



University of Colorado **Boulder**

Fundamentals of Data Communications

IP Routing

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Review

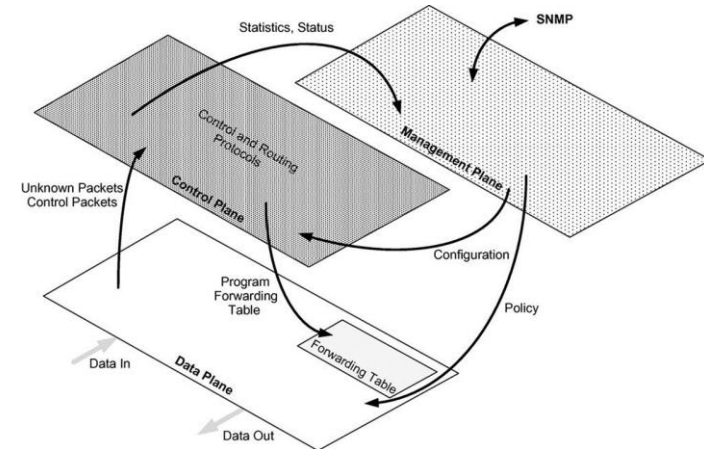
Static Routing

- **What is the purpose of routing?**
- **What is a static route?**
- **Static Routing vs Routing Protocols**
 - Manual
 - *Administrative Overhead*
 - *Scalability*
 - Static not dynamic
 - *(failover)*
- **Routing Table**

Router

- **Inside a router**

- Control plane: routing protocols
- Data plane: packet forwarding



- **Path selection**

- Minimum-hop/cost and shortest-path routing
- Algorithms: Link-state vs. Distance vector routing

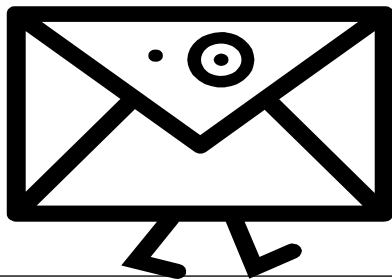
- **Topology change**

- Using beacons to detect topology changes
- Propagating topology information

What is Routing?

- **A famous quotation from RFC 791**

“A name indicates what we seek.
An address indicates where it is.
A route indicates how we get there.”
-- Jon Postel



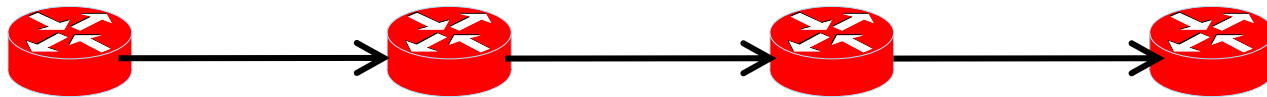
Routing vs. Forwarding

- **Routing: control plane**

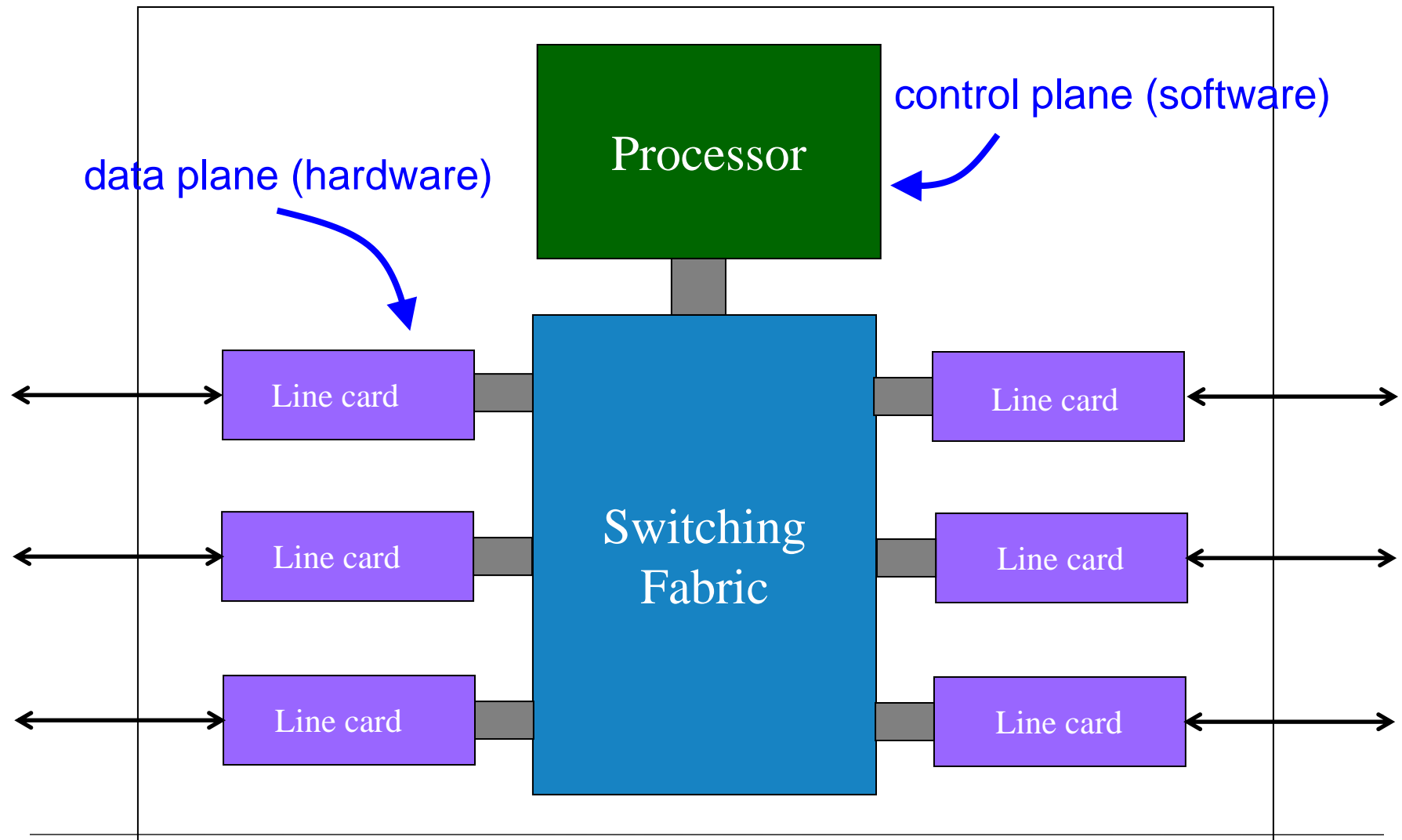
- Computing paths the packets will follow
- Routers talking amongst themselves
- Individual router *creating* a forwarding table

- **Forwarding: data plane**

- Directing a data packet to an outgoing link
- Individual router *using* a forwarding table



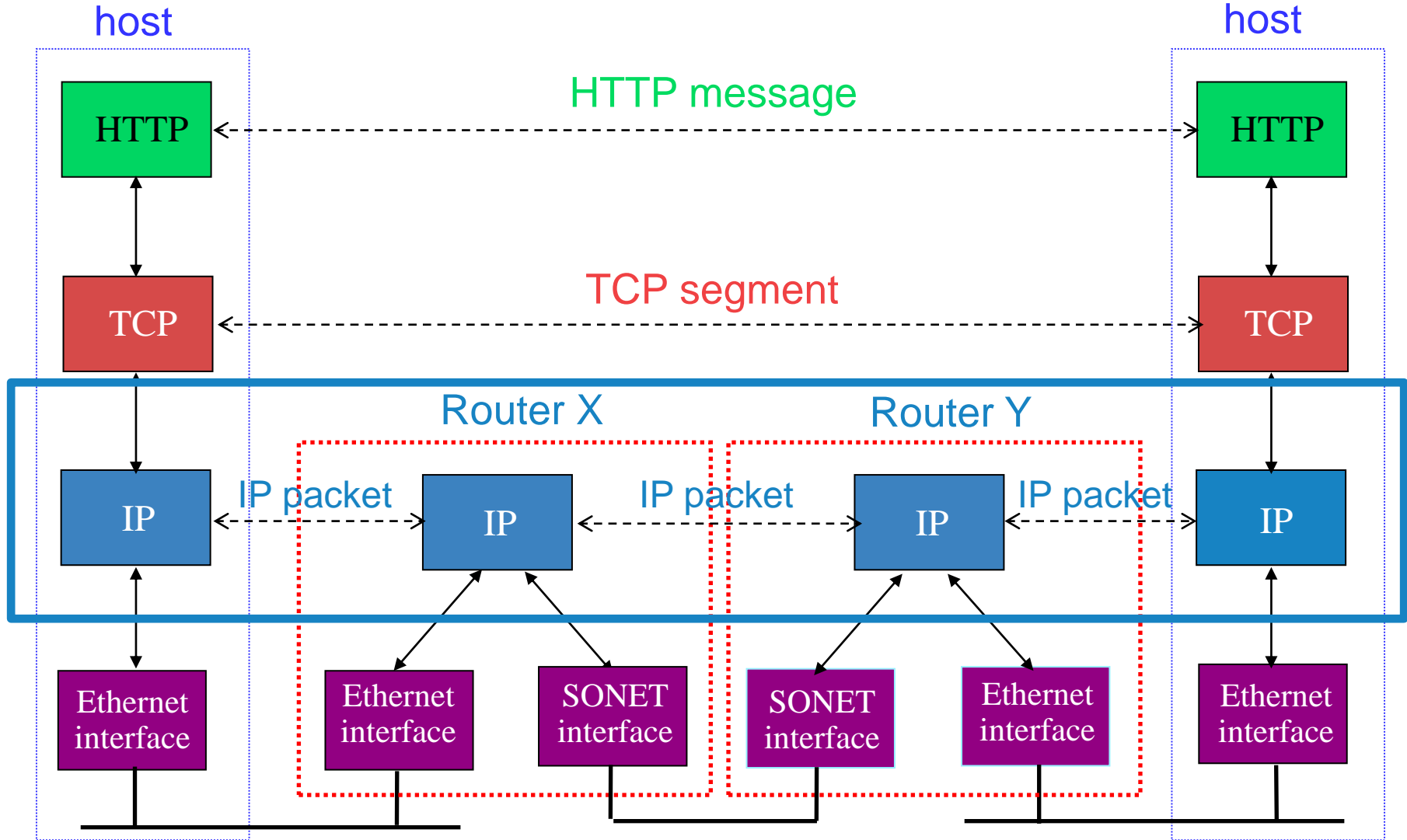
Data and Control Planes



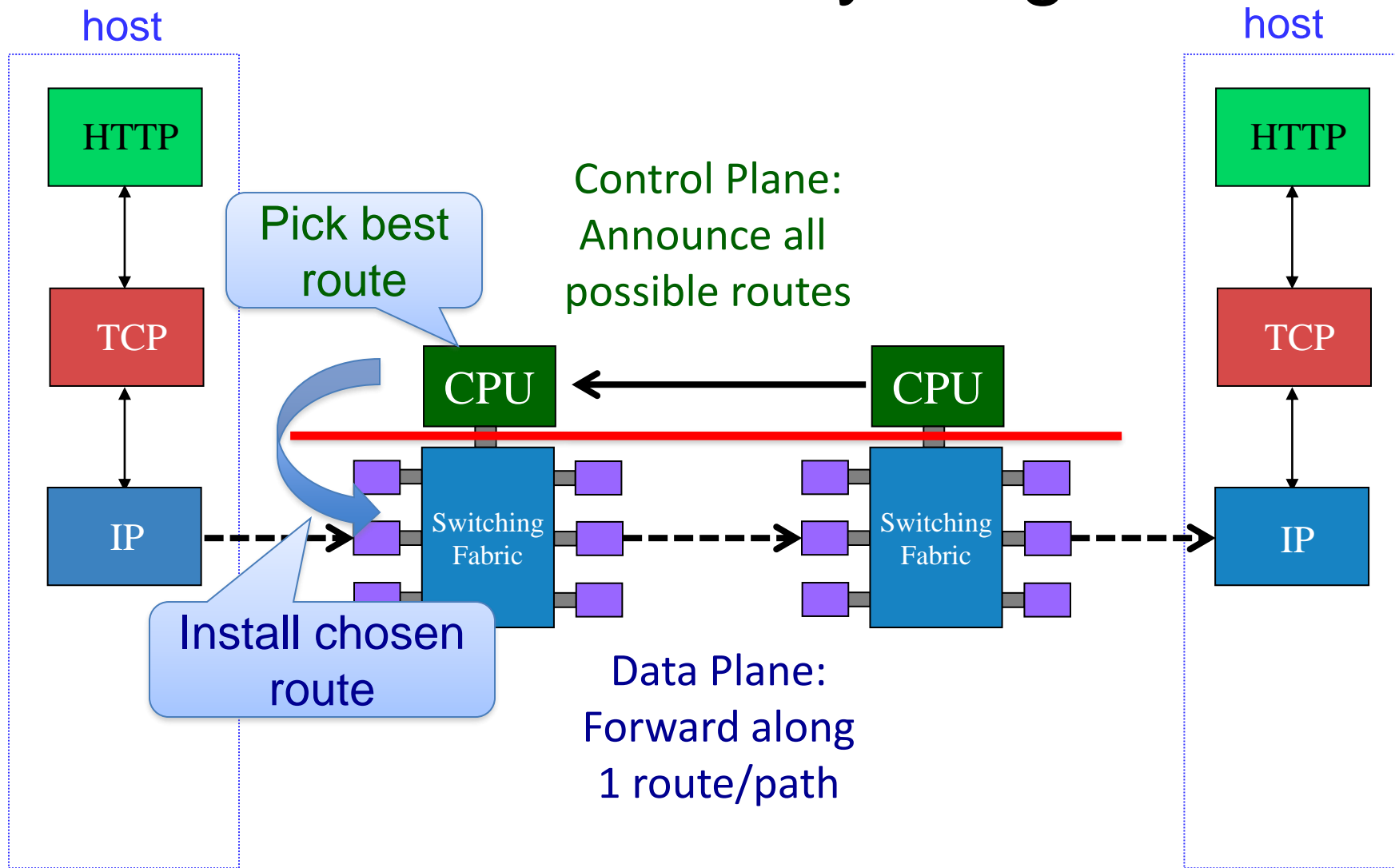
Where do Forwarding Tables Come From?

- **Routers have forwarding tables**
 - Map IP prefix (subnet) to outgoing link(s)
- **Entries can be statically configured**
 - E.g., “map 12.34.158.0/24 to fastethernet 0/0.1”
 - But, this doesn’t adapt
 - *To failures*
 - *To new equipment*
 - *To the need to balance load*
 - That is where routing protocols come in!

