



University of Colorado **Boulder**

Fundamentals of Data Communications

IP Routing

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Review

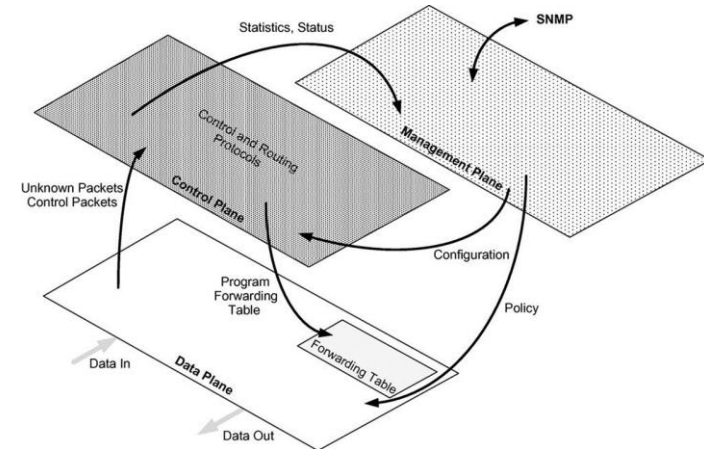
Static Routing

- **What is the purpose of routing?**
- **What is a static route?**
- **Static Routing vs Routing Protocols**
 - Manual
 - *Administrative Overhead*
 - *Scalability*
 - Static not dynamic
 - *(failover)*
- **Routing Table**

Router

- **Inside a router**

- Control plane: routing protocols
- Data plane: packet forwarding



- **Path selection**

- Minimum-hop/cost and shortest-path routing
- Algorithms: Link-state vs. Distance vector routing

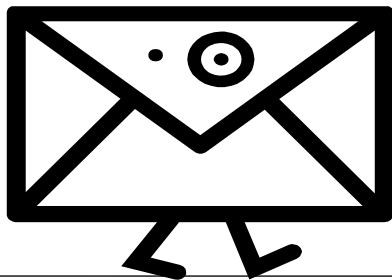
- **Topology change**

- Using beacons to detect topology changes
- Propagating topology information

What is Routing?

- **A famous quotation from RFC 791 (Internet Protocol 1981)**

“A name indicates what we seek.
An address indicates where it is.
A route indicates how we get there.”
-- Jon Postel



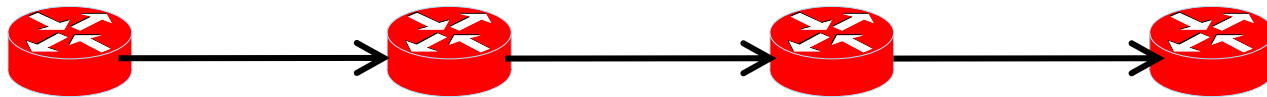
Routing vs. Forwarding

- **Routing: control plane**

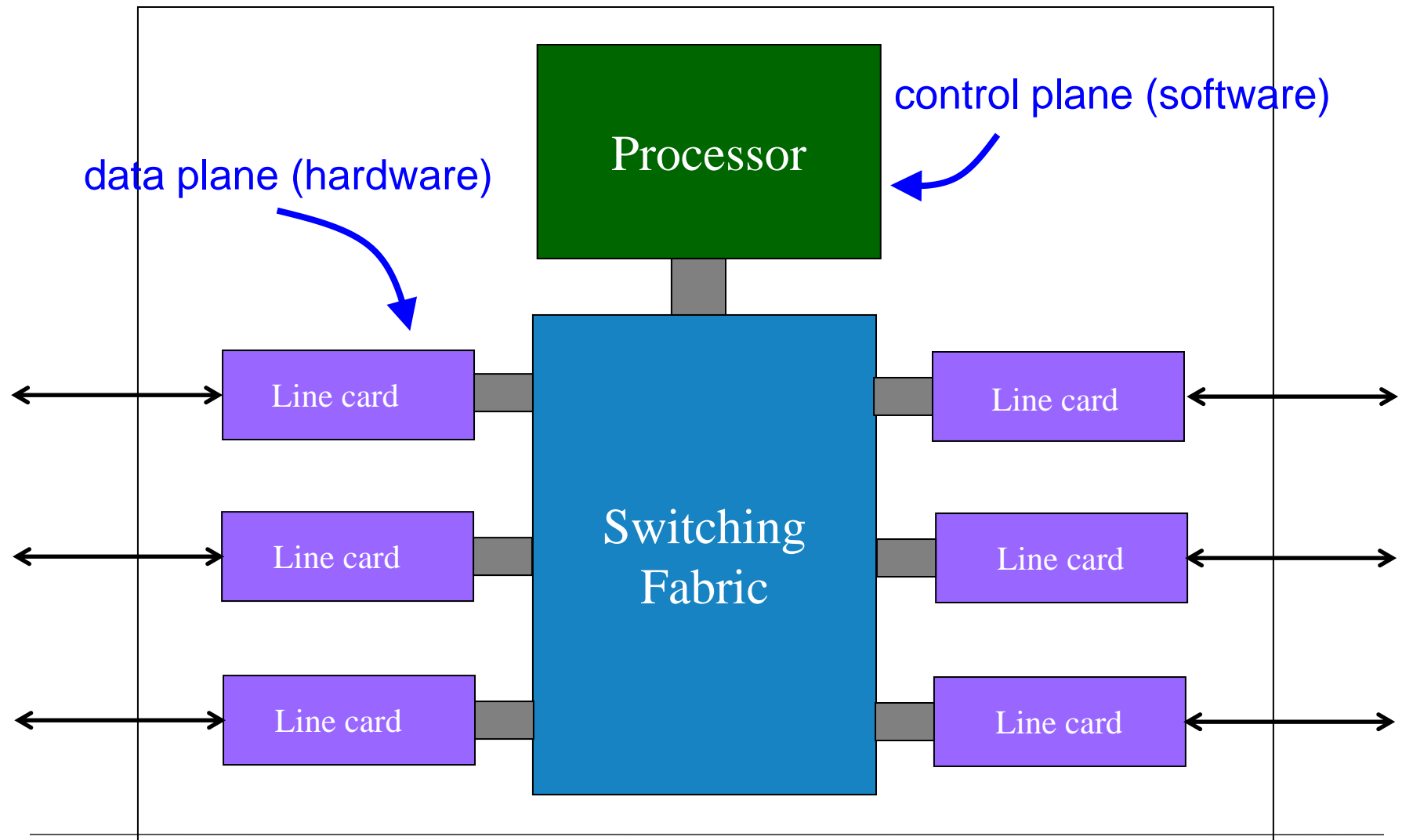
- Computing paths the packets will follow
- Routers talking amongst themselves
- Individual router creating a forwarding table

- **Forwarding: data plane**

- Directing a data packet to an outgoing link
- Individual router using a forwarding table



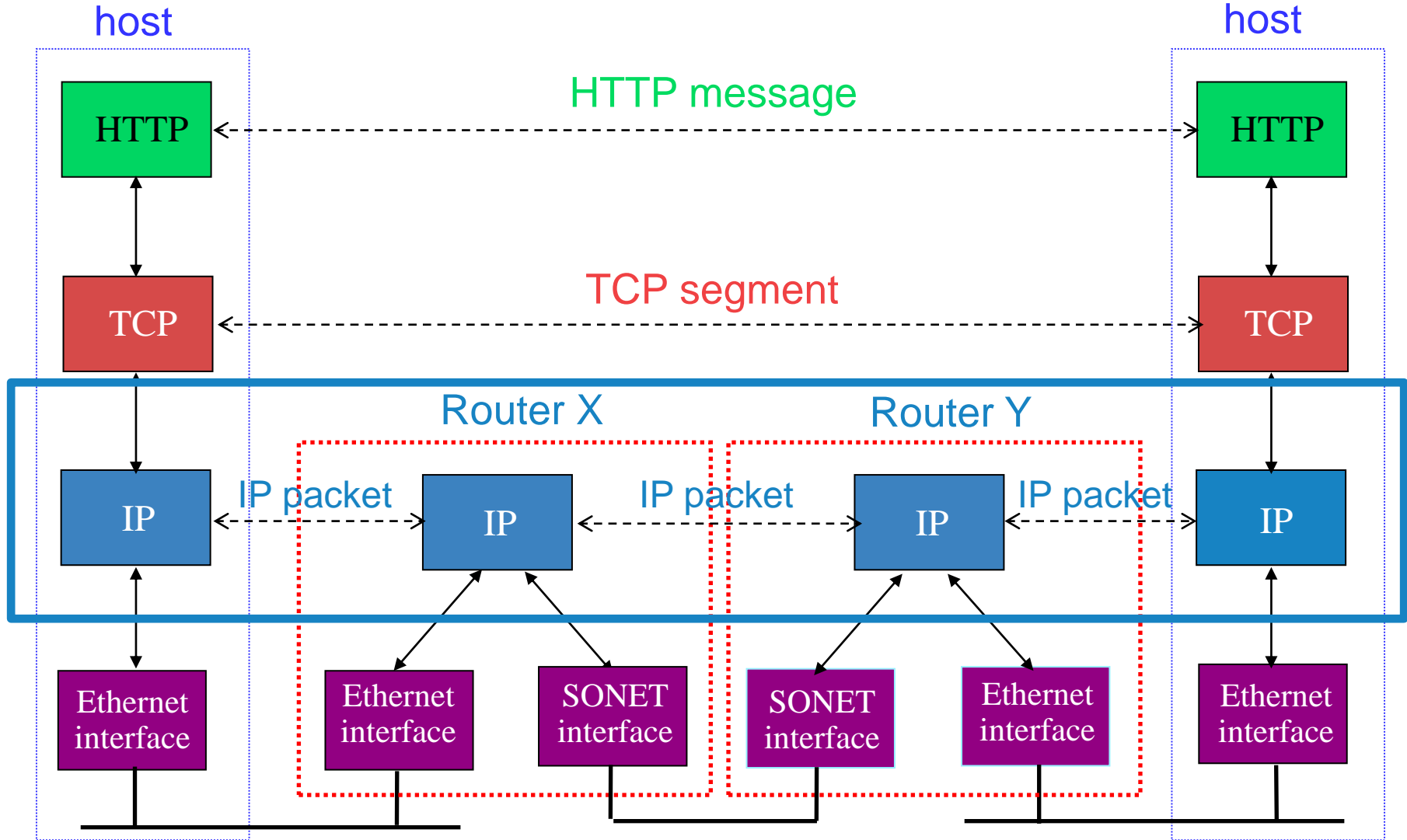
Data and Control Planes



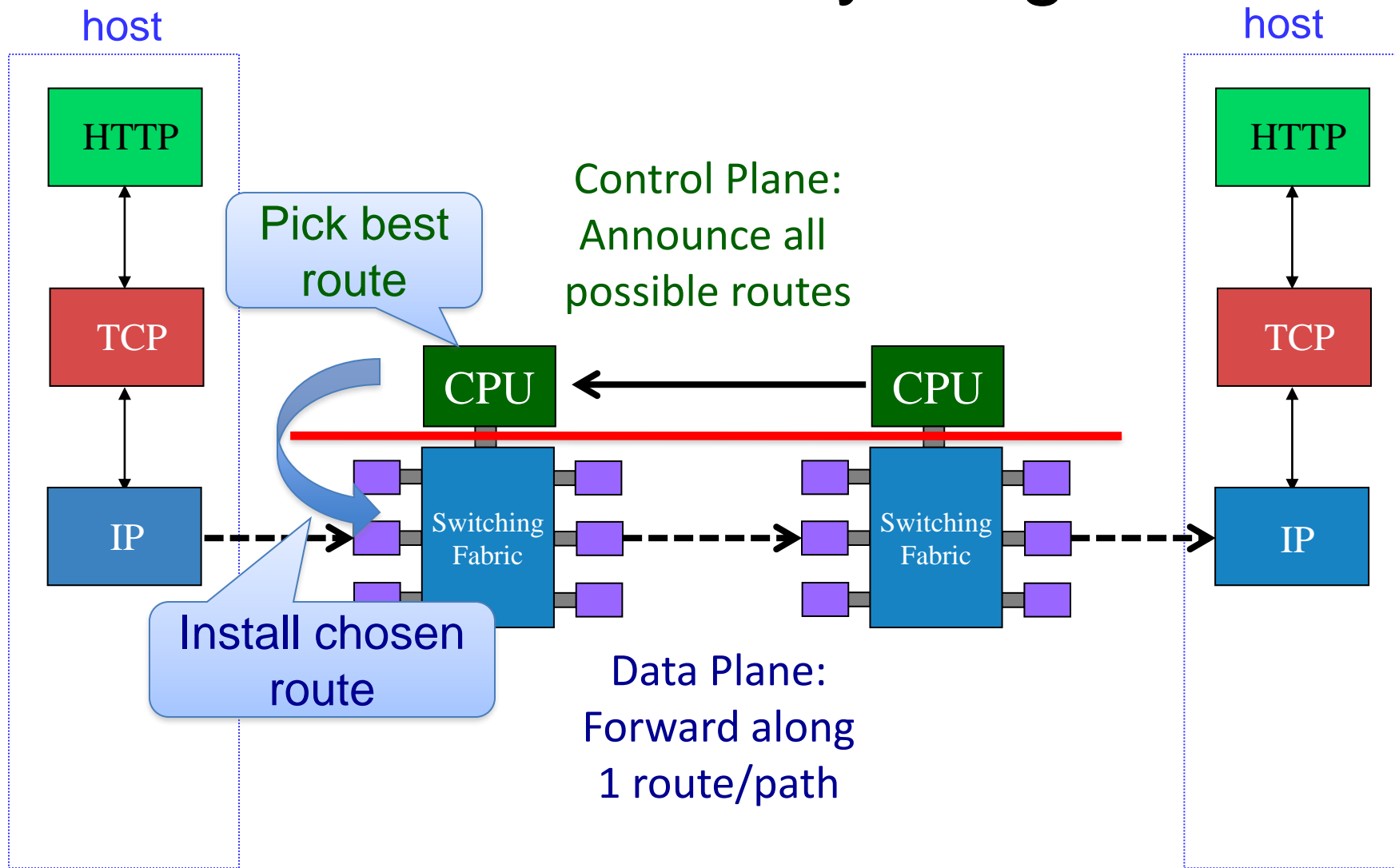
Where do Forwarding Tables come From?

- **Routers have forwarding tables**
 - Map IP prefix (subnet) to outgoing link(s)
- **Entries can be statically configured**
 - E.g., “map 12.34.158.0/24 to fastethernet 0/0.1”
 - But, this doesn’t adapt
 - *To failures*
 - *To new equipment*
 - *To the need to balance load*
 - That is where routing protocols come in!

