275. H-Index II July 18, 2019 | 14.1K views

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According to the definition of h-index on Wikipedia: "A scientist has index h if h of his/her N papers have at **least** h citations each, and the other N-h papers have **no more than** h citations each."

Example:

# Input: citations = [0,1,3,5,6]Output: 3

```
papers in total and each of them had
                received 0, 1, 3, 5, 6 citations respectively.
               Since the researcher has 3 papers with at least 3
   citations each and the remaining
               two with no more than 3 citations each, her h-index is 3.
Note:
If there are several possible values for h, the maximum one is taken as the h-index.
```

Could you solve it in logarithmic time complexity?

Explanation: [0,1,3,5,6] means the researcher has 5

- Solution

# that the number of articles whose citation number is higher than c would be n - i - 1. And together with the current article, there are n - i articles that are cited at least c times.

Approach 1: Linear search, O(k) time

citation[i] is greater or equal to n - i, i.e. c >= n - i. As we know that all the articles following i would be cited at least c times, so in total there are n - i articles that are cited at least c times. As a result, according to the definition, the H-Index of the author should be n - i.

number of citations

H-index = 3

Copy

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1 article idx 0 2 3 5 0 Following the above intuition, it is straightforward to implement the algorithm. We give some examples in class Solution: def hIndex(self, citations):

**Complexity Analysis** 

ullet Time complexity :  $\mathcal{O}(N)$  where N is the length of the input list, since in the worse case we would

# Approach 2: Binary Search, O(log N) time

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Algorithm

Java

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16 17 18 class Solution:

def hIndex(self, citations):

n = len(citations)

left, right = 0, n - 1

pivot = left + (right - left) // 2

elif citations[pivot] < n - pivot:</pre>

• Space complexity :  $\mathcal{O}(1)$ , it's a constant space solution.

ullet Time complexity :  $\mathcal{O}(\log N)$  since we apply binary search algorithm here.

if citations[pivot] == n - pivot:

return n - pivot

left = pivot + 1

right = pivot - 1

while left <= right:

:rtype: int

else:

return n - left

Rate this article: \* \* \* \* \*

:type citations: List[int]

Intuition

following task:

Java

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Given a sorted list citations of size n in ascending order, one is asked to find the first number citations[i] which meets the constraint of citations[i] >= n - i.

Following in the intuition we elaborated in the Approach 1, the problem is actually boiled down to the

half, which leads to a more optimal  $\mathcal{O}(\log N)$  time complexity comparing to the  $\mathcal{O}(N)$  of the linear search.

algorithm to solve the problem. In binary search algorithm, we recursively reduce the searching scope into

With the above formulation of the problem, it becomes clear that one could apply the binary search

1 article idx 0 2 0 3 5

First we pick a pivot element that is in the middle of the list, i.e. citations[pivot], which would divide the

H-index = 3

# original list into two sublists: citations[0 : pivot - 1] and citations[pivot + 1 : n]. Then comparing between the values of n - pivot and citations[pivot] element, our binary search algorithm breaks down to the following 3 cases: • citations[pivot] == n - pivot: We found the desired element! • citations[pivot] < n - pivot : Since the desired element should be greater or equal to n pivot , we then further look into the sublist on the right hand side, i.e. citations[pivot + 1 : n] . • citations[pivot] > n - pivot : We should look into the sublist on the left hand side, i.e. citations[0 : pivot-1]. A minor difference to the typical binary search algorithm, is that in this case we return the value of npivot as the result, rather than the value of the desired element. Implementation Python

O Previous Next **1** 

**Complexity Analysis** 

Comments: 14 Sort By ▼ Type comment here... (Markdown is supported) Preview Post diiikoo ★ 19 ② August 11, 2019 7:41 AM For the Binary search solution, I think the variable "idx" is of no use. sunxiaohua 🛊 3 🗿 October 21, 2019 7:49 AM Why there has a condition: If there are several possible values for h, the maximum one is taken as the h-index? Don't think multiple index could satisfied this. 3 A V C Share Reply **SHOW 2 REPLIES** I think as per the question condition should be c[i] == n - i the problem is poorly phrased. 2 A V C Share Share

In binary search I wonder how to decide between while(lo < hi)</pre> and while(lo <= hi)</pre> . I seem

Read More

"The h-index is defined as the maximum value of h such that the given author/journal has published h

lenchen1112 \* 1005 • December 3, 2019 4:52 PM For binary search approach, use pattern two will be easier to understand. class Solution: def hIndex(self, citations: List[int]) -> int:

to always get those wrong, any suggestions?

n = len(citations)

question description is either confusing or wrong.

2 A V C Share Reply

3 A V C Share Reply

damhajan92 ★ 0 ② 3 days ago

xI549 ★ 0 ② June 6, 2020 9:18 AM

**SHOW 1 REPLY** 

**SHOW 2 REPLIES** 

SHOW 1 REPLY

jaykpatel1996 ★1 ② July 10, 2020 10:41 AM I think this definition of H index from Wikipedia makes this question a lot easier.

papers that have each been cited at least h times." jaykpatel1996 ★ 1 ② July 5, 2020 11:01 AM This question should be marked as Hard.

zzznotsomuch \* 54 • June 19, 2020 9:10 AM In Java Solution, can someone explain why are we returning n-left from outside the loop? 

for the approach 2, i was wondering why we "return len - start" rather than "return len - end"

A Report

Intuition Thanks to the fact that the list of citation numbers is sorted in the ascending order, one could solve the problem in a single pass of iteration. Let's consider an article whose citation number c is index at i, i.e c = citations[i]. We would know

# Given the definition of H-Index, we just need to find the first article at i whose citation number c =

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3 the following. Implementation Python

:type citations: List[int]

if c >= n - idx:

for idx, c in enumerate(citations):

• Space complexity :  $\mathcal{O}(1)$ , it's a constant space solution.

pivot

return n - idx

:rtype: int

return 0

iterate the entire list.

n = len(citations)

researcher, write a function to compute the researcher's h-index.

ascending order.

Follow up: • This is a follow up problem to H-Index, where citations is now guaranteed to be sorted in

Given an array of citations sorted in ascending order (each citation is a non-negative integer) of a

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