Recall the Internet Layering Model



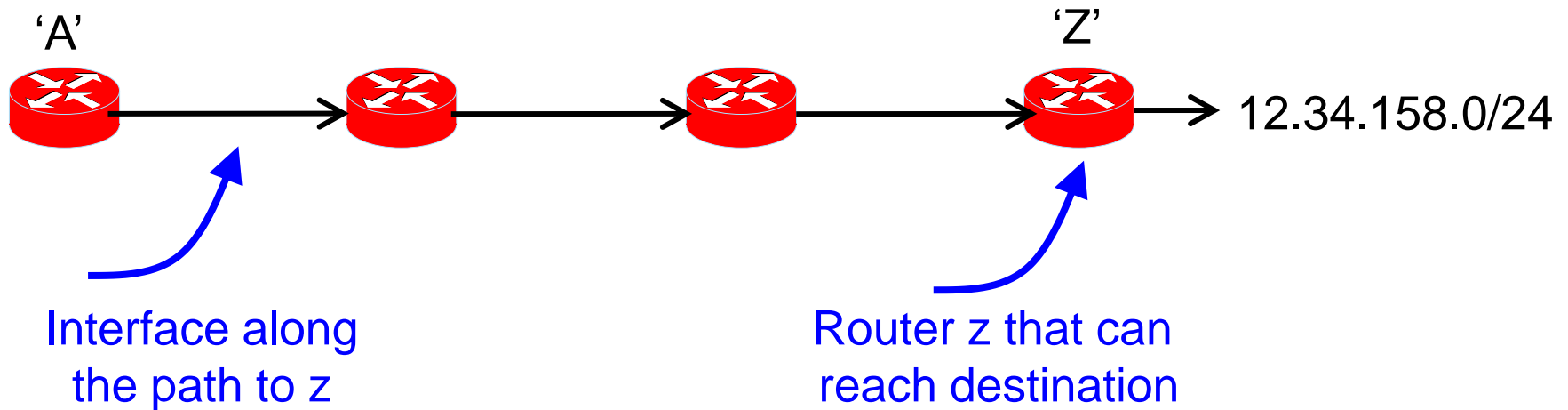
Recall the Internet Layering Model



Computing Paths Between Routers

- **Routers need to know two things:**

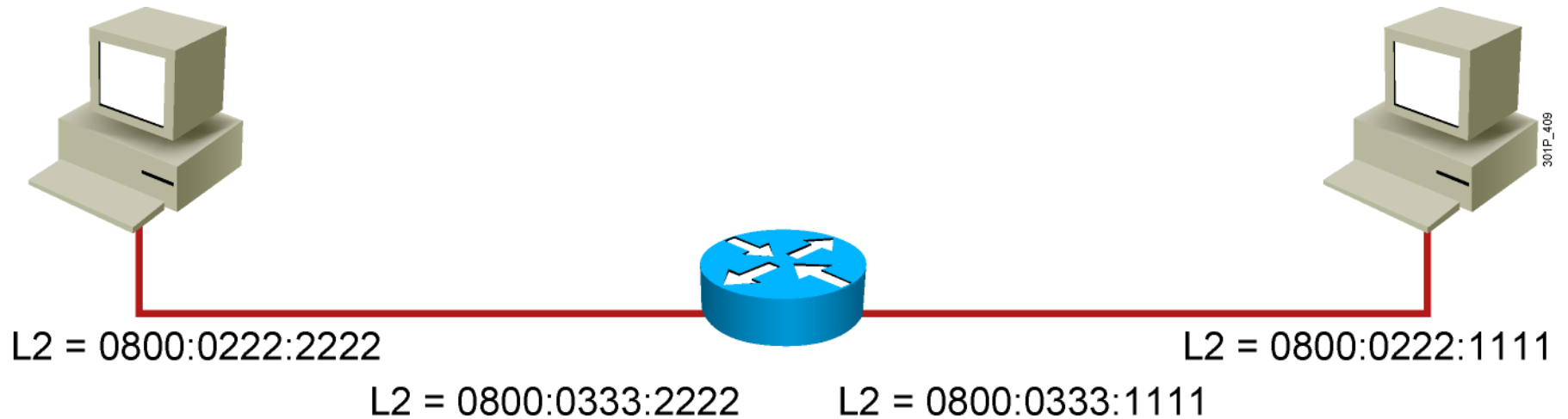
1. Which router to use to reach a destination prefix
2. Which outgoing interface to use to reach that router



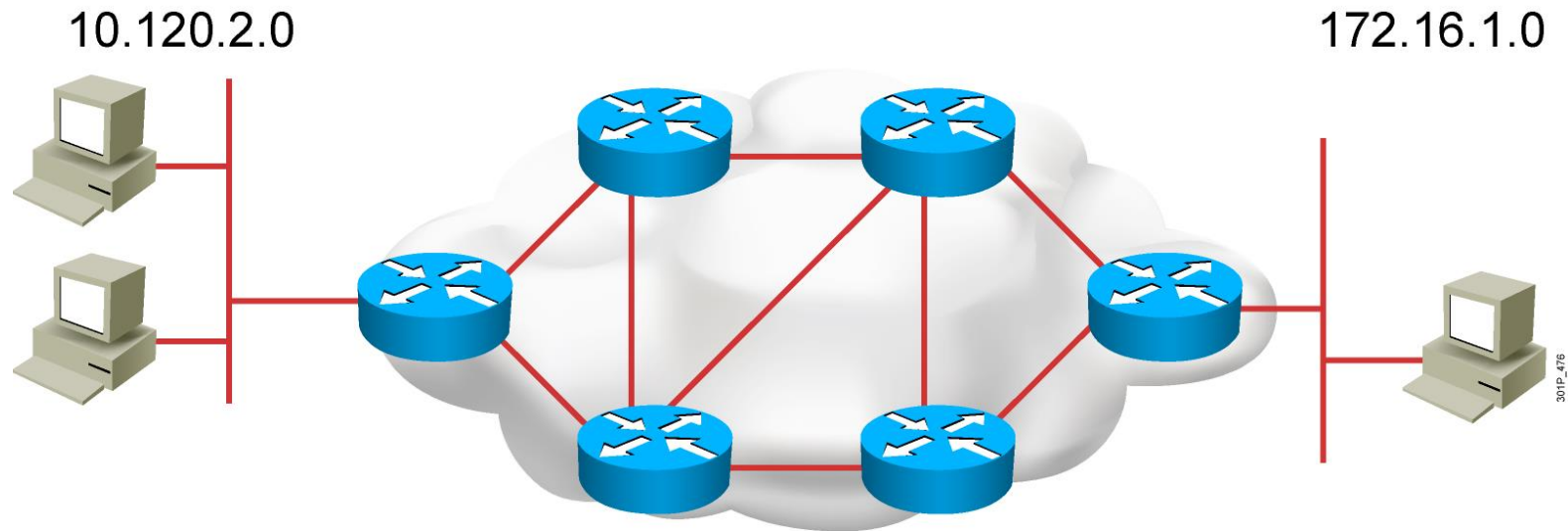
Routers and Routing

- **Routing**
 - The process by which data/traffic gets from one location to another
- **A router is the device used to route traffic**
- **Routers can forward packets over static routes or dynamic routes**
 - Based on the router configuration and network design
- **Static routes**
 - A route that a network administrator enters into the router manually
 - Unidirectional static routes must be configured **to** and **from** networks to allow bidirectional communications to occur
- **Dynamic routes**
 - A route that a network routing protocol adjusts automatically for topology or traffic changes

Layer 2 Addressing

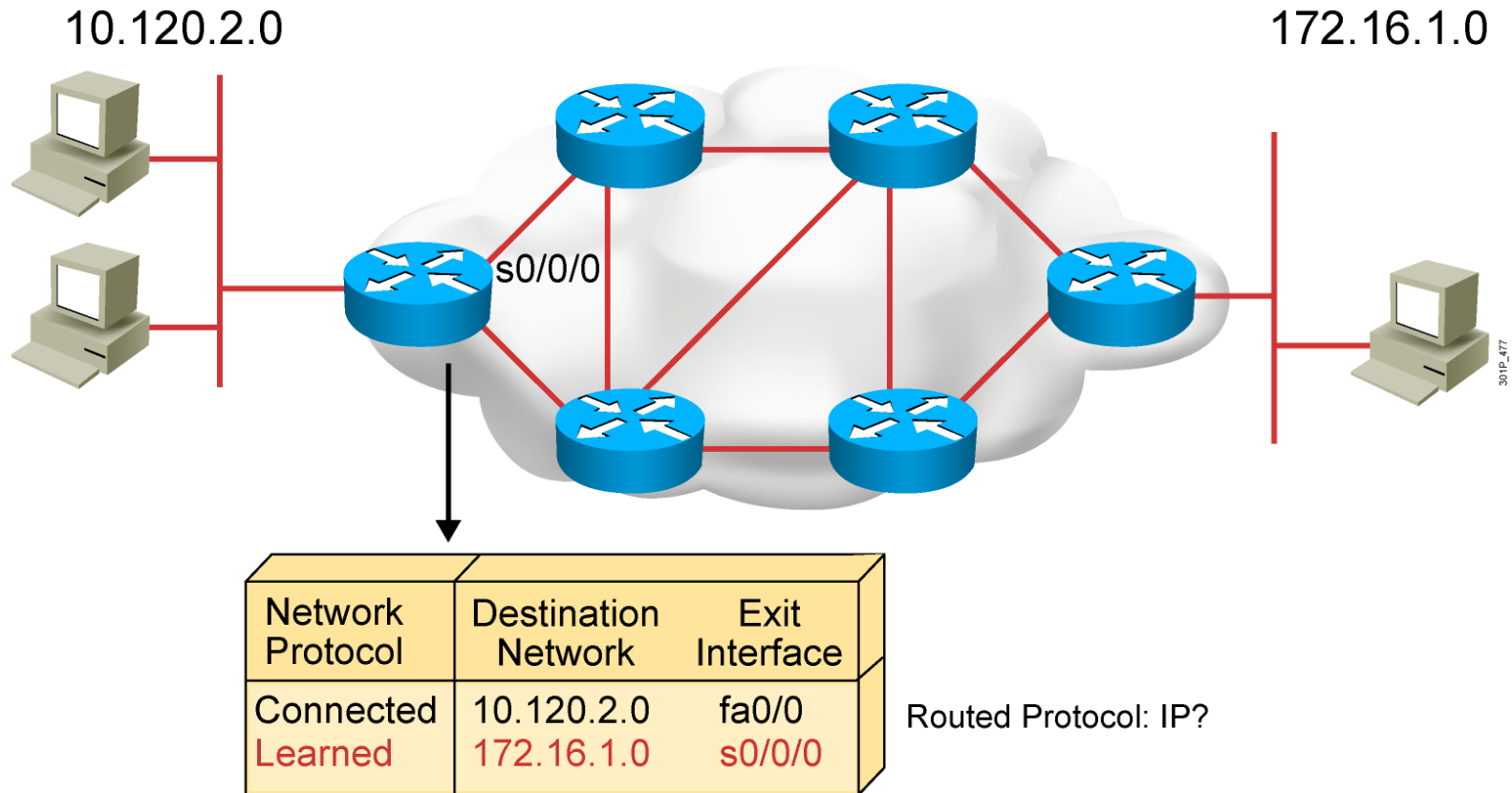


Router Operations



- **A router needs to do the following:**
 - Know the destination address
 - Identify the sources from which the router can learn
 - Discover possible routes to the intended destination
 - Select the best route
 - Maintain and verify routing information

Router Operations (Cont.)



- Routers must learn destinations that are not directly connected
 - Routers automatically learn connected routes (if the interface is up!)

Identifying Static and Dynamic Routes

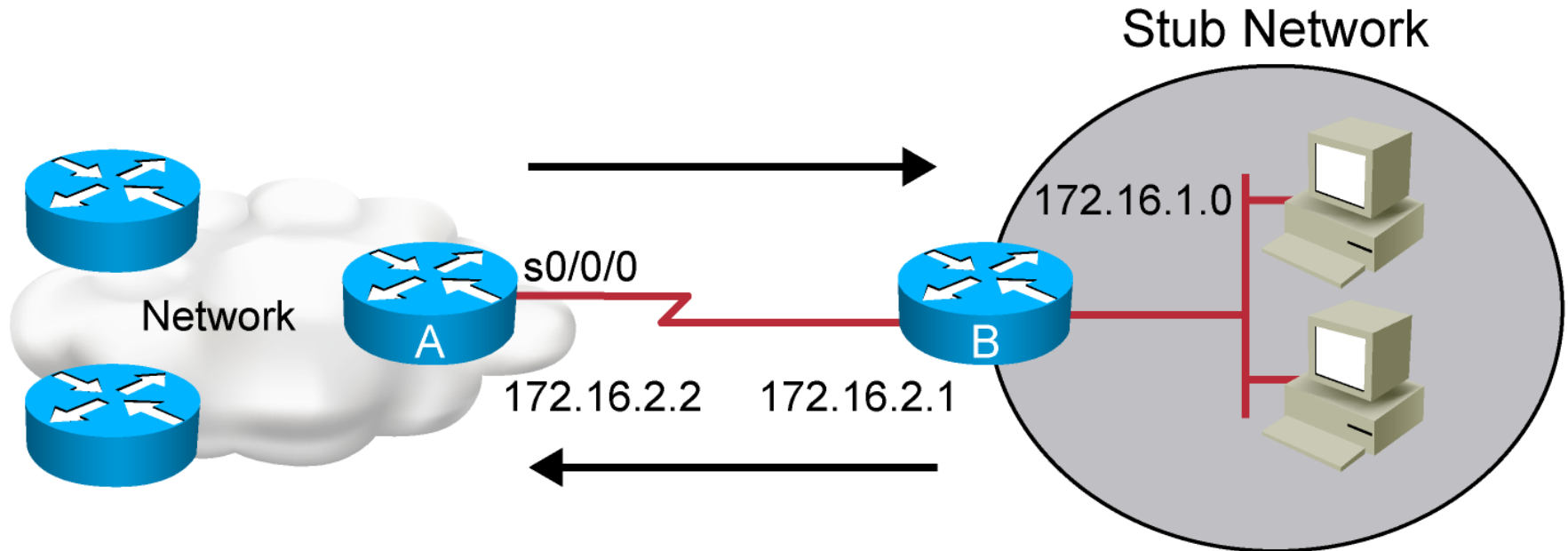
- **Static route**

- Uses a route that a network administrator enters into the router manually

- **Dynamic route**

- Uses a route that a network routing protocol adjusts automatically for topology or traffic changes

Static Routes



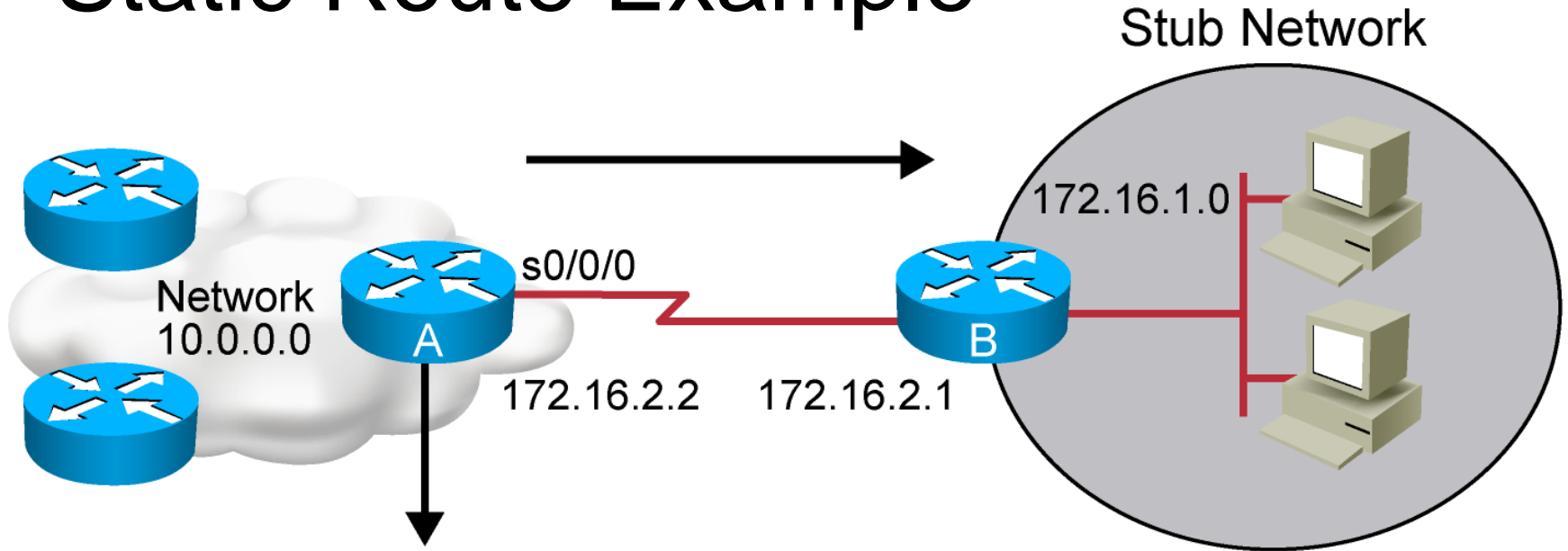
Configure unidirectional static routes **to** and **from** a network to allow communications to occur.