Recall the Internet Layering Model



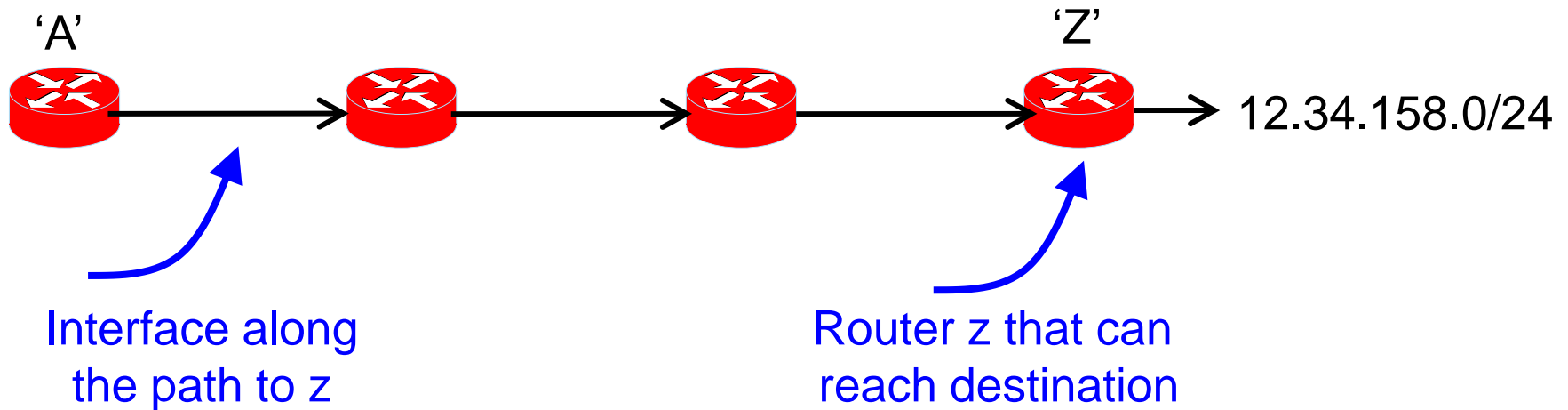
Recall the Internet Layering Model



Computing Paths Between Routers

- **Routers need to know two things:**

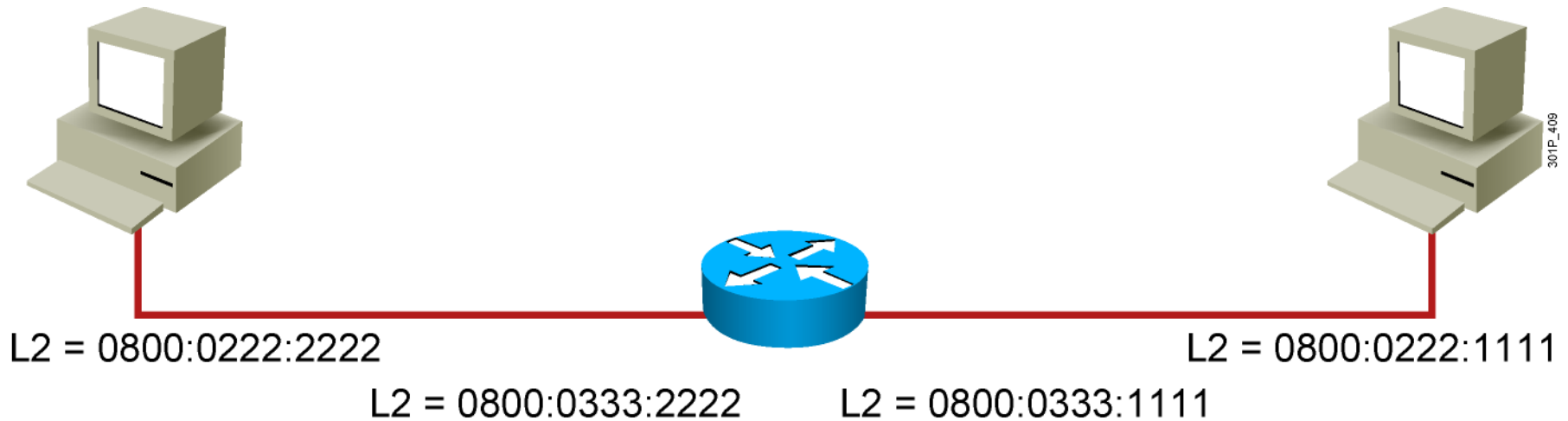
1. Which router to use to reach a destination prefix
2. Which outgoing interface to use to reach that router



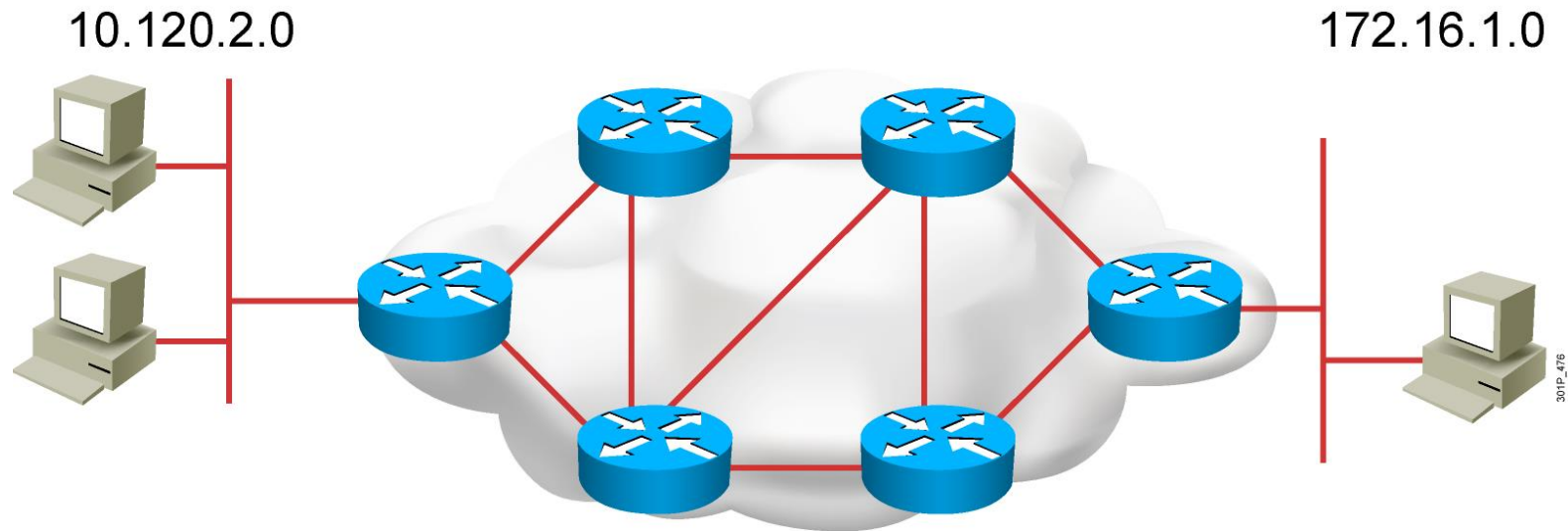
Routers and Routing

- **Routing**
 - The process by which data/traffic gets from one location to another
- **A router is the device used to route traffic**
- **Routers can forward packets over static routes or dynamic routes**
 - Based on the router configuration and network design
- **Static routes**
 - A route that a network administrator enters into the router manually
 - Unidirectional static routes must be configured **to** and **from** networks to allow bidirectional communications to occur
- **Dynamic routes**
 - A route that a network routing protocol adjusts automatically for topology or traffic changes

Layer 2 Addressing

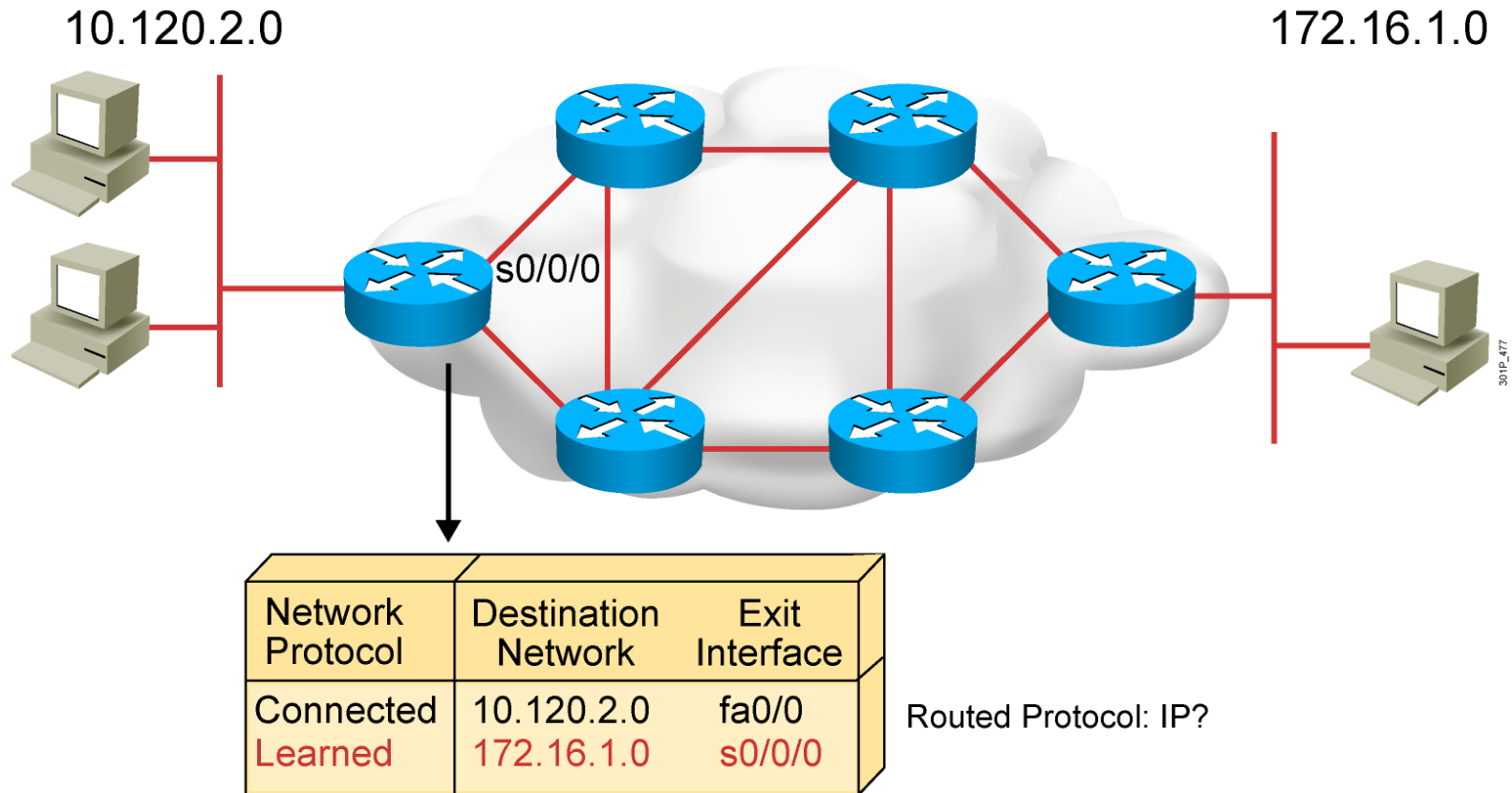


Router Operations



- **A router needs to do the following:**
 - Know the destination address
 - Identify the sources from which the router can learn
 - Discover possible routes to the intended destination
 - Select the best route
 - Maintain and verify routing information

Router Operations (Cont.)



- Routers must learn destinations that are not directly connected
 - Routers automatically learn connected routes (if the interface is up!)

Identifying Static and Dynamic Routes

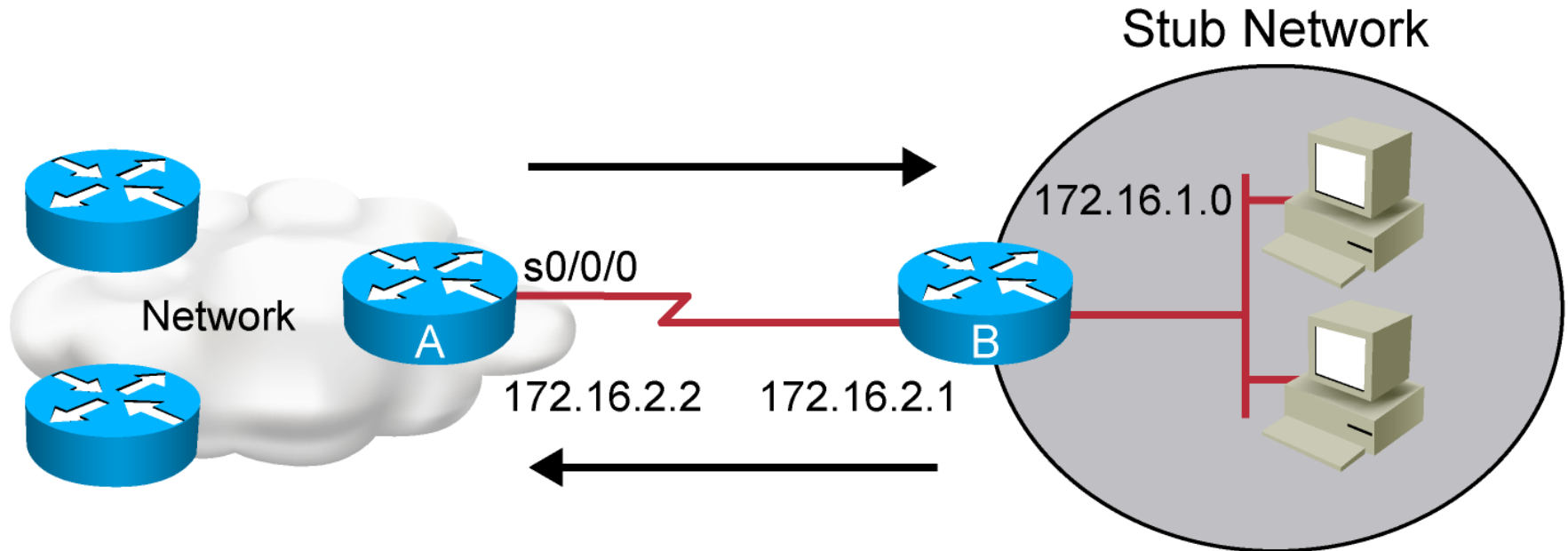
- **Static route**

- Uses a route that a network administrator enters into the router manually

- **Dynamic route**

- Uses a route that a network routing protocol adjusts automatically for topology or traffic changes

Static Routes



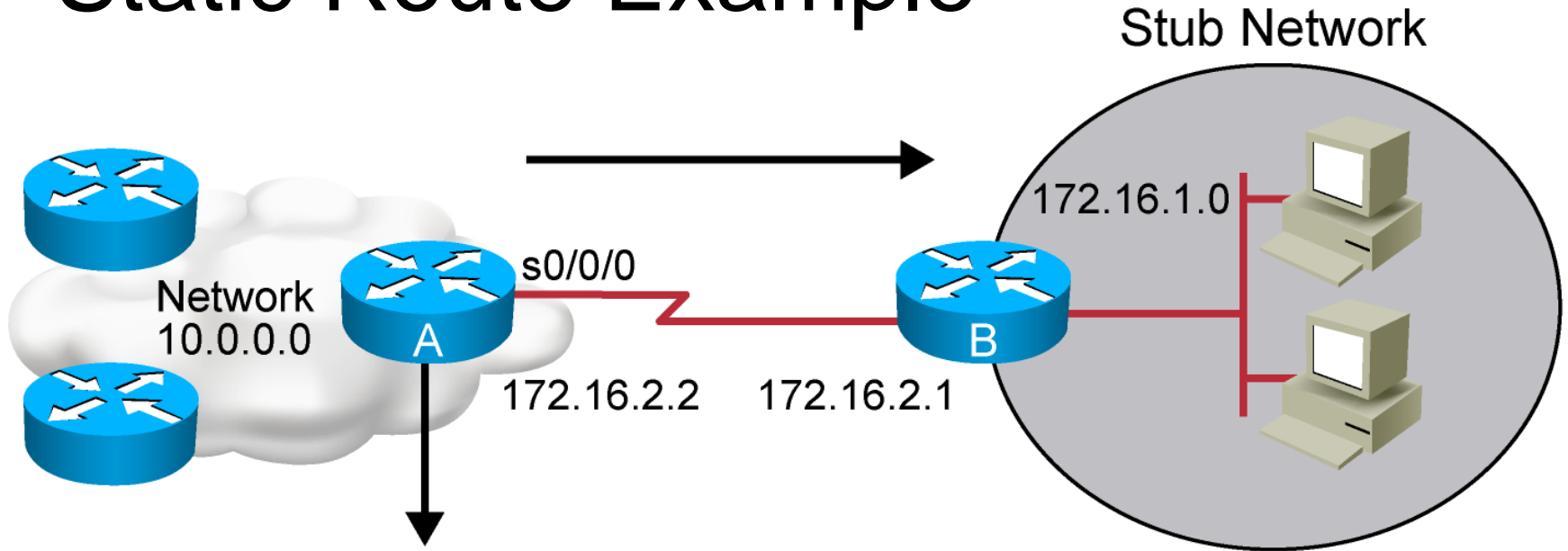
Configure unidirectional static routes **to** and **from** a network to allow communications to occur.

