

4ICT9

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# Cells and Cell Design Issues

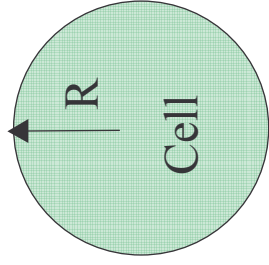
# Outline

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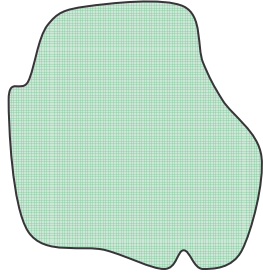
- Cell Shape
  - Actual cell/Ideal cell
- Signal Strength
- Handoff Region
- Cell Capacity
  - Traffic theory
  - Erlang B and Erlang C
- Cell Structure
- Frequency Reuse
- Reuse Distance
- Cochannel Interference
- Cell Splitting
- Cell Sectoring

# Cell Shape

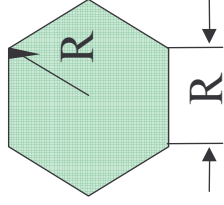
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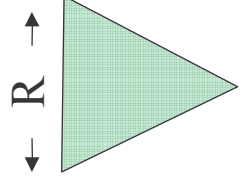
(a) Ideal cell



