CS450

Structure of Higher Level Languages

Lecture 12: Reduction, thunks

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Today we will learn...



- 1. Optimizing code to be tail-recursive
- 2. List reduction (append, fold1)
- 3. Learn about delayed evaluation (thunks)

Acknowledgment: Today's lecture is partially inspired by Professor Dan Grossman's wonderful lecture in CSE341 from the University of Washington: (Video 1) (Video 2)

Making map tail-recursive



```
(define (map f 1)
  (cond [(empty? 1) 1]
      [else (cons (f (first 1)) (map f (rest 1)))]))
```





```
(map f (list 1 2 3)) =
; First, build the pipeline accumulator
(define (accum0 x) x) (map-iter accum0 (list 1 2 3)) =
(define (accum1 x) (accum0 (cons (f 1) x))) (map-iter accum1 (list 2 3)) =
(define (accum2 x) (accum1 (cons (f 2) x))) (map-iter accum2 (list 3)) =
(define (accum3 x) (accum2 (cons (f 3) x))) (map-iter accum3 (list)) =
; Second, run the pipeline accumulator
(accum3 (list)) =
(accum2 (list (f 3))) =
(accum1 (list (f 2) (f 3))) =
(accum0 (list (f 1) (f 2) (f 3))) =
(list (f 1) (f 2) (f 3)))
```

Tail-recursive optimization pattern



To summarize, when a value has base case and an inductive case, we identified the following pattern for a tail-recursive optimization:

Unoptimized

```
(define (rec v)
  (cond
   [(base-case? v) (base v)]
    [else (step v (rec (dec v)))]))
```

Optimized

```
(define (rec v)
  (define (rec-aux accum v)
    (cond
      [(base-case? v) (accum (base v))]
      lelse
        (rec-aux
          (lambda (x) (accum (step v x)))
          (dec v))))))
  (rec-aux (lambda (x) x) v)
```





```
(define (map f 1)
  (define (map-iter accum 1)
     (cond [(empty? 1) (accum 1)]
        [else (map-iter (lambda (x) (accum (cons (f (first 1)) x))) (rest 1))]))
  (map-iter (lambda (x) x) 1))
```

Scanning

Remove zeros from a list



Spec

```
(require rackunit)
(check-equal? (list 1 3 4) (remove-0 (list 0 1 3 0 4)))
(check-equal? (list 1 2 3) (remove-0 (list 1 2 3)))
```

Remove zeros from a list



Spec

```
(require rackunit)
  (check-equal? (list 1 3 4) (remove-0 (list 0 1 3 0 4)))
  (check-equal? (list 1 2 3) (remove-0 (list 1 2 3)))

Solution

(define (remove-0 1)
    (cond
       [(empty? 1) 1]
       [(not (equal? (first 1) 0)) (cons (first 1) (remove-0 (rest 1)))]
       [else (remove-0 (rest 1))]))
```

Can we generalize this functional pattern?



Original

```
(define (remove-0 1)
  (cond
     [(empty? 1) 1]
     [(not (equal? (first 1) 0))
        (cons (first 1) (remove-0 (rest 1)))]
     [else (remove-0 (rest 1))])
```

Generalized

```
(define (filter to-keep? 1)
  (cond
    [(empty? 1) 1]
    [(to-keep? (first 1))
     (cons (first 1)
           (filter1 to-keep? (rest 1)))]
    [else (filter to-keep? (rest 1))]))
;; Usage example
(define (remove-0 1)
  (filter
    (lambda (x) (not (equal? x 0))) 1))
```

Is this function tail-recursive?

Can we generalize this functional pattern?



Original

```
(define (remove-0 1)
  (cond
     [(empty? 1) 1]
     [(not (equal? (first 1) 0))
        (cons (first 1) (remove-0 (rest 1)))]
     [else (remove-0 (rest 1))]))
```

Generalized

```
(define (filter to-keep? 1)
  (cond
    [(empty? 1) 1]
   [(to-keep? (first 1))
     (cons (first 1)
           (filter1 to-keep? (rest 1)))]
    [else (filter to-keep? (rest 1))]))
;; Usage example
(define (remove-0 1)
  (filter
    (lambda (x) (not (equal? x 0))) 1))
```

Is this function tail-recursive? **No.** Function cons is a tail-call; filter is not.

Tail-recursive filter



Revisiting the tail call optimization

Function **filter** has very similar shape than function **map**, so we can apply the same optimization pattern.

```
(define (filter to-keep? 1)
 (define (filter-aux accum 1)
   (cond
     [(empty? 1) (accum 1)]; same as before
      lelse
        (define hd (first 1)); cache the head of the list
        (define tl (rest l)); cache the tail of the list
        (cond
          [(to-keep? hd) (filter-aux (lambda (x) (accum (cons hd x))) tl)]
          [else (filter-aux accum tl)])]))
 (filter-aux (lambda (x) x) 1))
```