



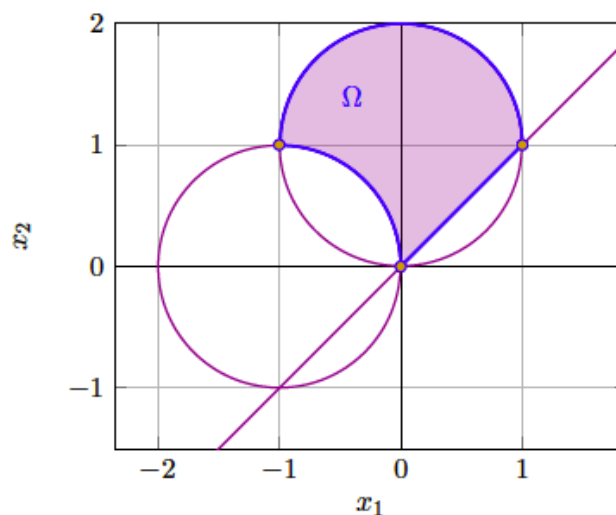
## MAT 3007 – Optimization

### Final Exam – Sample

#### Exercise 1 (KKT Conditions and Constrained Problems):

(20 points)

a) The following sketch shows the feasible region  $\Omega$



b) Based on our sketch in part a), we immediately see that the feasible region  $\Omega$  is not convex. Hence, problem (1) is not a convex program.

c) We have  $g_1(\bar{x}) = g_2(\bar{x}) = g_3(\bar{x}) = 0$  and hence, it follows  $\mathcal{A}(\bar{x}) = \{1, 2, 3\}$ .

d) We first calculate several derivatives:

$$\nabla f(x) = \begin{bmatrix} x_2 - 2x_1x_2^2 \\ \frac{1}{1+x_2} + x_1 - 2x_1^2x_2 \end{bmatrix}, \quad \nabla g_1(x) = \begin{bmatrix} 2x_1 \\ 2x_2 - 2 \end{bmatrix},$$

$$\nabla g_2(x) = \begin{bmatrix} -2x_1 - 2 \\ -2x_2 \end{bmatrix}, \quad \nabla g_3(x) = \begin{bmatrix} 1 \\ -1 \end{bmatrix}.$$

Inserting  $x = \bar{x}$ , we obtain

$$\nabla f(\bar{x}) = \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \quad \nabla g_1(\bar{x}) = \begin{bmatrix} 0 \\ -2 \end{bmatrix}, \quad \nabla g_2(\bar{x}) = \begin{bmatrix} -2 \\ 0 \end{bmatrix}, \quad \nabla g_3(\bar{x}) = \begin{bmatrix} 1 \\ -1 \end{bmatrix}.$$

We need to find  $\lambda \in \mathbb{R}^3$  such that

$$\nabla f(\bar{x}) + \nabla g_1(\bar{x})\lambda_1 + \nabla g_2(\bar{x})\lambda_2 + \nabla g_3(\bar{x})\lambda_3 = \begin{bmatrix} -2\lambda_2 + \lambda_3 \\ 1 - 2\lambda_1 - \lambda_3 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}.$$

Therefore, we have  $\lambda_3 = 1 - 2\lambda_1$  and  $\lambda_2 = \frac{1}{2} - \lambda_1$ . Since the multiplier need to be nonnegative, we need to have  $\lambda_1 \leq \frac{1}{2}$ . Furthermore, due to  $\mathcal{A}(\bar{x}) = \{1, 2, 3\}$ , the complementarity conditions are automatically satisfied. Thus,  $\bar{x}$  is KKT point and all associated multiplier are given by  $(\lambda, \frac{1}{2} - \lambda, 1 - 2\lambda)$  with  $\lambda \in [0, \frac{1}{2}]$ . This also shows that the multiplier is not unique in this case.

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**Exercise 2 (Convexity):**

(15 points)

Consider the function  $f(x, y) = -\log(x + e)\log(y + e)$  defined on the region  $\Omega := \mathbb{R}_+^2 = \{(x, y) : x \geq 0, y \geq 0\}$ . (Here,  $\log$  denotes the natural logarithm and  $e$  is Euler's number).