(b) Actual cell



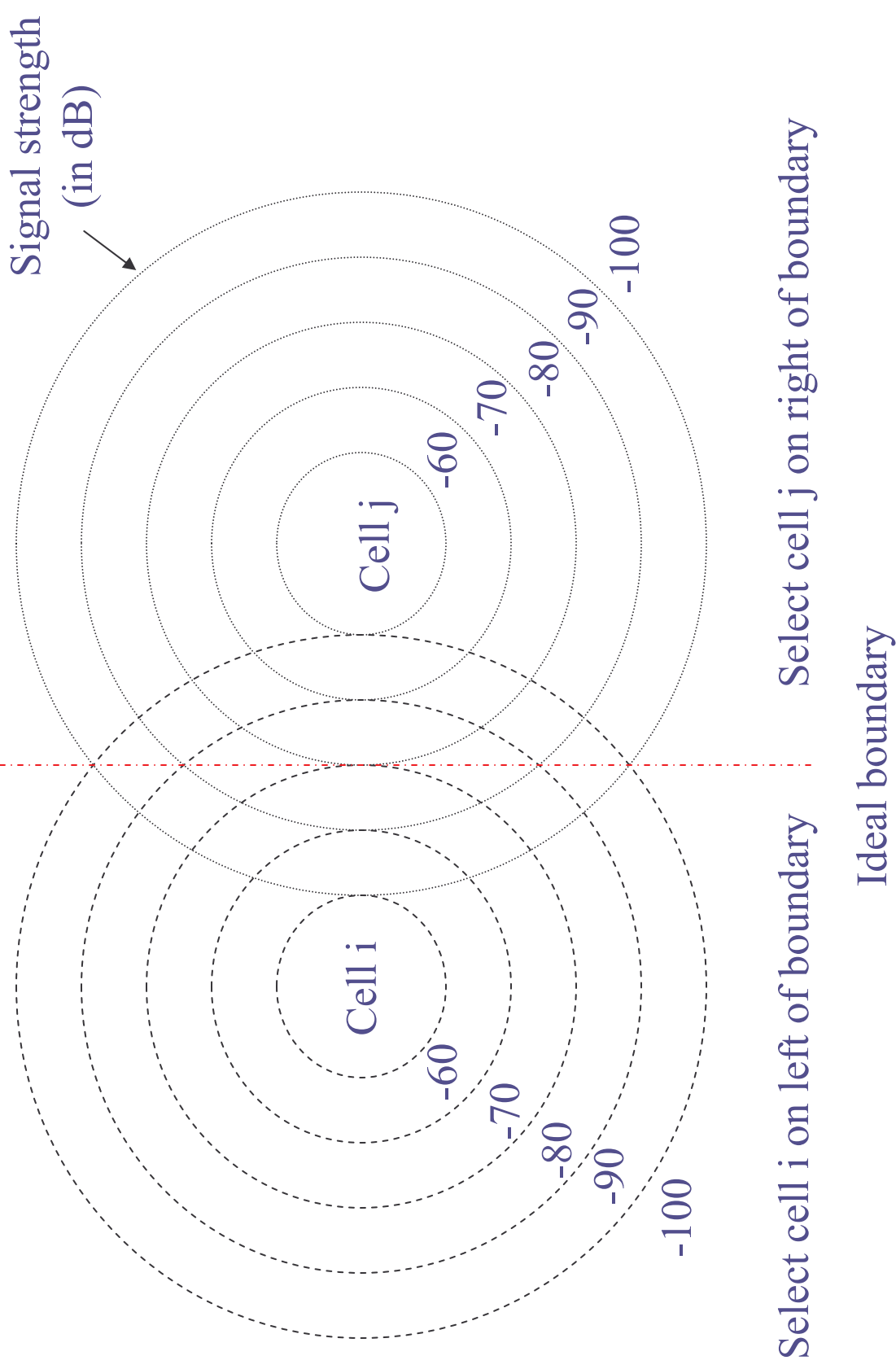
(c) Different cell models



# Impact of Cell Shape and Radius on Service Characteristics

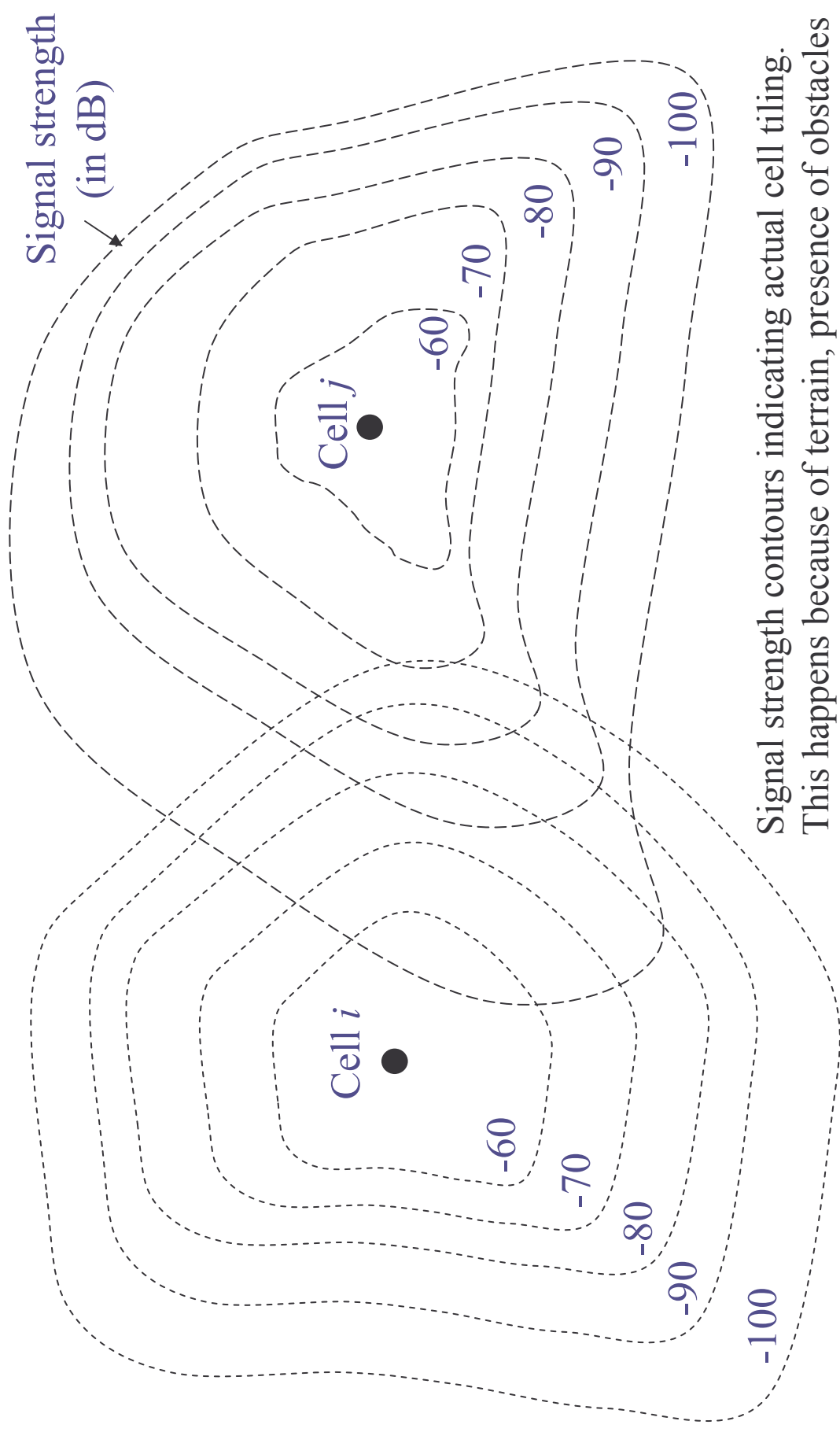
Shape of the Cell	Area	Boundary	Boundary Length/Unit Area	Channels/Unit Area with $N$ Channels/Cells	Channels/Unit Area when Number of Channels Increased by a Factor $K$	Channels/Unit Area when Size of Cell Is Reduced by a Factor $M$
Square cell (side = $R$ )	$R^2$	$4R$	$\frac{4}{R}$	$\frac{N}{R^2}$	$\frac{KN}{R^2}$	$\frac{K^2N}{R^2}$
Hexagonal cell (side = $R$ )	$\frac{3\sqrt{3}}{2}R^2$	$6R$	$\frac{4}{\sqrt{3}R}$	$\frac{N}{1.5\sqrt{3}R^2}$	$\frac{KN}{1.5\sqrt{3}R^2}$	$\frac{K^2N}{1.5\sqrt{3}R^2}$
Circular cell (radius = $R$ )	$\pi R^2$	$2\pi R$	$\frac{2}{R}$	$\frac{N}{\pi R^2}$	$\frac{KN}{\pi R^2}$	$\frac{K^2N}{\pi R^2}$
Triangular cell (side = $R$ )	$\frac{\sqrt{3}}{4}R^2$	$3R$	$\frac{4\sqrt{3}}{R}$	$\frac{4\sqrt{3}N}{3R^2}$	$\frac{4\sqrt{3}KN}{3R^2}$	$\frac{4\sqrt{3}K^2M^2N}{3R^2}$

