CS420

Introduction to the Theory of Computation

Lecture 11: Regular expressions

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Today we will learn...



- Language equivalence theorems
- Summary of Lang. Examples
- Regular expressions



• $L_1 \cup (L_2 \cup L_3) = (L_1 \cup L_2) \cup L_3$



- $L_1 \cup (L_2 \cup L_3) = (L_1 \cup L_2) \cup L_3$
- $\bullet \ (L_1 \cup L_2) = (L_2 \cup L_1)$



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- $(L_1 \cup L_2) = (L_2 \cup L_1)$
- $(L \cdot \emptyset) = (\emptyset \cdot L) = \emptyset$



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- ullet $(L_1\cdot L_3)\cup (L_2\cdot L_3)=ig((L_1\cup L_2)\cdot L_3ig)$

Lang.v examples



(Live coding...)

Regular expressions

Regular Expressions: Input validation



Regular Expressions: Input validation

HTML includes regular expressions to perform client-side form validation.

```
<input id="uname" name="uname" type="text"
    pattern="_([a-z]|[A-Z]|[0-9])+" minlength="4" maxlength="10">
```

- _[a-zA-Z0-9]+
- [a-zA-Z0-9] means any character beween a and z, or between A and Z, or between 0 and 9
- R+ means repeat R one or more times
- In this case, the username must start with an underscore _, and have one or more letters/numbers
- minlength and maxlength further restrict the string's length

Regular Expressions: Text manipulation



Regular Expressions: Text manipulation

Programming languages include regular expressions for fast and powerful text manipulation.

Example (JS)

```
let txt1 = "Hello World!";
let txt2 = txt1.replace(/[a-zA-Z]+/, "Bye"); // Replaces the first word by "Bye"
console.log(txt2);
// Bye World!
```

Regular expressions



A theoretical motivation

- What languages can we specify with the following operators?
 - Void
 - Nil
 - Char
 - App
 - Union
 - Star

Regular expressions



A theoretical motivation

- What languages can we specify with the following operators?
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Idea: specify a datatype that represents all possible expressions





We define regular expression R as either:

Notation	Meaning
Ø	Rejects all words
ϵ	Only accepts the empty string
c	Only accepts a string with a single character $oldsymbol{c}$
R_1R_2	Accepts a word from R_1 concatenated with a word from R_2
$R_1 R_2$	Accepts a word from R_1 or a word from R_2
R^{\star}	Accepts 0 or more copies of words from R



Expression: foo || bar

Is the following input accepted or rejected?

• Input: [f,o,o]



Expression: foo || bar

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• Input: [f,o,o]
Accept



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• **Input:** [b,a,r]



Expression: foo || bar

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- Input: [f,o,o]
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- Input: [b,a,r]
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- **Input:** [f,o]



Expression: foo || bar

Is the following input accepted or rejected?

- Input: [f,o,o]

 Accept
- Input: [b,a,r]
 Accept
- Input: [f,o]
 Reject
- Input: []



Expression: foo || bar

Is the following input accepted or rejected?

- Input: [f,o,o]

 Accept
- Input: [b,a,r]
 Accept
- Input: [f,o]
 Reject
- Input: []
 Reject



Expression: $(foo||bar)^*$

• Input: [f,o,o]



Expression: $(foo||bar)^*$

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Expression: $(foo||bar)^*$

- Input: [f,o,o]
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Expression: $(foo||bar)^*$

- Input: [f,o,o]
 Accept
- Input: [b,a,r]
 Accept
- **Input:** [f,o]



Expression: (foo||bar)*

- Input: [f,o,o]
 Accept
- Input: [b,a,r]
 Accept
- Input: [f,o]
 Reject
- Input: []



```
Expression: (foo||bar)*
```

- Input: [f,o,o]
 Accept
- Input: [b,a,r]
 Accept
- Input: [f,o]
 Reject
- Input: []
 Accept
- **Input:** [f,o,o,b,a,r]



```
Expression: (foo||bar)*
```

- Input: [f,o,o]
 Accept
- Input: [b,a,r]
 Accept
- Input: [f,o]
 Reject
- Input: []
 Accept
- Input: [f,o,o,b,a,r]

 Accept



Expression

$$(\emptyset \cdot \mathtt{c}) \mid\mid \mathtt{aa} \mid\mid \mathtt{a} \cdot \epsilon$$

• [b]



$$(\emptyset \cdot \mathtt{c}) \mid\mid \mathtt{aa} \mid\mid \mathtt{a} \cdot \epsilon$$

- [b] REJECT
- [b,c,a]



$$(\emptyset \cdot c) \mid\mid aa \mid\mid a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT



$$(\emptyset \cdot c) \mid\mid aa \mid\mid a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT
- [C]



$$(\emptyset \cdot c) \mid\mid aa \mid\mid a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b]



$$(\emptyset \cdot c) \mid\mid aa \mid\mid a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b] REJECT
- []



$$(\emptyset \cdot c) \mid\mid aa \mid\mid a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b] REJECT
- [] ACCEPT
- [a,a]



$$(\emptyset \cdot c) \mid\mid aa \mid\mid a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b] REJECT
- [] ACCEPT
- [a,a] ACCEPT
- [a]



$$(\emptyset \cdot c) \mid\mid aa \mid\mid a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b] REJECT
- [] ACCEPT
- [a,a] ACCEPT
- [a] ACCEPT

Examples



Source: <u>regexlib.com</u>

This expression matches a hyphen separated US phone number, of the form ANN-NNN-NNNN, where A is between 2 and 9 and N is between 0 and 9.

• [2-9][0-9]{2}-[0-9]{3}-[0-9]{4}

Examples



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Breaking it down:

- [2-9] corresponds to 2 || 3 || 4 || 5 || 7 || 8 || 9
- [0-9]{2} corresponds to the power of 2, thus pattern [0-9][0-9]
- -
- [0-9]{3}
- -
- [0-9]{4}