**Spring 2019 COMP 3511 Homework Assignment #4**

**Handout Date: April 24, 2019 Due Date: May 8, 2019**

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**Please read the following instructions carefully before answering the questions:**

* You should finish the homework assignment **individually**.
* There are total of **4** questions.
* When you write your answers, please try to be precise and concise.
* Fill in your name, student ID, email and Section number at the top of each page.
* Please fill in your answers in the space provided.
* **Homework Collection:** the homework is submitted to **assignment #4** on **CASS**

1. [20 points] Multiple choices
   1. Which of the following is true for compaction?

A) It can be done at assembly, load, or execution time.

B) It is used to solve the problem of internal fragmentation.

C) It cannot shuffle memory contents.

D) It is possible only if relocation is dynamic and done at execution time.

**Answer**: \_D\_

* 1. Assume a system has a TLB hit ratio of 90%. It requires 15 nanoseconds to access the TLB, and 85 nanoseconds to access main memory. What is the effective memory access time in nanoseconds for this system?

A) 108.5

B) 100

C) 22

D) 176.5

**Answer**: \_A\_

* 1. Consider a logical address with 18 bits used to represent an entry in a page table. How many entries are in the conventional page table?

A) 262144

B) 1024

C) 1048576

D) 18

**Answer**: \_A\_

* 1. Given the logical address 0xAEF9 (in hexadecimal) with a page size of 256 bytes, what is the page number?

A) 0xAE

B) 0xF9

C) 0xA

D) 0x00F9

**Answer**: \_A\_

* 1. Which of the following is a benefit of allowing a program that is only partially in memory to execute?

A) Programs can be written to use more memory than is available in physical memory.

B) CPU utilization and throughput is increased.

C) Less I/O is needed to load or swap each user program into memory.

D) All of the above

**Answer**: \_D\_

* 1. Given the reference string of page accesses: 1 2 3 4 2 3 4 1 2 1 1 3 1 4 and a system with three frames, what is the final configuration of the three frames after the LRU algorithm is applied?

A) 1, 3, 4

B) 3, 1, 4

C) 4, 1, 2

D) 1, 2, 3

**Answer**: \_B\_

* 1. \_\_\_\_\_ allows the parent and child processes to initially share the same pages, but when either process modifies a page, a copy of the shared page is created.

A) copy-on-write

B) zero-fill-on-demand

C) memory-mapped

D) virtual memory fork

**Answer**: \_A\_

* 1. Consider a disk queue holding requests to the following cylinders in the listed order: 116, 22, 3, 11, 75, 185, 100, 87. Using the SSTF scheduling algorithm, what is the order that the requests are serviced, assuming the disk head is at cylinder 88 and moving upward through the cylinders?

A) 116 - 22 - 3 - 11 - 75 - 185 - 100 - 87

B) 100 - 116 - 185 - 87 - 75 - 22 - 11 - 3

C) 87 - 75 - 100 - 116 - 185 - 22 - 11 - 3

D) 100 - 116 - 185 - 3 - 11 - 22 - 75 - 87

**Answer**: \_C\_

* 1. Which of the following disk head scheduling algorithms does not take into account the current position of the disk head?

A) FCFS

B) SSTF

C) SCAN

D) LOOK

**Answer**: \_A\_

* 1. A volume control block \_\_\_\_.

A) can contain information needed by the system to boot an operating system from that partition

B) is a directory structure used to organize the files

C) contains many of the file's details, including file permissions, ownership, size, and location of the data blocks

D) contains information such as the number of blocks in a partition, size of the blocks, and free-block and FCB count and pointers

**Answer**: \_D\_

1. [20 points] Memory management
   1. Consider a logical address space of 32 pages of 1KB each, mapped onto a physical memory of 128 frames. How many bits are there in the logical address and in the physical address, respectively? (4 points)

**Answer**:

Logical address = 32 \* 1000 = 2^5 \* 2^10 = 2^15 = 15bits

Physical address = 2^15 / 128 = 2^15 / 2^7 = 2^8 = 8bits

* 1. Suppose an OS has a 24-bit virtual address and a 20-bit physical address, and a 4KB page size. How many entries are there in the page table? What is the size of the physical memory? (4 points)

**Answer**:

2^24 / 2^12 = 2^12 entries

Size of physical memory = 2^20 = 1MB

* 1. Page translation. In a 32-bit machine we subdivide the virtual address into 3 segments as follows:

page number page offset

|  |  |  |
| --- | --- | --- |
| 8-bit | 12-bit | 12-bit |

We use a two-level page table (in memory) such that the first 8 bits of an address is an index into the first level page table and the next 12 bits are an index into a second level page table. Each page table entry is 32 bits in size. (12 points)

1. What is the page size in such a system? (2 points)

**Answer**:

2^12 = 4KBs

1. How many entries are in the 1st level page table? How many entries are in one 2nd level page table? (2 points)

**Answer**:

1st level = 2^8 entries

2nd level = 2^12 entries

1. How much space is occupied in memory by the page tables for a process that has 256MB of actual virtual address space allocated? Show your work with detailed explanation. (8 points)

**Answer**:

256MB = 2^28

2^28 / 2^12 = 2^16 pages = 2^16 entries

we need 16 second level and 2^12 entries each, that is 64KB to store 256MB.