- For fixed  $y$ , is the mapping  $f(x, y)$  a convex function of  $x$  for  $x \geq 0$ ? Explain your answer.
- For fixed  $x$ , is the mapping  $f(x, y)$  a convex function of  $y$  for  $y \geq 0$ ? Explain your answer.
- Is  $f$  a convex function of  $(x, y)$  on  $\Omega$ ? Explain your answer!

**Solution :**

- We have  $f_x(x, y) = -\frac{\log(y+e)}{x+e}$  and  $f_{xx}(x, y) = \frac{\log(y+e)}{(x+e)^2}$ . For  $y \geq 0$ , we have  $\log(y+e) \geq 1$ . Hence  $f_{xx}(x, y) \geq 0$  for all  $x$ . This shows convexity of  $x \mapsto f(x, y)$ .
- Similarly, we have  $f_y(x, y) = -\frac{\log(x+e)}{y+e}$  and  $f_{yy}(x, y) = \frac{\log(x+e)}{(y+e)^2}$ . For  $x \geq 0$ , it holds that  $\log(x+e) \geq 1$ . Hence  $f_{yy}(x, y) \geq 0$  for all  $y$ . This establishes convexity of  $y \mapsto f(x, y)$ .
- It holds that

$$\nabla f(x, y) = \begin{pmatrix} -\frac{\log(y+e)}{x+e} \\ -\frac{\log(x+e)}{y+e} \end{pmatrix} \quad \text{and} \quad \nabla^2 f(x, y) = \begin{pmatrix} \frac{\log(y+e)}{(x+e)^2} & -\frac{1}{(x+e)(y+e)} \\ -\frac{1}{(x+e)(y+e)} & \frac{\log(x+e)}{(y+e)^2} \end{pmatrix}.$$

The trace of the Hessian is positive and its determinant is given by

$$\frac{\log(x+e)\log(y+e)}{(x+e)^2(y+e)^2} - \frac{1}{(x+e)^2(y+e)^2}$$

This expression is also nonnegative due to  $\log(x+e) \geq 1$  and  $\log(y+e) \geq 1$ . Hence, the eigenvalues of  $\nabla^2 f$  need to be nonnegative which implies that  $f$  is convex on  $\mathbb{R}_+^2$ .

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**Exercise 3 (Integer Programming Formulation):**

(15 points)

A company wishes to put together an academic “package” for an executive training program. The package will consist of 6 courses. There are 4 fields and the 6 courses must cover all the 4 fields. There are 3 colleges, each offering one course in each of the 4 fields. The tuition (basic charge) assessed when at least one course is taken, at college  $j$  is  $T_j$  (independent of the number of courses taken). Moreover, each college imposes an additional charge (covering course materials, instructional aids, and so forth) for each course, the charge for taking course  $i$  (course in the field  $i$ ) at college  $j$  is  $c_{ij}$ . Formulate an integer program that will provide the company with the minimum amount it must spend to meet the requirements of the program.

**Solution :** Let  $y_j \in \{0, 1\}$  denote whether to take any course in college  $j$  and let  $x_{ij} \in \{0, 1\}$  denote whether to take course  $i$  and college  $j$ .

Then the integer programming formulation for this problem is:

$$\begin{aligned}
 & \text{minimize} && \sum_{j=1}^3 T_j y_j + \sum_{j=1}^3 \sum_{i=1}^4 c_{ij} x_{ij} \\
 & \text{subject to} && \sum_{j=1}^3 x_{ij} \geq 1 && \forall i = 1, \dots, 4 \\
 & && \sum_{j=1}^3 \sum_{i=1}^4 x_{ij} = 6 \\
 & && x_{ij} \leq y_j && \forall i, j \\
 & && x_{ij}, y_j \in \{0, 1\}.
 \end{aligned}$$


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**Exercise 4 (Branch-and-Bound Algorithm):**

(20 points)

Consider the following integer program:

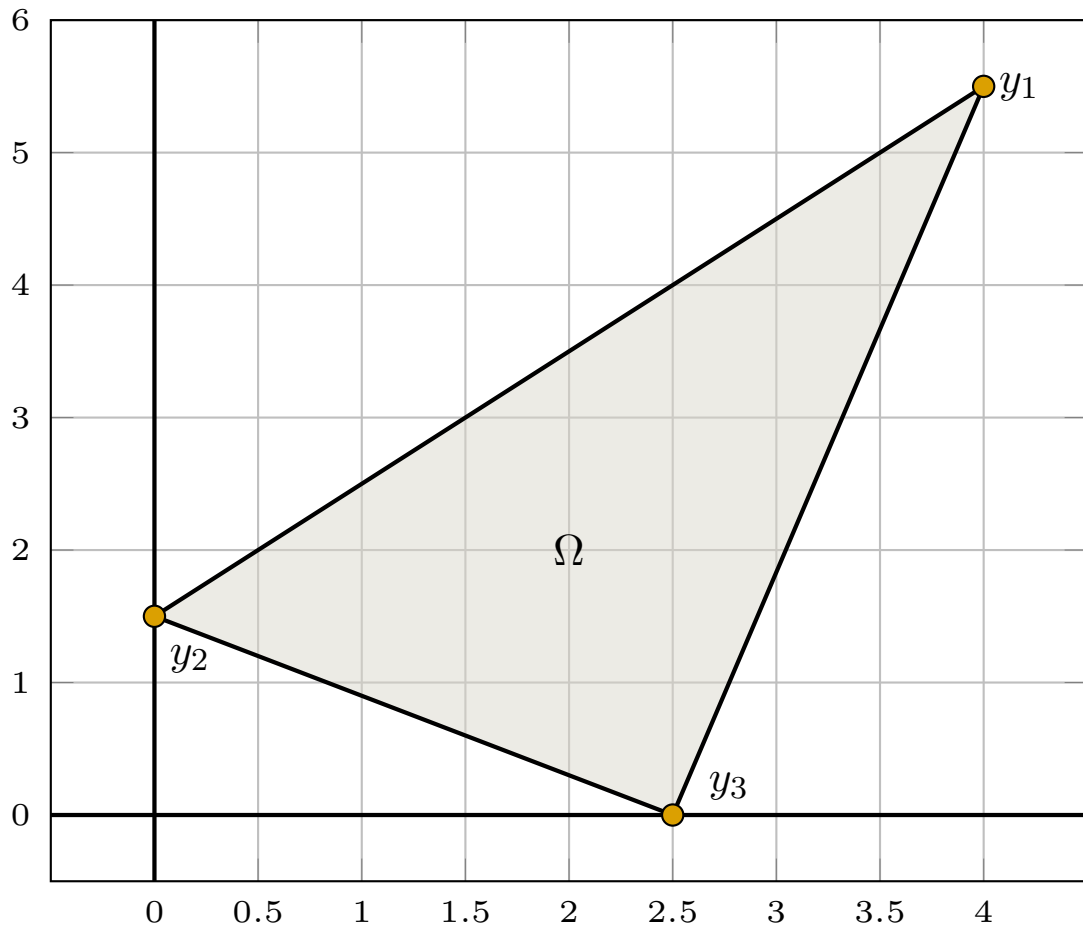
$$\begin{aligned}
 & \text{maximize} && x_1 && - && x_2 \\
 & \text{subject to} && -x_1 && + && x_2 && \leq && 1.5 \\
 & && -6x_1 && - && 10x_2 && \leq && -15 \\
 & && 22x_1 && - && 6x_2 && \leq && 55 \\
 & && x_1, && && x_2 && \in && \mathbb{Z}.
 \end{aligned}$$

Use the branch-and-bound method to solve the problem. Draw the branch-and-bound tree and mark the results on each node.