Static Route Configuration

```
(config)# ip route network [mask]  
{address | interface}[distance] [permanent]
```

- Defines a path to an IP destination network or subnet or host
- Address = IP address of the next hop router
- Interface = outbound interface of the local router
- Example:
 - ***ip route 192.168.1.0 255.255.255.0 192.168.222.254***
 - ****pro tip***

Static Route Example



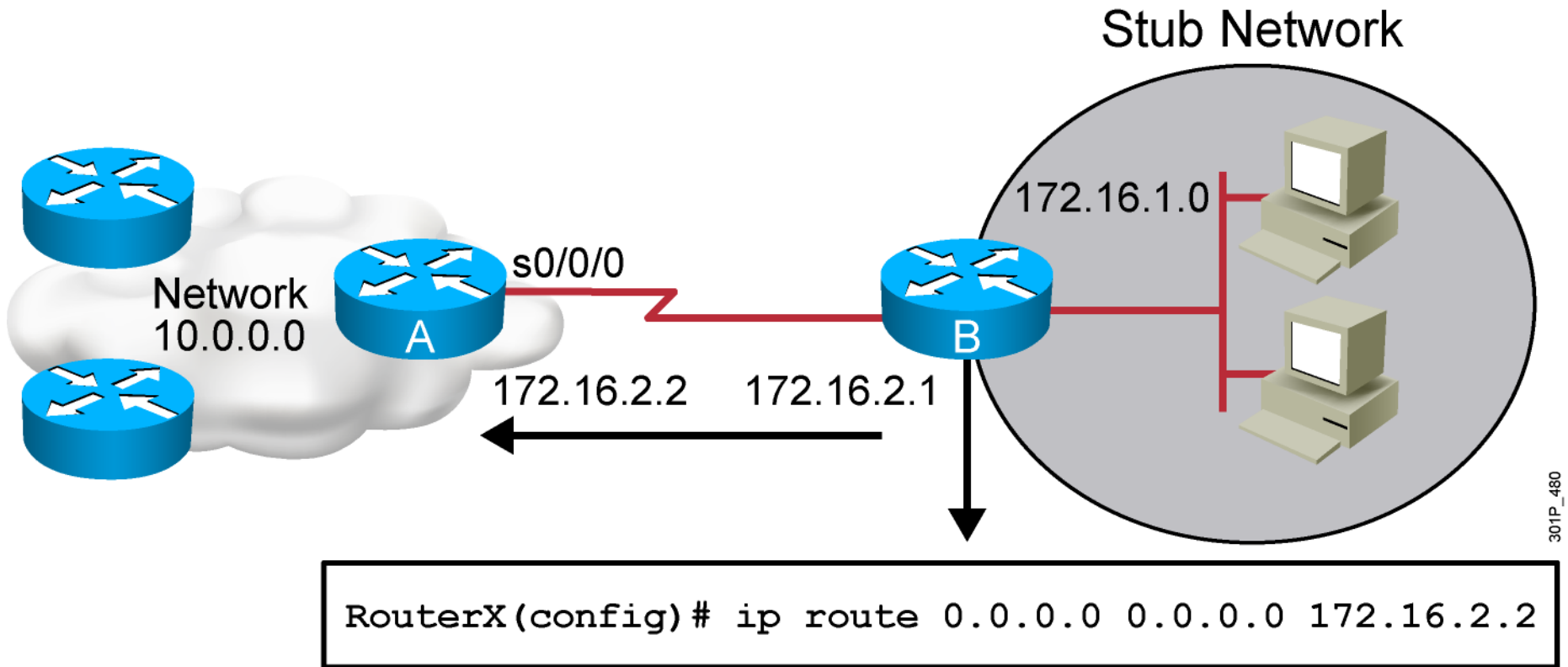
```
(config)# ip route 172.16.1.0 255.255.255.0 172.16.2.1
```

or

```
(config)#ip route 172.16.1.0 255.255.255.0 s0/0/0
```

- This is a unidirectional route.
- You must have a route configured in the opposite direction for traffic to return.

Default Routes



- This route allows the “stub network” (behind RouterB) to reach all networks beyond Router A.

Verifying the Static Route Configuration

- **Route Table**

```
RouterX# show ip route
```

```
Codes: C - connected, S - static, I - IGRP, R - RIP, M - mobile, B - BGP  
       D - EIGRP, EX - EIGRP external, O - OSPF, IA - OSPF inter area  
       E1 - OSPF external type 1, E2 - OSPF external type 2, E - EGP  
       i - IS-IS, L1 - IS-IS level-1, L2 - IS-IS level-2, * - candidate default  
       U - per-user static route
```

```
Gateway of last resort is 0.0.0.0 to network 0.0.0.0
```

```
    10.0.0.0/8 is subnetted, 1 subnets
```

```
C        10.1.1.0 is directly connected, Serial0/0/0
```

```
S*    0.0.0.0/0 is directly connected, Serial0
```

P2P with Internet

Floating Static Routes

Dynamic Routing

Overview and Terminology

- Think beyond directly connected devices
- “Network Map” and Routing tables
- Static Routes (type the Internet)
 - *Nearly 1 million routes!*
- **Routing protocols:** Propagate information about available networks
- **Routed protocols:** L3 protocols that can be routed (IP)

Overview and Terminology

- ***Routing type***: Link State vs. Distance Vector vs. Path Vector
- ***Exterior Routing Protocols (EGP)***: Inter-corporate route distribution
- ***Interior Routing Protocols (IGP)***: Intra-corporate route exchange
- ***Autonomous System (AS)***: Network under single management control (ISP)

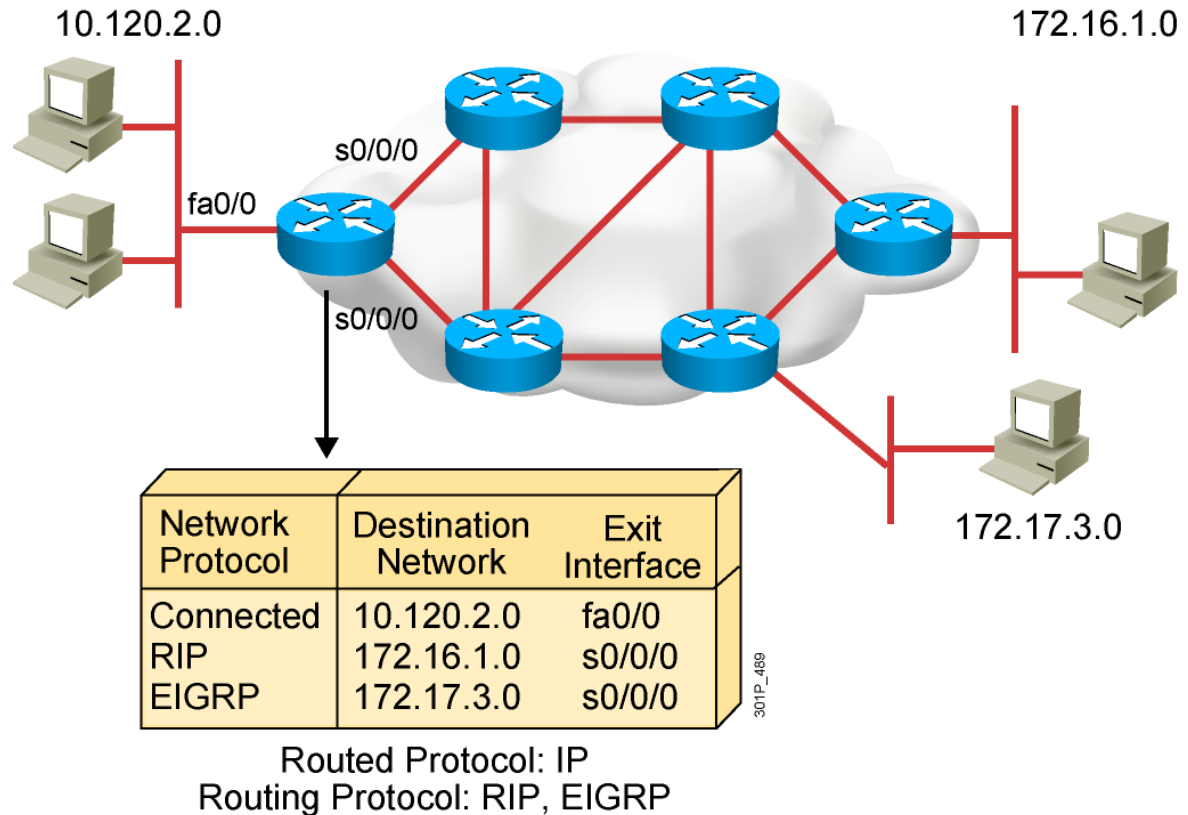
Routing Objectives

- Build a routing table (learn & propagate)
- Pick best routes (if more than one available)
- Remove invalid routes
- Replace routes if better advertisement received
- Restore routes fast (convergence time)
- Prevent loops



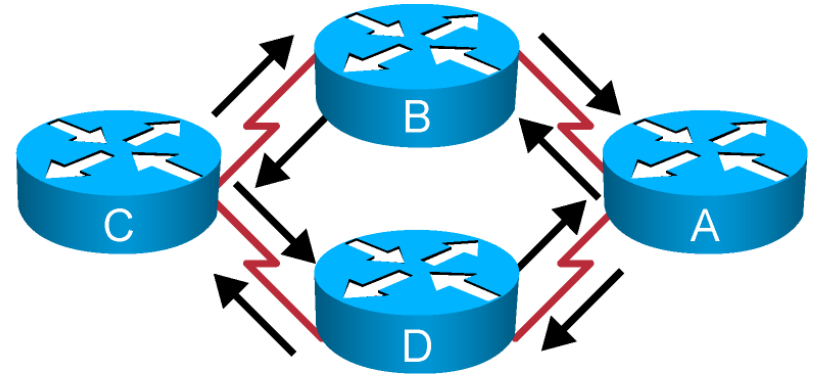
What is a Routing Protocol?

- **Routing** protocols are used between routers to determine paths and maintain routing tables.
- After the path is determined, a router can route a **routed** protocol.

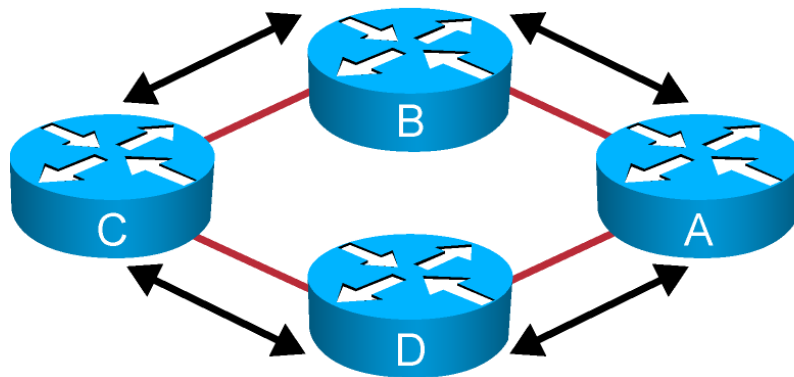


Classes of Routing Protocols

Distance Vector



Hybrid Routing



Link-State

Interior Gateway Routing Protocols (IGP)

- Several available in the market. Protocols have evolved over time to respond to changing network conditions
- Considerations while choosing a protocol:
 - **Type of routing protocol**
 - *Distance vector*
 - *Link-state*
 - *Hybrid*
 - **Update process**
 - *Full updates*
 - *Partial updates*
 - **Convergence time**
 - **Metric**
 - *Measure of link quality / preference*
 - **VLSM support**
 - *Subnet size variations*
 - *Permit better address allocation*
 - **Classless vs. Classful**
 - *TX or Not TX mask info*
 - *If TX then VLSM supported*

Distance Vector Protocols

- **Distance Vector Logic**

- Add directly connected
- Send updates out include directly connected and learned routes
 - *“Routing by Rumor”*
- Listen for routing updates
- Routing info: subnet & metric
- Use broadcast or multicast for updates
- Chose best of multiple routes
- Send/expect periodic full updates
- If updates no longer received, remove routes learned from such neighbor
- Assume that the advertising router is the next hop for a route

Distance Vector Protocols

- **RIPv1**

- Metric: hop count (smaller better) Infinite metric = 16
- Full updates every 30s (subnet # & metric)
- Convergence 3-5 minutes (depending on size of routing table)
- Classful (no VLSM)
- Broadcasts updates (255.255.255.255)
- “Route-poisoning” and split horizon
 - *Infinite metric of 16*
 - *Prevent advertising back to the interface it was received*

- **RIPv2**

- Same features as RIPv1
- Adds: VLSM support (updates include subnet #, mask & metric)
- MD5 and plain text authentication
- Includes next hop router IP on updates
- Uses external route tags (redistribution)
- Multicast routing updates (224.0.0.9)

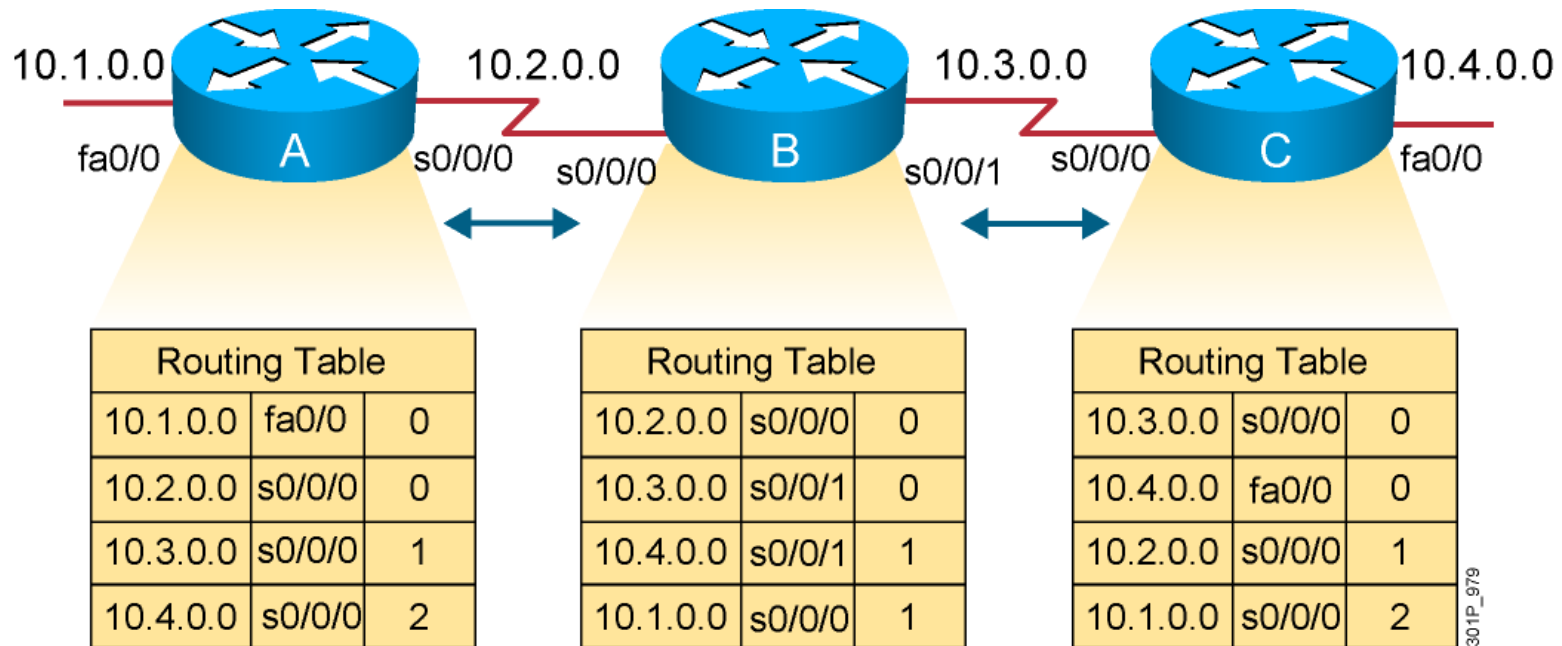


Distance Vector Protocols

- **IGRP**

- Proprietary
- Metric: bandwidth + delay (default), reliability, load, MTU
- Metric value (1 - 4 billion), smaller better.
- Metrics are cumulative: multi-hop path adds links delays
- Metrics are configurable
- Full updates every 90s
- No VLSM supported
- Infinite metric = 4,294,967,295 (smaller better)

Sources of Information and Discovering Routes



- Routers discover the best path to destinations from each neighbor.



