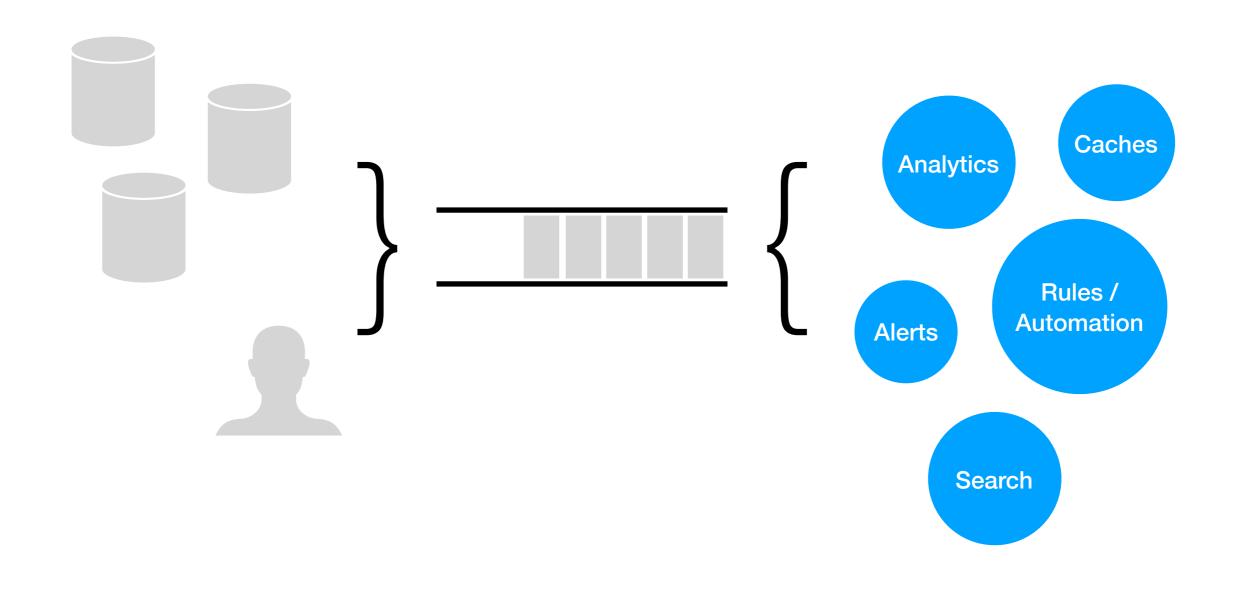
Declarative Differential Dataflows

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Producers Log Consumers

Benefits

- + decoupled systems (coordinate w/ data, stateless consumers)
- + operational scalability (O(n) pipes instead of O(n^2))
- + no data silos
- + proven in practice (Kafka, AWS Kinesis)

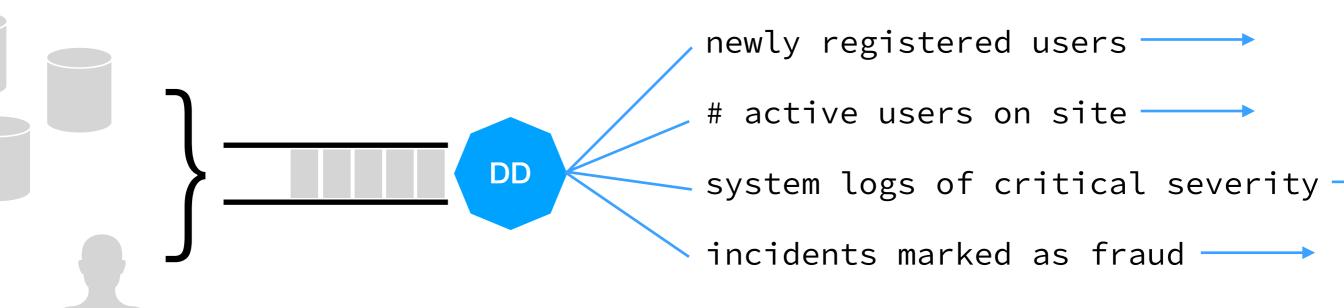
Challenges

- redundant filter logic in consumers (=> "Streaming SQL")
- change detection (diffing) handled by consumers
- have to keep up w/ producers (=> "Consumer Groups")
- weak processing primitives (no iteration)

Rephrased problem: Efficiently routing data change to many, heterogenous clients.

Approach

- consumers declare interest in data via Datalog queries
- utilise Differential Dataflow to run queries over the log
- outputs already diffs, ideal for maintaining derived views



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Differential Dataflow

```
/// BFS
let nodes = roots.map(|x| (x, 0));
nodes.iterate(|inner| {
 let edges = edges.enter(&inner.scope());
 let nodes = nodes.enter(&inner.scope());
 inner.join_map(&edges, |_k,l,d| (*d, l+1))
      .concat(&nodes)
      .group(|_{-}, s, t| t.push((*s[0].0, 1)))
})
```

Challenge I: Synthesise data flows w/o compilation?

```
/// BFS
let nodes = (roots.nap(|x| (x, 0));
nodes.iterate(|inner| {
 let édges = (edges enter(&inner.scope());
 let nodes enter(&inner.scope());
 inner.join_map(&edges, |_k,l,d| (*d, l+1))
      .concat(&nodes)
      .group(|_{-}, s, t| t.push((*s[0].0, 1)))
```

Differential collections are statically typed and too flexible for synthesis.

