

1Q57 D \Rightarrow Yet Another Monster Killing Problem

$n \rightarrow a_1, a_2, a_3, a_4, a_5, a_6$

$H \rightarrow \begin{matrix} p_1 & p_2 & \dots \\ s_1 & s_2 & \dots \end{matrix}$

p_i kills till $i \leq s_i$
&
 $\max(a_1, a_2, \dots, a_i) \leq p_i$

Approach

- Suppose we have p_k, s_k such that
 $p_k \geq p_j$ & $s_k \geq s_j \rightarrow$ remove $p_j, s_j \rightarrow$ why?

proof \rightarrow

Let $\max(a_1, a_2, \dots, a_j) \leq p_j$

Since $p_k > p_j$ it's obvious that

$$\max(a_1, a_2, \dots, a_j) \leq p_j \leq p_k$$

Now $j \leq s_j$ where $s_j \leq s_k$

so $j \leq s_k \rightarrow$ it's obvious that we always pick (p_k, s_k) over (p_j, s_j) {greedy}

- ① Delete all the entries such that $(p_j, s_j) < (p_k, s_k)$ where $p_j \leq p_k$ & $s_j \leq s_k$

We will be left with entries \rightarrow

(p_1, s_1)

(p_2, s_2)

(p_3, s_3)

\vdots

(p_n, s_n)

such that if $p_2 > p_1$

then $s_1 > s_2$

Ex - 1, 6

3, 4

8, 2

10, 1

\rightarrow valid list of heroes

② Pick a hero that can kill maximum monsters in one go

$m \rightarrow 6 \ 2 \ 3 \ 8 \ 4 \ 11 \ 2 \ 1 \ 5$

$h \rightarrow 6 \ 2 \ 12$

11 1

5 4

2 6

6 2 3 8 4 11 2 1 5

$\uparrow \uparrow$
 $i \ j$

do we have a guy whose

$p \geq \max(6, 2) = 6$ &

$s \geq 2$?

\hookrightarrow if yes, we are able to kill till j , increment it

6 2 3 8 4 11 2 1 5
 $\uparrow \quad \uparrow$
 $i \quad j$

$p \geq \max(6, 2, 3) = 6$

& $s \geq 3 \rightarrow$ no!

\rightarrow delete till $j-1$ &

$i \rightarrow j \quad j \rightarrow i+1$

sum += 1

3 8 4 11
 $\uparrow \quad \uparrow$
 $i \quad j$

2 1 5
 $\xrightarrow{\text{del by } p=5}$

ans = 6 Similarly