## 346. Moving Average from Data Stream 🗗 Oct. 6, 2019 | 19.5K views



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window. Example:

Given a stream of integers and a window size, calculate the moving average of all integers in the sliding

```
MovingAverage m = new MovingAverage(3);
m.next(1) = 1
m.next(10) = (1 + 10) / 2
m.next(3) = (1 + 10 + 3) / 3
m.next(5) = (10 + 3 + 5) / 3
```

# Intuition

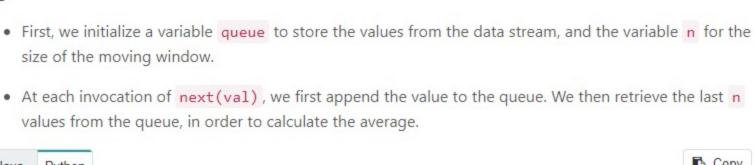
Solution

### structure of Array or List. Then from the data structure, later we retrieve the necessary elements to calculate the average.

Approach 1: Array or List

window 10 5 3 8 2 queue 11

Following the description of the problem, we could simply keep track of all the incoming values with the data



### Python Java 1 class MovingAverage:

def next(self, val: int) -> float:

size, queue = self.size, self.queue

Algorithm

- def \_\_init\_\_(self, size: int): self.size = size self.queue = []
- queue.append(val) # calculate the sum of the moving window 10 window\_sum = sum(queue[-size:])

```
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             return window_sum / min(len(queue), size)
Complexity
   ullet Time Complexity: \mathcal{O}(N) where N is the size of the moving window, since we need to retrieve N
     elements from the queue at each invocation of next(val) function.
   • Space Complexity: \mathcal{O}(M), where M is the length of the queue which would grow at each invocation
     of the next(val) function.
Approach 2: Double-ended Queue
```

First of all, one might notice that we do not need to keep all values from the data stream, but rather

By definition of the moving window, at each step, we add a new element to the window, and at the same

complexity  $(\mathcal{O}(1))$  to add or remove an element from both its ends. With the deque, we could reduce the

deque

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Secondly, to calculate the sum, we do not need to reiterate the elements in the moving window.

### time we remove the oldest element from the window. Here, we could apply a data structure called doubleended queue (a.k.a deque) to implement the moving window, which would have the constant time

Algorithm

Java Python

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Complexity

1 from collections import deque

Intuition

space complexity down to  $\mathcal{O}(N)$  where N is the size of the moving window.

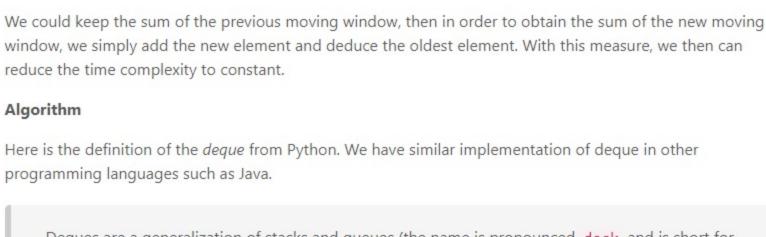
We could do better than the first approach in both time and space complexity.

the last **n** values which falls into the moving window.

pop left append right

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2 class MovingAverage: def \_\_init\_\_(self, size: int): self.size = size self.queue = deque() # number of elements seen so far self.window\_sum = 0

```
• Space Complexity: \mathcal{O}(N), where N is the size of the moving window.
Approach 3: Circular Queue with Array
Intuition
Other than the deque data structure, one could also apply another fun data structure called circular
queue, which is basically a queue with the circular shape.
                                            circular queue
                                              head tail
```

No need to resort to any library, one could easily implement a circular queue with a fixed-size array. The key to the implementation is the correlation between the index of head and tail elements, which we could summarize in the following formula:

element.

Python

class MovingAverage:

def \_\_init\_\_(self, size: int): self.size = size

self.count = 0

self.count += 1

Preview

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acoolguy \* 33 O November 16, 2019 6:28 AM

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fliess \*3 O October 9, 2019 12:37 PM

Simple Java Solution with a fixed sized array.

class MovingAverage {

0 ∧ ∨ ♂ Share ← Reply

int[] arr:

self.queue = [0] \* self.size self.head = self.window\_sum = 0 # number of elements seen so far

def next(self, val: int) -> float:

# move on to the next head

self.queue[self.head] = val

tail = (self.head + 1) % self.size

self.head = (self.head + 1) % self.size

# calculate the new sum by shifting the window

self.window\_sum = self.window\_sum - self.queue[tail] + val

return self.window\_sum / min(self.size, self.count)

