

In HW 3, successive characters in an input string are grouped into coherent units, called tokens. In the resulting tabToken, a series of tokens and their types are stored from the input string. HW4 processes tokens in the token table and produces a Symbol Table (symTab).

In a Symbol Table, variable names used in the program and their values are saved. Each entry in the symbol table consists of (a) the TOKEN (maximum 8 bytes or two words), (b) its VALUE (one word), and (c) the status (one word).

The value of a variable is **the address** of the input string where the variable is used. The LOC (location counter) keeps track of the addresses. LOC is initialized to 0x0400 and is incremented by 4 each time a new input string is processed.

The status of a symbol has three flags: F – the first occurrence of the symbol; A – already defined; D – a new definition. A symbol is defined from the current input string, if the symbol is found in the label field of the input string. The flag A keeps track if the symbol is already defined, and therefore its value in the symbol table has a valid address.

The status field uses only the last three bits: (First occurrence, Already defined, DEFN).

The basic structure of the assignment is as follows.

```

        LOC = 0x0400
nextLine:
        Read an input string and save token and type in tabToken [HW 3]
        If (tabToken[0][0] == '#') goto exit

        i=0;                                // index to TOKENS[][]
nextTok:  if (tabToken[i+1][0] != ':') goto operator
        TOKEN = tabToken[i][0]
        valVar = VAR (TOKEN, 1) // store label in symTab
        i +=2                        // skip ':'

operator: i++                                // do nothing with operator

        isComma = true
chkVar:  if (tabToken[i][0] == '#') goto dump // '#'
        if (!isComma || tabToken[i][1]!=2) goto nextVar
        TOKEN = tabToken[i][0]
        valVar = VAR (TOKEN, 0)

nextVar: isComma = (tabToken[i][0] == ',')
        i++
        goto chkVar

dump:   clear inBuf
        clear tabToken
        print symTab
        LOC +=4
        goto nextLine

exit:   stop the program

```

VAR () is passed the current TOKEN (or a pointer (index) to the current TOKEN), and a flag DEFN. VAR () returns a value as determined by the outcome of symbol table action routines, symACT0(), symACT1(), ..., symACT5().

VAR () first calls LookUp() to search for the TOKEN in **the symTab**. LookUp() returns -1 if the token is not found, the index of the token in the symbol table, otherwise.

If the token is not found, the newStatus shall be (F,A,D)=(1, 0, DEFN), and you save the token in the symbol table. In either case, the index to the symbol table entry of the token is saved in symIndex. If the token is found, however, the newStatus is dependent on the previous values of FAD in the symbol table.

```
VAR (TOKEN, DEFN){
    symIndex = srchSymTab(TOKEN)           // -1 if not found; index if found
    if (symIndex < 0) {                     // First occurrence
        newStatus = 0x4 | DEFN;
        symIndex = saveSymTab(TOKEN, newStatus);
    }
    else {
        oldStatus = symTab[symIndex][2];
        // determine the new status (F flag = 0)
        newStatus = oldStatus & 0x2 | ((oldStatus & 0x1) << 1); // set A flag
        newStatus = newStatus | DEFN;           // set D flag
        symTab[symIndex][2]=newStatus;
    }
    retVal = symACTS(newStatus, symIndex);    // from the jump table
    retVar: return retVal
}
```

Based on the newStatus, the assembler updates the symbol table and return a value according to the following.

4's (F)	2's (A)	1's (D)	action
0	0	0	symACT0: This is a forward reference. Store LOC in VALUE field of the symbol table.
0	0	1	symACT1: A previously used TOKEN is defined for the first time. Store LOC in VALUE field of the symbol table.
0	1	0	symACT2: A previously defined TOKEN is used.
0	1	1	symACT3: A previously defined TOKEN is defined again. Print "Double Definition Error."
1	0	0	symACT4: TOKEN seen for the first time as a forward reference. Store LOC in VALUE field.
1	0	1	symACT5: TOKEN seen for the first time as a definition. Store LOC in VALUE field of the symbol table.

1	1	0	N/A
1	1	1	N/A

Namely, the newStatus value triggers a function call indexed by the value. Since there are six statuses to consider, a jump table is often used to process actions corresponding to the new status value.

```

symACTS:    b    symACT0
            b    symACT1
            b    symACT2
            b    symACT3
            b    symACT4
            b    symACT5

```

The last instruction in each symACTs should be a branch instruction with **the target address of ret_VAR** to return from VAR ().

Each action routine will carry out what is described in the ‘action’ column of the table below. In order to call the right routine, the following MIPS instructions can be used, assuming that \$s1 has the new status value (you are free to use another register other than \$s1).

```

la    $s0, symACTS
sll   $s1, $s1, 2
add   $s0, $s0, $s1
jr    $s0

```

retSymAct:

Note

MIPS has a single level of function calls, as the return address is saved in a single register of \$ra. Thus, when VAR () calls LookUp(), for example, the return address from VAR () has to be saved prior to the call to LookUp().

Note

Your program will be tested according to the tokens generated by the following set of MIP instructions:

```

hex2char:    sw    $t0, saveReg($0)
            li    $t9, 3                # $t9: counter limit
            jal   hex2char
saveReg:     or    $t0, $t1, $0

```

the symbol table will be printed as follows.

```

hex2char:    sw    $t0, saveReg($0)

symTab:      hex2char    0x0400    0x5
            saveReg     0x0400    0x4
            li    $t9, 3                # $t9: counter limit

symTab:      hex2char    0x0400    0x5

```

```

                saveReg    0x0400    0x4
                jal        hex2char

symTab:         hex2char    0x0400    0x2
                saveReg    0x0400    0x4

                saveReg:    or      $t0, $t1, $0

symTab:         hex2char    0x0400    0x2
                saveReg    0x040C0x1

```

hex2char function

```

#
#   hex2char:
#       Function used to print a hex value into ASCII string.
#       Convert a hex in $a0 to char hex in $v0 (0x6b6a in $a0, $v0 should have 'a"6"b"6')
#
#       4-bit mask slides from right to left in $a0.
#       As corresponding char is collected into $v0,
#       $a0 is shifted right by four bits for the next hex digit in the last four bits
#
#       Make it sure that you are handling nested function calls in return addresses
#
                .data
saveReg:       .word    0:3

                .text
hex2char:
    # save registers
    sw        $t0, saveReg($0)    # hex digit to process
    sw        $t1, saveReg+4($0)  # 4-bit mask
    sw        $t9, saveReg+8($0)

    # initialize registers
    li        $t1, 0x0000000f     # $t1: mask of 4 bits
    li        $t9, 3              # $t9: counter limit

nibble2char:
    and       $t0, $a0, $t1       # $t0 = least significant 4 bits of $a0

    # convert 4-bit number to hex char
    bgt       $t0, 9, hex_alpha   # if ($t0 > 9) goto alpha
    # hex char '0' to '9'
    addi      $t0, $t0, 0x30       # convert to hex digit

```

```

b      collect

hex_alpha:
    addi    $t0, $t0, -10      # subtract hex # "A"
    addi    $t0, $t0, 0x61     # convert to hex char, a..f

    # save converted hex char to $v0
collect:
    sll     $v0, $v0, 8        # make a room for a new hex char
    or      $v0, $v0, $t0      # collect the new hex char

    # loop counter bookkeeping
    srl     $a0, $a0, 4        # right shift $a0 for the next digit
    addi    $t9, $t9, -1       # $t9--
    bgez    $t9, nibble2char

    # restore registers
    lw      $t0, saveReg($0)
    lw      $t1, saveReg+4($0)
    lw      $t9, saveReg+8($0)
    jr      $ra

```