Static Route Configuration

```
(config)# ip route network [mask]  
{address | interface}[distance] [permanent]
```

- Defines a path to an IP destination network or subnet or host
- Address = IP address of the next hop router
- Interface = outbound interface of the local router
- Example:
 - ***ip route 192.168.1.0 255.255.255.0 192.168.222.254***

Static Route Example



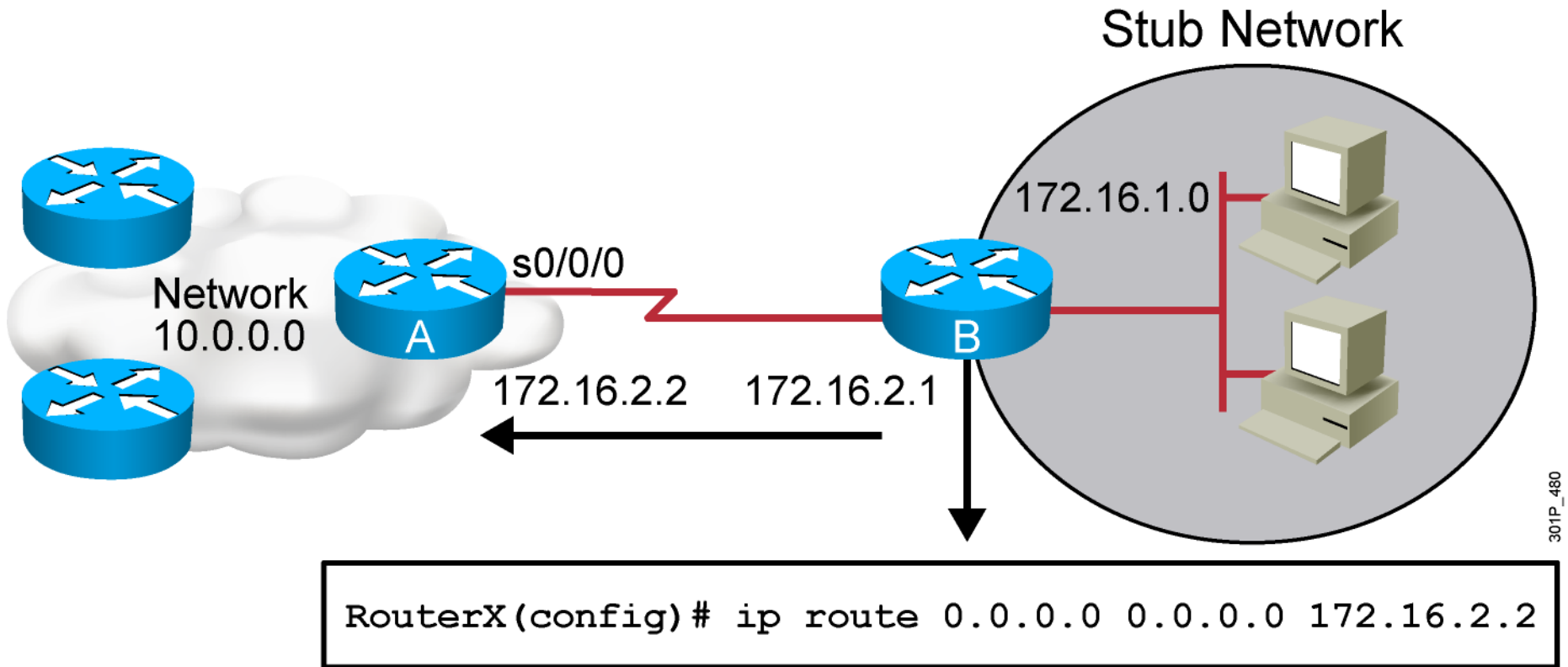
```
(config)# ip route 172.16.1.0 255.255.255.0 172.16.2.1
```

or

```
(config)#ip route 172.16.1.0 255.255.255.0 s0/0/0
```

- This is a unidirectional route.
- You must have a route configured in the opposite direction for traffic to return.

Default Routes



- This route allows the “stub network” (behind RouterB) to reach all networks beyond Router A.

Verifying the Static Route Configuration

- **Route Table**

```
RouterX# show ip route
```

```
Codes: C - connected, S - static, I - IGRP, R - RIP, M - mobile, B - BGP  
        D - EIGRP, EX - EIGRP external, O - OSPF, IA - OSPF inter area  
        E1 - OSPF external type 1, E2 - OSPF external type 2, E - EGP  
        i - IS-IS, L1 - IS-IS level-1, L2 - IS-IS level-2, * - candidate default  
        U - per-user static route
```

```
Gateway of last resort is 0.0.0.0 to network 0.0.0.0
```

```
    10.0.0.0/8 is subnetted, 1 subnets
```

```
C        10.1.1.0 is directly connected, Serial0/0/0
```

```
S*    0.0.0.0/0 is directly connected, Serial0
```

P2P with Internet

Floating Static Routes

Dynamic Routing

Overview and Terminology

- Think beyond directly connected devices
- “Network Map” and Routing tables
- Static Routes (type the Internet)
 - *Nearly 1 million routes!*
- **Routing protocols:** Propagate information about available networks
- **Routed protocols:** L3 protocols that can be routed (IP)

Overview and Terminology

- ***Routing type***: Link State vs. Distance Vector vs. Path Vector
- ***Exterior Routing Protocols (EGP)***: Inter-corporate route distribution
- ***Interior Routing Protocols (IGP)***: Intra-corporate route exchange
- ***Autonomous System (AS)***: Network under single management control (ISP)

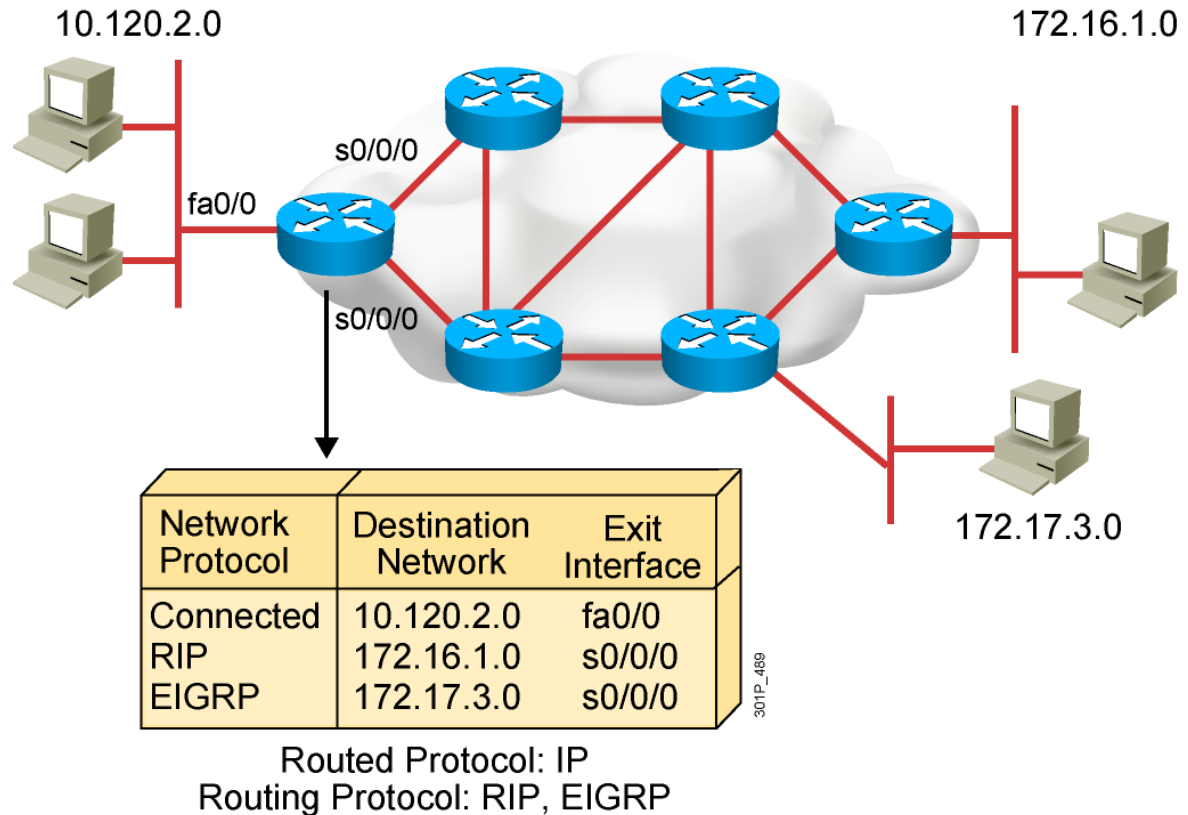
Routing Objectives

- Build a routing table (learn & propagate)
- Pick best routes (if more than one available)
- Remove invalid routes
- Replace routes if better advertisement received
- Restore routes fast (convergence time)
- Prevent loops



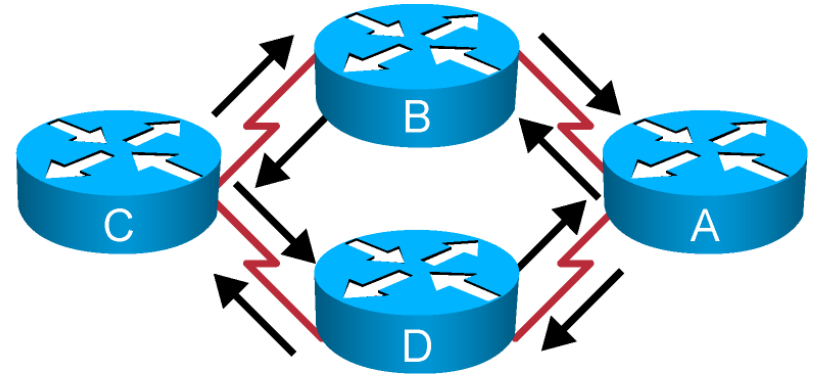
What is a Routing Protocol?

- **Routing** protocols are used between routers to determine paths and maintain routing tables.
- After the path is determined, a router can route a **routed** protocol.

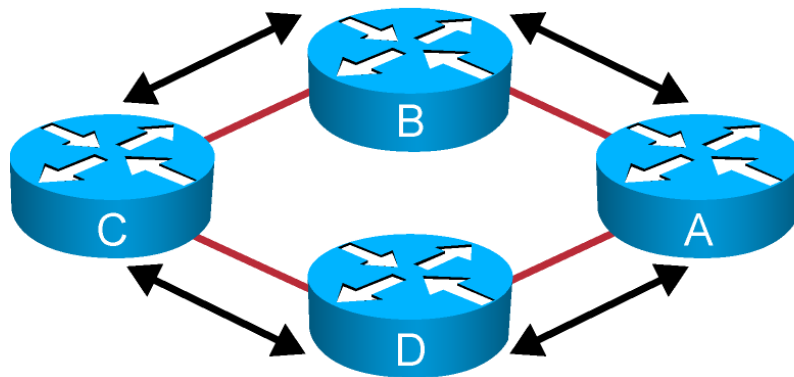


Classes of Routing Protocols

Distance Vector



Hybrid Routing



Link-State

Interior Gateway Routing Protocols (IGP)

- Several available in the market. Protocols have evolved over time to respond to changing network conditions
- Considerations while choosing a protocol:
 - **Type of routing protocol**
 - *Distance vector*
 - *Link-state*
 - *Hybrid*
 - **Update process**
 - *Full updates*
 - *Partial updates*
 - **Convergence time**
 - **Metric**
 - *Measure of link quality / preference*
 - **VLSM support**
 - *Subnet size variations*
 - *Permit better address allocation*
 - **Classless vs. Classful**
 - *TX or Not TX mask info*
 - *If TX then VLSM supported*

Distance Vector Protocols

- **Distance Vector Logic**

- Add directly connected
- Send updates out include directly connected and learned routes
 - *“Routing by Rumor”*
- Listen for routing updates
- Routing info: subnet & metric
- Use broadcast or multicast for updates
- Chose best of multiple routes
- Send/expect periodic full updates
- If updates no longer received, remove routes learned from such neighbor
- Assume that the advertising router is the next hop for a route

Distance Vector Protocols

- **RIPv1**

- Metric: hop count (smaller better) Infinite metric = 16
- Full updates every 30s (subnet # & metric)
- Convergence 3-5 minutes (depending on size of routing table)
- Classful (no VLSM)
- Broadcasts updates (255.255.255.255)
- “Route-poisoning” and split horizon
 - *Infinite metric of 16*
 - *Prevent advertising back to the interface it was received*

- **RIPv2**

- Same features as RIPv1
- Adds: VLSM support (updates include subnet #, mask & metric)
- MD5 and plain text authentication
- Includes next hop router IP on updates
- Uses external route tags (redistribution)
- Multicast routing updates (224.0.0.9)

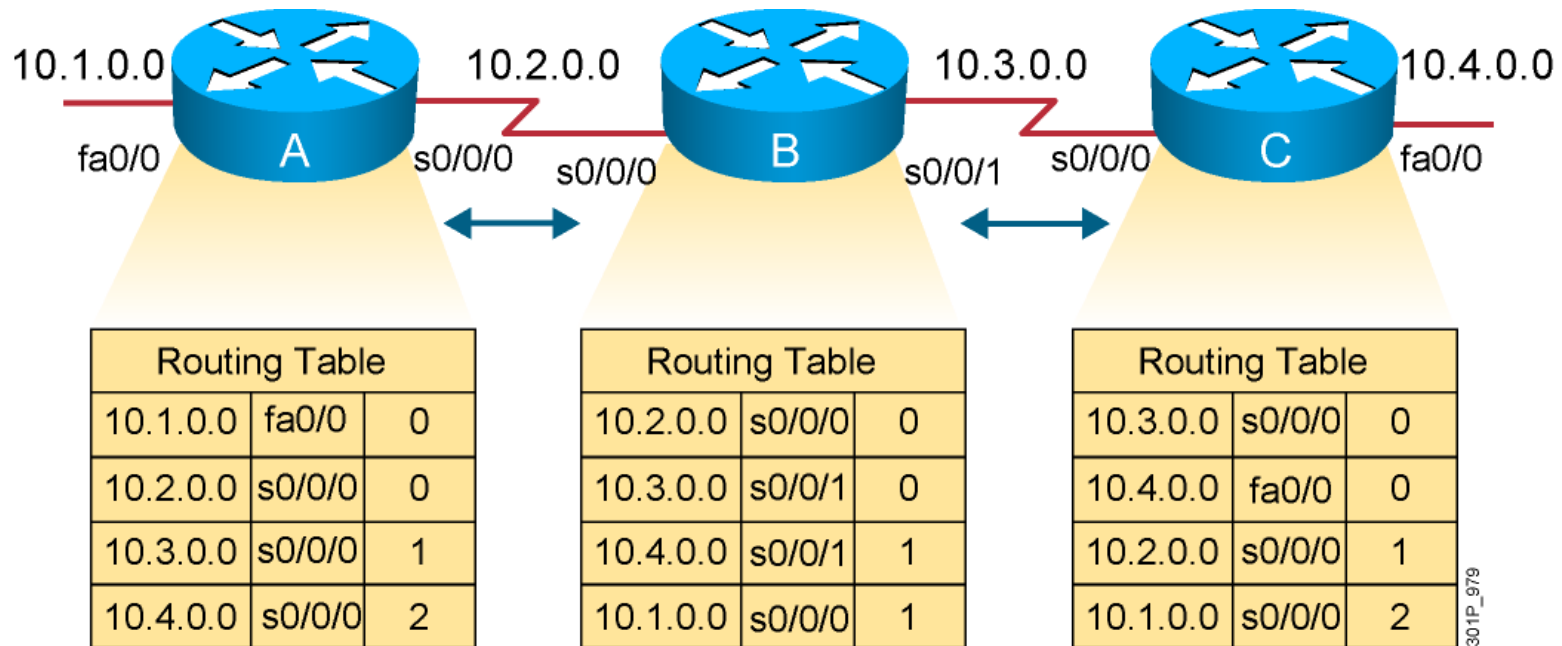


Distance Vector Protocols

- **IGRP**

- Proprietary
- Metric: bandwidth + delay (default), reliability, load, MTU
- Metric value (1 - 4 billion), smaller better.
- Metrics are cumulative: multi-hop path adds links delays
- Metrics are configurable
- Full updates every 90s
- No VLSM supported
- Infinite metric = 4,294,967,295 (smaller better)

Sources of Information and Discovering Routes



- Routers discover the best path to destinations from each neighbor.



Link State Protocols

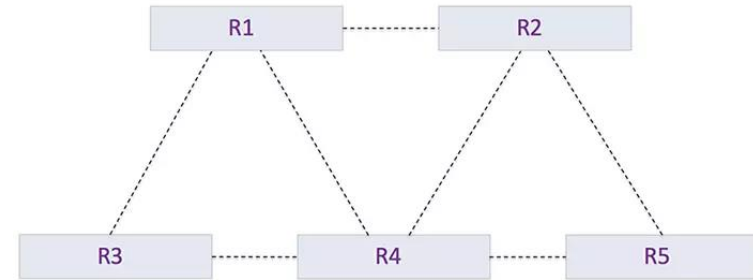
- **Link State logic**

- Add directly connected
- Don't transmit routing info until neighbors are discovered
- Routing info: topological information about the network, at the end every router has a complete map of the network.
- Use reliable protocol for route updates
- Calculate shortest path first (SPF) algorithm (Dijkstra) to determine best routes and next hop (also prevents loops)
- Full updates at start and after long periods of time (~30 min.)
- Partial updates when a link fails
- Fast convergence

Link State Protocols

- **OSPF**

- Most popular
- Discover neighbors, then exchange routes
- Reliable Transport Protocol
- Run SPF and store best routes
- More memory required, more processing
- Metric: “cost” based on bandwidth (smaller-better)
- Averages 10s convergence time
- VLSM supported
- Hello messages to confirm neighbor reachability
- Full updates every 30 min
- Partial updates when link fails



| Interface type | bandwidth | Cost |
|--------------------------|---------------------|------|
| Fast Ethernet and faster | 100 Mb/s and higher | 1 |
| Ethernet | 10 Mb/s | 10 |
| E1 | 2 Mb/s | 48 |
| T1 | 1.544 Mb/s | 64 |
| 128bps | 128bps | 781 |
| 64kbps | 64kbps | 1562 |
| 56kbps | 56kbps | 1785 |

Hybrid Protocols

- **EIGRP**

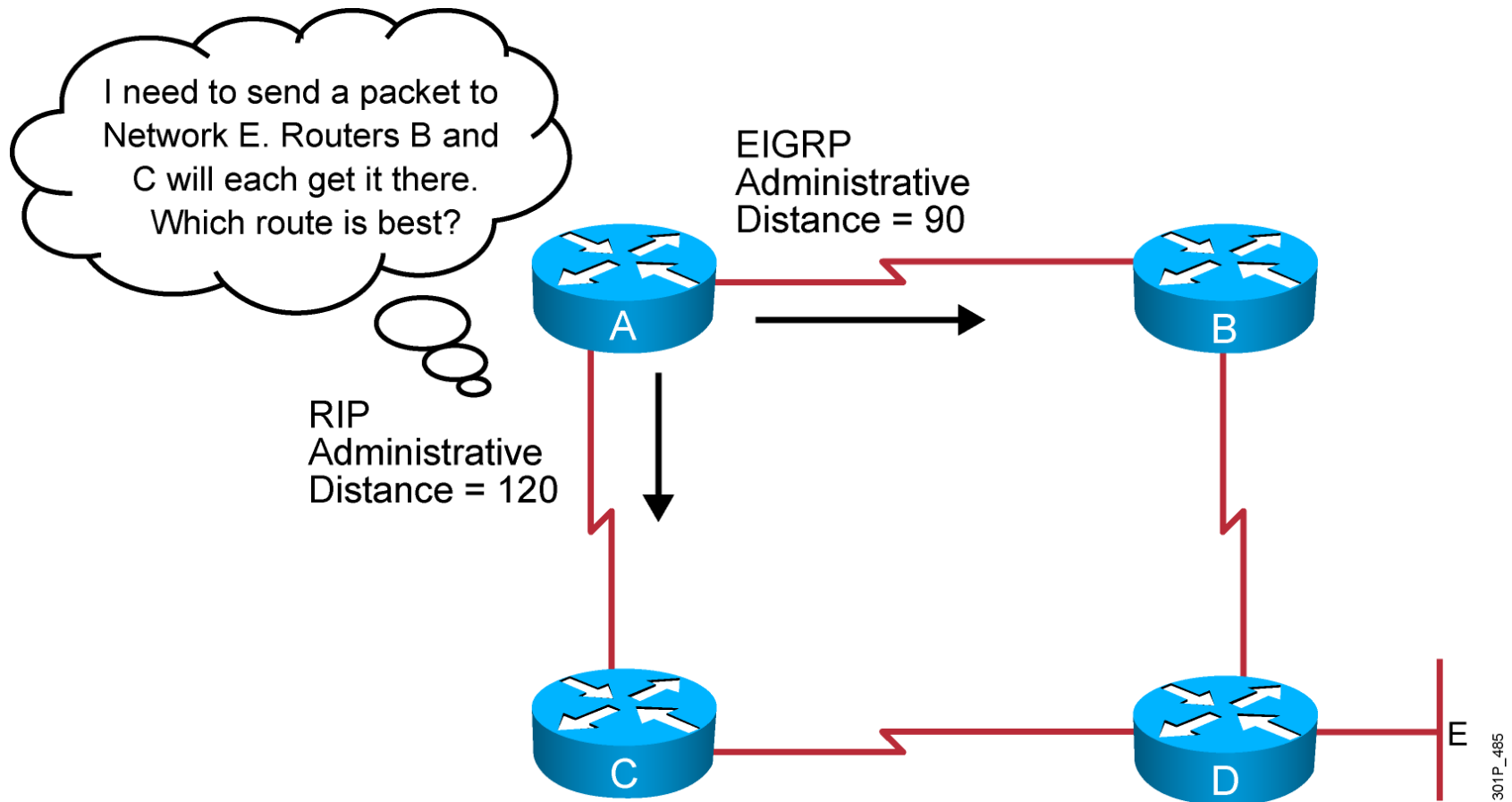
- Cisco Proprietary*
- Uses features from both link-state and distance vector protocols (balanced hybrid protocol: topology + routing tables)
- Diffused Update Algorithm (DUAL): exchange more info than distance vector but less than link-state
- Demands less computation
- Calculated best and “feasible successor routes” (alternatives), permits **immediate convergence**.
- Discovers neighbors like OSPF
- Metric based on bandwidth and delay (Idem IGRP * 256)
- Does not send periodic full updates

Loop Prevention in Distance Vector Protocols

- **Problems:**

- Multiple routes to same network
- Miscommunications on a single link
- Information loops through alternative paths
- Counting to infinity

Administrative Distance: Ranking Routes



Administrative Distance

- Applies to similar routes received from different protocols
- Prioritizes different routing protocols (which protocol do I believe first?)
- **The lower the AD number, the better**
- Examples:
 - ***Directly connected*** **0**
 - ***Static Routes*** **1**
 - ***eBGP*** **20**
 - ***EIGRP (internal)*** **90**
 - ***IGRP*** **100**
 - ***OSPF*** **110**
 - ***IS-IS*** **115**
 - ***RIP*** **120**
 - ***EIGRP (external)*** **170**
 - ***iBGP*** **200**

Classful Routing Protocol

- Classful routing protocols do not include the subnet mask with the route advertisement.
- Within the same network, consistency of the subnet masks is assumed.
- Summary routes are exchanged between foreign networks.
- These are examples of classful routing protocols:
 - ***RIPv1***
 - ***IGRP***

Classless Routing Protocol

- Classless routing protocols include the subnet mask with the route advertisement.
- Classless routing protocols support a variable-length subnet mask (VLSM). (aka subnetting)
- Summary routes can be manually controlled within the network.
- These are examples of classless routing protocols:
 - ***RIPv2***
 - ***EIGRP***
 - ***OSPF***
 - ***IS-IS***

Link-State Routing can be Problematic

- **Topology information is flooded**
 - High bandwidth and storage overhead
 - Forces nodes to divulge sensitive information
- **Entire path computed locally per node**
 - High processing overhead in a large network
- **Minimizes some notion of total distance**
 - Works only if policy is shared and uniform
- **Typically used only inside an AS**
 - E.g., OSPF and IS-IS

Distance Vector is on the Right Track

- **Advantages**

- Hides details of the network topology
- Nodes determine only “next hop” toward the destination

- **Disadvantages**

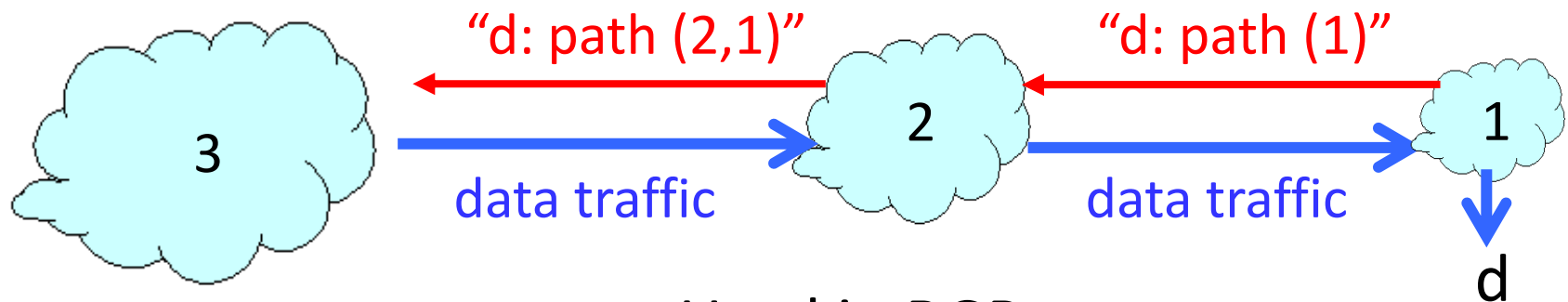
- Minimizes some notion of total distance, which is difficult in an interdomain setting
- Slow convergence due to the counting-to-infinity problem (“bad news travels slowly”)

- **Idea: extend the notion of a distance vector**

- To make it easier to detect loops

Path-Vector Routing

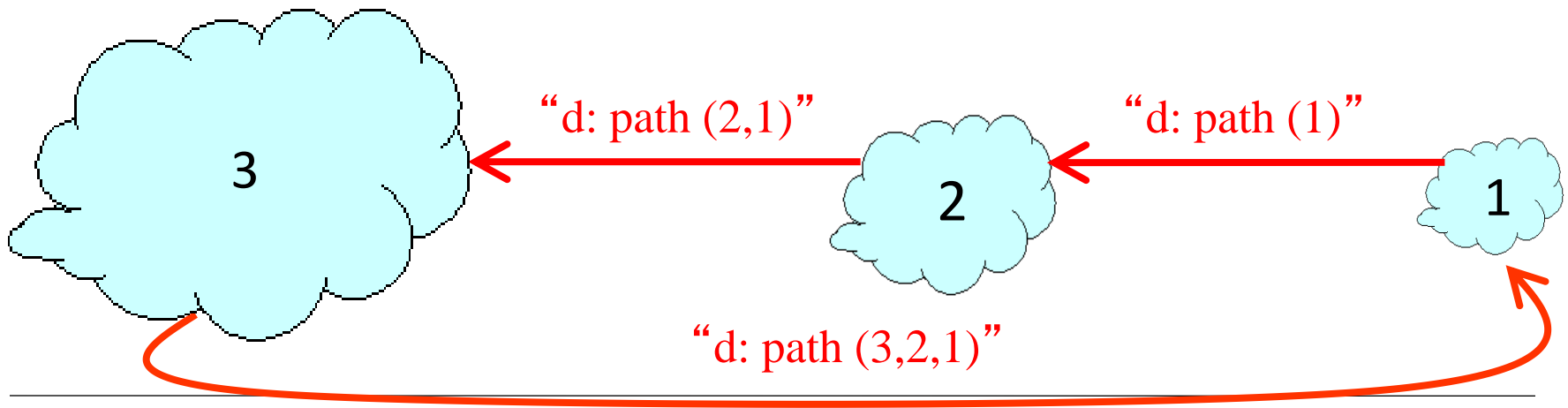
- **Extension of distance-vector routing**
 - Support flexible routing policies
- **Key idea: advertise the entire path**
 - Distance vector: send *distance metric* per dest d
 - Path vector: send the *entire path* for each dest d



Used in BGP

Faster Loop Detection

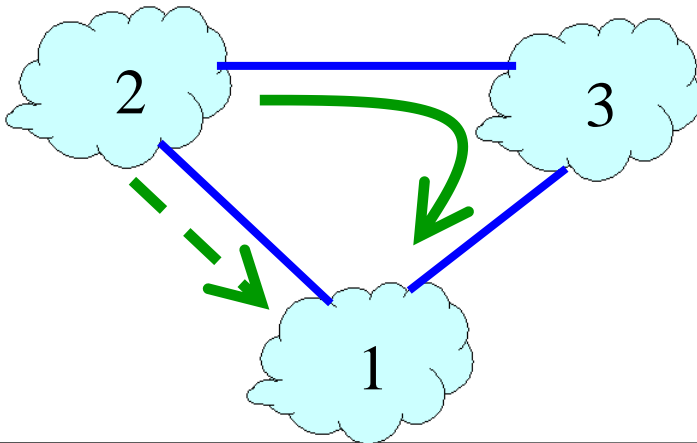
- **Node can easily detect a loop**
 - Look for its own node identifier in the path
 - E.g., node 1 sees itself in the path “3, 2, 1”
- **Node can simply discard paths with loops**
 - E.g., node 1 simply discards the advertisement



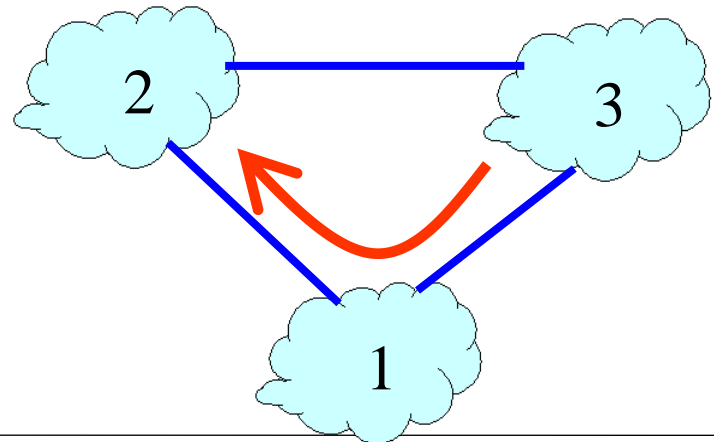
Path-Vector: Flexible Policies

- **Each node can apply local policies**
 - Path selection: Which path to use?
 - Path export: Which paths to advertise?

