

# 160. Intersection of Two Linked Lists

March 9, 2016 | 217.5K views

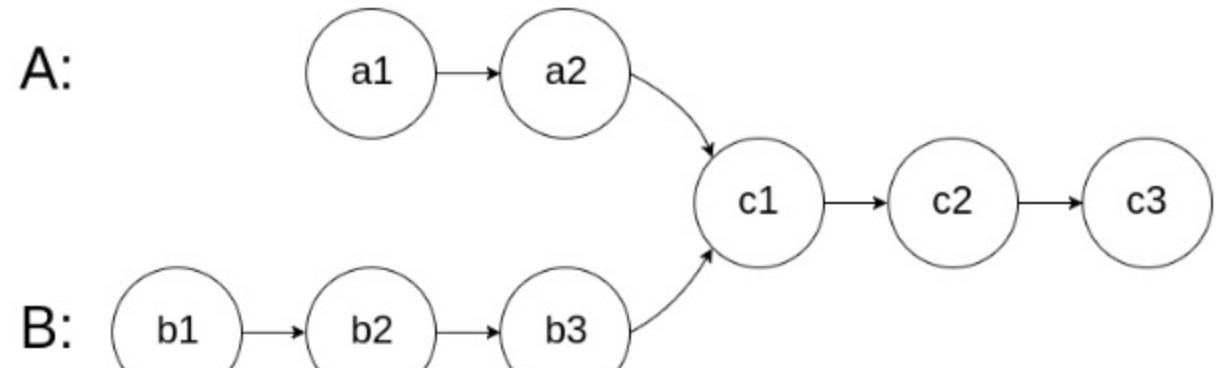
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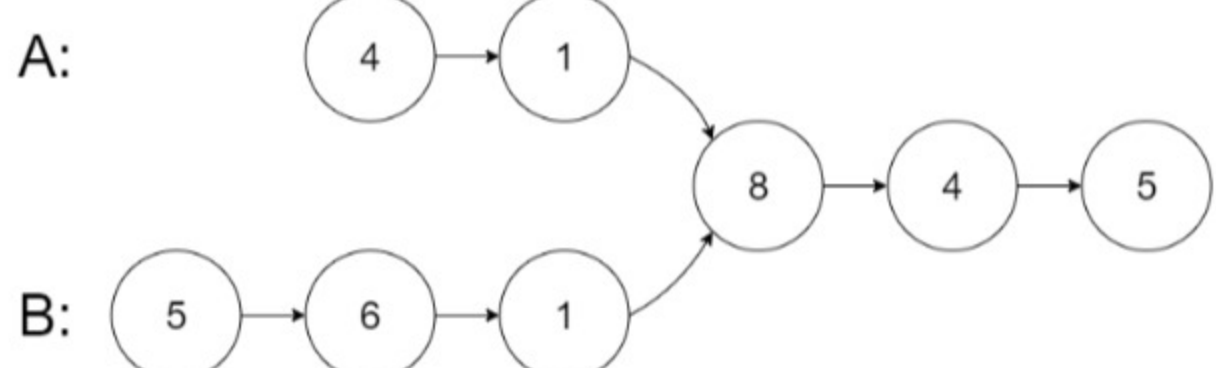
Write a program to find the node at which the intersection of two singly linked lists begins.

For example, the following two linked lists:



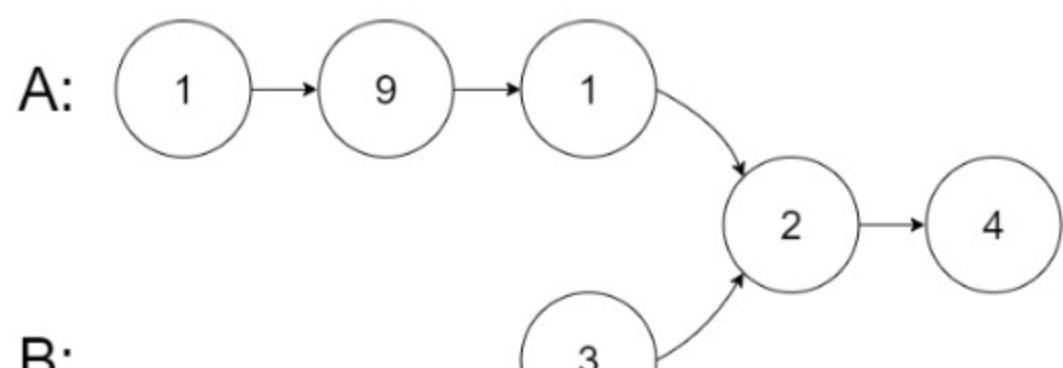
begin to intersect at node c1.

Example 1:



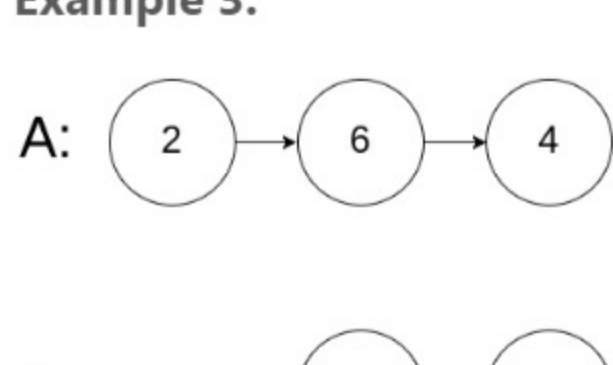
**Input:** intersectVal = 8, listA = [4,1,8,4,5], listB = [5,6,1,8,4,5], skipA = 2, skipB = 1  
**Output:** Reference of the node with value = 8  
**Input Explanation:** The intersected node's value is 8 (note that this must not be 0 if the lists intersect at node with value 8)

Example 2:



**Input:** intersectVal = 2, listA = [1,9,1,2,4], listB = [3,2,4], skipA = 3, skipB = 1  
**Output:** Reference of the node with value = 2  
**Input Explanation:** The intersected node's value is 2 (note that this must not be 0 if the lists intersect at node with value 2)

Example 3:



**Input:** intersectVal = 0, listA = [2,6,4], listB = [1,5], skipA = 3, skipB = 2  
**Output:** null  
**Input Explanation:** From the head of A, it reads as [2,6,4]. From the head of B, it reads as [1,5].  
**Explanation:** The two lists do not intersect, so return null.

