

# Convex MINLP formulations

August 6, 2020

More details are given in the paper Kronqvist and Misener (2020), and a preprint is available from [http://www.optimization-online.org/DB\\_HTML/2020/08/7957.html](http://www.optimization-online.org/DB_HTML/2020/08/7957.html)

## 1 Big-M formulation

Using the big-M formulation the problem can be written as

$$\begin{aligned}
 \min \quad & \sum_{i=1}^n \sum_{j=i+1}^n \sum_{k=1}^m d_k^{i,j} \\
 & d_k^{i,j} \geq p_k^i - p_k^j & \forall k \in D, \forall i \in P, \forall j \in P^i, \\
 & d_k^{i,j} \geq p_k^j - p_k^i & \forall k \in D, \forall i \in P, \forall j \in P^i, \\
 & \sum_{k=1}^m (p_k^i - c_k^r)^2 \leq 1 + M_r(1 - b_{i,r}) & \forall i \in P, \forall r \in B, \\
 & \sum_{r=1}^l b_{i,r} = 1 & \forall i \in P, \quad (1) \\
 & \sum_{i=1}^n b_{i,r} \leq 1 & \forall r \in B, \\
 & p_1^i \leq p_1^{i+1} & \forall i \in P \setminus n, \\
 & b_{i,r} \in \{0, 1\} & \forall i \in P, \forall r \in B, \\
 & \mathbf{d}^{i,j} \in [0, 10]^m & \forall i \in P, \forall j \in P^i, \\
 & \mathbf{p}^i \in [0, 10]^m & \forall i \in P,
 \end{aligned}$$

where  $M_r$  are sufficiently large constants. For these problems the smallest valid  $M_r$  is simply given by

$$M_r = \max_{i \in B} \left\{ (\|\mathbf{c}^r - \mathbf{c}^i\|_2 + 1)^2 - 1 \right\}. \quad (2)$$

Here,  $\mathbf{c}^i \in R^m$  denotes the center of ball  $i$  and  $c_1^i$  refers to the first coordinate of the center. Similarly,  $\mathbf{p}^i \in R^m$  refers to point  $i$  and  $p_1^i$  is the first coordinate

of the point. To simplify the notation, we use the sets  $D = \{1, 2, \dots, m\}$ ,  $P = \{1, 2, \dots, n\}$ ,  $P^i = \{i + 1, i + 2, \dots, n\}$ , and  $B = \{1, 2, \dots, l\}$ . The binary variable  $b_{i,r}$  selects if point  $i$  is assigned to ball  $r$ . The  $\ell_1$ -distance is represented by linear constraints and by the auxiliary variables  $\mathbf{d}^{i,j} \in R^m$ . As before,  $d_k^{i,j}$  refers to the  $k$ -th component of the vector  $\mathbf{d}^{i,j}$ . The problem formulation also contains an ordering constraint of the points along the first coordinate to reduce symmetries.

## 2 Convex hull formulation

Using the convex hull formulation presented by Sawaya and Grossmann (2007), the problem can then be formulated as

$$\begin{aligned}
\min \quad & \sum_{i=1}^n \sum_{j=i+1}^n \sum_{k=1}^m d_k^{i,j} \\
& d_k^{i,j} \geq p_k^i - p_k^j & \forall k \in D, \forall i \in P, \forall j \in P^i, \\
& d_k^{i,j} \geq p_k^j - p_k^i & \forall k \in D, \forall i \in P, \forall j \in P^i, \\
& ((1 - \epsilon)b_{i,r} + \epsilon) \left( \left\| \frac{\nu_r^{\mathbf{p}^i}}{(1 - \epsilon)b_{i,r} + \epsilon} - c^r \right\|_2^2 - 1 \right) \\
& - \epsilon \left( \|c^r\|_2^2 - 1 \right) (1 - b_{i,r}) \leq 0 & \forall i \in P, \forall r \in B, \\
& \sum_{r=1}^l b_{i,r} = 1 & \forall i \in P, \\
& \sum_{i=1}^n b_{i,r} \leq 1 & \forall r \in B, \\
& p_1^i \leq p_1^{i+1} & \forall i \in P \setminus n, \\
& \sum_{r=1}^l \nu_r^{\mathbf{p}^i} = \mathbf{p}^i & \forall i \in P \\
& \mathbf{0} \leq \nu_r^{\mathbf{p}^i} \leq \mathbf{10} \cdot b_{i,r} & \forall i \in P, \forall r \in B, \\
& b_{i,r} \in \{0, 1\} & \forall i \in P, \forall r \in B, \\
& \mathbf{d}^{i,j} \in [0, 10]^m & \forall i \in P, \forall j \in P^i, \\
& \mathbf{p}^i \in [0, 10]^m & \forall i \in P, \\
& \nu_r^{\mathbf{p}^i} \in [0, 10]^m & \forall i \in P, \forall r \in B.
\end{aligned} \tag{3}$$

Here we are using the same notation as in the big-M formulation, and the variables  $\nu_r^{\mathbf{p}^i} \in R^m$  are “copies” of the variables  $\mathbf{p}^i$ . Here we use  $\epsilon = 10^{-9}$  to avoid numerical difficulties.