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DIV: B/B1

ADBMS

Exp2

ADBMS Bliol22 Fx p 2. Aim: Penform operation like sawching, moutton, dela on B-Tree and B+ Tree. Theory: B-Tree B-Tree is a self-balancing would tree. B-Trees are useful when we are dealing with huge amounts of data that and he fitted in main memory. When the number of keys is high, the data is read from the disk in forms of blocks. Disk across time is very high compared to the main memory across time. The main idea of using B-Trees as to reduce the number of disk across where h is height of tree B-Tree is a fat tree, meaning, the height of B-Tree B-Tree is a fat tree, meaning, the height of B-Tree B-Tree is a fat tree, meaning, the height of B-Tree B-Tree is a fat tree, meaning, the height of B-Tree B-Tree is a fat tree is meaning.	
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require O(h) disk access where h is height of tree	
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THE IS IN THE LEW THE PARTY OF	c
is kept low by putting maximum possible keys in	
a 8-Tree node. Generally, the 8-Trees: node size is	
kept agent to disk block size. Since the height of	
B-Free is kept low so total disk arress for most a	+
the operations are reduced significantly compared t	0
balanted Brazy Search Trees like AVI Trees.	
House races at cled to fall values to count ast int	
Properties of B-Tree	
-> All leaves are at the same level.	
-> B-Tree is defined by term minimum degrees 't'. The	ne ne
value of 't' depends upon drak block size:	
-> Every node except the soot must contain atleast	
(t-1) keys. The root may contain minimum of 1 ke	10
CE-11 rays. The outer truly solitary trilling of 1 he	3
FOR EDUCATIONAL USE	-

- 11	
- 1	-> All nodes may contain atmost (2*t-1) keys.
	> No. of children of a node is equal to no. of keys
	in it plus 1
	-> All keys of a node are sorted in increasing order.
	> It grows & shoots from root
	Traceution hamens note at leaf Almle
	> Insertion happens only at Leaf Noole.
8	Time Camplexity of R-Type:
	Time Complexity of B-Tree: Scarch - O Clagn
1	Insort - O(logn)
	Delete - O (log n)
- 3	Secret Script of the secret se
	a consequent to the first the many date to anything
	· B+ Tree
	B+ Tree is an extension of B-Tree which allows
	effectent moutton, deleten and search operations. In 8 Tree,
Ī.	keys and records both can be stored in internal as well
	as leaf nodes. Whereas, in B+ tree, records can only
10	be stored on leaf nodes 8 internal nodes can only store
1	key values
	The leaf node of a B+ tree are linked together
	in the form of singly linked lists to make seanch
	queries more affectent. B+ trees one used to store the
	large amount of data which and be stored in main
363	memosy. Due to fact that, size of morn memosy is
	always Irrotted, the internal nodes of B+ tree one
	stored in main memory whereas, leafnodes are
1	stored in secondary memory.
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98	openties of Bit tree		1 4 1 2 1 2 1
->	All leasts are at same leve		Lameter in sensibility
\rightarrow	The soot has atleast two chil	0	Strollers week
\rightarrow	Each node except root can h	080	
1453	children and atleast m/2	ON	bildren
->	Each node can contain a r	ax	tonum of (m-1) keys
	and a minimum of [m/	2.7	-1 Keys
111-	Keys are used for indexing	30	a unforthiory to worth
-	> Data can be stored seque	ite	ally or directly
	Time Complexity		
	Search - O (lagn)		
Inomaci pr	Front - O (logn)	500	with towards a
lanaci pi	Jesent - O (logn) Delete - O (logn)	600	and to control of
	Texant - O (logn) Delete - O (logn)	- Y0	Simular Samuel Comment of the Commen
	Front - O (logn)	Ya.	and to an advertigation of the second
	B-Tree vs B+Tree	70	B+-Tree
•	B-Tree vs B+Tree B-Tree	1)	B+-Tree Redundant sound keys
•	B-Tree vs B+Tree B-Tree scanch keys can not be		
•	B-Tree vs B+Tree B-Tree vs B+Tree B-Tree	-	Redundant sowch keys
•	B-Tree vs B+Tree B-Tree vs B+Tree B-anch keys an not be repeatedly stored.	-	Redundant sounch keys an be present
7>	B-Tree vs B+Tree B-Tree vs B+Tree B-Tree B-	17>	Redundant sounch keys an be present Data can only be stored on the lost nodes.
•	B-Tree vs B+Tree B-Tree vs B+Tree B-Tree B-anch keys an not be repeatedly stored. Data an be stored in leaf node as well as internal nodes Southing for some data is	17>	Redundant sounch keys can be present Data can only be stored on the lost nodes. Searching is compariatively
7>	B-Tree vs B+Tree B-Tree Scanch keys can not be repeatedly stored. Dota can be stored in loaf node as well as internal nodes Scanching for some data is a slower process since data	17>	Redundant sourch keys an be present Data can only be stored on the lost nodes. Searching is compariatively faster as data an only be
7>	B-Tree vs B+Tree B-Tree vs B+Tree B-Tree B-anch keys an not be repeatedly stored. Data an be stored in leaf node as well as internal nodes Southing for some data is	17>	Redundant sounch keys can be present Data can only be stored on the lost nodes. Searching is compariatively