Link State Protocols

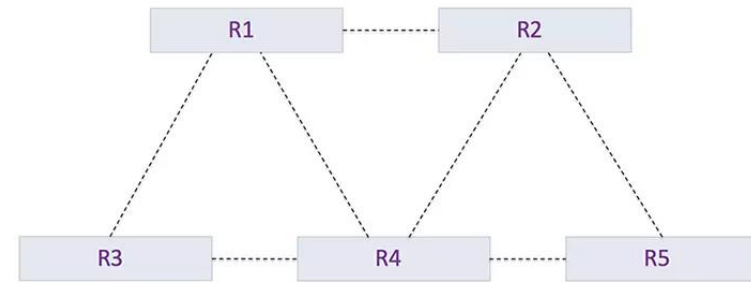
- **Link State logic**

- Add directly connected
- Don't transmit routing info until neighbors are discovered
- Routing info: topological information about the network, at the end every router has a complete map of the network.
- Use reliable protocol for route updates
- Calculate shortest path first (SPF) algorithm (Dijkstra) to determine best routes and next hop (also prevents loops)
- Full updates at start and after long periods of time (~30 min.)
- Partial updates when a link fails
- Fast convergence

Link State Protocols

- **OSPF**

- Most popular
- Discover neighbors, then exchange routes
- Reliable Transport Protocol
- Run SPF and store best routes
- More memory required, more processing
- Metric: “cost” based on bandwidth (smaller-better)
- Averages 10s convergence time
- VLSM supported
- Hello messages to confirm neighbor reachability
- Full updates every 30 min
- Partial updates when link fails



Interface type	bandwidth	Cost
Fast Ethernet and faster	100 Mb/s and higher	1
Ethernet	10 Mb/s	10
E1	2 Mb/s	48
T1	1.544 Mb/s	64
128bps	128bps	781
64kbps	64kbps	1562
56kbps	56kbps	1785

Hybrid Protocols

- **EIGRP**

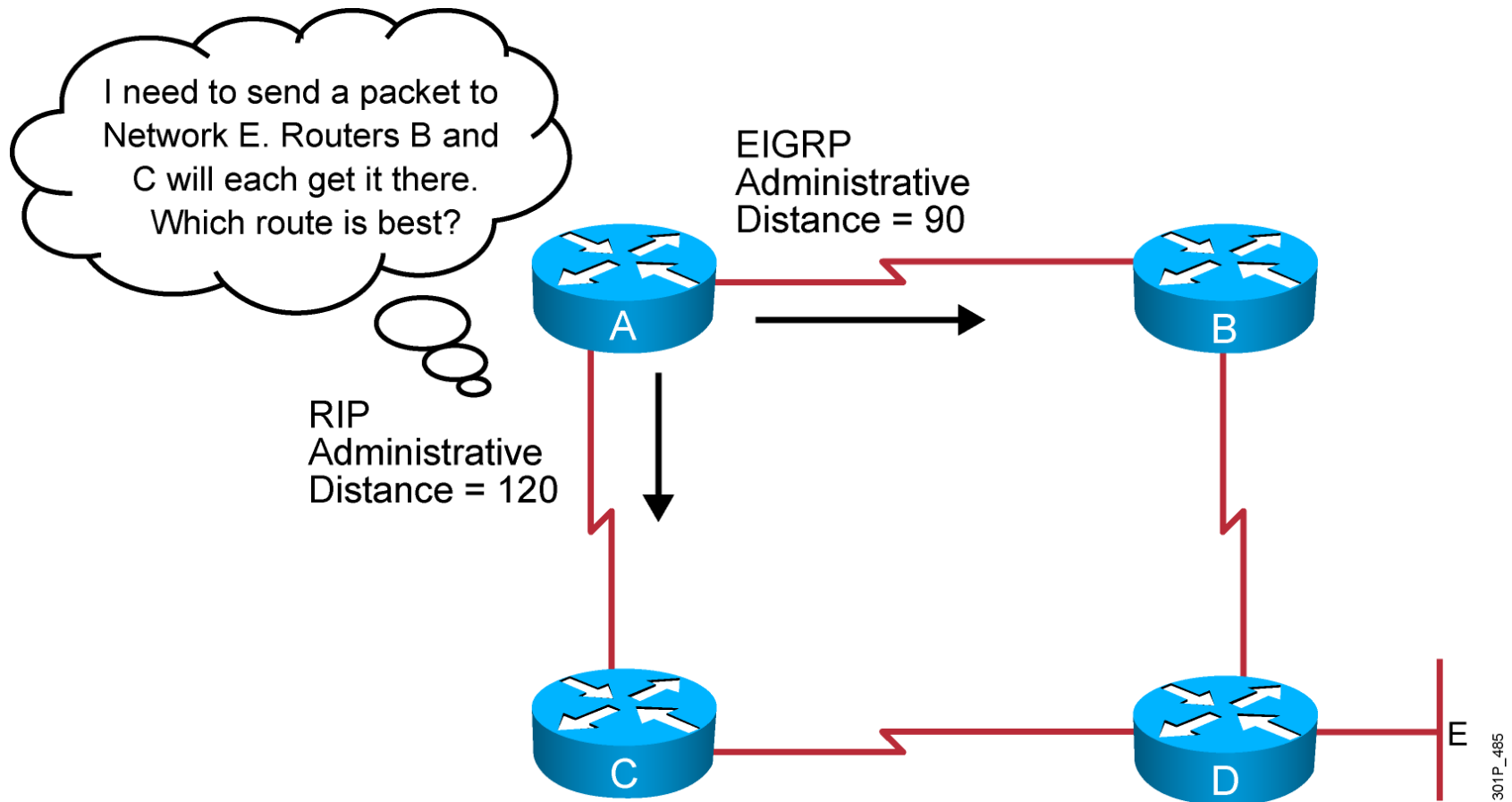
- Cisco Proprietary*
- Uses features from both link-state and distance vector protocols (balanced hybrid protocol: topology + routing tables)
- Diffused Update Algorithm (DUAL): exchange more info than distance vector but less than link-state
- Demands less computation
- Calculated best and “feasible successor routes” (alternatives), permits **immediate convergence**.
- Discovers neighbors like OSPF
- Metric based on bandwidth and delay (Idem IGRP * 256)
- Does not send periodic full updates

Loop Prevention in Distance Vector Protocols

- **Problems:**

- Multiple routes to same network
- Miscommunications on a single link
- Information loops through alternative paths
- Counting to infinity

Administrative Distance: Ranking Routes



Administrative Distance

- Applies to similar routes received from different protocols
- Prioritizes different routing protocols (which protocol do I believe first?)
- **The lower the AD number, the better**
- Examples:
 - ***Directly connected*** **0**
 - ***Static Routes*** **1**
 - ***eBGP*** **20**
 - ***EIGRP (internal)*** **90**
 - ***IGRP*** **100**
 - ***OSPF*** **110**
 - ***IS-IS*** **115**
 - ***RIP*** **120**
 - ***EIGRP (external)*** **170**
 - ***iBGP*** **200**

Classful Routing Protocol

- Classful routing protocols do not include the subnet mask with the route advertisement.
- Within the same network, consistency of the subnet masks is assumed.
- Summary routes are exchanged between foreign networks.
- These are examples of classful routing protocols:
 - ***RIPv1***
 - ***IGRP***

Classless Routing Protocol

- Classless routing protocols include the subnet mask with the route advertisement.
- Classless routing protocols support a variable-length subnet mask (VLSM). (aka subnetting)
- Summary routes can be manually controlled within the network.
- These are examples of classless routing protocols:
 - **RIPv2**
 - **EIGRP**
 - **OSPF**
 - **IS-IS**

Link-State Routing can be Problematic

- **Topology information is flooded**
 - High bandwidth and storage overhead
 - Forces nodes to divulge sensitive information
- **Entire path computed locally per node**
 - High processing overhead in a large network
- **Minimizes some notion of total distance**
 - Works only if policy is shared and uniform
- **Typically used only inside an AS**
 - E.g., OSPF and IS-IS

Distance Vector is on the Right Track

- **Advantages**

- Hides details of the network topology
- Nodes determine only “next hop” toward the destination

- **Disadvantages**

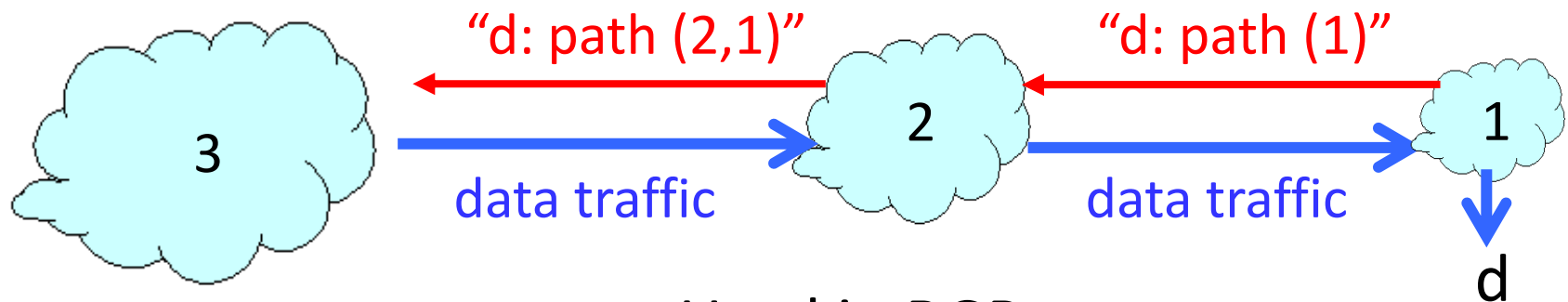
- Minimizes some notion of total distance, which is difficult in an interdomain setting
- Slow convergence due to the counting-to-infinity problem (“bad news travels slowly”)

- **Idea: extend the notion of a distance vector**

- To make it easier to detect loops

Path-Vector Routing

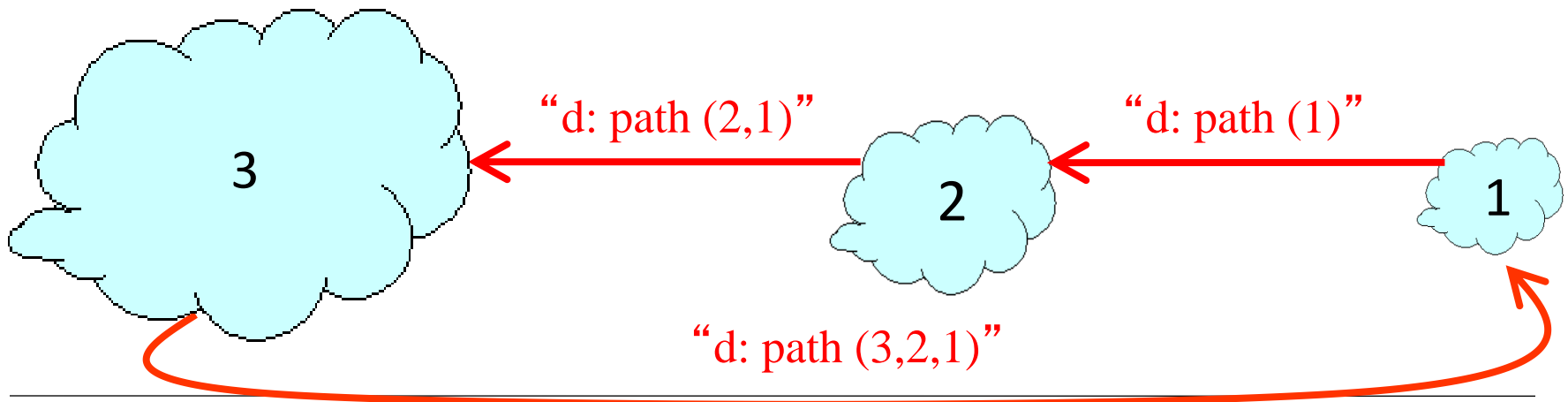
- **Extension of distance-vector routing**
 - Support flexible routing policies
- **Key idea: advertise the entire path**
 - Distance vector: send *distance metric* per dest d
 - Path vector: send the *entire path* for each dest d



Used in BGP

Faster Loop Detection

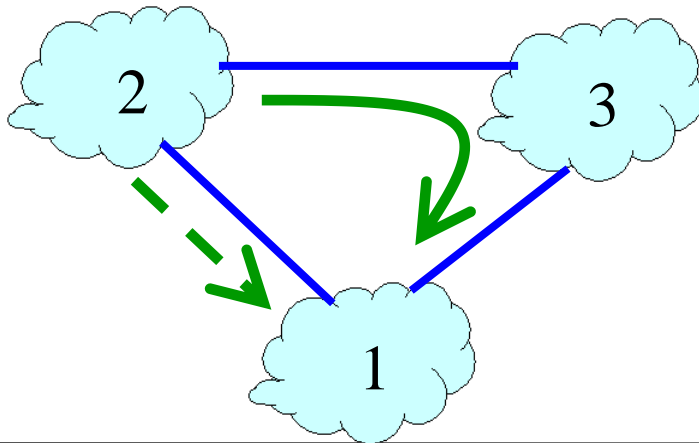
- **Node can easily detect a loop**
 - Look for its own node identifier in the path
 - E.g., node 1 sees itself in the path “3, 2, 1”
- **Node can simply discard paths with loops**
 - E.g., node 1 simply discards the advertisement



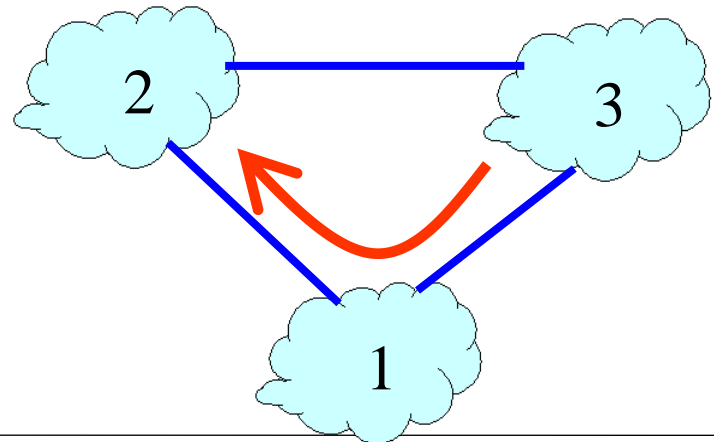
Path-Vector: Flexible Policies

- **Each node can apply local policies**
 - Path selection: Which path to use?
 - Path export: Which paths to advertise?

