

236. Lowest Common Ancestor of a Binary Tree

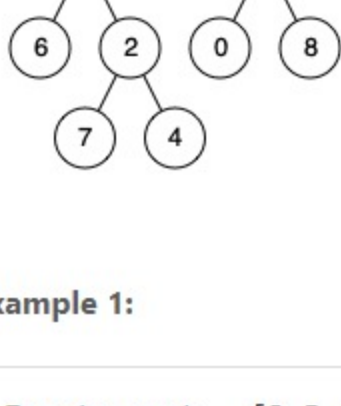
Nov. 15, 2018 | 215.2K views

Average Rating 4.09 (209 votes)

Given a binary tree, find the lowest common ancestor (LCA) of two given nodes in the tree.

According to the [definition of LCA on Wikipedia](#): "The lowest common ancestor is defined between two nodes p and q as the lowest node in T that has both p and q as descendants (where we allow **a node to be a descendant of itself**)."

Given the following binary tree: root = [3,5,1,6,2,0,8,null,null,7,4],



Example 1:

Input: root = [3,5,1,6,2,0,8,null,null,7,4], p = 5, q = 1
Output: 3
Explanation: The LCA of nodes 5 and 1 is 3.

Example 2:

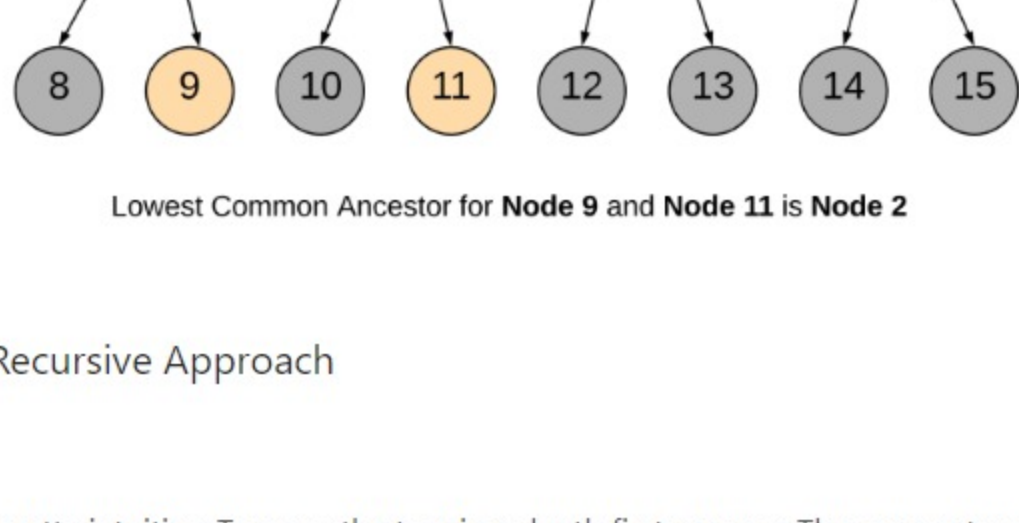
Input: root = [3,5,1,6,2,0,8,null,null,7,4], p = 5, q = 4
Output: 5
Explanation: The LCA of nodes 5 and 4 is 5, since a node can be a descendant of itself according to the LCA definition.

Note:

- All of the nodes' values will be unique.
- p and q are different and both values will exist in the binary tree.

Solution

First the given nodes p and q are to be searched in a binary tree and then their lowest common ancestor is to be found. We can resort to a normal tree traversal to search for the two nodes. Once we reach the desired nodes p and q , we can backtrack and find the lowest common ancestor.



