Previous Next

14. Longest Common Prefix

April 27, 2016 | 467.7K views

Average Rating: 4.66 (329 votes)

Write a function to find the longest common prefix string amongst an array of strings.

If there is no common prefix, return an empty string `""`.

Example 1:

```
Input: ["flower","flow","flight"]
Output: "fl"
```

Example 2:

```
Input: ["dog","racecar","car"]
Output: ""
Explanation: There is no common prefix among the input strings.
```

Note:

All given inputs are in lowercase letters `a-z`.

Solution

Approach 1: Horizontal scanning

Intuition

For a start we will describe a simple way of finding the longest prefix shared by a set of strings $LCP(S_1 \ldots S_n)$. We will use the observation that :

$LCP(S_1 \ldots S_n) = LCP(LCP(LCP(S_1, S_2), S_3), \ldots S_n)$

Algorithm

To employ this idea, the algorithm iterates through the strings $[S_1 \ldots S_n]$, finding at each iteration i the longest common prefix of strings $LCP(S_1 \ldots S_i)$ When $LCP(S_1 \ldots S_i)$ is an empty string, the algorithm ends. Otherwise after n iterations, the algorithm returns $LCP(S_1 \ldots S_n)$.

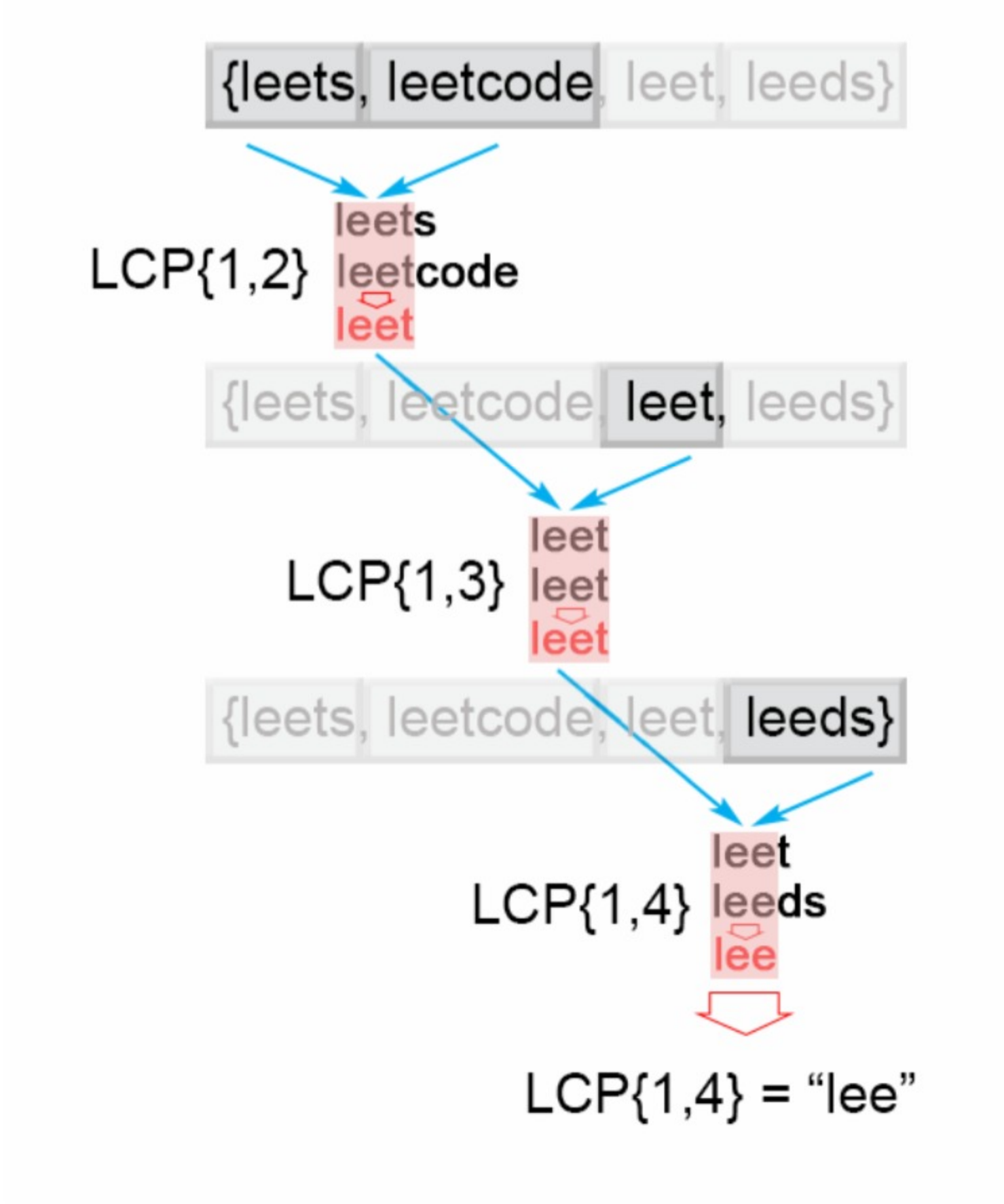

Figure 1. Finding the longest common prefix (Horizontal scanning)

Java

```java
public String longestCommonPrefix(String[] strs) {
    if (strs.length == 0) return "";
    String prefix = strs[0];
    for (int i = 1; i < strs.length; i++)
        while (strs[i].indexOf(prefix) != 0) {
            prefix = prefix.substring(0, prefix.length() - 1);
            if (prefix.isEmpty()) return "";
        }
    return prefix;
}
```

Complexity Analysis

- Time complexity : $O(S)$, where S is the sum of all characters in all strings.

  In the worst case all n strings are the same. The algorithm compares the string $S1$ with the other strings $[S_2 \ldots S_n]$ There are S character comparisons, where S is the sum of all characters in the input array.

- Space complexity : $O(1)$. We only used constant extra space.

Approach 2: Vertical scanning

Algorithm

