Circuit switching: In when switching, before transmitting data between two nodes, the vetwork establishes a dedicated end between the mode and veserves end-be-end amechan between the mode and veserves bendwidth in each hink to along the connection. The bendwidth in each hink to along the connection the reserved connection bandwidth is wasted whenever the reserved connection bandwidth is wasted whenever the concerned wodes woot are not sending date.

Pachet switching: when one node sends date to another mode, the fonding mode breaks the data into packets, and then sands each packet seperately. Here pachet is pointed wring the destination address when a packet is received, int uses this deprivation address to determine on which Link it should forward the partiet. This in packet switching in coming packet are forwarded to ontgoing live packet by packet. Also, packet savitching was tryically uses store and forward mechanism which receives and slaves the entire padret before forwarding. In store and forward, the pachet is stored until it has fully corrived so that checksum Can be competed from Bed. Then it forward to see appropriate ontgoing link atom for delivery.

the Frequency dinition multipliains and Time division multipleasing used multipleasing are two Common multipleasing used in Circuit of the Link is partitioned into frequency bondwidth of the Link is partitioned into frequency bonds. To TDM, partitions time in a link with bonds. To TDM, partitions time in a link with a revolving each Connection the Same slot in the revolving frame.

was at a constant bit grate. Would packet switching of or circuit switching be more appropriate in this case and why?

Cae and why?

Here is statistical multiplening gain is there to came by

Here in statistical multiplening gain is there to came by

Circuit-switching is more dedivable here to came by & Suppose that all of the network resources send data Usting whents, each connection will get a constant amount of Landberdthe that matches worits CBR. But circuit switching has more overhead in terms of Engualing meded to Set up the call, so there is on argument that packet-Julitching many absorbe preferable here fince there is no call setup overhead with packet-switching, If this is an oran question, either answer avoid be correct as long on y on provide wrect reasoning. Suppose that all of the network mesource are burstythat they only oceas, onally hance bats to send. Would packet-switching or wireint-smitching be more appropriate Padret-revitching is better her beganse then are Statistian multiplening, gain - when a sense does but have don't to land, it will but be allocated bondwidth, whith packet switching the sources.
Is then availed to the use by others sources. Suppose wers show a 1 mbps him. Assume each user weeks 500kbps when transmitting, but each wer transmit only 101. of a) of circuit-sulitating is used, how warmy wers can be surported? Am: Two meny each will get a dedicated 500 hbps. b) some of packet switching is med and there one three mens, find the fraction of time during which the queue grows? Ans: 4f two users and transmit Browltaneously the imput rate is 1 mbps. Since the line is 1 mbps, there mills be no quening delay. If three usen from mit Smultaneously, then the import rate is 1.5 mb 185. Since the output rate is Imbes, the queue and even at the rate of 0.5 mbps. The peop that a wer is transmilling = 0-1. Prob that All 3 users transmit simulaneous = (0.1)= 0.001. Thus the tradion is 0.001 during which queue strucked 3 trans. Simul.

Datagram and virtual circuit Dategram: packets are injected into the subject individually and nouted independently of each other. No advance Setup is veguired. In this content, the packets ove called datagrams and the subnet is called datagram subnit. Virtual aircuit: A Connection-oriented service is offered,

a path from the source nonter to the destination grønter minst be established before any data parliet can be sent. This connection is called virtual circuit and the subret is called virtual

would subvet.

| a Societ (| uphel. | 2 + B Wet |
|---------------------------|---|--|
| execut s | Cubnet | Virtual circuit subvet |
| Issue | pategram Subnet | to occurred. |
| circuit setup | Not needed | Each packet Contains |
| Advessing. | Not needed Each packet contains the full sowice and destination address. | eisenit number |
| | | te auch virtual circuit ons requires receiper connection table space per connection |
| State information | Routers do not hold sta Information about ameets | A P CO. |
| Roubing | anch padent is nouled | ristral circuit packets is set up; all packets pollow this obstructe |
| affect of ownter failures | None, except for packets lost during the exach | that passed through the failed router are terminated. |
| Quality of Service | e) Difficult to | Can be allo caled, in advance for each esthal |
| Congestion Control | Difficult | Carry if enough resources can be allocated In advance for each without advance for each without |
| | - | (Cu o Cu o |