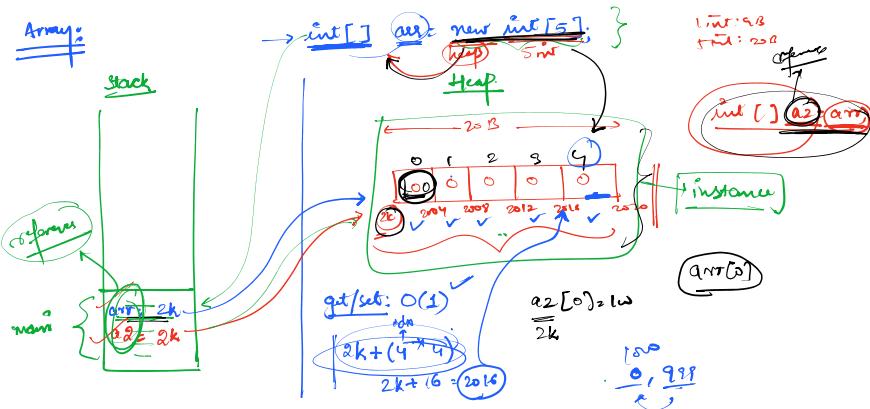


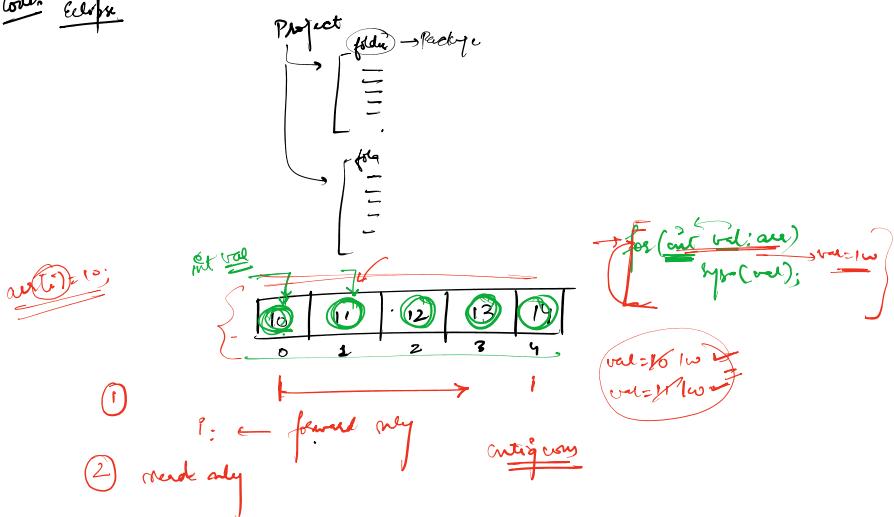
- Deviations
- Assignments
- Expectations



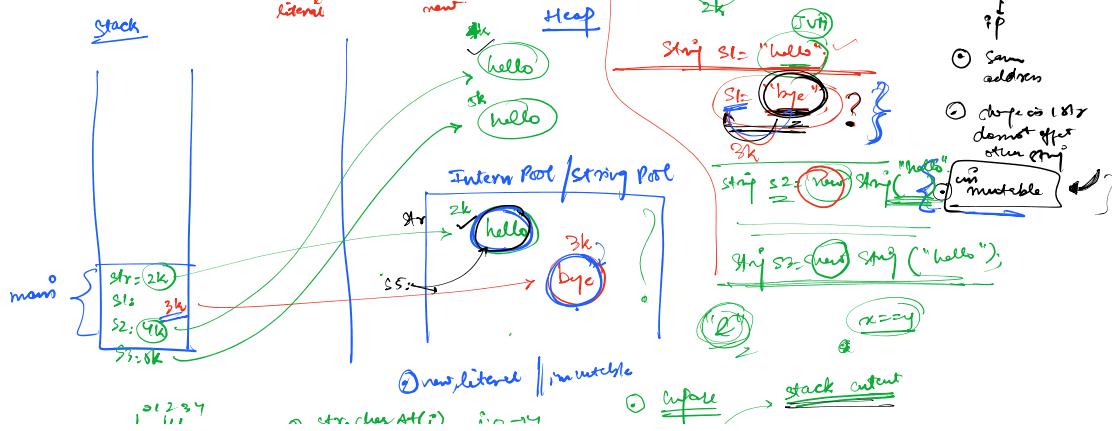
<u>Primitives:</u>	<u>References</u>
byte	
short	
int	
long	
double	
float	
char	
boolean	

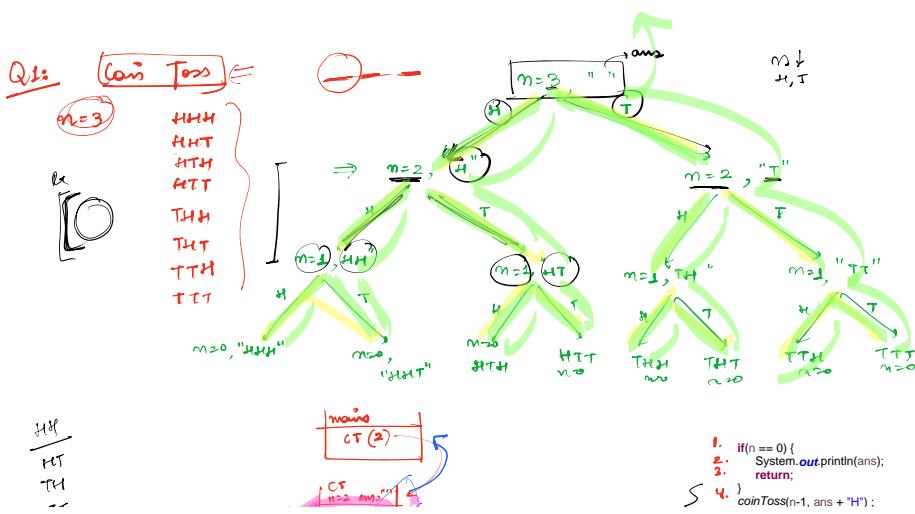
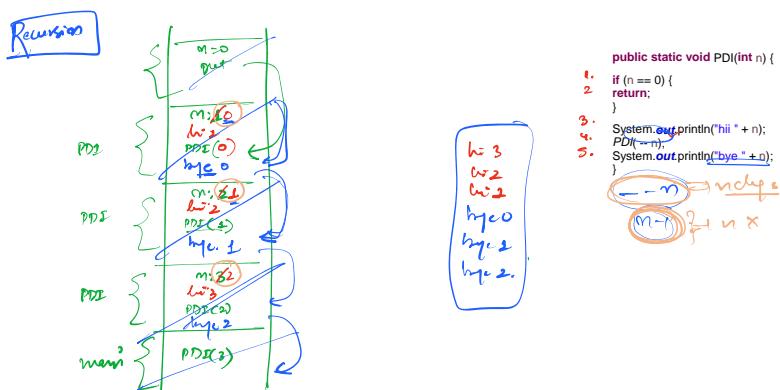
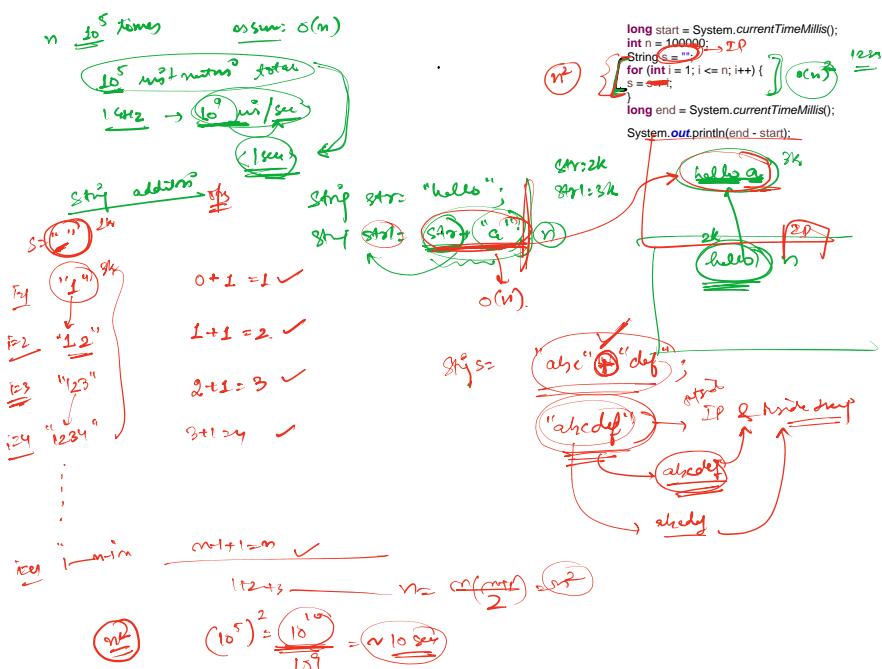
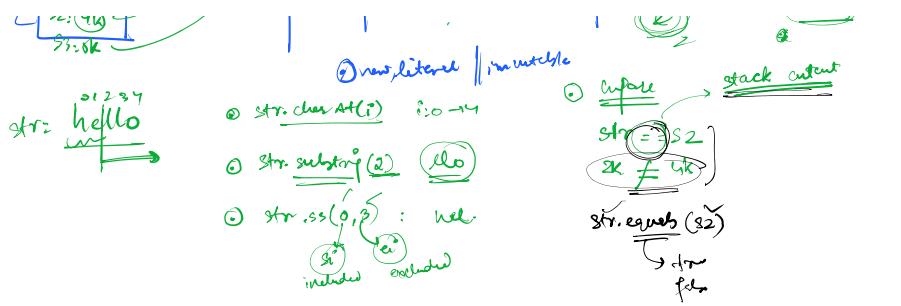


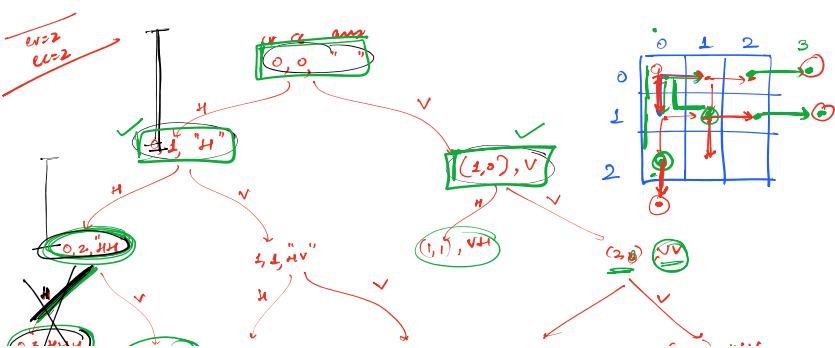
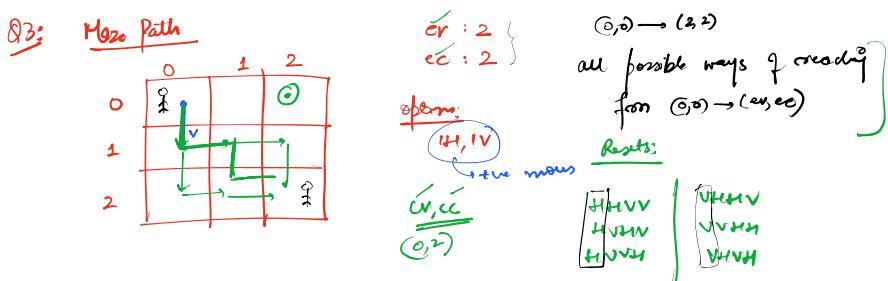
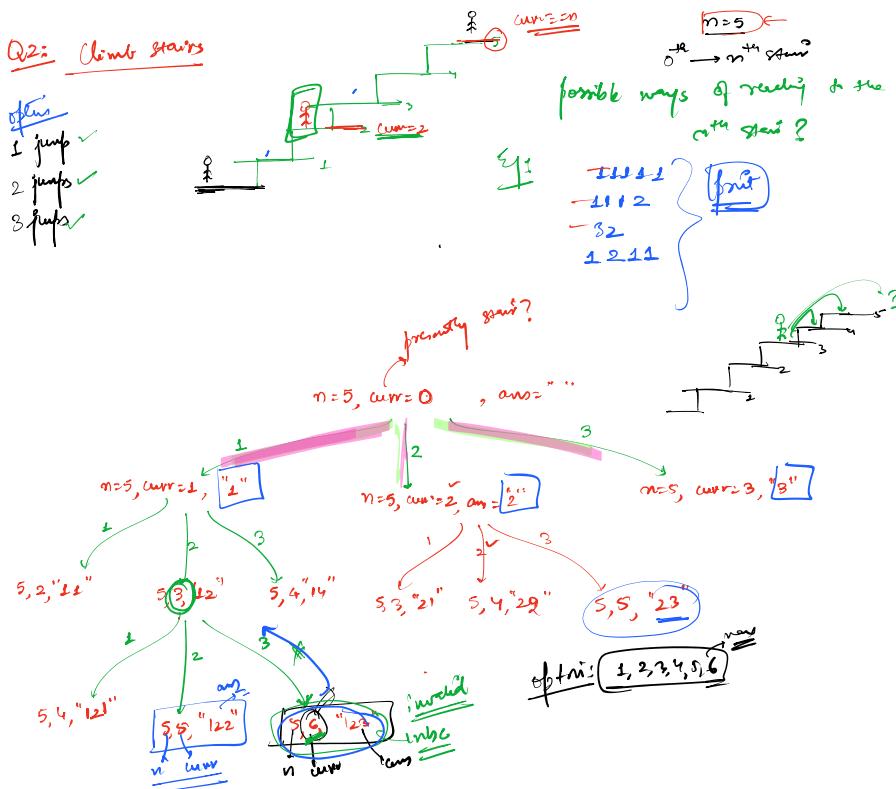
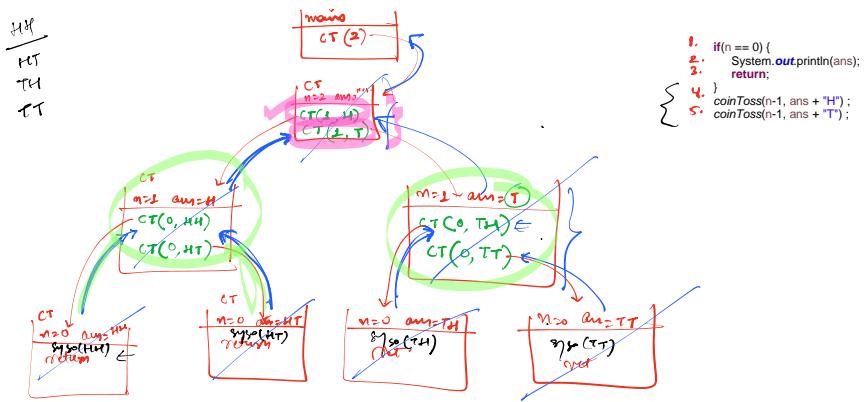
Looker Eclipses

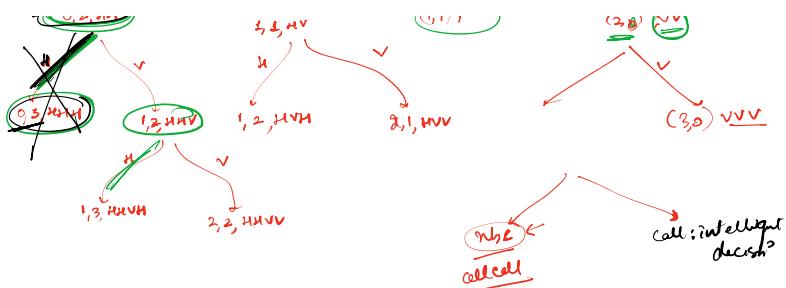


Strings :-









Q4: Valid Parenthesis

$$\boxed{m=3} \quad \begin{array}{l} \xrightarrow{\hspace{1cm}} 3 \text{ open brackets} \\ \downarrow \quad \quad \quad \downarrow \\ \xrightarrow{\hspace{1cm}} 3 \text{ close brackets} \end{array}$$

Brackets: ()

82x
57x

$$\frac{m=2}{(())}$$

100

乙

The diagram shows five configurations of three parentheses, each enclosed in a red bracket. The first four configurations are valid, indicated by a red checkmark above them. The fifth configuration is invalid, indicated by a red slash through it.