Node 2 prefers
“2, 3, 1” over “2, 1”

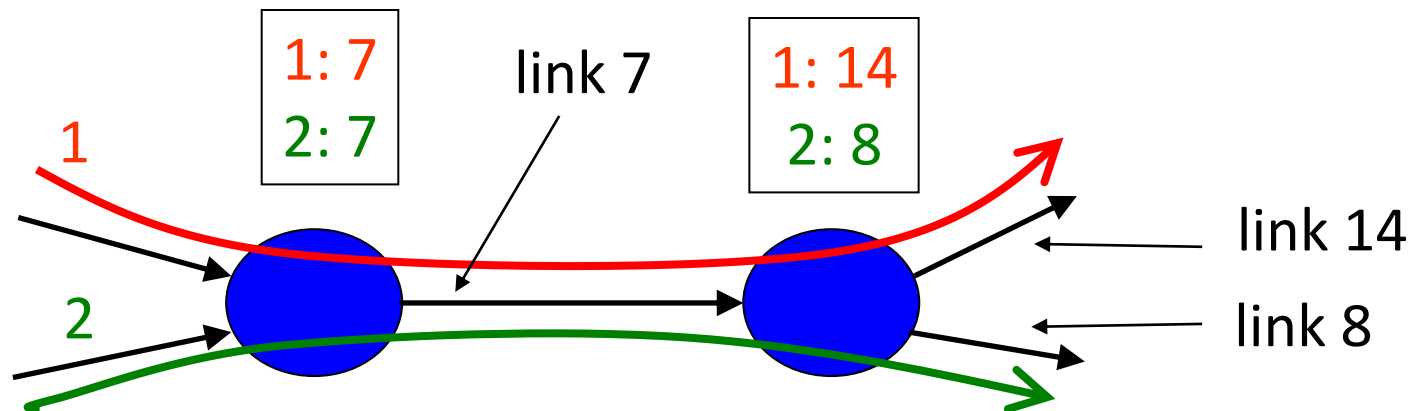


Node 1 doesn't let 3
hear the path “1, 2”



End-to-End Signaling

- **Establish end-to-end path in advance**
 - Learn the topology (as in link-state routing)
 - End host or router computes and signals a path
 - **Signaling:** *install entry for each circuit at each hop*
 - **Forwarding:** *look up the circuit id in the table*



Used in MPLS with RSVP

Distance Vector vs. Path Vector

- **Distance-vector routing**

- Pro: Less information and computation than link state
- Con: Slower convergence (e.g., count to infinity)

- **Path-vector routing**

- Share entire path, not distance: faster convergence
- More flexibility in selecting paths

- **Different goals / metrics if inter- or intra-domain**



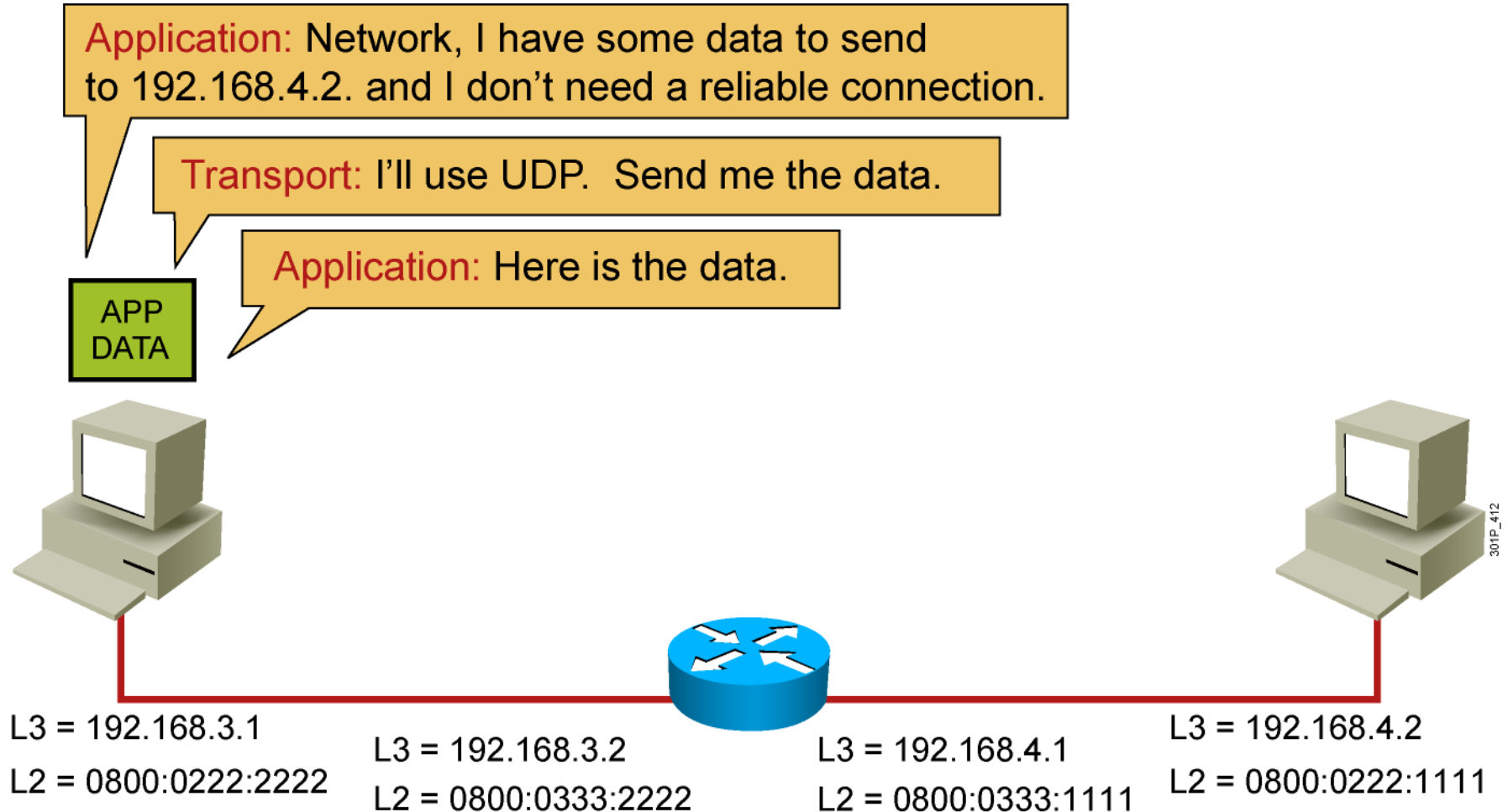
Questions?



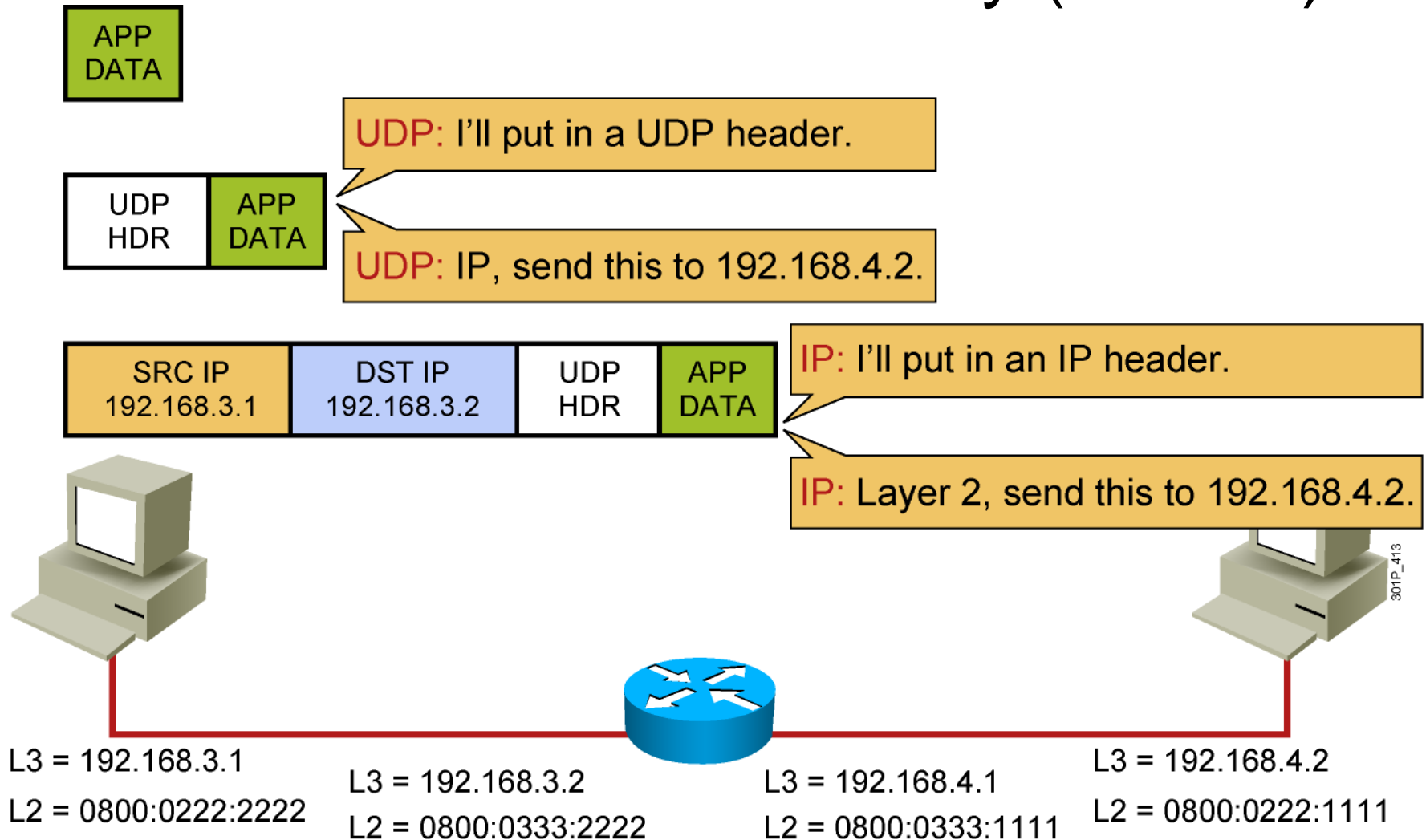
Lab

Appendix A

Host-to-Host Packet Delivery (1 of 17)



Host-to-Host Packet Delivery (2 of 17)

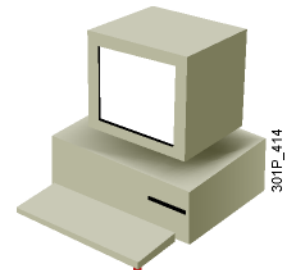
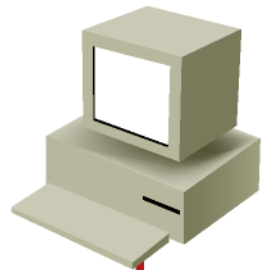


Host-to-Host Packet Delivery (3 of 17)

Layer 2: ARP, do you have a mapping for 192.168.4.2?

ARP: No, Layer 2 will have to hold the packet while I resolve the addressing.

| | | | |
|-----------------------|-----------------------|------------|-------------|
| SRC IP 192.168.3.1 | DST IP 192.168.4.2 | UDP HDR | APP DATA |
|-----------------------|-----------------------|------------|-------------|



