

1. Average seek time = 10ms
 Average rotational latency = 5ms
 Transfer time = 1ms for 4KB
 MM = 320 blocks
 # of read/write blocks = 10,000,000
 $320 \times 4 \text{ KB} = 1.28$

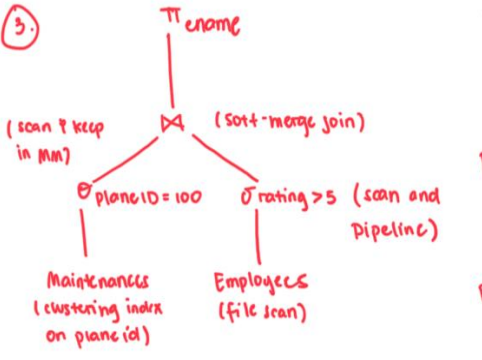
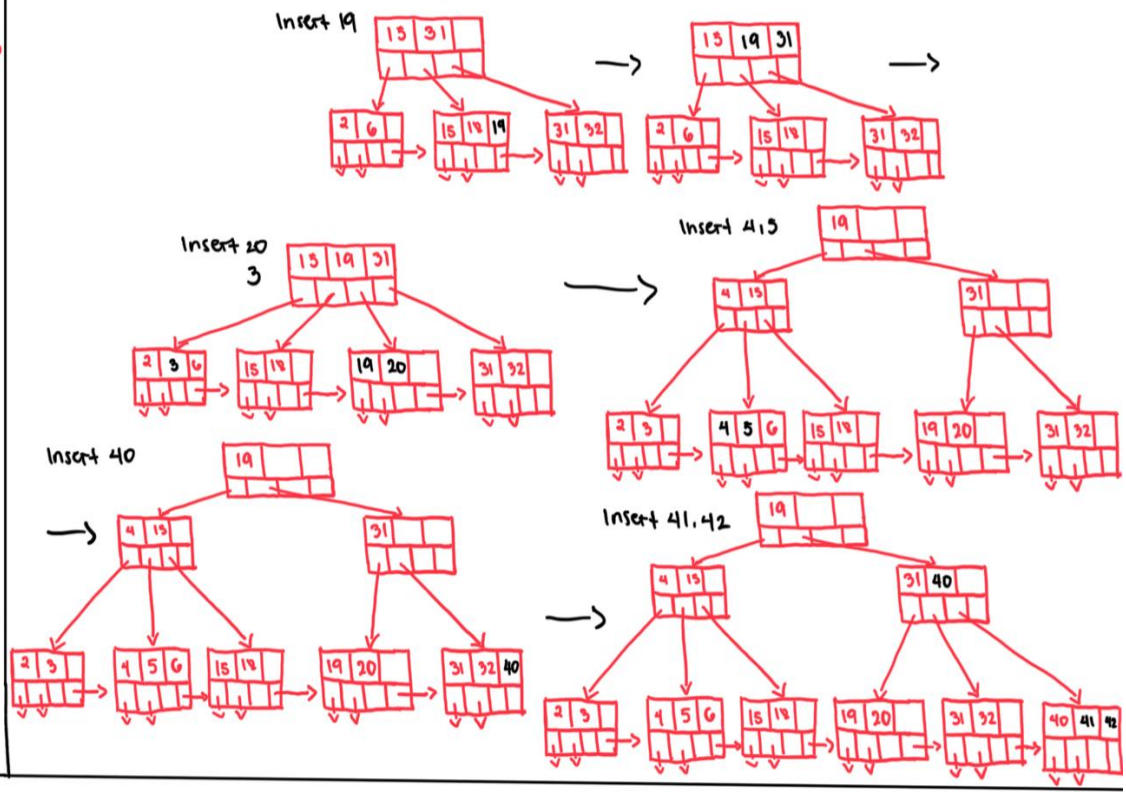
Phase 1: $\frac{10,000,000}{320} = 31,250$ sorted runs

Phase 2: Merge: 319-way merges, so
 $\frac{31,250}{319} = 98$ sorted sublists

Phase 3: 2PMMS for the total # of I/O's
 $3 \times 2 \times 10^7 = 6 \times 10^7$
 too big

Total: $6 \times 10^7 \times 16 \text{ms} = 960,000,000 \text{ms}$
 or 11 days

2. $n = 3$ keys
~~2, 15, 31, 32, 4, 19, 19, 20, 3, 4, 5, 40, 41, 42~~



$12 + 75 + 40 + 35 = 162 \text{ I/Os}$

Pushing Selection II: cost of applying $\sigma_{\text{planeid}=100}$
 Maintenance = 1000 blocks / 100 = 10
 Planeid clustering = 2 I/Os + 10 I/Os = 12

Pushing Selection III: cost of applying rating > 5
 $50 + 25 = 75 \text{ I/Os}$
 for scan and pipeline & MM

Pushing Selection IV: sort-merge join
 2PMMS cost of sorting
 $2 \times 2 \times 10 = 40$
 $10 + 25 = 35$
 maintenance employees



