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Given an unsorted array of integers, find the length of the longest consecutive elements sequence.

128. Longest Consecutive Sequence 💆

Your algorithm should run in O(n) complexity.

Input: [100, 4, 200, 1, 3, 2]

Example:

```
Output: 4
Explanation: The longest consecutive elements sequence is [1, 2, 3,
4]. Therefore its length is 4.
```

The brute force algorithm does not do anything clever - it just considers each number in nums, attempting to count as high as possible from that number using only numbers in nums. After it counts too high (i.e.

currentNum refers to a number that nums does not contain), it records the length of the sequence if it is

# larger than the current best. The algorithm is necessarily optimal because it explores every possibility.

**С**ору Python Java class Solution: def longestConsecutive(self, nums): longest\_streak = 0 for num in nums: current\_num = num

```
while current_num + 1 in nums:
  10
                     current_num += 1
                     current_streak += 1
 11
 12
                 longest_streak = max(longest_streak, current_streak)
 13
 14
             return longest_streak
Complexity Analysis
   • Time complexity : O(n^3).
     The outer loop runs exactly n times, and because currentNum increments by 1 during each iteration
     of the while loop, it runs in O(n) time. Then, on each iteration of the while loop, an O(n) lookup
     in the array is performed. Therefore, this brute force algorithm is really three nested O(n) loops, which
     compound multiplicatively to a cubic runtime.
```

# the sequence (i.e. nums[i] == nums[i-1] + 1). If it does, then we add to our current count and continue.

Algorithm

Otherwise, the sequence is broken, so we record our current sequence and reset it to 1 (to include the number that broke the sequence). It is possible that the last element of **nums** is part of the longest sequence, so we return the maximum of the current sequence and the longest one.

Python

class Solution:

if not nums:

nums.sort()

return 0

longest\_streak = 1 current\_streak = 1

def longestConsecutive(self, nums):

for i in range(1, len(nums)):

if nums[i] != nums[i-1]:

if nums[i] == nums[i-1]+1:

current\_streak += 1

Approach 3: HashSet and Intelligent Sequence Building

Java

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Here, an example array is sorted before the linear scan identifies all consecutive sequences. The longest sequence is colored in red.

18 return max(longest\_streak, current\_streak) 19

```
15
                    else:
                        longest_streak = max(longest_streak, current_streak)
 16
                        current streak = 1
 17
Complexity Analysis
   • Time complexity : O(nlgn).
     The main for loop does constant work n times, so the algorithm's time complexity is dominated by
     the invocation of sort , which will run in O(nlgn) time for any sensible implementation.
   • Space complexity : O(1) (or O(n)).
     For the implementations provided here, the space complexity is constant because we sort the input
     array in place. If we are not allowed to modify the input array, we must spend linear space to store a
     sorted copy.
```

### It turns out that our initial brute force solution was on the right track, but missing a few optimizations necessary to reach O(n) time complexity. Algorithm

Intuition

This optimized algorithm contains only two changes from the brute force approach: the numbers are stored in a <code>HashSet</code> (or <code>Set</code> , in Python) to allow O(1) lookups, and we only attempt to build sequences from

would necessarily be part of a longer sequence.

 $num_set = set(nums)$ 

for num in num\_set:

if num - 1 not in num\_set: current\_num = num current\_streak = 1

while current\_num + 1 in num\_set:

longest\_streak = max(longest\_streak, current\_streak)

current\_num += 1

current\_streak += 1

Python Java class Solution: def longestConsecutive(self, nums): longest streak = 0

### 16 return longest\_streak 17

6

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**Complexity Analysis** • Time complexity : O(n).

solution.