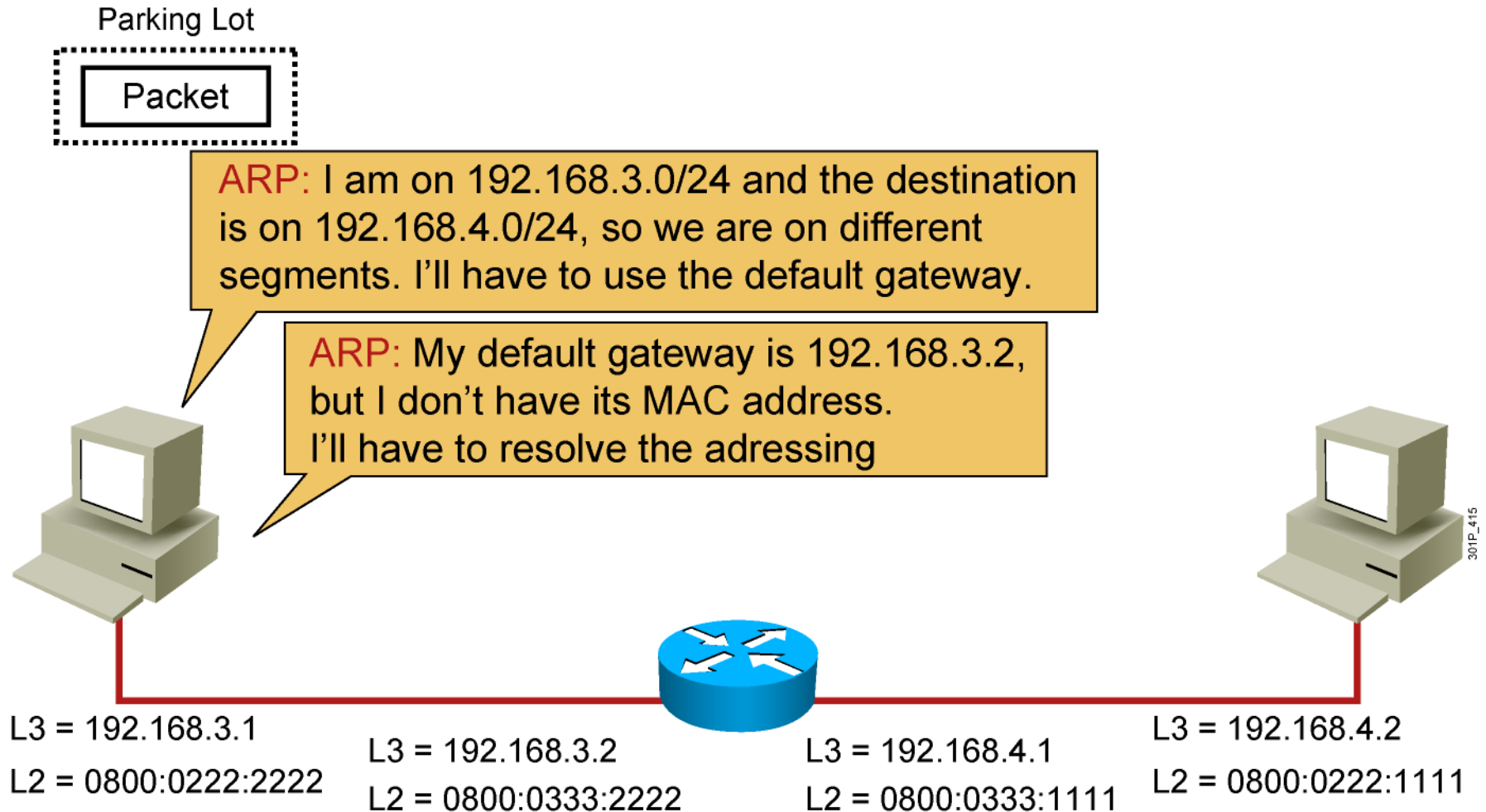
L3 = 192.168.3.1
L2 = 0800:0222:2222

L3 = 192.168.3.2
L2 = 0800:0333:2222

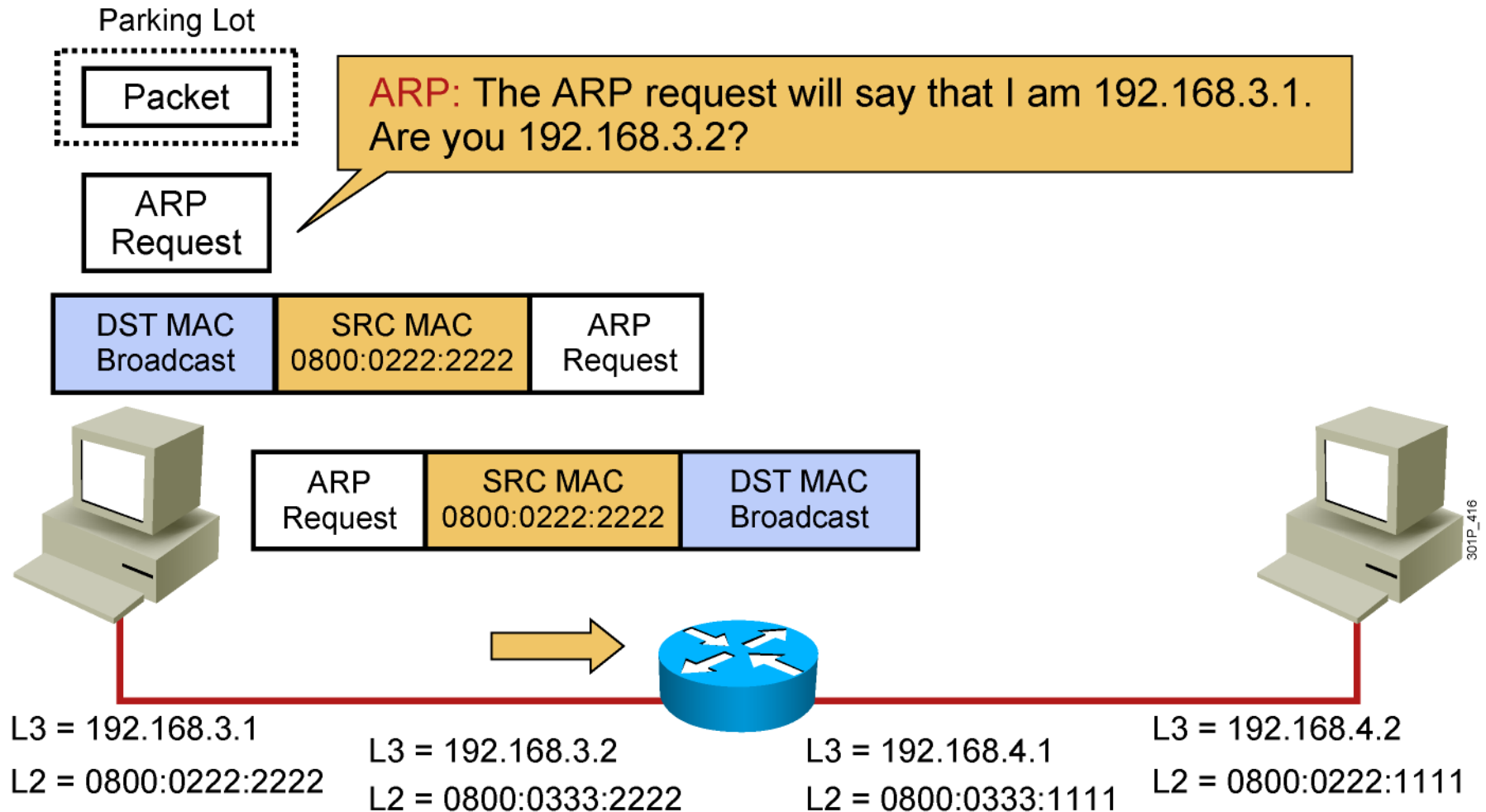
L3 = 192.168.4.1
L2 = 0800:0333:1111

L3 = 192.168.4.2
L2 = 0800:0222:1111

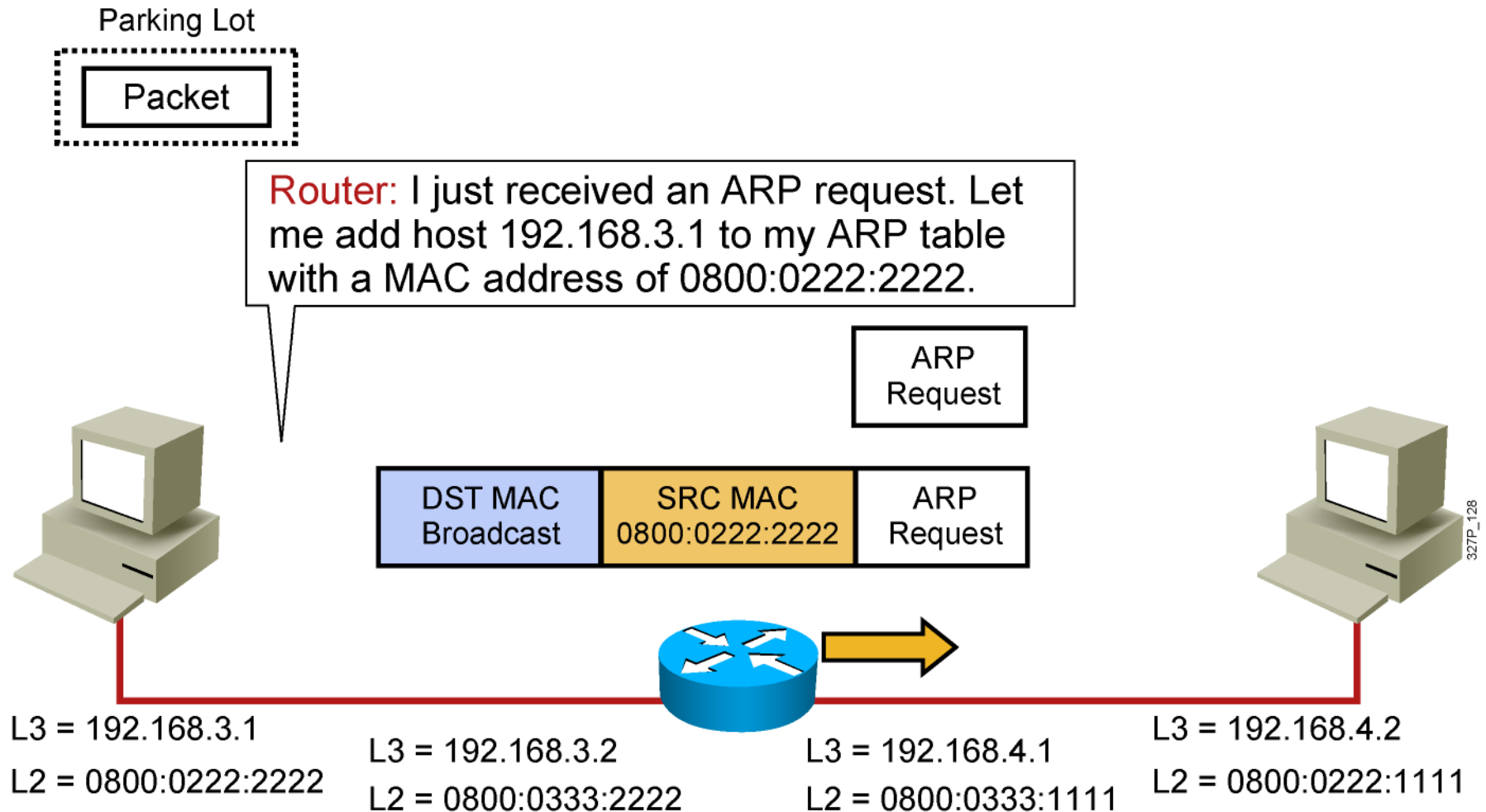
Host-to-Host Packet Delivery (4 of 17)



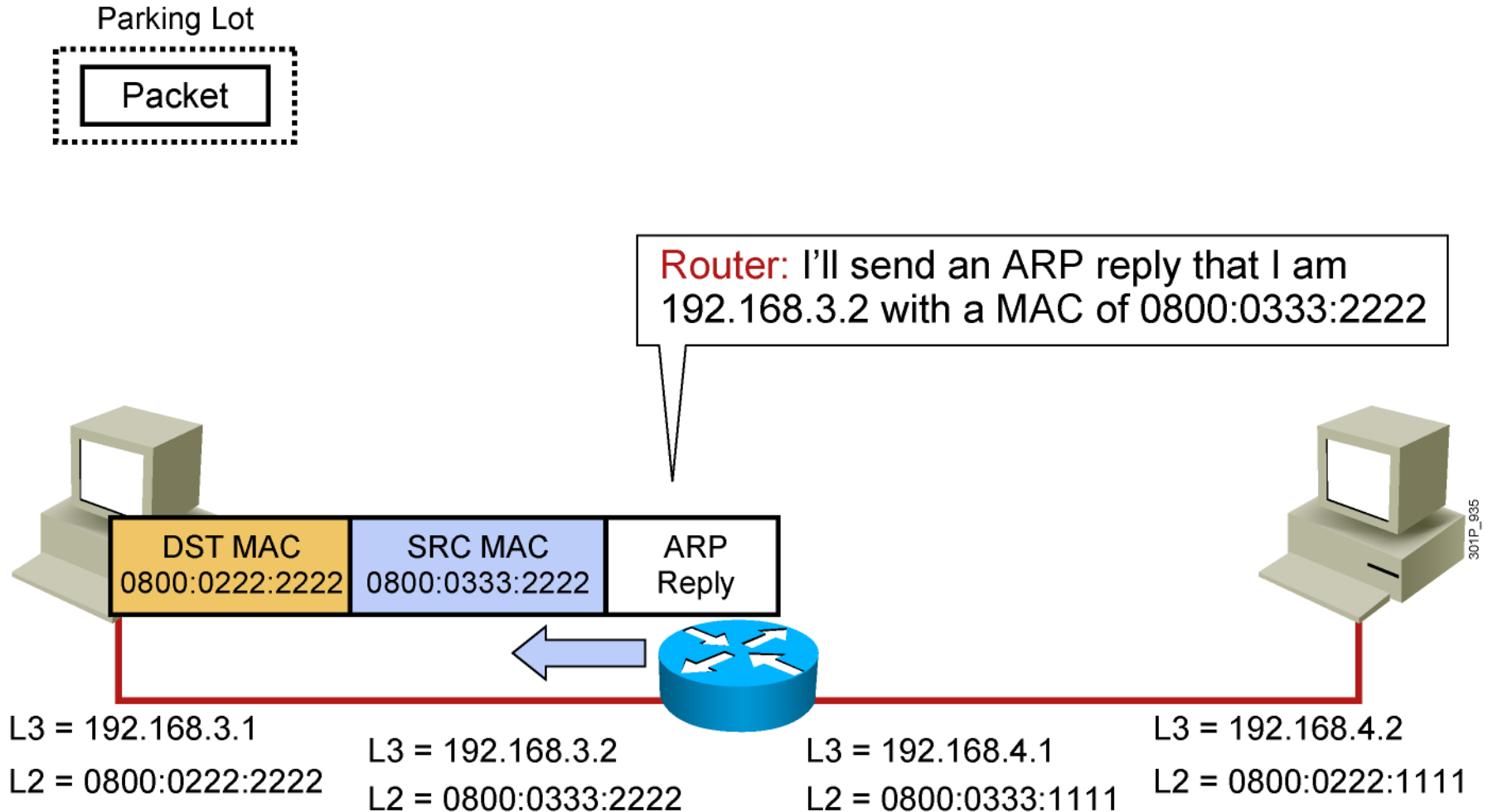
Host-to-Host Packet Delivery (5 of 17)



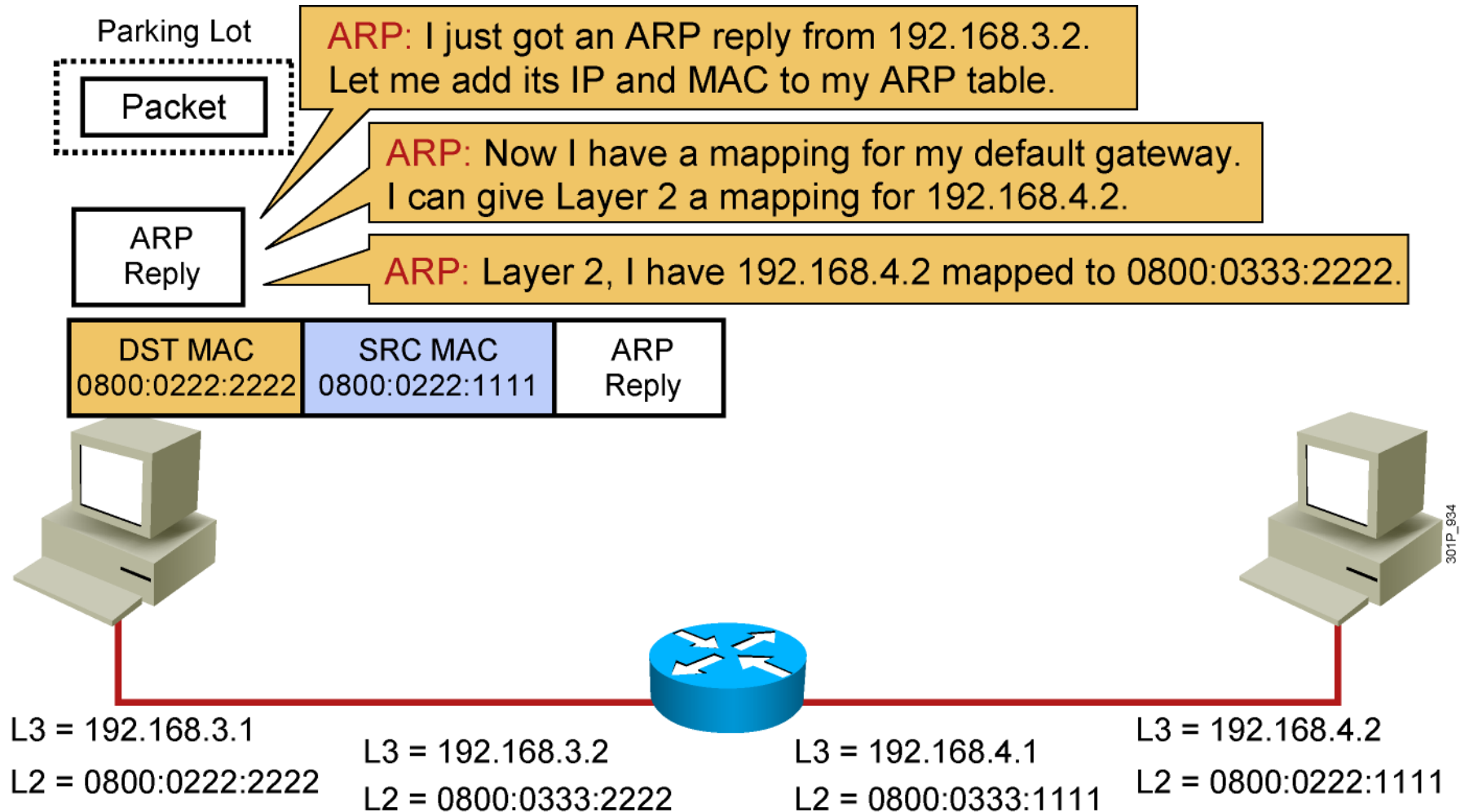
Host-to-Host Packet Delivery (6 of 17)



Host-to-Host Packet Delivery (7 of 17)

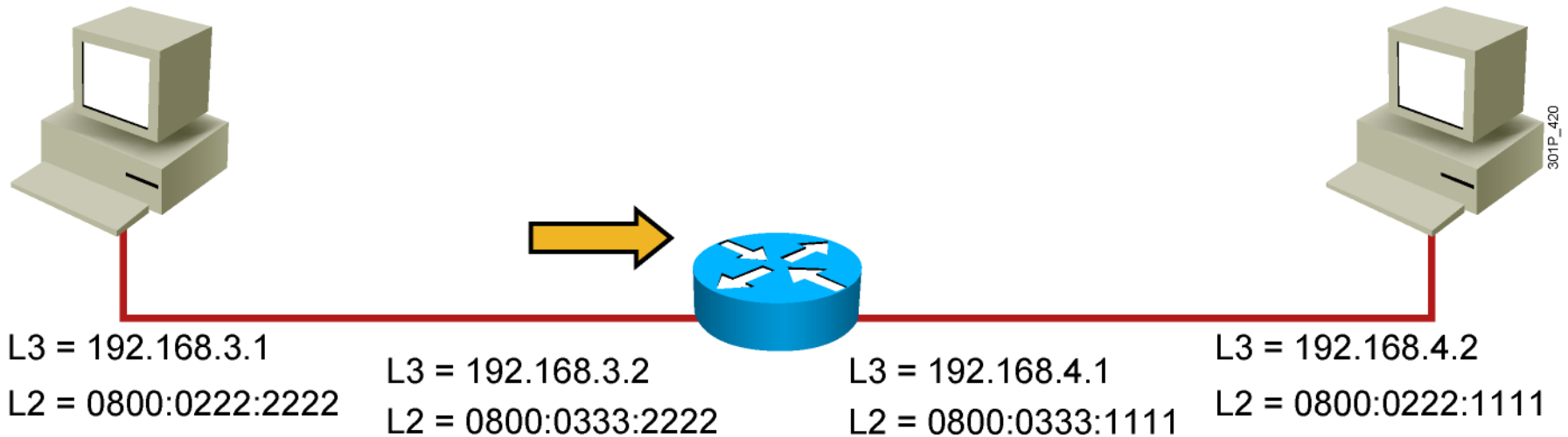


Host-to-Host Packet Delivery (8 of 17)



Host-to-Host Packet Delivery (9 of 17)

Layer 2: I can send out that pending frame.

