Sorting

Heapsort

Heapsort is a time- and space-efficient sorting algorithm

- its best-, average-, and worst-case complexity are all O(n log n)
- it has O(I) space complexity (better than either recursive algorithm we saw... why?)

Heapsort works by first transforming the data into a heap

- a heap is a tree-based data structure that makes finding the largest (or smallest) value a constant-time operation
- in order to understand heapsort, it's important to understand exactly what a heap is

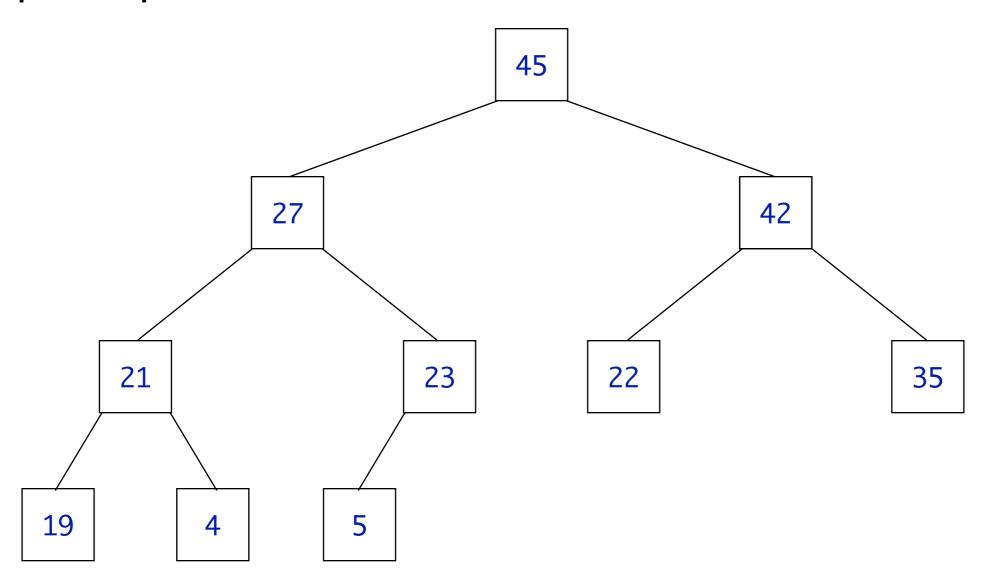
A heap is:

- a binary tree where values are organized using a strict weak ordering (see below)
- the value contained by a node is always larger than the values in the node's children (for a max heap)
- the tree is a complete binary tree (every level except the deepest hast as many nodes as possible, and all nodes at the bottom level are as far left as possible)

Strict weak ordering (common-sense orderings):

- x < x is never true (irreflexivity)</pre>
- if x != y, then x < y and y < x cannot both be true (asymmetric)
- if x < y and y < z, then x < z (transitivity)

Example heap:



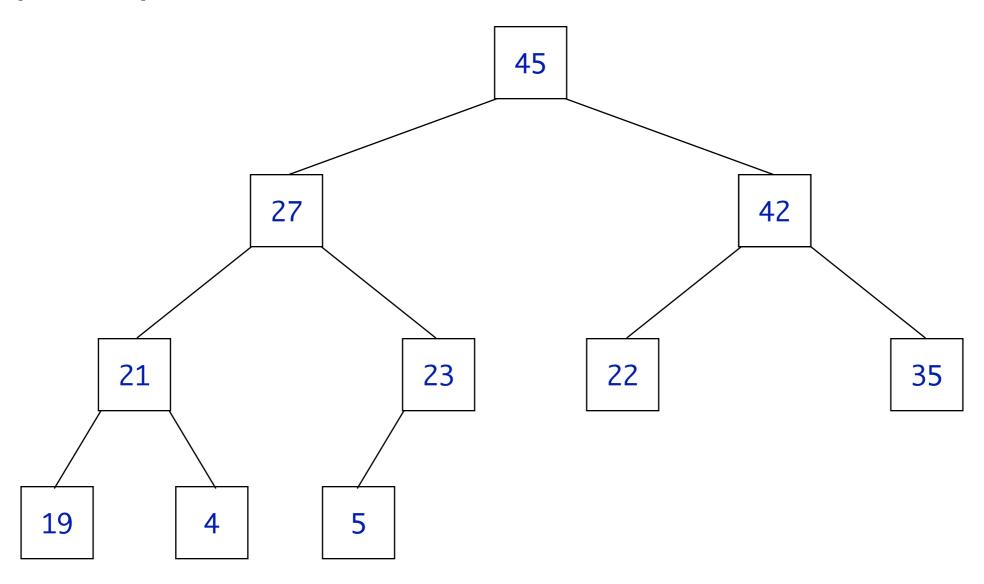
Heaps are complete binary trees

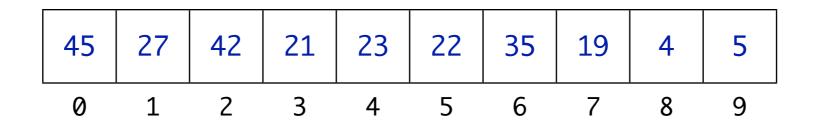
- all levels but the deepest contain the maximum possible
- nodes in the deepest level are as far left as possible

We've seen an efficient array representation for complete trees:

- data at the root of the tree is stored at index 0
- the left child of the element at index i is stored at index 2i+1
- the right child of the element at index i is stored at index 2i+2
- the parent of the element at index i is stored at index (i-1)/2 (using int division)

Example heap:



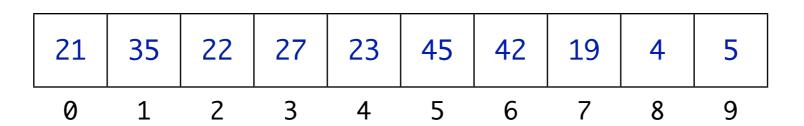


The first step in heapsort is to make the array into a heap (surprise!)

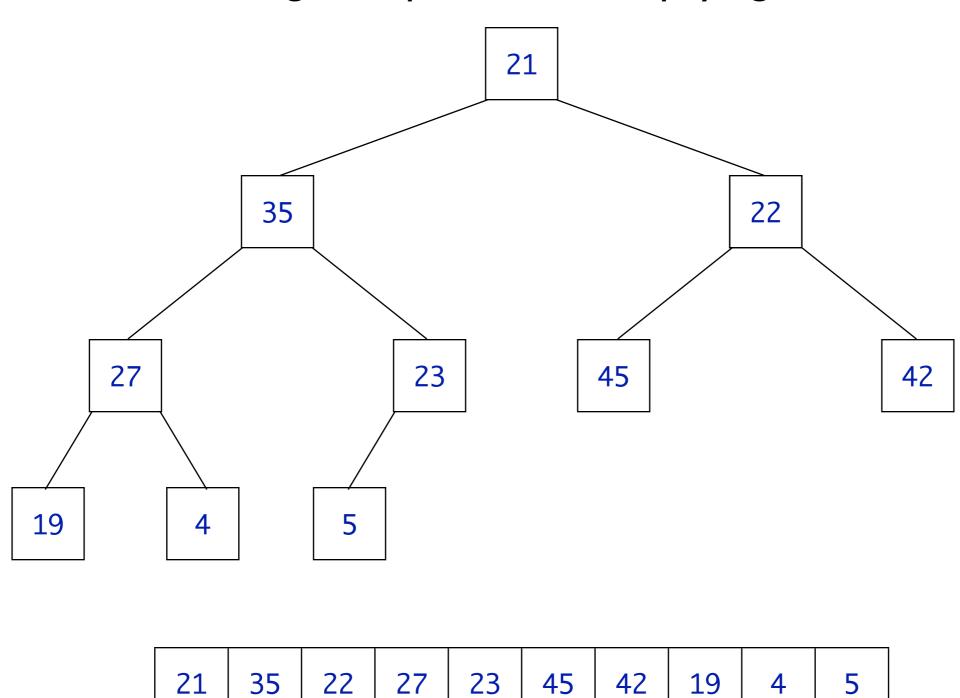
- the value contained in a node is always larger than that of either of its children
- it must be treated as a complete binary tree (easy, given our representation)

Consider the following array:

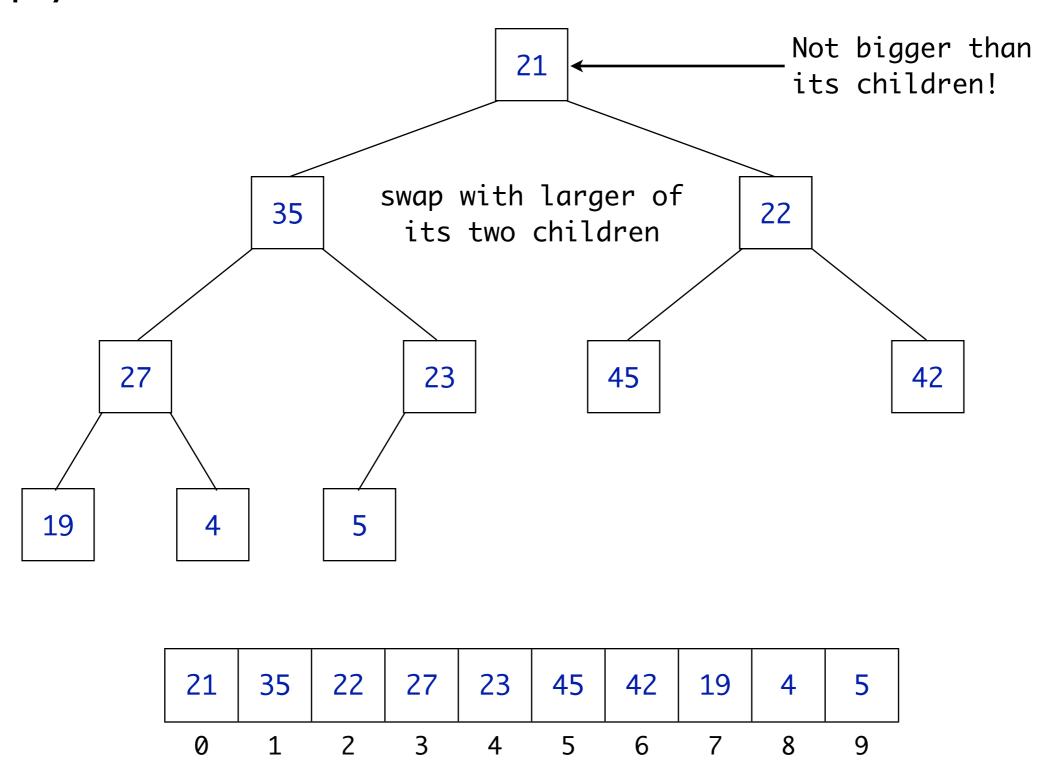
- it's not yet a heap, but—like we've seen—it can be treated as a complete binary tree
- we need to heapify the tree so that each node is larger than its children



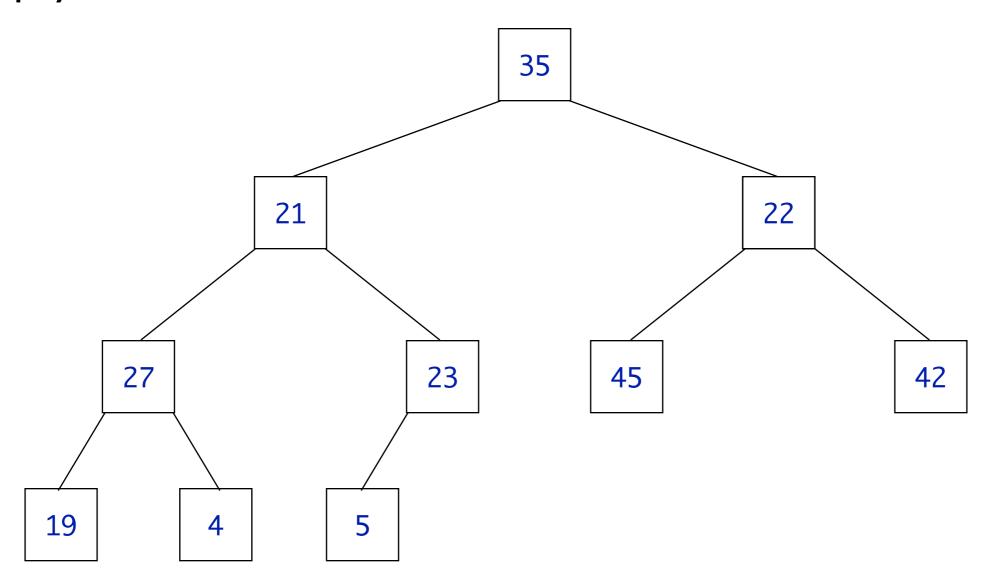
The process of making a heap is called "heapifying"—clever, huh?

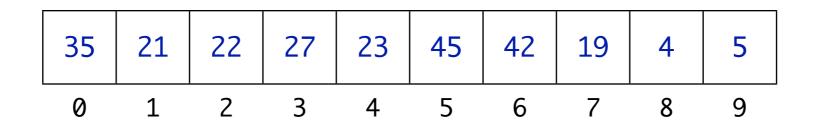


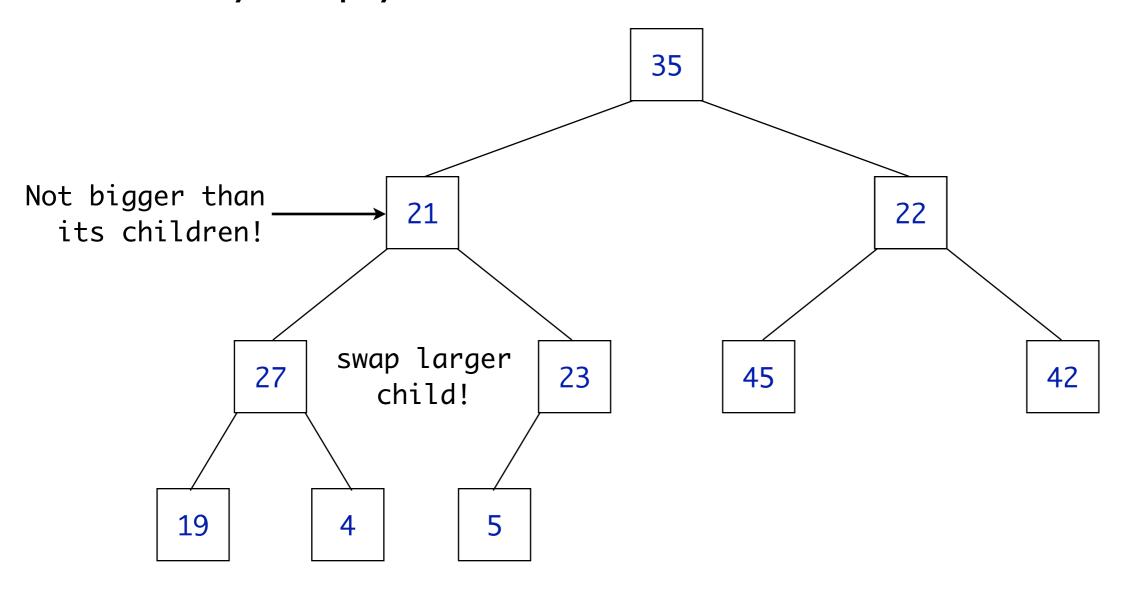
"Heapify" the entire tree:

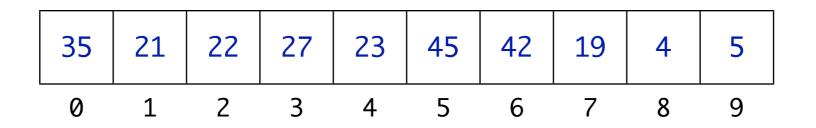


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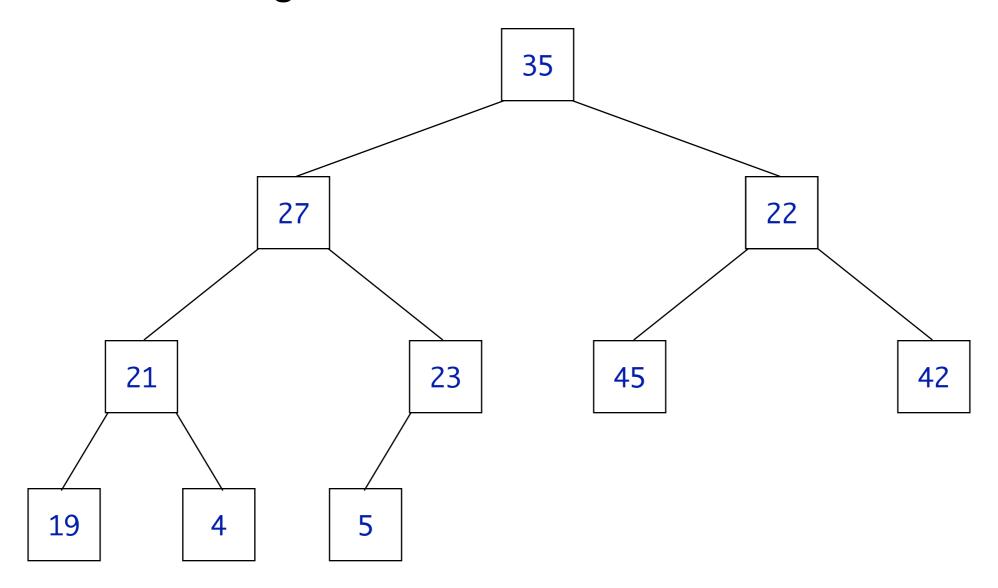