1. [30 points] Virtual Memory
   1. The following page table illustrates a system with 12-bit virtual and physical addresses and 256-byte pages. Free page frames are to be allocated in the order 9, F, D. A dash for a page frame indicates that the page is not in memory. (4 points)

|  |  |
| --- | --- |
| Page | Page-Frame |
| 0 | 0x4 |
| 1 | 0xB |
| 2 | 0xA |
| 3 | - |
| 4 | - |
| 5 | 0x2 |
| 6 | - |
| 7 | 0x0 |
| 8 | 0xC |
| 9 | 0x1 |

Convert the following virtual addresses to their equivalent physical addresses in hexadecimal. All numbers are given in hexadecimal. In the case of a page fault, you must use one of the free frames to update the page table and resolve the logical address to its corresponding physical address. (a) 0x2A1, (b) 0x4E6, (c) 0x94A, and (d) 0x316

**Answer**:

* + 1. 0xAA1
    2. 0XFE6
    3. 0x14A
    4. 0x916
  1. Suppose a replacement policy works as follows: it examines each page periodically and discards a page if it has not been used since the previous examination. How can you implement this algorithm? What is the advantage and disadvantage of such a replacement policy compared with LRU or second-chance replacement? (6 points) [Hint: use a reference bit]

**Answer**:

* 1. Suppose demand-paging is used. Memory-access time is 100 nanoseconds. It takes 8 milliseconds to service a page fault if an empty frame is available and if the replaced page is not modified, or it takes 20 milliseconds if the replaced page is modified. Assume that the page to be replaced is modified 70 percent of the time. What is the maximum acceptable page-fault rate for an effective access time of no more than 200 nanoseconds? (5 points)

**Answer**:

(1-p)\*100 + p\*(0.7\*20000000+0.3\*8000000)

=100 + p \* 16399900

p\*16399900 < 200 => p < 0.0000012195

* 1. Consider the following page reference string:

1, 2, 3, 4, 2, 1, 5, 6, 2, 1, 2, 3, 7, 6, 3, 2, 1, 2, 3, 6.

How many page faults would occur for the following replacement algorithms, assuming one, two, three, four, five, six, and seven frames? Remember that all frames are initially empty, so your first unique pages will cost one fault each. (1) LRU replacement (2) FIFO replacement (3) Optimal replacement. (15 points)

**Answer**:

|  |  |  |  |
| --- | --- | --- | --- |
| Number of frames | LRU | FIFO | Optimal |
| 1 | 20 | 20 | 20 |
| 2 | 18 | 18 | 15 |
| 3 | 16 | 16 | 11 |
| 4 | 10 | 14 | 9 |
| 5 | 8 | 10 | 7 |
| 6 | 7 | 10 | 7 |
| 7 | 7 | 7 | 7 |

1. [30 points] Secondary storage and file systems
   1. Please name three advantages that NVM such as SSDs has over HDDs. (3 points)

**Answer**:

No rotation latency

No noise

More reliable since no moving parts

* 1. Can a RAID level 1 organization achieve better performance for read requests than a RAID level 0 organization (with nonredundant striping of data)? If so, how? (4 points)

**Answer**:

* 1. Disk scheduling. Suppose that a disk drive has 5,000 cylinders, numbered 0 to 4999. The drive is currently serving a request at cylinder 143, and the previous request was at cylinder 125. The queue of pending requests, in FIFO order, is:

86, 1470, 913, 1774, 948, 1509, 1022, 1750, 130

Starting from the current head position (143), what is the total distance (in cylinders) that the disk arm moves to satisfy all the pending requests for each of the following disk-scheduling algorithms? (12 points)

1. FIFO
2. SSTF
3. SCAN
4. LOOK
5. C-SCAN
6. C-LOOK

**Answer**:

1. 7081 cylinders
2. 1745 cylinders
3. 1917 cylinders
4. 1745 cylinders
5. 4987 cylinders
6. 3363 cylinders
   1. The open-file table is used to maintain information about files that are currently open. The operating system usually maintains one table (system-wide file open table) that contains references to files that are currently being accessed by all users, instead of a separate table for each user. Please explain why. [Hint: consider two processes access the same file] (5 points)

**Answer**:

* 1. Consider a file system that uses *inodes* to represent files. Disk blocks are 8 KB in size, and a pointer to a disk block requires 4 bytes. This file system has 12 direct disk blocks, as well as single, double, and triple indirect disk blocks. What is the maximum size of a file that can be stored in this file system? (6 points)

**Answer**:

pointer of each disk block = 2^13 / 2^2 = 2^11 = 2048

max size = 12 \* 8 + 2048 \* 8 + 2048 \* 2048 \* 8 + 2048 \* 2048 \* 2048 \* 8 = 64TB