**Hint:** In order to solve the LP relaxations you can use a graphical approach or check the corresponding extreme points. The following sketch shows the feasible set

$$\Omega := \{x \in \mathbb{R}^2 : -x_1 + x_2 \leq 1.5, -6x_1 - 10x_2 \leq -15, 22x_1 - 6x_2 \leq 55\}.$$

The extreme points of  $\Omega$  are given by  $y_1 = (4, 5.5)^\top$ ,  $y_2 = (0, 1.5)^\top$ , and  $y_3 = (2.5, 0)^\top$ .



**Solution :** We denote the original problem by (S0). We first solve the LP relaxation of (S0) which is  $\max_{x \in \Omega} f(x)$ ,  $f(x) := x_1 - x_2$ . Using the sketch of  $\Omega$  or by checking the objective function values

$$f(y_1) = -1.5, \quad f(y_2) = -1.5, \quad f(y_3) = 2.5,$$

it is easy to see that  $y_3$  is the optimal solution. Since the optimal value needs to be an integer number, this implies that the optimal objective function value is bounded by 2.

We branch on  $x_1$ . We consider the two branches:

- (S1):  $x_1 \leq 2$ .
- (S2):  $x_1 \geq 3$ .

For (S1), we calculate the two new extreme points  $x_1 = 2$  and  $-12 - 10x_2 = -15$  and  $-2 + x_2 = 1.5$ . This gives  $(2, 0.3)^\top$  and  $(2, 3.5)^\top$  with function values 1.7 and  $-1.5$ , respectively.

We need to further branch on  $x_2$ . We consider the two branches:

- (S3):  $x_2 \leq 0$ .

- (S4):  $x_2 \geq 1$ .

Problem (S3) is infeasible. For (S4), we need to additionally the two new extreme points  $(2, 1)^\top$  and  $x_2 = 1, -6x_1 - 10 = -15$  which gives  $(\frac{5}{6}, 1)^\top$ . The corresponding objective function values are 1 and  $-\frac{1}{6}$ . Hence, the optimal of (S4) is  $(2, 1)^\top$  with optimal value 1. This is an integer solution and we obtain the lower bound 1 for the entire problem.

For (S2), using the sketch, the optimal solution must satisfy  $x_1 = 3$  and  $66 - 6x_2 = 55$ . This gives  $(3, \frac{11}{6})^\top$  with optimal value  $\frac{18-11}{6} = \frac{7}{6} < 2$ .

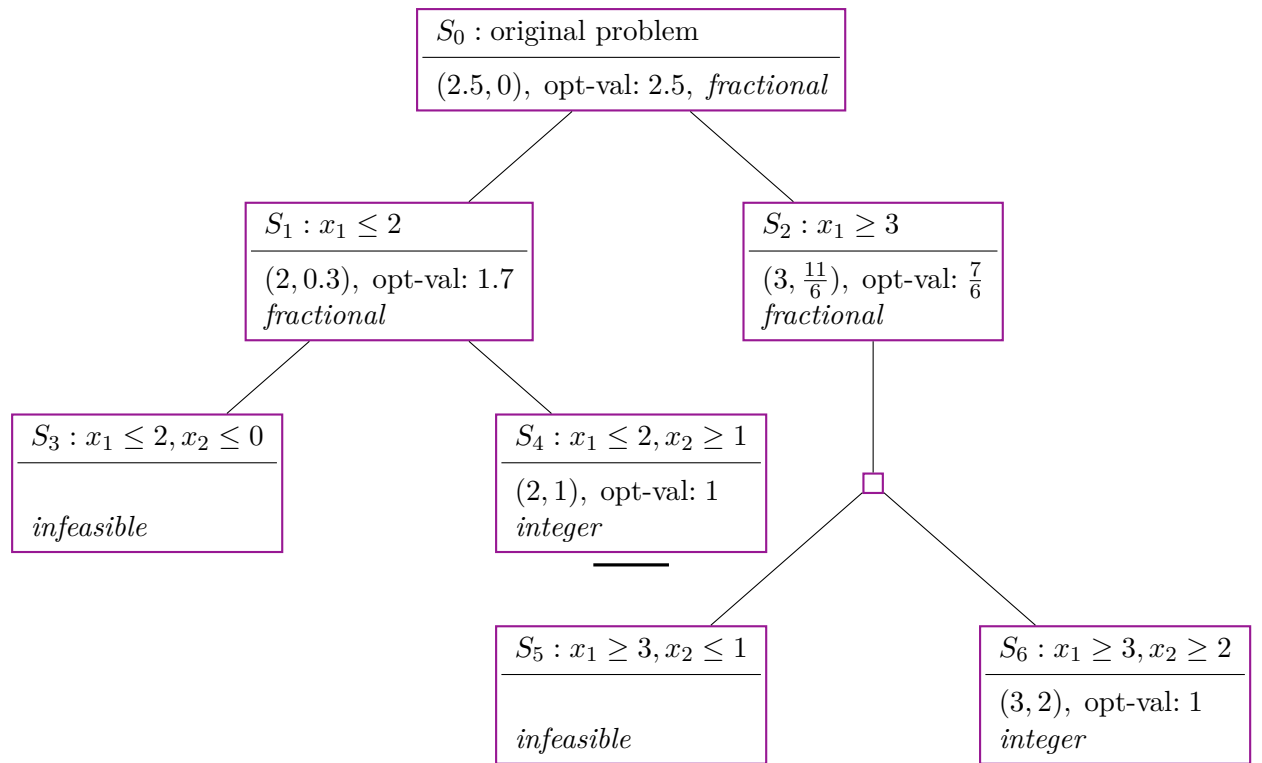
We can infer that the branch (S2) does not contain an optimal integer solution with function value greater than 1. We can stop here with the optimal solution  $x^* = (2, 1)^\top$ .