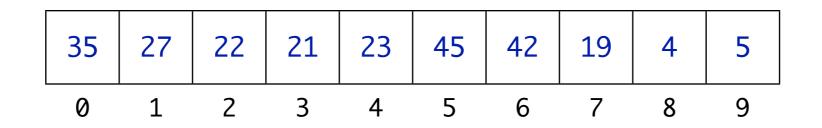


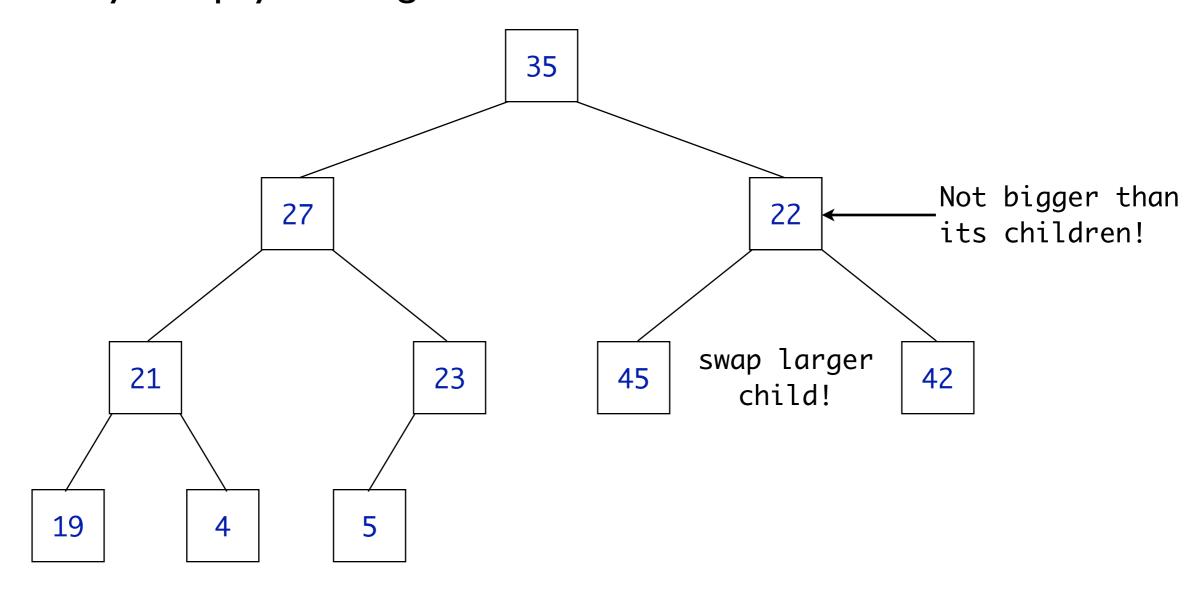


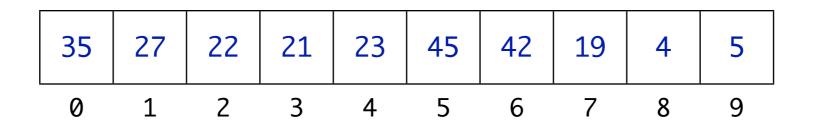


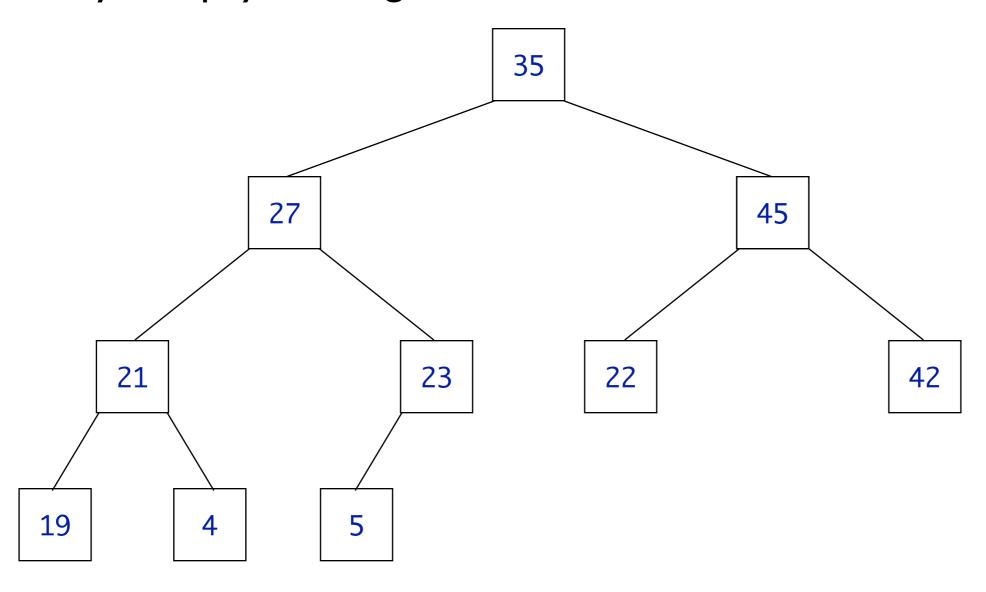
Left subtree is now good!

