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Example 1:

Given a linked list, rotate the list to the right by k places, where k is non-negative.

## Input: 1->2->3->4->5->NULL, k = 2

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```
Output: 4->5->1->2->3->NULL
 Explanation:
 rotate 1 steps to the right: 5->1->2->3->4->NULL
 rotate 2 steps to the right: 4->5->1->2->3->NULL
Example 2:
 Input: 0 - > 1 - > 2 - > NULL, k = 4
```

Explanation:

Output: 2->0->1->NULL

```
rotate 1 steps to the right: 2->0->1->NULL
 rotate 2 steps to the right: 1->2->0->NULL
 rotate 3 steps to the right: 0->1->2->NULL
 rotate 4 steps to the right: 2->0->1->NULL
Solution
```

## The nodes in the list are already linked, and hence the rotation basically means • To close the linked list into the ring.

Approach 1:

## • To break the ring after the new tail and just in front of the new head.

than n.

**Algorithm** 

Implementation

Python

class Solution:

# base cases if not head:

return None

return head

while old\_tail.next:

old\_tail.next = head

# close the linked list into the ring

old\_tail = old\_tail.next

if not head.next:

old\_tail = head

n += 1

n = 1

Java

1

2

3

4 5

6

7

8 9

10

11

12

13 14

15

16

17

18

19

20 21

22

23

24 25

26 27

The algorithm is quite straightforward:

Intuition

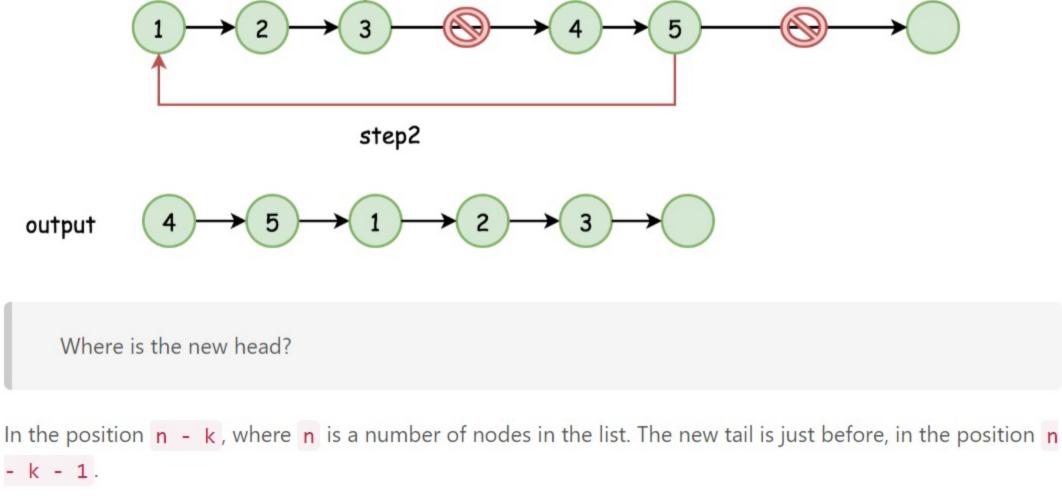
input

k = 2

step3

step1

step4