Approach 1: Recursive Approach

Intuition

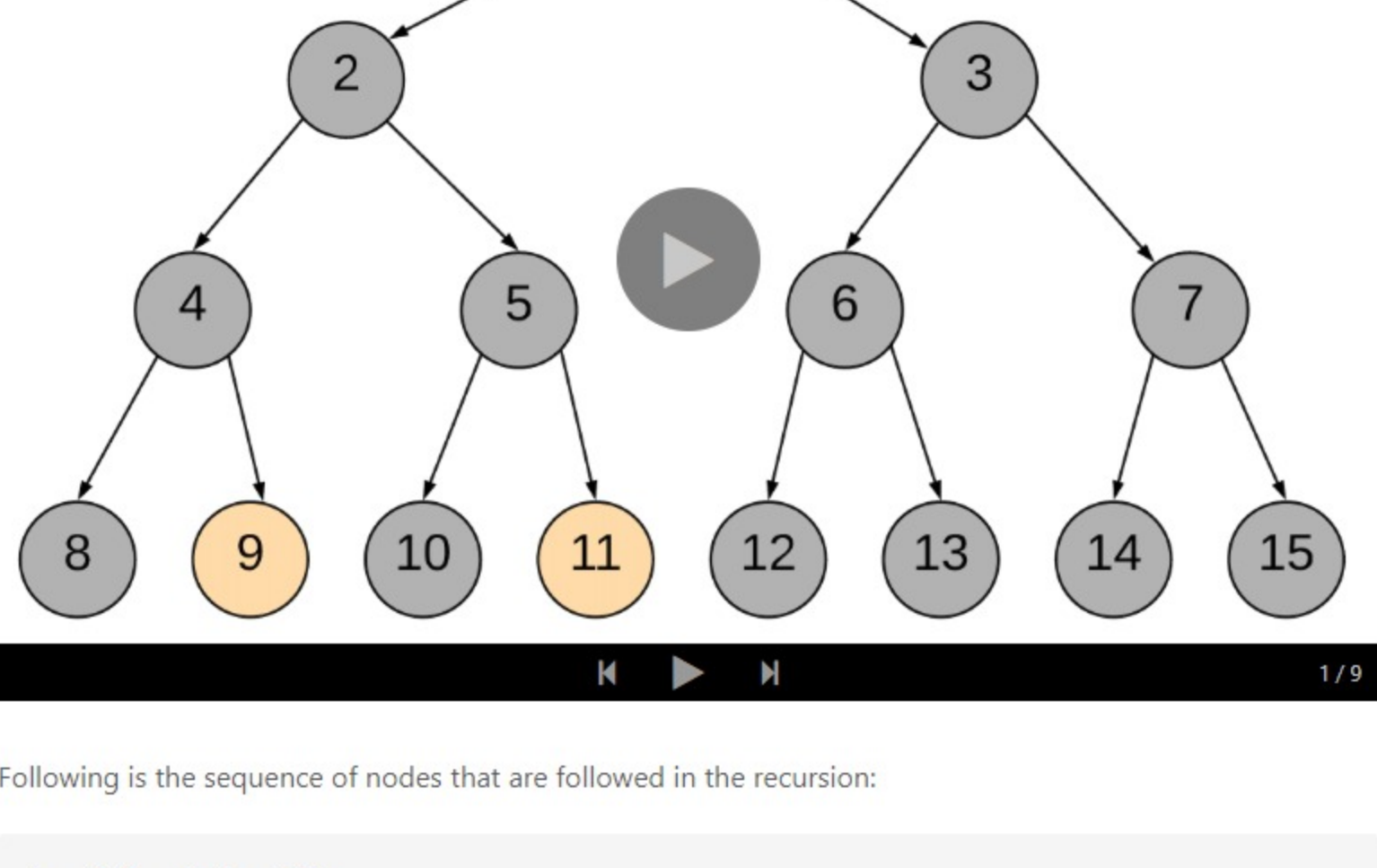
The approach is pretty intuitive. Traverse the tree in a depth first manner. The moment you encounter either of the nodes p or q , return some boolean flag. The flag helps to determine if we found the required nodes in any of the paths. The least common ancestor would then be the node for which both the subtree recursions return a **True** flag. It can also be the node which itself is one of p or q and for which one of the subtree recursions returns a **True** flag.

Let us look at the formal algorithm based on this idea.

Algorithm

- Start traversing the tree from the root node.
- If the current node itself is one of p or q , we would mark a variable **mid** as **True** and continue the search for the other node in the left and right branches.
- If either of the left or the right branch returns **True**, this means one of the two nodes was found below.
- If at any point in the traversal, any two of the three flags **left**, **right** or **mid** become **True**, this means we have found the lowest common ancestor for the nodes p and q .

Let us look at a sample tree and we search for the lowest common ancestor of two nodes 9 and 11 in the tree.



