

rock

★ 5276

Last Edit: August 25, 2019 6:09 PM

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Earlier sticks will be counted again. Therefore, always use current shortest two sticks till only one remains.

Java

```
public int connectSticks(int[] sticks) {
    PriorityQueue<Integer> pq = new PriorityQueue<>();
    for (int s : sticks) {
        pq.offer(s);
    }
    int sum = 0;
    while (pq.size() > 1) {
        int two = pq.poll() + pq.poll();
        sum += two;
        pq.offer(two);
    }
    return sum;
}
```

Python 3 heapq

```
class Solution:
    def connectSticks(self, sticks: List[int]) -> int:
        h = []
        for s in sticks:
            heapq.heappush(h, s)
        sum = 0
        while len(h) > 1:
            two = heapq.heappop(h) + heapq.heappop(h)
            sum += two
            heapq.heappush(h, two)
        return sum
```

Note: if the input list `sticks` is allowed to modify, we can just use `heapify(sticks)` to make the code shorter.

Analysis:

Time: $O(n\log n)$, space: $O(n)$, where $n = \text{sticks.length}$.

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what is the time complexity for this solution

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Amazing solution

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