

# 161. One Edit Distance

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Given two strings **s** and **t**, determine if they are both one edit distance apart.

## Note:

There are 3 possiblities to satisfy one edit distance apart:

1. Insert a character into **s** to get **t**
2. Delete a character from **s** to get **t**
3. Replace a character of **s** to get **t**

## Example 1:

**Input:** **s** = "ab", **t** = "acb"  
**Output:** true  
**Explanation:** We can insert 'c' into **s** to get **t**.

## Example 2:

**Input:** **s** = "cab", **t** = "ad"  
**Output:** false  
**Explanation:** We cannot get **t** from **s** by only one step.

## Example 3:

**Input:** **s** = "1203", **t** = "1213"  
**Output:** true  
**Explanation:** We can replace '0' with '1' to get **t**.

## Solution

### Approach 1: One pass algorithm

#### Intuition

First of all, let's ensure that the string lengths are not too far from each other. If the length difference is 2 or more characters, then **s** and **t** couldn't be one edit away strings.

**s** = "abcde**f**"  
**t** = "abcd" → couldn't be one edit away strings

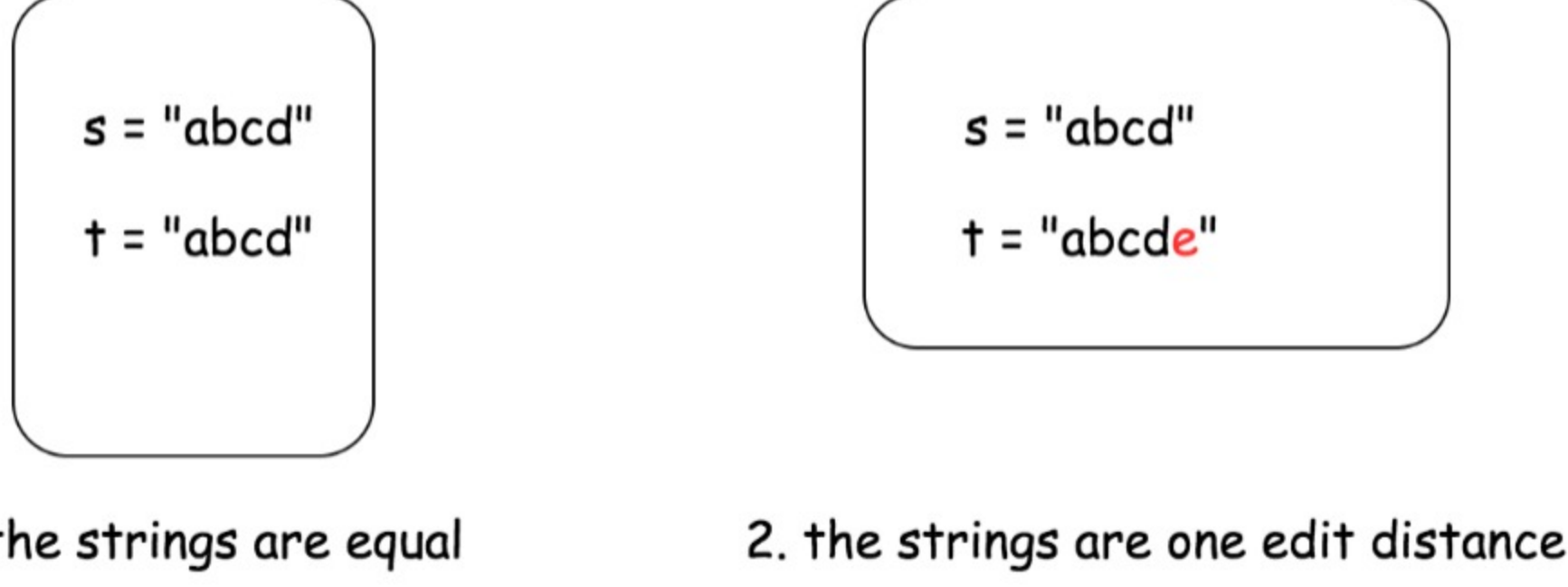
For the next let's assume that **s** is always shorter or the same length as **t**. If not, one could always call **isOneEditDistance(t, s)** to inverse the string order.

Now it's time to pass along the strings and to look for the first different character.

If there is no differences between the first **len(s)** characters, only two situations are possible :

- The strings are equal.
- The strings are one edit away distance.