- $((()))$
- $(())()$
- $((())()$
- $(()())$
- $(())()$

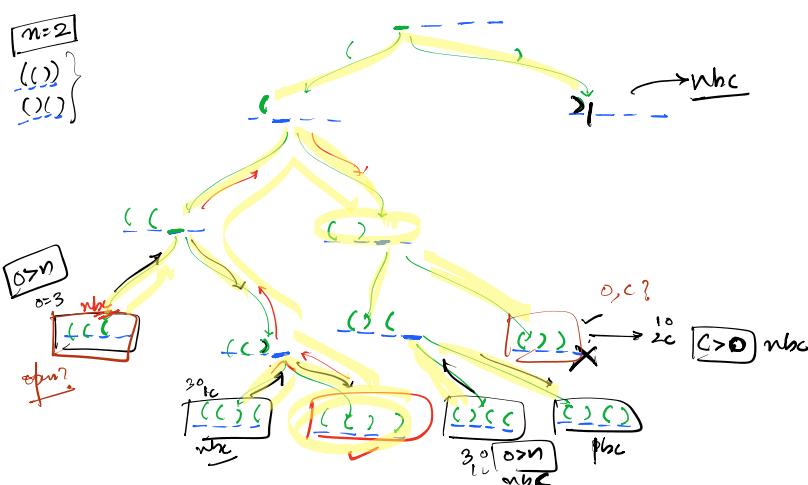
valid configuration
invalid!

5(1)
5(2)

6

$n=2$

()) }
 () () }



Homework:

1) Assignment: s1B ✓
2) Quiz: Recurs° ✓

Backtracking =

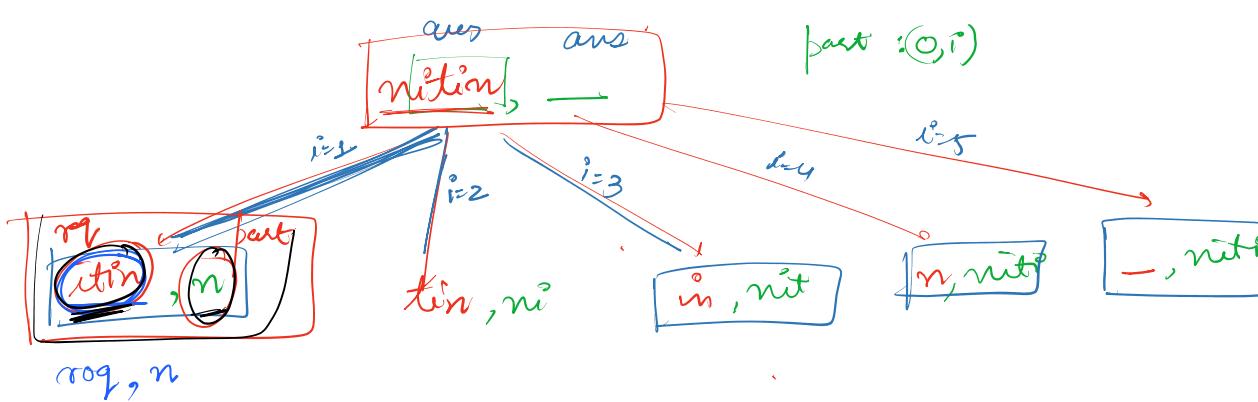
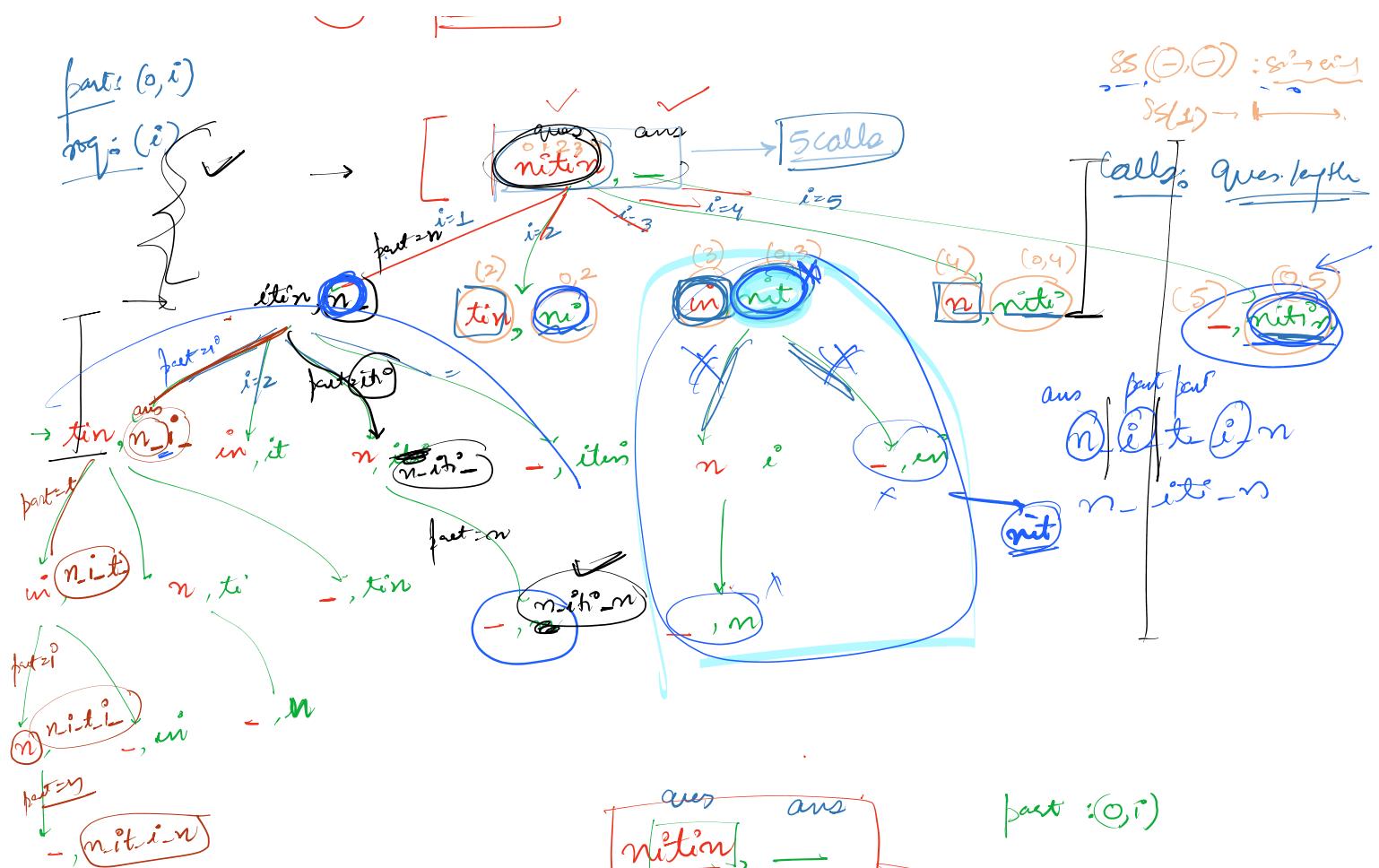
fact: $(0, i)$

① nitrogen
parts: five lone

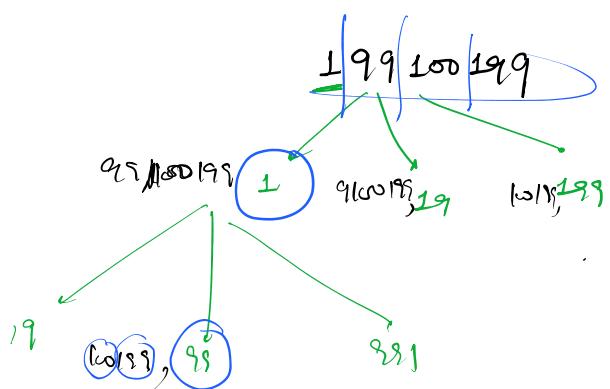
- 2 beat
- iti beat
- n beat

3 nitrogen

$$ss(\theta, \theta) : s^2 \rightarrow e^{i\gamma}$$



1 99 100 199 3

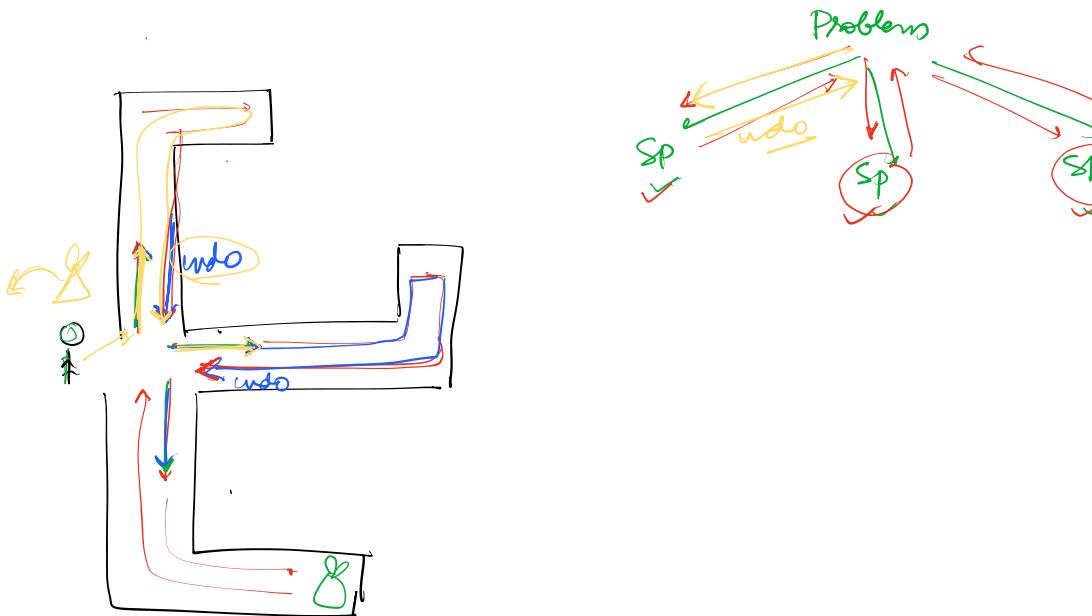


Backtracking:

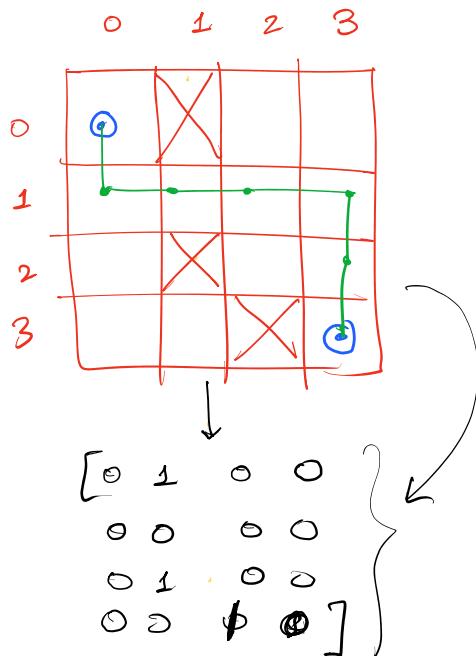
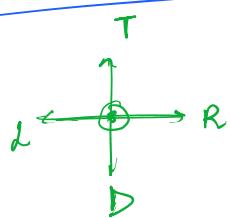
Problems



Backtracking:

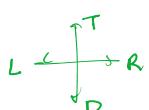


① Blocked Maze

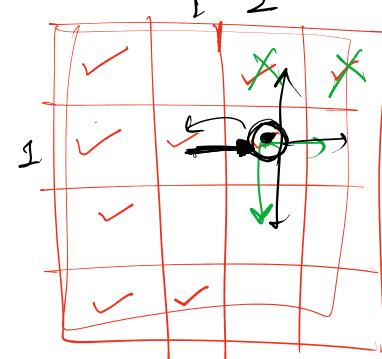
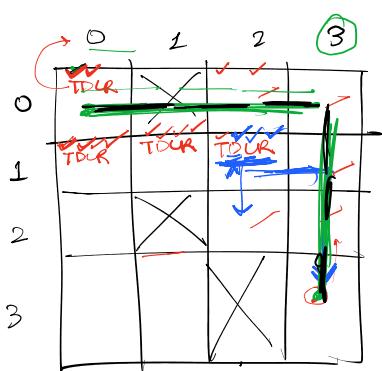


$(0,0) \rightarrow (3,3)$
ways?