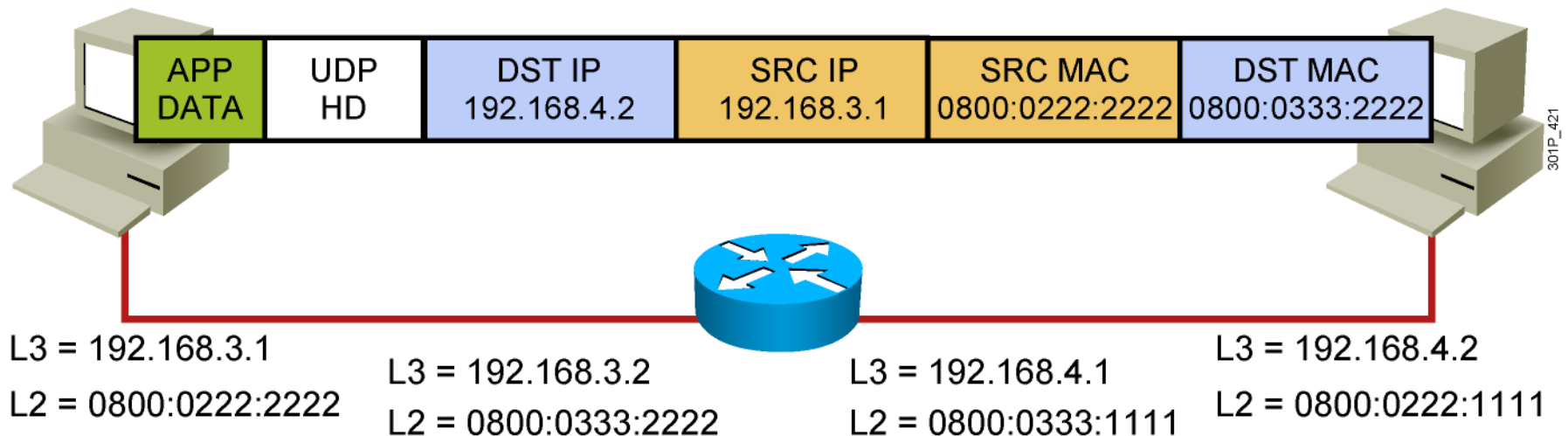
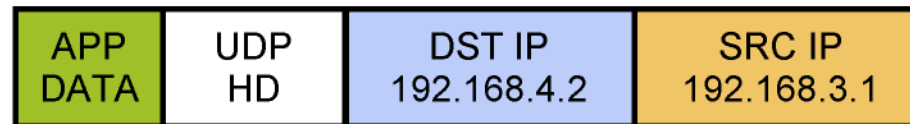


Host-to-Host Packet Delivery (10 of 17)

Router L2: I received a frame with my MAC address. I need to pass it to L3.

Router L3: This isn't my address. It needs to be routed.

Router L3: I need to forward this packet.

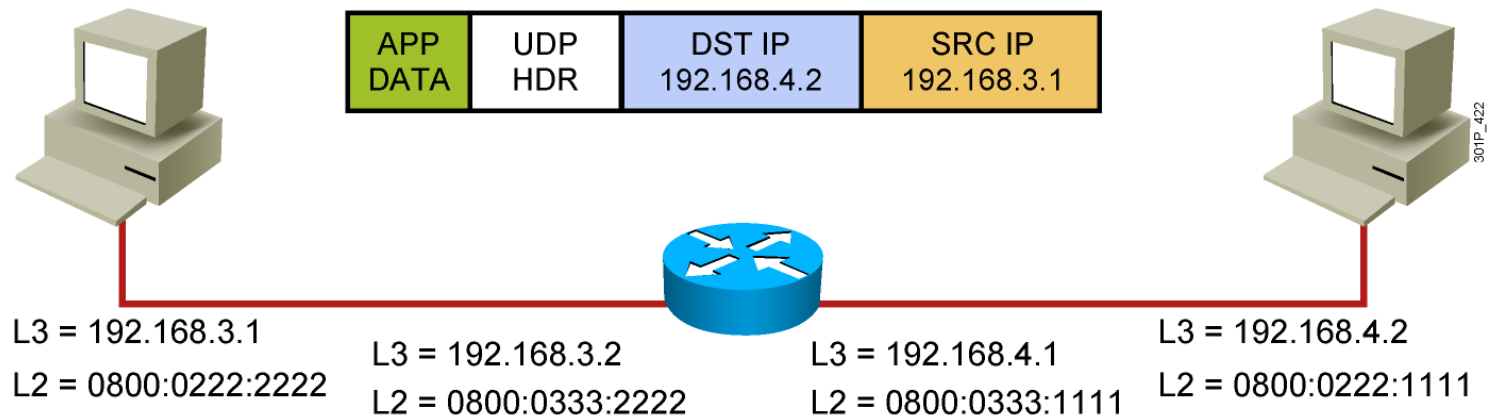


Host-to-Host Packet Delivery (11 of 17)

| Destination | Next Hop | Interface |
|----------------|-----------|-----------|
| 192.168.3.0/24 | Connected | fa 0/0 |
| 192.168.4.0/24 | Connected | fa 0/1 |

Router L3: I have an interface on the 192.168.4.0/24 segment. I can forward this packet directly to the host.

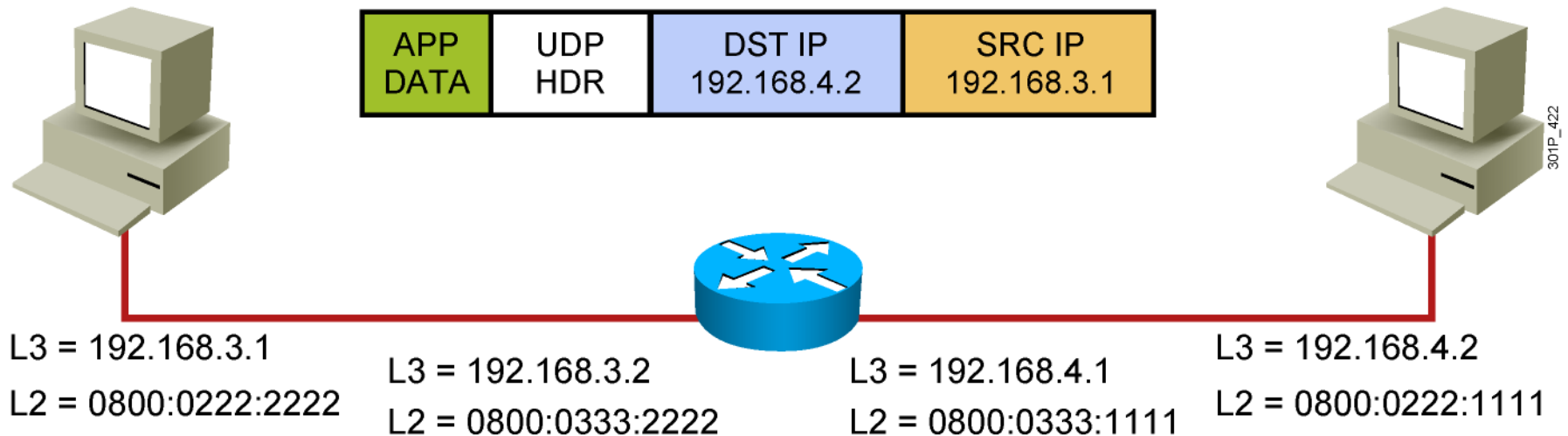
Router L3: L2, send this packet.



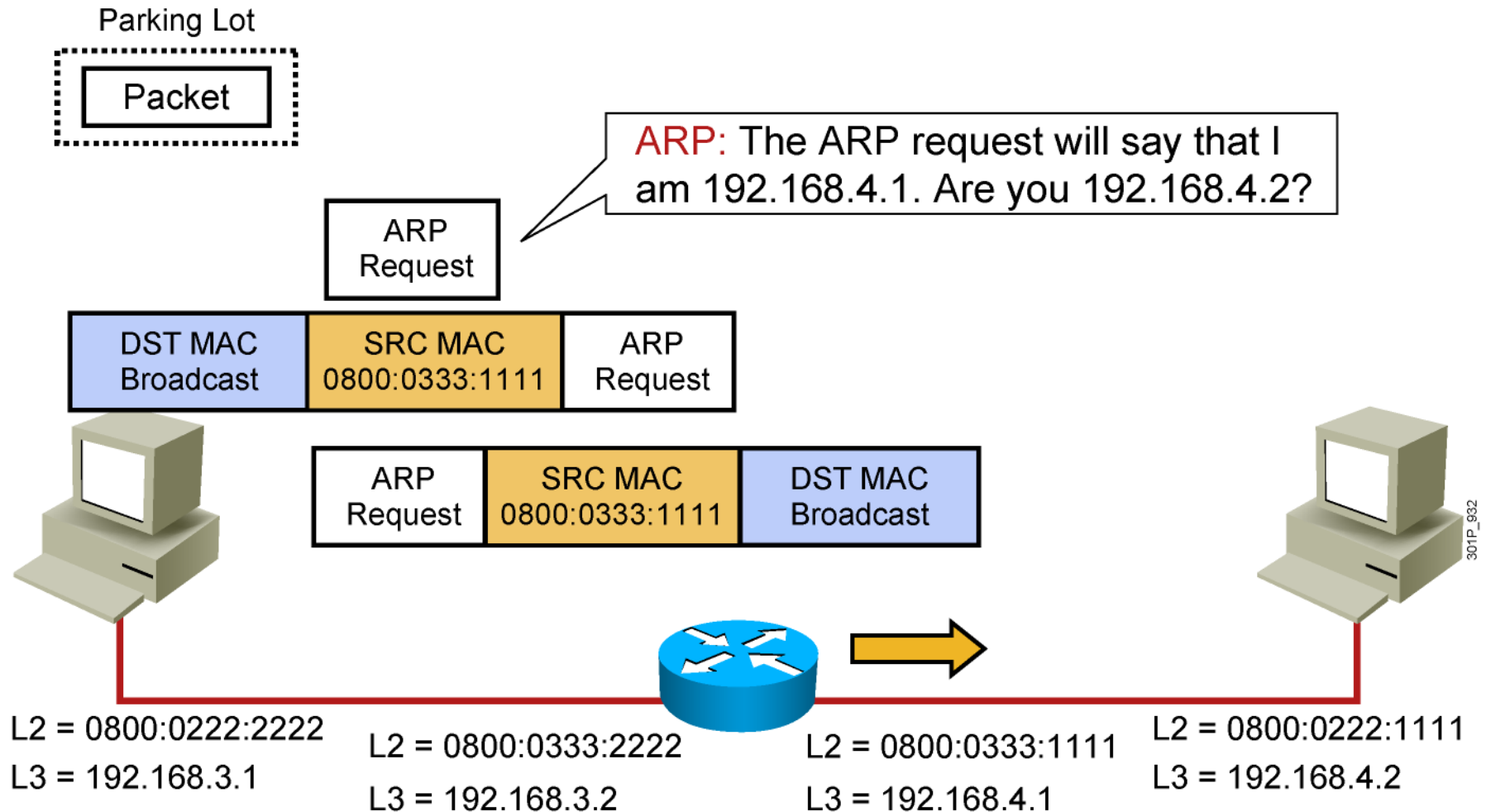
Host-to-Host Packet Delivery (12 of 17)

Router L3: I have an interface on the 192.168.4.0/24 segment. I can forward this packet directly to the host.