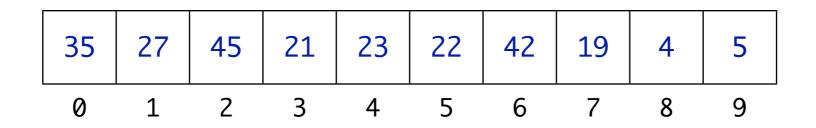


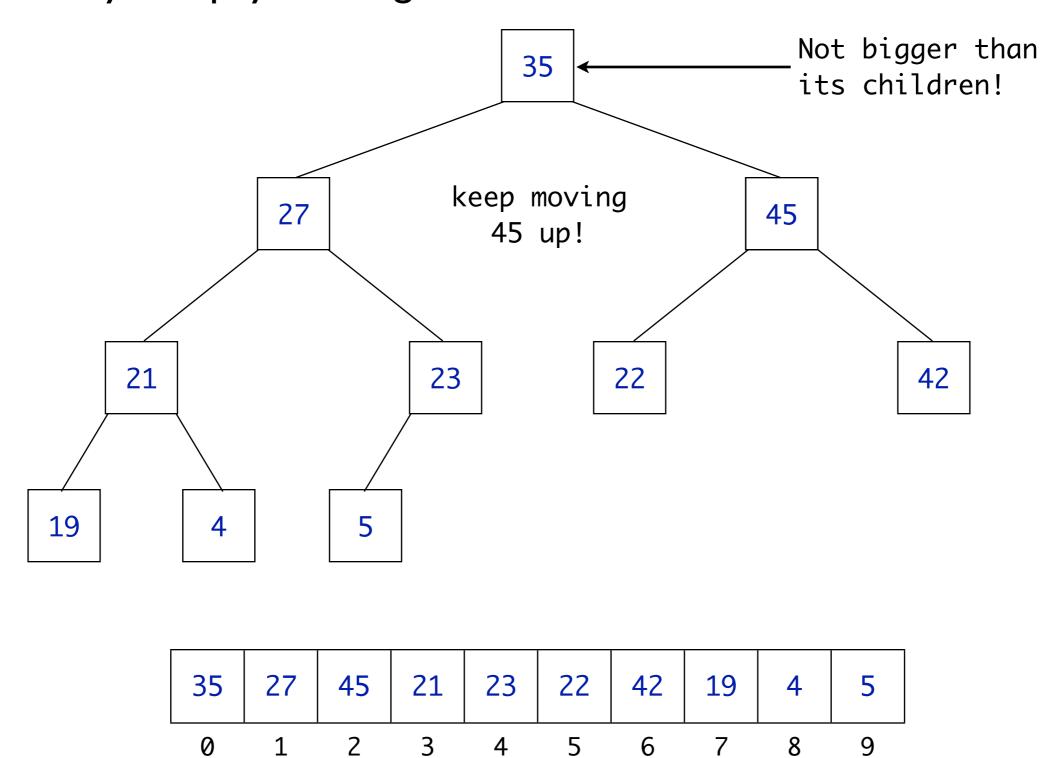


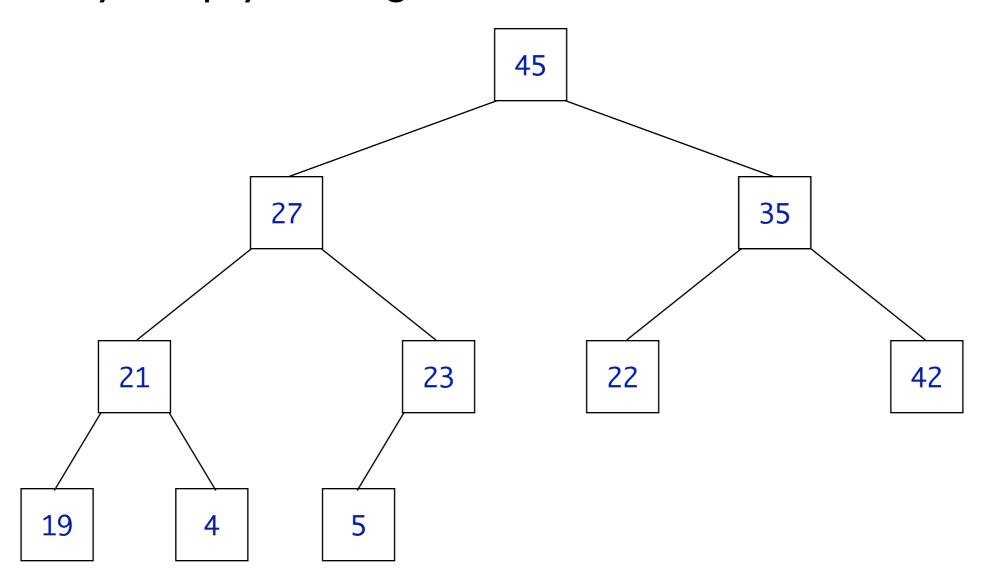


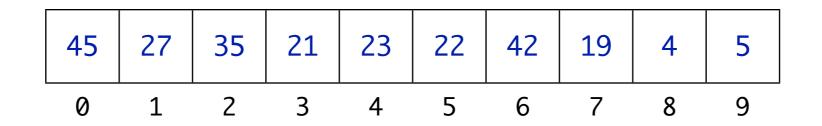


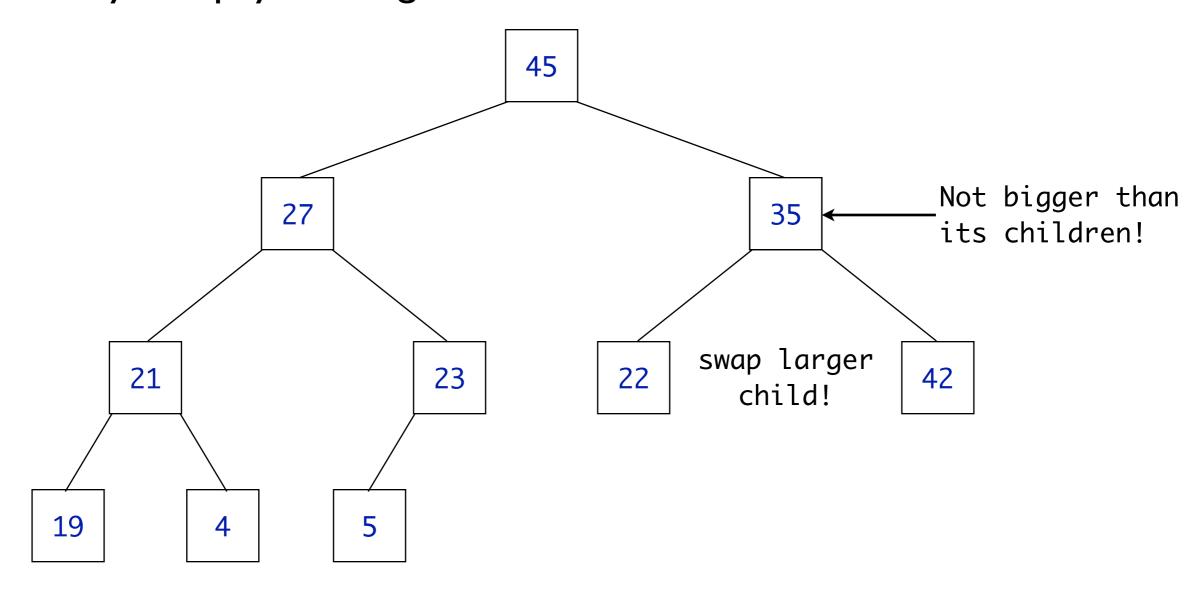


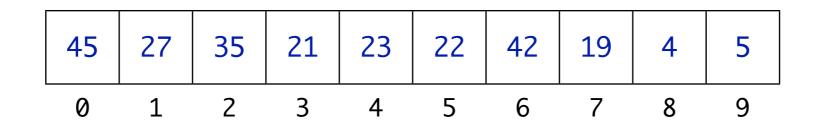




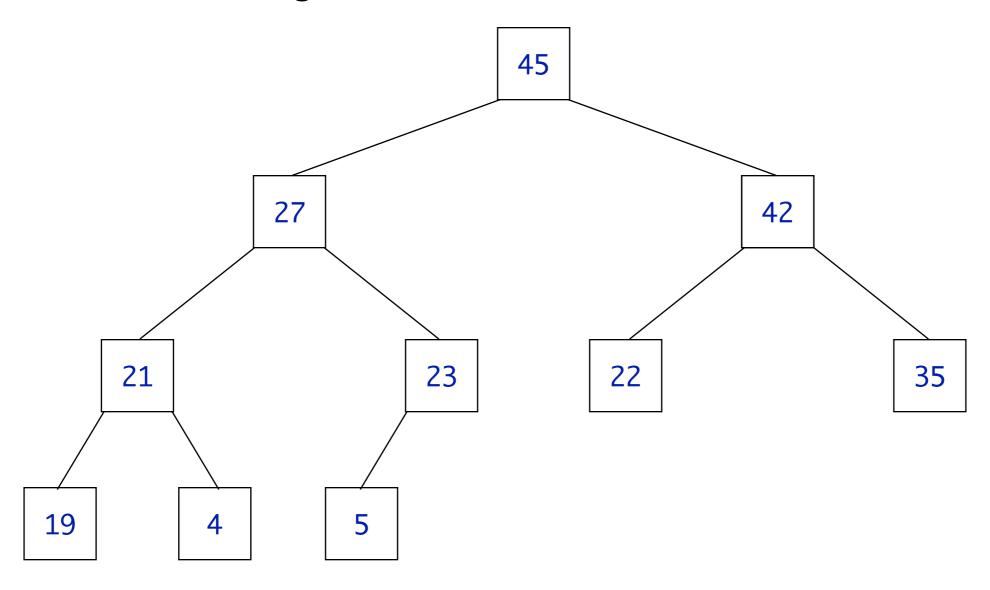


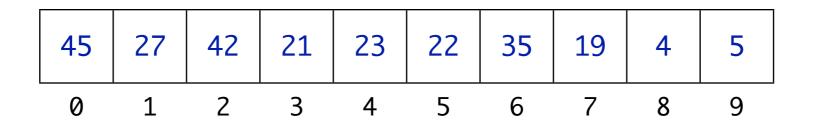






Right subtree is now good!



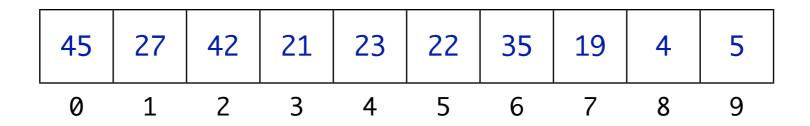


The first step in heapsort is to make the array into a heap (surprise!)

- the value contained in a node is always larger than that of either of its children
- it must be treated as a complete binary tree (easy, given our representation)

Consider the following array:

- now that it is a heap, we're ready to sort it...
- we know that the largest value is on the top, so we can put it in its correct place at the end of the array (via swapping)