Following is the sequence of nodes that are followed in the recursion:

```
1 --> 2 --> 4 --> 8
BACKTRACK 8 --> 4
4 --> 9 (ONE NODE FOUND, return True)
BACKTRACK 9 --> 4 --> 2
2 --> 5 --> 10
BACKTRACK 10 --> 5
5 --> 11 (ANOTHER NODE FOUND, return True)
BACKTRACK 11 --> 5 --> 2
```

2 is the node where we have left = True and right = True and hence it is the lowest cc

```
1 class Solution:
2
3     def __init__(self):
4         # Variable to store LCA node.
5         self.ans = None
6
7     def lowestCommonAncestor(self, root, p, q):
8         """
9         :type root: TreeNode
10        :type p: TreeNode
11        :type q: TreeNode
12        :rtype: TreeNode
13        """
14        def recurse_tree(current_node):
15
16            # If reached the end of a branch, return False.
17            if not current_node:
18                return False
19
20            # Left Recursion
21            left = recurse_tree(current_node.left)
22
23            # Right Recursion
24            right = recurse_tree(current_node.right)
25
26            # If the current node is one of p or q
27            mid = current_node == p or current_node == q
```

Complexity Analysis

- Time Complexity: $O(N)$, where N is the number of nodes in the binary tree. In the worst case we might be visiting all the nodes of the binary tree.
- Space Complexity: $O(N)$. This is because the maximum amount of space utilized by the recursion stack would be N since the height of a skewed binary tree could be N .

Approach 2: Iterative using parent pointers

Intuition

If we have parent pointers for each node we can traverse back from p and q to get their ancestors. The first common node we get during this traversal would be the LCA node. We can save the parent pointers in a dictionary as we traverse the tree.

Algorithm

- Start from the root node and traverse the tree.
- Until we find p and q both, keep storing the parent pointers in a dictionary.
- Once we have found both p and q , we get all the ancestors for p using the parent dictionary and add to a set called **ancestors**.
- Similarly, we traverse through ancestors for node q . If the ancestor is present in the ancestors set for p , this means this is the first ancestor common between p and q (while traversing upwards) and hence this is the LCA node.

```
1 class Solution:
2
3     def lowestCommonAncestor(self, root, p, q):
4         """
5         :type root: TreeNode
6         :type p: TreeNode
7         :type q: TreeNode
8         :rtype: TreeNode
9         """
10
11        # Stack for tree traversal
12        stack = [root]
13
14        # Dictionary for parent pointers
15        parent = {root: None}
16
17        # Iterate until we find both the nodes p and q
18        while p not in parent or q not in parent:
19
20            node = stack.pop()
21
22            # While traversing the tree, keep saving the parent pointers.
23            if node.left:
24                parent[node.left] = node
25            if node.right:
26                parent[node.right] = node
27
28            stack.append(node)
```

Complexity Analysis

- Time Complexity: $O(N)$, where N is the number of nodes in the binary tree. In the worst case we might be visiting all the nodes of the binary tree.
- Space Complexity: $O(N)$. In the worst case space utilized by the stack, the parent pointer dictionary and the ancestor set, would be N each, since the height of a skewed binary tree could be N .