Imagine a very short string is at the end of the array. The above approach will still do S comparisons. One way to optimize this case is to do vertical scanning. We compare characters from top to bottom on the same column (same character index of the strings) before moving on to the next column.

Java

```java
public String longestCommonPrefix(String[] strs) {
    if (strs == null || strs.length == 0) return "";
    for (int i = 0; i < strs[0].length() ; i++){
        char c = strs[0].charAt(i);
        for (int j = 1; j < strs.length; j ++) {
            if (i == strs[j].length() || strs[j].charAt(i) != c)
                return strs[0].substring(0, i);
        }
    }
    return strs[0];
}
```

Complexity Analysis

- Time complexity : $O(S)$, where S is the sum of all characters in all strings. In the worst case there will be n equal strings with length m and the algorithm performs $S = m \cdot n$ character comparisons. Even though the worst case is still the same as Approach 1, in the best case there are at most $n \cdot minLen$ comparisons where $minLen$ is the length of the shortest string in the array.
- Space complexity : $O(1)$. We only used constant extra space.

Approach 3: Divide and conquer

Intuition

The idea of the algorithm comes from the associative property of LCP operation. We notice that : $LCP(S_1 \ldots S_n) = LCP(LCP(S_1 \ldots S_k), LCP(S_{k+1} \ldots S_n))$, where $LCP(S_1 \ldots S_n)$ is the longest common prefix in set of strings $[S_1 \ldots S_n]$, $1 < k < n$

Algorithm

To apply the observation above, we use divide and conquer technique, where we split the $LCP(S_i \ldots S_j)$ problem into two subproblems $LCP(S_i \ldots S_{mid})$ and $LCP(S_{mid+1} \ldots S_j)$, where `mid` is $\frac{i+j}{2}$. We use their solutions `lcpLeft` and `lcpRight` to construct the solution of the main problem $LCP(S_i \ldots S_j)$. To accomplish this we compare one by one the characters of `lcpLeft` and `lcpRight` till there is no character match. The found common prefix of `lcpLeft` and `lcpRight` is the solution of the $LCP(S_i \ldots S_j)$.

