CS450

Structure of Higher Level Languages

Lecture 38: SimpleJS; translating LambdaJS to SimpleJS

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My goal with CS450 is to teach you ...



- 1. Fundamental concepts behind most programming languages
 - functional programming, delayed evaluation, control flow and exceptions, object oriented systems, monads, macros, pattern matching, variable scoping, immutable data structures
- 2. A framework to describe language concepts
 - λ -calculus and formal systems to specify programming languages
 - functional programming and monads to implement specifications
- 3. A methodology to understand complex systems
 - (formally) specify and implement each programming language feature separately
 - understand a complex system as a combination of smaller simpler systems
 - implement and test features independently

Today we will...



- Revise JavaScript's object system
- Introduce SimpleJS: S-Expression-based syntax and simpler JavaScript rules
- Introduce LambdaJS: λ -calculus + references + immutable objects
- Introduce translation from SimpleJS into LambdaJS

Why are we learning all SimpleJS and LambdaJS?

- You already know λ -calculus with references (heap)
- You already know how objects work (ie, a map with a lookup that work like frames and environments)
- I want to teach you the fundamentals of JavaScript by building it on top of concepts that you already know!
- I can introduce another kind of specifying the semantics of a system, by translating it into another system (denotational semantics)

Object prototypes



A.__proto__ = B links A object to B, if a field f is not available in A, then it is looked up in B (which works recursively until finding undefined).

```
a = {"x": 10, "y": 20}
b = {"x": 30, "z": 90, "__proto__": a}
b {x: 30, z: 90, *y: 20}
```

Functions are constructors

If we call a function A with new, then A is called as the constructor of a new object.

```
function C(x, y) { this.x = x; this.y = y }
c = new C(10, 20)
c {x: 10, y: 20}
```

Constructor's prototype

If A is a function, then A.prototype becomes the $__$ proto $__$ of every object created using A with new.

```
C.prototype = {"foo": true, "bar": 100}
d = new C(10, 20)
d {x: 10, y: 20, *foo: true, *bar: 100}
```

What is the name of the paper we are studying?

Quiz

SimpleJS

Introducing SimpleJS



- SimpleJS is just a simplification of JavaScript with fewer corner case, which is easier to learn.
- SimpleJS was created by your instructor for CS450 (yet close to what you have in The Essence of JavaScript)
- SimpleJS has a formal syntax (below) and also an S-expression syntax (hw8-util.rkt)
- Today we will **formally** describe SimpleJS in terms of how we can represent it in LambdaJS (defined in The Essence of JavaScript).

$$e ::= x \mid \mathtt{let} \ x = e \ \mathtt{in} \ e \mid x.y \mid x.y := e \mid x.y(e \cdots) \mid \mathtt{function}(x \cdots) \{e\} \mid \mathtt{new} \ e(e \cdots) \mid \mathtt{class} \ \mathtt{extends} \ e \ \{\mathtt{constructor}(x \cdots) \{e\} \ m \cdots \}$$

$$m ::= x(x \cdots)\{e\}$$

Writing Shape in SimpleJS



JavaScript

```
function Shape(x, y) {
    this.x = x;
    this.y = y;
let p = new Shape(10, 20);
Shape.prototype.translate =
    function(x, y) {
      this.x = this.x + x;
      this.y = this.y + y;
p.translate(1,2);
return p;
```

SimpleJS

```
(let Shape
  (function (x y)
    (begin (set! this.x x)
           (set! this.y y)))
 (let p (new Shape 10 20)
    (let Shape-proto Shape.prototype
      (begin
        (set! Shape-proto.translate
          (function (x y)
            (begin
              (set! this.x (! + this.x x))
              (set! this.y (! + this.y y)))))
        (p.translate 1 2)
        p))))
```



JavaScript

SimpleJS



JavaScript

SimpleJS

What are the possible problems of this form of inheritance?



JavaScript

SimpleJS

What are the possible problems of this form of inheritance?

How can we add a new method to Rectangle?



With the highlighted pattern we can safely mutate Rectangle.prototype. This is the same as Rectangle.prototype = {'_proto_': Shape.prototype }, but we have no syntax for such a pattern in SimpleJS.

JavaScript

```
function Rectangle(width, height) {
   this.x = 0;
   this.y = 0;
   this.width = width;
   this.height = height;
}
let p = function () {}
p.prototype = Shape.prototype;
Rectangle.prototype = new p();
let r1 = new Rectangle(10, 20);
return r1;
```

SimpleJS

LambdaJS

LambdaJS



Think Racket without define, without macros, with objects, and heap operations.

Expressions

$$e ::= v \mid x \mid \lambda x.e \mid e(e) \mid \{s \colon e\} \mid e[e] \mid e[e] \leftarrow e \mid \mathtt{alloc} \ e \mid e := e$$

Concrete LambdaJS S-expression syntax



Formal syntax	S-expression
$\lambda x.e$	(lambda (x) e)
$e_1(e_2)$	(e1 e2)
$\{ t "foo": 1+2, t "bar": x\}$	(object ["foo" (+ 1 2)] ["bar" x])
$o[exttt{"foo"}]$	(get-field o "foo")
$\verb"alloc" \{\}$	(alloc (object))
$x := \{\}$	(set! x (object))
x:=1;x	(begin (set! x 1) x)
$let\; x\; =\; 10\;in\; x+4$	(let ([x 10]) (+ x 4))

In Racket you can actually allocate a reference with (box e), which is equivalent to LambdaJS(alloc e), and update the contents of that reference with (set-box! b e), which is equivalent to LambdaJS (set! e).

