**f y in** 

Sept. 4, 2018 | 111.3K views

153. Find Minimum in Rotated Sorted Array 🗗 \*\*\* Average Rating: 4.40 (88 votes)

(i.e., [0,1,2,4,5,6,7] might become [4,5,6,7,0,1,2]).

Suppose an array sorted in ascending order is rotated at some pivot unknown to you beforehand.

Find the minimum element.

You may assume no duplicate exists in the array.

Example 1:

Output: 1

```
Input: [4,5,6,7,0,1,2]
Output: 0
```

## Intuition A very brute way of solving this question is to search the entire array and find the minimum element. The

A very cool way of solving this problem is using the **Binary Search** algorithm. In binary search we find out the mid point and decide to either search on the left or right depending on some condition.

## Since the given array is sorted, we can make use of binary search. However, the array is rotated. So simply

Algorithm

the left of mid.

Python

class Solution(object):

:rtype: int

def findMin(self, nums):

:type nums: List[int]

if len(nums) == 1:

# left pointer

left = 0

return nums[0]

if nums[right] > nums[0]:

# Find the mid element

mid = left + (right - left) / 2

return nums[0]

# Binary search way

while right >= left:

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**SHOW 6 REPLIES** 

Too many comparisons.

sufficient for binary search.

class Solution (

49 A V C Share Share

the-minimum-in-a-sorted-rotated-array/

jainrishabh ★ 73 ② September 9, 2018 7:50 PM

public static void main(String[] args) {

int arr[] = { 4, 4, 4, 4 };

4 A V C Share Reply

2 A V C Share Reply

class Solution {

chef29 ★ 2 ② September 5, 2018 10:08 AM

**SHOW 1 REPLY** 

A little simpler version:

DenisGubar ★ 62 ② September 6, 2018 11:50 AM

Java

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**Complexity Analysis** 

applying the binary search won't work here. In this question we would essentially apply a modified version of binary search where the **condition** that

time complexity for that would be O(N) given that  ${\tt N}$  is the size of the array.

decides the search direction would be different than in a standard binary search.

We want to find the smallest element in a rotated sorted array. What if the array is not rotated? How do we check that? If the array is not rotated and the array is in ascending order, then <code>last element > first element</code>.

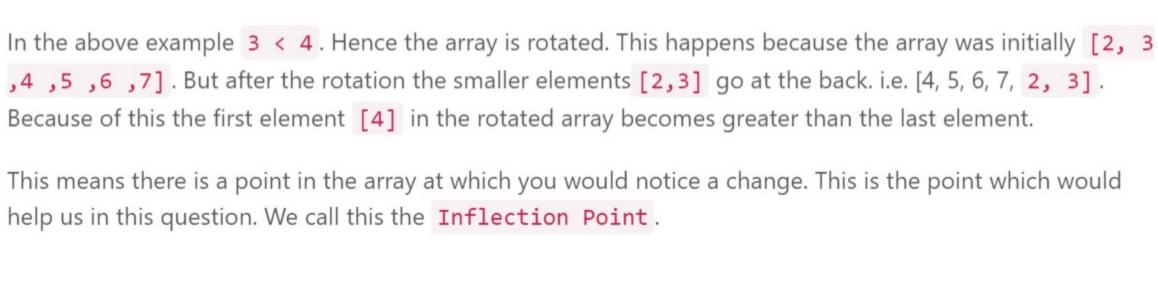
7 > 2

In the above example 7 > 2. This means that the array is still sorted and has no rotation.

7

2

3



**Rotated Array** 

In this modified version of binary search algorithm, we are looking for this point. In the above example notice

6 > 4

6

Mid

4. We stop our search when we find the inflection point, when either of the two conditions is satisfied:

2

3

Right

5

All the elements to the left of inflection point > first element of the array. All the elements to the right of inflection point < first element of the array. 1. Find the **mid** element of the array. 2. If mid element > first element of array this means that we need to look for the inflection point on the right of mid. 3. If mid element < first element of array this that we need to look for the inflection point on

In the above example mid element 6 is greater than first element 4. Hence we continue our search for the inflection point to the right of mid.

nums[mid] > nums[mid + 1] Hence, mid+1 is the smallest.

nums[mid - 1] > nums[mid] Hence, mid is the smallest.