DRRTRPDD
DRRDRD
DRRR DD



TDLR



Recursion Tree
Calls

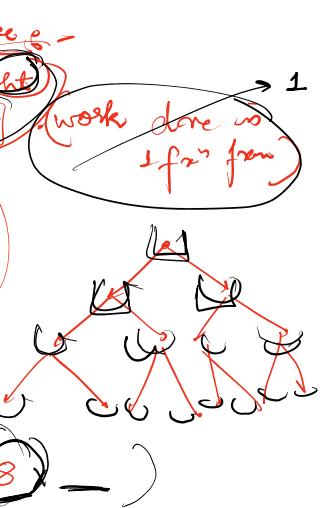
Boolean visited

Sudoku
solver

0 1 2 3 4 5 6 7 8 9

T	(0)	3	6	5	8	4			
	1	5	2						

1 (0,1)



$0,0$
 $(0,1)$
set

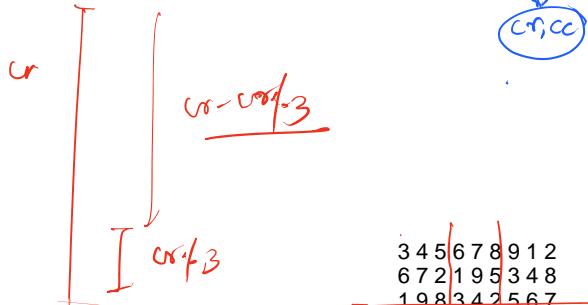
For $\underline{CC+1}$

Solution:-

cr = 7

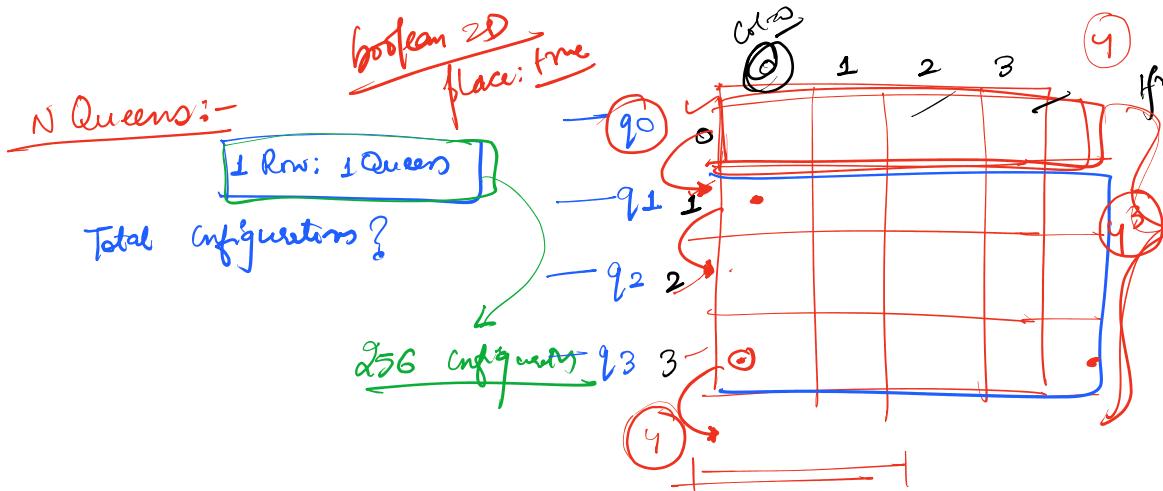
1 → 9

7 → 3



Y:
1 chocolates
3 students

3	4	5	6	7	8	9	1	2
6	7	2	1	9	5	3	4	8
1	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9	2	4	8	5	6
9	6	1	5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
5	3	4	2	8	6	1	7	9



kill

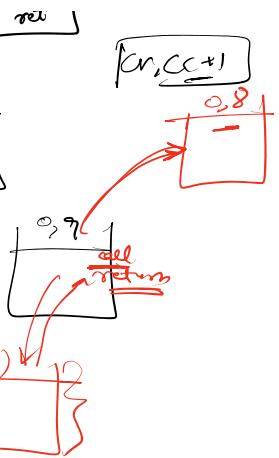
cr = 4

0 1 2 3

0 1 2 3

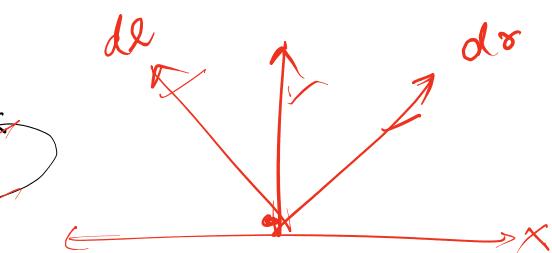
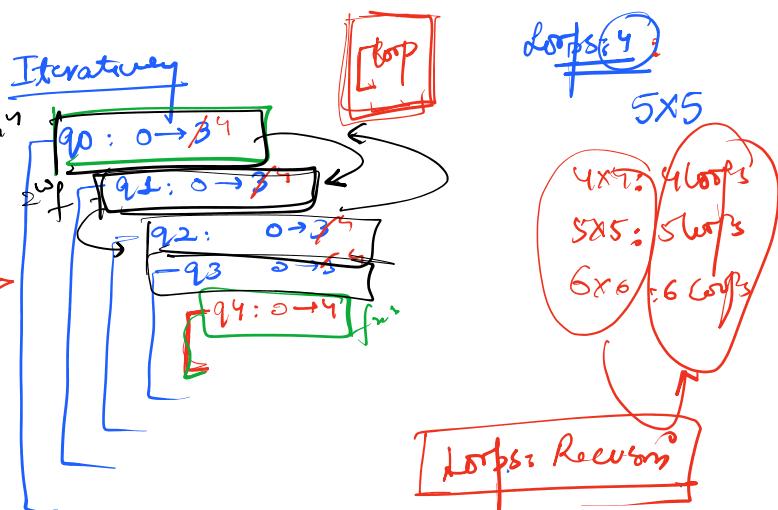
0 1 2 3

0 1 2 3

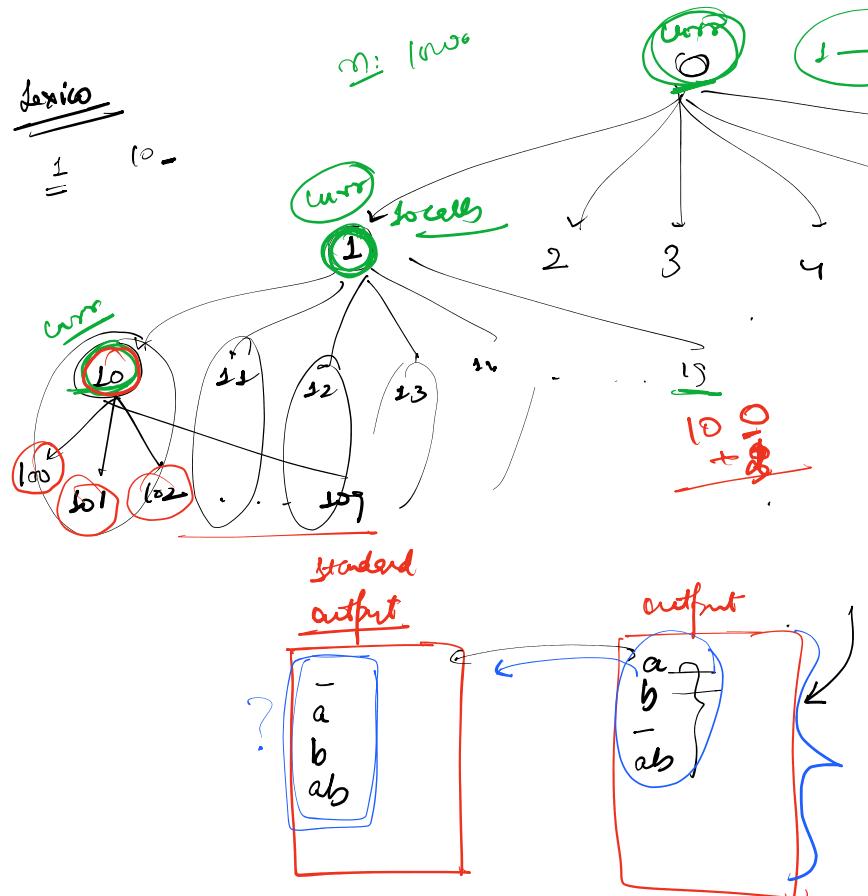


$$\frac{1}{3} = (2) \ominus$$

$$1 / 3 = 1$$



- \rightarrow Quiz: 5:30 pm; 20 mins. } \rightarrow MCQ
 \rightarrow Assignments: Thursday (DP)
 \rightarrow Thursday: 12:30 pm; DP

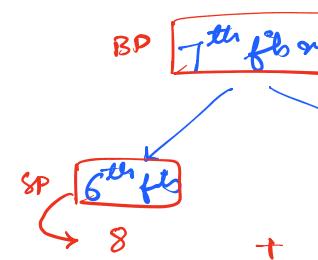


Dynamic Programming

Fibonacci

$$\begin{array}{ccccccccc} 0^{\text{th}} & 1^{\text{st}} & 2^{\text{nd}} & 3^{\text{rd}} & 4^{\text{th}} & 5^{\text{th}} & 6^{\text{th}} \\ \hline 0, & 1, & 1, & 2, & 3, & 5, & 8 \end{array}$$

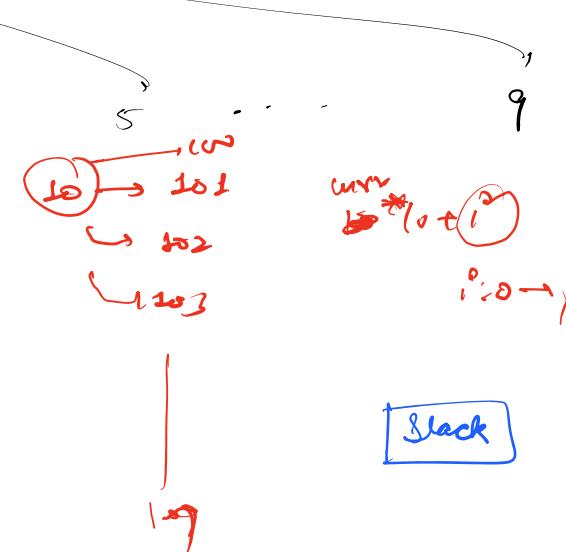
Recurse:



Slack



→ 9



$$\boxed{13} \xrightarrow{?} n : n^{\text{th}} \text{ fib no?}$$

\downarrow

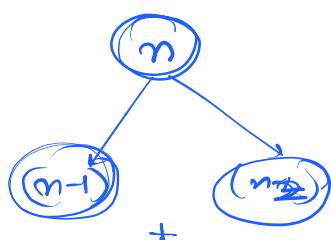
$$7 : 13$$

?

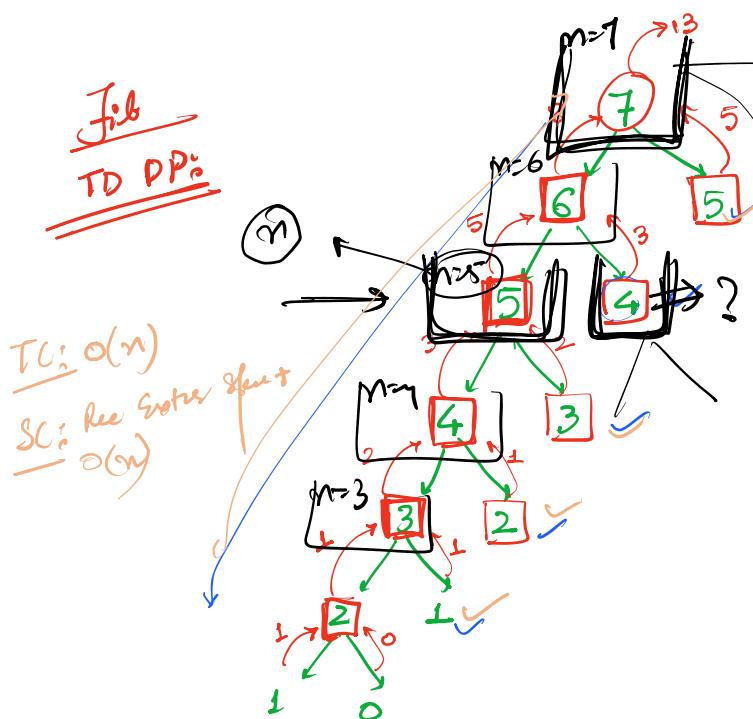
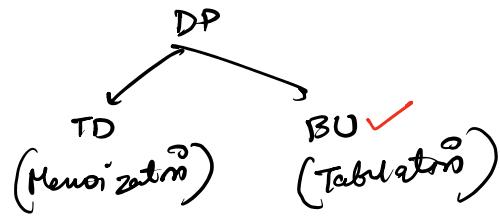
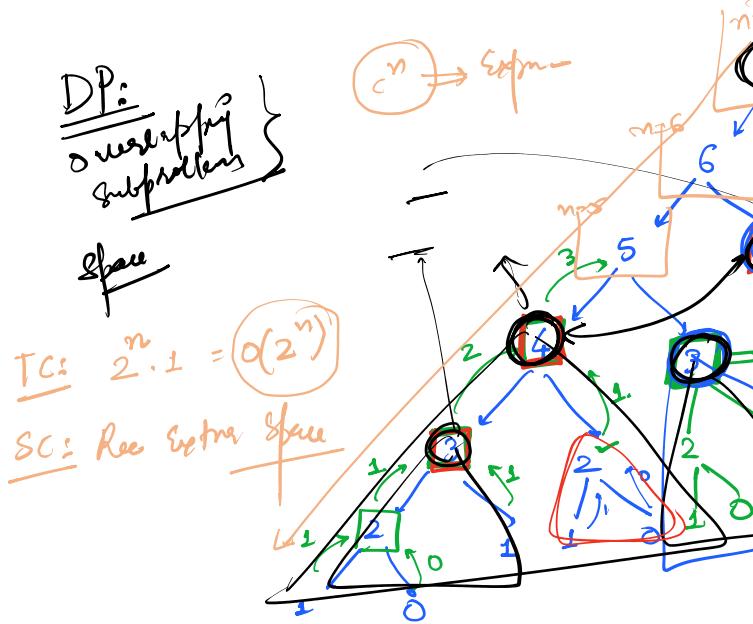
n^{th} ?

