Information Compression #6 Basics of Krylov subspace methods

Youhei Yamaji

Department of Applied Physics, The University of Tokyo

- 1. The largest and smallest eigenvelues
- 2. Sparse matrix generated by Hamiltonian
- 3. Krylov subspace method



Classification of Information Compression in Linear Algebra by Memory Costs

- (1) A matrix can be stored
- -SVD for dense matrix
- -Compressed sensing (so far)
- (2) Although a matrix cannot be stored, vectors can be stored
- -SVD for sparse matrix
- -Krylov subspace method
- (3) A vector cannot be stored
- -Matrix product/tensornetwork states

This Week's Information Compression Algorithm

Main focus:

Algorithms that calculate specified eigenvalues and eigenvectors of huge* sparse matrices

You may not store your matrix A or you may not pay $O(L^3)^$ cost

$$A \in \mathbb{R}^{L \times L}$$

Especially the largest and smallest eigenstates

Largest and Smallest Eigenvalues

1. Ground state of quantum many-body system

$$\langle O \rangle = \frac{\vec{u}^{\dagger} O \vec{u}}{\vec{u}^{\dagger} \vec{u}}$$

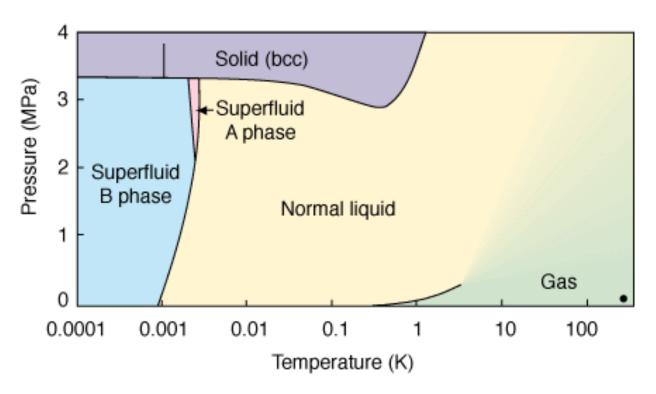
The ground state is important:

- -Room temperature is often enough low and well described by zero-temperature wave function
- -Interest in ground states (at zero temperature)

Low-temperature phase such as superfluid phase Zero-temperature phase transitions (quantum phase transition)

Low-Temperature Phases

Phase diagram of ³He



D. D. Osheroff, R. C. Richardson, and D. M. Lee, Phys. Rev. Lett. 28, 885 (1972).

Erkki Thuneberg http://ltl.tkk.fi/research/theory/helium.html

Largest and Smallest Eigenvalues

2. Principle component analysis for huge data Eigenvalue problem of covariance matrices

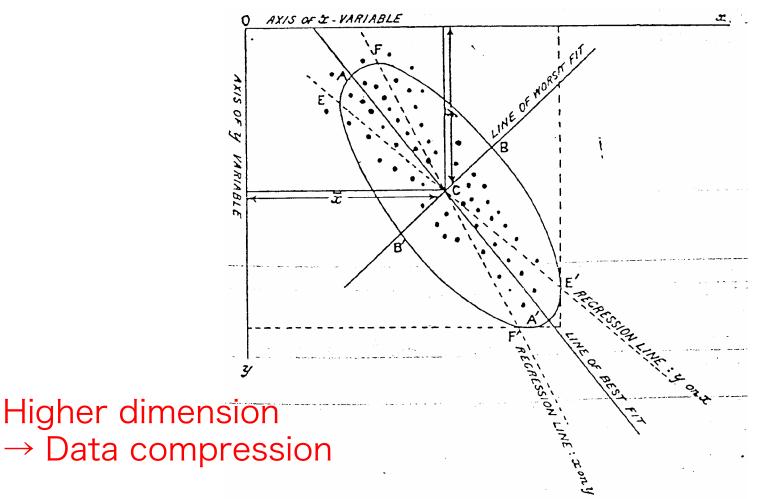
K. Pearson, Philosophical Magazine 2, 559 (1901)

$$\begin{bmatrix} \sum_{\ell} (x - \overline{x})^2 & \sum_{\ell} (x - \overline{x})(y - \overline{y}) \\ \sum_{\ell} (y - \overline{y})(x - \overline{x}) & \sum_{\ell} (y - \overline{y})^2 \end{bmatrix}$$