Solution: Unified data model

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```
;; unified data model

;; transactions

[123]: name "Alice" | Value
[123]: friend 456]
[456]: age 28]
[456]: authenticated? true]

***

**Retraction

**Retraction

**Retraction

**Retraction

**Assertion
```

Challenge II: Differential-compatible Query Language

```
/// BFS
let nodes = roots.map(|x| (x, 0));
nodes.iterate(|inner| {
 let edges = edges.enter(&inner.scope());
 let nodes = nodes.enter(&inner.scope());
 inner.join_map(&edges, |_k,l,d| (*d, l+1))
      .concat(&nodes)
      .group(|\_, s, t| t.push((*s[0].0, 1)))
})
```

```
;; Reachability

[(reaches ?x ?y) [?x :edge ?y]]
[(reaches ?x ?y) [?z :edge ?y] (reaches ?x ?z)]

[:find ?x ?y :where (reaches ?x ?y)]
```

```
;; Reachability
[(reaches ?x ?y) [?x :edge ?y]]
[(reaches ?x ?y) [?z :edge ?y] (reaches ?x ?z)]
[:find ?x ?y :where (reaches ?x ?y)]
JOIN
```

```
;; Reachability

[(reaches ?x ?y) [2x :edge ?y]]
[(reaches ?x ?y) [?z :edge ?y] (reaches ?x ?z)]

[:find ?x ?y :where (reaches ?x ?y)]
```

Challenge III: Constant bindings

```
;; all nodes reachable starting from node 100
[(reaches ?x ?y) [?x :edge ?y]]
[(reaches ?x ?y) [?z :edge ?y] (reaches ?x ?z)]
[:find ?y :where (reaches 100 ?y)]
```

Solution: Input transform

;; all nodes reachable starting from node 100

Challenge IV: Negation

:where

(and [?e :name "Alice"]

(not [?e :name "Alice"]))]

Solution? Negation ~ Set Difference

```
;; unbound negation
[:find ?e
 :where
 (not [?e :name "Alice"])]
;; contradictions
[:find ?e
 :where
 (and [?e :name "Alice"]
      (not [?e :name "Alice"]))]
```

Solution?? Negation ~ Anti-Join

```
;; unbound negation
[:find ?e
 :where
 (not [?e :name "Alice"])]
;; contradictions
[:find ?e
 :where
 (and [?e :name "Alice"]
      (not [?e :name "Alice"]))]
```

Solution??? Anti-Join against all tuples, push notclauses down.

```
;; unbound negation
                                   ;; tautologies
                                   [:find ?e
[:find ?e
 :where
                                    :where
 (not [?e :name "Alice"])]
                                    (or [?e :name "Alice"]
                                        (not [?e :name "Alice"]))]
;; contradictions
[:find ?e
 :where
 (and [?e :name "Alice"]
      (not [?e :name "Alice"]))]
```

Set Difference in Differential?

```
;; given unbound query...
[:find ?e
  :where
  (not [?e :name "Alice"])]

;; ...and some data

[[+1 100 :name "Alice"]]
```

Set Difference in Differential?

```
;; given unbound query...
[:find ?e
 :where
 (not [?e :name "Alice"])]
;; ...and some data
[[+1 100 :name "Alice"]]
                                    [[[{"Eid": 100}], -1]]
```

Challenge V: Union-compatibility

Solution: Explicit unification scope (add a projection)

Non-Challenge: Recursion

```
[:find ?x ?y :where (label ?x ?y)]
[[(label ?x ?y) [?x :label ?y]]
[(label ?x ?y) [?z :edge ?y] (label ?x ?z)]]
```

Non-Challenge: Recursion

```
[:find ?x ?y :where (label ?x ?y)]
     [[(label)?x ?y) [?x :label ?y]]
       (label ?x ?y) [?z :edge ?y] (label ?x ?z)]]
                                 everything else stays the same!
struct NamedRelation {
    variable: Variable<Vec<Value> ...>
    tuples: Collection<Vec<Value> ...>
}
```

3DF in Context

SOURCE OF TRUTH push new data read from log W_1 Datalog Query Plan Register Server W_2 push output diffs Control W_{i} **CLIENT** 3DF / K-Pg

Evaluation: Label Propagation

```
[:find ?x ?y :where (label ?x ?y)]

[[(label ?x ?y) [?x :label ?y]]
  [(label ?x ?y) [?z :edge ?y] (label ?x ?z)]]
```

System	Cores	HTTPD (169M)	POSTGRES (639M)
SociaLite Graspan Hand-rolled	4		00M 8628s 37s
our system	 1	151s	440s

Conclusions

with querying and diffing handled by Differential...

- + simple, declarative consumers
- + consumer performance must match throughput of relevant novelty, not throughput of overall log
- Differential's powerful processing primitives available for further processing