Approach 3: Iterative without parent pointers

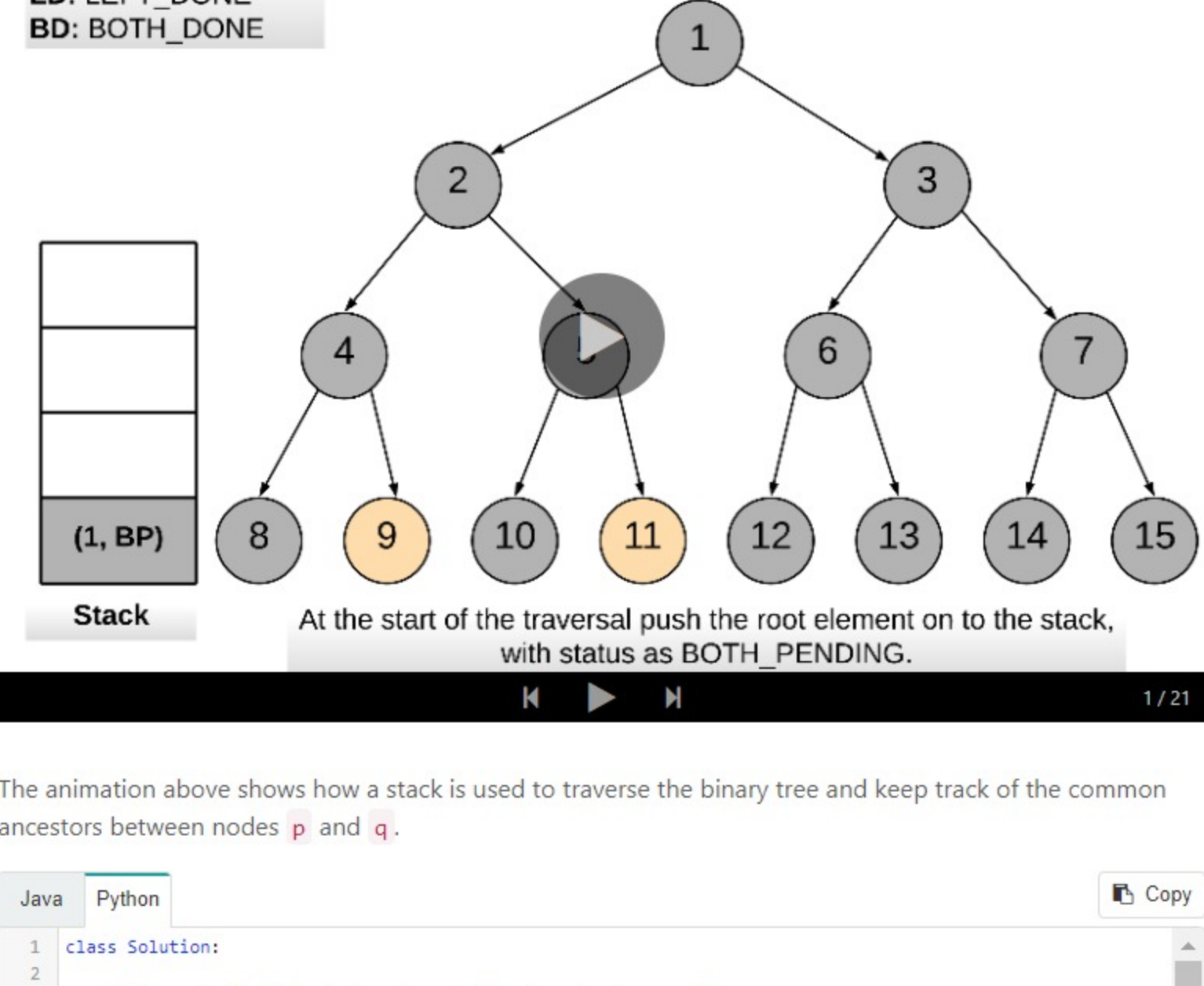
Intuition

In the previous approach, we come across the LCA during the backtracking process. We can get rid of the backtracking process itself. In this approach we always have a pointer to the probable LCA and the moment we find both the nodes we return the pointer as the answer.

Algorithm

- Start with root node.
- Put the (**root**, **root_state**) on to the stack. **root_state** defines whether one of the children or both children of **root** are left for traversal.
- While the stack is not empty, peek into the top element of the stack represented as (**parent_node**, **parent_state**).
- Before traversing any of the child nodes of **parent_node** we check if the **parent_node** itself is one of p or q .
- First time we find either of p or q , set a boolean flag called **one_node_found** to **True**. Also start keeping track of the lowest common ancestors by keeping a note of the top index of the stack in the variable **LCA_index**. Since all the current elements of the stack are ancestors of the node we just found.
- The second time **parent_node == p** or **parent_node == q** it means we have found both the nodes and we can return the **LCA_node**.
- Whenever we visit a child of a **parent_node** we push the (**parent_node**, **updated_parent_state**) onto the stack. We update the state of the parent since a child/branch has been visited/processed and accordingly the state changes.
- A node finally gets popped off from the stack because the state becomes **BOTH_DONE** implying both left and right subtrees have been pushed onto the stack and processed. If **one_node_found** is **True** then we need to check if the top node being popped could be one of the ancestors of the found node. In that case we need to reduce **LCA_index** by one. Since one of the ancestors was popped off.

Whenever both p and q are found, **LCA_index** would be pointing to an index in the stack which would contain all the common ancestors between p and q . And the **LCA_index** element has the **lowest** ancestor common between p and q .



