Lecture 2

Plan: 1) finish arborescence 2) matroid union

Spanning tree zource:

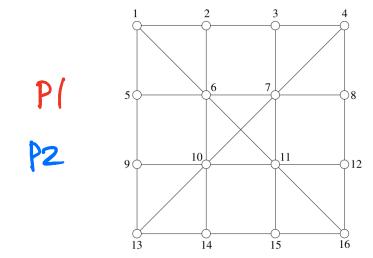
Given graph G, players alternate:

1)

2)

Pl wins if

eg. Pl win:



Recall: Pawinsit A)

Plwinsif B)

Today: with matroid union, show

Matroid Union:

Let M= (E, I) matroid.

Recall durl matroid

E.g. If M = MG for G = A, i.e.

Theorem The dual moutroid Mx

4

Proof Use

Fact: Com define a motrotel using

KI)

R2)

Then M: (C12)

工二

à

Thus

A)

8)

eg. asjoint spanning trees:

Ghas 2 disjoint spanning trees €>

and

moreover,

Theorem: G has two disjoint spanning trees =>

Proof Assume

· We only show (2);

Plan: use Minimax theorem for

· Let

· Ghas 2 edge disjt.

spaning trees (2)

· rn(F)=

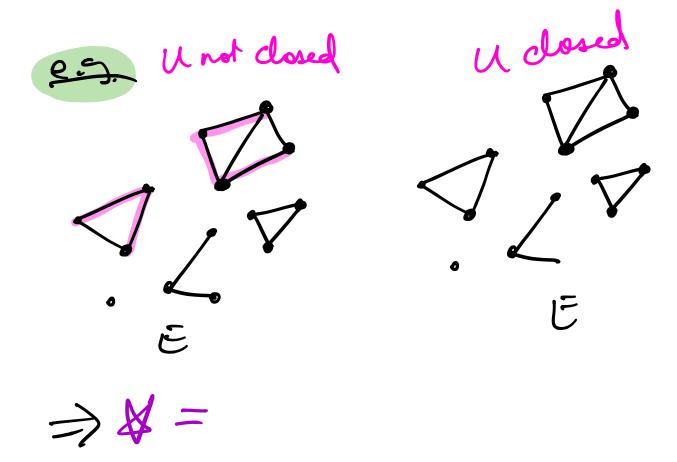
min-max:

· Matroid Intersection theorem:

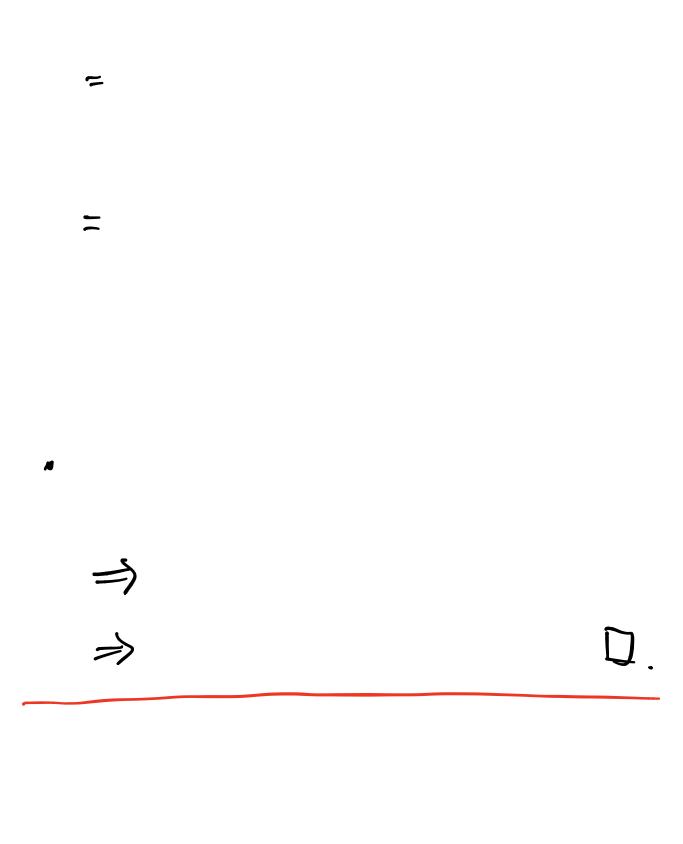
X

· Recall:

i.e.



=



(General) matroid union

· let M,=(E,I,), M,=(E,I), matroids.

Def. The matroid union

Careful:

Theorem: M,UMz is

Proof

Part 1: elté a matroid.

· Let

. Need to show

· Assume

· Since

1

=>

).

End part I \

Part 2: Rank function.

rm, um2(s)=

- · For >, use
- First

- · fet
- · May assure

· Then

(i.e.

LORMAR

- . I.e.
- · moutrood intersection theorem

rminns (E) =

7

=

=

endpart2 D.