

## Random Probing

Suppose that the hash table has  $b$  buckets ( $mTableSize$  is  $b$ ). In linear open addressing the buckets are examined in the order  $(f(k)+i) \% b$ ,  $0 \leq i < b$ , where  $k$  is the key of the element being searched for. In random probing a pseudo-random number generator is used to obtain a random sequence  $R(i)$ ,  $1 \leq i < b$  where  $R(1), R(2), \dots, R(b-1)$  is a permutation of  $[1, 2, \dots, b-1]$ . The buckets are examined in the order  $f(k), (f(k)+R(i)) \% b$ ,  $1 \leq i < b$ .

Note: the Random class is in **java.util**

The easiest way to generate  $R(i)$  is to instantiate a Random object (BEFORE the while loop in `findPos`), for which you pass  $f(k)$  (what `myHash` returns). Then if you iterate in the loop, call the Random object's (you already instantiated) `nextInt()` method to get the next "random" number (and similar to `myHash`, make sure that after you `% mTableSize`, that if it's  $< 0$ , you add `mTableSize`).

There is a possibility that after adding a particular random value (and after `%`, checking if  $< 0$ ), you may get the same location again. In a "safer" implementation, you should check if the new location is the same as the previous (and get the `nextInt` again if so).