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Oct. 31, 2019 | 7.6K views

Given a sorted array nums, remove the duplicates in-place such that duplicates appeared at most twice and return the new length.

Do not allocate extra space for another array, you must do this by modifying the input array in-place with O(1) extra memory.

Example 1:

Given nums = [1,1,1,2,2,3], Your function should return length = 5, with the first five elements of nums being 1, 1, 2, 2 and 3 respectively.

It doesn't matter what you leave beyond the returned length.

Example 2: Given nums = [0,0,1,1,1,1,2,3,3], Your function should return length = 7, with the first seven elements of nums being modified to 0, 0, 1, 1, 2, 3 and 3 respectively.

It doesn't matter what values are set beyond the returned length.

Clarification: Confused why the returned value is an integer but your answer is an array? Note that the input array is passed in by reference, which means modification to the input array will be

known to the caller as well.

// any modification to nums in your function would be known by the caller.

// using the length returned by your function, it prints the first len elements. for (int i = 0; i < len; i++) { print(nums[i]); }

```
Intuition
The input array is already sorted and hence, all the duplicates appear next to each other. The problem
statement mentions that we are not allowed to use any additional space and we have to modify the array in-
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place. The easiest approach for in-place modifications would be to get rid of all the unwanted duplicates. For

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Unwanted Duplicates

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every number in the array, if we detect > 2 duplicates, we simply remove them from the list of elements

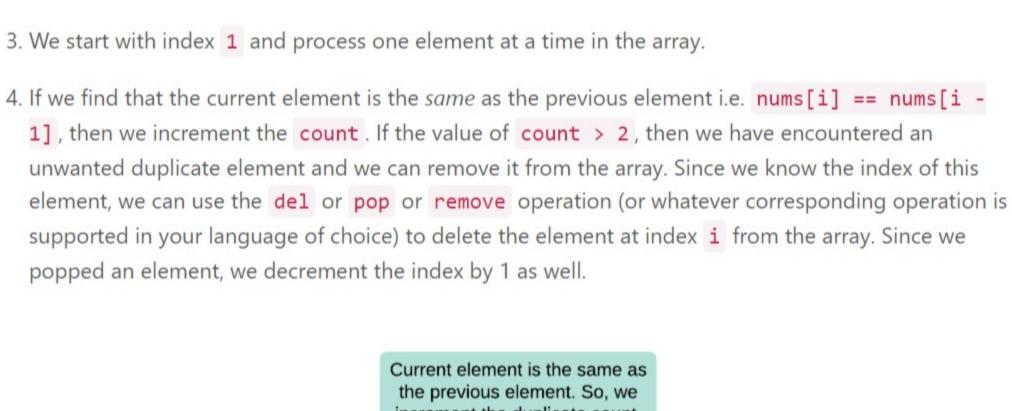
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and we do this for all the elements in the array.

1. The implementation is slightly tricky so to say since we will be removing elements from the array and iterating over it at the same time. So, we need to keep updating the array's indexes as and when we pop an element else we'll be accessing invalid indexes. 2. Say we have two variables, i which is the array pointer and count which keeps track of the count of a particular element in the array. Note that the minimum count would always be 1.



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count = **3** 1 1 1 2 3 3 3 Current element is the not the same as the previous element. We reset the count to 1 and continue.

If the count is more than 2, this is an # unwanted duplicate element and hence we # remove it from the array. if count > 2: nums.pop(i) # Note that we have to decrement the

We have to iterate over all the elements in the array. Suppose that the original array contains N

array index value to keep it consistent

• Time Complexity: Let's see what the costly operations in our array are:

elements, the time taken here would be O(N).

