162. Find Peak Element 2 May 29, 2017 | 153.5K views

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Given an input array nums, where $nums[i] \neq nums[i+1]$, find a peak element and return its index.

A peak element is an element that is greater than its neighbors.

The array may contain multiple peaks, in that case return the index to any one of the peaks is fine.

You may imagine that $nums[-1] = nums[n] = -\infty$. Example 1:

Input: nums = [1,2,3,1]

Output: 2 Explanation: 3 is a peak element and your function should return the index number 2.

Example 2: Input: nums = [1,2,1,3,5,6,4]Output: 1 or 5

Explanation: Your function can return either index number 1 where the peak element is or index number 5 where the peak element is 6.

Solution

Follow up: Your solution should be in logarithmic complexity.

nums[i] is the peak element. The reasoning behind this can be understood by taking the following three cases which cover every case into which any problem can be divided.

element. We start off by checking if the current element is larger than the next one. The first element satisfies this criteria, and is hence identified as the peak correctly. In this case, we didn't reach a point where we

0

5

3 3

Approach 1: Linear Scan

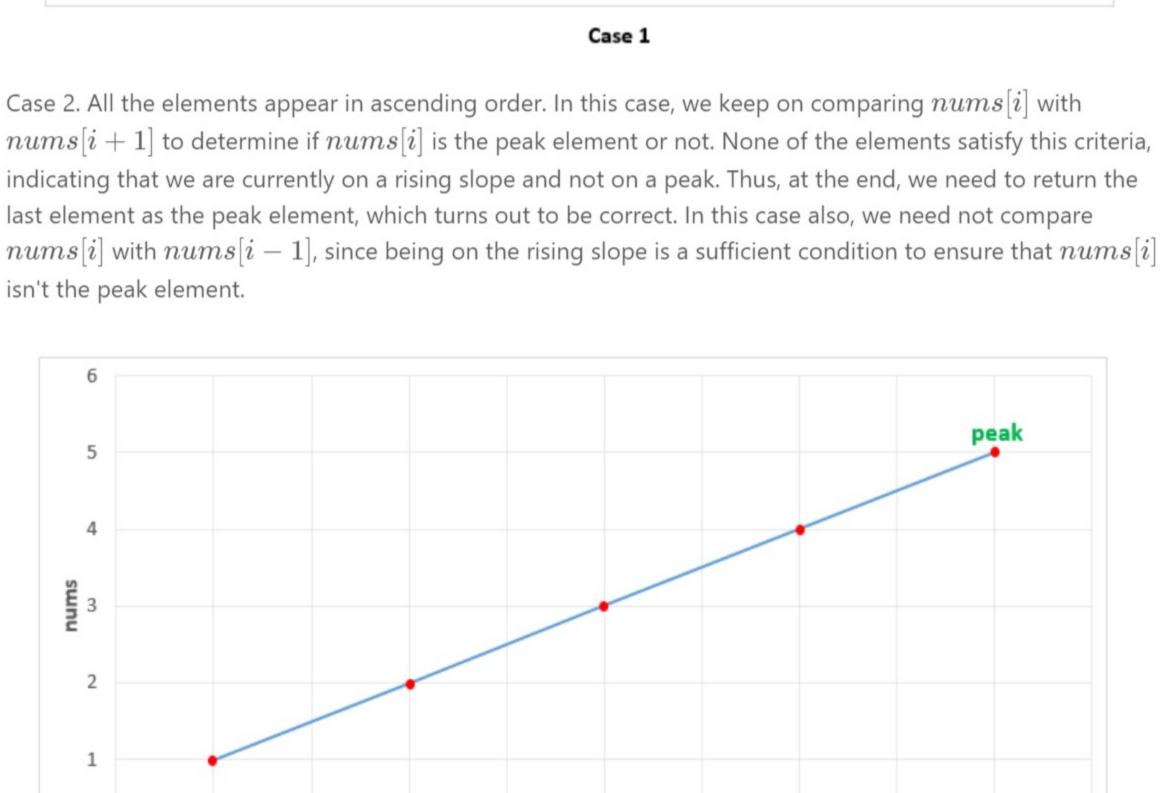
peak nums 3

In this approach, we make use of the fact that two consecutive numbers nums[j] and nums[j+1] are

never equal. Thus, we can traverse over the nums array starting from the beginning. Whenever, we find a

number nums[i], we only need to check if it is larger than the next number nums[i+1] for determining if