# Signal Strength



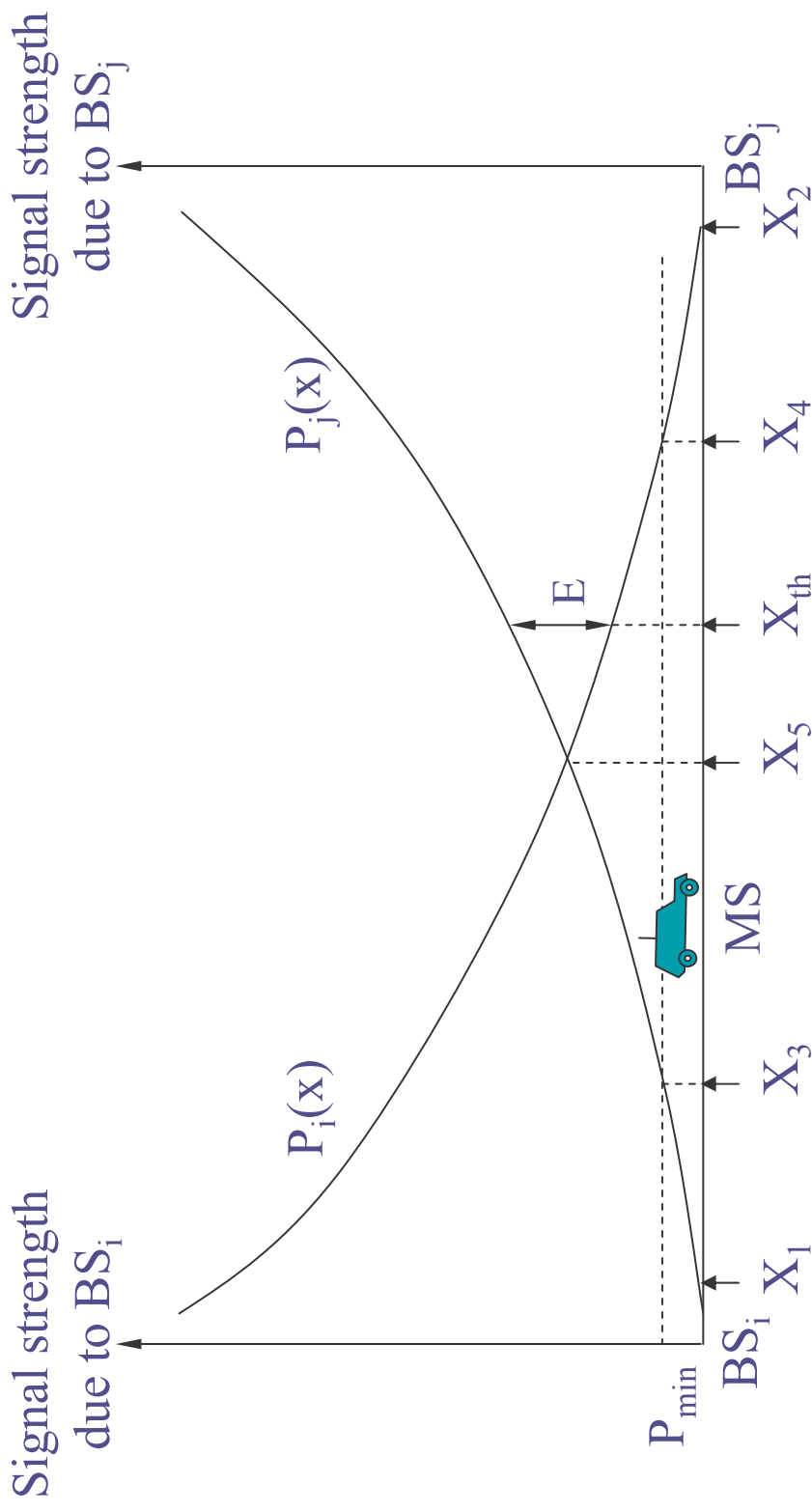
# Signal Strength

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Signal strength contours indicating actual cell tiling.  
This happens because of terrain, presence of obstacles  
and signal attenuation in the atmosphere.

# Handoff Region



- By looking at the variation of signal strength from either base station it is possible to decide on the optimum area where handoff can take place.

# Handoff Rate in a Rectangular

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Since handoff can occur at sides  $R_1$  and  $R_2$  of a cell

$$\lambda_H = R_1(X_1 \cos \theta + X_2 \sin \theta) + R_2(X_1 \sin \theta + X_2 \cos \theta)$$

where  $A = R_1 R_2$  is the area and assuming it constant, differentiate with respect to  $R_1$  (or  $R_2$ ) gives

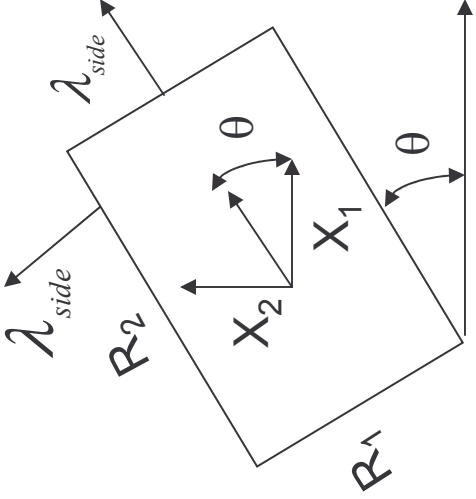
$$R_1^2 = A \frac{X_1 \sin \theta + X_2 \cos \theta}{X_1 \cos \theta + X_2 \sin \theta} \quad R_2^2 = A \frac{X_1 \cos \theta + X_2 \sin \theta}{X_1 \sin \theta + X_2 \cos \theta}$$

Total handoff rate is

$$\lambda_H = 2\sqrt{A(X_1 \cos \theta + X_2 \sin \theta)(X_1 \sin \theta + X_2 \cos \theta)}$$

$\lambda_H$  is minimized when  $\theta=0$ , giving

$$\lambda_H = 2\sqrt{AX_1X_2} \quad \text{and} \quad \frac{R_1}{R_2} = \frac{X_1}{X_2}$$





# Cell Capacity

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- Average number of MSs requesting service (Average arrival rate):  $\lambda$
  - Average length of time MS requires service (Average holding time):  $T$
  - Offered load:  $a = \lambda T$
- e.g., in a cell with 100 MSs, on an average 30 requests are generated during an hour, with average holding time  $T=360$  seconds.

Then, arrival rate  $\lambda=30/3600$  requests/sec.

A channel kept busy for one hour is defined as one Erlang (a), i.e.,

$$a = \frac{30 \text{ Calls}}{3600 \text{ Sec}} \cdot \frac{360 \text{ Sec}}{\text{call}} = 3 \text{ Erlangs}$$

# Cell Capacity

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- Average arrival rate during a short interval  $t$  is given by  $\lambda t$
- Assuming Poisson distribution of service requests, the probability  $P(n, t)$  for  $n$  calls to arrive in an interval of length  $t$  is given by

$$P(n, t) = \frac{(\lambda t)^n}{n!} e^{-\lambda t}$$

- Assuming  $\mu$  to be the service rate, probability of each call to terminate during interval  $t$  is given by  $\mu t$ .  
Thus, probability of a given call requires service for time  $t$  or less is given by

$$S(t) = 1 - e^{-\mu t}$$

# Erlang B and Erlang C

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- Probability of an arriving call being blocked is

$$B(S, a) = \frac{a^S}{S!} \cdot \frac{1}{\sum_{k=0}^S a^k}, \quad \longleftarrow \text{Erlang B formula}$$

where  $S$  is the number of channels in a group.

- Probability of an arriving call being delayed is

$$C(S, a) = \frac{a^S \overline{(S-1)!(S-a)}}{a^S \overline{(S-1)!(S-a)} + \sum_{i=0}^{S-1} a^i}, \quad \longleftarrow \text{Erlang C formula}$$

where  $C(S, a)$  is the probability of an arriving call being delayed with  $a$  load and  $S$  channels.

# Efficiency (Utilization)

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$$\text{Efficiency} = \frac{\text{Traffic nonblocked}}{\text{Capacity}}$$

$$= \frac{\text{Erlangs} \times \text{portions of nonrouted traffic}}{\text{Number of trunks (channels)}}$$

- Example: for previous example, if  $S=2$ , then

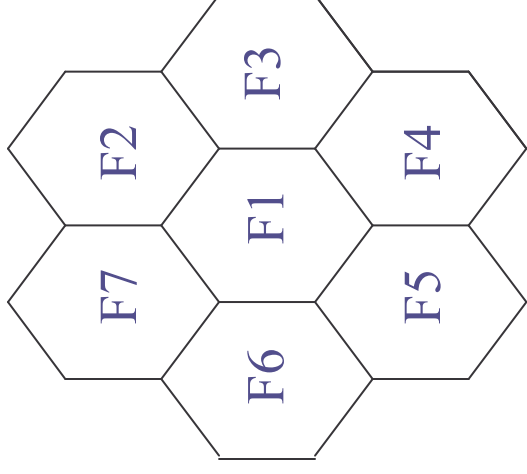
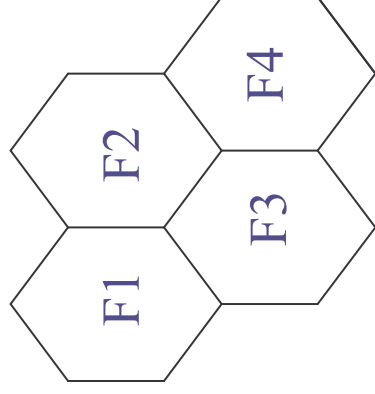
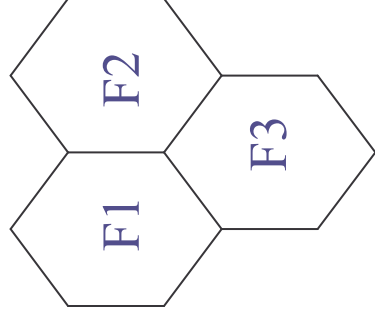
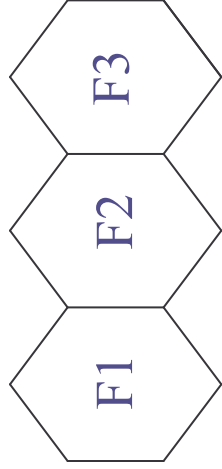
$B(S, a) = 0.6$ , ----- Blocking probability,  
i.e., 60% calls are blocked.

Total number of rerouted calls =  $30 \times 0.6 = 18$

$$\text{Efficiency} = 3(1-0.6)/2 = 0.6$$

# Cell Structure

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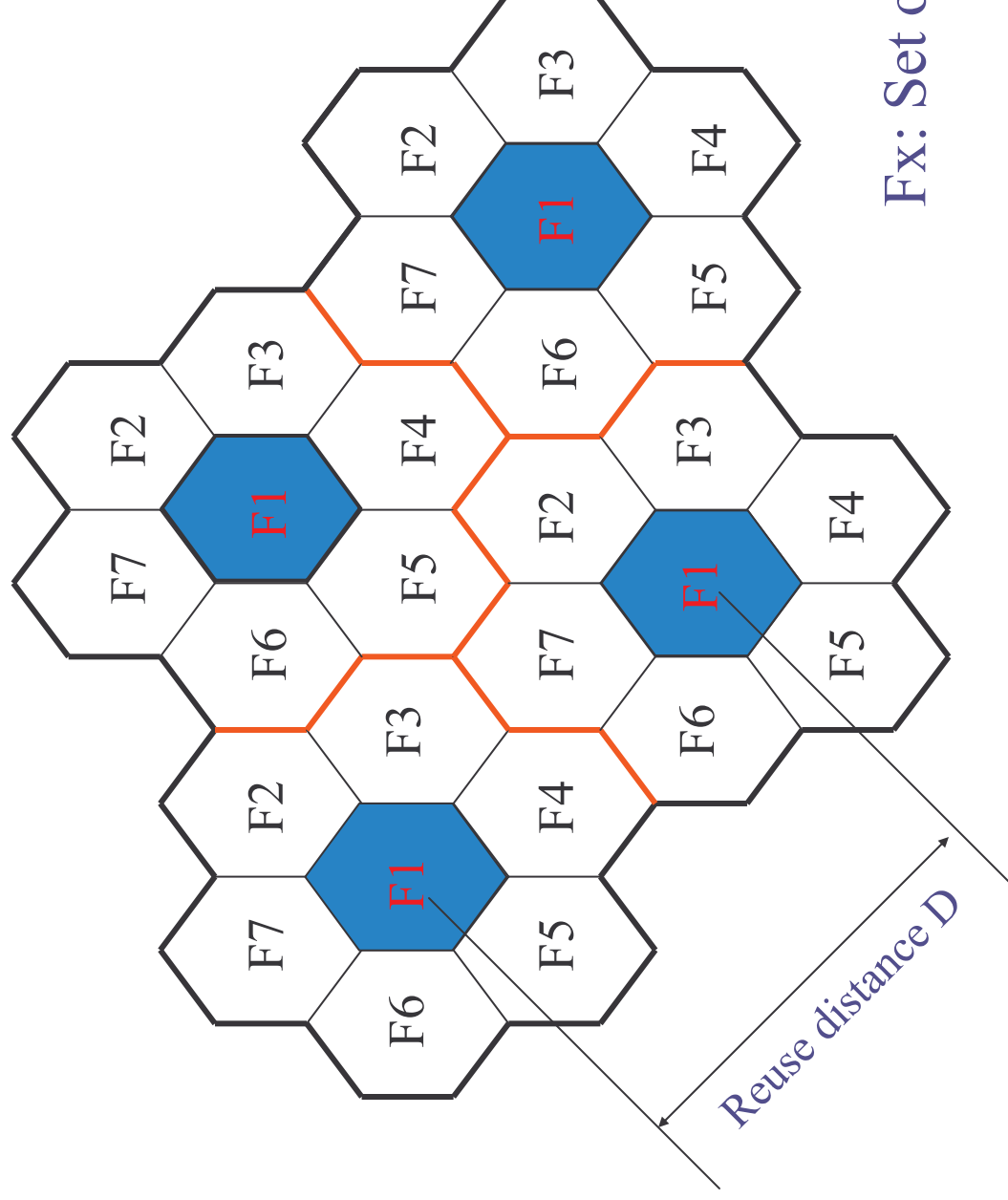


