

# Matrix-Tree Theorem

Dr. Danny Nguyen

*Transcribed by Thomas Cohn*

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**Defn:** Given  $G = (V, E)$ , the adjacency matrix  $A = (a_{ij})$  with  $a_{ij} = \begin{cases} 1 & (v_i, v_j) \in E \\ 0 & \text{otherwise} \end{cases}$

**Defn:** The laplacian  $L = (\ell_{ij})$  is defined by  $\ell_{ij} = \begin{cases} \deg(v_i) & i = j \\ -a_{ij} & i \neq j \end{cases}$

Observe:  $\det(L) = 0$ . This is because  $L \cdot \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ \vdots \\ 0 \end{pmatrix}$ , so  $\ker L \neq 0$ .

**Defn:** For each  $1 \leq i \leq n$ , the minor  $L_{i,j}$  is the determinant of  $L$  where the  $i$ -th row and  $j$ -th column are deleted.

**Ex:**

$$L = \begin{bmatrix} 2 & -1 & -1 & 0 \\ -1 & 3 & -1 & -1 \\ -1 & -1 & 3 & -1 \\ 0 & -1 & -1 & 2 \end{bmatrix} \quad L_{11} = \det \begin{bmatrix} 3 & -1 & -1 \\ -1 & 3 & -1 \\ -1 & -1 & 2 \end{bmatrix} = 8 \quad L_4 = \det \begin{bmatrix} 2 & -1 & -1 \\ -1 & 3 & -1 \\ -1 & -1 & 3 \end{bmatrix} = 8$$

Notice that every minor of the Laplacian is the same.

**Thm:** (Matrix Tree) Given any graph  $G$ , the number of spanning trees in  $G = L_{i,i}$  for any  $1 \leq i \leq n$ . ( $L$  is the Laplacian.)

Proof: Let  $e_1, \dots, e_m$  be all the edges in  $G$ . Orient each edge  $e_i$  arbitrarily (denote the oriented edge  $\vec{e}_i$ ).

Define the incidence matrix  $N \in \mathbb{R}^{n \times m}$  by  $n_{ij} = \begin{cases} 0 & x_i \notin \vec{e}_j \\ 1 & x_i = \text{head}(\vec{e}_j) \\ -1 & x_i = \text{tail}(\vec{e}_j) \end{cases}$

Observe that  $NN^T = L$ .

Proof:  $\ell_{i,i} = r_i(N) \cdot r_i(N) = \deg(v_i)$ . For  $i \neq j$ ,  $\ell_{i,j} = r_i(N) \cdot r_j(N) = \begin{cases} -1 & (i,j) \in E \\ 0 & (i,j) \notin E \end{cases} \quad \square$

So  $\det(L_{1,1}) = \det(N_1 N_1^T)$ , where  $N_1$  is  $N$  with row 1 deleted.

**Lemma:** (Cauchy-Binet) Let  $A \in \mathbb{R}^{\ell \times m}$ ,  $B \in \mathbb{R}^{m \times \ell}$ , with  $\ell \leq m$ .

$$\det(AB) = \sum_{I=\{i_1, \dots, i_\ell\} \subset [m]} \det(A^I) \det(B^I)$$

where  $A^I$  is  $A$  with cols  $i_1, \dots, i_\ell \in I$  and  $B^I$  is  $B$  with rows  $i_1, \dots, i_\ell \in I$ . This is proved in the textbook.

So we now have  $\det(L_{1,1}) = \det(N_1 N_1^T) = \sum_{I=\{i_1, \dots, i_{n-1}\} \subset [m]} \det(N_1^I) \det((N_1^I)^T) = \sum_{I=\{i_1, \dots, i_{n-1}\} \subset [m]} \det(N_1^I)^2$ .

**Lemma:** For each  $I = \{i_1, \dots, i_{n-1}\} \subset [m]$ ,  $\det(N_1^I) = \begin{cases} 0 & \{e_{i_1}, \dots, e_{i_{n-1}}\} \text{ is not a tree} \\ \pm 1 & \{e_{i_1}, \dots, e_{i_{n-1}}\} \text{ is a tree} \end{cases}$

Proof: If  $\{e_{i_1}, \dots, e_{i_{n-1}}\}$  is disconnected, then there is a cycle in one of the connected components. Adding up the columns for that cycle and multiplying by  $\pm 1$  as needed gives us a column which is  $\vec{0}$ . So the determinant of the matrix is 0.

If  $\{e_{i_1}, \dots, e_{i_{n-1}}\}$  is a tree, then  $\exists$  some leaf  $v \neq 1$  because the first row is gone. In the row of  $v$ , there exists a single  $\pm 1$  entry; perform Laplace Expansion along that row to calculate the determinant, and recurse.  $\square$

$\square$