$\star \boxed{5 \text{ fib}}$

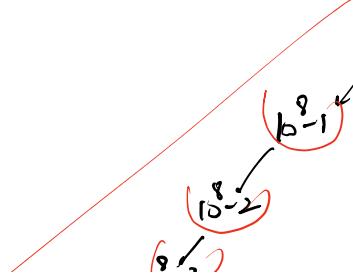
$5 = 13$

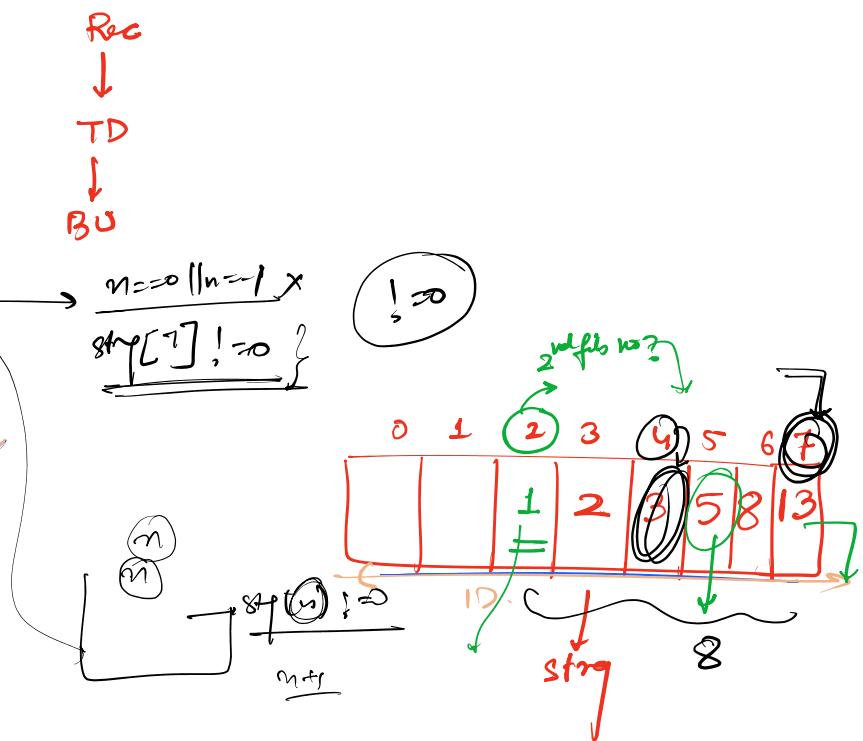
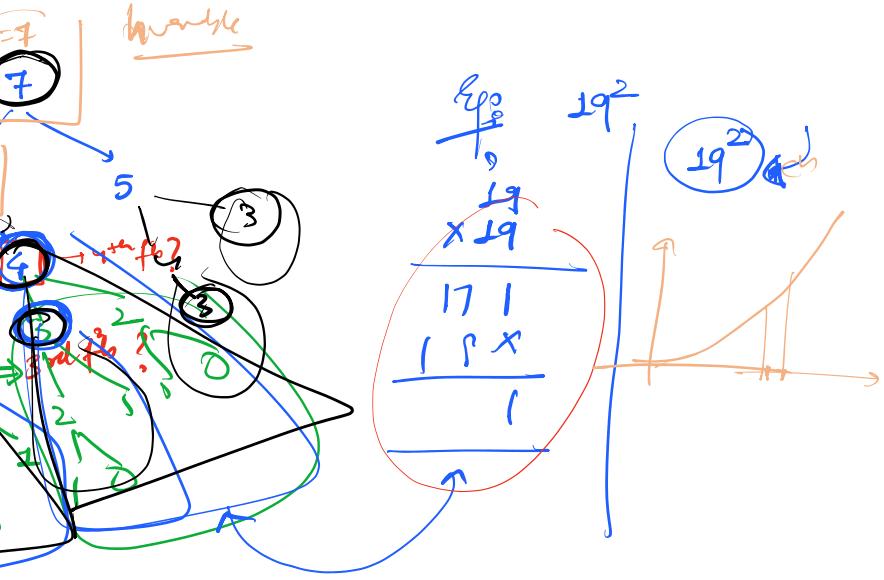


≡ 1 | Answer



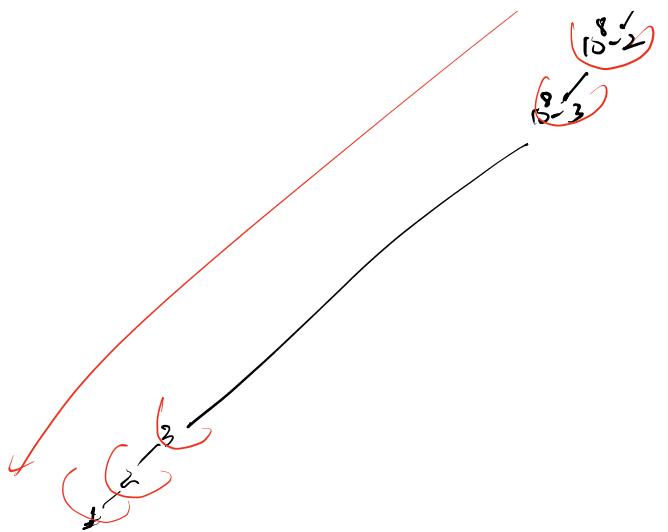
Properties





$$1 \times 10^8$$

$$n = \sqrt{10^8}$$



Bottom up

cell meaning? : TD ?

TC: $O(n)$
SC: $O(n)$

→ array size? : $n+1$

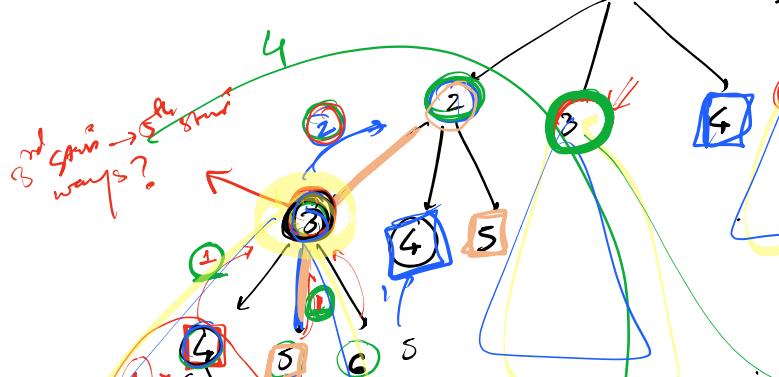
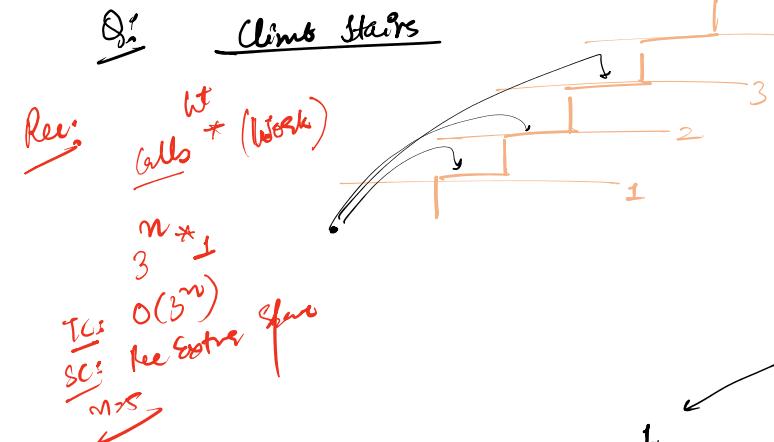
→ filling dir?

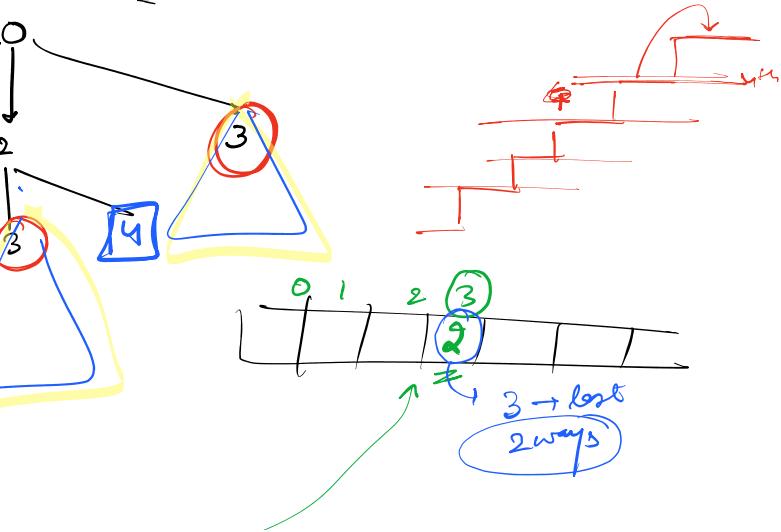
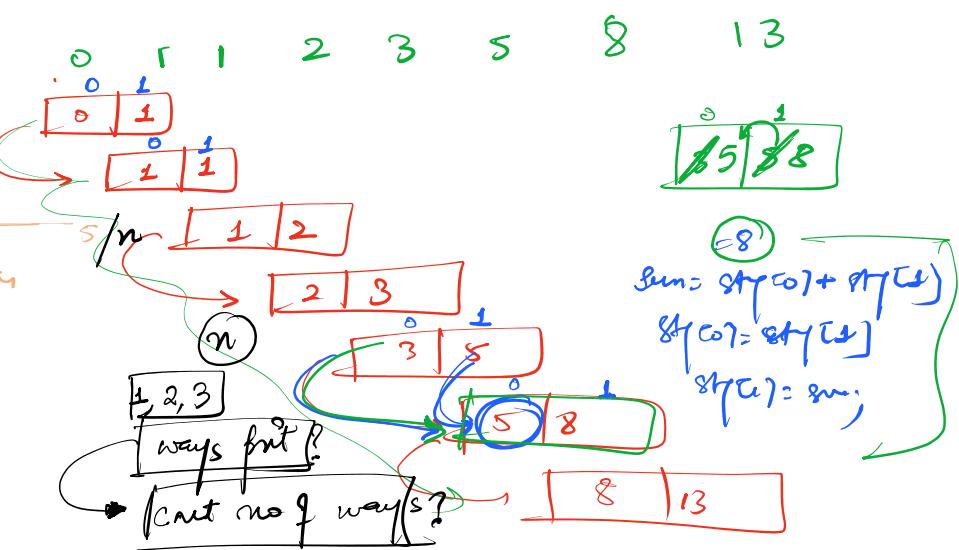
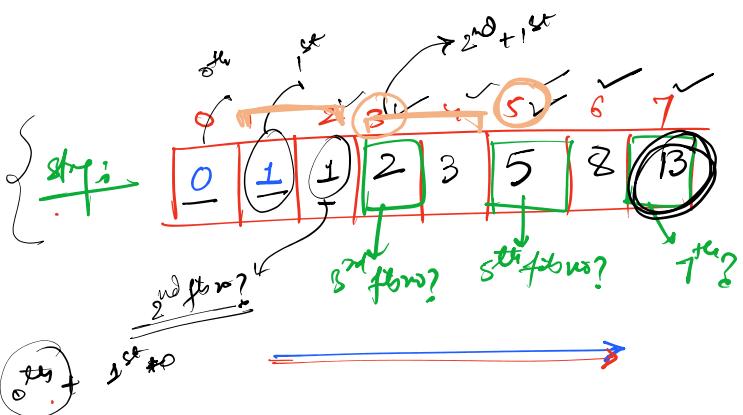
→ TD: BC ✓
BU: fib ✓

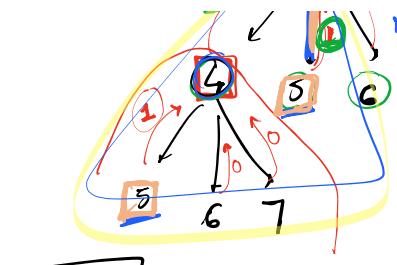
→ fill logic?

~~BU space effect~~

TC: $O(n)$
SC: $O(1)$

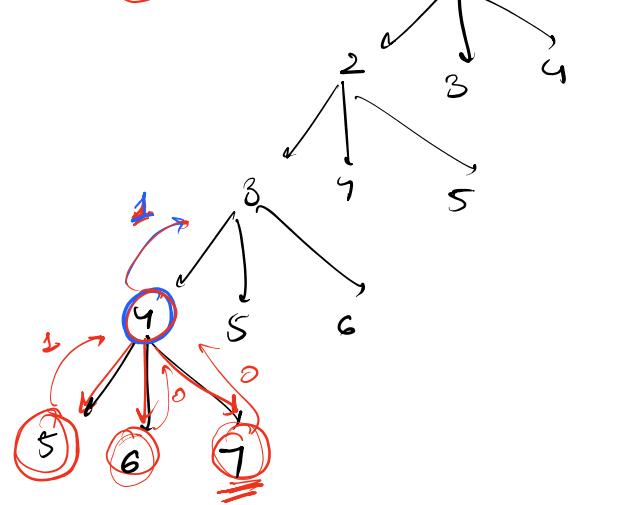






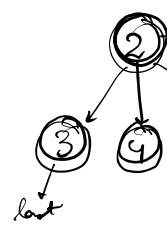
$m=5$

②



Bottom up:

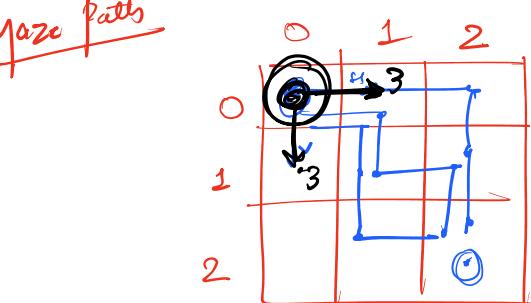
- cell meaning? no. of ways
- array size?
- filling direction?
- TD BC = BU fill
- fill logic?



BU \Rightarrow $SC: O(1)$

$TC: O(n^2)$

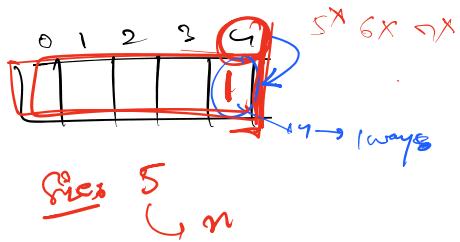
Maze paths



$ev=2$
 $ec=2$

6 paths

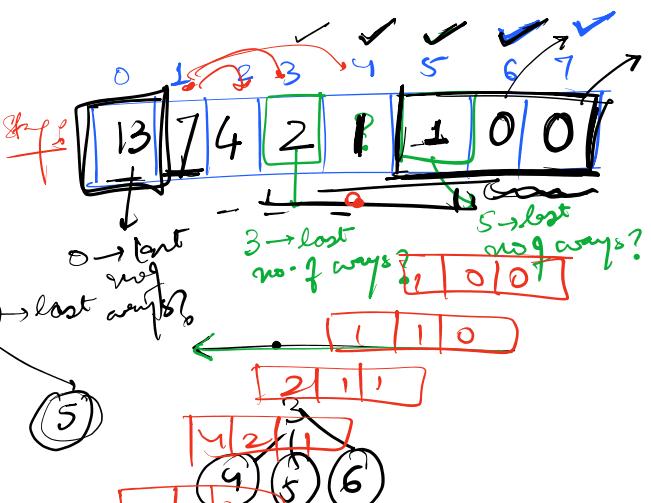
(2ways)



TQ: n unique problem \Rightarrow work

$$n^k = w$$

sci. see state of n



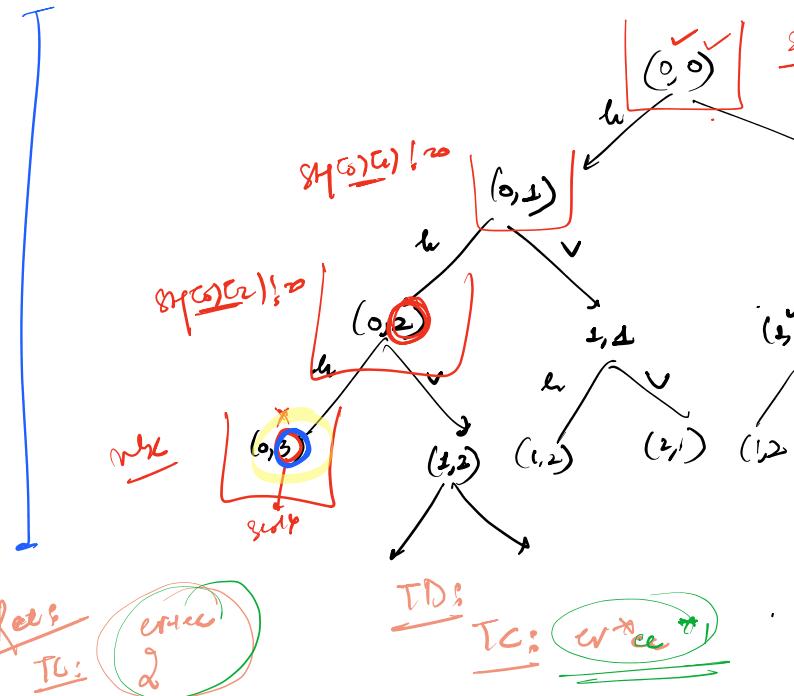
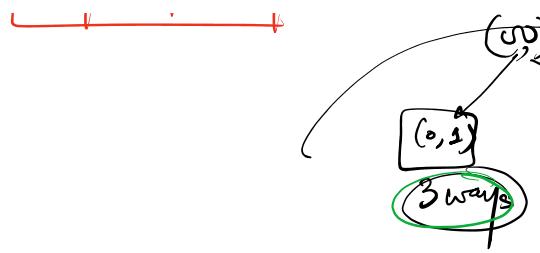
(n)

ways?