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loop, closer inspection reveals it to be linear. Because the while loop is reached only when currentNum marks the beginning of a sequence (i.e. currentNum-1 is not present in nums), the while loop can only run for n iterations throughout the entire runtime of the algorithm. This means that despite looking like  $O(n \cdot n)$  complexity, the nested loops actually run in O(n+n) = O(n)time. All other computations occur in constant time, so the overall runtime is linear. • Space complexity : O(n). In order to set up O(1) containment lookups, we allocate linear space for a hash table to store the O(n) numbers in  $rac{ extsf{nums}}{ extsf{nums}}$  . Other than that, the space complexity is identical to that of the brute force

Comments: 59 Sort By ▼ Type comment here... (Markdown is supported) Preview Post arboleda1298 ★ 182 ② March 14, 2019 9:07 PM Why is the tag for this Union-Find? 150 ∧ ∨ ♂ Share ★ Reply SHOW 11 REPLIES cryptoandy 🖈 164 🗿 December 6, 2018 3:05 AM For those who got confused by if the last solution is  $O(n^2)$  or O(n), please take a close look at the entering of the logic: if(!num\_set.contains(num-1)). That means, for example, 6,5,4,3,2,1 input, only the value 1 is valid for the loop(all other values have its value - 1 in the set), that is O(n).

Another corner example, 2, 5, 6, 7, 9, 11. All of these numbers are the "entrance" for the logic but the

I made a video if anyone is having trouble understanding the solution (clickable link)

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If the list is small, then O(nlogn) shouldn't make much of a difference either. It's a stupid question.

**ABattle** ★ 48 ② January 17, 2019 3:16 AM Can be reduced to a simple dfs/bfs

42 A V Share Seply

https://youtu.be/xdMyL--dOqE

20 \land 🗸 🗗 Share 🦘 Reply

public class Solution {

wcarvalho ★ 36 ② January 9, 2018 2:19 AM

**AnishSGM** ★ 23 **②** May 22, 2018 7:37 PM

Albert2522 ★ 31 ② January 1, 2018 1:51 PM

myself in particular) it wasn't obvious that we spend O(n).

I prefer using a HashMap here in the following way:

25 ∧ ∨ ♂ Share ★ Reply

num, it have to go through the entire set so far?

terrible\_whiteboard 🛊 627 🗿 May 19, 2020 6:24 PM

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public int longestConsecutive(int[] nums) {

Set (Integer) seen = new HashSet ()().

21 A V 🗗 Share 🦘 Reply **SHOW 2 REPLIES** 

**SHOW 7 REPLIES** 

Read More 9 🔨 🗸 Chare 👆 Reply defunator ★ 5 ② February 2, 2019 12:45 PM Excuse me, but as I know, HashSet works in AVERAGE O(1) for lookup. So it is not O(n), it is expected O(n)(in some cases the constant is too big, so it is more appropriate to

1. Firstly, put every element in a HashMap and mark it as unexplored.

use set or something else with log but more determined). 5 A V C Share Reply SHOW 1 REPLY

anmingyu11 ★ 425 ② December 25, 2017 12:35 PM class Solution { public int longestConsecutive(int[] nums) { final HashSet mySet = new HashSet(); for (int i : nums) mySet.add(i);

3 A V C Share Reply (123456)

Why is this hard? Easier than many medium questions

afsf11 ★ 10 ② July 24, 2019 6:00 PM

4 A V C Share Reply

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Approach 1: Brute Force Intuition Because a sequence could start at any number in nums, we can exhaust the entire search space by building as long a sequence as possible from every number. Algorithm

current\_streak = 1

• Space complexity : O(1). The brute force algorithm only allocates a handful of integers, so it uses constant additional space. Approach 2: Sorting Intuition If we can iterate over the numbers in ascending order, then it will be easy to find sequences of consecutive numbers. To do so, we can sort the array.

Before we do anything, we check for the base case input of the empty array. The longest sequence in an empty array is, of course, 0, so we can simply return that. For all other cases, we sort nums and consider each number after the first (because we need to compare each number to its previous number). If the current number and the previous are equal, then our current sequence is neither extended nor broken, so we simply move on to the next number. If they are unequal, then we must check whether the current number extends

[9, 1, 4, 7, 3, -1, 0, 5, 8, -1, 6]

**С**ору

numbers that are not already part of a longer sequence. This is accomplished by first ensuring that the number that would immediately precede the current number in a sequence is not present, as that number **С**ору

Although the time complexity appears to be quadratic due to the while loop nested within the for

Next 👀

wouldn't [100,99, 98, ..., 1] (i.e. an array in reverse order) make this run in \$\$O(n^2)\$\$ because for every

If the list is very long, then the O(n) hash-set solution is a drain on memory so, not an elegant solution.

I think, using a HashSet is a bit confusing in terms of the time complexity. For many people (and to

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