```
Approach 2: Overwriting unwanted duplicates
Intuition
The second approach is really inspired by the fact that the problem statement asks us to return the new
length of the array from the function. If all we had to do was remove elements, the function would not really
ask us to return the updated length. However, in our scenario, this is really an indication that we don't need
to actually remove elements from the array. Instead, we can do something better and simply overwrite the
duplicate elements that are unwanted.
     We won't be able to achieve this using a single pointer. We will be using a two-pointer approach
      where one pointer iterates over the original set of elements and another one that keeps track of the
      next "empty" location in the array or the next location that can be overwritten in the array.
Algorithm
   1. We define two pointers, i and j for our algorithm. The pointer i iterates of the array processing
     one element at a time and j keeps track of the next location in the array where we can overwrite an
     element.
   2. We also keep a variable count which keeps track of the count of a particular element in the array.
     Note that the minimum count would always be 1.
   3. We start with index 1 and process one element at a time in the array.
   4. If we find that the current element is the same as the previous element i.e. nums[i] == nums[i -
```

1], then we increment the count. If the value of count > 2, then we have encountered an

5. However, if the count is $\langle = 2 \rangle$, then we can move the element from index \mathbf{i} to index \mathbf{j} .

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unwanted duplicate element. In this case, we simply move forward i.e. we increment i but not j.

For the first two cases, both "i" and "j" progress at the same rate

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At this point, we don't

unwanted duplicate.

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We return the value of "j"

from the function since

that is the **new length of**

the array

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Next 👀

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6. If we encounter that the current element is *not* the same as the previous element i.e. nums[i] != nums[i - 1], then it means we have a new element at hand and so accordingly, we update count = 1 and also move this element to index j.

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def removeDuplicates(self, nums):

:type nums: List[int]

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At this point, the count is

3 and we don't copy

anything over since this

is an unwanted duplicate

element.

the array.

Python

class Solution(object):

Java

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overwriting the element at index "j" in the array if count <= 2: nums[j] = nums[i]**Complexity Analysis** • Time Complexity: O(N) since we process each element exactly once. • Space Complexity: O(1). Rate this article: * * * * * Comments: 9 Type comment here... (Markdown is supported) Preview soumyajitchatterjee73 🛊 15 🗿 November 1, 2019 9:10 AM using two-pointer: O(n) solution `private static int getRemoveDuplicates(int [] a) {

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- Internally you can think of this: // nums is passed in by reference. (i.e., without making a copy) int len = removeDuplicates(nums);
- Solution Approach 1: Popping Unwanted Duplicates
- Algorithm count = 11 1 1 2 3 3 3

1.

Python

Java

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Complexity Analysis

increment the duplicate count. count = 2 2 3 1 1 1 3 3

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Сору

count = 3

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- Pop the unwanted duplicate element thus reducing the size of the array by 1. 5. If we encounter that the current element is not the same as the previous element i.e. nums[i] != nums[i - 1], then it means we have a new element at hand and so accordingly, we update count =
- 6. Since we are removing all the unwanted duplicates from the original array, the final array that remains after process all the elements will only contain the valid elements and hence we simply return the length of this array. class Solution(object): def removeDuplicates(self, nums): :type nums: List[int] :rtype: int # Initialize the counter and the array index. i, count = 1, 1 # Start from the second element of the array and process # elements one by one. while i < len(nums): # If the current element is a duplicate, # increment the count. if nums[i] == nums[i - 1]: count += 1
- o Next, for every unwanted duplicate element, we will have to perform a delete operation and deletions in arrays are also O(N). o The worst case would be when all the elements in the array are the same. In that case, we would be performing N-2 deletions thus giving us $O(N^2)$ complexity for deletions \circ Overall complexity = $O(N) + O(N^2) \equiv O(N^2)$. • Space Complexity: O(1) since we are modifying the array in-place.

1 1 2 1 progress "j" since this is an count = 3

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count = 1

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count = 2

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7. It goes without saying that whenever we copy a new element to nums[j], we have to update the

value of j as well since j always points to the location where the next element can be copied to in

When we process the element 2, the count is reset to 1 and we copy it to the array location pointed to by "j". This is the first of our in-place overwrites

- :rtype: int # Initialize the counter and the second pointer. j, count = 1, 1 # Start from the second element of the array and process # elements one by one. for i in range(1, len(nums)): # If the current element is a duplicate, # increment the count. if nums[i] == nums[i - 1]: count += 1 else: # Reset the count since we encountered a different element # than the previous one count = 1 # For a count <= 2, we copy the element over thus
- Clean code :(. class Solution: def removeDuplicates(self, nums: List[int]) -> int: length = 0Read More 1 A V C Share Reply

int m = 0, i=1, count=1, lastVisit = a[0];

a[m] = a[0]:

ngoc_lam ★ 42 ② March 23, 2020 4:59 PM

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def removeDuplicates(self, nums):

:rtvpe: int

1 A V C Share Reply

:type nums: List[int]

2 A V C Share Reply

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Why cannot use two pointer like this? class Solution: def removeDuplicates(self, nums: List[int]) -> int: if len(nums) <= 2: Read More

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- var removeDuplicates = function(nums) { if (nums.length <= 2) { return nums.length: Read More
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 - def removeDuplicates(self, nums: List[int]) -> int: # only sorted Read More
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