CS450

Structure of Higher Level Languages

Lecture 2: Branching and function definitions

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# Logic

## Values

- Numbers
- Void
- Booleans
- Lists
- ..



## Boolean, numeric comparisons

```
value = number | boolean | · · ·
boolean = #t | #f
```

- False: #f
- True: anything that is not #f
- Logical negation: function (not e) negates the boolean result of expression e
- Numeric comparisons: <, >, <=, >=, =

To avoid subtle bugs, avoid using non-#t and non-#f values as true. In particular, **contrary to C** the number  $\emptyset$  corresponds to true. **Tip:** There is no numeric inequality operator. Instead, use (not (= x y))



# Logical and/or

```
expression = value | variable | function-call | or | and | · · ·
or = (or expression*)
and = (and expression*)
```

- Logical-and with short-circuit: and (0 or more arguments, 0-arguments yield #t)
- Logical-or with short-circuit: or (0 or more arguments, 0-arguments yield #f)



## Boolean examples

**Operations** and/or accept multiple parameters. Rectangle intersection:

```
(and (< a-left b-right)
    (> a-right b-left)
    (> a-top b-bottom)
    (< a-bottom b-top))</pre>
```

As an example of **short-circuit** logic, the expression

```
(or #t (f x y z))
```

evaluates to #t and does **not** evaluate (f x y z). Recall that and also short-circuits.



# Branching

# Branching with cond

cond evaluates each branch sequentially until the first branch's condition evaluates to true.

```
expression = value | variable | function-call | or | and | cond
cond = ( cond branch )
branch = [ condition expression ]
condition = expression | else
```

#### Example

If x is greater than 3 returns 100, otherwise if x is between 1 and 3 return 200, otherwise returns 300:

```
(cond [(> x 3) 100]
        [(> x 1) 200]
        [else 300])
```



# Creating variables

### Variable definition

A definition **binds** a variable to the result of evaluating an expression down to a value.

```
( define variable expression )
```

#### Examples

```
#lang racket
(define pi 3.14159)
pi
(* pi 2)
```

```
$ racket def-val.rkt
3.14159
6.28318
```



# Revisiting the language specification

A *program* consists of zero or more terms.

```
#lang racket term*
```

A term is either an expression or a definition.

```
term = expression | definition
```



If everything evaluates down to a value, then what does define evaluate to?

### Void

Definitions evaluate to #<void>, which is the only value that is not printed to the screen.

```
(define pi 3.14159) <-- A definition evaluates to --> #<void>
```

The void value cannot be created directly. Another way of getting a void value #<void> is by calling function (void).

Try running this program and confirm that its output is empty:

```
#lang racket
(void)
```



# Evaluating variable definition

When we execute a Racket program, we have an **environment** to bookkeep each variable, that is a map from variable names to values.

```
(define pi 3.14159)
(* pi 2)
```

```
; pi = 3.14159
#<void>
;^^^^- Eval define
(* pi 2)
```

```
; pi = 3.14159
; Prints #<void>
(* pi 2)
```



## Evaluating variable definition

When we execute a Racket program, we have an **environment** to bookkeep each variable, that is a map from variable names to values.

```
(define pi 3.14159)
(* pi 2)
```

```
; pi = 3.14159
#<void>
;^^^^- Eval define
(* pi 2)
```

```
; pi = 3.14159
; Prints #<void>
(* pi 2)
```

```
; pi = 3.14159
(* 3.14159 2)
; ^^^^^ Subst pi
```

```
; pi = 3.14159
6.28318
;^^^^- Eval func
```

```
; pi = 3.14159
; Print 6.28318
```



### Do variables evaluate?

Variables are considered expressions, so the runtime must lookup the value **bound** to a variable as one step of the evaluation.



### Beware of re-definitions

The following is legal Racket code:

```
#lang racket
(define pi 3.14159)
(* pi 2)
(define + #f)
(+ pi 2)
```

#### Redefinitions lead to subtle errors!

- Redefinitions produce subtle side-effects and may void existing assumptions
- As we will see, redefinitions also complicate the semantics and code analysis



# Exercises and recap

### Exercise: conditionals

- The modulo operator (%) in Racket is function modulo
- The equality operator in Racket is equal?

Translate the following code to Racket:

```
n = 16
if n % 15 == 0:
    return "fizzbuzz"
if n % 3 == 0:
    return "fizz"
if n % 5 == 0:
    return "buzz"
return n
```



## Exercise: conditionals (solution)

```
(define n 16)
(cond [(equal? (modulo n 15) 0) "fizzbuzz"]
       [(equal? (modulo n 3) 0) "fizz"]
       [(equal? (modulo n 5) 0) "buzz"]
       [else n])
```



```
#/ How many evaluation steps should we expect? |#
(define x (* 10 2))
(+ x (* 4 2))
```

Solution



```
#/ How many evaluation steps should we expect? |#
(define x (* 10 2))
(+ x (* 4 2))
```

