Given an input string , reverse the string word by word.

### Example:

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```
Input: ["t","h","e"," ","s","k","y"," ","i","s"," ","b","l","u","e"]
Output: ["b","l","u","e"," ","i","s"," ","s","k","y"," ","t","h","e"]
```

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### Note:

- A word is defined as a sequence of non-space characters.
- The input string does not contain leading or trailing spaces.
- The words are always separated by a single space.

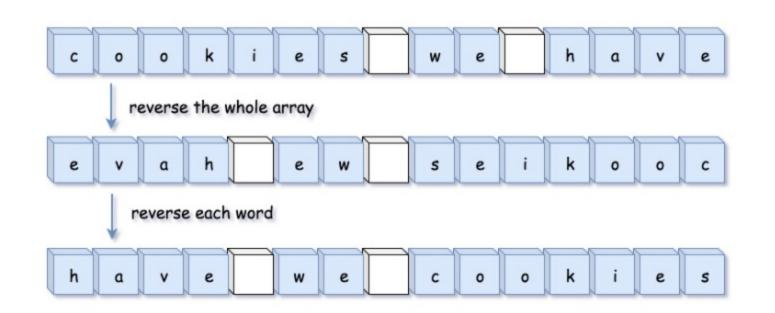
Follow up: Could you do it in-place without allocating extra space?

# Solution

## Approach 1: Reverse the Whole String and Then Reverse Each Word

To have this problem in Amazon interview is a good situation, since input is a mutable structure and hence one could aim  $\mathcal{O}(1)$  space solution without any technical difficulties.

The idea is simple: reverse the whole string and then reverse each word.



## Algorithm

Let's first implement two functions:

right pointers. C++ users could directly use built-in std::reverse.
• reverse\_each\_word(1: list), which uses two pointers to mark the boundaries of each word and

• reverse(1: list, left: int, right: int), which reverses array characters between left and

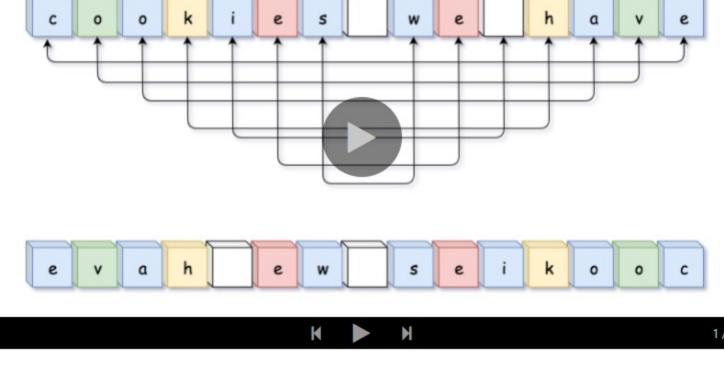
Previous function to reverse it.

Now reverseWords(s: List[str]) implementation is straightforward:

Reverse the whole string: reverse(s, 0, len(s) - 1).

- Reverse each word: reverse\_each\_word(s).
- Implementation

1. Reverse the whole array using two-pointers swap



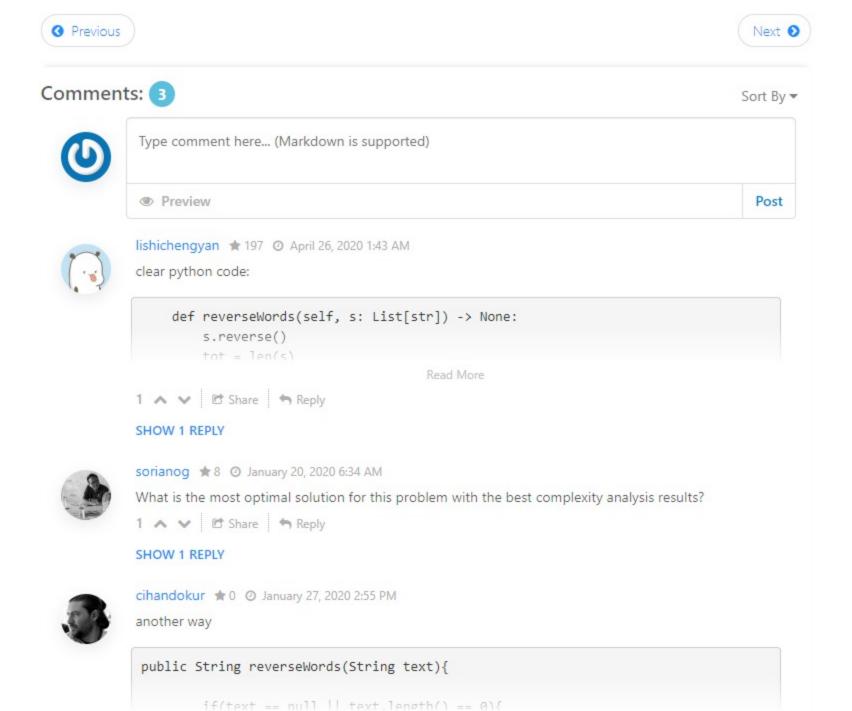
```
Сору
         Java Python
   1 class Solution {
       public:
       void reverseWords(vector<char>& s) {
         // reverse the whole string
         reverse(s.begin(), s.end());
         int n = s.size();
   8
         int idx = 0;
  9
         for (int start = 0; start < n; ++start) {
  10
          if (s[start] != ' ') {
            // go to the beginning of the word
  11
           if (idx != 0) s[idx++] = ' ';
  13
             // go to the end of the word
 14
 15
             int end = start;
             while (end < n && s[end] != ' ') s[idx++] = s[end++];
  16
  17
             // reverse the word
  18
             reverse(s.begin() + idx - (end - start), s.begin() + idx);
 19
 20
 21
             // move to the next word
 22
             start = end;
 23
 24
 25
       }
 26 };
Complexity Analysis
```

# • Time complexity: $\mathcal{O}(N)$ , it's two passes along the string.

- ullet Space complexity:  $\mathcal{O}(1)$ , it's a constant space solution.

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