Largest and Smallest Eigenvalues

2. Principle component analysis for huge data

K. Pearson, Philosophical Magazine 2, 559 (1901)



Category of Numerical Linear Algebra

You need to choose algorithm depending on whether

- your matrix is 1) sparse/dense
 - 2) stored/not stored in memory

For a matrix that is dense and stored, you can find standard subroutines with $O(L^3)^*$ cost in LAPACK

*L is the linear dimension of your matrix A $A \in \mathbb{R}^{L \times L}$

Largest and Smallest Eigenvalues

Ground state of quantum many-body system

Typically, sparse and not stored

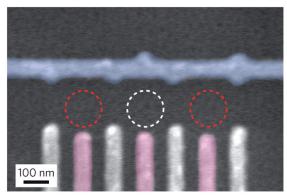
Principle component analysis for huge data Eigenvalue problem of covariance matrices Dense/sparse and stored/not stored

-Partial SVD/low-rank approximation will discussed in 8th lecture

Sparse Matrix Generated by Hamiltonian

Quantum dots

F. R. Braakman, et al., Nat. Nano. 8, 432 (2013)

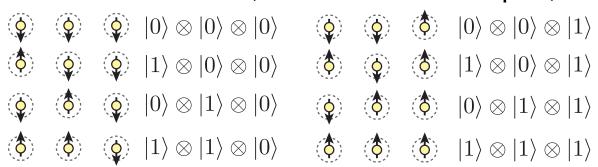


Quantum dot:

- -A quantum box can confine a single electron
- -Utilized for single electron transistor, quantum computers

Three-body problem:

 \rightarrow Number of states = 2^3 (factor 2 from spin)



superposition

States represented by superposition
$$\mathcal{F} = \{ \sum_{n_0=0,1} \sum_{n_1=0,1} \sum_{n_2=0,1} C_{n_0n_1n_2} |n_0\rangle \otimes |n_1\rangle \otimes |n_2\rangle : C_{n_0n_1n_2} \in \mathbb{C} \}$$

Mutual Interactions



1. Operators acting on a single qubit

A two dimensional representation of Lie algebra SU(2)

$$\begin{split} & [\hat{S}_{j}^{x}, \hat{S}_{j}^{y}] = i \hat{S}_{j}^{z} \\ & [\hat{S}_{j}^{y}, \hat{S}_{j}^{z}] = i \hat{S}_{j}^{x} \\ & [\hat{S}_{j}^{z}, \hat{S}_{j}^{x}] = i \hat{S}_{j}^{y} \end{split}$$

$$\hat{S}^y_j \ \hat{S}^y_j \ \hat{c}^z$$

-Commutator
$$[\hat{A},\hat{B}] \equiv \hat{A}\hat{B} - \hat{B}\hat{A}$$

$$\hat{S}_j^x|0\rangle = \frac{1}{2}|1\rangle$$

$$\hat{S}_j^x|1\rangle = \frac{1}{2}|0\rangle$$

$$\hat{S}_j^y|0\rangle = \frac{i}{2}|1\rangle$$

$$\hat{S}_j^y|1\rangle = -\frac{i}{2}|0\rangle$$

$$\hat{S}_j^z|1\rangle = \frac{1}{2}|1\rangle$$

$$\hat{S}_j^z|0\rangle = -\frac{1}{2}|0\rangle$$

Mutual Interactions

$$0 \quad 1 \quad N-1$$

Fock space of N qubits:

$$\mathcal{F} = \{ \sum_{n_0=0,1} \sum_{n_1=0,1} \cdots \sum_{n_{N-1}=0,1} C_{n_0 n_1 \cdots n_{N-1}} |n_0\rangle \otimes |n_1\rangle \otimes \cdots \otimes |n_{N-1}\rangle \}$$

$$(C_{n_0 n_1 \cdots n_{N-1}} \in \mathbb{C})$$

2. Operators acting on N-quibit Fock space:

$$\hat{S}_{j}^{a}, \hat{S}_{j}^{a} \hat{S}_{j+1}^{a} : \mathcal{F} \to \mathcal{F}$$

$$\hat{S}_{j}^{a} \doteq 1 \otimes \cdots \otimes 1 \otimes \hat{S}_{j}^{a} \otimes 1 \otimes \cdots \otimes 1$$

$$\hat{S}_{j}^{a} \hat{S}_{j+1}^{a} \doteq 1 \otimes \cdots \otimes 1 \otimes \hat{S}_{j}^{a} \otimes \hat{S}_{j+1}^{a} \otimes 1 \otimes \cdots \otimes 1$$