Case 1. All the numbers appear in a descending order. In this case, the first element corresponds to the peak



index

Case 2

Case 3. The peak appears somewhere in the middle. In this case, when we are traversing on the rising edge,

as in Case 2, none of the elements will satisfy nums[i] > nums[i+1] . We need not compare nums[i]

condition nums[i] > nums[i+1] is satisfied. We again, need not compare nums[i] with nums[i-1].

nums[i+1] failed for the previous($(i-1)^{th}$ element, indicating that nums[i-1] < nums[i]. Thus, we

with nums[i-1] on the rising slope as discussed above. When we finally reach the peak element, the

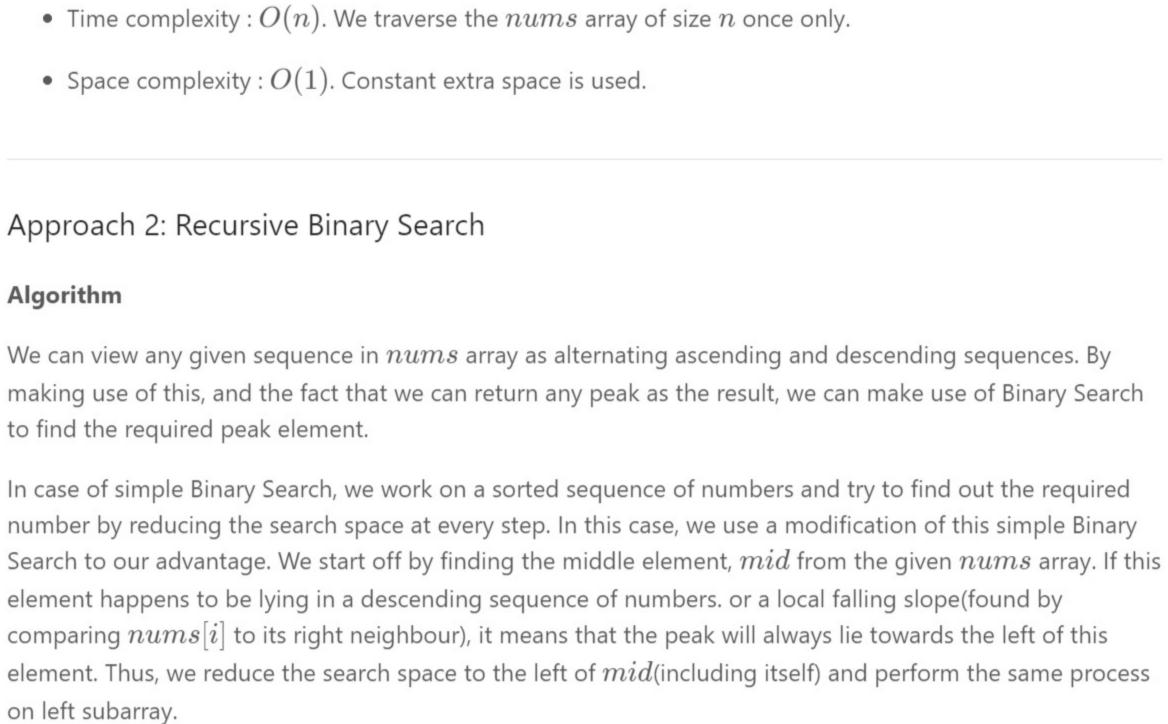
This is because, we could reach nums[i] as the current element only when the check nums[i] >

are able to identify the peak element correctly in this case as well.

3

peak

0 0 3 index Case 3 **С**ору Java public class Solution {



If the middle element, mid lies in an ascending sequence of numbers, or a rising slope(found by comparing

Thus, we reduce the search space to the right of mid and perform the same process on the right subarray.

In this way, we keep on reducing the search space till we eventually reach a state where only one element is

Case 1. In this case, we firstly find 3 as the middle element. Since it lies on a falling slope, we reduce the

slope, reducing the search space to [1] only. Thus, 1 is returned as the peak correctly.

search space to [1, 2, 3]. For this subarray, 2 happens to be the middle element, which again lies on a

falling slope, reducing the search space to [1, 2]. Now, 1 acts as the middle element and it lies on a falling

nums[i] to its right neighbour), it obviously implies that the peak lies towards the right of this element.

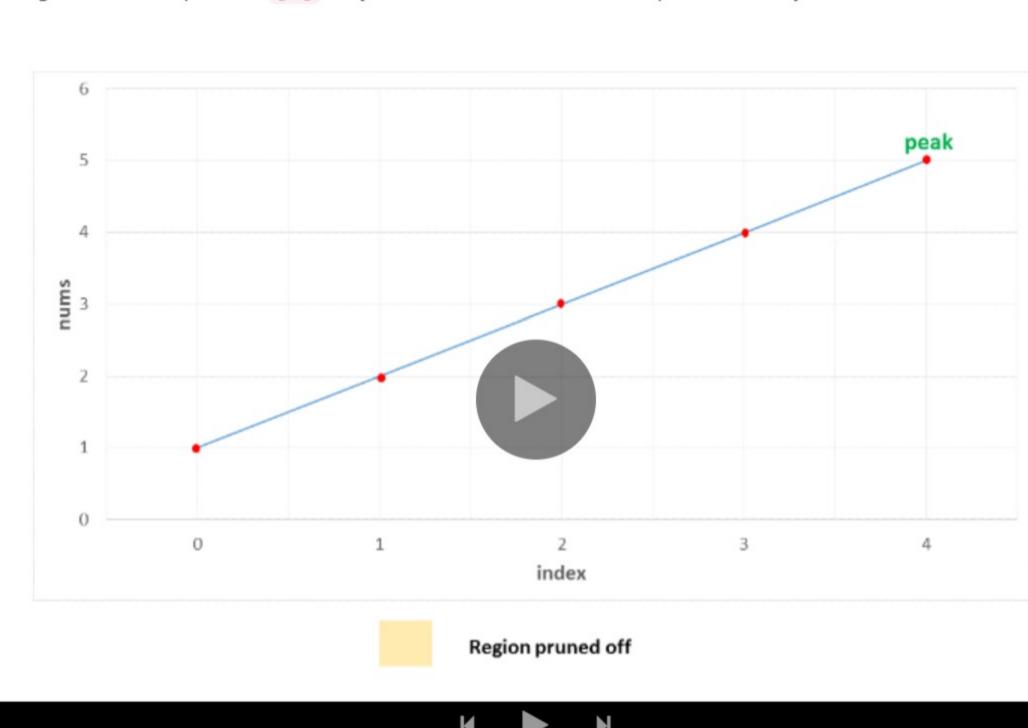
peak

nums

remaining in the search space. This single element is the peak element.

To see how it works, let's consider the three cases discussed above again.

3 index Region pruned off



Case 3. In this case, the peak lies somewhere in the middle. The first middle element is 4. It lies on a rising

slope, indicating that the peak lies towards its right. Thus, the search space is reduced to [5, 1]. Now, [5, 1]

happens to be the on a falling slope(relative to its right neighbour), reducing the search space to [5] only.

public int findPeakElement(int[] nums) { return search(nums, 0, nums.length - 1); public int search(int[] nums, int l, int r) { if (1 == r)return 1; int mid = (1 + r) / 2;if (nums[mid] > nums[mid + 1]) 10 return search(nums, 1, mid);

ullet Time complexity : $Oig(log_2(n)ig)$. We reduce the search space in half at every step. Thus, the total search

search space will be consumed in $log_2(n)$ steps. Thus, the depth of recursion tree will go upto $log_2(n)$