Translating SimpleJS into LambdaJS

Overview

Translating SimpleJS into LambdaJS



- 1. A SimpleJS object is represented as a reference to an immutable LambdaJS object
- 2. A SimpleJS function is represented as an object with two fields: (a) a lambda-function that represents the code, a prototype field which points to an empty SimpleJS object
- 3. Create an object with new expects a SimpleJS function as argument and must create a new object, initialize its prototype, and call the constructor function (see point 2)
- 4. Method invocation corresponds to accessing a SimpleJS function and passing the implicit this object to it (see 2)

Objectives of the translation

- Explicit this
- Functions are not objects: convert function into an object+lambda
- Explicit memory manipulation
- No method calls: use function calls

Translating a function



JavaScript

```
function Shape(x, y) {
  this.x = x;
  this.y = y;
};
```

Step 1: only objects and lambdas

```
Shape = {
    '$code': (obj, x, y) ⇒ {
      obj.x = x;
      obj.y = y;
    },
    'prototype' = {}
};
```

Translating a function



JavaScript

```
function Shape(x, y) {
  this.x = x;
  this.y = y;
};
```

Step 1: only objects and lambdas

```
Shape = {
  '$code': (obj, x, y) ⇒ {
    obj.x = x;
    obj.y = y;
  },
  'prototype' = {}
};
```

Step 2: explicit references

```
Shape = alloc {'$code': (this, x, y) \Rightarrow {
    this = (deref this)["x"] \leftarrow x; // In LambdaJS we have to replace the whole object
    this = (deref this)["y"] \leftarrow y;},
    'prototype': alloc {}};
```

Translating new



JavaScript

```
p1 = new Shape(0, 1);
```

Step 1: only objects and lambdas; no implicit this

```
p1 = {"__proto__": Shape.prototype};
Shape["$code"](p1, 0, 1);
```

Translating new



JavaScript

```
p1 = new Shape(0, 1);
```

Step 1: only objects and lambdas; no implicit this

```
p1 = {"__proto__": Shape.prototype};
Shape["$code"](p1, 0, 1);
```

Step 2: explicit references

```
p1 = alloc {"__proto__": (deref Shape)["prototype"]}};
(deref Shape)["$code"](p1, 0, 1);
```

Translating method invocation



JavaScript

```
p1.translate(10, 20);
```

Step 1: only objects and lambdas; no implicit this

```
m = p1["translate"];  // get object method
m["$code"](p1, 10, 20); // get code for method
```

Translating method invocation



JavaScript

```
p1.translate(10, 20);
```

Step 1: only objects and lambdas; no implicit this

```
m = p1["translate"];  // get object method
m["$code"](p1, 10, 20); // get code for method
```

Step 2: explicit references

Formally

```
m = (deref p1)["translate"];
(deref m)["$code"](p1, 10, 20);
```

SimpleJS

```
(let ([m (get-field (deref p1) "translate")])
  ((get-field (deref m) "$code") p1 10 20))
```

Translating SimpleJS into LambdaJS



Before

```
Shape.prototype.translate = function(x, y)
    this.x += x; this.y += y;
};
p1 = new Shape(0, 1);
p1.translate(10, 20);
```

After

```
// 1. Function declaration
Shape = alloc {
  "code": (this, x, y) \Rightarrow { ... },
  "prototype" = alloc {}};
p = (deref Shape)["prototype"];
(deref p)["translate"] = alloc {
  "code": (this, x, v) \Rightarrow { ... }
  "prototype": alloc {}}:
// 2. new
p1 = alloc {"__proto__":
           (deref Shape)["prototype"]};
(deref Shape)["$code"](p1, 0, 1);
// 3. method call
f = (deref p1)["translate"];
(deref f)["$code"](p1, 10, 20):
```

Translation function

Translation function



- Field lookup
- Field update
- Function declaration
- The new keyword
- Method call
- Class declaration

Field lookup

Field lookup



$$\mathbf{J}[\![x.y]\!] = (\mathsf{deref}\ x)["y"]$$

SimpleJS

 λ -JS

this.x

(get-field (deref this) "x")

Field update

Field update



In JavaScript, assigning an expression e into a field, returns the evaluation of e. However, in LambdaJS assignment returns the reference being mutated.

```
\mathbf{J} \llbracket x.y := e 
rbracket = 
\mathsf{let} \; \mathsf{data} = \mathbf{J} \llbracket e 
rbracket \; \mathsf{in}
x := (\mathsf{deref} \; \mathsf{obj}) ["y"] \leftarrow \mathsf{data};
\mathsf{data}
```

SimpleJS

 λ -JS

```
(set! this.x x)
```

```
(let [(data x)]
  (begin
     (set! this
           (update-field (deref this) "x" data))
     data)))
```

Free variables and bound variables



$$\mathbf{J}\llbracket x.y := e
rbracket = \mathsf{let} \; \mathsf{data} = \mathbf{J}\llbracket e
rbracket \; \mathsf{in} \; \mathsf{x} := (\mathsf{deref} \; \mathsf{x})["y"] \leftarrow \mathsf{data}; \mathsf{data}$$

SimpleJS

 λ -JS

(set! data.x 10)

What happened here?

