Circus-LATEX style explained Community Z Tools (CZT)

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* PARSER DOES NOT YET SUPPORT LOGICAL CONSTANTS

A Environments

A.1 Informal argument

To typeset an informal argument, you write in LATEX

```
\lambda \text{legin{argue}} \ S \dres (T \dres R) \\ \t1 = \id S \comp \id T \comp R \\ \t1 = \id (S \cap T) \comp R & law about $\id$ \\ \t1 = (S \cap T) \dres R. \end{argue} \text{which corresponds to} \  S \lhd (T \lhd R) \  = \id S \\( \circ \) id T \\( \circ \) R \  = \id (S \cap T) \dres R. \end{argue} \  \text{legin about id} \  = (S \cap T) \dres R. \end{argue} \  \text{legin about id} \  = (S \cap T) \dres R. \end{argue} \  \text{legin about id} \ \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \  \text{legin about id} \} \}
```

B Reference card

B.1 Special Circus symbols

Refinement

```
\begin{array}{lll} n \Vdash S = I & \text{n \circassertref S = I} \\ n \Vdash S \preccurlyeq I & \text{n \circassertref S \circsimulates I} \\ n \Vdash S \sqsubseteq I & \text{n \circassertref S \circrefines I} \\ n \Vdash S \sqsubseteq Tr\,I & \text{n \circassertref S \circrefines Tr}^{\sim} I \\ n \Vdash S \sqsubseteq SFl\,I & \text{n \circassertref S \circrefines SFl}^{\sim} I \\ n \Vdash S \sqsubseteq FlDv\,I & \text{n \circassertref S \circrefines FlDv}^{\sim} I \end{array}
```

B.2 *Circus* channels and name sets

```
channel e
                               \circchannel e
channel c:T
                               \circchannel c : T
\mathbf{channel}[X]c:X
                               \circchannel [X] c : X
{\bf channel from}\, S
                               \circchannelfrom S
channelfrom[X]S[X]
                               \circchannelfrom [X] S[X]
channelset n == \{ |c| \}
                               \circchannelset n == \lchanset c \rchanset
channelset n == CSRef
                               \circchannelset n == CS
\mathbf{channelset}[X]n == \mathit{CSRef}
                               \circchannelset [X] n == CS
nameset n == \{x\}
                               \circnameset n == \{ x ^{x} \}
\mathbf{nameset} \; n == \mathit{NSRef}
                               \circnameset n == NS
```

B.3 Circus actions

Action definition

```
n \mathrel{\widehat{=}} A n \circdef A n \mathrel{\widehat{=}} x : T \bullet A n \circdef x: T \circspot A
```

Basic actions Prefixing action

```
e \longrightarrow A
                                                                         e \then A
                                                   c.0 \longrightarrow A
                                                                         c.0 \then A
Skip
                \Skip
                                                   c!v \longrightarrow A
                                                                         c!v \then A
Stop
                \Stop
                                                   c?x \longrightarrow A
                                                                         c?x \then A
Chaos
                \Chaos
                                                   c?x!y?z \longrightarrow A
                                                                        c?x!y?z \then A
```

Prefixing action (extra)

```
\begin{array}{l} c?x:(P)!(f\,x)\longrightarrow A\\ \\ c?x\prefixcolon\prefix(P)!(f~x)\prefix(P)!(f~x)\prefix(P)\to A\\ \\ c?x\prefixcolon\prefix(P)\to A\\ \\ c[\mathbb{N}\times\mathbb{P}\prefix(P)]?x!(dom\prefix(P)\to A\\ \\ c[\text{$\mathbb{N}\times\mathbb{P}\prefix(P)]$}?x!(dom\prefix(P)\to A\\ \\ c[\text{$\mathbb{N}\times\mathbb{P}\prefix(P)$}]?x!(dom\prefix(P)\to A\\ \\ c[\text{$\mathbb{N}\times\mathbb{P}\prefix(P)$}]
```

Unary actions

 $(S) \qquad \qquad \text{\lschexpract S \lschexpract} \\ \mu X \bullet A \qquad \qquad \text{\lcircmu X \lcircspot A} \\ A \backslash CS \qquad \qquad \text{\lcirchide CS} \\ (P) \otimes A \qquad \qquad \text{\lcircguard P \lcircguard \lcircguard A} \\ \label{eq:lcircguard}$

