

Categories in Context - Ch. 1.7 Notes

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1 Functor Categories

In this section, we begin to build the intuition for working in "higher dimensions" of categories, noting some very interesting properties of natural transformations, functors, and introducing one of the most important building blocks in category theory: the functor category.

Definition 1.1 (Functor category). *For any pair of categories, there is an associated category $\mathbf{D}^{\mathbf{C}}$, whose data is as follows:*

- *Objects are functors $F : \mathbf{C} \rightarrow \mathbf{D}$*
- *Morphisms are natural transformations $\alpha : \mathbf{C} \Rightarrow \mathbf{D}$*
- *The identity transformation for a functor F is denoted $1_F : F \Rightarrow F$, whose components are defined to be $(1_F)_c := 1_{F_c}$.*
- *Along with component-wise "vertical" composition of transformations (see next lemma.).*

Composition of natural transformations may occur in more than one way, hence the distinction made above. In fact, it is this "multi-dimensional" character of their composition that belies the choice in terminology of "n-category": in a two category, the "2-morphisms" may be composed in 2 ways.

Lemma 1.2. *(vertical composition) Suppose $\alpha : F \Rightarrow G$ and $\beta : G \Rightarrow H$ are natural transformations between the functors $F, G, H : \mathbf{C} \rightarrow \mathbf{D}$. Then there is a natural transformation, syntactically denoted $\beta\alpha$ or $\beta \cdot \alpha : F \Rightarrow H$, whose components are defined as follows:*

$$(\beta \cdot \alpha)_c := \beta_c \cdot \alpha_c \tag{1}$$

Proof. The naturality of α and β implies that for any $f : c \rightarrow d$, then the following composite rectangle commutes:

$$\begin{array}{ccccc} Fc & \xrightarrow{\alpha_c} & Gc & \xrightarrow{\beta_c} & Hc \\ Ff \downarrow & & \downarrow Gf & & \downarrow Hf \\ Fd & \xrightarrow{\alpha_d} & Gd & \xrightarrow{\beta_d} & Hd \end{array}$$

It is necessary and sufficient to check that associativity and unitality of components, and this follows immediately from composition in \mathcal{D} . \square

The name "vertical composition" is suggestive: why vertical? Component-wise composition would certainly lend itself more to being called "horizontal". However, we must keep the big picture in mind. Here is the mnemonic you can use to determine which is which in your head:

- Drawing functors horizontally, a "vertically composable" pair of natural transformations fits into what we call a *pasting diagram*:

$$\begin{array}{ccc} & F & \\ & \curvearrowright & \\ C & \xrightarrow{G} & D \\ & \curvearrowleft & \\ & H & \end{array} \quad \Downarrow \begin{array}{c} \alpha \\ \beta \end{array} \quad = \quad \begin{array}{ccc} & F & \\ & \curvearrowright & \\ C & \xrightarrow{\beta \cdot \alpha} & D \\ & \curvearrowleft & \\ & H & \end{array}$$

- As the terminology suggests, there is a **horizontal composition**, which composes horizontally along pasting diagrams as defined above:

$$\begin{array}{ccc} & F & \\ & \curvearrowright & \\ C & \xrightarrow{G} & D \\ & \curvearrowleft & \\ & G & \end{array} \quad \Downarrow \alpha \quad \begin{array}{ccc} & H & \\ & \curvearrowright & \\ D & \xrightarrow{K} & E \\ & \curvearrowleft & \\ & K & \end{array} \quad = \quad \begin{array}{ccc} & HF & \\ & \curvearrowright & \\ C & \xrightarrow{KG} & E \\ & \curvearrowleft & \\ & KG & \end{array} \quad \Downarrow \beta * \alpha$$

Let's prove the lemma asserted above that this indeed defines a valid composition of transformations.

Lemma 1.3. (*horizontal composition*)

Given a pair of natural transformations $\alpha : F \Rightarrow G$ and $\beta : H \Rightarrow K$, then there is a natural transformation $\beta * \alpha : HF \Rightarrow KG$ whose component at $c \in \mathcal{C}$ is defined as the composite of the following square:

$$\begin{array}{ccc} HFc & \xrightarrow{\beta_{Fc}} & KFc \\ H\alpha_c \downarrow & \searrow (\beta * \alpha)_c & \downarrow K\alpha_c \\ HGc & \xrightarrow{\beta_{Gc}} & KGc \end{array}$$

Proof. The square defined above commutes by naturality of β applied to the component α_c . To prove that the components of $\beta * \alpha$ are natural, we must show that, given $f : c \rightarrow d$ in \mathcal{C} , the following naturality condition holds:

$$KGf \cdot (\beta * \alpha)_c = (\beta * \alpha)_d \cdot HFf \quad (2)$$

This relation can be summarized in the following diagram, and is sufficient as a proof:

$$\begin{array}{ccccc} HFc & \xrightarrow{H\alpha_c} & HGc & \xrightarrow{\beta_{Gc}} & KGc \\ HFf \downarrow & & \downarrow HGf & & \downarrow KGf \\ HFd & \xrightarrow{H\alpha_d} & HGd & \xrightarrow{\beta_{Gd}} & KGd \end{array}$$