4. T_1, T_2, T_3

1. $r_1(A) \ r_2(B) \ r_3(C) \ r_1(B) \ r_2(C) \ r_3(D) \ w_1(A) \ w_2(B) \ w_3(C)$

i. $T_1 \quad T_2 \quad T_3$
 $x_{l_1}(A); r_1(A)$
 $x_{l_2}(B); r_2(B)$
 $x_{l_3}(C); r_3(C)$
 $s_{l_1}(B); \text{denied}$
 $s_{l_2}(C); \text{denied}$
 $s_{l_3}(D); r_3(D)$
 $w_3(C); u_3(C); u_3(D);$
 $s_{l_2}(C); r_2(C)$
 $w_2(B); u_2(B); u_2(C);$
 $r_{l_1}(B); r_1(B)$
 $w_1(A); u_1(A); u_1(B)$

The first three reads were granted, and the next two requests were denied because both T_1 and T_2 aren't granted the shared lock due to the exclusive locks on B & C respectively. T_3 , however, is granted a shared lock, so it releases the locks on C & D. From there, T_2 & T_1 gets a shared lock for some X and releases the lock.

ii. $T_1 \quad T_2 \quad T_3$
 $s_{l_1}(A); r_1(A)$
 $s_{l_2}(B); r_2(B)$
 $s_{l_3}(C); r_3(C)$
 $s_{l_1}(B); r_1(B)$
 $s_{l_2}(C); r_2(C)$
 $s_{l_3}(D); r_3(D)$
 $x_{l_1}(A); w_1(A); u_1(A); u_1(B)$
 $x_{l_2}(B); w_2(B); u_2(B); u_2(C);$
 $x_{l_3}(C); w_3(C); u_3(C); u_3(D);$

The first 6 read actions and shared locks were granted. No requests are denied by the scheduler and update locks were upgraded to exclusive locks.

iii. $T_1 \quad T_2 \quad T_3$
 $u_{l_1}(A); r_1(A);$
 $u_{l_2}(B); r_2(B);$
 $u_{l_3}(C); r_3(C);$
 $s_{l_1}(B); \text{denied}$
 $s_{l_2}(C); \text{denied}$
 $s_{l_3}(D); r_3(D);$
 $w_3(C); u_3(C); u_3(D)$
 $s_{l_2}(C); r_2(C);$
 $w_2(B); u_2(B); u_2(C);$
 $s_{l_1}(B); r_1(B);$
 $w_1(A); u_1(A); u_1(B)$

The first three reads were granted. The request for a shared lock on T_1 & T_2 were denied because T_2 & T_3 already has an update lock. A scheduler will not grant a shared lock when there is an update lock.

2. $r_1(A); r_2(B); r_3(C); r_1(B); r_2(C); r_3(A); w_1(A); w_2(B); w_3(C);$

i. $T_1 \quad T_2 \quad T_3$
 $x_{l_1}(A); r_1(A)$
 $x_{l_2}(B); r_2(B)$
 $x_{l_3}(C); r_3(C)$
 $s_{l_1}(B); \text{denied}$
 $s_{l_2}(C); \text{denied}$
 $s_{l_3}(D); \text{denied}$

The first three read actions and exclusive locks were granted. But, requests for a shared lock results to a deadlock & no transaction proceeds. This is because the requests for some $s_{l_i}(X)$ already has an exclusive lock for some transaction T_i

ii. $T_1 \quad T_2 \quad T_3$
 $s_{l_1}(A); r_1(A)$
 $s_{l_2}(B); r_2(B)$
 $s_{l_3}(C); r_3(C)$
 $s_{l_1}(B); r_1(B)$
 $s_{l_2}(C); r_2(C)$
 $s_{l_3}(D); r_3(D)$
 $x_{l_1}(A); \text{denied}$
 $x_{l_2}(B); \text{denied}$
 $x_{l_3}(C); \text{denied}$

The first 6 shared locks and read actions were granted. The next requests were denied and resulted to a deadlock. The request for an upgrade to exclusive is denied by the scheduler because the request for some lock X already has a shared lock for some transaction T_i

iii. $T_1 \quad T_2 \quad T_3$
 $u_{l_1}(A); r_1(A);$
 $u_{l_2}(B); r_2(B);$
 $u_{l_3}(C); r_3(C);$
 $s_{l_1}(B); \text{denied}$
 $s_{l_2}(C); \text{denied}$
 $s_{l_3}(A); \text{denied}$

The first three read actions were granted, but the next requests resulted to a deadlock. This is because request for some $s_{l_i}(X)$ gets denied due to an update lock on the same X for some transaction T_i

e.g., T_2 is granted a shared lock on C that then unlocks B & C. T_1 is then granted with a lock on B that then unlocks A & B.

5.

a.) $R(A, B, C, D)$

$AB \rightarrow C$

$B \rightarrow D$ violation!

$CD \rightarrow A$

$AD \rightarrow B$

$R_1(B, D), R_2(B, A, C)$

$A^+ = A$

$B^+ = D, B$

$C^+ = C$

$D^+ = D$

$AB^+ = A, B, C, D$

$BC^+ = B, C, D, A$

$CD^+ = C, D, A, B$

$AD^+ = A, D, B, C$

$A \ B \ C \ D$

$AD = BC$

$AB = CD$

$CD = AB$

b.) No BCNF violations

$A \rightarrow B$

$B \rightarrow C$

$C \rightarrow D$

$D \rightarrow A$

$R(A, B, C, D)$