cont? no 9 ways

ways

6 ways

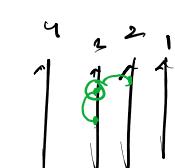
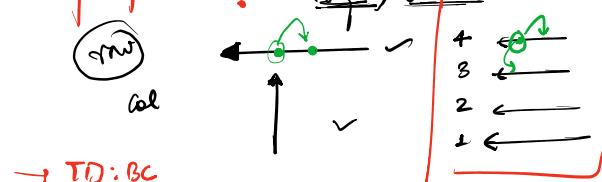


~~Bottom 5 p's~~

→ cell meaning? : count or no. of ways to reach

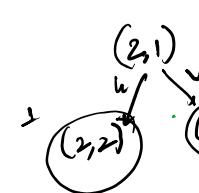
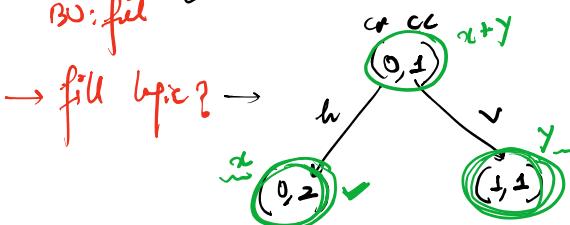
→ array size? $\rightarrow (n+2)(m+2)$

→ filling direction? right, bottom



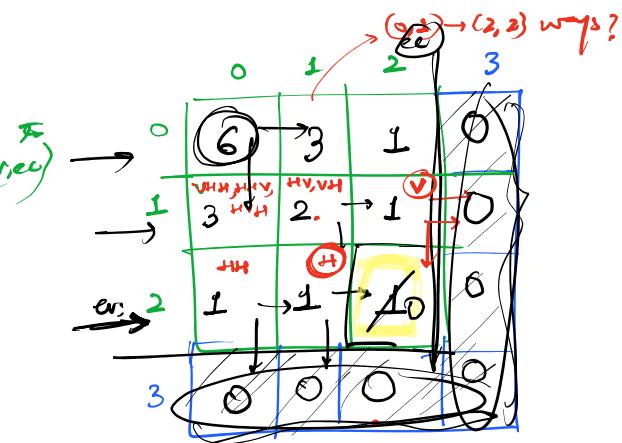
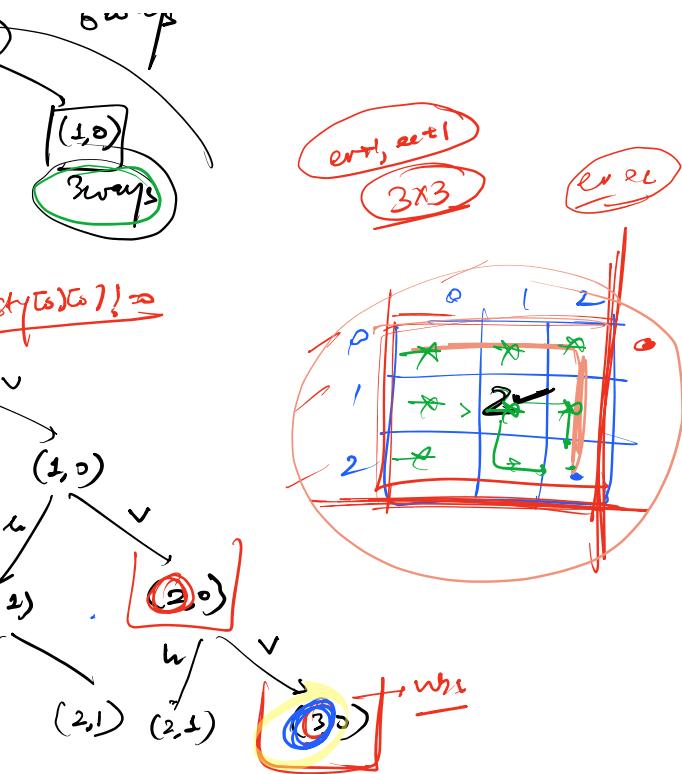
→ TD: BC
BC: fil

→ fill logic?



~~Wine Problem:~~

	w_1	w_2	w_3
$y_{r=1}$	1	4	2
	✓	6	1
..			



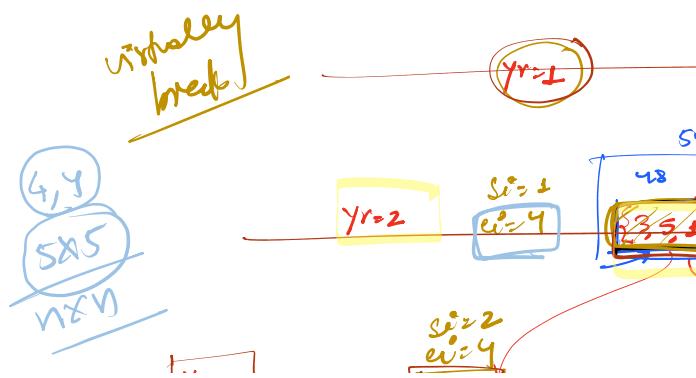
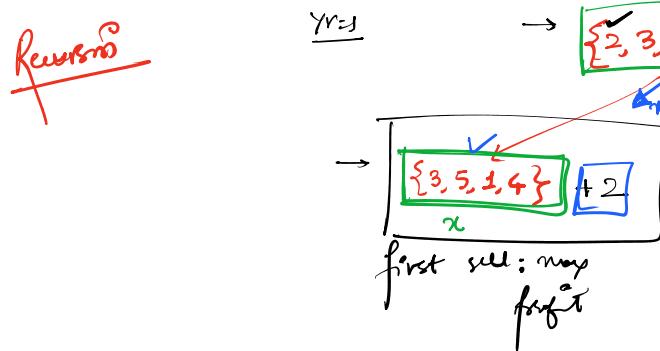
$(3,1)$

$$\begin{array}{c}
 \text{max profit} \\
 \text{Constraints} \\
 1 + 6 + 6 + 16 = 29
 \end{array}$$

Greedy Apprach
 [Best Soln]
 [2nd Soln]
 [3rd Soln]

	$\text{yr}=1$	$\text{yr}=2$	$\text{yr}=3$	$\text{yr}=4$
	(1)	2	3	4
	X	<u>8</u>		
	X	12	(6)	
	X	16	(8)	

yr	w_1	w_2	w_3	w_4
$\text{yr}=1$	(2)	3	5	
$\text{yr}=2$	4	6	10	
$\text{yr}=3$	6	9	15	
$\text{yr}=4$	8	12	20	
$\text{yr}=5$	10	15	25	



6
X
X

$$\frac{1+6+6+16}{13} = 29$$

Ant Soln
Bee -

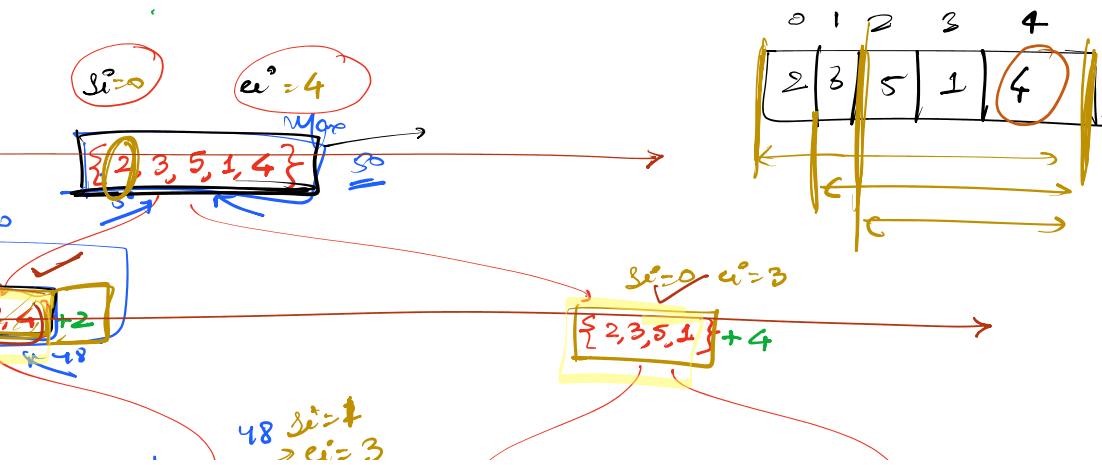
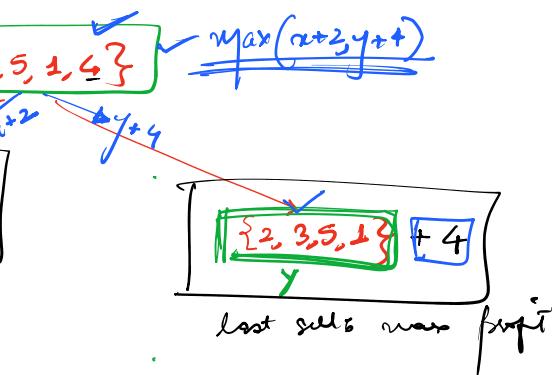
Globally Best Soln

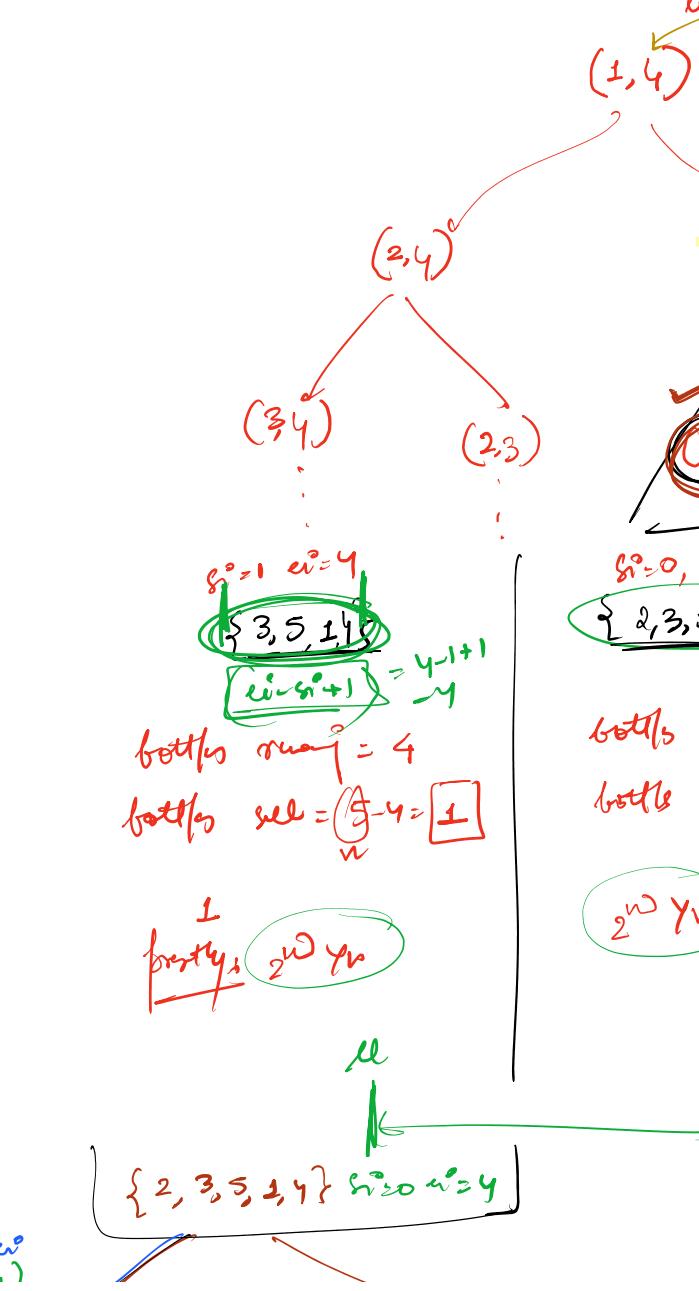
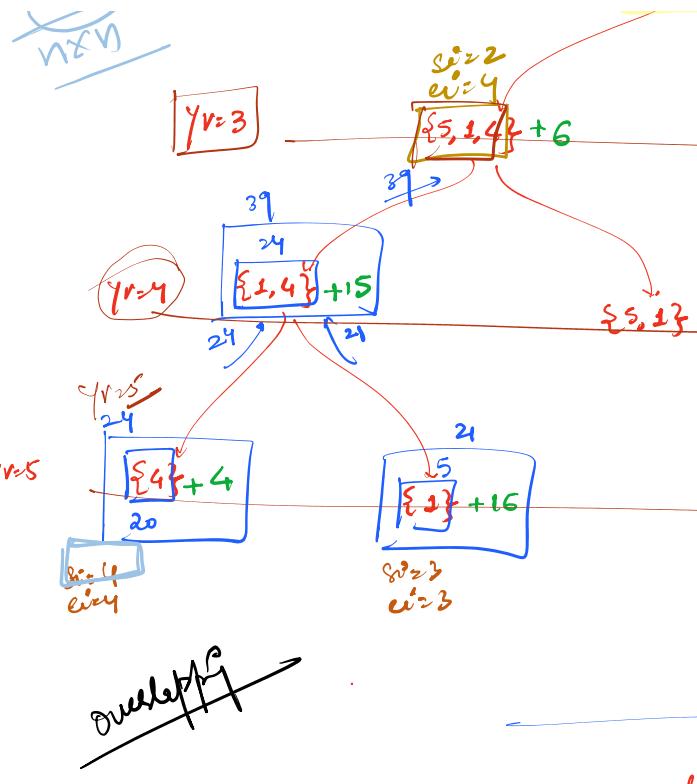
04 w5
1 4
2 8
3 12
4 16
5 20