Notes:

- If the two linked lists have no intersection at all, return `null`.
- The linked lists must retain their original structure after the function returns.
- You may assume there are no cycles anywhere in the entire linked structure.
- Each value on each linked list is in the range `[1, 10^9]`.
- Your code should preferably run in  $O(n)$  time and use only  $O(1)$  memory.

## Solution

### Approach 1: Brute Force

For each node  $a_i$  in list A, traverse the entire list B and check if any node in list B coincides with  $a_i$ .

**Complexity Analysis**

- Time complexity :  $O(mn)$ .
- Space complexity :  $O(1)$ .

### Approach 2: Hash Table

Traverse list A and store the address / reference to each node in a hash set. Then check every node  $b_i$  in list B: if  $b_i$  appears in the hash set, then  $b_i$  is the intersection node.

**Complexity Analysis**

- Time complexity :  $O(m + n)$ .
- Space complexity :  $O(m)$  or  $O(n)$ .

### Approach 3: Two Pointers

- Maintain two pointers  $pA$  and  $pB$  initialized at the head of A and B, respectively. Then let them both traverse through the lists, one node at a time.
- When  $pA$  reaches the end of a list, then redirect it to the head of B (yes, B, that's right!); similarly when  $pB$  reaches the end of a list, redirect it to the head of A.
- If at any point  $pA$  meets  $pB$ , then  $pA/pB$  is the intersection node.
- To see why the above trick would work, consider the following two lists: A = {1,3,5,7,9,11} and B = {2,4,9,11}, which are intersected at node '9'. Since B.length (=4) < A.length (=6),  $pB$  would reach the end of the merged list first, because  $pB$  traverses exactly 2 nodes less than  $pA$  does. By redirecting  $pB$  to head A, and  $pA$  to head B, we now ask  $pB$  to travel exactly 2 more nodes than  $pA$  would. So in the second iteration, they are guaranteed to reach the intersection node at the same time.
- If two lists have intersection, then their last nodes must be the same one. So when  $pA/pB$  reaches the end of a list, record the last element of A/B respectively. If the two last elements are not the same one, then the two lists have no intersections.

**Complexity Analysis**

- Time complexity :  $O(m + n)$ .
- Space complexity :  $O(1)$ .


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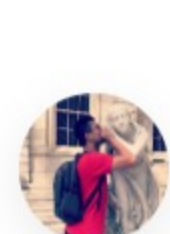
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 **TrafalgarZZZ** ★ 165 | August 28, 2018 3:07 PM  
Here's my solution which is like approach 3 but a little different. I store the size of ListA and ListB as len1 and len2. Then I reset the pointers to headA and headB and find the difference between len1 and len2, and then let the pointer of the longer list proceed by the difference between len1 and len2. Finally, traverse through the lists again, the intersection node can be easily found.

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 **jianchao-li** ★ 14342 | July 30, 2018 3:35 PM  
My accepted codes for the three approaches. Maybe someone could optimize them further and add them to the Solution? Thank you.

### Approach 1: Brute Force (527ms)

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 **lenchen1112** ★ 976 | September 13, 2018 8:21 PM

Solution 3 in Python:

```
# Definition for singly-linked list.
# class ListNode(object):
#     def __init__(self, x):
#         self.val = x
#         self.next = None
```

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 **Marcus\_X** ★ 125 | March 14, 2019 5:48 AM

You can do this by simply counting the lengths of both lists. Pass one: count lengths while searching to see if they have the same tail. Diff tails? Return NULL. Pass two: In an ideal world the lists would be the same length and simply walking them forward until they had the same pointer would be enough. To deal with uneven lists skip the longer list ahead X nodes, where X = abs(lenA-lenB). Now if you walk both lists forward one at a time they meet at the intersection point. This solution is still O(N) time and

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 **Anow** ★ 26 | March 19, 2019 11:57 AM

Is there any problem with example 1? Why not 1 but 8 ?

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 **terrible\_whiteboard** ★ 627 | May 19, 2020 6:20 PM

I made a video if anyone is having trouble understanding the solution (clickable link)

<https://youtu.be/c7dOI-hDa2Q>



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 **jkvavadiya** ★ 17 | October 19, 2018 1:16 AM

```
public ListNode getIntersectionNode(ListNode headA, ListNode headB) {
    ListNode pA = headA, pB = headB;
    while (pA != pB) {
        pA = pA != null ? pA.next : headB;
        pB = pB != null ? pB.next : headA;
    }
    return pA;
}
```

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 **qiyu8290** ★ 45 | October 27, 2018 11:54 PM

similar idea with counting A, B length:

```
public class Solution {
    public ListNode getIntersectionNode(ListNode headA, ListNode headB) {
        // ...
    }
}
```

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 **lu8** ★ 9 | December 19, 2017 1:02 PM

this question does not cover the following case:  
if headA and headB has no intersection, the solution will go to dead loop.

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 **RubbiShThird** ★ 9 | August 29, 2018 10:21 AM

reverse either list, if two list have the intersection, we can reach the other list from the other list.

```
class Solution(object):
    def getIntersectionNode(self, headA, headB):
        # ...

```

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