Node 2 prefers
“2, 3, 1” over “2, 1”

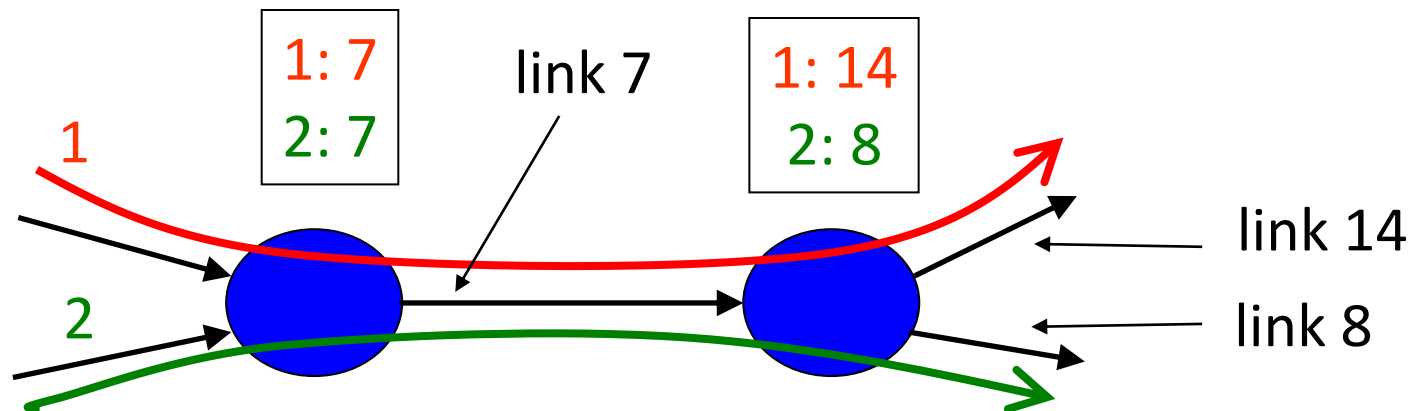


Node 1 doesn't let 3
hear the path “1, 2”



End-to-End Signaling

- **Establish end-to-end path in advance**
 - Learn the topology (as in link-state routing)
 - End host or router computes and signals a path
 - **Signaling:** *install entry for each circuit at each hop*
 - **Forwarding:** *look up the circuit id in the table*



Used in MPLS with RSVP

Distance Vector vs. Path Vector

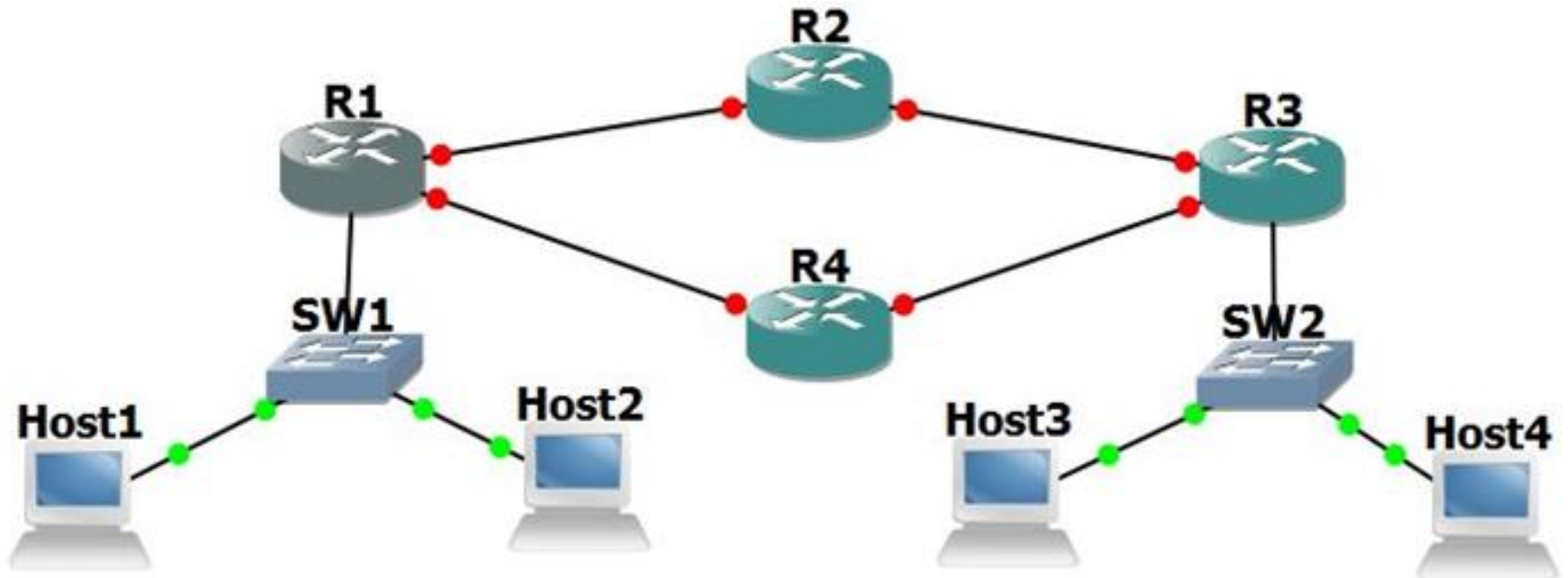
- **Distance-vector routing**

- Pro: Less information and computation than link state
- Con: Slower convergence (e.g., count to infinity)

- **Path-vector routing**

- Share entire path, not distance: faster convergence
- More flexibility in selecting paths

- **Different goals / metrics if inter- or intra-domain**



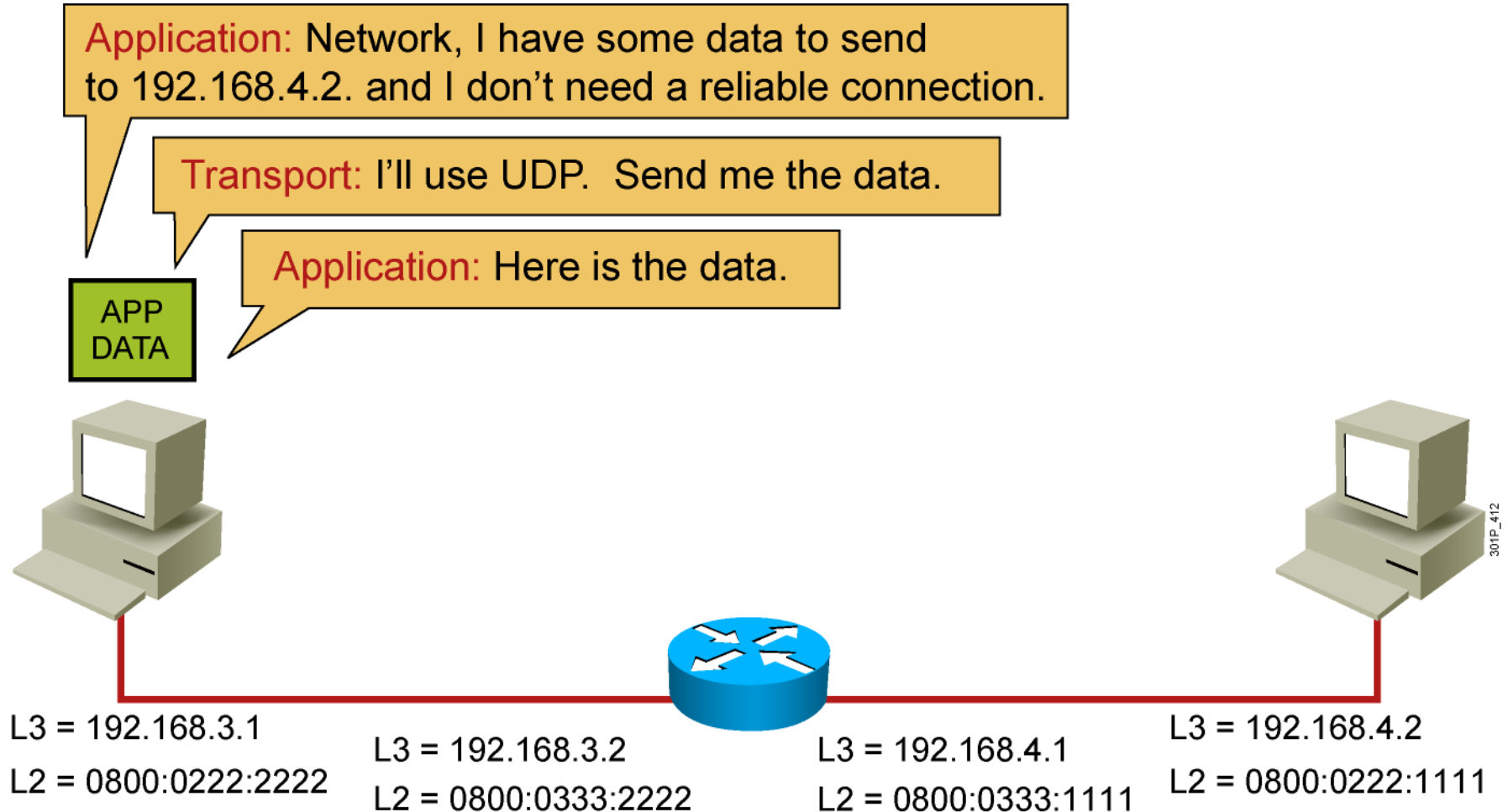
Questions?



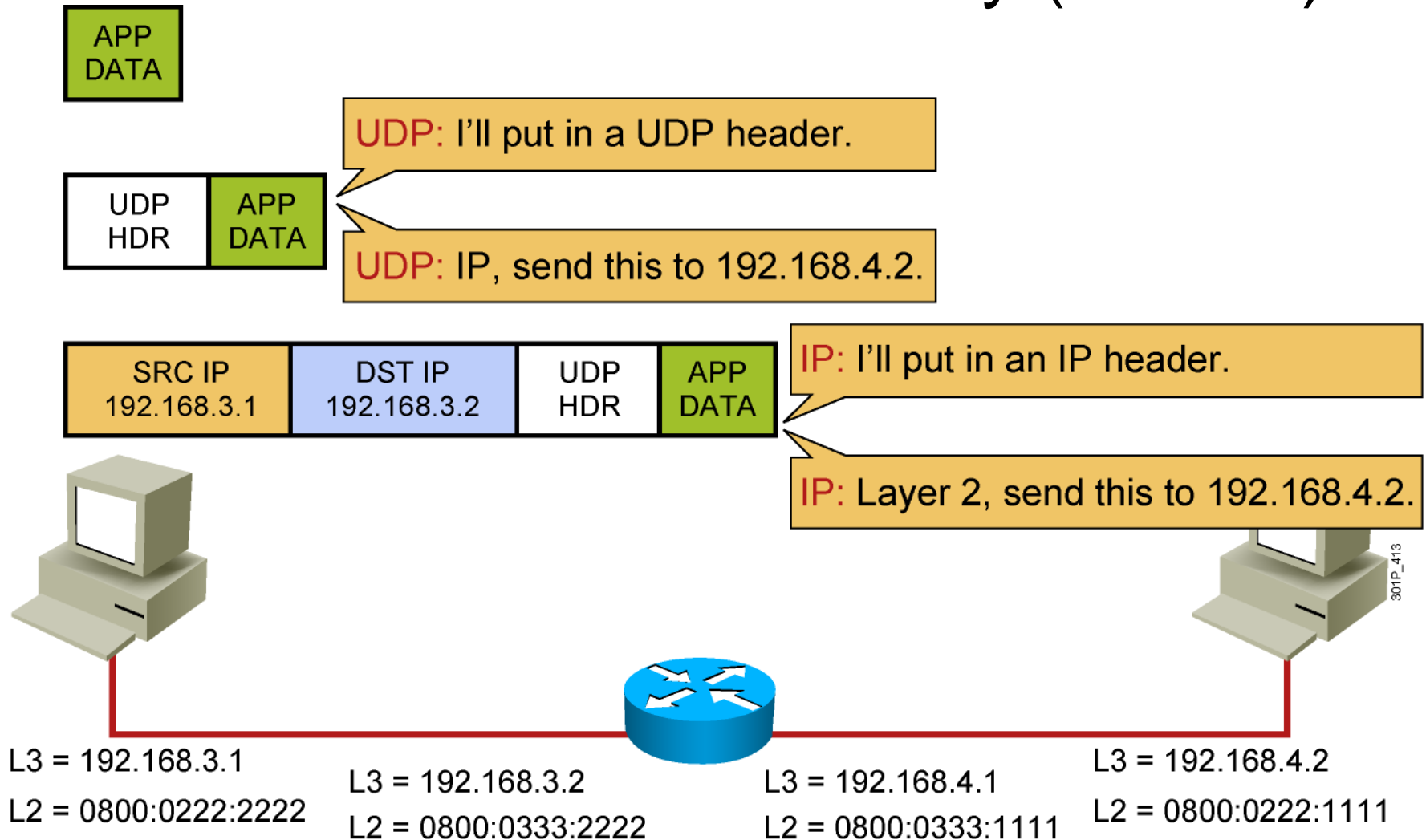
Lab

Appendix A

Host-to-Host Packet Delivery (1 of 17)



Host-to-Host Packet Delivery (2 of 17)

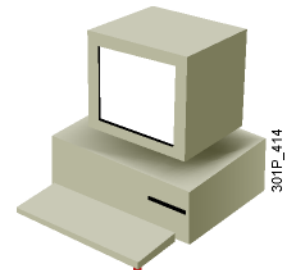
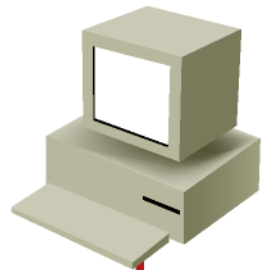


Host-to-Host Packet Delivery (3 of 17)

Layer 2: ARP, do you have a mapping for 192.168.4.2?

ARP: No, Layer 2 will have to hold the packet while I resolve the addressing.

SRC IP 192.168.3.1	DST IP 192.168.4.2	UDP HDR	APP DATA
-----------------------	-----------------------	------------	-------------



