EXERCISE 2.9

a. (a)

•
$$\mathbf{E}[\max(X_1, X_2)] = \sum_{n=1}^k n \times \Pr(\max(X_1, X_2) = n)$$

$$= \sum_{n=1}^k \sum_{m=1}^n \Pr(\max(X_1, X_2) = n)$$

$$= \sum_{m=1}^k \sum_{n=m}^k \Pr(\max(X_1, X_2) = n)$$

$$= \sum_{m=1}^k \Pr(\max(X_1, X_2) \ge m)$$

$$= \sum_{m=1}^k (1 - \Pr(\max(X_1, X_2) < m))$$

$$= \sum_{m=1}^k \left(1 - \frac{(m-1)^2}{k^2}\right)$$

$$= k - \frac{1}{k^2} \times \frac{k(k-1)(2k-1)}{6}$$

$$= \frac{4k^2 + 3k - 1}{6k}$$

$$\begin{split} \mathbf{E}[\min(X_1, X_2)] &= \sum_{n=1}^k n \times \Pr(\min(X_1, X_2) = n) \\ &= \sum_{n=1}^k \sum_{m=1}^n \Pr(\min(X_1, X_2) = n) \\ &= \sum_{m=1}^k \sum_{n=m}^k \Pr(\min(X_1, X_2) = n) \\ &= \sum_{m=1}^k \Pr(\min(X_1, X_2) \ge m) \\ &= \sum_{m=1}^k \frac{(k+1-m)^2}{k^2} \\ &= \frac{1}{k^2} \times \frac{k(k+1)(2k+1)}{6} \\ &= \frac{2k^2 + 3k + 1}{6k} \end{split}$$

b. (b)

 • From (a), $\mathbf{E}[\max(X_1,X_2)] + \mathbf{E}[\min(X_1,X_2)] = k+1$

.
$$\begin{split} \mathbf{E}[X_1] &= \mathbf{E}[X_2] = \sum_{n=1}^k n \times \Pr(X_1 = n) \\ &= \sum_{n=1}^k n \times \frac{1}{k} \\ &= \frac{1}{k} \times \frac{k(k+1)}{2} \\ &= \frac{k+1}{2} \end{split}$$

- $\mathbf{E}[X_1] + \mathbf{E}[X_2] = k + 1$
- $\div \mathbf{E}[\max(X_1, X_2)] + \mathbf{E}[\min(X_1, X_2)] = \mathbf{E}[X_1] + \mathbf{E}[X_2]$

c. (c)

- $\max(X_1, X_2) + \min(X_1, X_2) = X_1 + X_2$
- $\mathbf{E}[\max(X_1, X_2) + \min(X_1, X_2)] = \mathbf{E}[X_1 + X_2]$ = $\mathbf{E}[\max(X_1, X_2)] + \mathbf{E}[\min(X_1, X_2)] = \mathbf{E}[X_1] + \mathbf{E}[X_2]$

2.15

- Let X be the number of total coin flips until the k-th head
- Let X_i be the number of coin flips after the i-th head until the i+1-th head. (Excluding the coin flip that results in i-th head, and including the coin flip that results in i+1-th head)
 - X_i all follows geometric distribution with parameter p, $\mathbf{E}[X_i] = \frac{1}{p}$
 - $X = \sum_{i=0}^{k-1} X_i$
- $\mathbf{E}[X] = \sum_{i=0}^{k-1} \mathbf{E}[X_i] = \frac{k}{p}$
- The expected number of coin flips until the k-th head is $\frac{k}{p}$

2.18

- After the *n*-th item,
- In order for the k-th item to be stored in memory
 - It should have replaced the item when it appeared $(\frac{1}{k})$
 - And it should not have been replaced afterwards $\left(\prod_{i=k+1}^n \frac{i-1}{i}\right)$
- Therefore, the overall probability of k-th item being stored in memory is

$$\frac{1}{k} \times \prod_{i=k+1}^n \frac{i-1}{i} = \frac{1}{k} \times \frac{k}{k+1} \times \dots \times \frac{n-1}{n} = \frac{1}{n}$$

- ullet There is an equal chance for each of n items encountered to be the item stored in memory
- : The desired property is achieved