Binary actions

 $A \parallel NSa \mid NSb \parallel B$ A \linter NSa | NSb \rinter B $A \parallel \mid B$ A \interleave B $A \parallel NSa \mid CS \mid NSb \parallel B$ A \lpar NSa | CS | NSb \rpar B $A \parallel CS \parallel B$ A \lpar CS \rpar B $A[NSb \mid CSa \mid CSb \mid NSb]B$ A [NSb | CSa | CSb | NSb] B $A[CSa \mid CSb]B$ A [CSa | CSb] B $A \sqcap B$ A \intchoice B $A \square B$ A \extchoice B A; BA \circseq B ANameAName AName(x, y)AName(x, y)AName[new/old, x/y]AName[new/old, x/y]

Replicated actions

Parenthesised actions

```
(A)
                              (A)
(x:T\bullet A)
                              (x : T \circspot A)
(\mathbf{val}\,x:T\bullet A)
                              (\circval x : T \circspot A)
(\mathbf{res}\,x:T\bullet A)
                              (\circres x : T \circspot A)
(\mathbf{vres}\,x:T\bullet A)
                              (\circvres x : T \circspot A)
(x:T\bullet A)(v)
                              (x : T \circspot A)(v)
(\mathbf{val}\,x:\,T\bullet A)(v)
                              (\circval x : T \circspot A)(v)
(\mathbf{res} \ x : T \bullet A)(v)
                              (\circres x : T \circspot A)(v)
(\mathbf{vres}\,x:\,T\bullet A)(v)
                              (\circvres x : T \circspot A)(v)
(\mu X \bullet x : T \bullet A)(v)
                              (\circmu X \circspot x : T \circspot A)(v)
(\mu X \bullet (x : T \bullet A))(v)
                              (\circmu X \circspot (x : T \circspot A))(v)
(\mu X \bullet \mathbf{val} x : T \bullet A)(v) (\circmu X \circspot \circval x : T \circspot A)(v)
(\mu X \bullet \mathbf{res} \, x : T \bullet A)(v) (\circmu X \circspot \circres x : T \circspot A)(v)
(\mu X \bullet \mathbf{vres} \, x : T \bullet A)(v) (\circmu X \circspot \circvres x : T \circspot A)(v)
```

B.4 Circus command definitions

Guarded commands

```
x, y := v1, v2
                                x, y := v1, v2
x, y : [P, Q]
                                x, y \prefixcolon [~ P, Q ~]
:[P,Q]
                                \prefixcolon [~ P, Q ~]
{ P }
                               \{~ P ~\}
[P]
                                [~ P ~]
if P \longrightarrow A \parallel B fi
                                \circif P \circthen A \circelse B \circfi
\operatorname{\mathbf{do}} P \longrightarrow A \operatorname{\mathbf{od}}
                                \circdo P \circthen A \circod
\mathbf{con}\,X\bullet A
                                \circcon X \circspot A
\mathbf{var}\,x:T\bullet A
                                \circvar x : T \circspot A
```

Parameterised commands

 $\operatorname{val} x: T \bullet A$ \circval x : T \circspot A $\operatorname{res} x: T \bullet A$ \circspot A

```
\mathbf{vres}\,x:\,T\,\bullet\,A\qquad \qquad \texttt{\circvres}\,\,\mathtt{x}\,:\,\mathtt{T}\,\,\texttt{\circspot}\,\,\mathtt{A}
```

B.5 Circus processes

Process definition

```
\begin{array}{ll} \mathbf{process} \ n \ \widehat{=} \ PD & \text{ \  \  } \ \text{circprocess n \  \  } \ \text{ \  \  } \ \mathbf{PD} \\ \mathbf{process}[X] \ n \ \widehat{=} \ PD & \text{ \  \  } \ \mathbf{EX}] \ n \ \mathbf{PD} \end{array}
```

Basic process

```
\operatorname{process} n \stackrel{\frown}{=} \operatorname{begin} \dots BP \dots \operatorname{end}
```

