Today's agenda

- understanding Sorting
- Problems on Sorting
- Sorting techniques

**Sorting:** Arranging data in increasing / decreasing.
↓
on what Parameter.

Ex1: 2 4 10 15 27 → true

Ex2: 20 7 3 -5 -8 → dec. order → true

Ex3: 1 2 3 7 4 9 6 → Not sorted on the basis of value.
→ sorted on the basis of factor count.

#factor: 1 2 2 2 3 3 4

- bubble sort
- selection sort
- insertion sort
- merge sort
- quick sort
- bucket sort
- radix sort

etc.

Steps to Solve Problem:

1 2 3 4

Sort the array
↓
inbuilt sorting function.
↓↓
Arrays.sort(arr); → inc. order
↓
T.C: O(NlogN) S.C: O(N)

Worst T.C: $O(N^2)$

↳ How in levelup.

## Q) Order of Removal

↳ Given N elements at every step remove an array element. Cost to remove element = Sum of array elements present. Find min cost to remove all elements.

Note: Add cost first and then remove.

Ex1: arr[3] = {3 2 5}

remove 3: 10
remove 2: 7
remove 5: 5
22

arr[3] = {3 2 5}

remove 2: 10
remove 5: 8
remove 3: 3
21

Ex2: arr[4] = {4 6 2 7}

| | |
|---|---|
| remove 7: | 4 + 6 + 2 + 7 |
| remove 2: | 4 + 6 + 2 |
| remove 6: | 4 + 6 |
| remove 4: | 4 |

max Contribution: 0th index → min
2nd max Contri: 1st index → 2nd min
⋮

Array should be sorted in inc. order & remove from last.

arr[4] = { 4 6 2 7 }



{ 2 4 6 7 }

remove 7 : 2 + 4 + 6 + 7

remove 6 : 2 + 4 + 6 (6*2)

remove 4 : 2 + 4 (4*3)

remove 2 : 2

2+4

adding frequency = N - i

//Pseudo code

```
int orderofremoval(int arr[N]){
    Arrays.sort(arr);
    int ans = 0;
    for(int i=0; i<N; i++){
        int temp = arr[i] * (N-i);
        ans = ans + temp;
    }
    return ans;
}
```

T.C : O(NlogN) + O(N)

= O(NlogN + N)

= O(NlogN)

S.C : O(N)

Tracing

```
int orderofremoval(int arr[N]){
    Arrays.sort(arr);
    int ans = 0;
    for(int i=0; i<N; i++){
        int temp = arr[i] * (N-i);
        ans = ans + temp;
    }
    return ans;
}
```

arr[4] = { 4 6 2 7 }

(indices: 0 1 2 3)

↓

| 0 | 1 | 2 | 3 |
|---|---|---|---|
| 2 | 4 | 6 | 7 |

ans = 0 + 8 + 12 + 12 + 7 = 39

| i | temp |
|---|---|
| 0 | 2*4 |
| 1 | 4*3 |
| 2 | 6*2 |
| 3 | 7*1 |

Q) Good Integer { only distinct }

↳ Given arr[n], Calculate no. of good integers.

An element is said to be good if

{ No. of element < ele :: ele itself }

i :: arr[i]

Ex1 : { -1 -4 3 5 -15 4 } ⇒ ans: 3

| | 0 | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|---|
| | -1 | -4 | 3 | 5 | -15 | 4 |
| # less : | 2 | 1 | 3 | 5 | 0 | 4 |

Ex1 : { -1 -4 3 5 -15 4 }

↓ sort in inc. order

| | 0 | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|---|
| | -15 | -4 | -1 | 3 | 4 | 5 |

Count of elements smaller arr[i] = i

element itself = arr[i]

//Pseudo code

```
int goodintegers (int arr[N]){
    Arrays.sort(arr);
    int count=0;
    for (int i=0; i<N; i++){
        if (arr[i]==i) {count++;}
    }
    return count;
}
```

T.C: O(NlogN + N)
= O(NlogN)

S.C: O(N)

{No. of element < ele :: ele itself}
↓ i = arr[i]

Good Integers : { Data can repeat } → last idea won't work directly.

