

MIT World Peace University
Analysis of Algorithms

Unit 1

NAMAN SONI ROLL No. 10

Contents

1	Divide and Conquer	2
1.1	Control Abstraction	2
1.2	Time Complexity of the general algorithm	2
1.3	Methods for Solving recurrences	2
1.4	Math you need to Review	2
2	Divide-and-Conquer Examples	3
3	Proof Techniques	3
3.1	Proof by contradiction	3
3.2	Proof by Mathematical Induction	3
3.3	Direct Proof	4

1 Divide and Conquer

1.1 Control Abstraction

```
1  DANDC (P)
2  {
3      if SMALL (P) then return S (p);
4      else
5      {
6          divide p into smaller instances p1, p2,...Pk, k>=1;
7          apply DANDC to each of these sub problems;
8          return (COMBINE (DANDC (p1), DANDC (P2),...,DANDC (pk)));
9      }
10 }
```

1.2 Time Complexity of the general algorithm

- A recurrence is an equation or inequality that describes a function in terms of its value on smaller inputs.
- Special techniques are required to analyze the space and time required.
- $T(n) = \frac{aT(\frac{n}{b}+1)(n)+c(n)}{O(1)}$
- Time Complexity (recurrence relation): (
 - where D(n): time for splitting
 - C(n): time for conquer
 - c: a constant)

1.3 Methods for Solving recurrences

1. Substitution method: This method involves guessing a solution and then proving that it is correct.
2. Recurrence tree method: This method involves constructing a tree diagram that represents the recursive calls and their relationship to each other.
3. Master theorem: This is a general theorem that provides a method for solving recurrences of a specific form.

1.4 Math you need to Review

Properties of Logarithms:

- $\log_b(xy) = \log_b(x) + \log_b(y)$
- $\log_b(\frac{x}{y}) = \log_b(x) - \log_b(y)$
- $\log_b xa = a \log_b x$
- $\log_b a = \frac{\log_x a}{\log_x b}$

Properties of exponentials:

- $a^{(b+c)} = a^b a^c$
- $a^{bc} = (a^b)^c$

- $\frac{a^b}{a^c} = a^{(b-c)}$
- $b = a^{\log_a b}$
- $b^c = a^{c \cdot \log_a b}$

2 Divide-and-Conquer Examples

- Sorting: mergesort and Quick Sort
- Binary tree traversals
- Binary search
- Multiplication of Large integers
- Matrix Multiplication: Strassen's algorithm
- Closest-pair and convex-hull algorithms

3 Proof Techniques

- Proof is a kind of demonstration to convince that the given mathematical statement is true.
- The statement which is to be proved is called theorem. Once a particular theorem is proved then it can be used to prove further statements.
- The theorem is also called Lemma.
- The proof can be a deductive proof or inductive proof.
- The deductive proof consists of a sequence of statements given with logical reasoning.
- The inductive proof is a recursive kind of proof which consists of a sequence of parameterized statements that use the statement itself or the statement with lower values of its parameters/.

3.1 Proof by contradiction

- In this type of proof, for the statement of this form is A and the B.
- Prove by contradiction. There exist two irrational numbers x and y such that x^y is rational.
 - An irrational number is any number that cannot be expressed as a/b where a and b are integers and value b is non zero. To prove that x^y is rational when x and y are rational numbers.

3.2 Proof by Mathematical Induction

- Prove $1 + 2 + 3 + \dots + n = n(n+1)/2$
 - 1) Basis of induction
 - Assume, $n = 1$ then
 - LHS = $n = 1$
 - RHS = $n(n+1)/2 = 1(1+1)/2 = 2/2 = 1$
 - 2) Induction hypothesis
 - Now we will assume $n = K$ and will obtain the result for it. The equation then becomes,
 - $1 + 2 + 3 + \dots + K = K(K+1)/2$

- 3) Inductive step
- Now we assume that equation is true for $n = K$ and we will then check if it is also true for $n = K + 1$ or not.
- Consider the equation assuming $n = K + 1$
- $LHS = 1 + 2 + 3 + \dots + K + K + 1$
- $= K(K + 1)/2 + K + 1$
- $= K(K + 1) + 2(K + 1)/2$
- $= (K + 1)(K + 2)/2$

3.3 Direct Proof

- In direct Proof, the intended proof can be proved by basic principle or axiom.
- Example:- Prove that the negative of any even integer is even.
- Solution: to prove this, let n be any positive even number. Hence we can write n as
- $n = 2m$ where m can be any number
- if we multiply both side by -1 , we get
- $-n = -2m$
- $-n = 2(-m)$
- Multiplying any number by 2 makes it an even number.
- Hence, $-n$ is even.
- Thus proves that the negative of any even integer is even.

3.4 Proof by Counter-Example

- This is a technique of proof in which $a \rightarrow b$ is true if $\neg a \rightarrow \neg b$
- If negative statement of given statement is true then the given statement becomes automatically true.
- Example: prove contraposition that $x + 8$ is odd
- Solution:
 - Step 1: we assume that x is not odd
 - Step 2: that means x is even. By definition of even numbers $2 * \text{any number} = \text{even number}$.
 - $x = 2 * m$ where m can be any number
 - we can write $x + 8$ as $2 * m + 8 = 2(m + 4) = \text{even number}$
 - thus, $x + 8$ is even