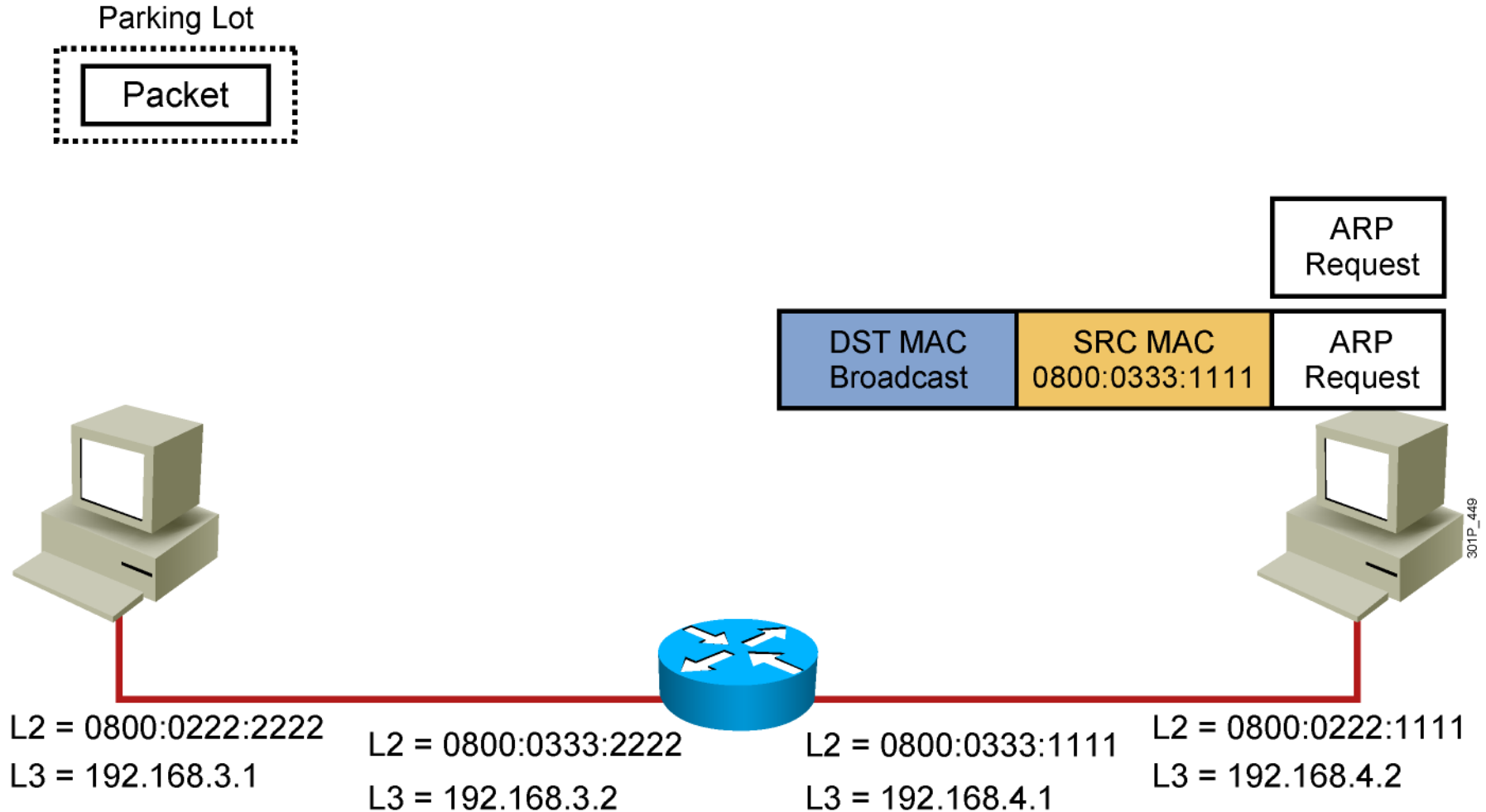
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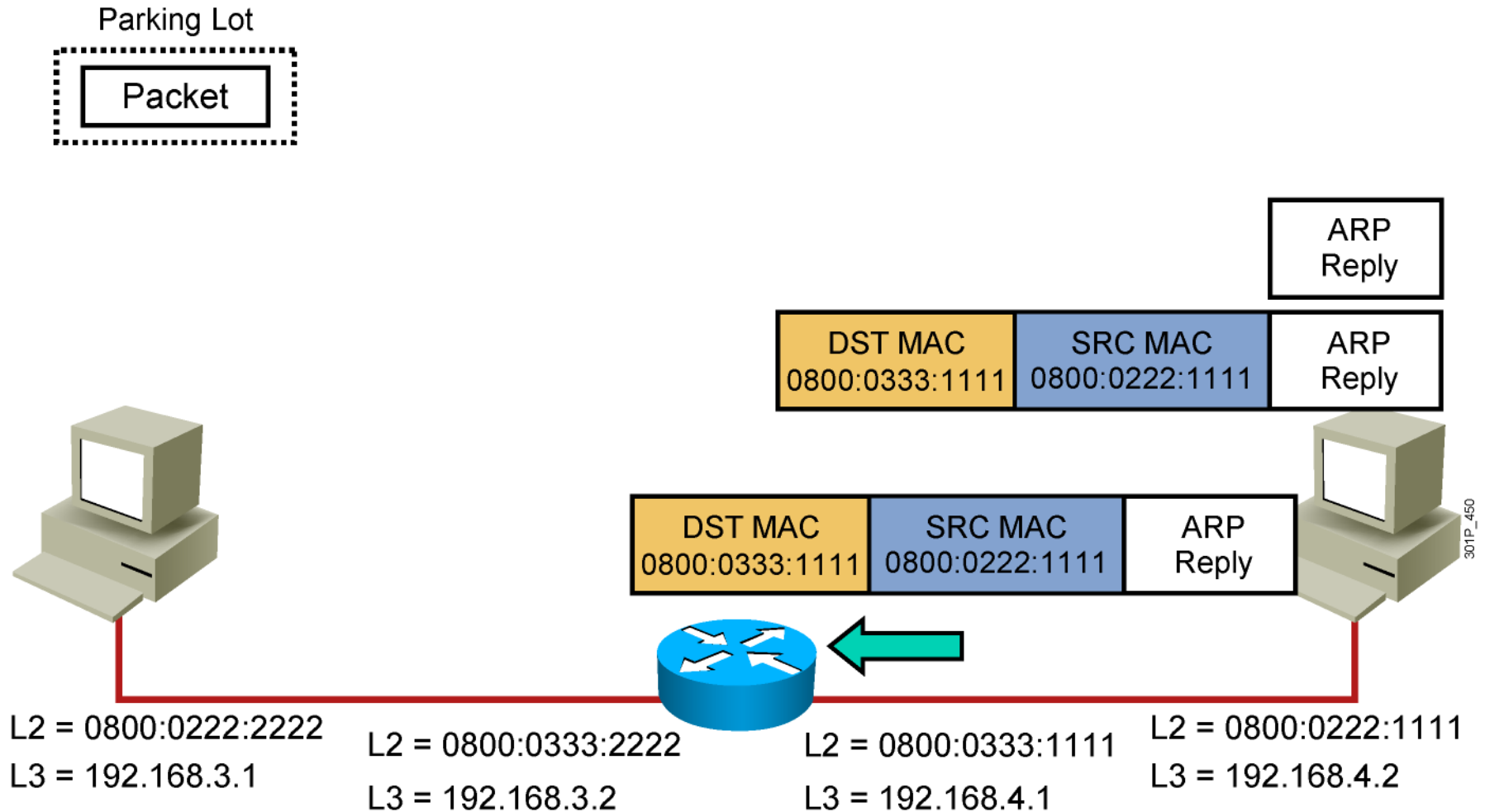
Host-to-Host Packet Delivery (13 of 17)



Host-to-Host Packet Delivery (14 of 17)



Host-to-Host Packet Delivery (15 of 17)



Host-to-Host Packet Delivery (16 of 17)

Router ARP: I just got an ARP reply from 192.168.4.2. Let me add its IP and MAC to my ARP table.

Router ARP: Now I have a mapping. I can give Layer 2 a mapping for 192.168.4.2.

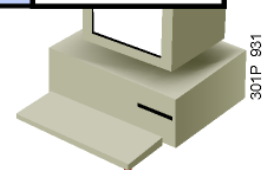
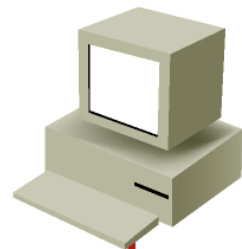
Router ARP: Layer 2, I have 192.168.4.2 mapped to 0800:0222:1111.

ARP Reply

DST MAC
0800:0333:1111

SRC MAC
0800:0222:1111

ARP Reply



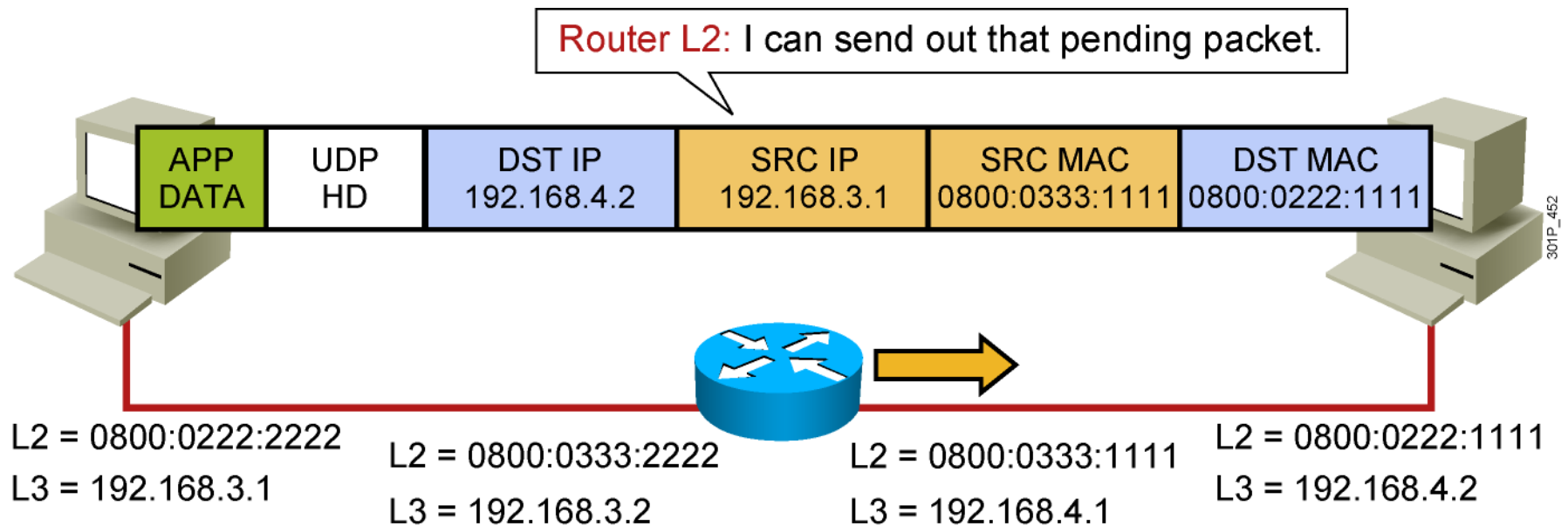
L2 = 0800:0222:2222
L3 = 192.168.3.1

L2 = 0800:0333:2222
L3 = 192.168.3.2

L2 = 0800:0333:1111
L3 = 192.168.4.1

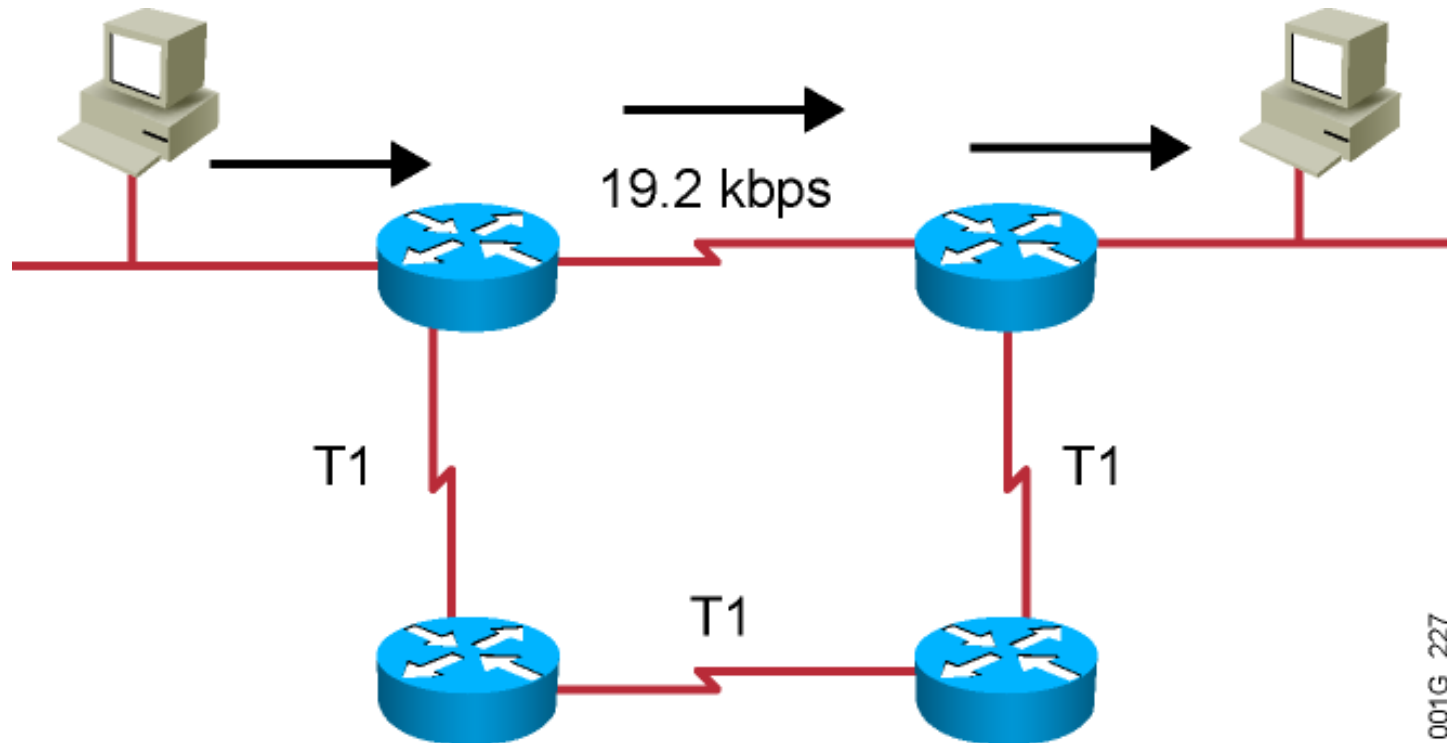
L2 = 0800:0222:1111
L3 = 192.168.4.2

Host-to-Host Packet Delivery (17 of 17)



Appendix B

RIP Overview



001G_227

- Maximum is 16 equal-cost paths (default = 4)
- Hop-count metric selects the path
- Routes update every 30 seconds



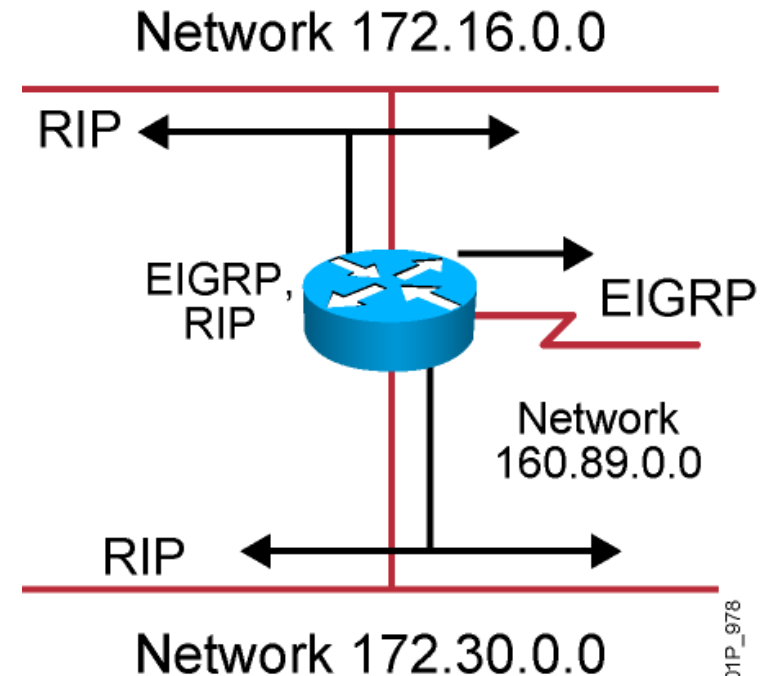
RIPv1 and RIPv2 Comparison

| | RIPv1 | RIPv2 |
|--|-----------|---------------------------|
| Routing protocol | Classful | Classless |
| Supports variable-length subnet mask? | No | Yes |
| Sends the subnet mask along with the routing update? | No | Yes |
| Addressing type | Broadcast | Multicast |
| Defined in ... | RFC 1058 | RFCs 1721, 1722, and 2453 |
| Supports manual route summarization? | No | Yes |
| Authentication support? | No | Yes |



IP Routing Configuration Tasks

- Router configuration
 - ***Select routing protocols***
 - ***Specify networks or interfaces***



301P_978

RIP Configuration

```
RouterX(config)# router rip
```

- Starts the RIP routing process

```
RouterX(config-router)# version 2
```

- Enables RIP version 2

```
RouterX(config-router)# network network-number
```

- Selects participating attached networks
- Requires a major classful network number



RIP Configuration Example

