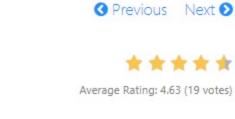
57. Insert Interval

May 14, 2019 | 34.5K views



6 9 6

You may assume that the intervals were initially sorted according to their start times.

Given a set of non-overlapping intervals, insert a new interval into the intervals (merge if necessary).

Example 1:

```
Input: intervals = [[1,3],[6,9]], newInterval = [2,5]
 Output: [[1,5],[6,9]]
Example 2:
```

Input: intervals = [[1,2],[3,5],[6,7],[8,10],[12,16]], newInterval = [4,8]

method signature.

```
Output: [[1,2],[3,10],[12,16]]
  Explanation: Because the new interval [4,8] overlaps with [3,5],[6,7],[8,10].
NOTE: input types have been changed on April 15, 2019. Please reset to default code definition to get new
```

Approach 1: Greedy. **Greedy algorithms**

Solution

Greedy problems usually look like "Find minimum number of something to do something" or "Find maximum number of something to fit in some conditions", and typically propose an unsorted input.

globally optimal solution.

The idea of greedy algorithm is to pick the locally optimal move at each step, that will lead to the

ullet Figure out how to sort the input data ($\mathcal{O}(N\log N)$ time). That could be done directly by a sorting or indirectly by a heap usage. Typically sort is better than the heap usage because of gain in space. • Parse the sorted input to have a solution ($\mathcal{O}(N)$ time).

Please notice that in case of well-sorted input one doesn't need the first part and the greedy solution could have $\mathcal{O}(N)$ time complexity, here is an example.

The standard solution has $\mathcal{O}(N \log N)$ time complexity and consists of two parts:

How to prove that your greedy algorithm provides globally optimal solution?

Intuition

Let's consider the following intervals

complexity should be $\mathcal{O}(N)$.

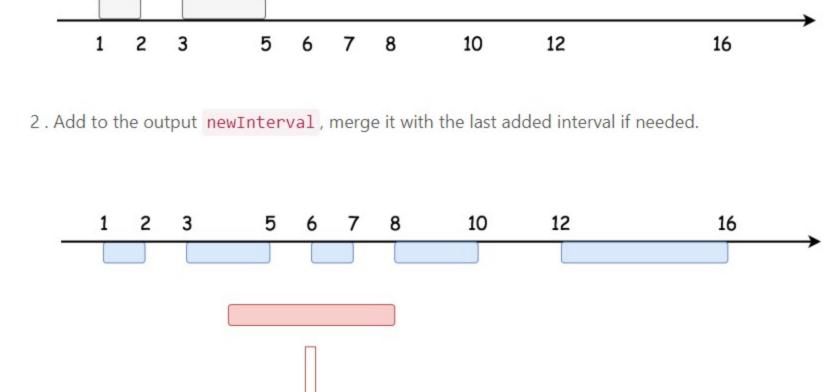
Usually you could use the proof by contradiction.

10 12 16

Here we have an example of a greedy problem with a well-sorted input, and hence the algorithm time

The straightforward one-pass strategy could be implemented in three steps.

1 . Add to the output all the intervals starting before newInterval . 12 10 16



16

16

Сору

5 10 12 2 3 8 1

8

3 . Add the next intervals one by one, merge if needed.

