

ZLive: A new Z Animator

The Community Z Tools Project
(CZT)

<http://czt.sourceforge.net>

Mark Utting, Petra Malik, Rohit Bansal



Agenda

- What is a Z animator?
- What use is a Z animator?
- Why build another Z animator?
- The Proposed Architecture
- Examples
- Questions and Comments

What is a Z animator?

- ♦ Goal: *Evaluate some Z expressions*

Search for one/all
concrete solutions

Not all Z expressions
are executable!
Some are infinite...

ISO Z
Standard (2002)

Set expressions,
Predicates ($\exists \forall \dots$),
Schemas (=Operations)

What use is a Z animator?

- Forward execution:

Rapid Prototyping

- *Init ; Incr ; Add[in? := 3] ; Sqrt ; ...*

- Backwards/sideways evaluation.

Finite theorem proving
or model checking:

- *Sqrt[value' := 4]*
 - *[Sqrt | status! ≠ ok]*

- Searching for a solution (\exists)
- Looking at ALL solutions (\forall)
- Searching for counter examples

- Testing an operation (provide inputs+outputs)

- *Sqrt[value := 16, value' := 4]*

- Playing around (students)

- $\{ x, y : 0..N \mid x * y = N \}$

Conclusion:

Validation
Verification
Experimentation

Why build another Z animator?

- ♦ Existing animators: limitations + unsupported
 - ♦ Evaluation time unpredictable and unbounded
 - ♦ Do not handle all Z constructs, especially Z Std.
- ♦ We want a Java-based, open-source animator that fits into CZT.
- ♦ Jaza was evaluated as best Z animator.
 - ♦ Goal: more *sophistication*, more *coverage*.
- ♦ It's a lot of fun trying to build one!



The Proposed Architecture

1) Read Z Specification

2) Read target Expr

3) Unfold defns (schema ops, toolkit etc.)

4) Flatten to atomic predicates

5) Calculate modes of each predicate

6) Reorder predicates into cheapest order

7) Enumerate all solutions (lazily)

repeat



Example 1

1) $\text{Pop} == [q, q' : \text{seq } \mathbb{N}; \text{out!} : \mathbb{N} \mid q = \langle \text{out!} \rangle \hat{\ } q']$

2) $\text{Pop}[q := \langle 10, 11, 12 \rangle]$

3) $\{ \text{Pop}$
 $\mid q = \{(1, 10), (2, 11), (3, 12)\}$
 $\bullet \ll q' == q', \text{out!} == \text{out!} \gg \}$

Then unfold Pop and 'normalise'

$\{ q, q' : \mathbb{P}(\mathbb{Z} \times \mathbb{Z}); \text{out!} : \mathbb{Z}$
 $\mid q \in \text{seq } \mathbb{N} \wedge q' \in \text{seq } \mathbb{N} \wedge \text{out!} \in \mathbb{N}$
 $\wedge q = \{(1, 10), (2, 11), (3, 12)\} \wedge q = \{(1, \text{out!})\} \hat{\ } q'$
 $\bullet \ll q' == q', \text{out!} == \text{out!} \gg \}$

Eg 1: Flatten

$$\begin{aligned} & \{ q, q': \mathbb{P}(\mathbb{Z} \times \mathbb{Z}); \text{out!} : \mathbb{Z} \\ & \mid q \in \text{seq } \mathbb{N} \wedge q' \in \text{seq } \mathbb{N} \wedge \text{out!} \in \mathbb{N} \\ & \wedge q = \{(1, 10), (2, 11), (3, 12)\} \wedge q = \{(1, \text{out!})\} \cap q' \\ & \bullet \ll q' == q', \text{out!} == \text{out!} \gg \} \end{aligned}$$

After flattening expressions to one operator per predicate:

4)
$$\begin{aligned} & \{ q, q': \mathbb{P}(\mathbb{Z} \times \mathbb{Z}); \text{out!} : \mathbb{Z} \\ & \mid q \in \text{seq } \mathbb{N} \wedge q' \in \text{seq } \mathbb{N} \wedge \text{out!} \in \mathbb{N} \\ & \wedge q = \{v1, v2, v3\} \wedge v1 = (1, 10) \wedge v2 = (2, 11) \wedge v3 = (3, 12) \\ & \wedge q = v4 \cap q' \wedge v4 = \{v5\} \wedge v5 = (1, \text{out!}) \\ & \bullet \ll q' == q', \text{out!} == \text{out!} \gg \} \end{aligned}$$

An aside: (5) modes

- Mode == $\{\text{InputVars}\} \rightarrow \{\text{OutputVars}\} \text{ [Num.of.Solns]}$
- Eg. $q=\{v1,v2,v3\}$ has modes:
 - $\{v1,v2,v3\} \rightarrow \{q\} \text{ [1]}$
 - $\{q\} \rightarrow \{v1,v2,v3\} \text{ [0.5]}$
 - $\{q,v1,v2,v3\} \rightarrow \{\} \text{ [0.5]}$
- Eg. $q=v4 \cap q'$ has modes:
 - $\{v4,q'\} \rightarrow \{q\} \text{ [1]}$
 - $\{q\} \rightarrow \{v4,q'\} \text{ [#q + 1]}$
 - $\{q,q'\} \rightarrow \{v4\} \text{ [1]}$
 - $\{q,v4\} \rightarrow \{q\} \text{ [1]}$
 - $\{q,v4,q'\} \rightarrow \{\} \text{ [0.5]}$
- Eg. $out! \in \mathbb{N}$ has modes:
 - $\{out!\} \rightarrow \{\} \text{ [0.5]}$
 - $\{\} \rightarrow \{out!\} \text{ } [\infty] \text{ (not usable, except for partial searches)}$

Eg 1: Reorder and Enumerate

- 4) $\{ q, q': \mathbb{P}(\mathbb{Z} \times \mathbb{Z}); \text{out!} : \mathbb{Z}$
| $q \in \text{seq } \mathbb{N} \wedge q' \in \text{seq } \mathbb{N} \wedge \text{out!} \in \mathbb{N}$
 $\wedge q = \{v1, v2, v3\} \wedge v1 = (1, 10) \wedge v2 = (2, 11) \wedge v3 = (3, 12)$
 $\wedge q = v4 \wedge q' \wedge v4 = \{v5\} \wedge v5 = (1, \text{out!})$
• $\ll q' == q', \text{out!} == \text{out!} \gg \}$

6) After reordering (output vars in bold blue):

- $\{ q, q': \mathbb{P}(\mathbb{Z} \times \mathbb{Z}); \text{out!} : \mathbb{Z}$
| $\mathbf{v1} = (1, 10) \wedge \mathbf{v2} = (2, 11) \wedge \mathbf{v3} = (3, 12) \wedge \mathbf{q} = \{v1, v2, v3\}$
 $\wedge q = \mathbf{v4} \wedge \mathbf{q'} \wedge v4 = \{\mathbf{v5}\} \wedge v5 = (1, \mathbf{out!})$
 $\wedge q \in \text{seq } \mathbb{N} \wedge q' \in \text{seq } \mathbb{N} \wedge \text{out!} \in \mathbb{N}$
• $\ll q' == q', \text{out!} == \text{out!} \gg \}$

7) Enumerate Solutions:

This has 4 solutions, but only one passes the next predicate.

Questions and Comments?

- How much unfolding?
 - Eg. $s \cup t == \{ x:T \mid x \in s \vee x \in t \}$ (unfold this toolkit defn?)
 - But can we deduce these modes for $r = s \cup t$?
 - $\{s,t\} \rightarrow \{r\}$ [1]
 - $\{r\} \rightarrow \{s,t\}$ [$2^{\#s} \times 2^{\#t}$]
 - $\{r,s\} \rightarrow \{t\}$ [$2^{\#s}$]
 - $\{r,t\} \rightarrow \{s\}$ [$2^{\#t}$]
- How best to handle disjuncts? DNF?
 - Too big and causes duplicate searching
 - + More accurate mode analysis of each branch.