(a) Line Structure

(b) Plan Structure

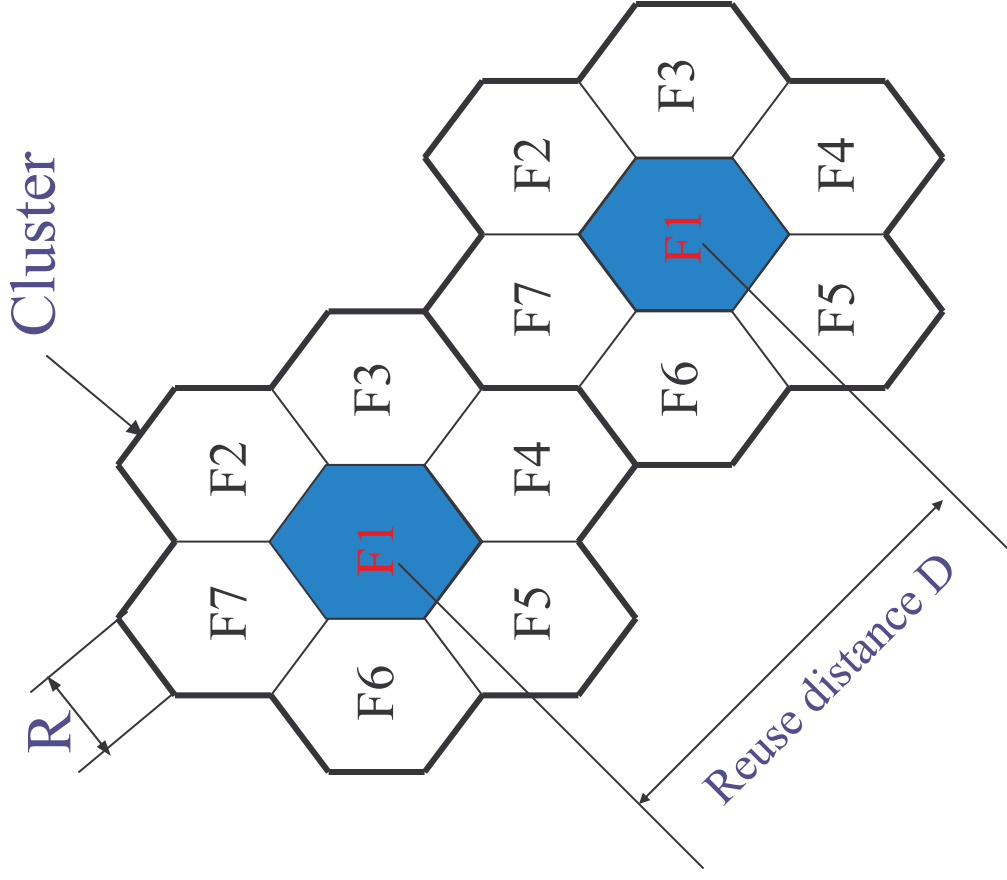
Note:  $F_x$  is set of frequency, i.e., frequency group.

# Frequency Reuse



7 cell reuse cluster

# Reuse Distance



- For hexagonal cells, the reuse distance is given by

$$D = \sqrt{3N}R$$

where  $R$  is cell radius and  $N$  is the reuse pattern (the cluster size or the number of cells per cluster).

- Reuse factor is

$$q = \frac{D}{R} = \sqrt{3N}$$

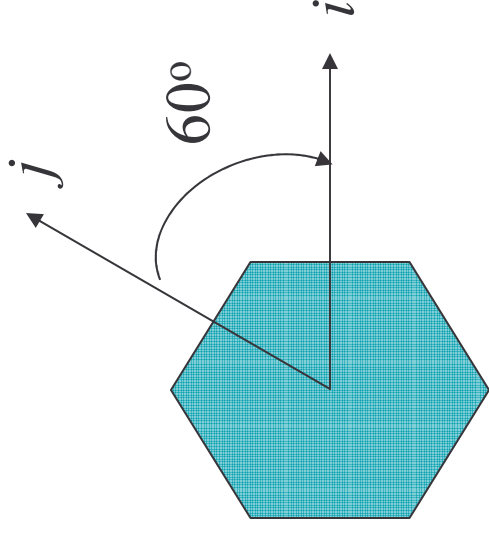
## Reuse Distance (Cont'd)

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- The cluster size or the number of cells per cluster is given by

$$N = i^2 + ij + j^2$$

where  $i$  and  $j$  are integers.