\circprocess n \circdef \circbegin \ldots BP \ldots \circend

 $\begin{array}{ll} \mathbf{state} \; n \; == \; \mathbf{S} \\ \mathbf{state} \; S \end{array} \qquad \begin{array}{ll} \mathbf{circstate} \; \mathbf{n} \; == \; \mathbf{S} \\ \mathbf{state} \; S \end{array}$

Unary processes

 $P \setminus CS$ P\circhide CS

PName PName

 $\begin{array}{ll} PName[\mathbb{N}] & \text{PName[nat]} \\ PName(x,y) & \text{PName(x,y)} \\ PName[\mathbb{N}](x,y) & \text{PName[nat](x,y)} \end{array}$

 $PName \lfloor x \rfloor$ PName \lcircindex x \rcircindex

 $PName[\mathbb{N}]\lfloor x\rfloor$ PName[\nat] \lcircindex x \rcircindex

PName[c,d:=e,f] PName \lcircrename c, d := e, f \rcircrename

 $PName[\mathbb{N}][c,d:=e,f]$ PName[\nat] \lcircrename c, d := e, f \rcircrename

Binary processes

Parameterised and indexed processes

Parenthesised processes

```
(P)
                           (P)
(x:T\bullet P)
                           (x : T \circspot P)
(x:T\odot P)
                           (x : T \circindex P)
(P)[c := d]
                           (P) \lcircrename c := d \rcircrename
(x:T \bullet P)[c:=d]
                           (x : T \circspot P) \lcircrename c := d \rcircrename
(x:T\odot P)[c:=d]
                           (x : T \circindex P) \lcircrename c := d \rcircrename
(x:T\bullet P)(v)
                           (x : T \circspot P)(v)
                           (x : T \circindex P)(v)
(x:T\odot P)(v)
[X](x:X \bullet P)[\mathbb{N}](1)
                           [X](x : X \circspot P)[\nat](1)
[X](x:X\odot P)[\mathbb{N}](1)
                           [X](x : X \circindex P)[\nat](1)
(\mu X \bullet x : T \bullet P)(v)
                           (\circmu X \circspot x : T \circspot P)(v)
(\mu X \bullet (x : T \bullet P))(v)
                           (\circmu X \circspot (x : T \circspot P))(v)
(\mu X \bullet \mathbf{val} x : T \bullet P)(v)
                           (\circmu X \circspot \circval x : T \circspot P)(v)
```

Replicated processes

```
||||x:T \bullet P| \qquad || \text{Interleave x: T \circspot P} \\ ||x:T [\![ CS]\!] \bullet P \qquad || \text{Parallel x: T \lpar CS \rpar \circspot P} \\ || x:T \bullet P \qquad || \text{Intchoice x: T \circspot P} \\ || x:T \bullet P \qquad || \text{Extchoice x: T \circspot P} \\ || x:T \bullet P \qquad || \text{Semi x: T \circspot P} \\ || x:T \bullet P \qquad || \text{Semi x: T \circspot P} \\ || x:T \bullet P \qquad || \text{Semi x: T \circspot P} \\ || x:T \bullet P \qquad || \text{Semi x: T \circspot P} \\ || x:T \bullet P \qquad || \text{Semi x: T \circspot P} \\ || x:T \bullet P \qquad || x:T \bullet P \qquad || x:T \land P \qquad
```

B.6 Mathematical toolkits $SS \lozenge S$ SS \dcap S

Circus prelude

circus prore	440	Circus Spivey's Z bag toolkit	
\mathbb{B}	\boolean		
\mathbb{U}	\universe	bag X	\bag~X
True	\true	$B \sharp n$	B \bcount n
False	\false	$n\otimes B$	n \otimes B
		$B \uplus C$	B \uplus C
Circus mod	el checking toolkit	$B \cup C$	B \uminus C
$SS\ \c)\ TT$	SS \gendj TT	$x\to B$	x \inbag B
$\otimes SS$	\regions SS	$B \sqsubseteq C$	B \subbageq C
$SS \searrow S$	SS \dsetminus S	$\llbracket x, y rbracket$	\lbag x, y \rbag