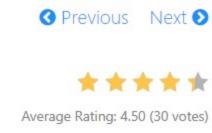
95. Unique Binary Search Trees II 🛂

Oct. 30, 2018 | 38.8K views



Сору

Right subtree : G (n - i) BST

Example:

Given an integer n, generate all structurally unique **BST's** (binary search trees) that store values 1 ... n.

Input: 3 Output:

```
[1,null,3,2],
   [3,2,null,1],
   [3,1,null,null,2],
   [2,1,3],
   [1,null,2,null,3]
 Explanation:
 The above output corresponds to the 5 unique BST's shown below:
    1
                  1 1 3
    2
Constraints:
```

Solution

Tree definition

0 <= n <= 8

Definition for a binary tree node.

Java

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Python

class TreeNode:

def __init__(self, x):

self.left = None

of possible BST is actually a Catalan number.

different right subtrees, where **G** is a Catalan number.

Elements available for left subtree : G (i - 1) BST

self.val = x

self.right = None

First of all, here is the definition of the TreeNode which we would use.

```
Approach 1: Recursion
First of all let's count how many trees do we have to construct. As you could check in this article, the number
```

Let's pick up number i out of the sequence 1 .. n and use it as the root of the current tree. Then there are i - 1 elements available for the construction of the left subtree and n - i elements available for the

class Solution:

def generateTrees(self, n):

:rtype: List[TreeNode]

if start > end:

all_trees = []

def generate_trees(start, end):

return [None,]

for i in range(start, end + 1): # pick up a root

left_trees = generate_trees(start, i - 1)

Type comment here... (Markdown is supported)

(12)

all possible left subtrees if i is choosen to be a root

:type n: int

1 2

> 3 4

5

6 7

8

9

10 11

12 13

14

15

Algorithm

000

right subtree. As we already discussed that results in G(i - 1) different left subtrees and G(n - i)

Let's follow the logic from the above article, this time not to count but to actually construct the trees.

Pick up as a root Now let's repeat the step above for the sequence 1 ... i - 1 to construct all left subtrees, and then for the sequence $i + 1 \dots n$ to construct all right subtrees. This way we have a root i and two lists for the possible left and right subtrees. The final step is to loop over both lists to link left and right subtrees to the root. **С**ору Python Java

```
# all possible right subtrees if i is choosen to be a root
  16
  17
                       right_trees = generate_trees(i + 1, end)
  18
  19
                       # connect left and right subtrees to the root i
                       for l in left_trees:
  20
  21
                            for r in right_trees:
  22
                                current_tree = TreeNode(i)
  23
                                current_tree.left = 1
  24
                                current_tree.right = r
  25
                                all_trees.append(current_tree)
  26
                   return all_trees
  27
Complexity analysis
   • Time complexity: The main computations are to construct all possible trees with a given root, that is
      actually Catalan number G_n as was discussed above. This is done rac{1}{n} times, that results in time
      complexity nG_n. Catalan numbers grow as \frac{4^n}{n^{3/2}} that gives the final complexity \mathcal{O}(\frac{4^n}{n^{1/2}}). Seems to be
      large but let's not forget that here we're asked to generate G_n \sim rac{4^n}{n^{3/2}} tree objects as output.
   • Space complexity : nG_n as we keep G_n trees with n elements each, that results in \mathcal{O}(\frac{4^n}{n^{1/2}})
      complexity.
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