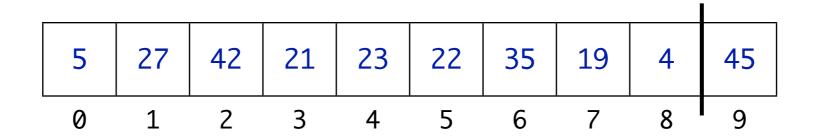


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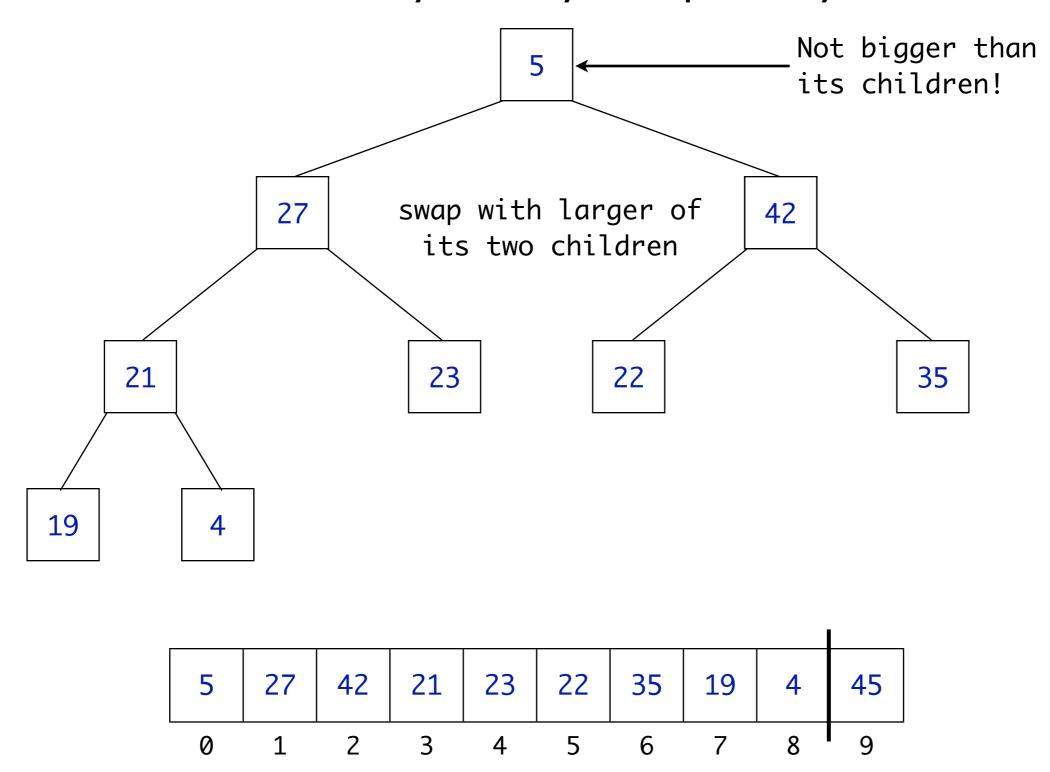
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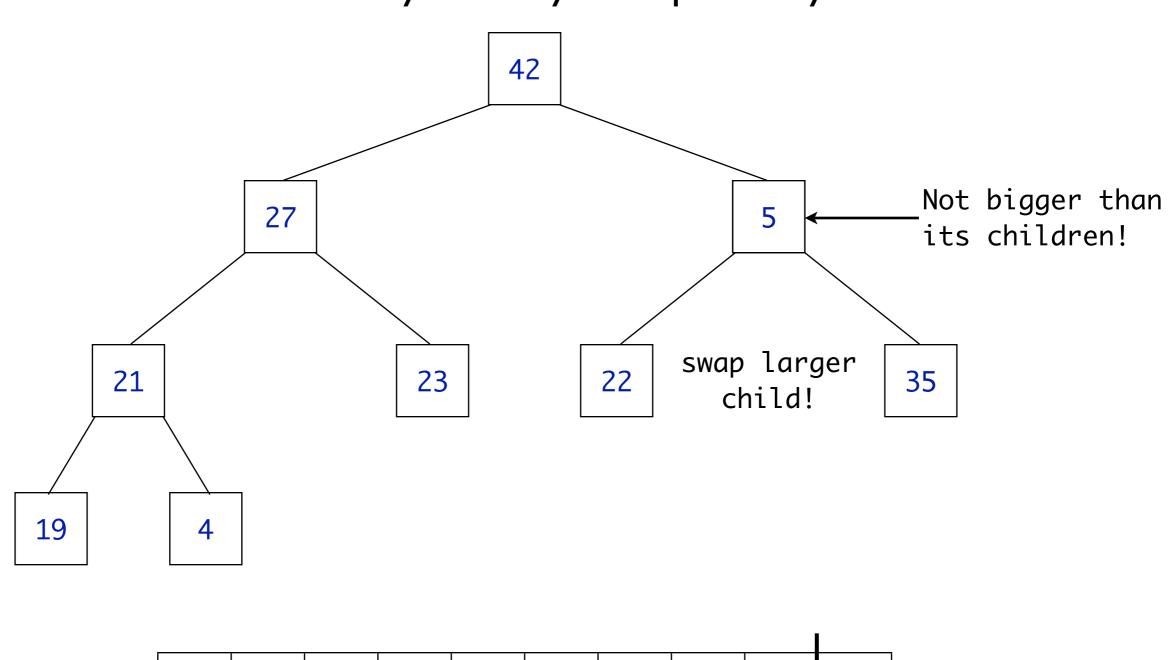
- now that it is a heap, we're ready to sort it...
- we know that the largest value is on the top, so we can put it in its correct place at the end of the array (via swapping)
- we now have a sorted array of size I (at the end)
- the unsorted array is down to 9 elements, and it's NEARLY A HEAP already!

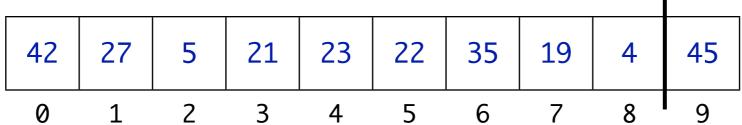


Nine-element unsorted array... Nearly a heap already!

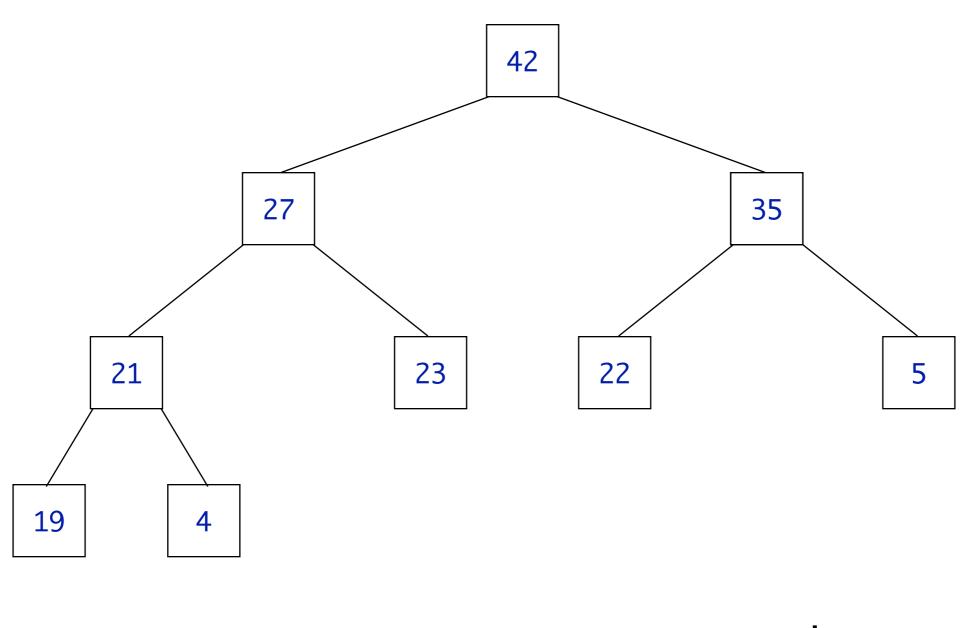


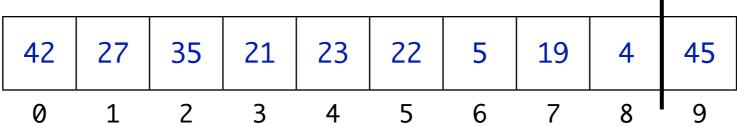
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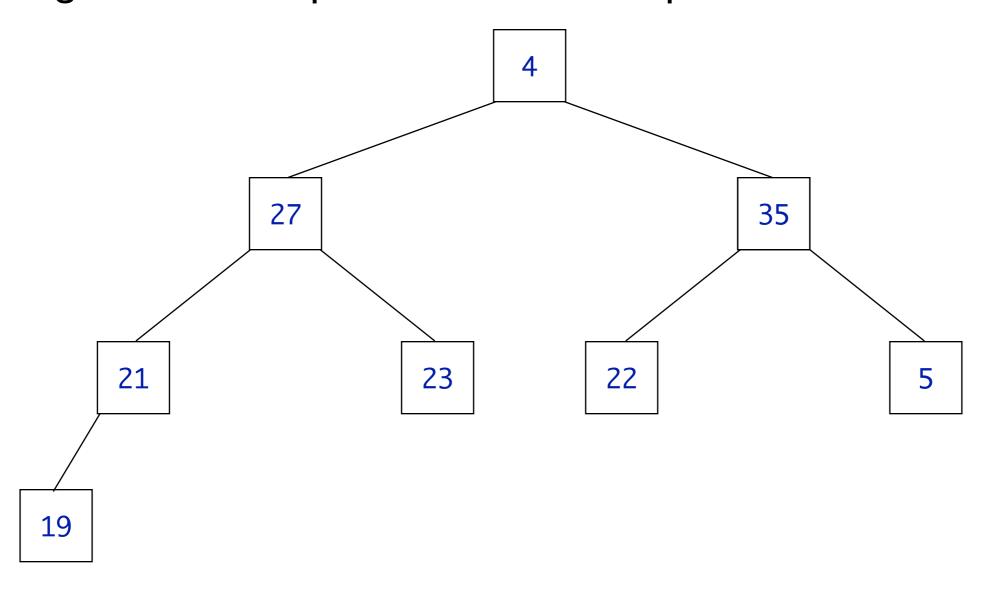


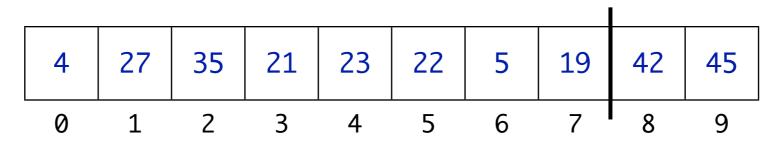
That was easy!