4

Left

before 2 is 7 and 7>2 i.e. nums[mid - 1] > nums[mid]. Thus we have found the point of inflection and 2 is the smallest element. 🖺 Сору

# If the list has just one element then return that element.

# Hence the smallest element is first element. A[0]

3 5 6 4 Right Left Mid In the above example. With the marked left and right pointers. The mid element is 2. The element just

# right pointer right = len(nums) - 1# if the last element is greater than the first element then there is no rotation. # e.g. 1 < 2 < 3 < 4 < 5 < 7. Already sorted array.

# This point would be the point of change. From higher to lower value.

# if the mid element is greater than its next element then mid+1 element is the smallest

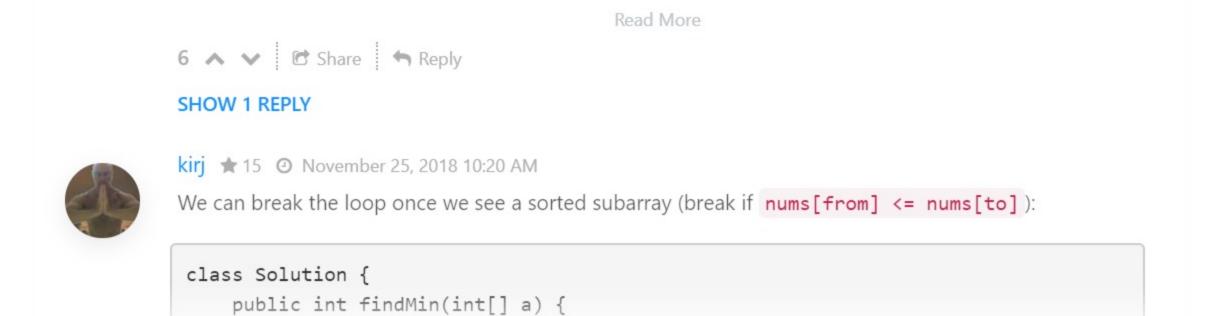
```
ullet Time Complexity : Same as Binary Search O(\log N)
  • Space Complexity : O(1)
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Comments: 55
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            Type comment here... (Markdown is supported)
             Preview
                                                                                              Post
            tomba ★ 141 ② February 17, 2019 2:05 PM
                 int findMin(vector<int>& nums) {
                     int lo = 0, hi = nums.size()-1;
                    while (lo < hi) {
                         int mid = 10 + (hi-10)/2:
                                                  Read More
            50 A V C Share   Reply
```



Great explanation! I had one here but it wasn't as concise: http://code.scottshipp.com/2018/06/27/find-

The problem's invariant for shifted cases is left element always greater than right. Maintaining it is

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Read More

Read More

2 A V C Share Reply **SHOW 6 REPLIES** kindernerd ★8 ② September 9, 2018 7:32 PM

the solution code seems to have a bug, returns -1 for [2,2,2,2,3,3,1,2]

public int findMin(int[] nums) { if(nums[0]<nums[nums.length-1]){</pre> return nums[0]: Read More 1 A V C Share Reply **SHOW 1 REPLY** 

Easiest way to solve this problem: Apply a filter of < last element to each element, and we get a boolean array [3,4,5,1,2] = filter( < last element) = > [F, F, F, F, T, T]

> SHOW 1 REPLY ( 1 2 3 4 5 6 )

0 ∧ ∨ ♂ Share ★ Reply

lukskywalkerrr ★ 10 ② June 11, 2020 1:44 PM

Input: [3,4,5,1,2] Example 2:

Solution Approach 1: Binary Search

> 2 3 5 6 4 No Rotation 3 > 4

> > 6

5