L3 = 192.168.3.1
L2 = 0800:0222:2222

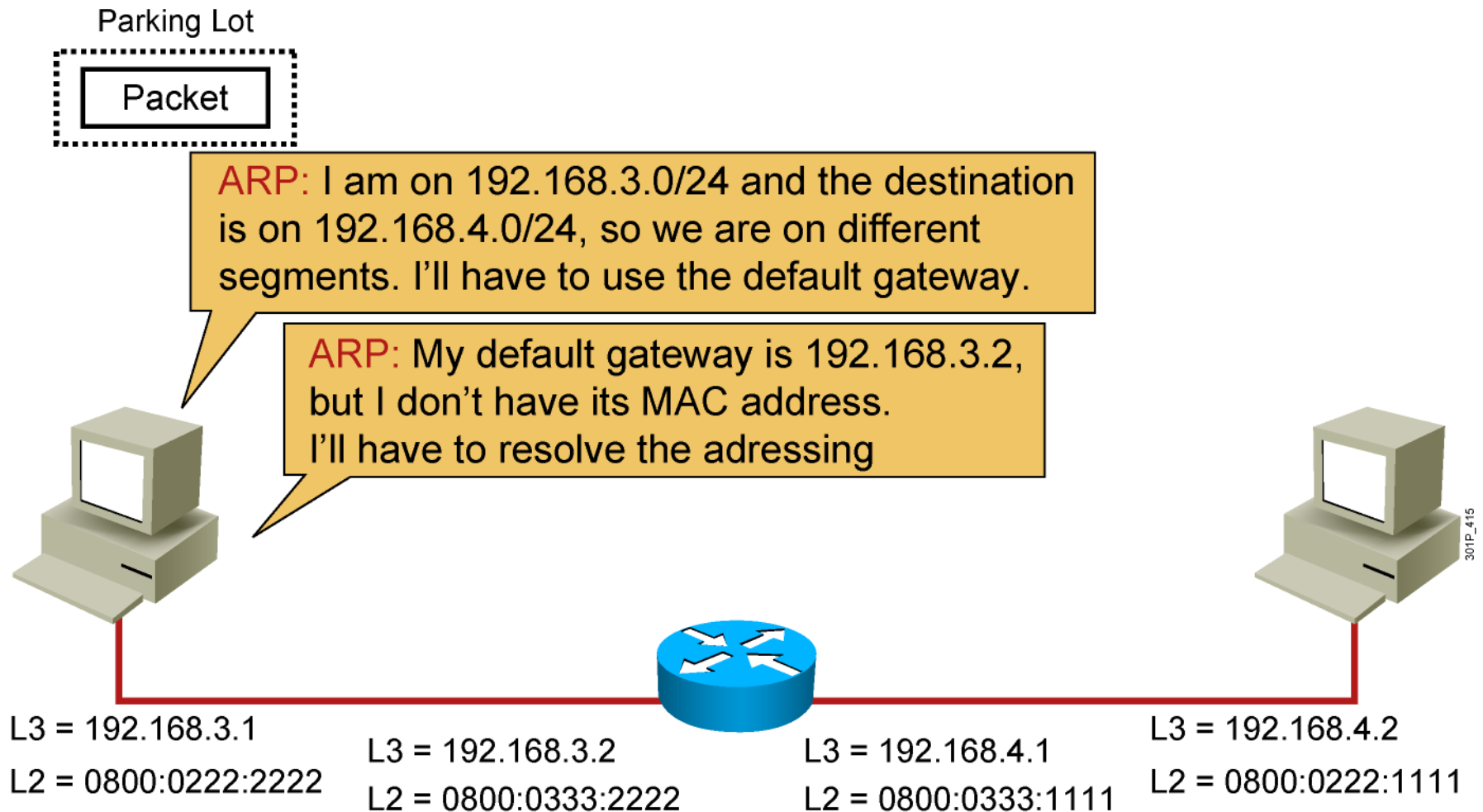
L3 = 192.168.3.2
L2 = 0800:0333:2222

L3 = 192.168.4.1
L2 = 0800:0333:1111

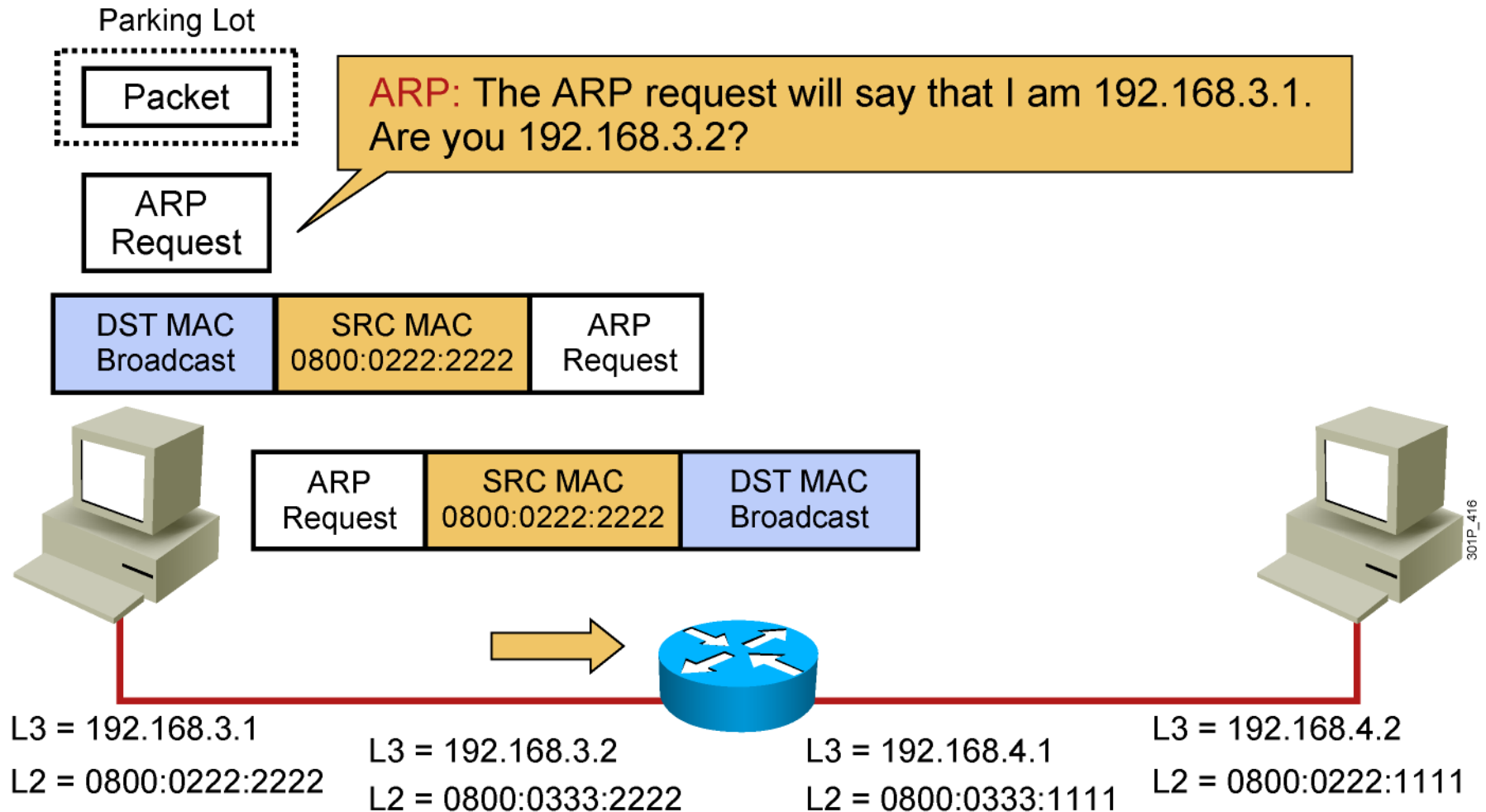
L3 = 192.168.4.2
L2 = 0800:0222:1111



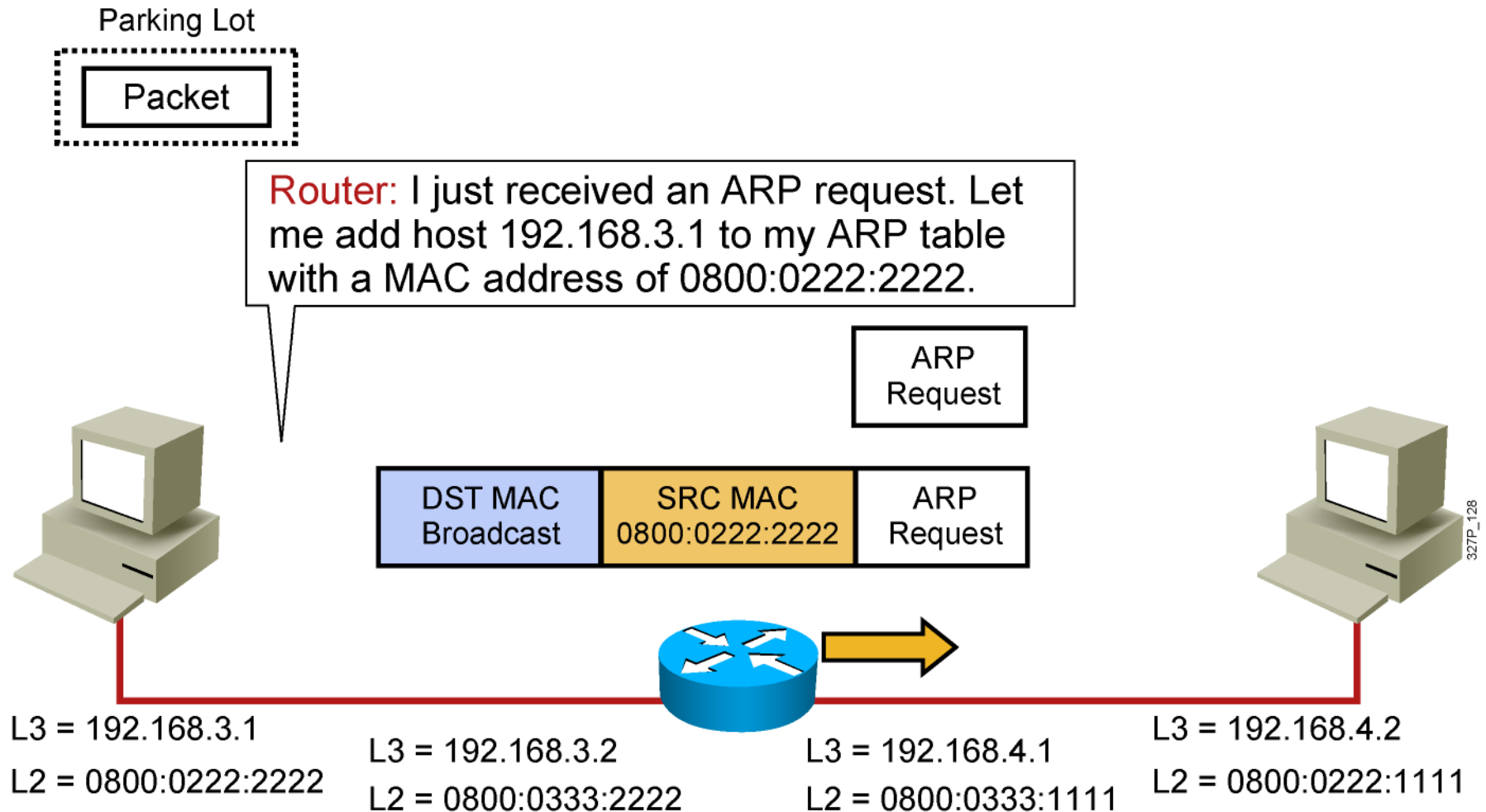
Host-to-Host Packet Delivery (4 of 17)



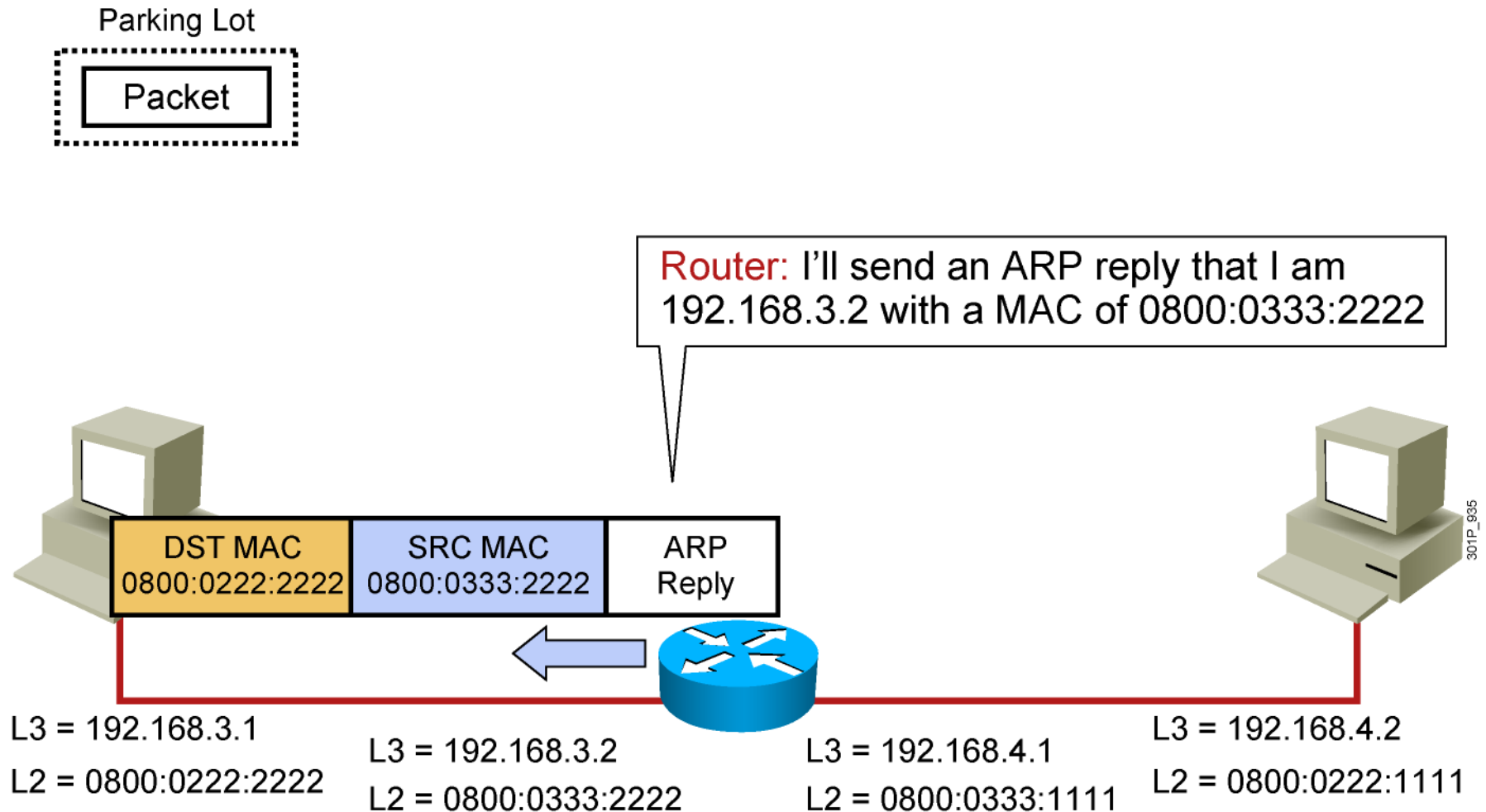
Host-to-Host Packet Delivery (5 of 17)



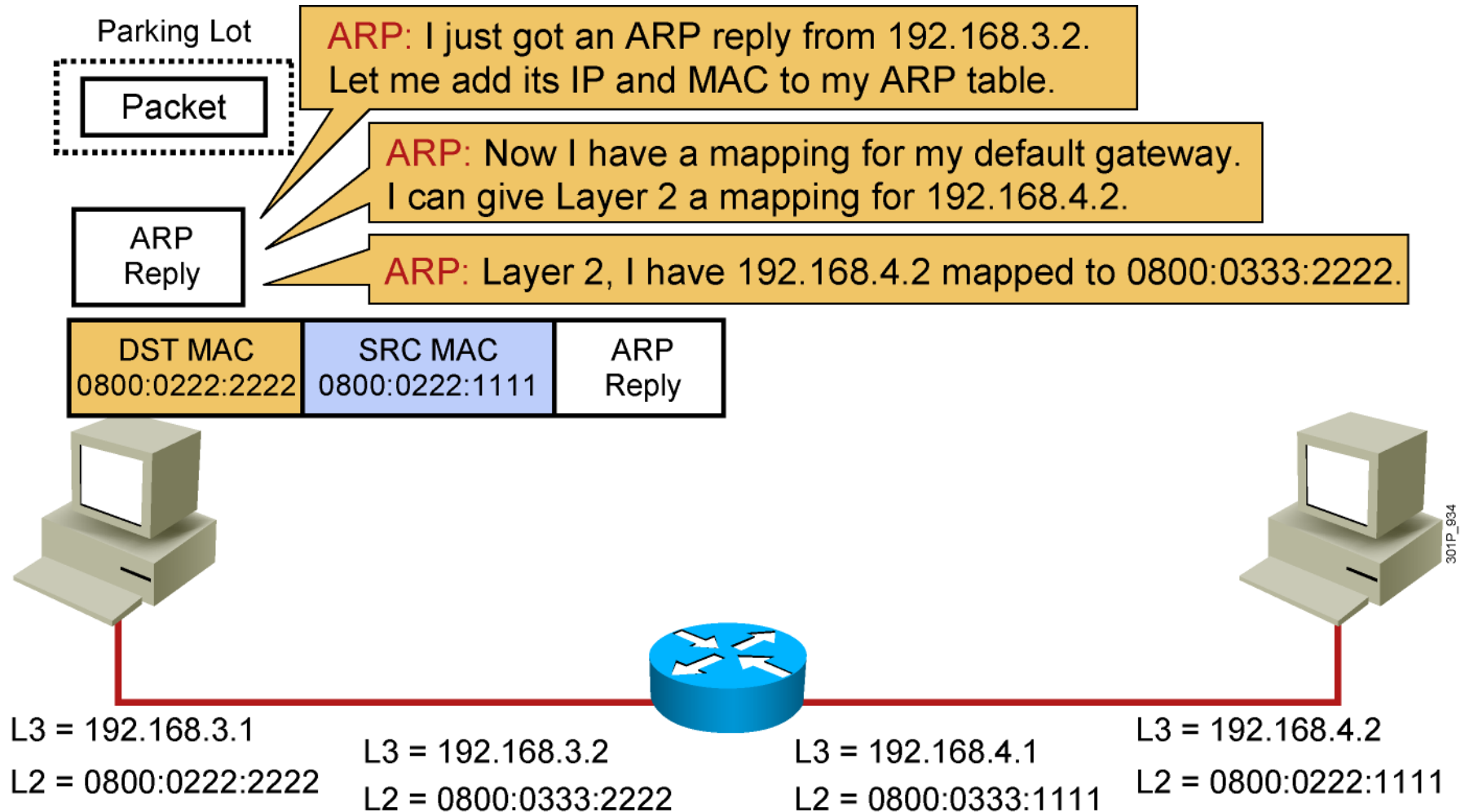
Host-to-Host Packet Delivery (6 of 17)