Java

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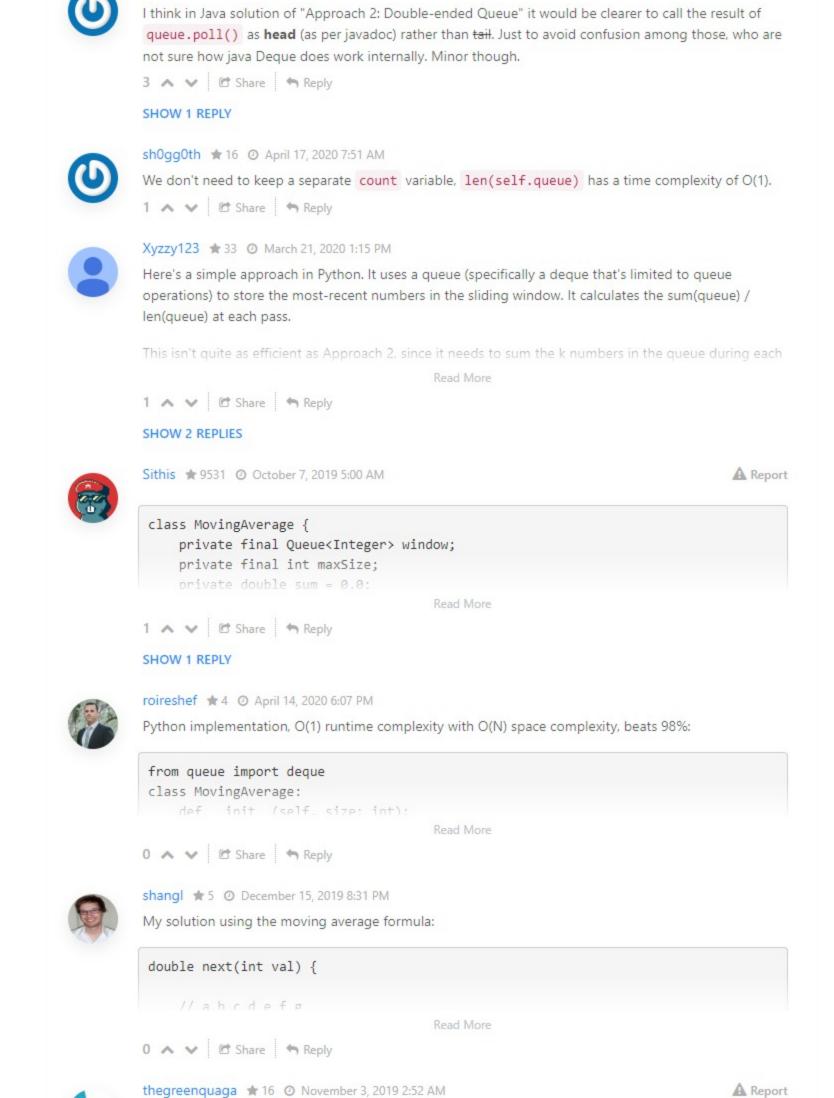
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Algorithm



I would mention a queue instead of a Deque here in the explanation. We aren't really using both the stack and queue aspects of this data structure. Instead we are simply using the deque as a queue.



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Deques are a generalization of stacks and queues (the name is pronounced deck and is short for double-ended queue ). Deques support thread-safe, memory efficient appends and pops from

either side of the deque with approximately the same O(1) performance in either direction.

calculate the sum of moving window in constant time.

self.window\_sum = self.window\_sum - tail + val

• Time Complexity:  $\mathcal{O}(1)$ , as we explained in intuition.

return self.window\_sum / min(self.size, self.count)

Follow the intuition, we replace the queue with the deque and add a new variable window\_sum in order to

8 self.count = 0 9 10 def next(self, val: int) -> float: 11 self.count += 1 12 # calculate the new sum by shifting the window self.queue.append(val) 13 14 tail = self.queue.popleft() if self.count > self.size else 0 15

tail = (head + 1) mod size

• The major advantage of circular queue is that by adding a new element to a full circular queue, it

Another advantage of circular queue is that a single index suffices to keep track of both ends of the

queue, unlike deque where we have to keep a pointer for each end.

automatically discards the oldest element. Unlike deque, we do not need to explicitly remove the oldest

In other words, the tail element is right next to the head element. Once we move the head forward, we would overwrite the previous tail element. Eating its own tail

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 $tail = (head + 1) \mod size$