CS450

Structure of Higher Level Languages

Lecture 3: Lists and code serialization

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# Today we will learn...

- Using and building lists
- Data-structures encoded as lists
- Serializing code and analyzing it



# Data structures

#### Data structures

When presenting each data structure we will introduce two sets of functions:

- Constructors: functions needed to build the data structure
- Accessors: functions needed to retrieve each component of the data structure. Also known as selectors.

Each example we discuss is prefaced by some unit tests. We are following a Test Driven Development methodology.



#### Constructors

- empty: creates an empty list
- list: creates a list by specifying its elements
- cons: add an element to the left-hand-side (the end) of the list



Constructor: empty

Empty lists are rendered as '()



Constructor: list

```
expression = \cdots | list list = (list expression*)
```

Function call list constructs a list with the evaluation of a possibly-empty sequence of expressions e1 up to en as values v1 up to vn which Racket prints as: '(v1 ... v2)

```
#lang racket
(list (+ 0 1) (+ 0 1 2) (+ 0 1 2 3))
(list)
```

```
$ racket list-ex1.rkt
'(1 3 6)
'()
```



Constructor: cons

```
(define 1 (list 1 2 3))
(cons 5 1)
; '(5 1 2 3)
```

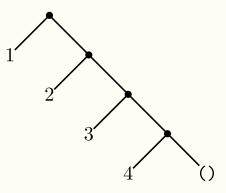
There is no direct constructor to add an element to the right, but we can implement one.



#### Constructors: cons vs list

#### cons vs list

#### Graphical representation



Textual representation

```
'(1 .
'(2 .
'(3 .
'(4 . '()))))
```



### Accessing lists

Accessor: empty?

You can test if a list is empty with function empty?. An empty list is printed as '().

```
#lang racket
(require rackunit)
(check-false (empty? (list (+ 0 1) (+ 0 1 2) (+ 0 1 2 3))))
(check-true (empty? (list)))
```



# Lists are linked-lists of pairs

Accessors: first, rest

Lists in Racket are implemented as a linked-list using pairs terminated by the empty list '().

- Function first returns the head of the list, given a nonempty list.
- Function rest returns the tail of the list, given a nonempty list.

```
(list 1 2 3 4)
```



# User data-structures

#### User data-structures

We can represent data-structures using pairs/lists. For instance, let us build a 3-D point data type.

```
(require rackunit)
(define p (point 1 2 3))
(check-true (point? p))
(check-equal? (list 1 2 3) p)
(check-equal? 1 (point-x p))
(check-equal? 2 (point-y p))
(check-equal? 3 (point-z p))
(check-true (origin? (list 0 0 0)))
(check-false (origin? p))
```



#### User data-structures

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(check-false (origin? p))
```



#### On data-structures

- We only specified immutable data structures
- The effect of updating a data-structure is encoded by **creating/copying** a data-structure
- This pattern is known as a <u>persistent data structure</u>



# Serializing code

# Quoting: a specification

Function (quote e) serializes expression e. Note that expression e is not evaluated.

- A variable x becomes a symbol 'x. You can consider a **symbol** to be a special kind of string in Racket. You can test if an expression is a symbol with function **symbol**?
- A function application  $(e_1 \cdots e_n)$  becomes a list of the serialization of each  $e_i$ .
- Serializing a (define x e) yields a list with: symbol 'define, the serialization of variable x, and the serialization of e. Serializing (define  $(x_1 \cdots x_n) e$ ) yields a list with symbol 'define followed by a nonempty list of symbols ' $x_i$  followed by serialized e.
- Serializing (lambda  $(x_1...x_n)$  e) yields a list with symbol 'lambda, followed by a possibly-empty list of symbols  $x_i$ , and the serialized expression e.
- Serializing a (cond  $(b_1 \ e_1) \cdots (b_n \ e_n))$  becomes a list with symbol 'cond followed by a serialized branch. Each branch is a list with two components: serialized expression  $b_i$  and serialized expression  $e_i$ .

### Quoting exercises:

- We can write 'term rather than (quote term)
- How do we serialize term (lambda (x) x) with quote?
- How do we serialize term (+ 1 2) with quote?
- How do we serialize term (cond [(> 10 x) x] [else #f]) with quote?
- Can we serialize a syntactically invalid Racket program?



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- How do we serialize term (cond [(> 10 x) x] [else #f]) with quote?
- Can we serialize a syntactically invalid Racket program? No! You would not be able to serialize this expression (. Quote only accepts a S-expressions (parenthesis must be well-balanced, identifiers must be valid Racket identifiers, number literals must be valid).
- Can we serialize an invalid Racket program?



## Quoting exercises:

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- Can we serialize a syntactically invalid Racket program? No! You would not be able to serialize this expression (. Quote only accepts a S-expressions (parenthesis must be well-balanced, identifiers must be valid Racket identifiers, number literals must be valid).
- Can we serialize an invalid Racket program? Yes. For instance, try to quote the term: (lambda)



### Quote example

