

Programming Language Concepts

CSCI-344
Term 20175

Programming 5
March 5, 2018

Type Systems Due: March 26, 2018

1 Introduction

In this programming assignment, you will extend the type checker for Typed Impcore (relatively easy) and build the type checker for Typed μ Scheme (relatively hard).

This programming assignment must be completed as a pair programming assignment; see [WK00] for useful guidelines on pair programming.

2 Description

Complete the following exercises. See Requirements and Submissions for important restrictions.

- A. (20pts) Complete Exercise 2 of Chapter 6 from *Programming Languages: Build, Prove, and Compare* (p. 459). The exercise asks you to complete the type checker for Typed Impcore by implementing the rules for array operations. You will modify the provided `timpcore.sml` interpreter.

You need only complete the four cases of the `ty` function (within the `typeof` function) that currently raise `LeftAsExercise` exceptions (lines 1292–1295). You should not need to modify any other part of the interpreter (but you will need to read other parts of the interpreter, particularly the `datatype ty` definition (line 1034), which gives the SML representation of Typed Impcore types). A reasonable solution is 20 to 50 lines of SML, depending on the quality of the error messages (but high quality error messages are not required for this assignment).

Tips, Advice, and Hints:

- Although most of the existing cases of the `ty` function use `eqType` to compare the type of a subexpression with the required type of the subexpression, this approach will not work if the required type of the subexpression is an array type and you do not know the type of the elements of the array. Instead, you will need to use `case` expressions to match the type of a subexpression against `ARRAYTY t` in order both check that the type of the subexpression is an array type and to extract the type of the elements.
- B. (145pts) Complete Exercises 22 and 23 of Chapter 6 from *Programming Languages: Build, Prove, and Compare* (pp. 465–468). The exercises ask you to implement renaming to avoid variable capture and to write the type checker for Typed μ Scheme. You will modify the provided `tuscheme.sml` interpreter.

You need to complete the `renameForallAvoiding` function that currently raises a `LeftAsExercise` exception, all cases of the `ty` function (within the `typeof` function) that currently raise `LeftAsExercise` exceptions and all cases of the `elabdef` function that currently raise `LeftAsExercise` exceptions. You should not need to modify any other part of the interpreter (but you will need to read other parts of the interpreter, particularly the `datatype tyex` definition (lines 1022–1026), which gives the SML representation of Typed μ Scheme type-level expressions, and the `datatype exp`, `datatype value`, and `datatype def` definitions (lines 1168–1178, lines 1181–1187, and lines 1192–1195), which give the SML representation of Typed μ Scheme expressions, values, and definitions). A reasonable solution is 120 to 150 lines of SML, depending on the quality of the error messages (but high quality error messages are not required for this assignment).

Tips, Advice, and Hints:

- Don't copy and paste the Typed Impcore `typeof` and `elabdef` functions into Typed μ Scheme; there are too many details to change to make it worthwhile. Use the structure as a guide, but start from scratch.
- Compile early and compile often!

- Test early and test often! Be sure to test with both examples that should type check and examples that should *not* type check.
- Build the type checker one piece at a time:
 - First, write code to type check LITERAL/NUM, LITERAL/BOOLV, and LITERAL/SYM.
 - Next, type check VAL. (Because EXP is just syntactic sugar for VAL, this will also type check EXP).
 - * At this point, “programs” like (val ans 1) and (val ans #t) should type check.
 - Next, type check IFX.
 - * At this point, “programs” like (val ans (if #t -1 1)) and (val ans (if #f -1 1)) should type check.
 - Next, type check VAR.
 - * At this point, “programs” like (val x #t) (val y -1) (val z 1) (val ans (if x y z)) should type check.
 - Next, type check SET, BEGIN, and WHILE.
 - Next, type check LETX/LET. Be careful with the environments; be sure to type check sub-expressions with the right environment.
 - Next, type check LETX/LETSTAR (by treating it as syntactic sugar for nested LETs).
 - Next, type check LAMDBA (which is quite similar to LETX/LET). To create a function type, use the FUNTY constructor of the tyex datatype.
 - Next, type check APPLY. Pattern match against the FUNTY constructor of the tyex datatype to simultaneously determine if a type is a function type and to extract the argument types and the result type of the function type.
 - Next, type check VALREC and DEFINE (remembering that DEFINE is syntactic sugar for VALREC). Don’t forget to check the form of the expression in VALREC.
 - Next, type check LETRECX. Don’t forget to check the form of the expressions in LETRECX.
 - Next, type check LITERAL/NIL and LITERAL/PAIR. Remember that the empty list literal is polymorphic, but non-empty list literals are monomorphic.
 - Next, type check TYLAMBDA. Don’t forget to check the restriction that a type-lambda may not abstract over a type variable that’s free in the type environment.
 - Finally, implement renameForallAvoiding and type check TYAPPLY.
- Reread Section 6.6 of *Programming Languages: Build, Prove, and Compare*, paying special attention to the useful functions and representations that are already implemented in the interpreter.
- Do not use SML’s polymorphic equality (the = operator) to compare types. SML’s polymorphic equality compares values for *structural equality*. Use the provided eqType function to compare types. More details about the eqType function may be found in Section 6.6.6 of *Programming Languages: Build, Prove, and Compare*.