Ex1:

| | 0 | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|---|
| | 0 | 2 | 2 | 3 | 3 | 8 |
| #less: | 0 | 1 | 1 | 3 | 3 | 5 |

→ ans = 3

Ex2:

| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| | -4 | -2 | 3 | 3 | 5 | 5 | 5 | 5 | 8 | 8 | 8 | 10 | 17 |
| #less: | 0 | 1 | 2 | 2 | 4 | 4 | 4 | 4 | 8 | 8 | 8 | 11 | 12 |

↳ ans: 3

Obs1: if element is the first occ. if (arr[i] != arr[i-1])

⇓

i == arr[i]

count of ele < ele

Obs2: if element is the repeat element.

↳ count of ele < ele will remain same as count of first occ.

Note: To check the first occ: arr[i] != arr[i-1]

//Pseudo code

```
int goodIntegers(int arr[N]){
    Arrays.sort(arr);
    int count=0;

    int less_count = 0;
    //H.W -> handle for 0th index
    for(int i=1; i<N; i++){
        if(arr[i] != arr[i-1]){
            less_count = i;
        }
        else{
            //Nothing
        }

        if(arr[i] == less_count){
            count ++;
        }
    }

}
```

T.C: O(N log N)

S.C: O(N)

```
int count=0;

int lesscount=0;

for(int i=1; i<N; i++){
    if(arr[i] != arr[i-1]){
        lesscount = i;
    }
    else{ x
        //Nothing
    }

    if(arr[i] == lesscount){
        count++;
    }
}
```

Ex2: { -4 -2 3 3 5 5 5 5 8 8 8 10 17 }

i: 0 1 2 3 4 5 6 7 8 9 10 11 12

#less: 0 1 2 2 4 4 4 4 8 8 8 11 12

count = 0

lesscount: 0

| i | lesscount |
|---|---|
| 1 | 1 |
| 2 | 2 |
| 3 | 2 |
| 4 | 4 |
| 5 | 4 |
| 6 | 4 |
| 7 | 4 |
| 8 | 8 |



Break till 9:50 PM

// Sorting techniques

## ① Bubble Sort

↳ Sort the array in asc. order but we can swap adjacent elements only.

arr[8]: { 5 7 5 4 10 -2 6 3 }
0 1 2 3 4 5 6 7

iter 0: { 5 7 5 4 10 -2 6 3 } → {0, N-2}
0 1 2 3 4 5 6 7
5 5 4 7 -2 6 3 10

{ 5 5 4 7 -2 6 3 10 }
i

iter 1: { 5 5 4 7 -2 6 3 10 } → {0, N-3}
0 1 2 3 4 5 6 7
4 5 -2 6 3 7

.
.
.
(
)

iter N-1:

↳ N-1 iterations

//Pseudo Code

```
void bubbleSort (int arr[n]){
    for(int i=0; i<n-1 ; i++){ //n-1 iterations
        for (int j=0; j<n-1-i;j++){
            if (arr[j] >arr[j+1]){
                int temp= arr[j];
                arr[j] = arr[j+1];
                arr[j+1] = temp;
            }
        }
    }
}
```

T.C : $O(N^2)$

S.C : $O(1)$

```
void bubblesort (int arr[n]){
    for (int i=0; i<n-1; i++){          // i<=n-2 ; n-1 iterations
        for (int j=0; j<n-1-i; j++){    // j<=n-2-i
            if (arr[j] > arr[j+1]){
                int temp = arr[j];
                arr[j] = arr[j+1];
                arr[j+1] = temp;
            }
        }
    }
}
```

arr[8] = { 5 4 6 4 -2 6 3 10 }

i = 0 ; j → 0 <= 6

i = 1 ; j → 0 <= 5

{ 5 4 5 -2 6 3 7 10 }

N elements

N-1 elements fix

Nth will be fixed auto

[0, N-2] → N-1 iterations

you didn't come this far only to come this far.