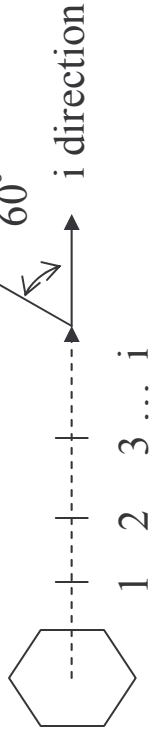


- $N = 1, 3, 4, 7, 9, 12, 13, 16, 19, 21, 28, \dots$ , etc.

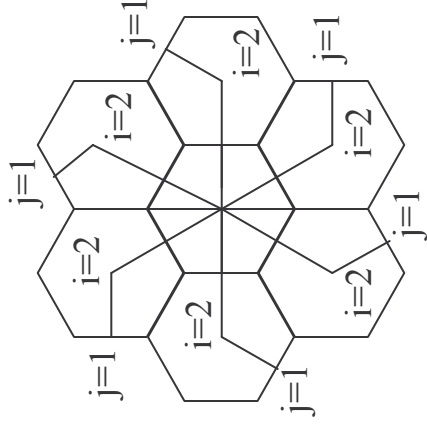
The popular value of  $N$  being 4 and 7.



# Reuse Distance (Cont'd)



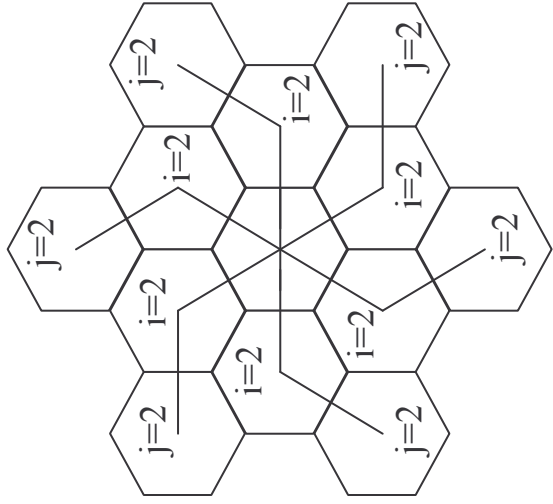
(a) Finding the center of an adjacent cluster using integers  $i$  and  $j$  (direction of  $i$  and  $j$  can be interchanged).



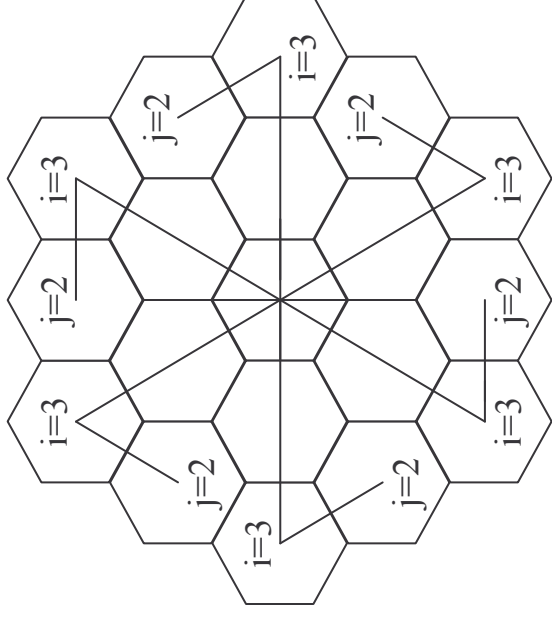
(b) Formation of a cluster for  $N = 7$  with  $i=2$  and  $j=1$

# Reuse Distance (Cont'd)

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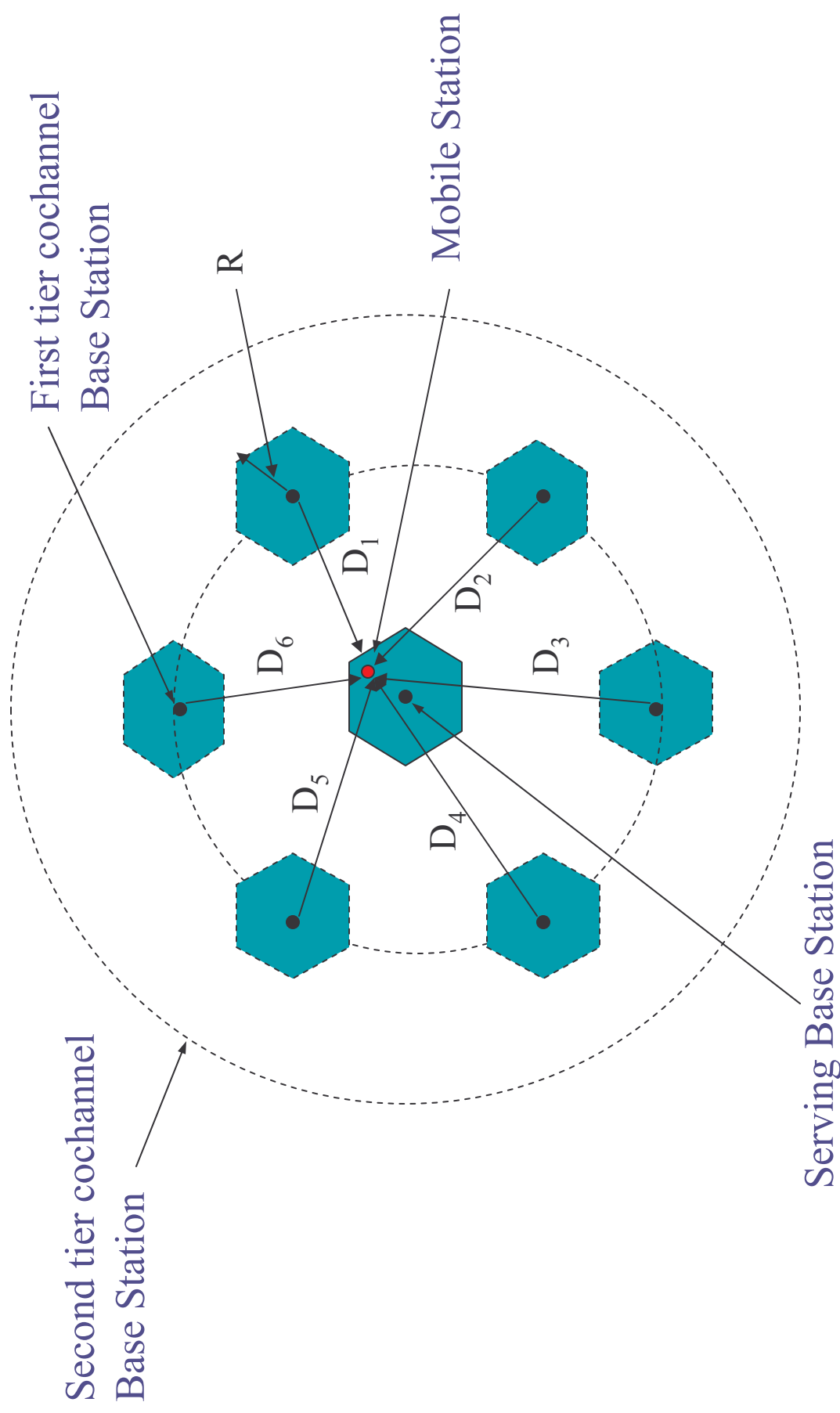