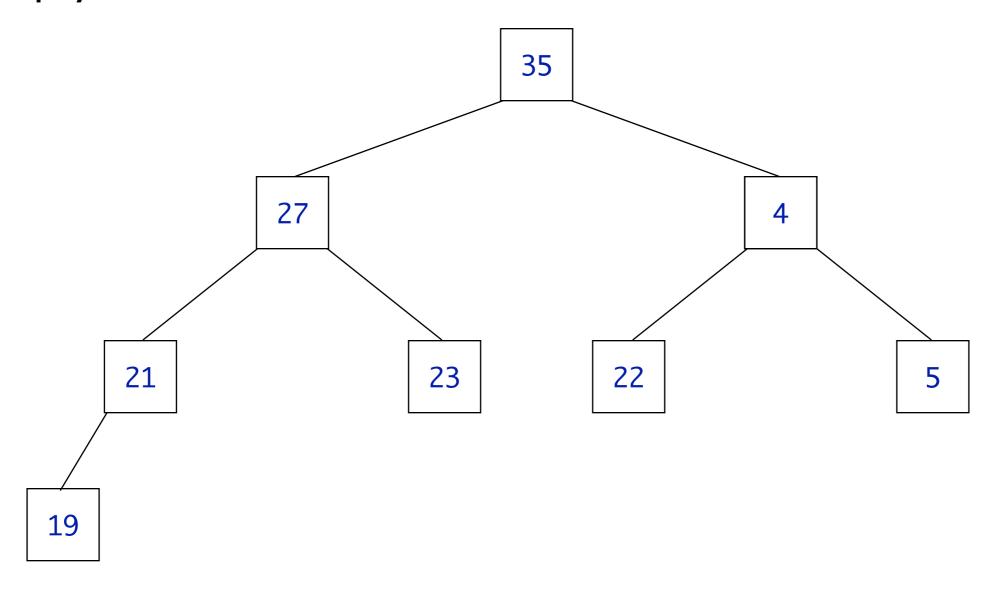


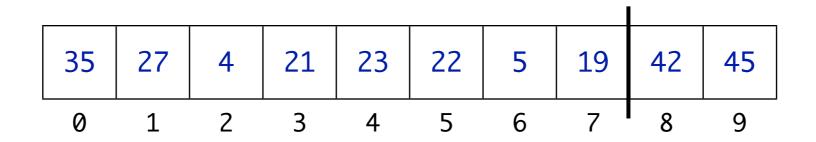


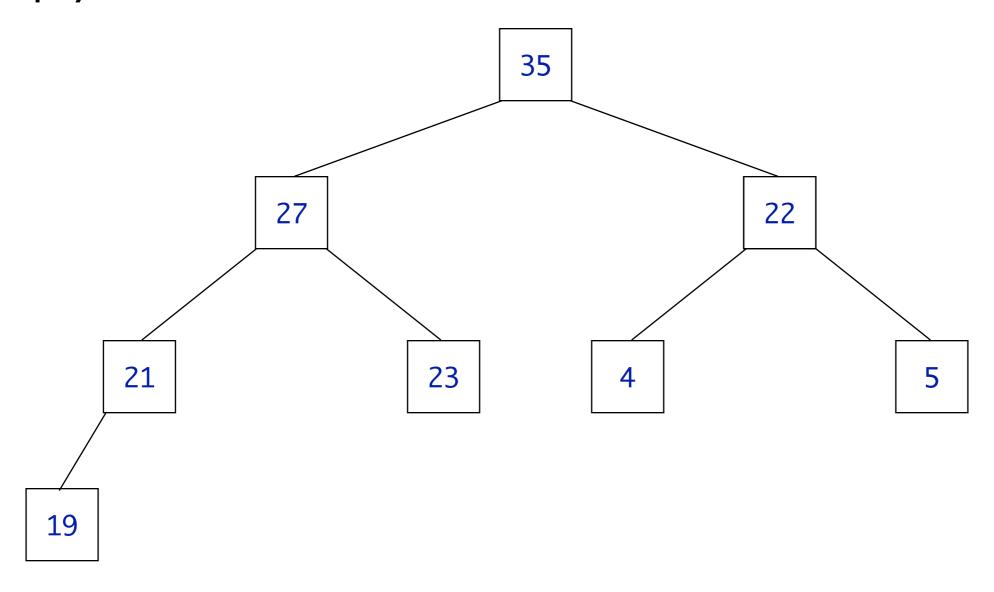
Take largest value and put it in its correct spot...