$$2 + \boxed{6 + 12 + 4 + 25} \\ = 14 + 10 = 29 \\ = 49$$

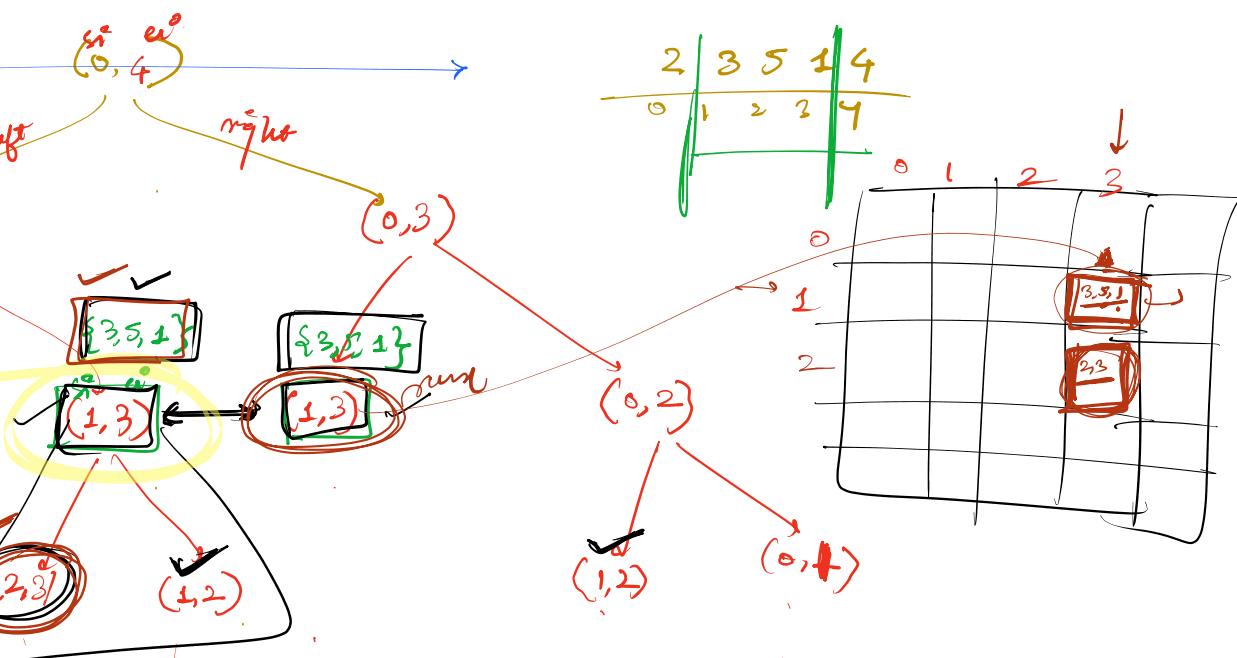
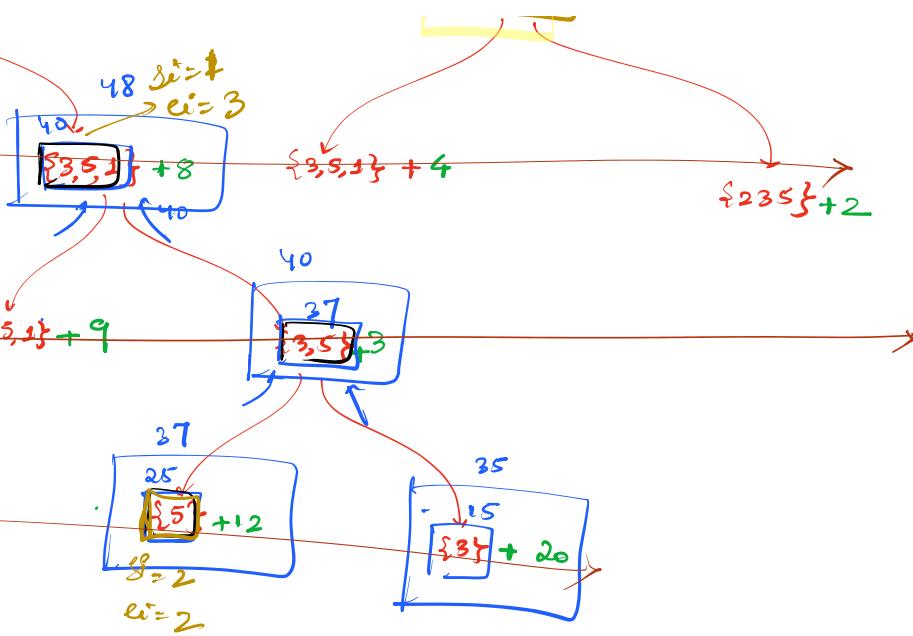
$$2 + 8 + 3 + 12 + 25 = 50$$

Greedy X





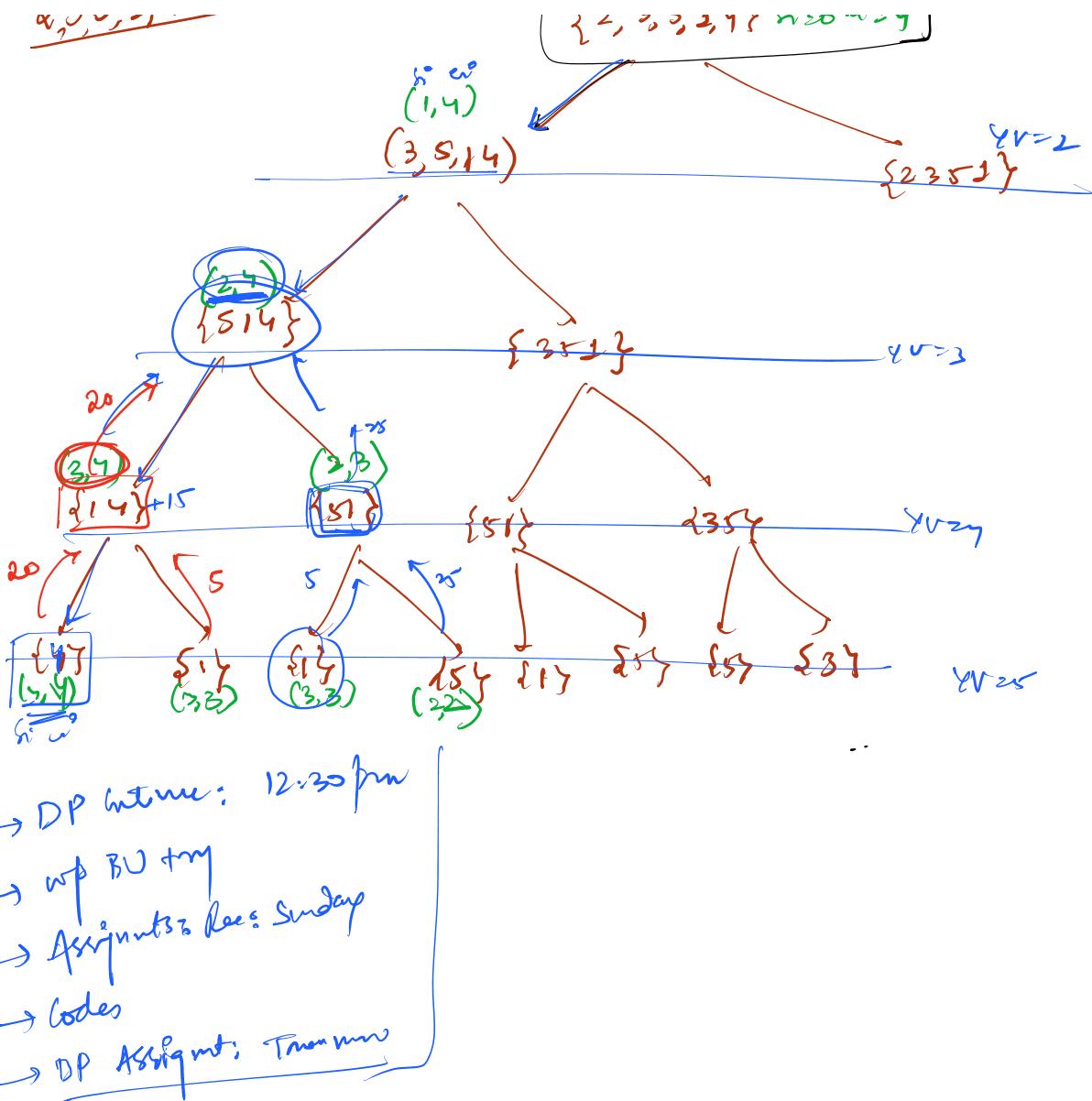
~~2, 3, 5, 1, 4~~



$$\begin{array}{c}
 \text{rowj} = 4 \Rightarrow 3 - 0 + 1 = 4 \\
 \text{sell} = 5 - 1 = 4 \rightarrow \text{lyr} \\
 \text{rowj} = 2 \Rightarrow 3 - 2 + 1 = 2 \\
 \text{sell} = 5 - 2 = 3 \rightarrow \text{lyr}
 \end{array}$$

u

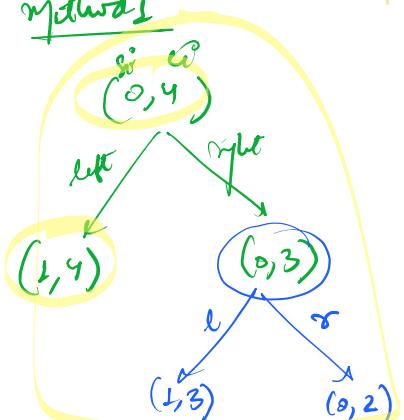
$$\begin{aligned}
 \text{bottles rowj} &= ei^o - si^o + 1 \\
 \text{sell} &= m - (ei^o - si^o + 1) \\
 p. \text{ lyr} &= m - (ei^o - si^o + 1) + 1
 \end{aligned}$$



Bottom up

- Cell meaning
- array size: $n \times n$
- filling down: diagonally

method 1



0	1	2	3	4
2	3	1	4	

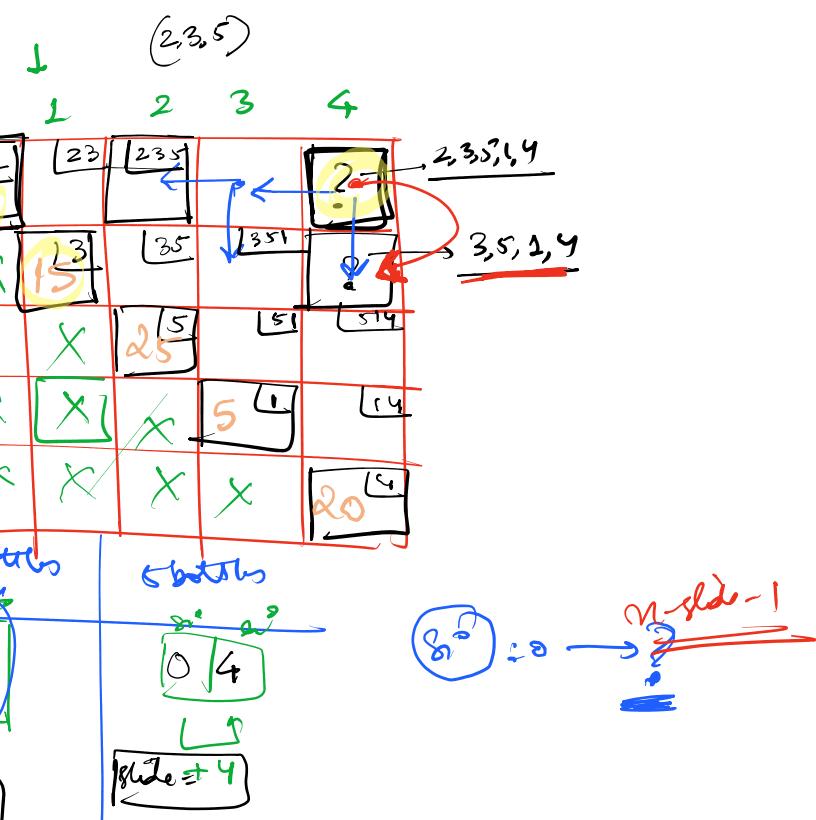
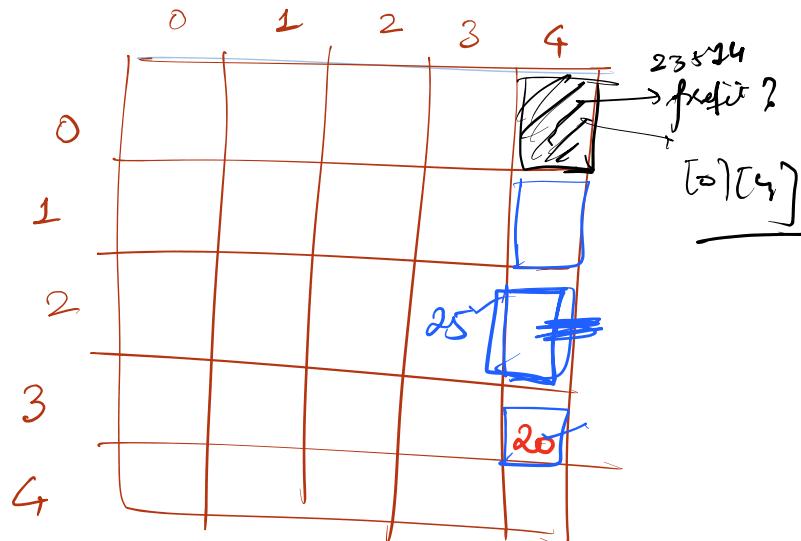
Si	ei	0
1	X	1
2	X	2
3	X	3
4	X	4

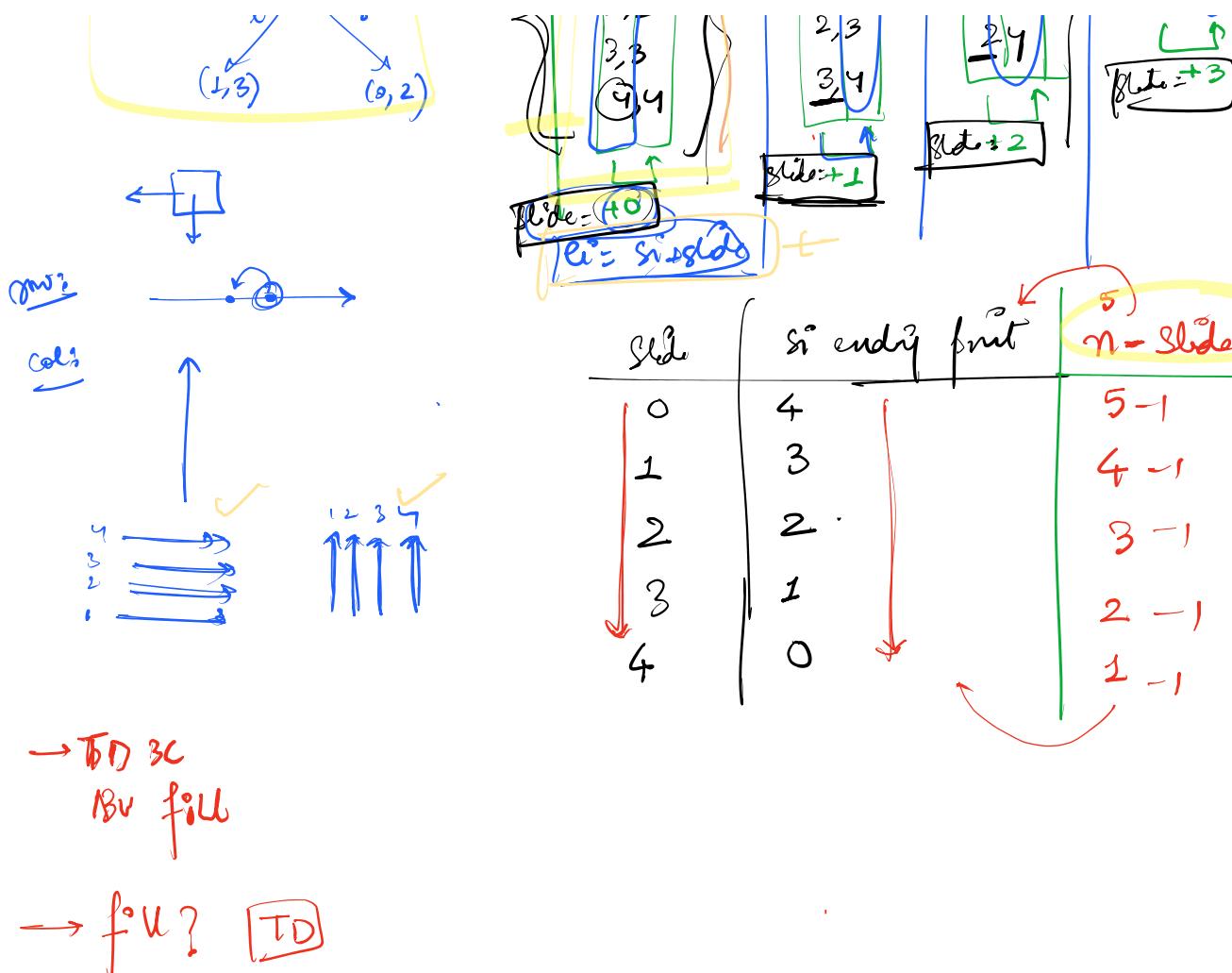
$Si > ei$

method 2		method 2	
1 bottle	2 bottles	3 bottles	4 bottles
0,0	0,1 1,0	0,2 1,1 2,0	0,3 1,2 2,1 3,0
1,1	1,2 2,1	1,3 2,2 3,1	1,4 2,3 3,2 4,1
2,2	2,3 3,2	2,4 3,3 4,2	2,5 3,4 4,3 5,2
3,3	3,4 4,3	3,5 4,4 5,3	3,6 4,5 5,4 6,3
4,4	4,5 5,4	4,6 5,5 6,4	4,7 5,6 6,5 7,4