Host-to-Host Packet Delivery (7 of 17)

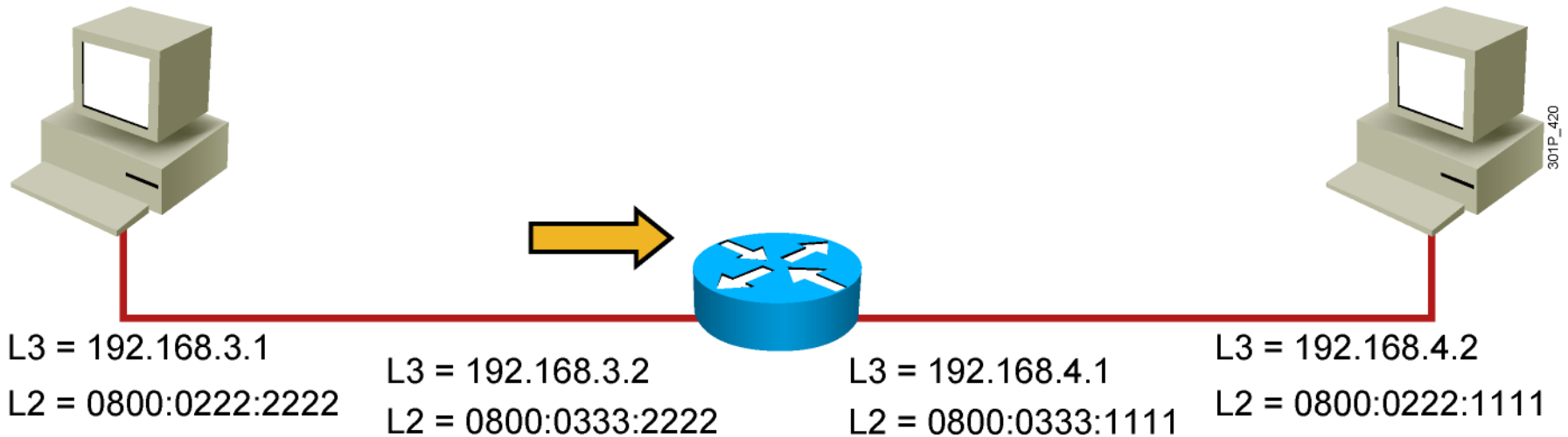


Host-to-Host Packet Delivery (8 of 17)



Host-to-Host Packet Delivery (9 of 17)

Layer 2: I can send out that pending frame.

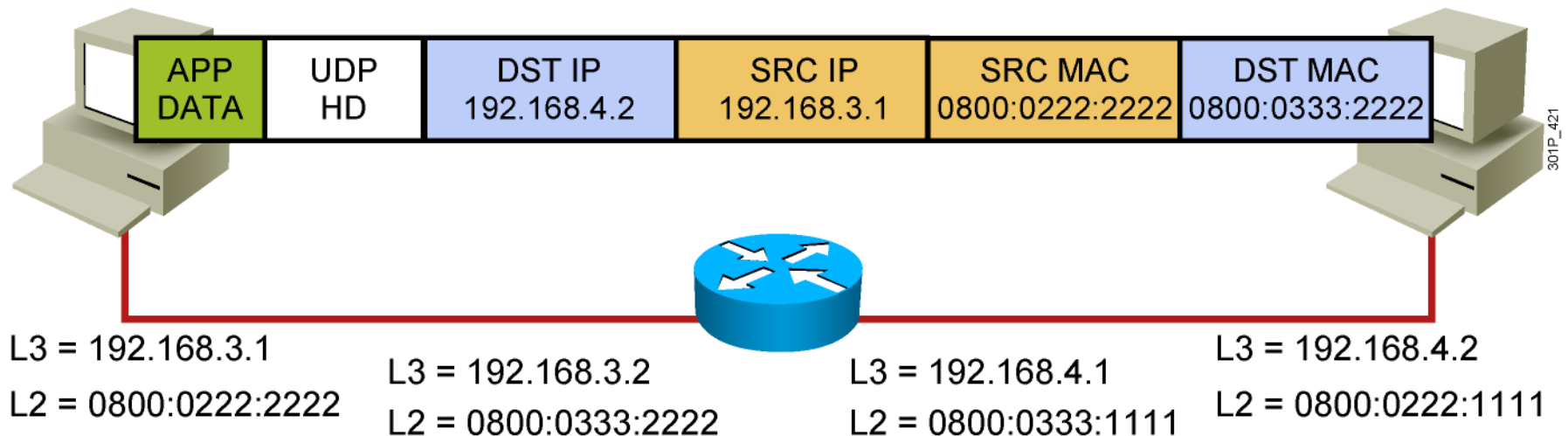
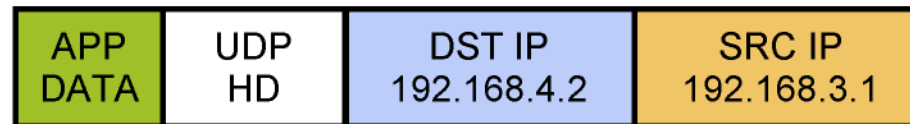


Host-to-Host Packet Delivery (10 of 17)

Router L2: I received a frame with my MAC address. I need to pass it to L3.

Router L3: This isn't my address. It needs to be routed.

Router L3: I need to forward this packet.

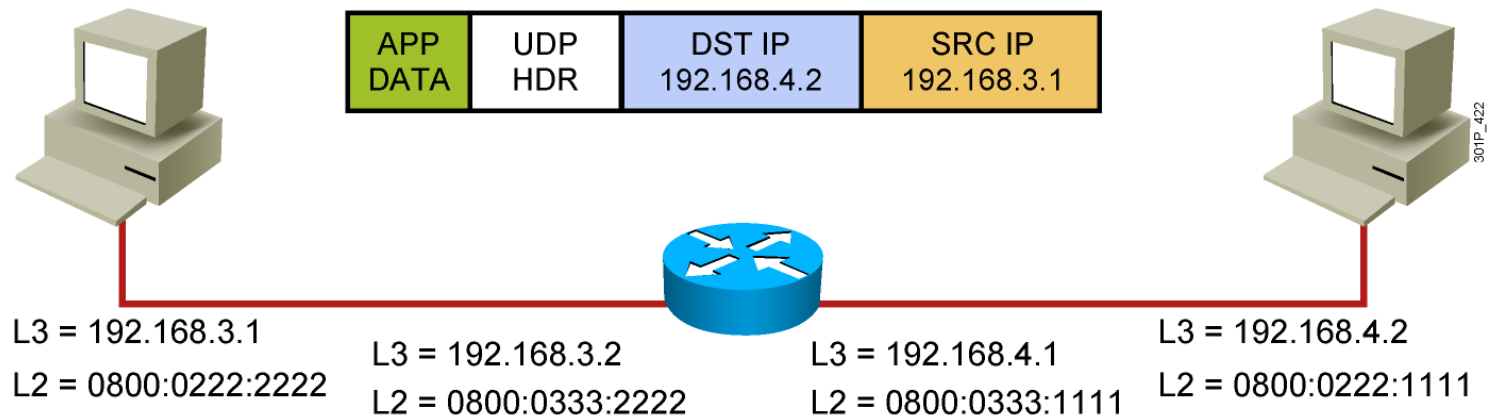


Host-to-Host Packet Delivery (11 of 17)

Destination	Next Hop	Interface
192.168.3.0/24	Connected	fa 0/0
192.168.4.0/24	Connected	fa 0/1

Router L3: I have an interface on the 192.168.4.0/24 segment. I can forward this packet directly to the host.

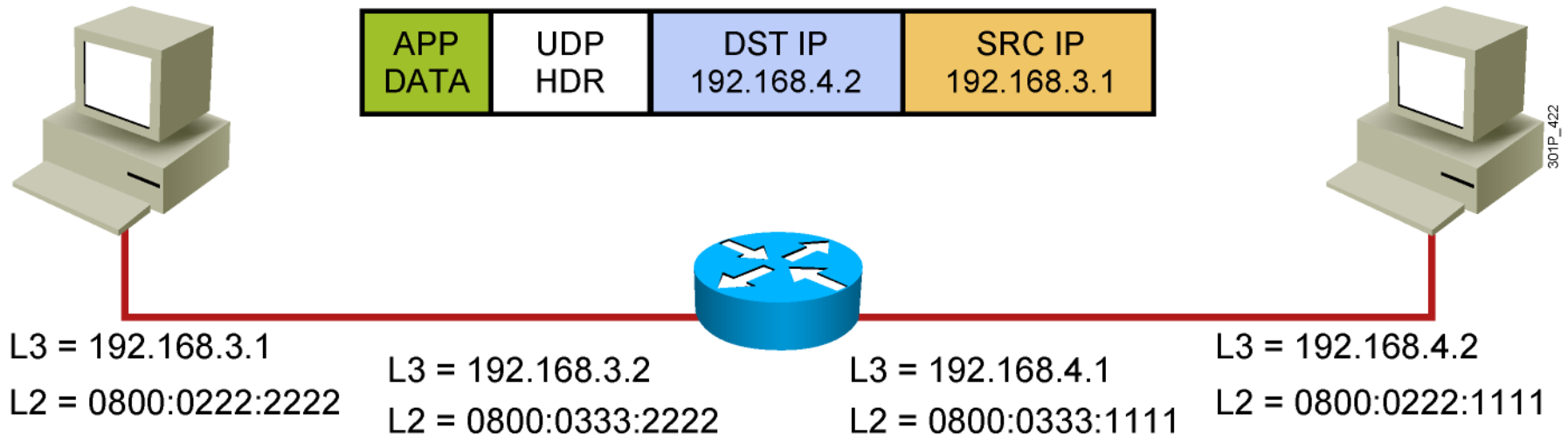
Router L3: L2, send this packet.



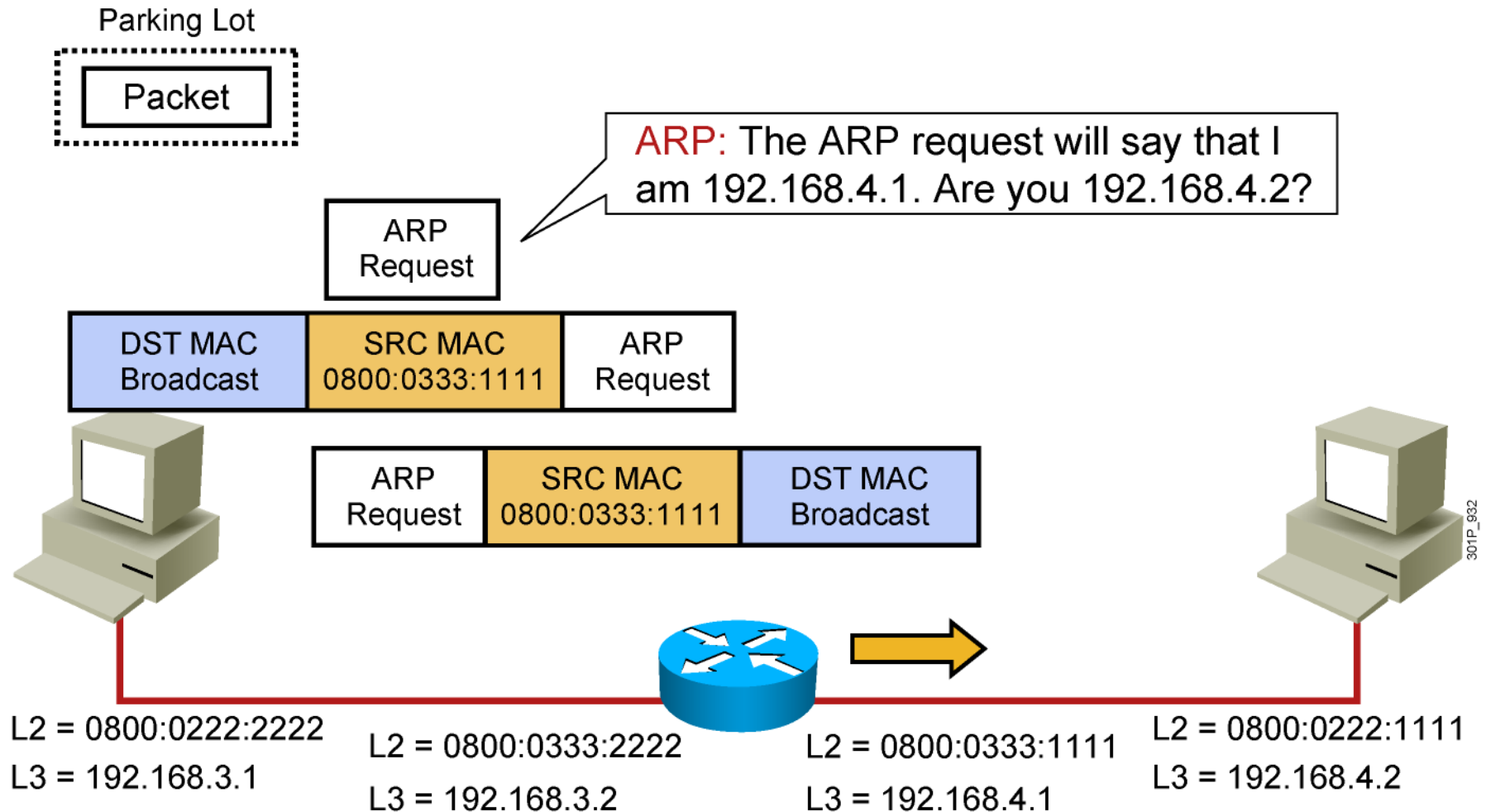
Host-to-Host Packet Delivery (12 of 17)

Router L3: I have an interface on the 192.168.4.0/24 segment. I can forward this packet directly to the host.