#### Solution

```
; Step 1: eval *
(define x 20)
(+ x (* 4 2))
```



```
#/ How many evaluation steps should we expect? |#
(define x (* 10 2))
(+ x (* 4 2))
```

#### Solution



```
#/ How many evaluation steps should we expect? |#
(define x (* 10 2))
(+ x (* 4 2))
```

#### Solution

```
; Step 1: eval *
(define x 20)
(+ x (* 4 2))
```

```
; Step 2: eval define
; x = 20
(+ x (* 4 2))
```

```
; Step 3: eval x
; x = 20
(+ 20 (* 4 2))
```



```
#/ How many evaluation steps should we expect? |#
(define x (* 10 2))
(+ x (* 4 2))
```

#### Solution

```
; Step 1: eval *
(define x 20)
(+ x (* 4 2))

; Step 4: eval *
(+ 20 8)
```

```
; Step 2: eval define
; x = 20
(+ x (* 4 2))
```

```
; Step 3: eval x
; x = 20
(+ 20 (* 4 2))
```



```
#/ How many evaluation steps should we expect? |#
(define x (* 10 2))
(+ x (* 4 2))
```

#### Solution

```
; Step 1: eval *
(define x 20)
(+ x (* 4 2))

; Step 4: eval *
(+ 20 8)
```

```
; Step 2: eval define
; x = 20
(+ x (* 4 2))
; Step 5: eval +
28
```

```
; Step 3: eval x
; x = 20
(+ 20 (* 4 2))
```



# Function declaration

### Function declaration

A function declaration creates an anonymous function and consists of:

- parameters: zero or more parameters (identifiers, known as symbols)
- body which consist of one or more terms

When calling a function we replace each argument by the parameter defined in the lambda. If the number of parameters is not the expected one, then we get an error. The return value of the function corresponds to the evaluation of the *last* term in the body (known as the **tail position**).

```
function-dec = ( lambda ( variable* ) term+)
```

We can define circumference as a function and parameterize the radius:

```
#lang racket
(define circumference (lambda (radius) (* 2 3.14159 radius)))
(circumference 2)
```

```
$ racket func.rkt
12.56636

UMass
Boston
```

# Evaluating a lambda

```
(define circ
                                          ; circ = lambda ...
                                                              ; circ = lambda ...
  (lambda (radius) (* 2 3.14159 radius)))
                                         #<void>
                                                              : Prints #<void>
                                          :^^^^- Eval define
                                                            (circ 2)
 (circ 2)
                                          (circ 2)
; circ = lambda ...
((lambda (radius) (* 2 3.14159 radius)) 2)
              ; circ = lambda ... ; circ = lambda ...
 ; circ = lambda ...
 (* 2 3.14159 2)
                                                              ; Prints 12.56636
                                          12.56636
 :^^^^^ - Applied func
```

For more information on evaluation, read Section 1.1.5 of SICP.



# Function definition

### Function definition

Racket introduces a shorthand notation for defining functions.

```
( define (variable+ ) term+ )
```

A function definition expects one or more variables (symbols). The first variable is the function variable. The remaining variables are the arguments of the function declaration. The one-or-more terms consist of the body of the function declaration.

Which is a short-hand for:

```
( define variable (lambda ( variable* ) term+ ))
```



### Exercise

The McCarthy 91 function was invented by computer scientist John McCarthy to motivate formal verification.

$$M(n)=n-10 ext{ if } n>100 \ M(n)=M(M(n+11)) ext{ if } n\leq 100$$

- Implement the function in Racket
- What is M(99)?



### Exercise

The McCarthy 91 function was invented by computer scientist John McCarthy to motivate formal verification.

$$M(n)=n-10 ext{ if } n>100 \ M(n)=M(M(n+11)) ext{ if } n\leq 100$$

- Implement the function in Racket
- What is M(99)?

The McCarthy 91 function is equivalent to

$$M(n)=n-10 \quad ext{if } n>100 \ M(n)=91 \quad ext{if } n\leq 100$$



# Homework example

# Template excerpt

The homework template is in our Directory page:

```
(define ex1 'todo)
(define ex2 'todo)
(define ex3 'todo)
;; ...
```



#### Autograder Results

Results

Code

#### Sanity check (0.0/1.0)

Are you using the homework template?

I could not find the following definitions:

- \* define-basic?
- \* define-func?
- \* define?
- \* apply-args
- \* apply-func
- \* apply?
- \* lambda-body
- \* lambda-paráms
- \* lambda?
- \* bst-insert
- \* tree-set-value
- \* tree-set-right \* tree-set-left
- \* tree-value
- \* tree-right \* tree-left
- \* tree-leaf \* tree
- \* ex3
- \* ex2
- \* ex1

Tip #1: try assigning a dummy value to each definition. For instance:

(define define-basic? #f)

Tip #2: ensure your definitions are made public. The first two lines of your file should be:

#lang racket (provide (all-defined-out)) STUDENT

AUTOGRADER SCORE

0.0 / 24.0

FAILED TESTS

Sanity check (0.0/1.0)

