

# Lecture 7

Plan:

- ✓ 0. Discuss pset
- ✓ 1. Finish LP duality (prev notes).
- 2. Faces of Polyhedra

# Faces of Polyhedra

Def:  $a^{(1)}, \dots, a^{(k)} \in \mathbb{R}^n$  are

affinely independent if

$$\sum_{i=1}^k \lambda_i a^{(i)} = 0$$

and  $\sum \lambda_i = 0$  imply  $\lambda_1 = \dots = \lambda_k = 0$ .

(w/out  $\sum \lambda_i = 0$ , is just linear indp.)

linear independent  $\Rightarrow$  affine independent.

Note:

$\{c(i)\}$  n. o. independent iff

$\left\{ \alpha \rightarrow \text{affinely independent} \right.$

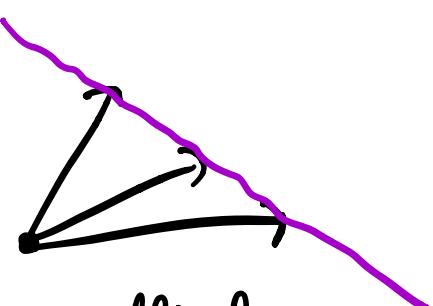
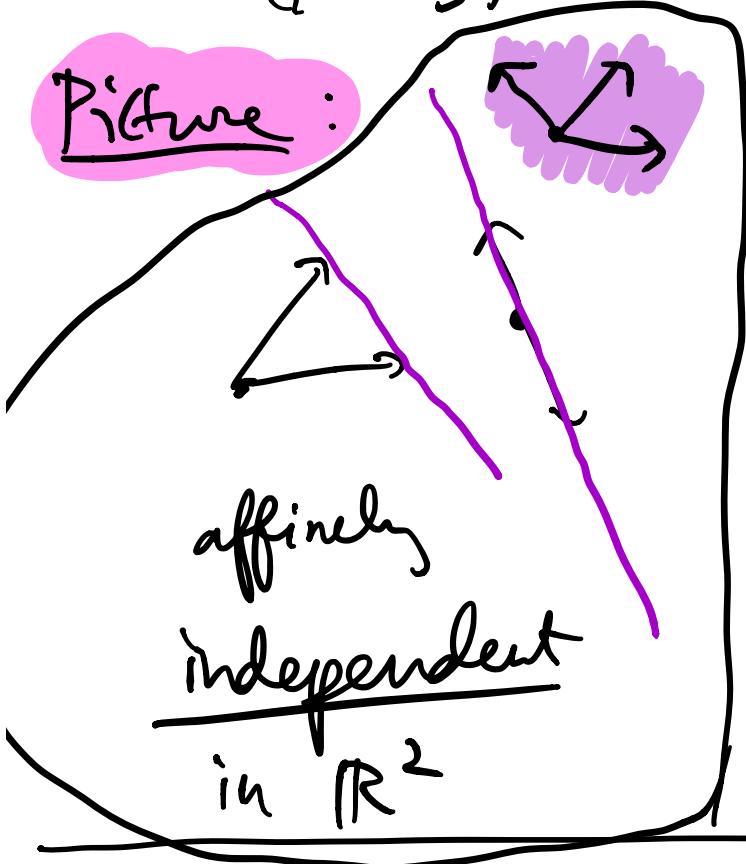
$$\left\{ \begin{bmatrix} \alpha^{(i)} \\ 1 \end{bmatrix} \right\}$$

linearly independent.

$\Leftrightarrow \text{aff}\left(\{\alpha^{(i)}\}\right)$  has dimension  $K-1$

# vectors.

Picture:



affinely  
dependent  
in  $\mathbb{R}^2$ .

Def Dimension  $\dim(P)$  of

polyhedron  $P$ :

$-1 + \max \# \text{affinely}$   
 $\text{independent points in } P.$

Equivalently, dimension of  
affine hull  $\text{aff}(P)$ .