• Space complexity : $O(log_2(n))$. We reduce the search space in half at every step. Thus, the total

space will be consumed in $log_2(n)$ steps. Here, n refers to the size of nums array.

```
public int findPeakElement(int[] nums) {
             int l = 0, r = nums.length - 1;
             while (l < r) {
                 int mid = (1 + r) / 2;
                 if (nums[mid] > nums[mid + 1])
                     r = mid;
                 else
                    l = mid + 1;
             return 1;
 13
Complexity Analysis
   • Time complexity : O(log_2(n)). We reduce the search space in half at every step. Thus, the total search
     space will be consumed in log_2(n) steps. Here, n refers to the size of nums array.
   • Space complexity : O(1). Constant extra space is used.
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```

SHOW 2 REPLIES xrssa * 1925 • November 18, 2019 10:28 AM Have a look at MIT 6006's 1st lecture on 'Algorithmic thinking and Peak Finding' (timestamp for divide

https://www.youtube.com/watch?v=HtSuA80QTyo&list=PLUI4u3cNGP61Oq3tWYp6V F-**SHOW 2 REPLIES**

miss the 2D version of peak finding in the lecture.

akhil1311 ★ 177 ② July 25, 2019 6:10 PM

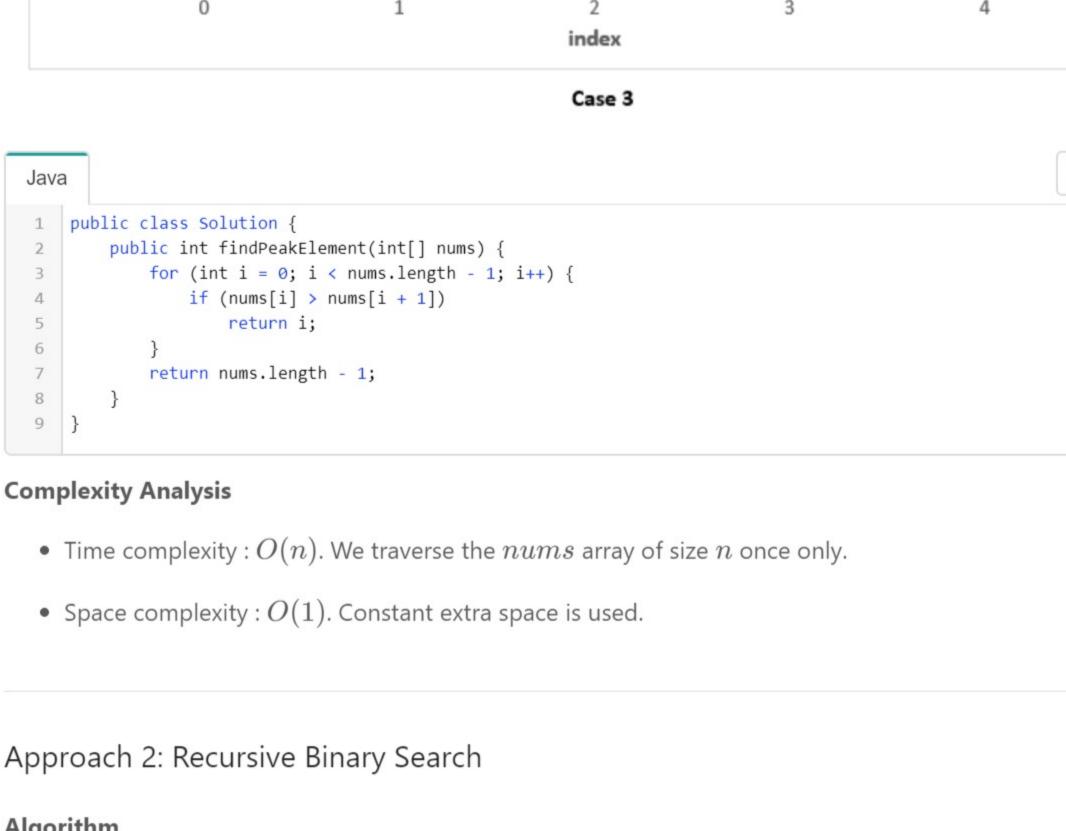
8 A V C Share Reply rohitniyer **†** 6 **②** March 24, 2018 4:42 PM

not greater than 6) then it goes to the else condition and low becomes mid +1 that is 5. and our search range reduces by half. we could have have a possible solution at index 1 where the element 4 is greater than its neighbours. Read More 6 ∧ ∨ ♂ Share ★ Reply **SHOW 5 REPLIES**

quanhoang91 ★ 9 ② June 24, 2020 3:46 PM Solution is wrong. Binary search solution always assume its either on left or right. But thats not the case. E.g. input is [2,3,2,2,2,2,2,2], expected output is 7 where it should be 1 **SHOW 3 REPLIES**

needed to compare nums[i] with nums[i-1] also, to determine if it is the peak element or not.

index



Thus, 5 is identified as the peak element correctly.

Java

11 12

13

Java

10

11 12

Complexity Analysis

Approach 3: Iterative Binary Search

- peak nums 3 index Region pruned off public class Solution {
- Algorithm The binary search discussed in the previous approach used a recursive method. We can do the same process in an iterative fashion also. This is done in the current approach. public class Solution {

Nice solution. How do I get such solutions by myself in 30 minutes? 22 \Lambda 🗸 🗗 Share 🦘 Reply

SHOW 3 REPLIES gary17 ★8 ② September 15, 2018 8:03 AM The explanation is not intuitive for me, it can be super simple: Instead of see looking for peak, try another word turn: it keeps going up, then it has to have a peak turn or reach to the end which by the definition is also what we are looking for. The word peak is still ok, I was given "local maximum/local minimum", totally confused/mislead me.

jasonz1121 ★ 11 ② September 4, 2018 6:14 AM Looks like sol2 and 3 are not right. Test with case [1,2,1,1,1,1] 10 A V C Share Reply **SHOW 2 REPLIES**

(123456 ... 89)

Joseph_Zhao ★ 3 **②** October 11, 2018 8:32 PM

Case 2. In this case, we firstly find 3 as the middle element. Since it lies on a rising slope, we reduce the search space to [4, 5]. Now, 4 acts as the middle element for this subarray and it lies on a rising slope, reducing the search space to [5] only. Thus, 5 is returned as the peak correctly.

return search(nums, mid + 1, r);

1/6

Сору

Сору

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Comments: 83 Type comment here... (Markdown is supported) Preview delta45724 ★ 165 ② January 14, 2019 10:45 AM better provide a strict proof why Binary Search works here 113 🔨 🖝 Share 🦘 Reply SHOW 11 REPLIES ygoOP ★ 187 ② July 9, 2018 6:47 PM This line can integer overflow "int mid = (l + r) / 2; 35 ∧ ∨ ☑ Share ¬ Reply SHOW 5 REPLIES

and conquer strategy for peak finding: 27.40 min). You will understand the correctness of the

algorithm(binary search in peak finding). Lecturer also showed the master theorem for this also. Never

ddepaolo 🖈 23 🗿 November 10, 2018 5:16 AM The solution assumes the test cases are sorted already. However they are not. So I think it's just wrong. 23 🔨 🔀 Share 👆 Reply

i guess binary search will not work for the case when we take an array nums[1,4,2,3,5,6,7,8,9,10] because if we calculate mid =(0+9)/2=4 and nums[4] is 5 and nums[4] is not greater than nums[5](that is 5 is

Tested with case [1,200,3,4,5,6,7,8,9,10]. However, both sol2 and 3 returned 9... 3 A V C Share Reply **SHOW 5 REPLIES**