2

4

9

10

11 12

13 14

15

16 17

18 19

20

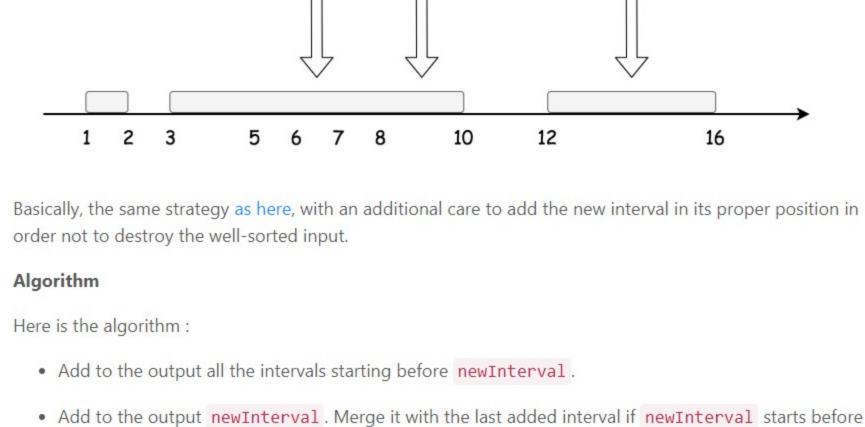
21 22

23 24

25

26

3



10

12

new_start, new_end = newInterval 5 idx, n = 0, len(intervals) 6 output = [] 7 # add all intervals starting before newInterval 8

while idx < n and new_start > intervals[idx][0]:

if there is no overlap, just add the interval if not output or output[-1][1] < new_start:

if there is an overlap, merge with the last interval

add next intervals, merge with newInterval if needed

if there is no overlap, just add an interval

 $output[-1][1] = max(output[-1][1], new_end)$

output.append(intervals[idx])

output.append(newInterval)

interval = intervals[idx]

Ummm... why is this a hard problem?

girishiitj 🖈 115 🗿 September 24, 2019 4:47 PM

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start, end = interval

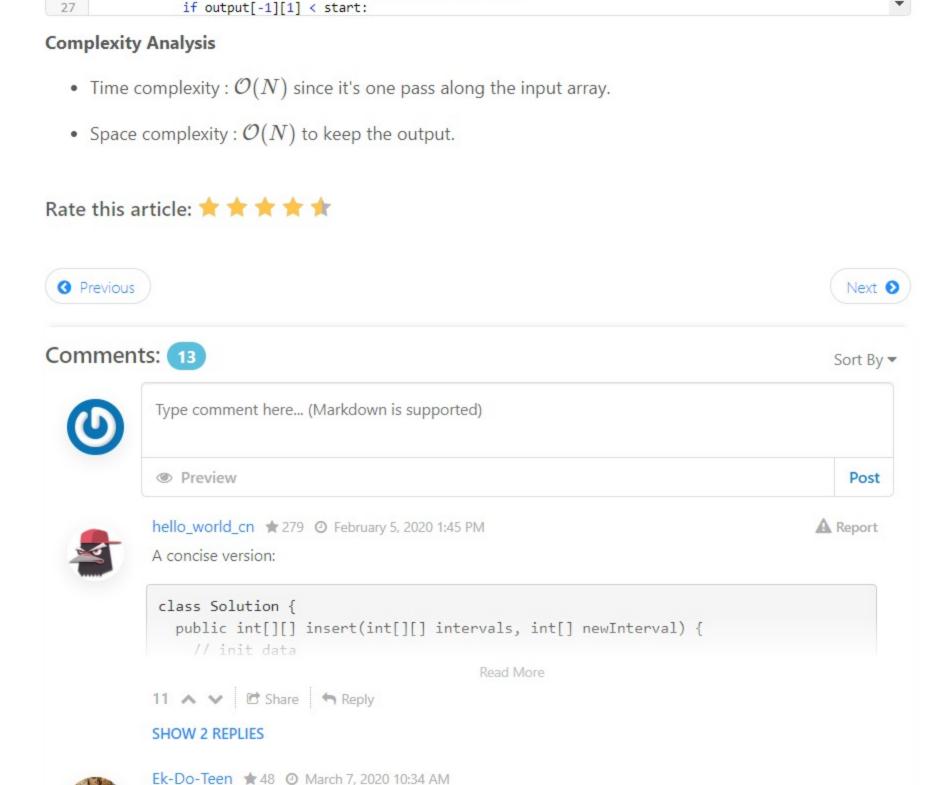
idx += 1

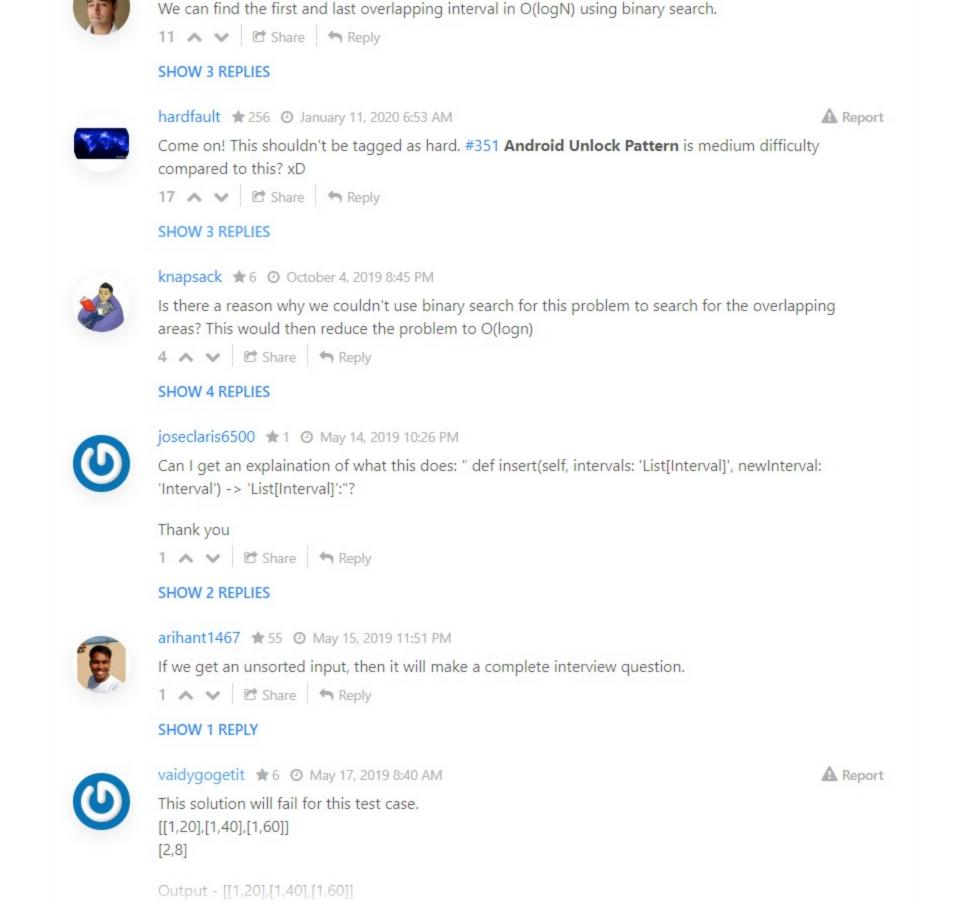
add newInterval

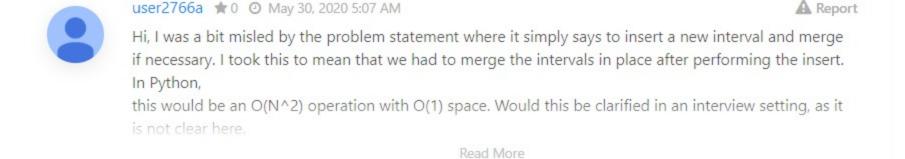
while idx < n:

idx += 1

the last added interval. Add the next intervals one by one. Merge with the last added interval if the current interval starts before the last added interval. Implementation Python Java class Solution: def insert(self, intervals: 'List[Interval]', newInterval: 'Interval') -> 'List[Interval]': 2 3 # init data







didn't get the hard tag on that, I assumed there exists some superclever thing behind the lock, and not

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this one.

Username1604 ★ 42 ② May 21, 2020 6:04 PM

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