Example:  $P = \emptyset, \dim(P) = -1$

$P = \text{singleton} \quad \cdot \quad \dim(P) = 0$

$P = \text{line segment} \quad \nearrow \quad \dim(P) = 1$

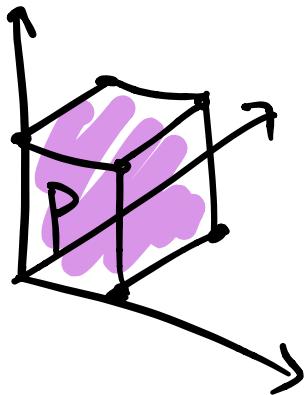
;

$$\text{aff}(P) = \mathbb{R}^n$$

$$\dim(P) = n;$$

P "full dimensional"

e.g. cube in  $\mathbb{R}^3$ :  $\{x : 0 \leq x_i \leq 1\}$



$$\dim P = 3$$

$$\dim \mathbb{R}^3 = 3$$

(as polyhedron).

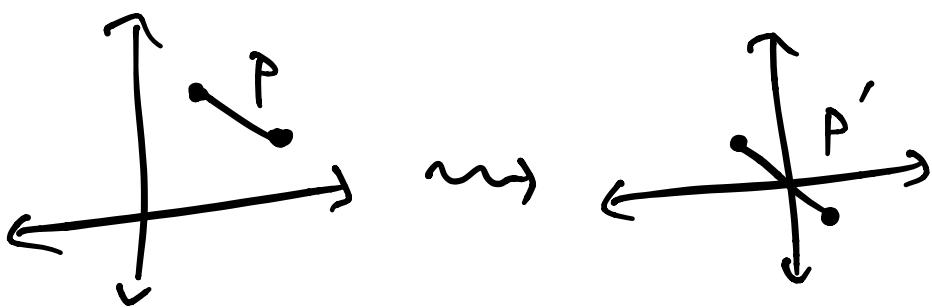
Why affine, not linear? affine

independence is translation  
invariant:

if I used max # lin indep points - 1

$$\dim(P) = 1$$

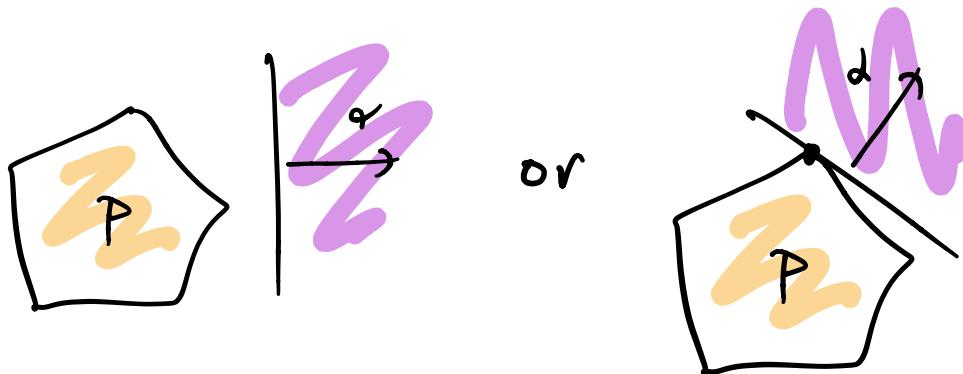
$$\dim(P') = 0$$



$$l = \dim(P) = \dim(P').$$

**Def:**  $\alpha^T x \leq \beta$  is a valid inequality

for  $P$  if  $\alpha^T x \leq \beta$  for all  $x \in P$ .