B-Tree	1 B+-Tree
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Conclusion: Thus, we successivations	In B-Tree and B+ Tree

B Tree

Code:

```
class BTreeNode:
  def __init__(self, leaf=False):
    self.leaf = leaf
    self.keys = []
    self.child = []
```

class BTree:

```
def __init__(self, t):
 self.root = BTreeNode(True)
 self.t = t
def insert(self, k):
 root = self.root
 if len(root.keys) == (2 * self.t) - 1:
  temp = BTreeNode()
  self.root = temp
  temp.child.insert(0, root)
  self.split_child(temp, 0)
  self.insert_non_full(temp, k)
 else:
  self.insert_non_full(root, k)
def insert_non_full(self, x, k):
 i = len(x.keys) - 1
 if x.leaf:
  x.keys.append((None, None))
  while i \ge 0 and k[0] < x.keys[i][0]:
    x.keys[i + 1] = x.keys[i]
   i = 1
  x.\text{keys}[i+1] = k
 else:
  while i \ge 0 and k[0] < x.keys[i][0]:
   i = 1
  i += 1
```

```
if len(x.child[i].keys) == (2 * self.t) - 1:
    self.split_child(x, i)
    if k[0] > x.keys[i][0]:
     i += 1
  self.insert_non_full(x.child[i], k)
def split_child(self, x, i):
 t = self.t
 y = x.child[i]
 z = BTreeNode(y.leaf)
 x.child.insert(i + 1, z)
 x.keys.insert(i, y.keys[t - 1])
 z.keys = y.keys[t: (2 * t) - 1]
 y.keys = y.keys[0: t - 1]
 if not y.leaf:
  z.child = y.child[t: 2 * t]
  y.child = y.child[0: t - 1]
def print_tree(self, x, l=0):
 print("Level ", l, " ", len(x.keys), end=":")
 for i in x.keys:
  print(i, end=" ")
 print()
 1 += 1
 if len(x.child) > 0:
  for i in x.child:
    self.print_tree(i, 1)
```

```
def search_key(self, k, x=None):
  if x is not None:
   i = 0
   while i < len(x.keys) and k > x.keys[i][0]:
     i += 1
   if i < len(x.keys) and k == x.keys[i][0]:
     return (x, i)
   elif x.leaf:
     return None
   else:
     return self.search_key(k, x.child[i])
  else:
   return self.search_key(k, self.root)
def main():
 B = BTree(3)
 for i in range(10):
  B.insert((i, 2 * i))
 B.print_tree(B.root)
 if B.search_key(8) is not None:
  print("\nFound")
```

```
else:
    print("\nNot Found")

if __name__ == '__main__':
    main()
```

Output:

```
Level 1 2:(3, 6) (4, 8)
Level 1 4:(6, 12) (7, 14) (8, 16) (9, 18)
Found
```

B+ Tree

Code:

import math

```
class Node:
    def __init__(self, order):
        self.order = order
        self.values = []
        self.keys = []
        self.nextKey = None
        self.parent = None
        self.check_leaf = False
```

def insert_at_leaf(self, leaf, value, key):

```
if (self.values):
       temp1 = self.values
       for i in range(len(temp1)):
          if (value == temp1[i]):
            self.keys[i].append(key)
             break
          elif (value < temp1[i]):
            self.values = self.values[:i] + [value] + self.values[i:]
            self.keys = self.keys[:i] + [[key]] + self.keys[i:]
             break
          elif(i + 1 == len(temp1)):
            self.values.append(value)
            self.keys.append([key])
             break
     else:
       self.values = [value]
       self.keys = [[key]]
class BplusTree:
  def __init__(self, order):
     self.root = Node(order)
     self.root.check_leaf = True
  def insert(self, value, key):
     value = str(value)
     old_node = self.search(value)
```

```
old_node.insert_at_leaf(old_node, value, key)
  if (len(old_node.values) == old_node.order):
    node1 = Node(old_node.order)
    node1.check leaf = True
    node1.parent = old_node.parent
    mid = int(math.ceil(old_node.order / 2)) - 1
    node1.values = old_node.values[mid + 1:]
    node1.keys = old\_node.keys[mid + 1:]
    node1.nextKey = old_node.nextKey
    old_node.values = old_node.values[:mid + 1]
    old_node.keys = old_node.keys[:mid + 1]
    old_node.nextKey = node1
    self.insert_in_parent(old_node, node1.values[0], node1)
def search(self, value):
  current node = self.root
  while(current_node.check_leaf == False):
    temp2 = current_node.values
    for i in range(len(temp2)):
       if (value == temp2[i]):
         current_node = current_node.keys[i + 1]
         break
       elif (value < temp2[i]):
         current_node = current_node.keys[i]
         break
       elif (i + 1 == len(current_node.values)):
```

```
current_node = current_node.keys[i + 1]
          break
  return current_node
def find(self, value, key):
  1 = self.search(value)
  for i, item in enumerate(l.values):
    if item == value:
       if key in l.keys[i]:
         return True
       else:
         return False
  return False
def insert_in_parent(self, n, value, ndash):
  if (self.root == n):
    rootNode = Node(n.order)
    rootNode.values = [value]
    rootNode.keys = [n, ndash]
     self.root = rootNode
    n.parent = rootNode
    ndash.parent = rootNode
    return
  parentNode = n.parent
  temp3 = parentNode.keys
  for i in range(len(temp3)):
```

```
if (temp3[i] == n):
       parentNode.values = parentNode.values[:i] + \
         [value] + parentNode.values[i:]
       parentNode.keys[:i +
                           1] + [ndash] + parentNode.keys[i + 1:]
       if (len(parentNode.keys) > parentNode.order):
         parentdash = Node(parentNode.order)
         parentdash.parent = parentNode.parent
         mid = int(math.ceil(parentNode.order / 2)) - 1
         parentdash.values = parentNode.values[mid + 1:]
         parentdash.keys = parentNode.keys[mid + 1:]
         value_ = parentNode.values[mid]
         if (mid == 0):
           parentNode.values = parentNode.values[:mid + 1]
         else:
           parentNode.values = parentNode.values[:mid]
         parentNode.keys = parentNode.keys[:mid + 1]
         for j in parentNode.keys:
           j.parent = parentNode
         for j in parentdash.keys:
           j.parent = parentdash
         self.insert_in_parent(parentNode, value_, parentdash)
def delete(self, value, key):
  node_ = self.search(value)
  temp = 0
```

```
for i, item in enumerate(node_.values):
     if item == value:
       temp = 1
       if key in node_.keys[i]:
          if len(node\_.keys[i]) > 1:
            node_.keys[i].pop(node_.keys[i].index(key))
          elif node_ == self.root:
            node_.values.pop(i)
            node_.keys.pop(i)
          else:
            node_.keys[i].pop(node_.keys[i].index(key))
            del node_.keys[i]
            node_.values.pop(node_.values.index(value))
            self.deleteEntry(node_, value, key)
       else:
          print("Value not in Key")
          return
  if temp == 0:
     print("Value not in Tree")
     return
def deleteEntry(self, node_, value, key):
  if not node_.check_leaf:
     for i, item in enumerate(node_.keys):
       if item == key:
```

```
node_.keys.pop(i)
            break
       for i, item in enumerate(node_.values):
         if item == value:
            node_.values.pop(i)
            break
     if self.root == node_ and len(node_.keys) == 1:
       self.root = node\_.keys[0]
       node_.keys[0].parent = None
       del node_
       return
     elif (len(node_.keys) < int(math.ceil(node_.order / 2)) and
node_.check_leaf == False) or (len(node_.values) < int(math.ceil((node_.order -</pre>
1)/2)) and node_.check_leaf == True):
       is_predecessor = 0
       parentNode = node_.parent
       PrevNode = -1
       NextNode = -1
       PrevK = -1
       PostK = -1
       for i, item in enumerate(parentNode.keys):
         if item == node_:
            if i > 0:
              PrevNode = parentNode.keys[i - 1]
              PrevK = parentNode.values[i - 1]
```

```
if i < len(parentNode.keys) - 1:
       NextNode = parentNode.keys[i + 1]
       PostK = parentNode.values[i]
if PrevNode == -1:
  ndash = NextNode
  value_ = PostK
elif NextNode == -1:
  is\_predecessor = 1
  ndash = PrevNode
  value = PrevK
else:
  if len(node_.values) + len(NextNode.values) < node_.order:</pre>
     ndash = NextNode
     value_ = PostK
  else:
    is\_predecessor = 1
     ndash = PrevNode
     value_ = PrevK
if len(node_.values) + len(ndash.values) < node_.order:</pre>
  if is_predecessor == 0:
    node_, ndash = ndash, node_
  ndash.keys += node_.keys
  if not node_.check_leaf:
    ndash.values.append(value_)
```

```
else:
    ndash.nextKey = node\_.nextKey
  ndash.values += node_.values
  if not ndash.check_leaf:
    for j in ndash.keys:
      j.parent = ndash
  self.deleteEntry(node_.parent, value_, node_)
  del node_
else:
  if is_predecessor == 1:
    if not node_.check_leaf:
       ndashpm = ndash.keys.pop(-1)
       ndashkm_1 = ndash.values.pop(-1)
       node_.keys = [ndashpm] + node_.keys
       node_.values = [value_] + node_.values
       parentNode = node_.parent
       for i, item in enumerate(parentNode.values):
         if item == value_:
           p.values[i] = ndashkm_1
           break
    else:
       ndashpm = ndash.keys.pop(-1)
       ndashkm = ndash.values.pop(-1)
       node_.keys = [ndashpm] + node_.keys
       node_.values = [ndashkm] + node_.values
```

```
parentNode = node_.parent
    for i, item in enumerate(p.values):
       if item == value_:
         parentNode.values[i] = ndashkm
         break
else:
  if not node_.check_leaf:
    ndashp0 = ndash.keys.pop(0)
    ndashk0 = ndash.values.pop(0)
    node_.keys = node_.keys + [ndashp0]
    node_.values = node_.values + [value_]
    parentNode = node_.parent
    for i, item in enumerate(parentNode.values):
       if item == value_:
         parentNode.values[i] = ndashk0
         break
  else:
    ndashp0 = ndash.keys.pop(0)
    ndashk0 = ndash.values.pop(0)
    node_.keys = node_.keys + [ndashp0]
    node_.values = node_.values + [ndashk0]
    parentNode = node_.parent
    for i, item in enumerate(parentNode.values):
       if item == value:
         parentNode.values[i] = ndash.values[0]
         break
```

```
if not ndash.check_leaf:
            for j in ndash.keys:
               j.parent = ndash
          if not node_.check_leaf:
            for j in node_.keys:
               j.parent = node_
          if not parentNode.check_leaf:
            for j in parentNode.keys:
               j.parent = parentNode
def printTree(tree):
  lst = [tree.root]
  level = [0]
  leaf = None
  flag = 0
  lev_leaf = 0
  node1 = Node(str(level[0]) + str(tree.root.values))
  while (len(lst) != 0):
     x = lst.pop(0)
     lev = level.pop(0)
     if (x.check_leaf == False):
       for i, item in enumerate(x.keys):
          print(item.values)
     else:
```

```
for i, item in enumerate(x.keys):
          print(item.values)
       if (flag == 0):
          lev_leaf = lev
          leaf = x
          flag = 1
record_len = 3
bplustree = BplusTree(record_len)
bplustree.insert('5', '33')
bplustree.insert('15', '21')
bplustree.insert('25', '31')
bplustree.insert('35', '41')
bplustree.insert('45', '10')
printTree(bplustree)
if(bplustree.find('45', '10')):
  print("Found")
else:
  print("Not found")
Output:
          '25']
          '45']
ound
```