### CS450

Structure of Higher Level Languages

Lecture 5: Lists; quoting

Tiago Cogumbreiro

### Today we will learn...



- Being successful in CS 450
- Defining user data-structures
- Serializing code with quote
- Exercises with lists

# User data-structures

### User data-structures



We can represent data-structures using pairs/lists. For instance, let us build a 3-D point data type.

```
(require rackunit)
(define p (point 1 2 3))
(check-true (point? p))
(check-equal? (list 1 2 3) p)
(check-equal? 1 (point-x p))
(check-equal? 2 (point-y p))
(check-equal? 3 (point-z p))
(check-true (origin? (list 0 0 0)))
(check-false (origin? p))
```

### User data-structures



We can represent data-structures using pairs/lists.

For instance, let us build a 3-D point data type.

```
(require rackunit)
(define p (point 1 2 3))
(check-true (point? p))
(check-equal? (list 1 2 3) p)
(check-equal? 1 (point-x p))
(check-equal? 2 (point-y p))
(check-equal? 3 (point-z p))
(check-true (origin? (list 0 0 0)))
(check-false (origin? p))
```

### On data-structures



- We only specified immutable data structures
- The effect of updating a data-structure is encoded by **creating/copying** a data-structure
- This pattern is known as a <u>persistent data structure</u>

# Serializing code

# Quoting: a specification



Function (quote e) serializes expression e. Note that expression e is not evaluated.

- A variable x becomes a symbol 'x. You can consider a **symbol** to be a special kind of string in Racket. You can test if an expression is a symbol with function symbol?
- A function application  $(e_1 \cdots e_n)$  becomes a list of the serialization of each  $e_i$ .
- Serializing a (define x e) yields a list with symbol 'define and the serialization of e. Serializing (define  $(x_1 \cdots x_n) e$ ) yields a list with symbol 'define followed by a nonempty list of symbols ' $x_i$  followed by serialized e.
- Serializing (lambda  $(x_1...x_n)$  e) yields a list with symbol 'lambda, followed by a possibly-empty list of symbols  $x_i$ , and the serialized expression e.
- Serializing a (cond  $(b_1 \ e_1) \cdots (b_n \ e_n))$  becomes a list with symbol 'cond followed by a serialized branch. Each branch is a list with two components: serialized expression  $b_i$  and serialized expression  $e_i$ .

## Quoting exercises:



- We can write 'term rather than (quote term)
- How do we serialize term (lambda (x) x) with quote?
- How do we serialize term (+ 1 2) with quote?
- How do we serialize term (cond [(> 10 x) x] [else #f]) with quote?
- Can we serialize a syntactically invalid Racket program?

## Quoting exercises:



- We can write 'term rather than (quote term)
- How do we serialize term (lambda (x) x) with quote?
- How do we serialize term (+ 1 2) with quote?
- How do we serialize term (cond [(> 10 x) x] [else #f]) with quote?
- Can we serialize a syntactically invalid Racket program? No! You would not be able to serialize this expression (. Quote only accepts a S-expressions (parenthesis must be well-balanced, identifiers must be valid Racket identifiers, number literals must be valid).
- Can we serialize an invalid Racket program?

## Quoting exercises:



- We can write 'term rather than (quote term)
- How do we serialize term (lambda (x) x) with quote?
- How do we serialize term (+ 1 2) with quote?
- How do we serialize term (cond [(> 10 x) x] [else #f]) with quote?
- Can we serialize a syntactically invalid Racket program? No! You would not be able to serialize this expression (. Quote only accepts a S-expressions (parenthesis must be well-balanced, identifiers must be valid Racket identifiers, number literals must be valid).
- Can we serialize an invalid Racket program? Yes. For instance, try to quote the term:
   (lambda)

### Quote example



```
#lang racket
(require rackunit)
(check-equal? 3 (quote 3)) ; Serializing a number returns the number itself
(check-equal? 'x (quote x)) ; Serializing a variable named x yields symbol 'x
(check-equal? (list '+ 1 2) (quote (+ 1 2))) ; Serialization of function as a list
(check-equal? (list 'lambda (list 'x) 'x) (quote (lambda (x) x)))
(check-equal? (list 'define (list 'x)) (quote (define (x))))
```

### Manipulating quoted terms



#### Specification

```
function-dec = ( lambda ( variable* ) term+)
```

- How do we get the parameter list?
- How do we get the body?
- What does variable\* mean?
- What does term+ mean?

#### On HW1 Q.4

- The input format of the quoted term are **precisely** described in the slides of Lecture 3
- You do **not** need to test recursively if the terms in the body of a function declaration or definition are valid.
- A list, with one symbol lambda followed by zero or more symbols, and one or more terms.

# Exercises with lists



Summation of all elements of a list

### Spec

```
(require rackunit)
(check-equal? 10 (sum-list (list 1 2 3 4)))
(check-equal? 0 (sum-list (list)))
```



Summation of all elements of a list

### Spec

```
(require rackunit)
(check-equal? 10 (sum-list (list 1 2 3 4)))
(check-equal? 0 (sum-list (list)))
```

#### Solution

```
#lang racket
; Summation of all elements of a list
(define (sum-list 1)
  (cond [(empty? 1) 0]
       [else (+ (first 1) (sum-list (rest 1)))]))
```



Returns a list from n down to 1

### Spec

```
(require rackunit)
(check-equal? (list) (count-down 0))
(check-equal? (list 3 2 1) (count-down 3))
```



Returns a list from n down to 1

### Spec

```
(require rackunit)
(check-equal? (list) (count-down 0))
(check-equal? (list 3 2 1) (count-down 3))
```

#### Solution

```
#lang racket (define (count-down n) (cond [(\le n \ 0) \ (list)] [else (cons n (count-down (- n 1)))]))
```



Point-wise pairing of two lists

#### Spec



Point-wise pairing of two lists



Point-wise pairing of two lists

#### Solution