Quantum entanglement

Example: Two qubits



- -Superposition
- -Utilized for quantum teleportation cf.) EPR "paradox"

Mutual interactions between two qubits

$$\hat{H} = J \sum_{a=x,y,z} \hat{S}_0^a \hat{S}_1^a \quad (J \in \mathbb{R}, J > 0)$$

→Superposition (♦) (♦)









$$|1\rangle \otimes |0\rangle - |0\rangle \otimes |1\rangle$$

Hamiltonian Matrix

N-qubit Fock space:

$$\mathcal{F} = \{ \sum_{n_0 = 0, 1} \sum_{n_1 = 0, 1} \cdots \sum_{n_{N-1} = 0, 1} C_{n_0 n_1 \cdots n_{N-1}} | n_0 \rangle \otimes | n_1 \rangle \otimes \cdots \otimes | n_{N-1} \rangle \}$$

$$(C_{n_0 n_1 \cdots n_{N-1}} \in \mathbb{C})$$

Mutual interactions among N qubits:

Hamiltonian operator

$$\hat{H}:\mathcal{F}
ightarrow\mathcal{F}$$

$$\hat{H} = J \sum_{j=0}^{N-1} \sum_{a=x,y,z} \hat{S}_{j}^{a} \hat{S}_{\text{mod}(j+1,N)}^{a}$$

Vectors in Fock Space

Correspondence between spin and bit

$$|\uparrow\rangle = |1\rangle$$

$$|\downarrow\rangle = |0\rangle$$

 2^{N} -dimensional Fock space:

$$\mathcal{F} = \{ \sum_{n_0=0,1} \sum_{n_1=0,1} \cdots \sum_{n_{N-1}=0,1} C_{n_0 n_1 \cdots n_{N-1}} |n_0\rangle \otimes |n_1\rangle \otimes \cdots \otimes |n_{N-1}\rangle \}$$

$$(C_{n_0 n_1 \cdots n_{N-1}} \in \mathbb{C})$$

Decimal representation of orthonormalized basis

$$|I\rangle_{\rm d} = |n_0\rangle \otimes |n_1\rangle \otimes |n_2\rangle \otimes \cdots \otimes |n_{N-1}\rangle \qquad I = \sum_{\nu=0}^{N-1} n_{\nu} \cdot 2^{\nu}$$

Wave function as a vector

$$|\phi\rangle = \sum_{n_0=0}^{1} \sum_{n_1=0}^{1} \cdots \sum_{n_{N-1}=0}^{1} C_{n_0 n_1 \cdots n_{N-1}} |n_0\rangle \otimes |n_1\rangle \otimes \cdots \otimes |n_{N-1}\rangle$$

$$v(I) = C_{n_0 n_1 \cdots n_{N-1}} \qquad v(0:2^N-1)$$

Vectors and Matrices in Fock Space

Inner product of vectors

$$(\langle n_0 | \otimes \langle n_1 | \otimes \cdots \otimes \langle n_{N-1} |) \times (|n'_0\rangle \otimes |n'_1\rangle \otimes \cdots \otimes |n'_{N-1}\rangle)$$

$$= \langle n_0 | n'_0\rangle \times \langle n_1 | n'_1\rangle \times \cdots \times \langle n_{N-1} | n'_{N-1}\rangle$$

$$\langle n | \times | n'\rangle = \langle n | n'\rangle = \delta_{n,n'}$$

$$\langle \phi' | \phi \rangle = \sum_{n_0=0}^{1} \sum_{n_1=0}^{1} \cdots \sum_{n_{N-1}=0}^{1} C'^*_{n_0 n_1 \cdots n_{N-1}} C_{n_0 n_1 \cdots n_{N-1}}$$

$$|\phi'\rangle = \sum_{n_0=0}^{1} \sum_{n_1=0}^{1} \cdots \sum_{n_{N-1}=0}^{1} C'_{n_0 n_1 \cdots n_{N-1}} |n_0\rangle \otimes |n_1\rangle \otimes \cdots \otimes |n_{N-1}\rangle$$

$$|\phi\rangle = \sum_{n_0=0}^{1} \sum_{n_1=0}^{1} \cdots \sum_{n_{N-1}=0}^{1} C_{n_0 n_1 \cdots n_{N-1}} |n_0\rangle \otimes |n_1\rangle \otimes \cdots \otimes |n_{N-1}\rangle$$

Hamiltonian matrix

$$H_{II'} = \langle I|\hat{H}|I'\rangle$$

Orthonomalized basis: $|I\rangle, |I'\rangle \in \mathcal{F}$ $\langle I|I'\rangle = \delta_{I,I'}$

$$|I\rangle, |I'\rangle \in \mathcal{F}$$

$$\langle I|I'\rangle=\delta_{I,I'}$$

Sparse Matrix

- Particle or orbital number: N
- Fock space dimension: exp[N x const.]
- # of terms in Hamiltonian: Polynomial of N
- \rightarrow # of matrix elements of Hamiltonian matrix: (Polynomial of M) x exp[N x const.]

For sufficiently large N, (Polynomial of M) x exp[N x const.] << (exp[N x const.])²

Then, the Hamiltonian matrix is sparse

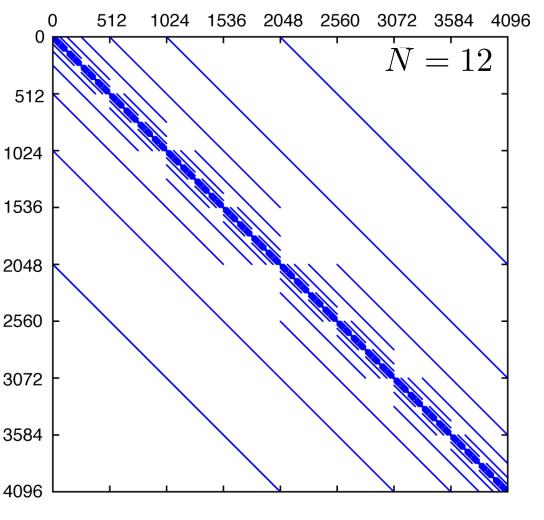
An Example of Hamiltonian Matrix

$$\hat{H} = J \sum_{i=0}^{N-1} \hat{S}_i^z \hat{S}_{i+1}^z - \Gamma \sum_{i=0}^{N-1} \hat{S}_i^x$$

-Non-commutative

$$\left[\sum_{i=0}^{N-1} \hat{S}_{i}^{z} \hat{S}_{i+1}^{z}, \sum_{i=0}^{N-1} \hat{S}_{i}^{x}\right] \neq 0$$

- →Quantum fluctuations or Zero point motion
- -Sparse # of elements $\propto O(2^N)$
- -Solvable
- -Hierarchical matrix?



Computational and Memory Costs

Matrix-vector product of dense matrix

$$v_i = \sum_{j=0}^{N_{\rm H}-1} A_{ij} u_j$$

Computational: O((Fock space dimension)²)

Memory: $O((\text{Fock space dimension})^2)$

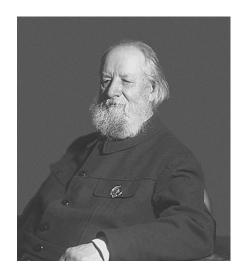
Matrix-vector product of large and sparse matrix

Computational: O(Fock space dimension)

Memory: O(Fock space dimension)

Hamiltonian is not stored in memory

Krylov Subspace Method for Sparse and Huge Matrices



Alexey Krylov
Aleksey Nikolaevich Krylov
1863-1945
Russian naval engineer and applied mathematician

Krylov subspace

$$A \in \mathbb{C}^{L \times L}$$

$$\mathcal{K}_n(A, \vec{b}) = \operatorname{span}\{\vec{b}, A\vec{b}, \dots, A^{n-1}\vec{b}\}$$

Numerical cost to construct K_n : $\mathcal{O}(\text{nnz}(A) \times n)$

Numerical cost to orthogonalize K_n : $\mathcal{O}(L \times n^2)$

Cornelius Lanczos 1950 Walter Edwin Arnoldi 1951 *nnz: Number of non-zero entries/elements

An Algorithm for Eigenvalue Problems of Large & Sparse Matrix: Power Method

Min. Eigenvalue of hermitian

Initial vector:
$$|v_1\rangle = \sum_{n=0}^{\infty} c_n |n\rangle$$

Prameter: $\max_{n} \{E_n\} \leq \Lambda$

$$\hat{H}|n\rangle = E_n|n\rangle$$

$$\langle n'|n\rangle = \delta_{n',n}$$

$$E_0 \le E_1 \le \cdots$$

$$\lim_{m \to +\infty} \frac{(\Lambda - \hat{H})^m |v_1\rangle}{\sqrt{\langle v_1 | (\Lambda - \hat{H})^{2m} |v_1\rangle}} = |0\rangle$$

$$(\Lambda - \hat{H})^m |v_1\rangle = \sum_n (\Lambda - E_n)^m c_n |n\rangle$$

$$\sum_{\substack{n > 0 \\ m \to +\infty}} (\Lambda - E_n)^{2m} |c_n|^2$$

$$\lim_{\substack{n > 0 \\ (\Lambda - E_0)^{2m} |c_0|^2}} = 0$$

Advanced Algorithm: Krylov Subspace Method

Krylov subspace method:

Finding approximate eigenstates in a Krylov subspace

$$\mathcal{K}_m(\hat{H}, |v_1\rangle) = \operatorname{span}\{|v_1\rangle, \hat{H}|v_1\rangle, \dots, \hat{H}^{m-1}|v_1\rangle\}$$

Construction and orthogonalization of Krylov subspaces

Shift invariance:

$$\mathcal{K}_m(\hat{H}, |v_1\rangle) = \mathcal{K}_m(\hat{H} + z\mathbf{1}, |v_1\rangle)$$

Krylov subspace method:

- -Lanczos method (symmetric/hermitian),
 - Arnoldi method (general matrix)
- -Conjugate gradient method (CG method) (many variation)

Initial:
$$\beta_1 = 0$$
, $|v_0\rangle = 0$
for $j = 1, 2, ..., m$ do
 $|w_j\rangle = \hat{H}|v_j\rangle - \beta_j|v_{j-1}\rangle$
 $\alpha_j = \langle w_j|v_j\rangle$
 $|w_j\rangle \leftarrow |w_j\rangle - \alpha_j|v_j\rangle$
 $\beta_{j+1} = \sqrt{\langle w_j|w_j\rangle}$
 $|v_{j+1}\rangle = |w_j\rangle/\beta_{j+1}$

$$\alpha_j = \langle v_j | \hat{H} | v_j \rangle$$
$$\beta_j = \langle v_{j-1} | \hat{H} | v_j \rangle = \langle v_j | \hat{H} | v_{j-1} \rangle$$

Orthogonalization

$$|v_{j}\rangle = \frac{\hat{H}|v_{j-1}\rangle - \sum_{\ell=1}^{j-1} |v_{\ell}\rangle \langle v_{\ell}|\hat{H}|v_{j-1}\rangle}{\langle v_{j}|\hat{H}|v_{j-1}\rangle}$$

$$\langle v_{\ell} | \hat{H} | v_{j-1} \rangle = \begin{cases} 0 & (\ell \le j - 3) \\ \beta_{j-1} & (\ell = j - 2) \\ \alpha_{j-1} & (\ell = j - 1) \end{cases}$$

Initial:
$$\beta_1 = 0$$
, $|v_0\rangle = 0$
for $j = 1, 2, ..., m$ do

$$|w_j\rangle = \hat{H}|v_j\rangle - \beta_j|v_{j-1}\rangle$$

$$\alpha_j = \langle w_j|v_j\rangle$$

$$|w_j\rangle \leftarrow |w_j\rangle - \alpha_j|v_j\rangle$$

$$\beta_{j+1} = \sqrt{\langle w_j|w_j\rangle}$$

$$|v_{j+1}\rangle = |w_j\rangle/\beta_{j+1}$$

$$\alpha_{j} = \langle v_{j} | \hat{H} | v_{j} \rangle$$

$$\langle v_{j} | v_{k} \rangle = \delta_{j,k}$$

$$\beta_{j} = \langle v_{j-1} | \hat{H} | v_{j} \rangle = \langle v_{j} | \hat{H} | v_{j-1} \rangle$$

Hamiltonian projected onto m D Krylov subsace

Eigenvalues of projected Hamiltonian

→ Approximate eigenvalues of original Hamiltonian

Lanczos Method: # of Vectors Required

Initial:
$$\beta_1 = 0$$
, $|v_0\rangle = 0$
for $j = 1, 2, ..., m$ do
$$|w_j\rangle \leftarrow \hat{H}|v_j\rangle - \beta_j|v_{j-1}\rangle \qquad |v_{j-1}\rangle \rightarrow |w_j\rangle, |v_j\rangle$$

$$\alpha_j = \langle w_j|v_j\rangle \qquad |w_j\rangle, |v_j\rangle$$

$$|w_j\rangle \leftarrow |w_j\rangle - \alpha_j|v_j\rangle \qquad |w_j\rangle, |v_j\rangle$$

$$\beta_{j+1} = \sqrt{\langle w_j|w_j\rangle} \qquad |w_j\rangle, |v_j\rangle$$

$$|v_{j+1}\rangle = |w_j\rangle/\beta_{j+1} \qquad |w_j\rangle \rightarrow |v_{j+1}\rangle, |v_j\rangle$$

Convergence of Lanczos Method

Yousef Saad, Numerical Methods for Large Eigenvalue Problems (2nd ed) The Society for Industrial and Applied Mathematics 2011

Assumption:
$$\lambda_1 > \lambda_2 > \cdots > \lambda_n$$

Convergence theorem for the largest eigenvalue

$$0 \leq \lambda_{1} - \lambda_{1}^{(m)} \leq (\lambda_{1} - \lambda_{n}) \left[\frac{\tan \theta(|v_{1}\rangle, |0\rangle)}{C_{m-1}(1+2\gamma_{1})} \right]^{2}$$

$$\sim 4(\lambda_{1} - \lambda_{n}) \left[\tan \theta(|v_{1}\rangle, |0\rangle) \right]^{2} e^{-4\sqrt{\gamma_{1}}m}$$

$$\gamma_{1} = \frac{\lambda_{1} - \lambda_{2}}{\lambda_{2} - \lambda_{n}}$$

$$C_{k}(t) = \frac{1}{2} \left[\left(t + \sqrt{t^{2} - 1} \right)^{k} + \left(t + \sqrt{t^{2} - 1} \right)^{-k} \right]_{29}$$

Distribution of Eigenvalues of Hermitian Matrices

An important relationship between distribution or density of states and statistical mechanics

$$P(E) = \frac{\rho(E)e^{-\beta E}}{\int dE' \rho(E')e^{-\beta E'}} \sim \frac{\exp[-(E - \langle E \rangle)^2/2CT^2]}{\sqrt{2\pi CT^2}}$$

$$k_{\rm B} = 1$$

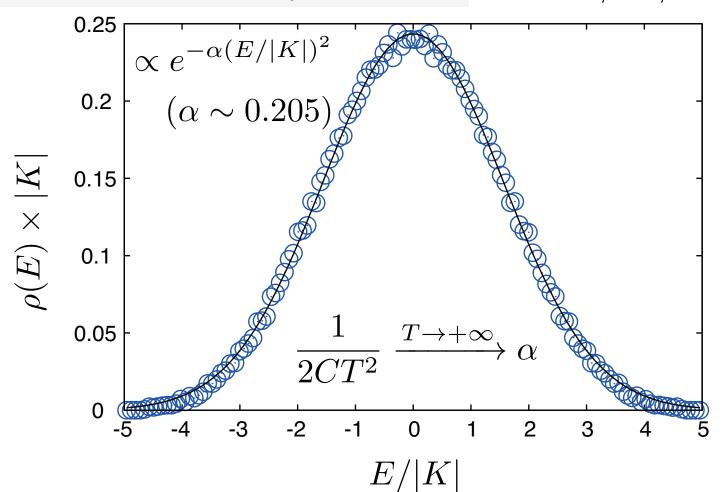
$$C = \frac{\langle \hat{H}^2 \rangle - \langle \hat{H} \rangle^2}{T^2}$$

$$\langle \hat{H}^m \rangle \sim \int E^m P(E) dE$$

An Example of Density of State

24 site cluster of Kitaev model (frustrated S = 1/2 spins)