(We can continue branching:

- (S5):  $x_2 \leq 1$ .
- (S6):  $x_2 \geq 2$ .

Problem (S5) is infeasible. For (S6), using the sketch, the optimal solution must satisfy  $x_1 = 3$  and  $x_2 = 2$ . This is an integer solution with optimal value 1. Hence, the problem has the two global solutions  $(2, 1)^\top$  and  $(3, 2)^\top$ .)

A complete picture of the procedure given by:



**Exercise 5 (Algorithms for Unconstrained Problems):**

(15 points)

We consider the unconstrained minimization problem

$$\min_{x \in \mathbb{R}^3} f(x) := \frac{1}{2}x_1^4 + (x_1^2 - 1)x_2^2 + 2x_3 + x_3^2.$$

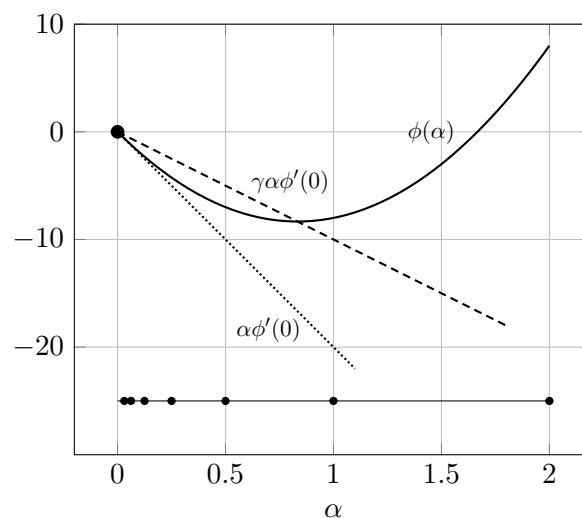
The gradient and Hessian of  $f$  are given by (*you don't need to verify this*):

$$\nabla f(x) = \begin{pmatrix} 2x_1^3 + 2x_1x_2^2 \\ 2(x_1^2 - 1)x_2 \\ 2 + 2x_3 \end{pmatrix}, \quad \nabla^2 f(x) = \begin{pmatrix} 6x_1^2 + 2x_2^2 & 4x_1x_2 & 0 \\ 4x_1x_2 & 2(x_1^2 - 1) & 0 \\ 0 & 0 & 2 \end{pmatrix}.$$