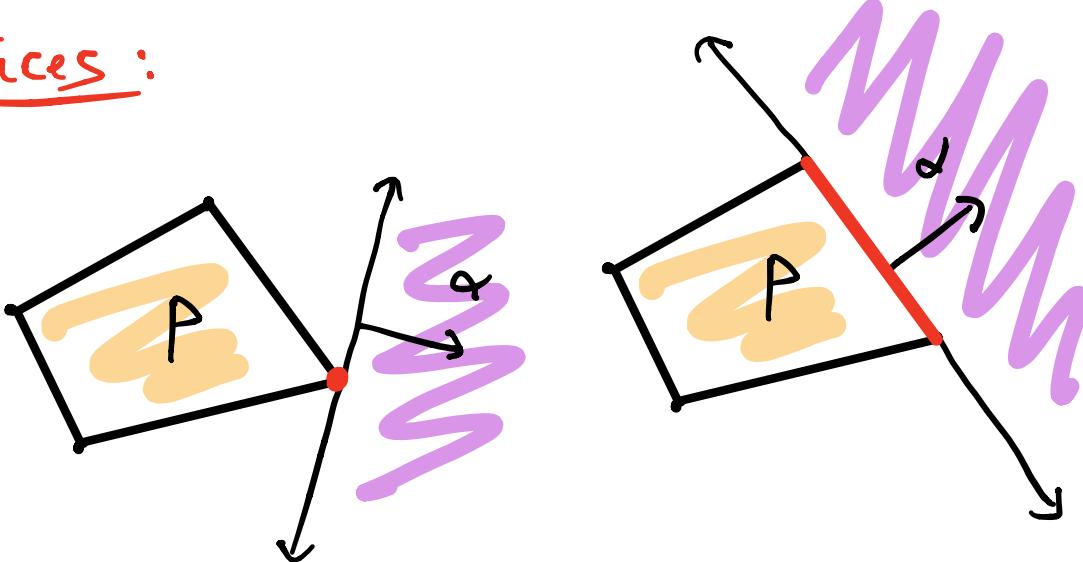


**Def** A face of a polyhedron

$P$  is  $\{x \in P : \}$  for

$Q^T x \leq \beta$  valid.

### Faces:



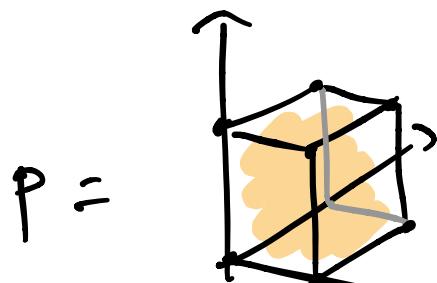
### Properties:

- Faces are polyhedra
- Empty face & entire P  
are called trivial faces)
- else F nontrivial  
 $\leq \dim(F) \leq$
- $F : \dim(F) = \dim(P) - 1$  called facets.  
in  $\mathbb{R}^n$  called vertices

~~• T - array for unitary vectors~~

Ex : list the 28 faces of the cube

$$P = \{x \in \mathbb{R}^3 : \quad \}$$



Fact :  $\infty$  many valid ineqs,  
but # faces finite!

Theorem : ("Faces" theorem)

Let  $A \in \mathbb{R}^{m \times n}$ ,

$$A = \left[ \begin{array}{c|c} \vdots & \\ -a_1^T & \hline \vdots & \end{array} \right]$$

Any nonempty face of  $P = \{x : Ax \leq b\}$   
is

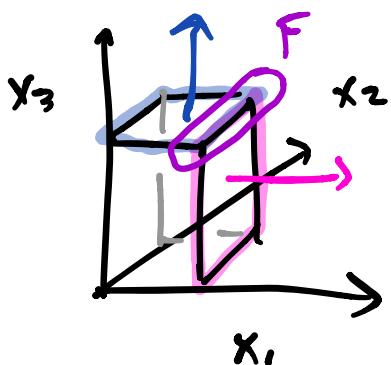
$$\left\{ x : \begin{array}{l} \\ \vdots \\ \end{array} \right.$$

for some set  $I \subseteq \{1, \dots, m\}$ .



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E.g. cube



$$F = \left\{ x : \begin{array}{l} \\ \vdots \\ \end{array} \right\}$$



Proof Consider valid inequality

$\alpha^T x \leq b$  giving nonempty face  $F$ .

- $F = \underline{\text{optimum}}$  solutions to bounded LP

$\max$   
 $(P)$  subject to

- Let  $y^*$  optimal solution to dual.

- ## • Complementary slackness:

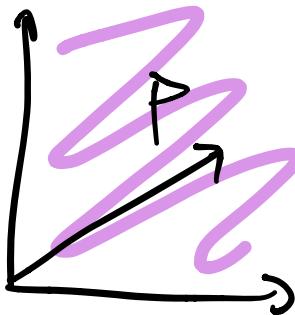
optimal solns F are

$\mathfrak{I}_V$  :

Thus we can take  $I = \{i : y_i^* > 0\}$ .  $\square$

Ex : positive orthant  $\{x \in \mathbb{R}^n : x_i \geq 0\} = P$   
has  $2^n + 1$  faces

- How many of  $\dim P$ ?



