CS420

Logical Foundations of Computer Science

Lecture 7: Mock mini-test 1

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Today we will learn...

- Existential operator
- Mock Mini-Test 1
- Formal language
- Language operators
- Language equivalence



From proposition to proof state

```
Goal forall (a b c:nat), a = b → b = c.
Proof.
intros.
```

What is the expected proof state?



From proposition to proof state

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What is the expected proof state?

Solution

- Each parameter of a theorem is an assumption
- Each variable in the forall is one parameter becomes an assumption
- Each pre-condition of an implication becomes an assumption
- Variables and pre-conditions are parameters
 Boston

You can name assumptions in a forall

```
Goal forall (a b c:nat) (eq_a_b: a = b),
  b = c.
Proof.
intros.
```

What is the expected proof state?



You can name assumptions in a forall

```
Goal forall (a b c:nat) (eq_a_b: a = b),
  b = c.
Proof.
intros.
```

What is the expected proof state?

Solution

```
1 subgoal
a, b, c : nat
eq_a_b : a = b
______(1/1)
b = c
```

- Implications are just anonymous parameters (name will be generated automatically)
- Think assert (x = y) versus assert
 (Ha: x = y)



From proof state to proposition:

What is the lemma that originates the following proof state?

```
a, b, c: nat

P, Q: Prop

H: P → a = b

H0: Q \/ P

H1: b = c

-----(1/1)

a = c
```



From proof state to proposition:

What is the lemma that originates the following proof state?

```
a, b, c: nat

P, Q: Prop

H: P → a = b

H0: Q \/ P

H1: b = c

------(1/1)

a = c
```

Solution 1:

```
Goal forall (a b c: nat) (P Q: Prop) (H: P \rightarrow a = b) (H0: Q \setminus/ P) (H1: b = c), a = c.
```

Solution 2:

Goal forall (a b c: nat) (P Q: Prop), (P
$$\rightarrow$$
 a = b) \rightarrow (Q \backslash / P) \rightarrow (b = c) \rightarrow a = c.



Existential quantification

 $\exists x.P$

Existential quantification

```
Inductive ex (A : Type) (P : A \rightarrow Prop) : Prop := | ex_intro : forall (x : A) (_ : P x), ex P.
```

Notation:

```
exists x:A, P x
```

• To conclude a goal exists x:A, P x we can use tactics exist x. which yields P x. Alternatively, we can use apply ex_intro.

```
forall n, exists z, z + n = n
```

• To use a hypothesis of type H:exists x:A, P x, you can use destruct H as (x,H), or inversion H

```
forall n, (exists m, m < n) \rightarrow n <> 0.
```



Defining arbitrary logical relations

Defining less-than-equal

Inductive definition of <

$$\frac{1}{n < n}$$
 le_n

```
rac{n \leq n}{n \leq n} le_n rac{n \leq m}{n \leq \mathtt{S}\,m} le_S
```

```
Inductive le : nat \rightarrow nat \rightarrow Prop :=
  | le_n : forall n:nat,
    le n n
  | le_S : forall (n m : nat),
    le n m \rightarrow
    le n (S m).
```

- Any pre-condition will appear above the line
- Preconditions are separated by whitespace



How do we know that less-than-equal was defined correctly?



How do we know that less-than-equal was defined correctly?

With theorems!

```
(* Simple tests *)
Goal 1 \leq 1. Proof. Admitted.
Goal 1 \leq 10. Proof. Admitted.
(* More interesting properties *)
Theorem le_is_reflexive: forall x,
 x \leq x
Proof. Admitted. (* Proved in class *)
Theorem le_is_anti_symmetric: forall x y,
 x \leq y \Rightarrow
 y \leq x \rightarrow
 x = y.
Proof. Admitted. (* Proved in class *)
Theorem le_is_transitive: forall x y z,
 x \leq y \rightarrow
  y \leq z \rightarrow
 X \leq 7.
Proof. Admitted.
```



Mock Mini-Test 1

All functions defined in Coq via Fixpoint must terminate on all inputs.