We want to apply Newton's method and the gradient descent method with backtracking to solve the problem  $\min_x f(x)$ . We choose the initial point  $x^0$  and the Armijo parameter as follows:

$$x^0 = (0, 1, 1)^\top, \quad \gamma = 0.5, \quad \sigma = 0.5.$$

- Compute the Newton direction  $d_n^0$  and the step  $x_n^1 = x^0 + d_n^0$ . Is  $d_n^0$  a descent direction of  $f$  at  $x^0$ ?
- We now choose  $d_g^0 = -\nabla f(x^0)$  and set  $\phi(\alpha) := f(x^0 + \alpha d_g^0) - f(x^0)$ . Compute the gradient iterate  $x_g^1$  and the stepsize  $\alpha_0$  using backtracking and the following plot:



- Is the function  $f$  coercive?

**Solution :**

- The Newton direction is given by

$$d_n^0 = -\nabla^2 f(x^0)^{-1} \nabla f(x^0) = - \begin{pmatrix} 6 & 0 & 0 \\ 0 & -2 & 0 \\ 0 & 0 & 2 \end{pmatrix}^{-1} \begin{pmatrix} 0 \\ -2 \\ 4 \end{pmatrix} = - \begin{pmatrix} 0 \\ 1 \\ 2 \end{pmatrix}.$$

Hence, the Newton step is  $x_n^1 = (0, 0, -1)^\top$ . Moreover, we have  $\nabla f(x^0)^\top d_n^0 = -(8-2) < 0$  which shows that  $d_n^0$  is a descent direction.

- b) The Armijo condition is satisfied whenever  $\phi(\alpha) \leq \gamma\alpha\phi'(0)$ . Since  $\sigma = \frac{1}{2}$ , this is the case for  $\alpha_0 = \frac{1}{2}$ . We then obtain  $x_g^1 = x^0 + \alpha_0 d_g^0 = (0, 2, -1)^\top$ .
- c) Consider the family of points  $x_\alpha = (0, \alpha, 0)^\top$  with  $\alpha \in \mathbb{R}$ . It follows  $f(x_\alpha) = -\alpha^2$  and hence, we obtain  $\|x_\alpha\| \rightarrow \infty$  and  $f(x_\alpha) \rightarrow -\infty$  as  $\alpha \rightarrow \infty$ . This implies that  $f$  is not coercive.

**Exercise 6 (True & False):**

(15 points)

State whether each of the following statements are *True* or *False*. For each part, only your answer, which should be either *True* or *False*, will be graded. Explanations will not be read.

- a) We consider a nonlinear program with  $f(x) = x_1^2 + x_2^2 - 2x_3^4$ ,  $g(x) = \|x\|^2 - 5$ , and  $h(x) = x_1^2 + x_2 + 2x_3$ . (The feasible set is given by  $\Omega = \{x : g(x) \leq 0, h(x) = 0\}$ ). This problem possesses a global solution.
- b) Consider the nonlinear program  $\min_{x \in \Omega} f(x)$  with linear inequality constraints  $\Omega := \{x : Ax \leq b\}$ . Let  $x^*$  be a local solution of this problem, then  $x^*$  satisfies the KKT conditions.
- c) Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  be a convex function and let us apply the gradient descent method with backtracking to solve  $\min_x f(x)$ . Suppose the method stops at iteration  $k$  with  $\nabla f(x^k) = 0$ . Then,  $x^k$  is a global minimizer of  $f$ .
- d) Consider an integer optimization problem and its LP relaxation. If the integer program is feasible, i.e., the feasible set is nonempty, then the LP relaxation also must be feasible.
- e) Consider the integer program

$$\min_x c^\top x \quad \text{s.t.} \quad Ax = b, \quad x \geq 0, \quad x \in \mathbb{Z}^n, \quad (1)$$

where  $A \in \mathbb{R}^{m \times n}$  is a totally unimodular matrix and  $b \in \mathbb{Z}^m$  is a given integer vector. We apply the interior-point method to solve the associated LP relaxation of problem (1). Then, the method is guaranteed to return an integer solution.

**Solution :**

- a) True.
- b) True.
- c) True.
- d) True.
- e) False.