□

1.1 Whiskering

There are some typing issues that need to be addressed in order to justify the syntax $L\beta F : LHF \Rightarrow LKF$ as seen in the book. In particular, the juxtaposition of L with β , and β with F would imply that one is "applying" L to β and β to F to produce a transformation. This is not the case. In fact there is a subtle conversion being done under the hood. In the case of $L\beta$, we must view this in terms of the application of L to components $\beta_{Hc} \in E$, such that $L\beta_{Hc} \in E$, and βF in terms of the application of components to F as it is applied to an object $\beta_{Fc}(Fc) \in D$ for a given object $c \in C$ (parens have been added for clarity). We shorthand this threading of information using the syntax $L\beta F$, and capture it in a diagram:

$$\begin{array}{ccccc} C & \xrightarrow{F} & D & \begin{array}{c} \xrightarrow{H} \\ \Downarrow \beta \\ \xrightarrow{K} \end{array} & E & \xrightarrow{L} & F \end{array}$$

Lemma 1.4. (*middle four interchange*)

Given functors and natural transformations

$$\begin{array}{ccccc} C & \xrightarrow{F} & D & \xrightarrow{J} & E \\ \Downarrow \alpha & & \Downarrow \gamma & & \\ C & \xrightarrow{G} & D & \xrightarrow{K} & E \\ \Downarrow \beta & & \Downarrow \delta & & \\ C & \xrightarrow{H} & D & \xrightarrow{L} & E \end{array}$$

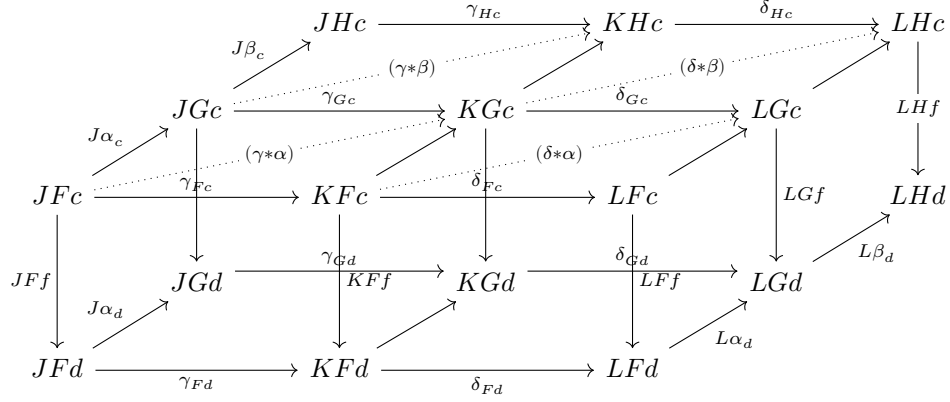
The transformations $JF \Rightarrow LH$ defined by first composing vertically and then composing horizontally equals the natural transformation defined by first composing horizontally, and then composing vertically:

$$\begin{array}{ccc} \begin{array}{ccc} C & \xrightarrow{F} & D \\ \Downarrow \beta \cdot \alpha & & \\ C & \xrightarrow{H} & D \end{array} & \begin{array}{ccc} D & \xrightarrow{J} & E \\ \Downarrow \delta \cdot \gamma & & \\ D & \xrightarrow{L} & E \end{array} & = & \begin{array}{ccc} C & \xrightarrow{JF} & E \\ \Downarrow \gamma * \alpha & & \\ C & \xrightarrow{KG} & E \\ \Downarrow \delta * \beta & & \\ C & \xrightarrow{LH} & E \end{array} \end{array}$$

Proof. (exercise 1.7.iv) We can show that the above lemma is true by showing that the following equality holds:

$$(\delta \cdot \gamma) * (\beta \cdot \alpha) = (\delta * \beta) \cdot (\gamma * \alpha) \quad (3)$$

For natural transformations, equality is defined by equality of components. That is to say $\alpha = \beta$ if and only if $\alpha_c = \beta_c$ for all $c \in C$. It suffices to show in this case, that the components yield precisely the same data for each side of the equality. Let $f : c \rightarrow d$ and behold my glory:



Note that every individual cube is a commutative cube, and all faces are commutative as well. Therefore the requisite paths through the cube are equal by naturality of their component shapes, and therefore

$$(\delta \cdot \gamma) * (\beta \cdot \alpha) = (\delta * \beta) \cdot (\gamma * \alpha) \quad (4)$$

□

1.2 2-categories

Definition 1.5. (2-category) A **2-category** is comprised of the following data:

- Objects, e.g. categories \mathcal{C}
- 1-morphisms between objects (e.g. functors $F : \mathcal{C} \rightarrow \mathcal{D}$)
- 2-morphisms between parallel 1-morphisms (e.g. natural transformations).

Such that

- The objects and 1-morphisms form a category. Identities are 1-morphisms.
- For each pair of objects \mathcal{C}, \mathcal{D} , the 1-morphisms and 2-morphisms form a category under vertical composition, with identities defined by the diagram

$$\begin{array}{ccc} & F & \\ \mathcal{C} & \begin{array}{c} \curvearrowright \\ \Downarrow 1_F \\ \curvearrowleft \end{array} & \mathcal{D} \\ & F & \end{array}$$

- There is also a category whose objects are the objects which a morphisms $\mathcal{C} \rightarrow \mathcal{D}$ is a "2-cell", which composes under the operation "horizontal composition".
- The law of middle four interchange holds

1.3 Size Matters

Caution is needed when talking about sizing with respect to functor categories. If \mathbf{C} and \mathbf{D} are small, then $\mathbf{D}^{\mathbf{C}}$ is small. However, if both are only locally small, then $\mathbf{D}^{\mathbf{C}}$ need not be (why? Discuss in the meetup - exercise 1.7.i).