(c) A cluster with  $N=12$  with  $i=2$  and  $j=2$



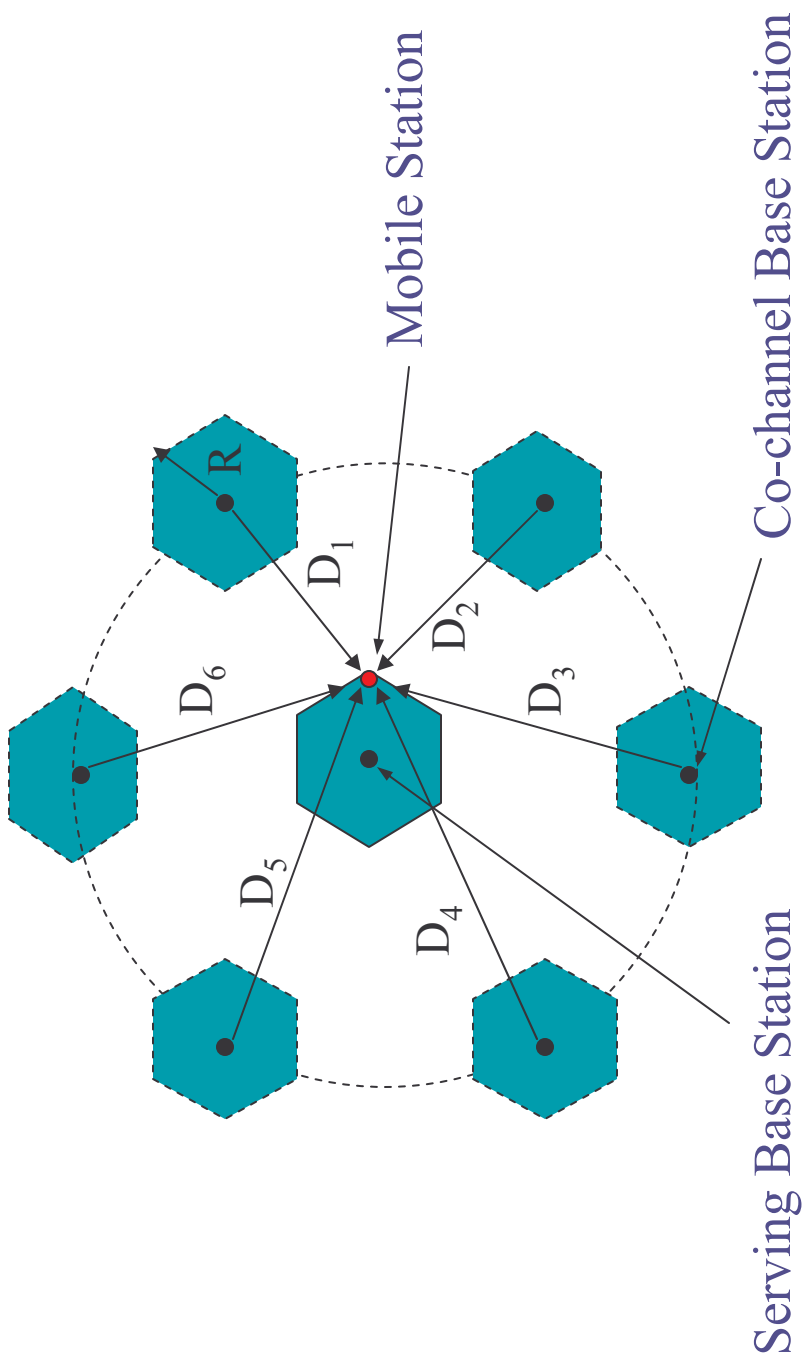
(d) A Cluster with  $N=19$  cells with  $i=3$  and  $j=2$

# Cochannel Interference



# Worst Case of Cochannel Interference

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# Cochannel Interference

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- Cochannel interference ratio is given by

$$\frac{C}{I} = \frac{\text{Carrier}}{\text{Interference}} = \frac{C}{\sum_{k=1}^M I_k}$$

where  $I$  is co-channel interference and  $M$  is the maximum number of co-channel interfering cells.