```
#lang racket
(require rackunit)
(check-equal? 3 (quote 3)) ; Serializing a number returns the number itself
(check-equal? 'x (quote x)); Serializing a variable named x yields symbol 'x
(check-equal? (list '+ 1 2) (quote (+ 1 2))); Serialization of function as a list
(check-equal? (list 'lambda (list 'x) 'x) (quote (lambda (x) x)))
(check-equal? (list 'define (list 'x)) (quote (define (x))))
```



## Manipulating quoted terms

#### Specification

```
function-dec = ( lambda ( variable* ) term+)
```

- How do we get the parameter list?
- How do we get the body?
- What does variable\* mean?
- What does term+ mean?

#### On HW1 Q.4

- The input format of the quoted term are **precisely** described in the slides of Lecture 3
- You do **not** need to test recursively if the terms in the body of a function declaration or definition are valid.
- A list, with one symbol lambda followed by zero or more symbols, and one or more terms.



### Racket spec

HW1: Question 4

```
program = #lang racket term*
term = definition | expression
definition = basic-def | function-def
basic-def = ( define identifier expression )
function-def = ( define (variable+ ) term+ )
expression = value | variable | function-call | function-decl | · · ·
value = number | · · ·
function-call = ( expression+ )
function-dec = ( lambda ( variable* ) term+)
```

# Tips for solving HW1

### HW1: Question 4

- 1. Do all parts except lambda?, define?, and define-func?.
- 2. Write lambda?
- 3. Write define-func?
- 4. Write define?

#### More tips

- Function application is simpler than it seems
- All acceptance-tests from define-func? should pass in define?



# Being successful in CS 450

# Forum questions policy

- 1. Private questions (Discord) have the **lowest** priority
- 2. Instructor/TAs cannot comment on why a student's submission is not working
- 3. If a student lists which test-cases have been used, then the instructor/TAs can give more inputs or test cases
- 4. Private questions regarding code must always be accompanied with the URL of latest Gradescope submission
- 5. Students cannot share their solutions (partial/full) in public posts



# The final grade is given by the instructor

(not by the autograder)

We are grading the correctness of a solution

The autograder only **approximates** your grade

- Students may request for manual grading
- Grading partial solutions automatically is hard:
  - Solution may be using disallowed functions
  - Solution may be tricking the autograder system



# Tip #1: avoid fighting the autograder

- 1. It's not personal: The autograder is not against you
- 2. It's not picky: The autograder is not against one specific solution
- 3. **Correlation is not causation:** Having a colleague with the same problem as you have, does **not** imply that the autograder is wrong
- 4. Spend your time wisely: don't spend it thinking the autograder is wrong

#### Instead, discuss

- 1. Use the autograder for your benefit: submit solution to test your hypothesis
- 2. **Think before resubmitting:** try explaining your solution to someone
- 3. Ask before resubmitting: write test cases and discuss those test cases with others

10% of your grade is participation, so discuss!



### Tip #2: participate

10% of your grade is participation

Software engineering and academic life is about *communication*: you are expected to interact to solve your homework assignments.

- 1. Exercises are explained succinctly on purpose: ask questions to know more
- 2. Exercises have few test cases on purpose: share test-cases to know more

Make time in your schedule to interact



### Tip #3: time management

Work on your homework assignment incrementally

- after each class you can solve a new exercise (with few exceptions)
- when you get stuck in an exercise: (1) ask in our forum, and while you are waiting
   (2) continue working on other exercises
- don't leave everything to the weekend before submission



## Tip #4: learn to ask questions

The better your formulate a question,

The faster you will get an answer

Ask yourself

- 1. Which slides do you think the exercise relates to?
- 2. Which test-cases have you tried that counter your intuition?

Asking question

- 1. Describe the problem you are having (relate exercise and lessons)
- 2. Explain your attempts at fixing the problem (list used tests)