All functions defined in Coq via Fixpoint must terminate on all inputs.

Solution: True

All functions must terminate.



If S(n + m) = n + Sm is the goal in the current proof state, then \code{reflexivity} will solve the goal.



If S(n + m) = n + Sm is the goal in the current proof state, then \code{reflexivity} will solve the goal.

Solution: False

```
Goal
  forall n m,
  S (n + m) = n + S m.
Proof.
  intros.
  Fail reflexivity.
Abort.
```



A *polymorphic* type is one that is parameterized by a type argument by using the universal quantifier forall. For instance: forall (X:Type), list $X \rightarrow list X$ is a polymorphic type.



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Solution: True



If E has type $beq_nat m n = true$, then E also has type m = n.



If E has type beq_nat $m \cdot n = true$, then E also has type m = n. Solution: False

```
Goal
  forall n m (E:Nat.eqb n m = true),
  m = n.
Proof.
  intros.
  Fail apply E.
Abort.
```



The proposition forall n, S n \Leftrightarrow n is provable in Coq.



The proposition forall n, S n \Leftrightarrow n is provable in Coq.

Solution: True

```
Goal
 forall n, S n <> n
Proof.
  intros.
 intros N.
  induction n. {
    inversion N.
  inversion N.
  apply IHn.
 assumption.
Qed.
```



What is the type of the following expression?

Nat.eqb 28



What is the type of the following expression?

Nat.eqb 28

Answer: nat → bool



What is the type of the following expression?



What is the type of the following expression?

14 = 68

Answer: Prop



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

forall n, n <> S n



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

forall n, n <> S n

Answer: induction



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

forall (n m:nat),
$$n = m \setminus / n <> m$$



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

forall (n m:nat), $n = m \setminus / n <> m$

Answer: BY INDUCTION



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

forall A B:Type, forall (f g: A
$$\rightarrow$$
 B), f = g \rightarrow forall x, f x = g x



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

```
forall A B:Type, forall (f g: A \rightarrow B), f = g \rightarrow forall x, f x = g x
```

Answer: EASY

```
Goal
  forall A B:Type, forall (f g: A → B), f = g → forall x, f x = g x.
Proof.
  intros.
  rewrite H.
  reflexivity.
Qed.
```



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

forall P : Prop, P



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

```
forall P : Prop, P

Answer: NOT PROVABLE

Goal
   forall P : Prop, P.
Proof.
   intros X.
   Fail apply X.
Abort.
```



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

forall n,
$$n+5 \le n+6$$



The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

forall n, $n+5 \le n+6$

Answer: INDUCTION



```
H: ~ ~ P

H0: P\/ ~ P

_____(1/1)

P
```



```
H: ~ ~ P
H0: P \/ ~ P
-----(1/1)
P

destruct H0. {
    assumption.
}
apply H in H0.
contradiction.
```



```
\begin{array}{l} H: P \rightarrow Q \\ H0: P \setminus / \sim P \\ \hline \sim P \setminus / Q \end{array} \tag{1/1}
```



```
H : P \rightarrow Q
H0: P\/~P
                                         _{-}(1/1)
~ P \/ Q
destruct H0. {
  apply H in H0.
  right.
  assumption.
left.
assumption.
```



```
P, Q: Prop

PQ: P \rightarrow Q

NQ: ~ Q

HP: P

_______(1/1)

False
```



```
P, Q: Prop

PQ: P \rightarrow Q

NQ: \sim Q

HP: P

______(1/1)

False

apply PQ in HP. contradiction.
```



```
forall (A:Type) (1:list A), 1 = [] \rightarrow 1 = []
```



```
forall (A:Type) (1:list A), 1 = [] \rightarrow 1 = [] intros. assumption.
```