Free variables and bound variables



$$\mathbf{J}\llbracket x.y := e \rrbracket = \mathsf{let} \; \mathsf{data} = \mathbf{J}\llbracket e \rrbracket \; \mathsf{in} \; \mathsf{x} := (\mathsf{deref} \; \mathsf{x})["y"] \leftarrow \mathsf{data}; \mathsf{data}$$

SimpleJS

 λ -JS

(set! data.x 10)

What happened here?

- 1. Variable data is used in the generated code
- 2. We must ensure that data is not captured (free) in the generated code!

What problem occurs when generating code?

(One sentence is enough.)

Quiz

Function declaration

Function declaration



Field **prototype** can be accessed by the user, so we declare it as a reference. Field **\$code** does not actually exist in JavaScript, so we prefix it with a dollar sign (\$) to visually distinguish artifacts of the translation.

SimpleJS

```
(function (x y)
  (begin
        (set! this.x x)
        (set! this.y y)))
```

λ -JS

The keyword

The keyword



```
\begin{split} \mathbf{J}\llbracket \mathtt{new} \ e_f(e\cdots) \rrbracket = \\ \mathsf{let} \ \mathsf{ctor} = \mathsf{deref} \ \mathbf{J}\llbracket e_f \rrbracket \ \mathsf{in} \\ \mathsf{let} \ \mathsf{obj} = \mathsf{alloc} \ \{\texttt{"\$proto"} : \mathsf{ctor}[\texttt{"prototype"}] \} \ \mathsf{in} \\ \mathsf{ctor}[\texttt{"\$code"}] (\mathsf{obj}, \mathbf{J}\llbracket e \rrbracket \cdots); \\ \mathsf{obj} \end{split}
```

SimpleJS

 λ -JS

(new Shape 0 1)

```
(let [(ctor (deref Shape))
        (o (alloc (object "$proto" (get-field ctor "prototype"))))]
    (begin
        ((get-field ctor "$code") o 0 1)
        o))
```

Method invocation

Method invocation



```
\mathbf{J}[\![x.y(e\cdots)]\!]] = \Big( \big( \mathsf{deref} \, \, \big( (\mathsf{deref} \, \, x)["\mathtt{y}"] \big) \big) ["\$ \mathsf{code}"] \Big) (x, \mathbf{J}[\![e\cdots]\!])
```

SimpleJS

 λ -JS

(p1.translate 10 20)

```
((get-field
  (deref (get-field (deref p1) "translate"))
  "$code")
  p1 10 20)
;; In Racket pseudo code
(define p1:obj (deref p1)); 1. get obj from ref
(define translated:m (get-field p1:obj "translate")); 2. get fiel
(define translated:o (deref translated:m)); 3. get object from re
(define translated:f (get-field translated:o "$code"); 4. get fun
(translated:f p1 10 20); 5. call fun pass this (p1)
```

Function call



We will not be implementing function calls in Homework Assignment 8.

```
egin{aligned} & \mathbf{J}\llbracket e_f(e\cdots)
rbracket = & \mathbf{J}\llbracket e_o
rbracket = & \mathbf{J}\llbracket e_o
rbracket = & \mathbf{J}\llbracket e\cdots
bracket = & \mathbf{J}\llbracket e\cdots 
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```

Example 1

```
class Foo {
  constructor() { this.x = 0; }
  bar() { this.x++; }
}
var foo = new Foo();
foo["bar"](); // foo.bar();
// Caveat: foo.bar() ≠ (foo.bar)()
```

Example 2

```
class Foo {
  constructor() { this.x = 0; }
  bar() { this.x++; }
}
var foo = new Foo();
var bar = foo["bar"];
bar(); // TypeError: this is undefined
```

Class declaration

Class declaration



To allow dynamically dispatching to X's methods, the first four lines instantiate X without calling its constructor. This way, we can safely mutate the **cls**'s prototype without affecting X and any changes to X are visible to **cls** via lookup.

```
\mathbb{C}[ \text{class extends } X \{body\} ] =
                        let parent = \mathbb{C}[X] in
                 let parent' = function (){} in
             parent'.prototype := parent.prototype
                     let proto = new parent' in
                let cls = function (x \cdots) \{e_c\} in
                      cls.prototype := proto;
            	extstyle{\mathsf{proto.m}} := 	extstyle{\mathsf{function}}(y\cdots)\{e_m\};\cdots
                                    cls
where body = \mathtt{constructor}(x\cdots)\{e_c\}\ m(y\cdots)\{e_m\}\cdots
```