For polytopes can bound # faces in terms of # vertices.

"upper bound theorem!"

Dehn-Sommerville equation

For extreme points (dimension 0 faces)  
Can just use equalities.

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## Theorem ("Vertex" theorem)

Let  $x^*$  extreme point for

$$P = \{x : Ax \leq b\}.$$

$$A = \begin{bmatrix} & & \\ -a_1^T & - & \\ & & \end{bmatrix}$$

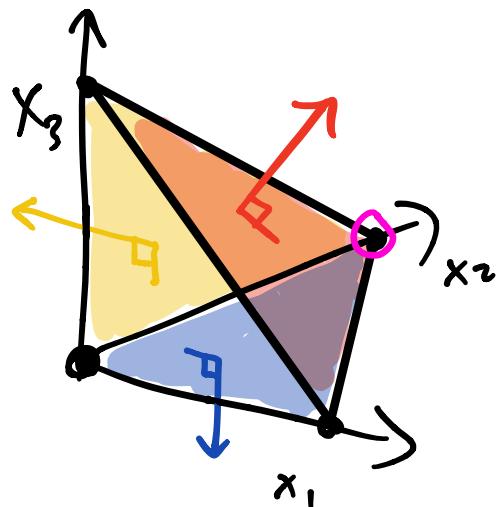
Then  $\exists I \subset \mathbb{N}$  s.t.  $x^*$  is the unique soln to



moreover, any such unique solution  $x^*$  is extreme.

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e.g. simplex  $(0,1,0)$  is intersection of  
3 constraints



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Proof: Given extreme point  $x^*$ ,

- define  $I = \{i : x_i^* > 0\}$ .
- Note for  $i \notin I$ ,  $x_i^* = 0$ .
- By "faces theorem",  $x^*$  uniquely defined by

(\*)

$i \in I$

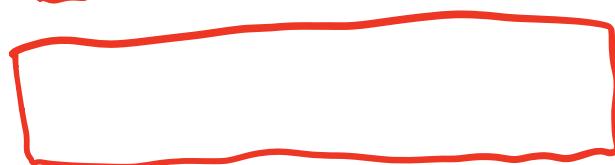
(\*\*)

$i \notin I$ .

- Suppose  $\exists$  other soln.  $\hat{x}$  to (\*).

- Because

for  $i \notin I$ ,



still satisfies (\*), (\*\*) for

- Contradicts  $F$  having only one point.  $\square$ .

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## Basic Feasible Solutions:

For  $P = \{ \quad , \quad \}$

can describe extreme points  
very explicitly.

(

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Corollary of Vertex Thm: Extreme pts. of

$P = \{ \quad \}$  come from setting

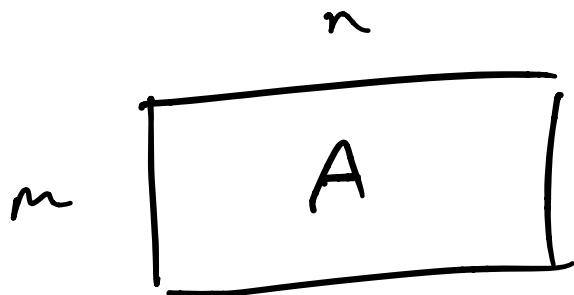


and finding unique solution to  
for remaining variables.

Can say more: Extreme points

of  $P = \{x : Ax = b, x \geq 0\}$  are  
the basic feasible solutions (BFS),  
feasible solns obtained as follows:

- Remove redundant rows  
from A ( )



- Choose  $m$  columns of  $A$ ,  $C$

$$\begin{matrix} & n \\ m & \boxed{\quad} \end{matrix} \quad \parallel \quad = \begin{bmatrix} b \end{bmatrix}$$

- Solve  $A_B X_B = 0$ ,  
set

$$x_i^* = \begin{cases} & i \in B \\ & \text{else} \end{cases}$$

$$\{ \text{bfcs} \} = \{ \text{extreme pts} \}.$$

## Corollary of Faces Theorem

facets are the maximal nontrivial faces of a nonempty polyhedron  $P$ .

Pf : Exercise.

## Corollary of vertex theorem

vertices are the minimal nontrivial

Facets .

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Pf: