

478. Generate Random Point in a Circle

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Given the radius and x-y positions of the center of a circle, write a function `randPoint` which generates a uniform random point in the circle.

Note:

- input and output values are in [floating-point](#).
- radius and x-y position of the center of the circle is passed into the class constructor.
- a point on the circumference of the circle is considered to be in the circle.
- `randPoint` returns a size 2 array containing x-position and y-position of the random point, in that order.

Example 1:

```
Input:
["Solution", "randPoint", "randPoint", "randPoint"]
[[1, 0, 0], [], [], []]
Output: [null, [-0.72939, -0.65505], [-0.78502, -0.28626], [-0.83119, -0.19803]]
```

Example 2:

```
Input:
["Solution", "randPoint", "randPoint", "randPoint"]
[[10, 5, -7.5], [], [], []]
Output: [null, [11.52438, -8.33273], [2.46992, -16.21705], [11.13430, -12.42337]]
```

Explanation of Input Syntax:

The input is two lists: the subroutines called and their arguments. `Solution`'s constructor has three arguments, the radius, x-position of the center, and y-position of the center of the circle. `randPoint` has no arguments. Arguments are always wrapped with a list, even if there aren't any.

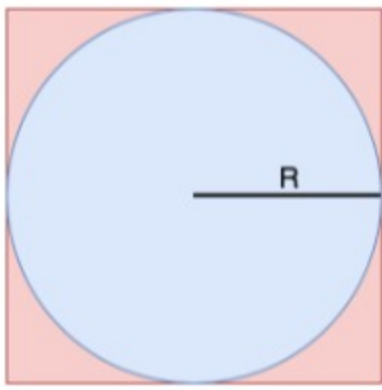
Solution

Approach 1: Rejection Sampling

Intuition

It is easy to generate a random point in a square. Could we use randomly generated points in a square to get random points in a circle? Which generated points could we use, and which ones would we need to toss away? How often would we generate points that we could use?

Algorithm



A square of size length $2R$ overlaid with a circle of radius R .

To get uniform random points in a circle C of radius R , we can generate uniform random points in the square S of side length $2R$, keeping all of the points which are at most [euclidean distance](#) R from the center, and rejecting all which are farther away than that. This technique is called [rejection sampling](#). Each possible location on the circle has the same probability of being generated, so the sampling of points will be uniformly distributed.

The area of the square is $(2R)^2 = 4R$ and the area of the circle is $\pi R \approx 3.14R$. $\frac{3.14R}{4R} = \frac{3.14}{4} = .785$. Therefore, we will get a usable sample approximately 78.5% of the time and the expected number of times that we will need to sample until we get a usable sample is $\frac{1}{.785} \approx 1.274$ times.

```
C++JavaCopy
1 class Solution {
2 public:
3     double rad, xc, yc;
4     //c++11 random floating point number generation
5     mt19937 rng(random_device{}());
6     uniform_real_distribution<double> uni(0, 1);
7
8     Solution(double radius, double x_center, double y_center) {
9         rad = radius, xc = x_center, yc = y_center;
10    }
11
12    vector<double> randPoint() {
13        double x0 = xc - rad;
14        double y0 = yc - rad;
15
16        while(true) {
17            double xg = x0 + uni(rng) * 2 * rad;
18            double yg = y0 + uni(rng) * 2 * rad;
19            if (sqrt(pow((xg - xc), 2) + pow((yg - yc), 2)) <= rad)
20                return {xg, yg};
21        }
22    }
23};
```

Complexity Analysis

- Time Complexity: $O(1)$ on average. $O(\infty)$ worst case. (per `randPoint` call)
- Space Complexity: $O(1)$.

Approach 2: Inverse Transform Sampling (Math)

Disclaimer

This solution relies on advanced math which is not expected knowledge for a coding interview. It is presented here only for educational purposes.

Algorithm

Assume that we are given a circle C of radius 1 that is centered at the [origin](#), and our task is to sample uniform random points on this circle.

Lets imagine another circle B of radius $\frac{1}{2}$ which is also centered at the origin.

The circumference of C is twice the circumference of B , because the circumference of a circle is [directly proportional](#) to the radius. Also, the probability of sampling a point from a subregion in circle C is directly proportional to the area of the subregion. Therefore, the probability of sampling a point on the perimeter of C is twice that of sampling a point on the perimeter of B .

More generally, what is implied is that the sampling probability is directly proportional to the distance from the origin, from 0 up to R . This can be modeled as a [probability density function](#) f , where x is the distance from the origin and $f(x)$ is the relative sampling probability at x .

The area under any probability density function curve must be 1. Therefore, the equation must be $f(x) = 2x$.

Using our probability density function f , we can compute the [cumulative distribution function](#) F , where $F(x)$ is the probability of sampling a point within a distance of x from the origin.

The cumulative distribution function is the integral of the probability density function.

$$F(x) = \int f(x) = \int 2x = x^2$$

Lastly, we can use our cumulative distribution function F to compute the inverse cumulative distribution function F^{-1} , which accepts uniform random value between 0 and 1 and returns a random distance from origin in accordance with f .^[1]

$$\begin{aligned} F^{-1}(F(x)) &= x \\ F^{-1}(x^2) &= x \\ F^{-1}(x) &= \sqrt{x} \end{aligned}$$

Now, to generate a uniform random point on C , we just need to compute a random distance D from origin using F^{-1} and a uniform random angle θ over the range $[0, 2 \cdot \pi)$.

The points will be generated as [polar coordinates](#). To convert to [cartesian coordinates](#), we can use the following formulas.

$$\begin{aligned} X &= D \cdot \cos(\theta) \\ Y &= D \cdot \sin(\theta) \end{aligned}$$

```
C++JavaCopy
1 class Solution {
2 public:
3     double rad, xc, yc;
4     //c++11 random floating point number generation
5     mt19937 rng(random_device{}());
6     uniform_real_distribution<double> uni(0, 1);
7
8     Solution(double radius, double x_center, double y_center) {
9         rad = radius, xc = x_center, yc = y_center;
10    }
11
12    vector<double> randPoint() {
13        double d = rad * sqrt(uni(rng));
14        double theta = uni(rng) * (2 * M_PI);
15        return {d * cos(theta) + xc, d * sin(theta) + yc};
16    }
17};
```

Complexity Analysis

- Time Complexity: $O(1)$ per `randPoint` call.
- Space Complexity: $O(1)$

Footnotes


- ^[1] This technique of using the inverse cumulative distribution function to sample numbers at random from the corresponding probability distribution is called [inverse transform sampling](#).
- This solution is inspired by [this](#) answer on Stack Overflow.

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