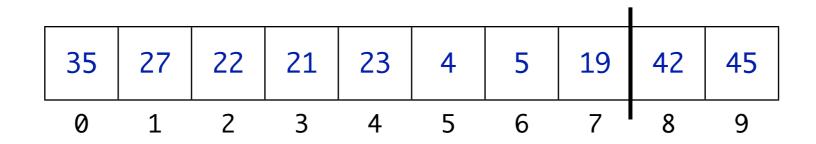




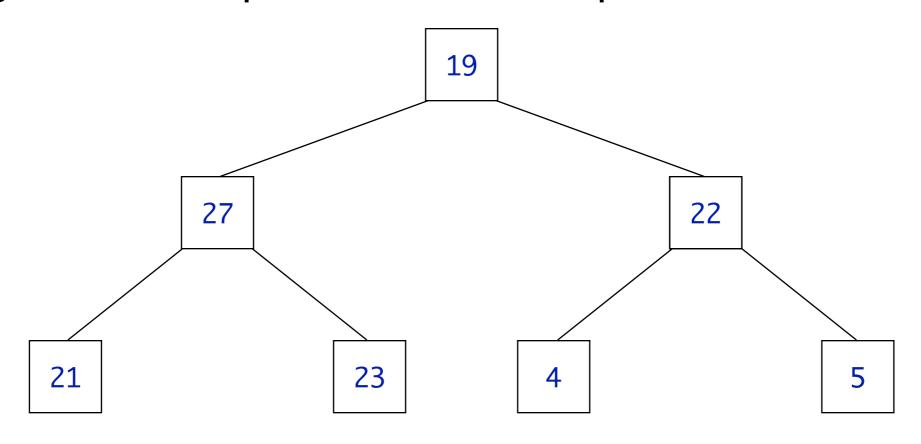


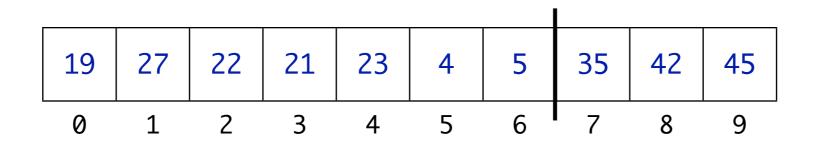


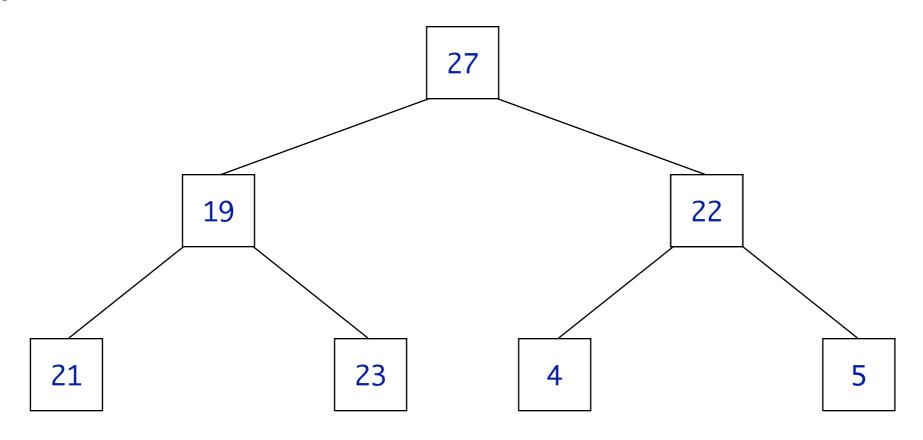


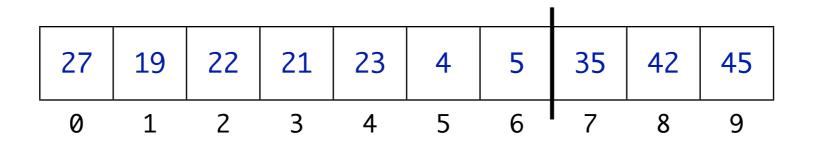


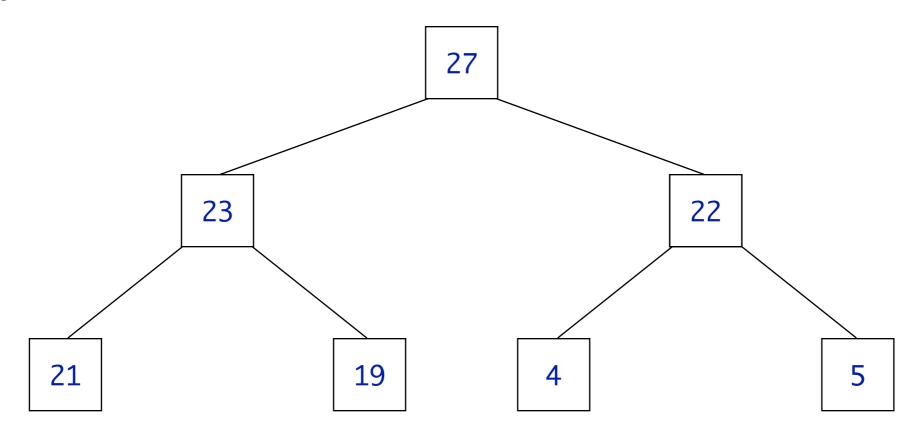
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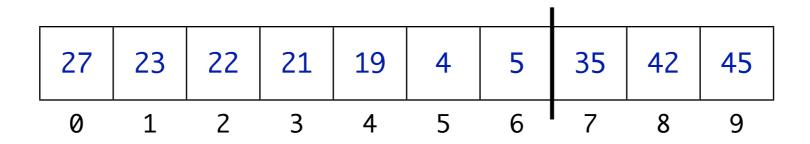




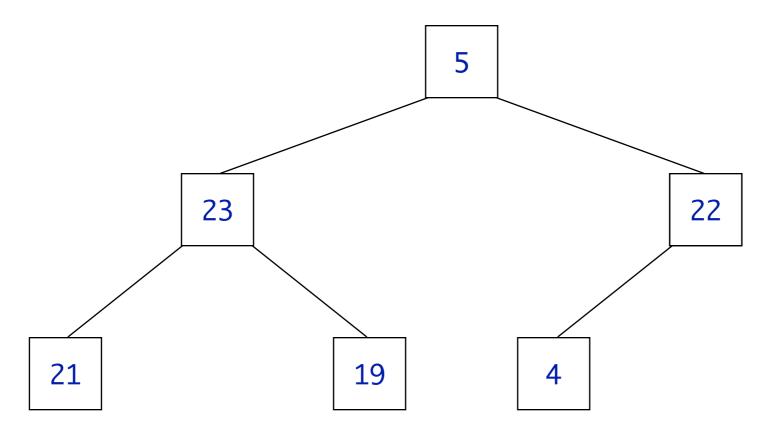


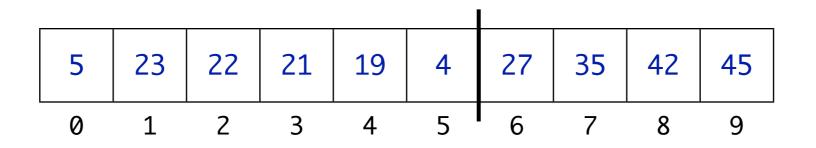


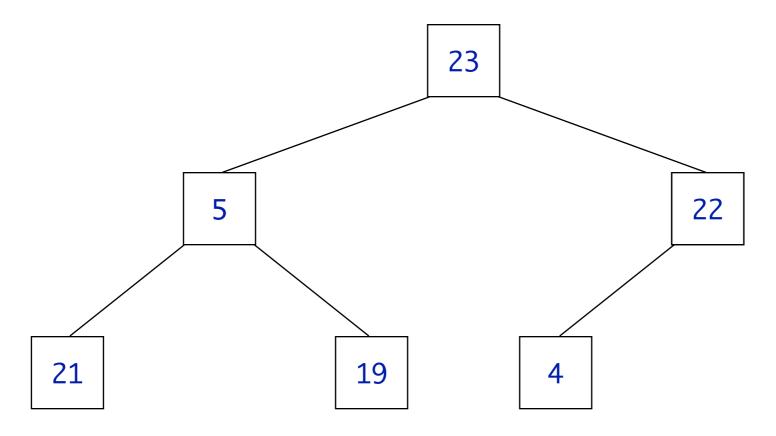


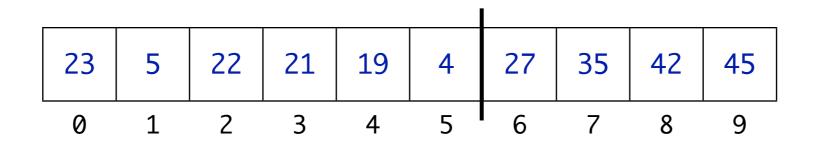


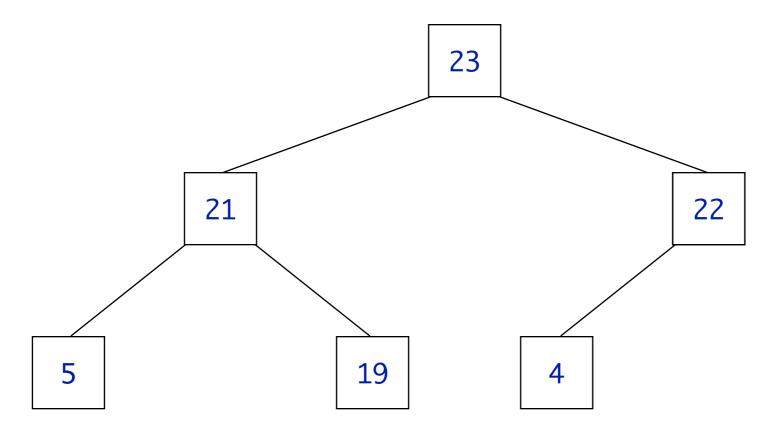
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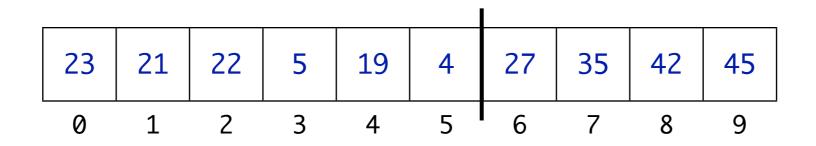




