\*\*\*

n x m - 1

74. Search in 2D Matrix

April 14, 2019 | 29.3K views

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properties: • Integers in each row are sorted from left to right.

Write an efficient algorithm that searches for a value in an  $m \times n$  matrix. This matrix has the following

- The first integer of each row is greater than the last integer of the previous row.
- Example 1:

Input:

```
matrix = [
   [1, 3, 5, 7],
   [10, 11, 16, 20],
   [23, 30, 34, 50]
 target = 3
 Output: true
Example 2:
```

### matrix = [

Input:

```
[1, 3, 5, 7],
  [10, 11, 16, 20],
  [23, 30, 34, 50]
target = 13
Output: false
```

## Intuition

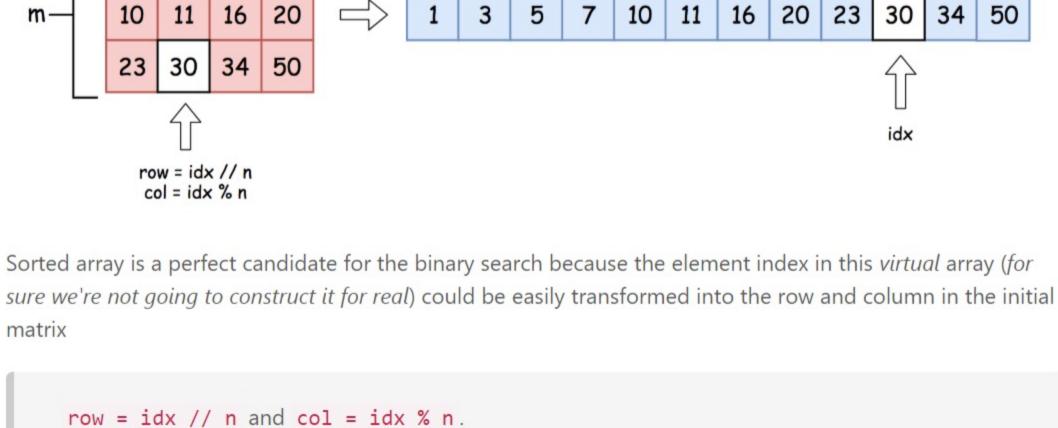
Solution

# n

Approach 1: Binary search

5 7

One could notice that the input matrix  $\mathbf{m} \times \mathbf{n}$  could be considered as a sorted array of length  $\mathbf{m} \times \mathbf{n}$ .



**Algorithm** The algorithm is a standard binary search: • Initialise left and right indexes left = 0 and right = m x n - 1.

```
• While left < right :
     • Pick up the index in the middle of the virtual array as a pivot index : pivot_idx = (left +
```

• The index corresponds to row = pivot\_idx // n and col = pivot\_idx % n in the initial

matrix, and hence one could get the pivot\_element. This element splits the virtual array in two

target = 3

target < pivot\_element --> search on the left

10

11

pivot\_idx

20

16

23

30

34

50

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parts. • Compare pivot\_element and target to identify in which part one has to look for target.

right) / 2.

n

3

11

30

pivot\_element

1

10

23

m-

15 16

17

18

19

20

**Complexity Analysis** 

O Previous

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5

16

34

else:

return False

• Space complexity :  $\mathcal{O}(1)$ .

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else:

if target < pivot\_element:</pre>

right = pivot\_idx - 1

left = pivot\_idx + 1

• Time complexity :  $\mathcal{O}(\log(mn))$  since it's a standard binary search.

7

20

50

**Implementation** 

right left



### witcher3 \* 43 • June 3, 2019 11:40 PM we can do this in O(log(m)) + O(log(n)) if we do two binary searches. One for row and another for col.

m = (h + 1) / 2Read More 25 A V C Share Reply

def binSearch(self, matrix, target, l, h):

while 1 <= h:

dfsbfs \* 12 • August 23, 2019 9:31 PM

who on earth uses '//' as division sign.

2 A V C Share Reply

O(logm + logn) solution for Java:

We start the left-bottom value V,

**SHOW 2 REPLIES** 

thanks!

Space: O(1)

s961206 ★ 734 ② July 11, 2019 7:15 AM

overflow issue if both left and right is very large.

sddlpeter ★ 18 ② January 14, 2020 5:25 AM

I have a easy-understanding solution without binary search

Type comment here... (Markdown is supported)

laplachar ★4 ② May 7, 2020 12:10 AM I don't understand the intuition of this solution. Why is that you can extract the row and column from the pivot index by doing pivotIndex/columns and pivotIndex % columns? **SHOW 1 REPLY** 

For "(left + right) // 2", using "left + (right - left) // 2" might be better since left+right might cause

- public boolean searchMatrix(int[][] matrix, int target) { if(matrix.length == 0 || matrix[0].length == 0) return false; Read More 2 A V C Share Share
- if (target == V) return true; else if (target < v) we move upward; else if (target > v) we move rightward (we can early stop in this case: once we meet the value > target, Read More
- 1 A V C Share Reply haoyangfan 🖈 897 🧿 July 27, 2019 10:32 AM two binary search, 0ms, java Time: O(log(m) + log(n)
- class Solution Read More 1 A V C Share Reply SHOW 1 REPLY Reputation

I implemented 2D binary search and felt so stupid after I looked at the solution.

A Report sainvamshid **★**6 **②** May 17, 2020 11:31 AM bool searchMatrix(vector<vector<int>>& matrix, int target) { if (!matrix.size()) return false; int i = 0. i = matrix[0].size()-1:

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suthee \*1 ② July 2, 2020 3:43 AM