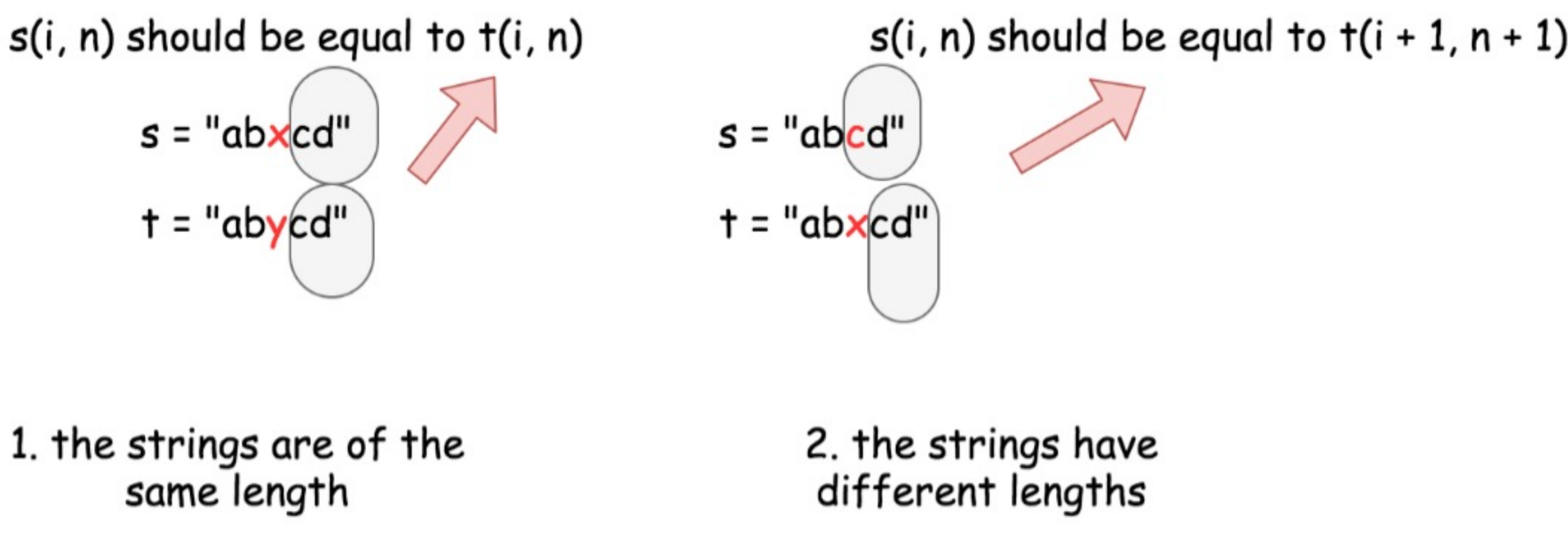
If the first len(s) characters are the same :



Now what if there is a different character so that **s[i] != t[i]**.

- If the strings are of the same length, *all* next characters should be equal to keep one edit away distance. To verify it, one has to compare the substrings of **s** and **t** both starting from the **i + 1** th character.
- If **t** is one character longer than **s**, the additional character **t[i]** should be the only difference between both strings. To verify it, one has to compare a substring of **s** starting from the **i** th character and a substring of **t** starting from the **i + 1** th character.

**s[i] != t[i] :**



#### Implementation

JavaPythonCopy

```
1 class Solution:
2     def isOneEditDistance(self, s: 'str', t: 'str') -> 'bool':
3         ns, nt = len(s), len(t)
4
5         # Ensure that s is shorter than t.
6         if ns > nt:
7             return self.isOneEditDistance(t, s)
8
9         # The strings are NOT one edit away distance
10        # if the length diff is more than 1.
11        if nt - ns > 1:
12            return False
13
14        for i in range(ns):
15            if s[i] != t[i]:
16                # if strings have the same length
17                if ns == nt:
18                    return s[i + 1:] == t[i + 1:]
19                # if strings have different lengths
20                else:
21                    return s[i:] == t[i + 1:]
22
23        # If there is no diffs on ns distance
24        # the strings are one edit away only if
25        # t has one more character.
26        return ns + 1 == nt
```

#### Complexity Analysis

- Time complexity :  $O(N)$  in the worst case when string lengths are close enough **abs(ns - nt) <= 1**, where **N** is a number of characters in the longest string.  $O(1)$  in the best case when **abs(ns - nt) > 1**.
- Space complexity :  $O(N)$  because strings are immutable in Python and Java and to create substring costs  $O(N)$  space.

#### Problem generalization : Edit distance

Given two words **word1** and **word2**, find the minimum number of operations required to convert **word1** to **word2**.

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**closewen** ★27 · June 30, 2019 11:19 PM  
This solution is not O(1) memory, each s.substring will create a new string object.  
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**milinthosani** ★31 · September 20, 2019 5:35 AM  
My 2 pointer approach:  
Time complexity: O(max(m,n))  
Space complexity: O(1)  

```
class Solution {
    public boolean isOneEditDistance(String s, String t) {
        // ...
    }
}
```

  
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**renjunyao** ★32 · August 12, 2019 12:08 PM  
We can use two pointers, instead of substrings to make it O(1) space.  
4 · · Share · Reply

**beauji\_1116** ★2 · April 14, 2020 10:19 AM  
Why output should be false when input is "c" and "c"? It didn't say u cannot use the same character to replace the old one, so why i cannot use a letter "c" to replace itself?  
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**trkx48** ★2 · March 25, 2020 7:40 AM  
Logic is like this

- Get the length of both s & t strings as l1 and l2
- If they differ more than 1, they are more than one edit distance apart, so straight return false.
- Have two pointers i and j, starting with 0th index in s & t strings respectively and loop until they

  
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**Aj007** ★6 · June 16, 2019 11:58 AM  
Correct me if I am wrong but since a substring is created in the approach 1, this algo is O(N) space too since in Java, C# strings are immutable.  
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**atulagrawal91** ★19 · February 12, 2019 5:02 PM  

```
if (Math.abs(s.length() - t.length()) > 1)
    return false;
```

```
int countChanges = 0;
```

  
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**mn002** ★1 · November 13, 2019 11:29 AM  
I would say that identical strings should return true. I hate it when these problems aren't fully specified.  
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**ssl\_91** ★28 · August 5, 2019 8:14 AM  

```
return true
```

 works fine for me. Curious to know when **(ns + 1 == nt)** will be executed?  
In any, match or not-match case, it should be already returned in the for loop, AFAIK.  
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**user1999Z** ★0 · June 22, 2020 10:47 PM  
2 pointer method: One pass (no string comparison **s == t**)  
Time: O(min(m, n))  
Space: O(1)  

```
def isOneEditDistance(s: String, t: String): Boolean = {
    // ...
}
```

  
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