Total = +3

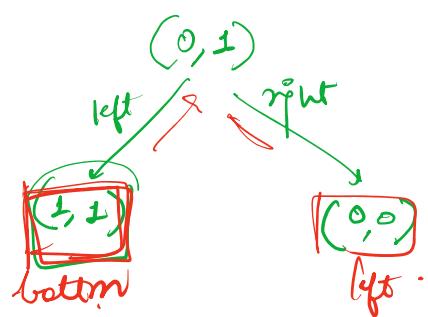
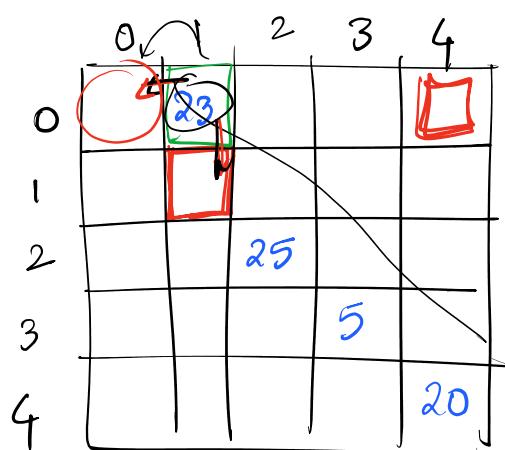
$$P = \frac{1}{N^2} \sum_{i=1}^N \sum_{j=1}^N \left(\text{exp}(-\gamma |x_i - x_j|) + \text{exp}(-\gamma |y_i - y_j|) \right)$$





$\{2, 3, 5, 1, 4\}$

for (i)



slide = 4

Si endy front = $n - \text{slide} - 1$

```
int slide = 0; slide <= n - 1; slide++) {  
    for (int si = 0; si <= n - slide - 1; si++) {  
        int ei = si + slide;  
        int yr = n - (ei - si + 1) + 1; // 4  
  
        if (si == ei) {  
            strg[si][ei] = arr[si] * yr; // 1  
        } else {  
            int left = strg[si + 1][ei] + arr[si] * yr; // bottom + 2 * 4 = 15 + 8 = 23  
            int right = strg[si][ei - 1] + arr[ei] * yr; // 10 + 12 = 22  
            int res = Math.max(left, right);  
            strg[si][ei] = res;  
        }  
    }  
}
```

Wildcard Pattern Matching

Given a text and a wildcard pattern, implement wildcard pattern matching algorithm that finds if wildcard pattern is matched with text. The matching should cover the entire text (not partial text).

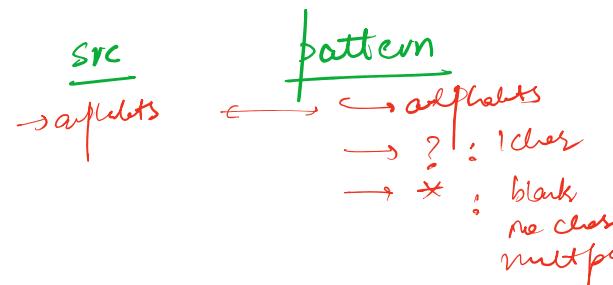
The wildcard pattern can include the characters '?' and '*'.

'?' – matches any single character

'*' – Matches any sequence of characters (including the empty sequence)

For example,

```
Text = "baaabab",
Pattern = "*****ba*****ab", output : true
Pattern = "baaa?ab", output : true
Pattern = "ba*a?", output : true
Pattern = "a*ab", output : false
```



```
Input =      ba aab ab
Pattern = *****ba*****ab
Output : true

Input =      baaabab
Pattern =     a * ab
Output : false

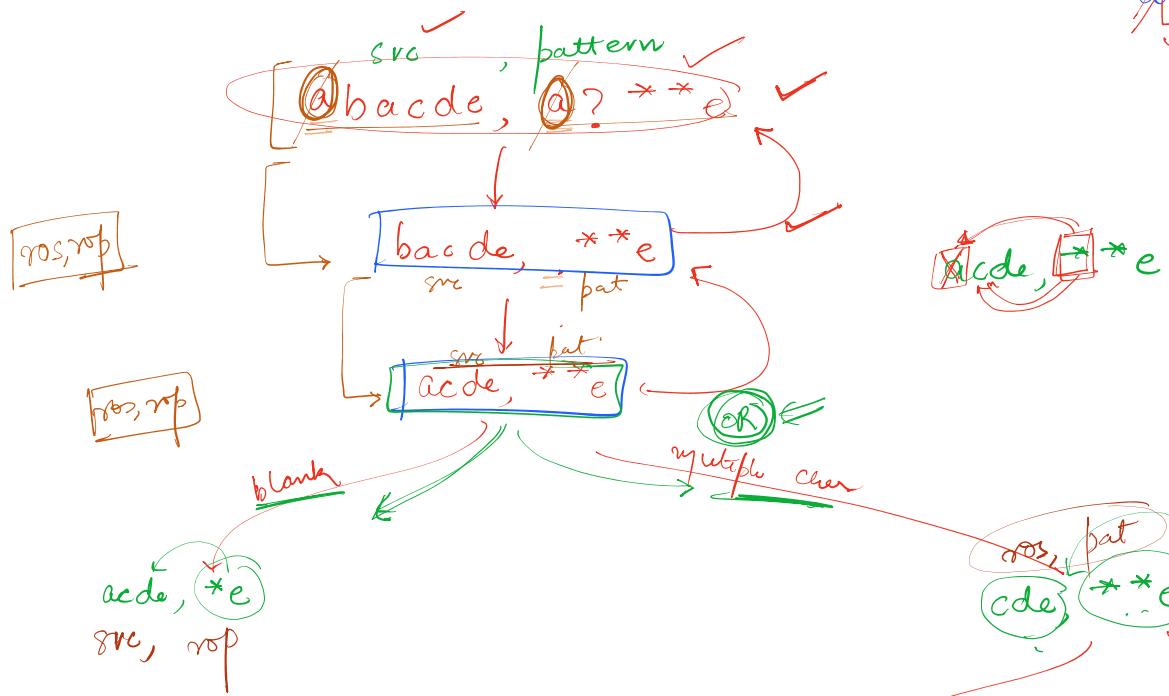
Input =      ba aab ab
Pattern =     ba * a?
Output : true
```

Legend:

- ? → 1 char
- * → blank
- * → 1 char
- * → m. c.

Each occurrence of '?' character in wildcard pattern can be replaced with any other character and each occurrence of '*' with a sequence of characters such that the wildcard pattern becomes identical to the input string after replacement.

Recursion



cher

1

s
b c d e
r o s

*: blank
~~super class~~
m. cher

3

m.c.

src, np

char match, ? : src, np : aligned
* : (src, np) : blank right bottom
src, pat) : np.c.
: false

cde, *e

src pat
-, - : true

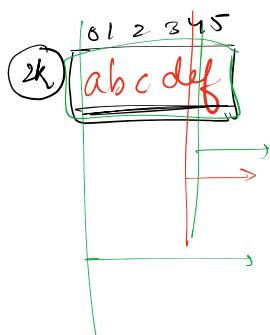
abc, d*

abc, - : false

{
-, ab*? : false
-, blank : true}



Initially



idx = 2 : cdef

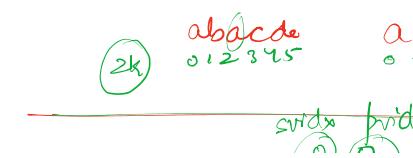
idx = 4 : ef

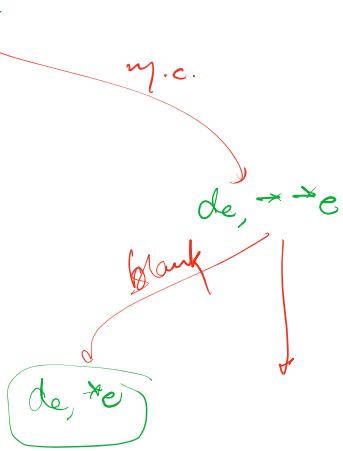
idx = 0 : abcdef

idx = 5 : f

idx = 6 : -

abacde → a? * * e }



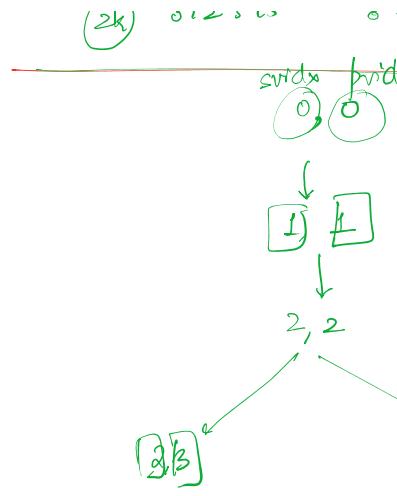
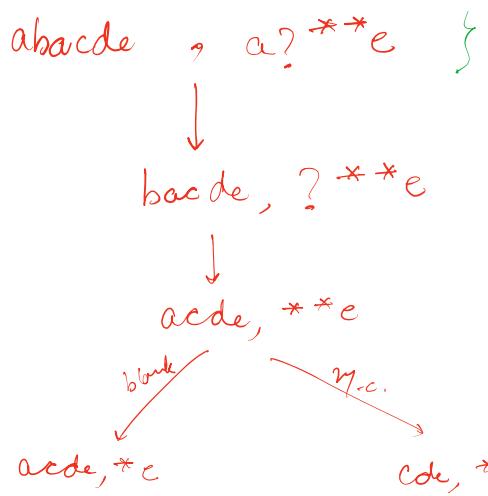


$\text{read} \rightarrow \text{end}$

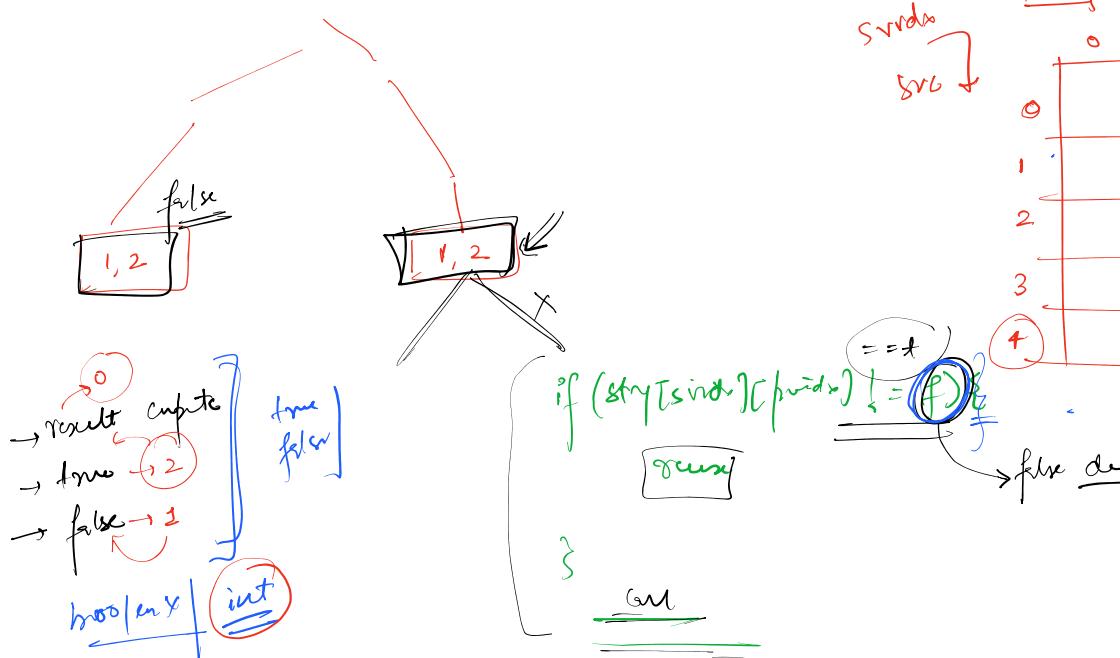
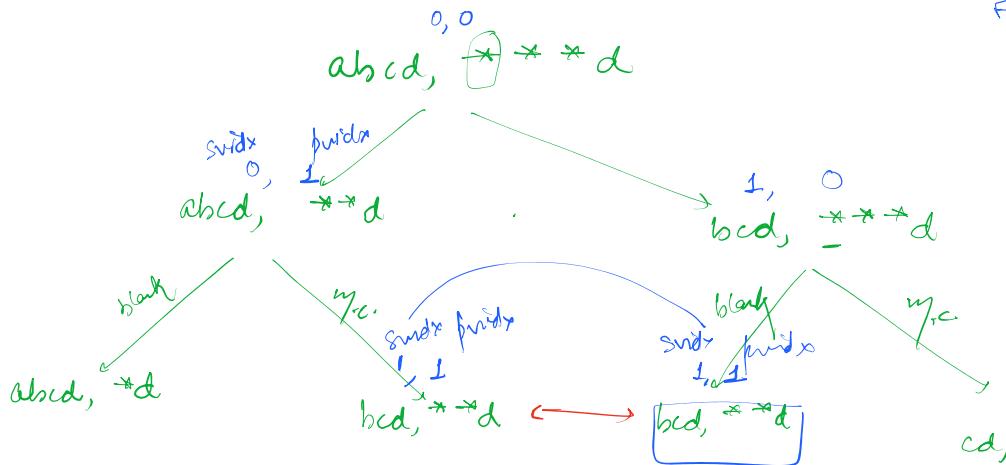
$$\frac{\text{rdx} = \text{S1 next}}{\text{break}}$$

? $\oplus^* e$ $\oplus k$

$x \rightarrow$



~~overleffig~~



~~mention up~~

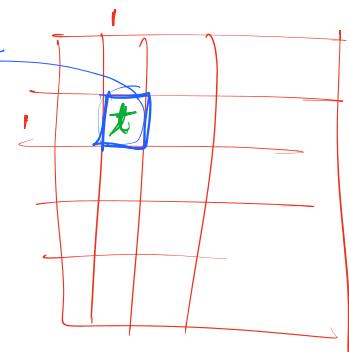
$$01 \overline{)23} \quad 01 \quad 2 \left\{ \begin{matrix} 3 \\ 45 \end{matrix} \right. \\ \cancel{2} \cancel{1} \cancel{2} \quad 2 \cancel{*} \quad 2 \left\{ \begin{matrix} \cancel{3} \\ \cancel{4} \cancel{5} \end{matrix} \right.$$

pat . (

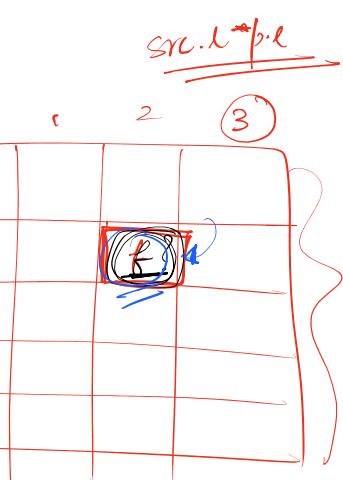
2 2 2 -1

x →

→
3,2.



