102. Binary Tree Level Order Traversal

March 15, 2019 | 95.8K views
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Given a binary tree, return the *level order* traversal of its nodes' values. (ie, from left to right, level by level).

For example:
Given binary tree `[3,9,20,null,null,15,7]`.

```
  3
 / \
9  20
  /  \
 15   7
```

return its level order traversal as:

```
[
  [3],
  [9,20],
  [15,7]
]
```

Solution

How to traverse the tree

There are two general strategies to traverse a tree:

- *Depth First Search* (`DFS`)

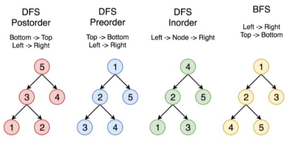
  In this strategy, we adopt the `depth` as the priority, so that one would start from a root and reach all the way down to certain leaf, and then back to root to reach another branch.

  The DFS strategy can further be distinguished as `preorder`, `inorder`, and `postorder` depending on the relative order among the root node, left node and right node.

- *Breadth First Search* (`BFS`)

  We scan through the tree level by level, following the order of height, from top to bottom. The nodes on higher level would be visited before the ones with lower levels.

On the following figure the nodes are numerated in the order you visit them, please follow `1-2-3-4-5` to compare different strategies.



Here the problem is to implement split-level BFS traversal : `[[1], [2, 3], [4, 5]]`.

Approach 1: Recursion

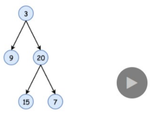
Algorithm

The simplest way to solve the problem is to use a recursion. Let's first ensure that the tree is not empty, and then call recursively the function `helper(node, level)`, which takes the current node and its level as the arguments.

This function does the following :

- The output list here is called `levels`, and hence the current level is just a length of this list `len(levels)`. Compare the number of a current level `len(levels)` with a node level `level`. If you're still on the previous level - add the new one by adding a new list into `levels`.
- Append the node value to the last list in `levels`.
- Process recursively child nodes if they are not `None` : `helper(node.left / node.right, level + 1)`.

Implementation



levels = []

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```
class Solution:
    def levelOrder(self, root):
        """
        :type root: TreeNode
        :rtype: List[List[int]]
        """
        levels = []
        if not root:
            return levels

        def helper(node, level):
            # start the current level
            if len(levels) == level:
                levels.append([])

            # append the current node value
            levels[level].append(node.val)

            # process child nodes for the next level
            if node.left:
                helper(node.left, level + 1)
            if node.right:
                helper(node.right, level + 1)

        helper(root, 0)
        return levels
```

Complexity Analysis

- Time complexity : $O(N)$ since each node is processed exactly once.
- Space complexity : $O(N)$ to keep the output structure which contains N node values.

Approach 2: Iteration

Algorithm

The recursion above could be rewritten in the iteration form.

Let's keep nodes of each tree level in the queue structure, which typically orders elements in a FIFO (first-in-first-out) manner. In Java one could use `LinkedList` implementation of the `Queue` interface. In Python using `Queue` structure would be an overkill since it's designed for a safe exchange between multiple threads and hence requires locking which leads to a performance loose. In Python the queue implementation with a fast atomic `append()` and `popleft()` is `deque`.

The zero level contains only one node `root`. The algorithm is simple :

- Initiate queue with a `root` and start from the level number `0` : `level = 0`.
- While queue is not empty :
  - Start the current level by adding an empty list into output structure `levels`.
  - Compute how many elements should be on the current level : it's a queue length.
  - Pop out all these elements from the queue and add them into the current level.
  - Push their child nodes into the queue for the next level.
  - Go to the next level `level++`.

Implementation

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```
from collections import deque
class Solution:
    def levelOrder(self, root):
        """
        :type root: TreeNode
        :rtype: List[List[int]]
        """
        levels = []
        if not root:
            return levels

        level = 0
        queue = deque([root,])
        while queue:
            # start the current level
            levels.append([])
            # number of elements in the current level
            level_length = len(queue)

            for i in range(level_length):
                node = queue.popleft()
                # fulfill the current level
                levels[level].append(node.val)

                # add child nodes of the current level
                # in the queue for the next level
                if node.left:
```

Complexity Analysis

- Time complexity : $O(N)$ since each node is processed exactly once.
- Space complexity : $O(N)$ to keep the output structure which contains N node values.

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yoursunggin ★ 34 · March 30, 2019 10:21 PM

For the solution 2, you don't need the level variable.
you can go with

```
levels.get(levels.size()-1).add(node.val);
```

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axelramar9 ★ 87 · April 16, 2019 11:24 AM

Approach 1 can be simplified as:

```
class Solution(object):
    def levelOrder(self, root):
        levels = []
```

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starsheep ★ 7 · August 3, 2019 12:28 AM

For both approaches, it was my understanding that the output structure is not taken into account when calculating space complexity in these problems. Am I wrong? If not, should the correct space complexity be O(H) where H is the height of the tree (recursion stack)?

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amster ★ 7 · March 10, 2020 6:24 AM

Isn't the recursive solution dfs and not bfs? It travels depth first preorder down the tree, not breadth first.

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meowlicious99 ★ 476 · January 8, 2020 10:37 AM

this question should be marked easy, imo.

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divyansarakb ★ 0 · October 15, 2019 9:58 AM

The iterative solution (**for Java at least**) is incorrect and I don't know what the correct solution would be

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rhenze06 ★ 0 · July 25, 2019 7:51 PM

Python Stack Implementation -- follows the same methodology as the levels in Approach 1, albeit probably higher complexity due to level_dict scan at the end (equal to number of levels)

```
class Solution:
    def levelOrder(self, root: TreeNode) -> List[List[int]]:
```

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kohok47 ★ 1 · April 9, 2019 10:02 PM

Is it possible to create a recursive function that does a breadth-first binary tree traversal in JavaScript? It seems like it should be similar to the approach above, but my attempt is not working thus far, and this post (https://stackoverflow.com/questions/33703019/breadth-first-traversal-of-a-tree-in-javascript) claims that it is not possible: "DFS is easy to implement recursively because you can use the call stack as the stack. You can't do that with BFS - because you need a queue."

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zhang-peter ★ 28 · March 28, 2019 11:40 AM

as for python solution of iteration, why I use list as a queue is much slower than using deque?

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YungHogaFire ★ 0 · 2 days ago

Why is this labeled as medium and Binary Tree Level Order Traversal II as easy?

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