ such that no

Polindromes of length 5 or greater exist

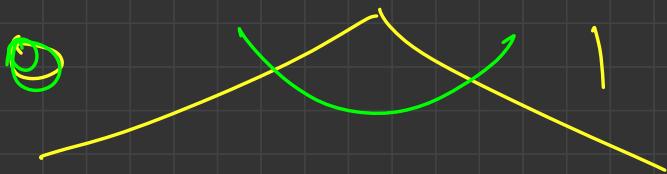




$$\delta f[i][s] \xrightarrow{\text{check for palindromes}} \delta f[i+1][s - \text{substr}(1, y) + \lVert y \rVert]$$

$d\varphi(i \gamma(s))$

tone



$$d\varphi(i+1) \left[s \cdot \text{suft}(1, 4) + \eta_0^{(1)} \right] d\varphi(i+1) \left[s \cdot \text{suft}_8(1, 4) + \eta^{(1)} \right]$$



bool possible (int index, last, s)

if (index == n)

return true

if (diff(index)(last) != -1)

return false

else → check for palindromes

possible(index+1, last.substring(1, n)
+ s(index), s)

$\begin{bmatrix} 0 & 0 & | & 0 & 0 & 0 & ; \end{bmatrix}$

$\begin{bmatrix} 1 & 0 & | & 0 & 0 & 0 & ; \end{bmatrix}$

$\begin{bmatrix} 0 & 0 & | & 0 & 1 & 0 & 0 & ; \end{bmatrix}$

$\begin{bmatrix} 0 & 1 & | & 0 & 1 & 0 & ; \end{bmatrix}$

$\begin{bmatrix} 1 & 0 & | & 0 & 1 & 0 & ; \end{bmatrix}$

$\begin{bmatrix} 1 & 1 & | & 0 & 1 & 0 & 0 & ; \end{bmatrix}$

Cycling DP states

- What happens when your current state is dependent on itself?
 - $dp[i]$ depending on $dp[i]$ itself

Problem 1

- Given a positive integer $N \leq 1e6$, at every step the following 3 things can happen to N with equal probability.
 - $N = N / 2$
 - $N = N - 1$
 - N remains unchanged
- Find expected number of steps it will take to convert for N to become 0

Problem 1

state:

$dp[n]$ = expected number of steps to convert n to 0

transition:

$dp[n] = \frac{1}{3} (1 + dp[n - 1] + 1 + dp[n / 2] + 1 + dp[n])$

$dp[n] = \frac{1}{3} (3 + dp[n - 1] + dp[n / 2] + dp[n])$

$dp[n] = 1 + dp[n - 1] / 3 + dp[n / 2] / 3 + dp[n] / 3$

$\frac{2}{3} * dp[n] = 1 + dp[n - 1] / 3 + dp[n / 2] / 3$

$dp[n] = \frac{3}{2} + (dp[n - 1] + dp[n / 2]) / 2$

base case:

$dp[0] = 0$

final subproblem = $dp[n]$

Problem 2: Link