→ → → d



wts [boolean] → H/H

Boolean → null
class

→ → → 3 → 4 5 → G → 1 2 0 - 0

Bottom up

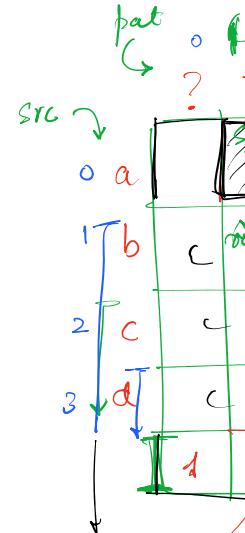
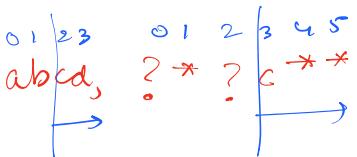
→ cell meaning

→ array size

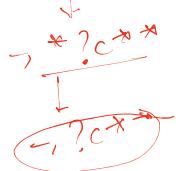
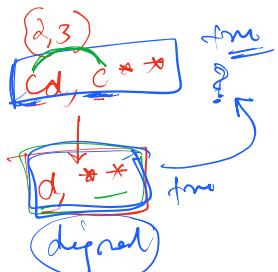
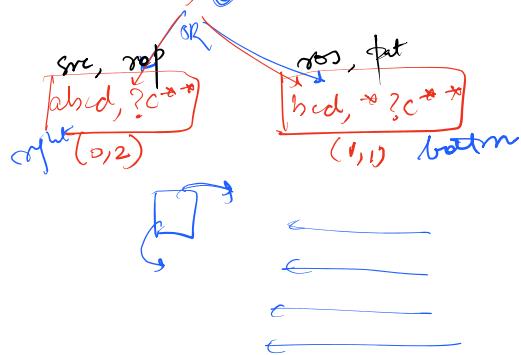
$$4 \times 6$$

$$5 \times 7$$

$$\rightarrow (s \cdot l + 1) \times (p \cdot l + 1)$$

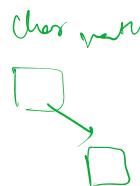
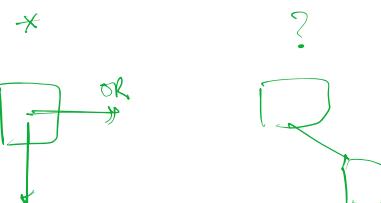


→ filling direction



→ TD BC
B to fill

→ filling logic



	0	1	2	3	4	5	6
	?	*	?	c	*	*	-
0 a	x	x	f	f	x	x	f
1 b	f	t	t	f	t	t	f
2 c	f	f	f	t	t	t	f
3 d	f	f	f	f	t	t	f
4 -	f	f	f	f	*	t	t

The first few Catalan numbers for $n = 0, 1, 2, 3, \dots$ are 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862



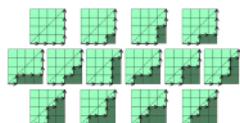
false

Applications :

1. Number of possible Binary Search Trees with n keys.
2. Number of expressions containing n pairs of parentheses which are correctly matched. For n = 3, possible expressions are ((())), ()(), ()(), (())(), (())().
3. Number of ways a convex polygon of n+2 sides can split into triangles by connecting vertices.



4. Number of full binary trees (A rooted binary tree is full if every vertex has either two children or no children) with n+1 leaves.
5. Number of different Unlabelled Binary Trees can be there with n nodes.
6. The number of paths with $2n$ steps on a rectangular grid from bottom left, i.e., $(n-1, 0)$ to top right $(0, n-1)$ that do not cross above the main diagonal.



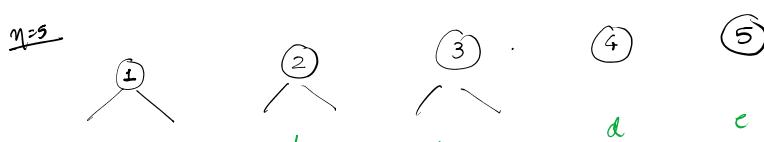
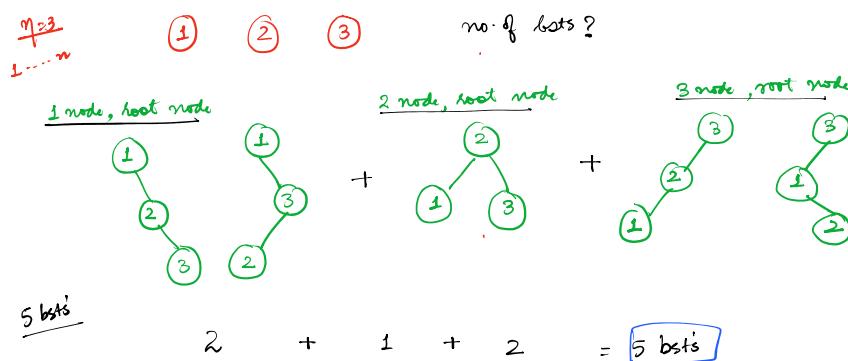
7. Number of ways to insert n pairs of parentheses in a word of n+1 letters, e.g., for n=2 there are 2 ways: $((ab)c)$ or $(a(bc))$. For n=3 there are 5 ways, $((ab)(cd))$, $((ab)c)d)$, $((a(bc))d)$, $(a((bc)d))$, $(a(b(cd)))$.
8. Number of Dyck words of length $2n$. A Dyck word is a string consisting of n X's and n Y's such that no initial segment of the string has more Y's than X's. For example, the following are the Dyck words of length 6: XXXYYY XYXXYY XYXYXY XXYYXY XXYXYY.
9. Number of ways to tile a staircase shape of height n with n rectangles. The following figure illustrates the case n = 4:

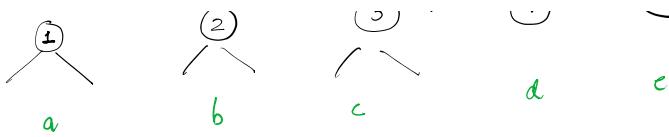


10. Number of ways to form a "mountain ranges" with n upstrokes and n down-strokes that all stay above the original line. The mountain range interpretation is that the mountains will never go below the horizon.

$n = 0$	*	1 way
$n = 1$	/	1 way
$n = 2$	/ \ / \ \ /	2 ways
$n = 3$	/ \ / \ / \ / \ \ / \ \ \ /	5 ways

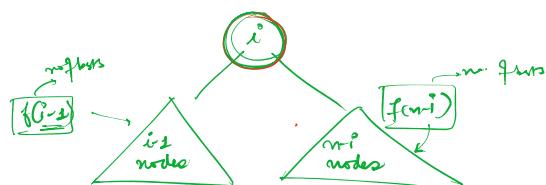
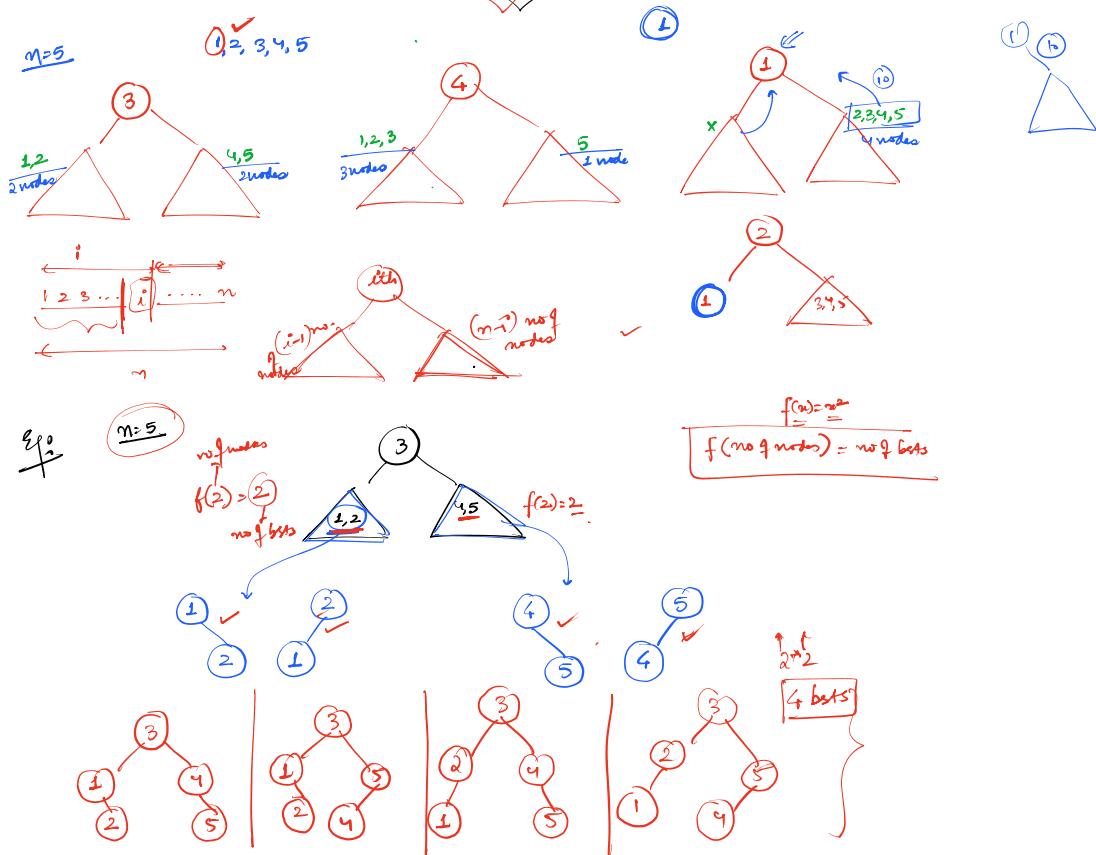
Mountain Ranges





total lists: no of lists obtained by having 1 node as root node → a
 +
2 nodes → b
 +
3rd node → c
 +
4th node → d
 +
5th node → e
 = a+b+c+d+e

$$\text{total bests} = \sum_{i=2}^n \text{no of bests obtained by having } i^{\text{th}} \text{ node as root node}$$



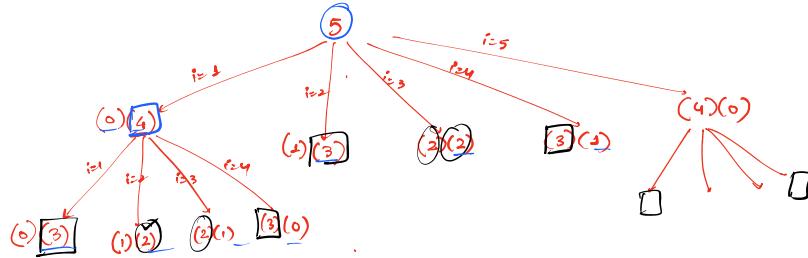
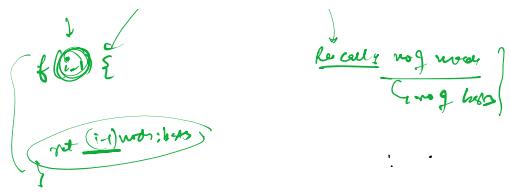
one of lists obtained by having a node as root node: $f(m_1) \rightarrow f(m_2)$

$$\text{no. of books with } \\ \text{at most } i = \sum_{j=1}^n f(j) * f(n-i)$$

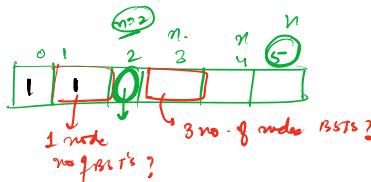
$$\frac{BC_3}{4} = \frac{6!}{3!3!1!4!}$$

$$= \frac{11 \times 10 \times 9 \times 8}{12 \times 11 \times 10 \times 9 \times 8 \times 7}$$

Recall + neg move



Bottom up
 → cell means?
 → array size? $10 \rightarrow n+1$
 → fully dev
 → TD BC, BU fill
 → fill logic?
 0 1 2 3 4 5 6 7 8 9
 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900, 2674440
Catalan Series!



→ DP Assignment
 → Predecessor Day: Catalyst given: $\frac{15 \text{ pm}}{12 \text{ min}}$ $\frac{12 \text{ min}}{12 \text{ min}}$ ✓
 → Catalyst 2D's: $\frac{\text{Catalyst} + \text{DP}}{12 \text{ min}}$ $\frac{12 \text{ min}}{12 \text{ min}}$ $\frac{12 \text{ min}}{12 \text{ min}}$ $\frac{12 \text{ min}}{12 \text{ min}}$
 → Monday: $\frac{\text{Catalyst} + \text{BU}}{12 \text{ min}}$ $\frac{12 \text{ min}}{12 \text{ min}}$
 → Tuesday: $\frac{\text{Catalyst} + \text{BU}}{12 \text{ min}}$ $\frac{12 \text{ min}}{12 \text{ min}}$
 → Wednesday: $\frac{\text{Catalyst} + \text{BU}}{12 \text{ min}}$ $\frac{12 \text{ min}}{12 \text{ min}}$
 → Thursday: $\frac{\text{Catalyst} + \text{BU}}{12 \text{ min}}$ $\frac{12 \text{ min}}{12 \text{ min}}$