A. Kitaev, Annals Phys. 321, 2 (2006). $2^{24} = 16,777,216$



Example of Dense Matrix: Random Symmetric Matrices

Eugene P. Wigner, Annals of Mathematics, 2nd Series, 67, 325 (1958)

Wigner's random matrix
$$(A)_{ij}=a_{ij}$$
 (Not necessarily sparse)
$$\int p_{ij}(a)da=1$$
 $p_{ij}(+a)=p_{ij}(-a)$ $\langle a_{ij}^n\rangle=\int p_{ij}(a)a^nda\leq B_n$ $\langle a_{ij}^2\rangle=\int p_{ij}(a)a^2da=1$

Example of Dense Matrix: Random Symmetric Matrices

Eugene P. Wigner, Annals of Mathematics, 2nd Series, 67, 325 (1958)

Density of states of $L \times L$ symmetric random matirx

$$A\vec{v} = E\vec{v}$$

$$\sigma(E) = \begin{cases} \frac{\sqrt{4L - E^2}}{2\pi L} & (E^2 < 4L) \\ 0 & (E^2 > 4L) \end{cases}$$

Comment:

Sparse matrices in quantum many-body problems show smaller density of states than random matrices around the both ends of the distribution

- → Sparse around maximum/minimum eigenvalues
- → Lanczos method may work well

Approximate SVD by Krylov Subspace Method

Low-rank approximation by block Krylov subspace

C. Musco & C. Musco, NIPS'15 Proceedings of 28th International Conference on Neural Information Processing Systems 1, 1396 (2015)

$$\|A - ZZ^TA\|_2 \le (1+\epsilon)\|A - A_k\|_2 \qquad \text{Operator norm defined by 2-norm (Spectral norm)}$$

$$A \in \mathbb{R}^{L \times M} \ Z \in \mathbb{R}^{L \times k} \ \text{rank } k \leq L, M$$

 $q = \mathcal{O}(\ln d/\sqrt{\epsilon})$

random matrix $\Pi \in \mathbb{R}^{M \times k}$

$$\mathcal{K}_{q+1} = \operatorname{span}\{A\Pi, (AA^T)A\Pi, \dots, (AA^T)^qA\Pi\}$$

 $Q \in \mathbb{R}^{N imes qk}$ Orthogonalized basis set of the block Krylov subspace

$$M = Q^T A A^T Q \in \mathbb{R}^{qk \times qk}$$

 U_k : the top k singular vectors of M

$$Z = QU_k$$

$$(\Pi)_{ij}$$
: Random number generated by $e^{-x^2/2}/\sqrt{\pi}$

Next Week

第7回: 物質科学における情報圧縮 7 Information compression in condensed matter physics

Important References

Yousef Saad, Numerical Methods for Large Eigenvalue Problems (2nd ed) The Society for Industrial and Applied Mathematics 2011

Next Week

第1回: 現代物理学における巨大なデータ

第2回: 現代物理学と情報圧縮

第3回: 情報圧縮の数理1 (線形代数の復習)

第4回: 情報圧縮の数理2 (特異値分解と低ランク近似)

第5回: 情報圧縮の数理3 (スパース・モデリングの基礎)

第6回: 情報圧縮の数理4 (クリロフ部分空間法の基礎)

11/15 第7回: 物質科学における情報圧縮

11/15 #7 Information compression in materials science

第8回: データ解析の高速化:スパース・モデリングの物質科学への応用

第9回: データ空間の圧縮:クリロフ部分空間法の物質科学への応用

第10回: 高度なデータ圧縮:情報のエンタングルメントと行列積表現

第11回: 行列積表現の固有値問題への応用

第12回: テンソルネットワーク表現への発展

第13回: テンソルネットワーク繰り込みによる情報圧縮