C. (bonus 20pts) Complete Exercises 25 and 26 of Chapter 6 from *Programming Languages: Build, Prove, and Compare* (p. 469). The exercises asks you to implement the pair and sum type constructors and the polymorphic functions pair, fst, snd, left, right, and either *without* changing the abstract syntax, values, type checker, or evaluator of Typed μ Scheme. You will modify the provided tuscheme.sml interpreter.

You need only add to the kinds (Δ), types (Λ), and values (ρ) components of the val initialBasis definition (lines 2220–2464).

2.1 Interpreter Source Code

Source code for the interpreters is available on the CS Department file system at:

```
/usr/local/pub/mtf/plc/programming/prog05-typesys
```

and packaged as an archive at:

```
/usr/local/pub/mtf/plc/programming/prog05-typesys.tar
```

Note that this source code contains just the SML code from the textbook, with simple comments identifying page numbers. There is a `Makefile` for building the interpreters and running tests.

Copy the interpreter source code to a local directory and make modifications to your local copy; for example, executing

```
$ tar xvf /usr/local/pub/mtf/plc/programming/prog05-typesys.tar
```

will copy the interpreter source code to a new local directory named `prog05-typesys`.

2.2 Tests

The `prog05-typesys/tests` directory contains a number of tests for the Typed Impcore and Typed μ Scheme type checkers, as well as scripts for running the tests.

For example, executing

```
$ make timpcore-tests.out
```

from your `prog05-typesys` directory will build your Typed Impcore interpreter and run the Typed Impcore tests; the test results will be echoed to the terminal and also saved in the `timpcore-tests.out` file.

For each `test.imp` file, if the test has no type errors, then there is a `test.soln.tychk` file containing the output of the reference interpreter (the name and/or value defined by each declaration and the type of the defined name and/or value; see *Programming Languages: Build, Prove, and Compare* (pp. 406–408)). If the test has type errors, then there is a `test.soln.tyerr` file containing the output of the reference interpreter, concluding with the type error message reported by the reference solution type checker.

Similarly, executing

```
$ make tuscheme-tests.out
```

from your `prog05-typesys` directory will build your Typed μ Scheme interpreter and run the Typed μ Scheme tests; the test results will be echoed to the terminal and also saved in the `tuscheme-tests.out` file.

For each `test.scm` file, if the test has no type errors, then there is a `test.soln.tychk` file containing the output of the reference interpreter (the name and/or value defined by each declaration and the type of the defined name and/or value; see *Programming Languages: Build, Prove, and Compare* (p. 436) and *Programming Languages: Build, Prove, and Compare* (p. 453)). If the test has type errors, then there is a `test.soln.tyerr` file containing the output of the reference interpreter, concluding with the type error message reported by the reference solution type checker.

2.3 Reference Interpreters

Reference Typed Impcore and Typed μ Scheme interpreters are available on the CS Department Linux systems (e.g., `glados.cs.rit.edu` and `queeg.cs.rit.edu` and ICLs 1 and 2) at:

```
/usr/local/pub/mtf/plc/bin/timpcore
```

and

```
/usr/local/pub/mtf/plc/bin/tuscheme
```

Use the reference interpreters to develop tests and to check your interpreters.

3 Requirements and Submission

Always use pattern matching to inspect and deconstruct values. The use of `null`, `hd`, `tl`, `#1`, `#2`, etc., will result in zero credit for the assignment.

Your modified interpreters must be a valid Standard ML program. In particular, they must compile with Moscow ML or MLton without any error messages. If your submission produces error messages (e.g., syntax errors or type errors), then your submission will not be tested and will result in zero credit for the assignment.

Write a `README.txt` file. Your `README.txt` file should be formatted as follows:

<code>Names:</code> <code>Time spent on assignment:</code> <code>Additional Collaborators:</code>

Submit `README.txt`, `timpcore.sml`, and `tuscheme.sml` to the Programming 05 Dropbox on MyCourses by the due date. Only one submission is required per group.

References

[WK00] Laurie A. Williams and Robert R. Kessler. All I Really Need to Know About Pair Programming I Learned in Kindergarten. *Communications of the ACM*, 43(5):108–114, May 2000. <http://doi.acm.org/10.1145/332833.332848>.