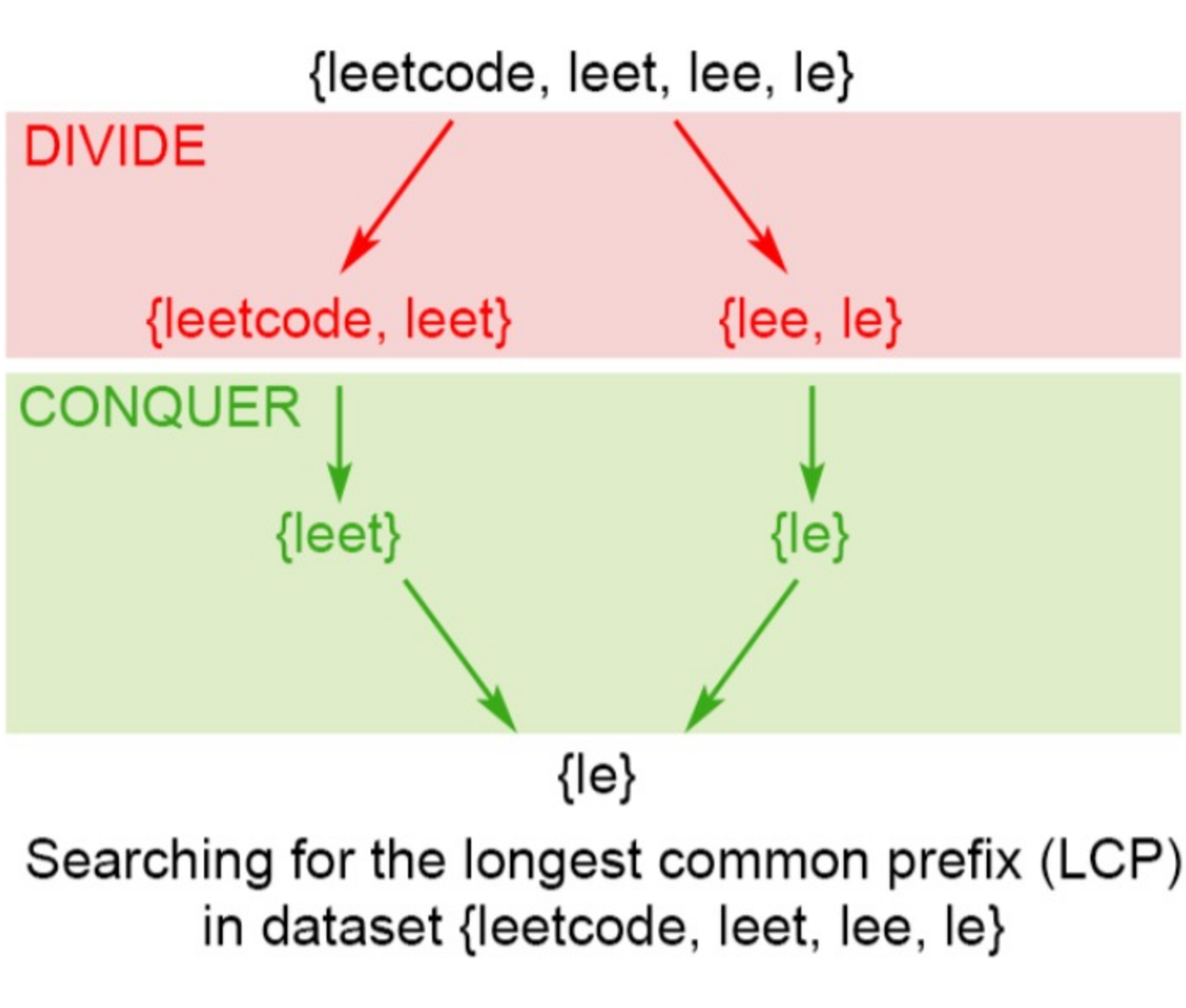

Figure 2. Finding the longest common prefix of strings using divide and conquer technique