The animation above shows how a stack is used to traverse the binary tree and keep track of the common ancestors between nodes p and q .

```
1 class Solution:
2
3     # Three static flags to keep track of post-order traversal.
4
5     # Both left and right traversal pending for a node.
6     BOTH_PENDING = 2
7     # Left traversal done.
8     LEFT_DONE = 1
9     # Both left and right traversal done for a node.
10    # Indicates the node can be popped off the stack.
11    BOTH_DONE = 0
12
13    def lowestCommonAncestor(self, root, p, q):
14        """
15        :type root: TreeNode
16        :type p: TreeNode
17        :type q: TreeNode
18        :rtype: TreeNode
19        """
20
21        # Initialize the stack with the root node.
22        stack = [(root, Solution.BOTH_PENDING)]
23
24        # This flag is set when either one of p or q is found.
25        one_node_found = False
26
27        while stack:
```

Complexity Analysis


- Time Complexity: $O(N)$, where N is the number of nodes in the binary tree. In the worst case we might be visiting all the nodes of the binary tree. The advantage of this approach is that we can prune backtracking. We simply return once both the nodes are found.
- Space Complexity: $O(N)$. In the worst case the space utilized by stack would be N since the height of a skewed binary tree could be N .

Rate this article: ★★★★★


PreviousNext

Comments: 53

Sort By ▾

- 

Type comment here...(Markdown is supported)


PreviewPost
- 

Hi Guys,

I have been monitoring the response the article is receiving in the form of ratings and I see a downward trend. As an author, I feel it is my responsibility to make sure the article is easily understandable and has all the optimal solutions. I would love to have some sort of feedback from you, the readers as to what can be improved in the article and what problems do you see currently. Thank you!

Read More


95 ▾ ▾ ▾ Share ▾ Reply

SHOW 22 REPLIES
- 

```
class Solution(object):
    def lowestCommonAncestor(self, root, p, q):
        if root is None:
            return None
```

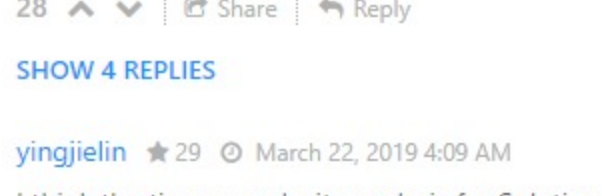
Read More

20 ▾ ▾ ▾ Share ▾ Reply

SHOW 1 REPLY
- 


NideeshT ★ 589 ▾ July 16, 2019 6:23 AM

Java Code + Whiteboard Youtube Video Explanation accepted -<https://www.youtube.com/watch?v=uKhLoNaG9LI> (clickable link)



Read More


28 ▾ ▾ ▾ Share ▾ Reply

SHOW 4 REPLIES
- 

yingjieln ★ 29 ▾ March 22, 2019 4:09 AM


I think the time complexity analysis for Solution 1 is not accurate. Since the recursiveTree function only returns a boolean, there is no way to tell whether both p and q are found in any subtree. Therefore, even if both p and q are found in the left subtree of a node, full efforts are still made to check the right subtree. Therefore, the time complexity is still O(N), but this is not the worst case, but every case.

18 ▾ ▾ ▾ Share ▾ Reply

SHOW 3 REPLIES
- 

```
class Solution(object):
    def lowestCommonAncestor(self, root, p, q):
        if root is None:
            return None
```


Read More

15 ▾ ▾ ▾ Share ▾ Reply
- 

softwareshortcut ★ 436 ▾ May 2, 2019 6:18 AM

1st solution is pretty well thought. It looks so simple but might be hard to come up with during an interview in only 10-15 minutes unless you've seen it before. Using the addition of **mid+left+right** groups a bunch of test cases together. My solution had the same approach but a lot more test cases.


9 ▾ ▾ ▾ Share ▾ Reply

SHOW 2 REPLIES
- 

```
def __init__(self):
    """
    Definition for a binary tree node.
    """
    public class TreeNode {
```

Read More


7 ▾ ▾ ▾ Share ▾ Reply

SHOW 2 REPLIES
- 

blueglaucus ★ 3 ▾ January 15, 2020 3:48 PM

Is it possible to do solution 1 without using an instance variable to store the LCA node?


1 ▾ ▾ ▾ Share ▾ Reply

SHOW 1 REPLY
- 

professor19 ★ 147 ▾ June 20, 2019 12:45 PM

I think the solution 2 approach when we add element in set it is log(n) rather than constant, so complexity is O(N logN)

1 ▾ ▾ ▾ Share ▾ Reply

SHOW 2 REPLIES
- 

kansalho06 ★ 18 ▾ July 2, 2019 12:17 PM

Just an edge case:
Solution 1 doesn't work if p and q both are same nodes, it will return NULL but is expected to return the same node.

0 ▾ ▾ ▾ Share ▾ Reply

SHOW 3 REPLIES