• Find the old tail and connect it with the head old\_tail.next = head to close the ring. Compute the length of the list **n** at the same time.

def rotateRight(self, head: 'ListNode', k: 'int') -> 'ListNode':

We were assuming that k < n. What about the case of k >= n?

• Find the new tail, which is (n - k % n - 1) th node from the head and the new head, which is (n - k % n) th node. • Break the ring new\_tail.next = None and return new\_head.

k = 2

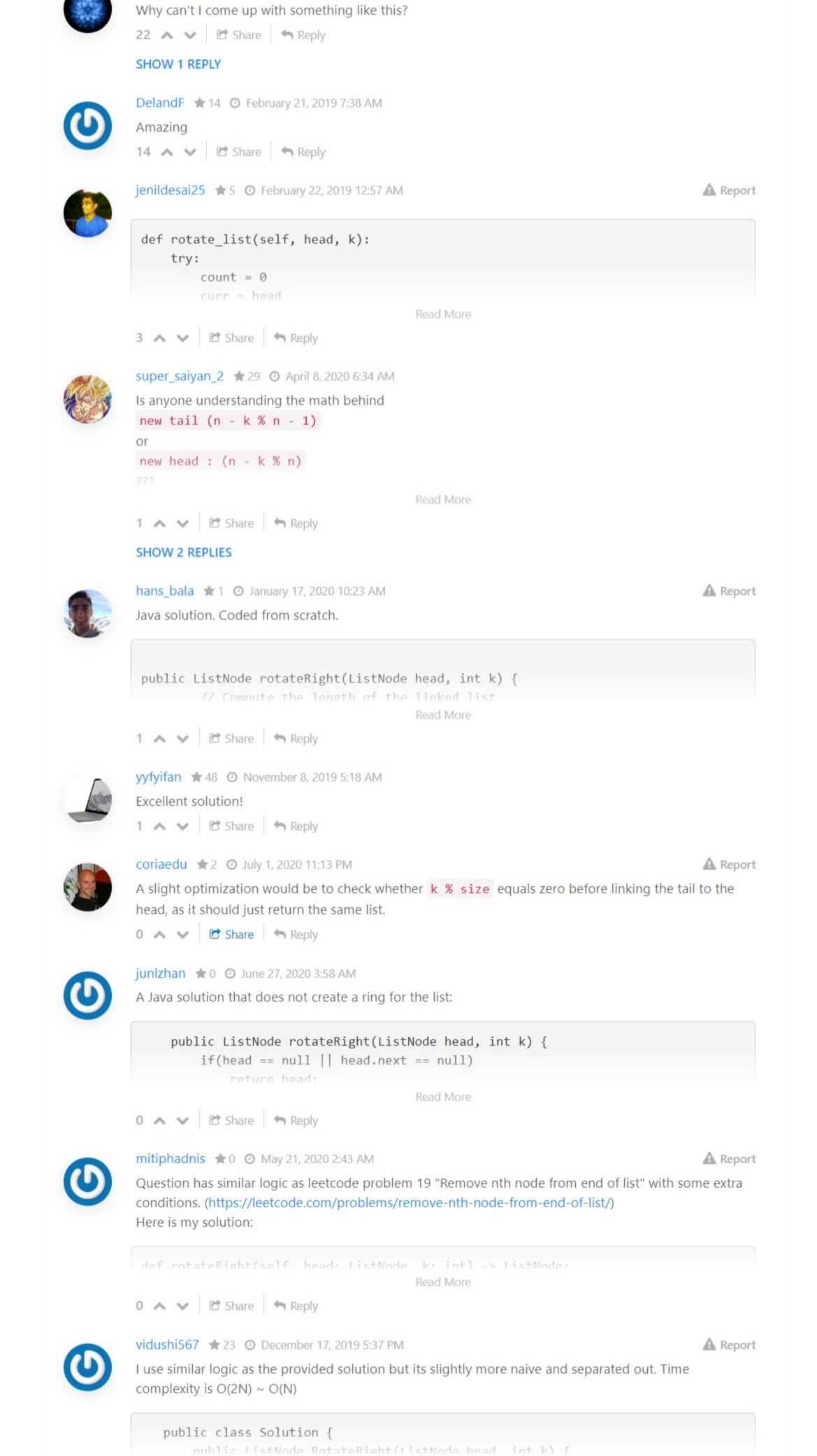
k could be rewritten as a sum k = (k // n) \* n + k % n, where the first term doesn't result in any

rotation. Hence one could simply replace k by k % n to always have number of rotation places smaller

1. Close the ring and compute number of nodes : n = 0

Copy

# find new tail : (n - k % n - 1)th node # and new head : (n - k % n)th node new\_tail = head for i in range(n - k % n - 1): new\_tail = new\_tail.next new\_head = new\_tail.next # break the ring new\_tail.next = None return new\_head **Complexity Analysis** ullet Time complexity :  $\mathcal{O}(N)$  where N is a number of elements in the list. • Space complexity :  $\mathcal{O}(1)$  since it's a constant space solution. Rate this article: \* \* \* \* \* Next O Previous Comments: 11 Sort By ▼ Type comment here... (Markdown is supported) Preview Post yanshengjia 🛊 61 🗿 October 8, 2019 5:34 AM



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