Java

```java
public String longestCommonPrefix(String[] strs) {
    if (strs == null || strs.length == 0) return "";
        return longestCommonPrefix(strs, 0 , strs.length - 1);
}

private String longestCommonPrefix(String[] strs, int l, int r) {
    if (l == r) {
        return strs[l];
    }
    else {
        int mid = (l + r)/2;
        String lcpLeft =   longestCommonPrefix(strs, l , mid);
        String lcpRight =  longestCommonPrefix(strs, mid + 1,r);
        return commonPrefix(lcpLeft, lcpRight);
   }
}

String commonPrefix(String left,String right) {
    int min = Math.min(left.length(), right.length());
    for (int i = 0; i < min; i++) {
        if ( left.charAt(i) != right.charAt(i) )
            return left.substring(0, i);
    }
    return left.substring(0, min);
}
```

Complexity Analysis

In the worst case we have n equal strings with length m

- Time complexity : $O(S)$, where S is the number of all characters in the array, $S = m \cdot n$ Time complexity is $2 \cdot T\left(\frac{n}{2}\right) + O(m)$. Therefore time complexity is $O(S)$. In the best case this algorithm performs $O(minLen \cdot n)$ comparisons, where $minLen$ is the shortest string of the array

- Space complexity : $O(m \cdot \log n)$

  There is a memory overhead since we store recursive calls in the execution stack. There are $\log n$ recursive calls, each store need m space to store the result, so space complexity is $O(m \cdot \log n)$

Approach 4: Binary search

The idea is to apply binary search method to find the string with maximum value `L`, which is common prefix of all of the strings. The algorithm searches space is the interval $(0 \ldots minLen)$, where `minLen` is minimum string length and the maximum possible common prefix. Each time search space is divided in two equal parts, one of them is discarded, because it is sure that it doesn't contain the solution. There are two possible cases: `S[1...mid]` *is not a common string. This means that for each* `j > i S[1..j]` *is not a common string and we discard the second half of the search space.* `S[1...mid]` is common string. This means that for for each `i < j S[1..i]` is a common string and we discard the first half of the search space, because we try to find longer common prefix.

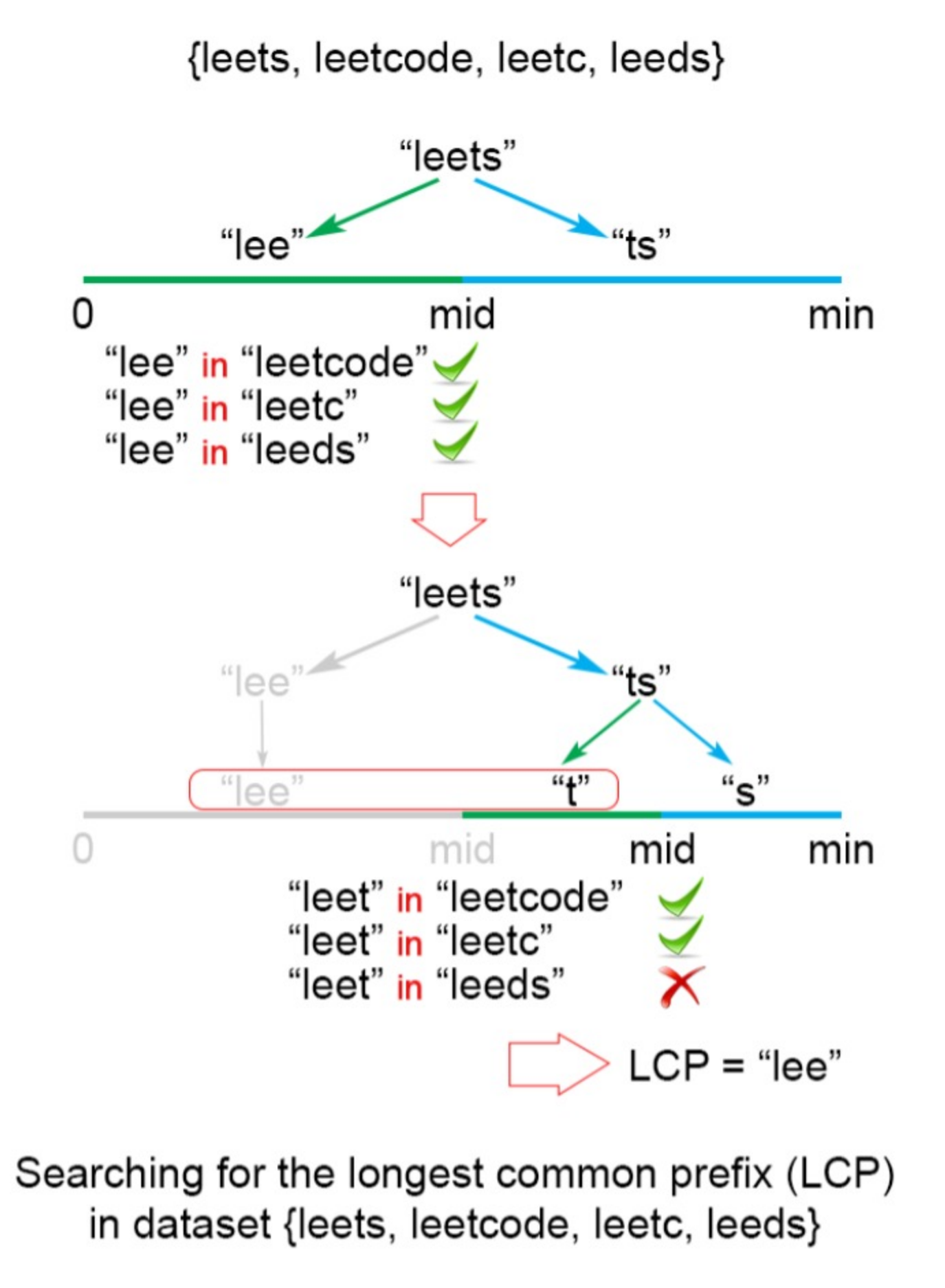

Figure 3. Finding the longest common prefix of strings using binary search technique

Java

```java
public String longestCommonPrefix(String[] strs) {
    if (strs == null || strs.length == 0)
        return "";
    int minLen = Integer.MAX_VALUE;
    for (String str : strs)
        minLen = Math.min(minLen, str.length());
    int low = 1;
    int high = minLen;
    while (low <= high) {
        int middle = (low + high) / 2;
        if (isCommonPrefix(strs, middle))
            low = middle + 1;
        else
            high = middle - 1;
    }
    return strs[0].substring(0, (low + high) / 2);
}

private boolean isCommonPrefix(String[] strs, int len){
    String str1 = strs[0].substring(0,len);
    for (int i = 1; i < strs.length; i++)
        if (!strs[i].startsWith(str1))
            return false;
    return true;
}
```

Complexity Analysis

In the worst case we have n equal strings with length m

- Time complexity : $O(S \cdot \log m)$, where S is the sum of all characters in all strings.

  The algorithm makes $\log m$ iterations, for each of them there are $S = m \cdot n$ comparisons, which gives in total $O(S \cdot \log m)$ time complexity.

- Space complexity : $O(1)$. We only used constant extra space.

Further Thoughts / Follow up

Let's take a look at a slightly different problem:

> Given a set of keys S = $[S_1, S_2 \ldots S_n]$, find the longest common prefix among a string `q` and S. This LCP query will be called frequently.

We could optimize LCP queries by storing the set of keys S in a Trie. For more information about Trie, please see this article Implement a trie (Prefix trie). In a Trie, each node descending from the root represents a common prefix of some keys. But we need to find the longest common prefix of a string `q` and all key strings. This means that we have to find the deepest path from the root, which satisfies the following conditions: *it is prefix of query string* `q` each node along the path must contain only one child element. Otherwise the found path will not be a common prefix among all strings. * the path doesn't comprise of nodes which are marked as end of key. Otherwise the path couldn't be a prefix a of key which is shorter than itself.

Algorithm

The only question left, is how to find the deepest path in the Trie, that fulfills the requirements above. The most effective way is to build a trie from $[S_1 \ldots S_n]$ strings. Then find the prefix of query string `q` in the Trie. We traverse the Trie from the root, till it is impossible to continue the path in the Trie because one of the conditions above is not satisfied.

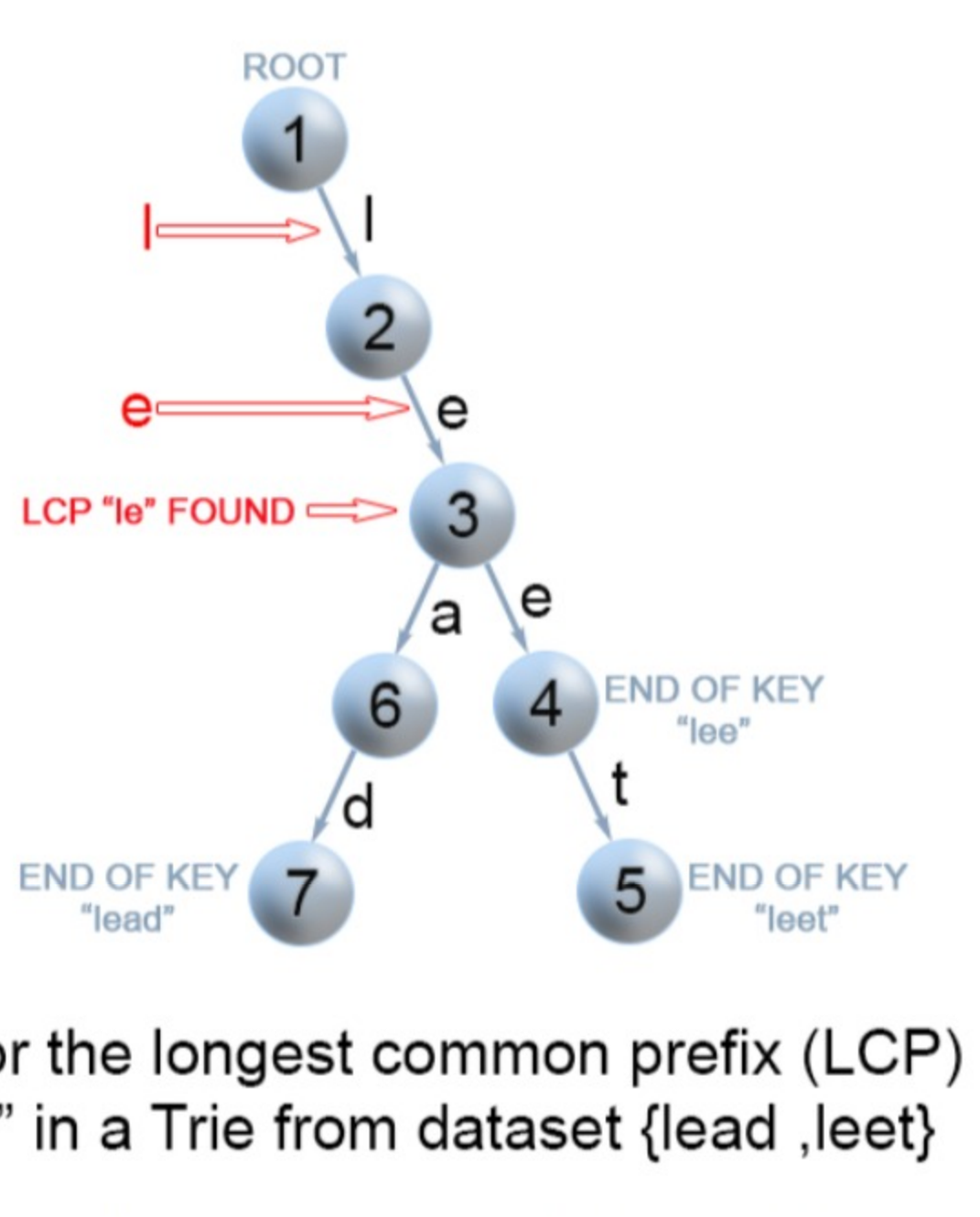

Figure 4. Finding the longest common prefix of strings using Trie

Java

```java
public String longestCommonPrefix(String q, String[] strs) {
    if (strs == null || strs.length == 0)
         return "";
    if (strs.length == 1)
         return strs[0];
    Trie trie = new Trie();
    for (int i = 1; i < strs.length ; i++) {
        trie.insert(strs[i]);
    }
    return trie.searchLongestPrefix(q);
}

class TrieNode {

    // R links to node children
    private TrieNode[] links;

    private final int R = 26;

    private boolean isEnd;

    // number of children non null links
    private int size;
    public void put(char ch, TrieNode node) {
        links[ch -'a'] = node;
        size++;
    }
```

Complexity Analysis

In the worst case query q has length m and it is equal to all n strings of the array.

- Time complexity : preprocessing $O(S)$, where S is the number of all characters in the array, LCP query $O(m)$.

  Trie build has $O(S)$ time complexity. To find the common prefix of q in the Trie takes in the worst case $O(m)$.

- Space complexity : $O(S)$. We only used additional S extra space for the Trie.

Rate this article: ★★★★★

Previous Next

Comments: 217 Sort By

Type comment here... (Markdown is supported)

Preview Post

desileetcoders ★337 April 15, 2018 2:00 AM
I pay premium cost which is 4 * netflix per month and leetcode people don't even vet solutions, this is beyond ridiculous
260 Share Reply
SHOW 14 REPLIES

pudgeywudgey ★66 October 22, 2018 10:37 AM
So which is the most efficient solution?
66 Share Reply
SHOW 4 REPLIES

grl_pwr ★56 September 5, 2018 12:14 PM Report
def longestCommonPrefix(self, strs):
"""
:type strs: List[str]
:rtype: str
Read More
53 Share Reply
SHOW 3 REPLIES

asofiamferreira ★102 January 28, 2019 9:22 PM Report
Python3
Runtime: 32 ms, faster than 100%
class Solution:
def longestCommonPrefix(self, strs):
Read More
102 Share Reply
SHOW 26 REPLIES

pye ★73 March 10, 2019 7:36 AM
Python 20ms
def longestCommonPrefix(self, strs):
prefix=""
if len(strs)==0: return prefix
Read More
41 Share Reply
SHOW 7 REPLIES

azimbabu ★133 August 2, 2018 12:15 PM Report
I think time complexity of binary search approach is not correct. It seems to me that it's O(S * log(m)) in the worst case when we have n strings and each one has length m. The algorithm runs binary search on length [1, 2, ... m].
31 Share Reply
SHOW 3 REPLIES

terrible_whiteboard ★626 May 19, 2020 5:58 PM
I made a video if anyone is having trouble understanding the solution (clickable link)
https://youtu.be/g5kH8EX4l-U
Read More
20 Share Reply

zhengzhicong ★294 November 8, 2018 6:02 PM Report
python3:
class Solution:
def longestCommonPrefix(self, strs):
"""
Read More
27 Share Reply
SHOW 9 REPLIES

ScottLee1314 ★16 August 19, 2018 8:41 PM Report
Question of the complexity 'Approach 3: Divide and conquer'
T(n) = 2·T(n/2)+O(m)
T(n) = 2^xT(n/2^x) + xO(m)
let n = 2^x, x = log(2)(n)= log(n)
Read More
15 Share Reply
SHOW 7 REPLIES

jerry563 ★19 March 18, 2019 7:52 PM Report
I'm really confused of this test case:
Input
["ca","a"]
Output
"a"
Read More
19 Share Reply
SHOW 5 REPLIES

‹ 1 2 3 4 5 6 ... 21 22 ›