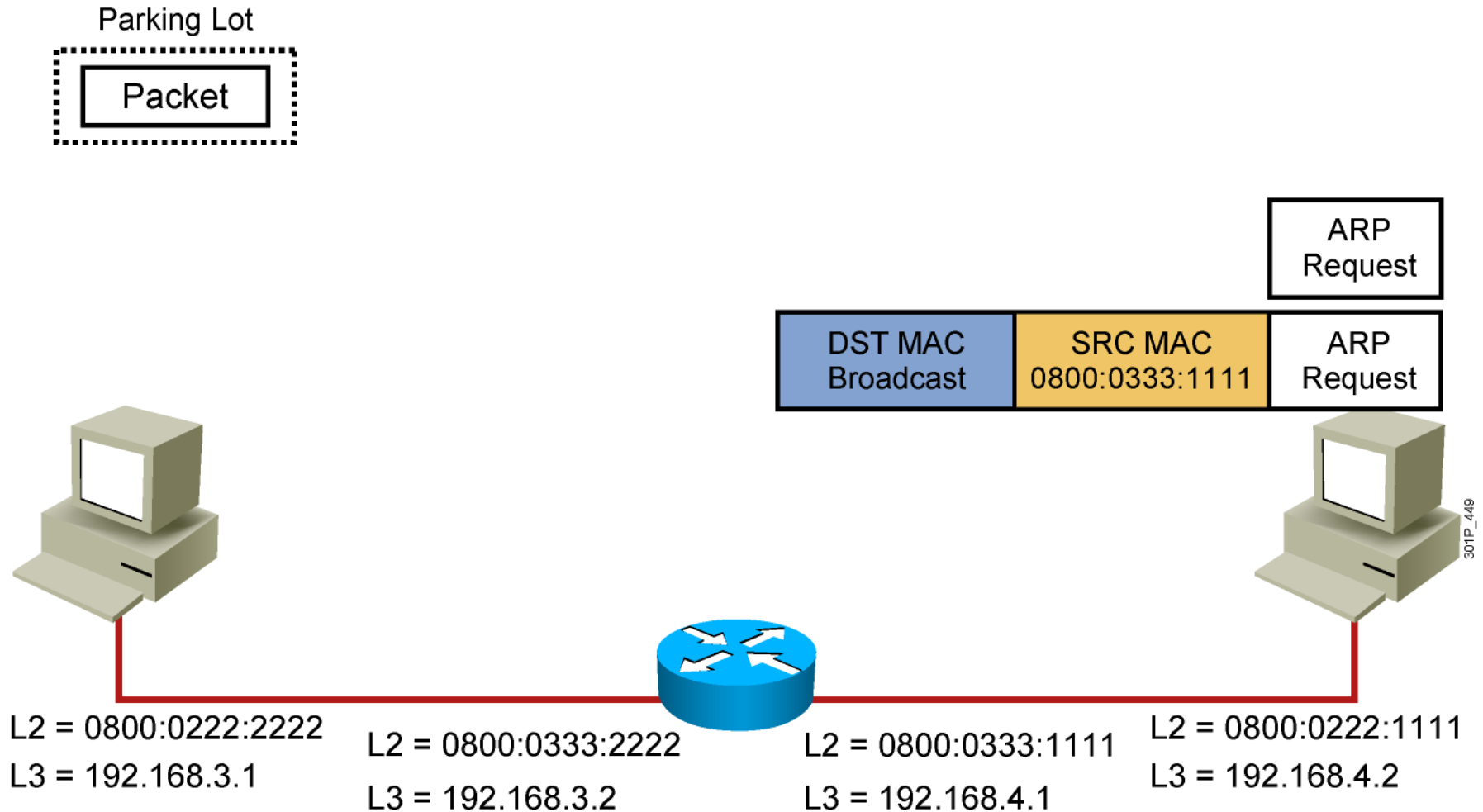
Router L3: L2, send this packet.



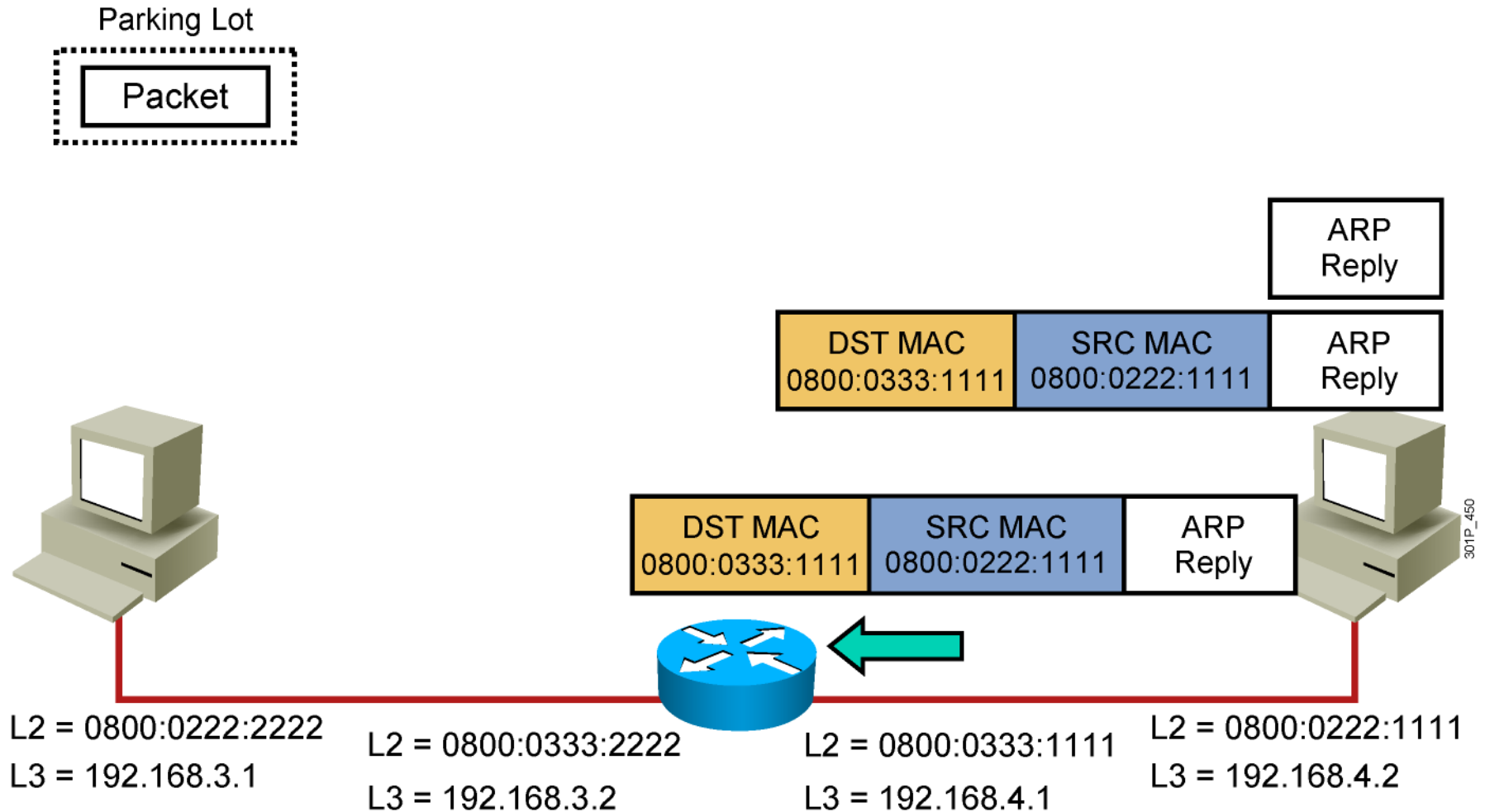
Host-to-Host Packet Delivery (13 of 17)



Host-to-Host Packet Delivery (14 of 17)



Host-to-Host Packet Delivery (15 of 17)



Host-to-Host Packet Delivery (16 of 17)

Router ARP: I just got an ARP reply from 192.168.4.2. Let me add its IP and MAC to my ARP table.

Router ARP: Now I have a mapping. I can give Layer 2 a mapping for 192.168.4.2.

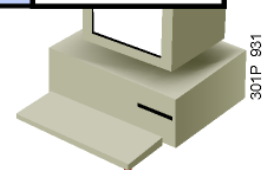
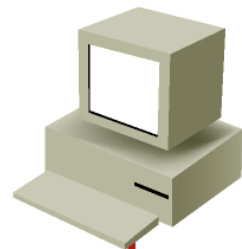
Router ARP: Layer 2, I have 192.168.4.2 mapped to 0800:0222:1111.

ARP Reply

DST MAC
0800:0333:1111

SRC MAC
0800:0222:1111

ARP Reply



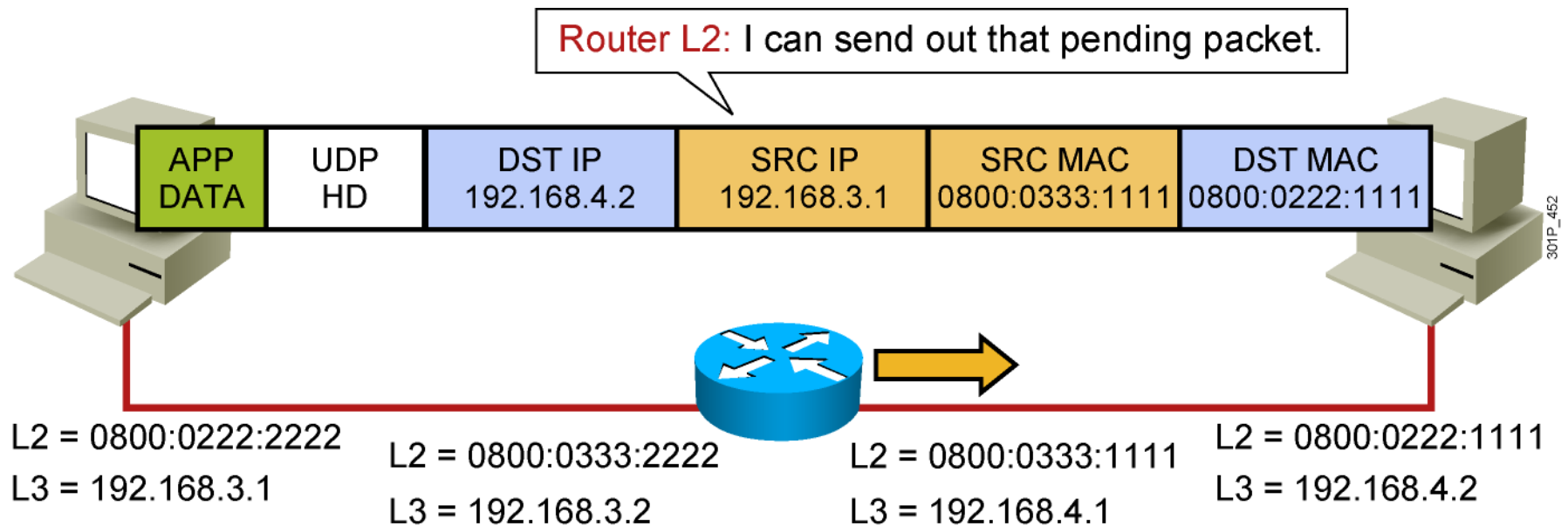
L2 = 0800:0222:2222
L3 = 192.168.3.1

L2 = 0800:0333:2222
L3 = 192.168.3.2

L2 = 0800:0333:1111
L3 = 192.168.4.1

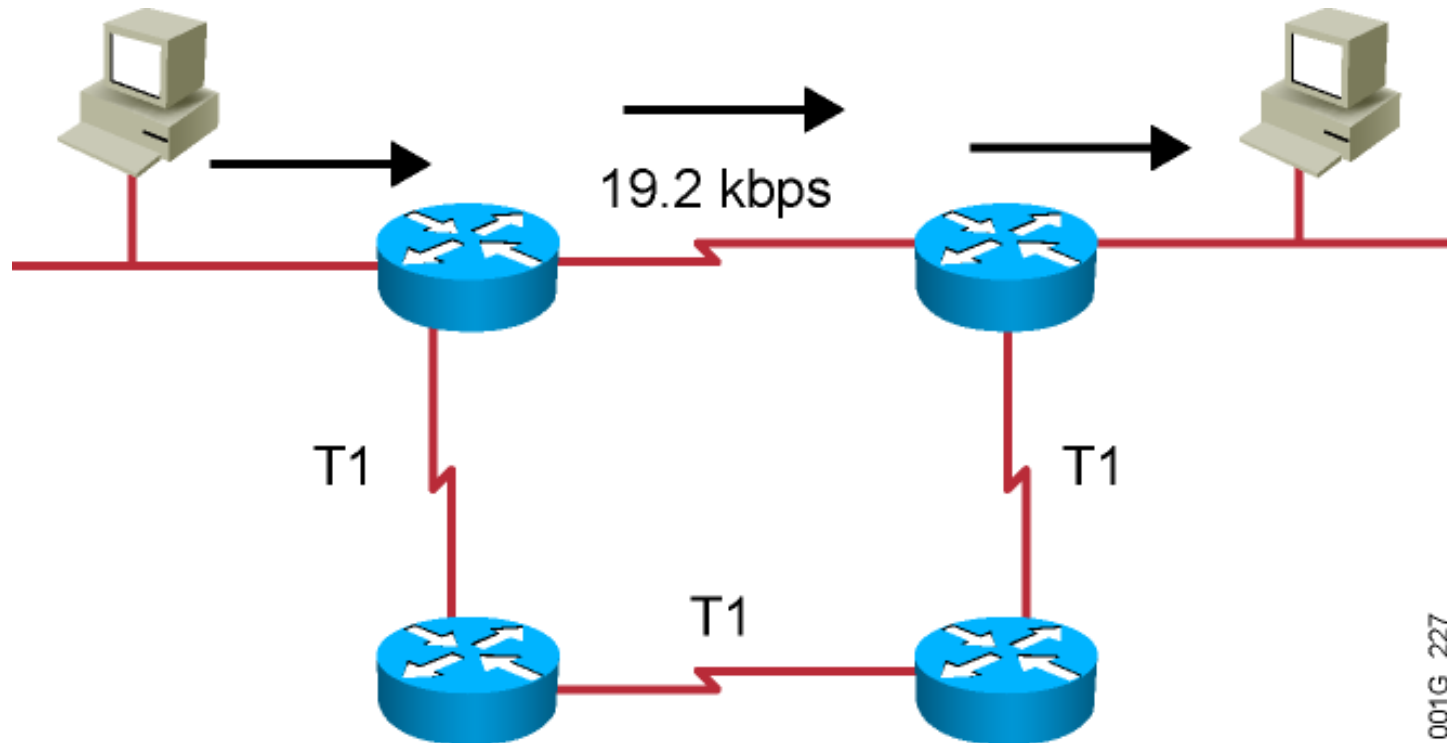
L2 = 0800:0222:1111
L3 = 192.168.4.2

Host-to-Host Packet Delivery (17 of 17)



Appendix B

RIP Overview



001G_227

- Maximum is 16 equal-cost paths (default = 4)
- Hop-count metric selects the path
- Routes update every 30 seconds



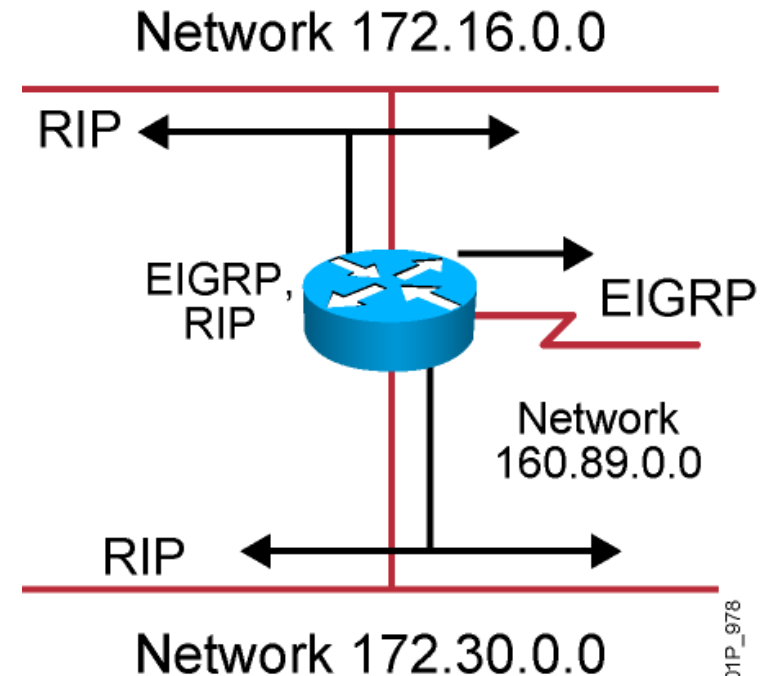
RIPv1 and RIPv2 Comparison

	RIPv1	RIPv2
Routing protocol	Classful	Classless
Supports variable-length subnet mask?	No	Yes
Sends the subnet mask along with the routing update?	No	Yes
Addressing type	Broadcast	Multicast
Defined in ...	RFC 1058	RFCs 1721, 1722, and 2453
Supports manual route summarization?	No	Yes
Authentication support?	No	Yes



IP Routing Configuration Tasks

- Router configuration
 - ***Select routing protocols***
 - ***Specify networks or interfaces***



301P_978



RIP Configuration

```
RouterX(config)# router rip
```

- Starts the RIP routing process

```
RouterX(config-router)# version 2
```

- Enables RIP version 2

```
RouterX(config-router)# network network-number
```

- Selects participating attached networks
- Requires a major classful network number



RIP Configuration Example