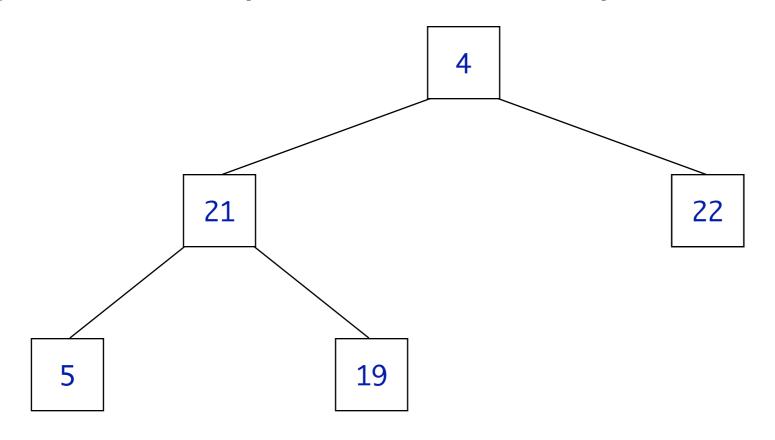


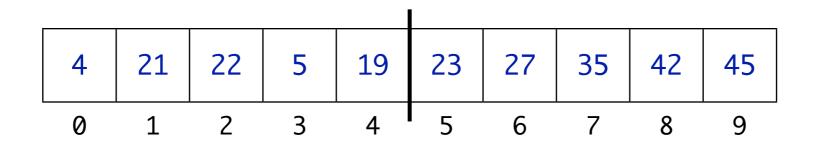


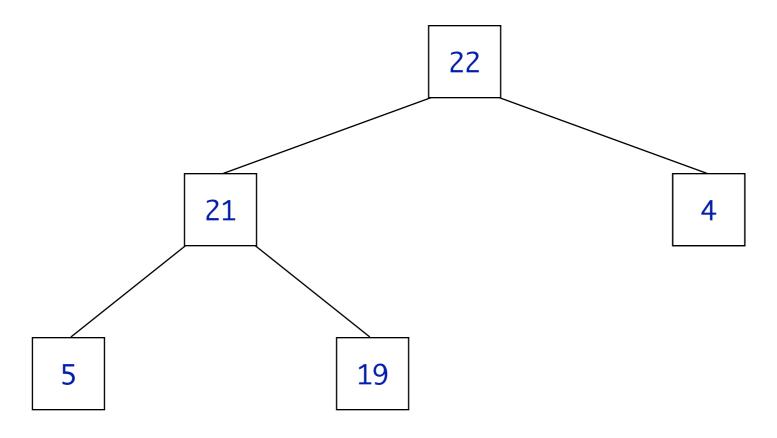


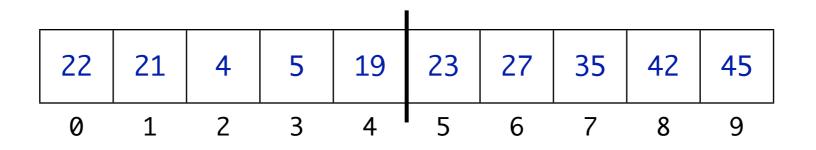


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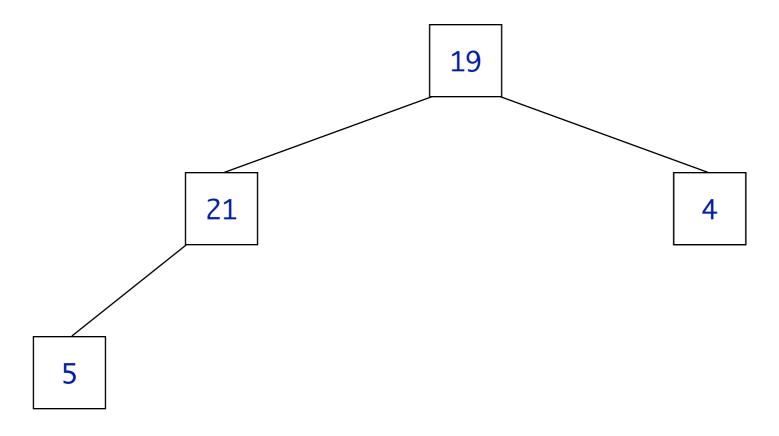


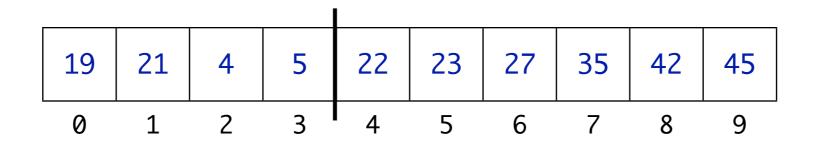


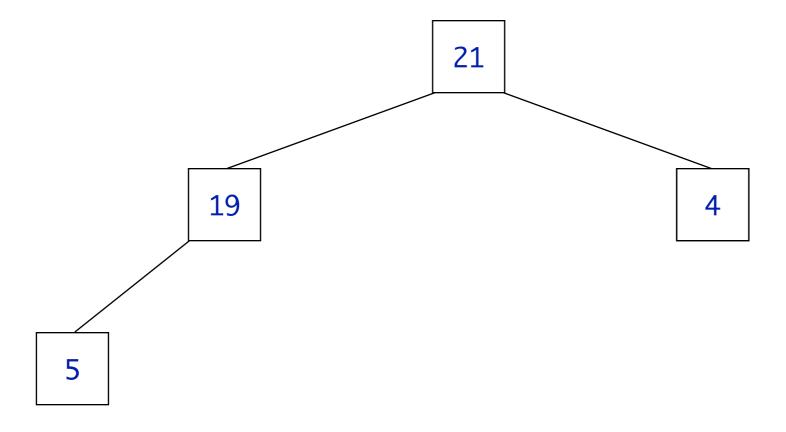


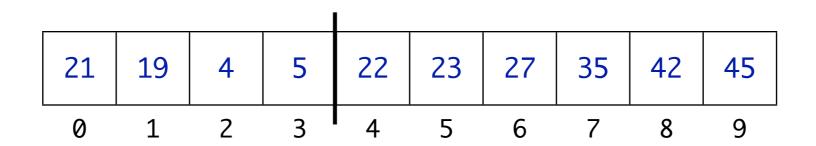


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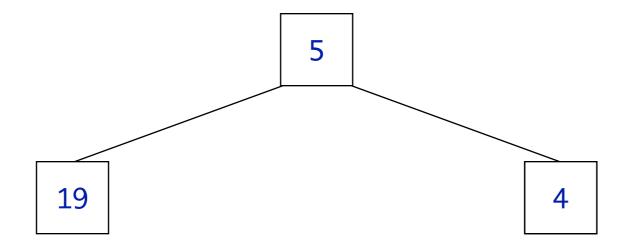


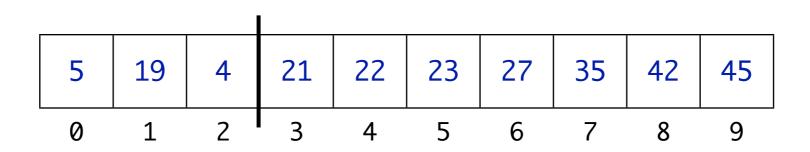


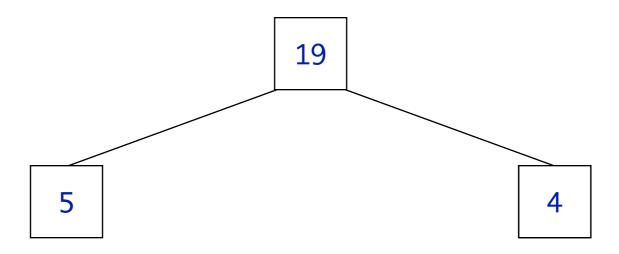


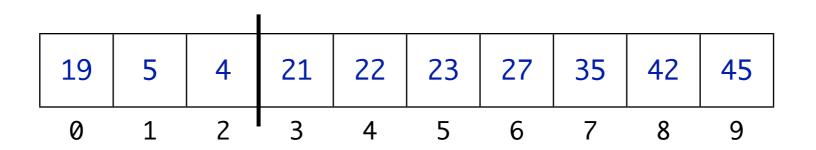


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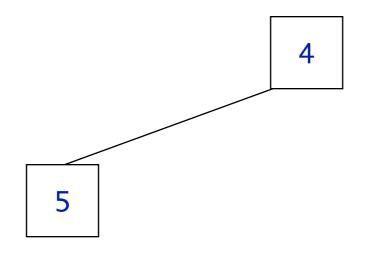


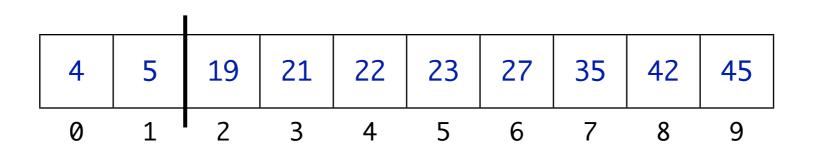


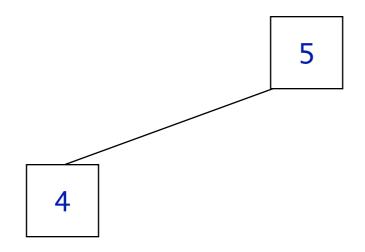


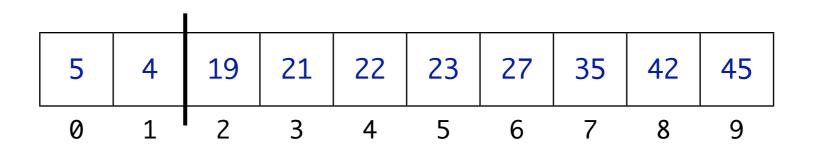


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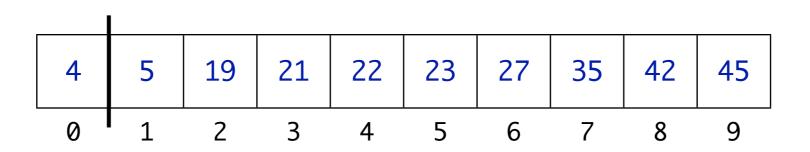




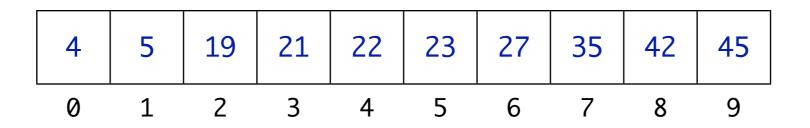


Take largest value and put it in its correct spot...

4



One value left...All done!



An implementation:

```
template <typename Item, typename SizeType>
void heapsort(Item data[], SizeType size) {
    make_heap(data, size);
    while (size-- > 1) {
        swap(data[0], data[size]);
        reheapify_down(data, size);
```

To make the heap:

```
template <typename Item, typename SizeType>
void make_heap(Item data[], SizeType size) {
    for (SizeType i = 1; i < size; i++) {</pre>
        SizeType k = i;
        while (k != 0 && data[k] > data[parent(k)]) {
            swap(data[k], data[parent(k)]);
            k = parent(k);
```