2.22

a. BubbleSort mutation count

- Proposition: The number of mutations in the bubble sort is equal to the total number of inverted pair in the initial permutation
 - Proposition: A bubble sort mutation decreases the number of inverted pair by 1
 - Let
 - The permutation before the mutation be $a_1,...,a_k,a_{k+1},...,a_n$
 - And the permutation after the mutation be $b_1,...,b_k,b_{k+1},...,b_n$
 - Let's say the mutation swapped a_k and a_{k+1}
 - $\bullet \ \, \forall i \neq k, k+1, b_i = a_i, b_k = a_{k+1}, b_{k+1} = a_k$
 - After the permutation:
 - $\left(b_{k},b_{k+1}\right)$ is not inverted, while $\left(a_{k},a_{k+1}\right)$ was inverted
 - $\forall i, j \text{ such that } i, j \neq k, k+1 \land i < j$, the invertedness of the pair (b_i, b_j) is equal to (a_i, a_j)
 - $\forall i < k$, the invertedness of the pair (b_i, b_k) is equal to (a_i, a_{k+1}) , and the invertedness of the pair (b_i, b_{k+1}) is equal to (b_i, b_k) , leaving the total number of inverted pairs the same.
 - Likewise, $\forall j > k+1$, the invertedness of the pair (b_k, b_j) is equal to (a_{k+1}, a_j) , and the invertedness of the pair (b_{k+1}, b_j) is equal to (a_k, a_j) , leaving the total number of inverted pairs the same.
 - Therefore, the number of inverted pairs is descreased by exactly one after the mutation.
 - Since the result of the bubble sort is a sorted array with no inverted pairs,
 - The number of inverted pairs in the initial permutation is the number of mutations that happened in the process.

b. Expected Number of Inverted Pairs in Any Permutation

- We can compute the expected number of bubble sort mutation for any permutation by computing the expected number of inverted pair in any permutation
- For any permutation $a_1,...,a_n$, there exists a permutation where the order of elements are completely reversed $b_1=a_n,...,b_n=a_1$
 - Exactly one of (a_i, a_j) and (b_i, b_j) is a inverted pair
 - Therefore, among the two permutations, there is exactly $\binom{n}{2}=\frac{n(n-1)}{2}$ inverted pairs
 - The average number of inverted pair among the two permutations is $\frac{n(n-1)}{4}$
- Since every permutation can be coupled like that, the average number, or the expected number of inverted pairs for every permutation is $\frac{n(n-1)}{4}$
- The expected number of bubble sort mutation for any permutation is also $\frac{n(n-1)}{4}$

EXERCISE 2.27

$$\mathbf{E}[X] = \sum_{k=1}^{\infty} k \Pr(X = k)$$
$$= \sum_{k=1}^{\infty} k \frac{6}{\pi^2} \frac{1}{k^2}$$
$$= \frac{6}{\pi^2} \sum_{k=1}^{\infty} \frac{1}{k}$$

• The expected value of X diverges

EXERCISE 2.32

a. (a)

• For $i \leq m$, trivially $\Pr(E_i) = 0$

- For i>m, In order for E_i to happened

• The *i*-th candidate must be the best $(\frac{1}{n})$

- We must not hire m+1-th to i-1-th candidate

• The best candidate among the first i-1 candidate should have been in the first m candidate $\left(\frac{m}{i-1}\right)$

$$\Pr(E_i) = \begin{cases} 0 \text{ if } i \leq m \\ \frac{1}{n} \times \frac{m}{i-1} \text{ otherwise} \end{cases}$$

$$\Pr(E) = \sum_{i=1}^{n} \Pr(E_i) = \frac{m}{n} \sum_{i=1}^{n} \frac{1}{i-1}$$

b. (b)

$$\begin{split} & \forall j \text{ such that } i-1 < j < i, \frac{1}{j} < \frac{1}{i-1} < \frac{1}{j-1} \\ & \sum_{i=m+1}^n \int_{i-1}^i \frac{1}{j} \mathrm{d}j < \sum_{i=m+1}^n \int_{i-1}^i \frac{1}{i-1} \mathrm{d}j < \sum_{i=m+1}^n \int_{i-1}^i \frac{1}{j-1} \mathrm{d}j \\ & \int_m^n \frac{1}{j} \mathrm{d}j < \sum_{i=m+1}^n \frac{1}{i-1} < \int_m^n \frac{1}{j-1} \mathrm{d}j \\ & [\ln|j] \rfloor_m^n < \sum_{i=m+1}^n \frac{1}{i-1} < [\ln|j-1] \rfloor_m^n \\ & \ln n - \ln m < \sum_{i=m+1}^n \frac{1}{i-1} < \ln(n-1) - \ln(m-1) \\ & \frac{m}{n} (\ln n - \ln m) < \Pr(E) < \frac{m}{n} (\ln(n-1) - \ln(m-1)) \end{split}$$

c. (c)

•
$$\frac{\mathrm{d}}{\mathrm{d}m} \left(\frac{m}{n} (\ln n - \ln m) \right) = \frac{1}{n} (\ln n - \ln m) - \frac{m}{n} \frac{1}{m}$$
$$= \frac{1}{n} (\ln n - \ln m - 1) = 0$$
$$\Leftrightarrow m = \frac{n}{e}$$

- + $\frac{m}{n}(\ln n \ln m)$ is maximized when $m = \frac{n}{e}$
- For this m,
- $\Pr(E) > \frac{m}{n}(\ln n \ln m) = \frac{1}{e}$