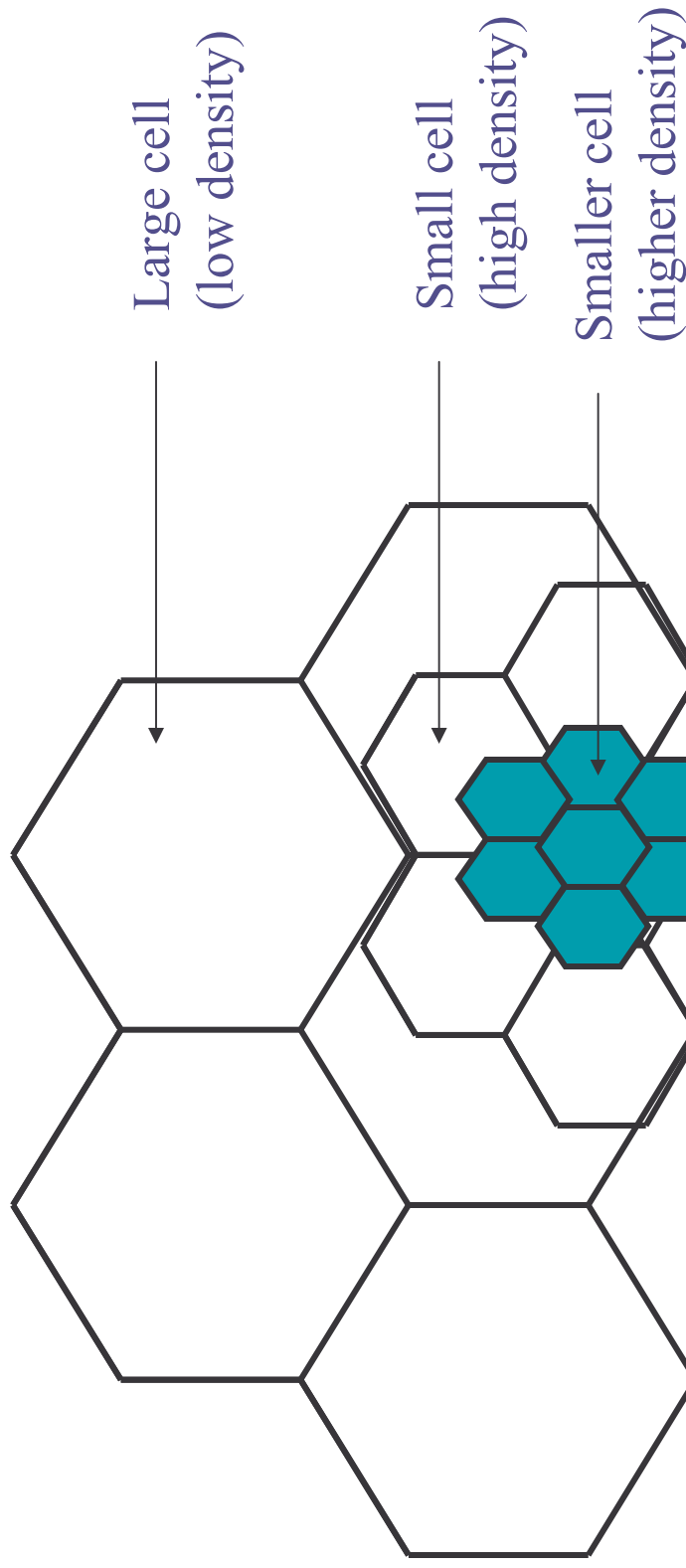
For  $M = 6$ ,  $C/I$  is given by

$$\frac{C}{I} = \frac{C}{\sum_{k=1}^M \left( \frac{D_k}{R} \right)^{-\gamma}}$$

where  $\gamma$  is the propagation path loss slope  
and  $\gamma = 2 \sim 5$ .

# Cell Splitting

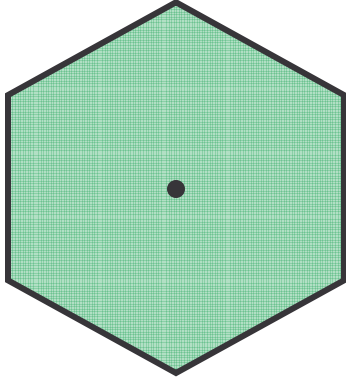
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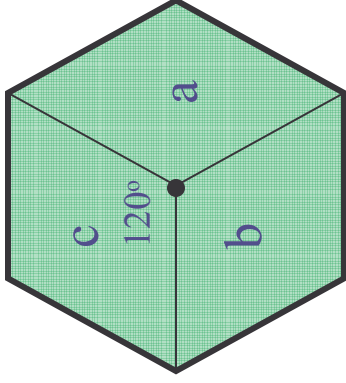
Depending on traffic patterns the smaller cells may be activated/deactivated in order to efficiently use cell resources.

# Cell Sectoring by Antenna Design

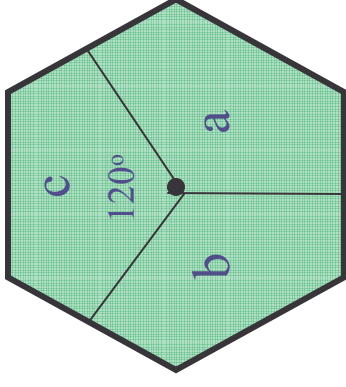
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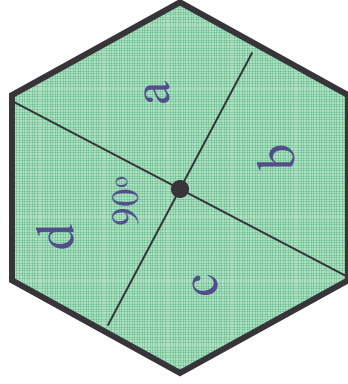
(a). Omni



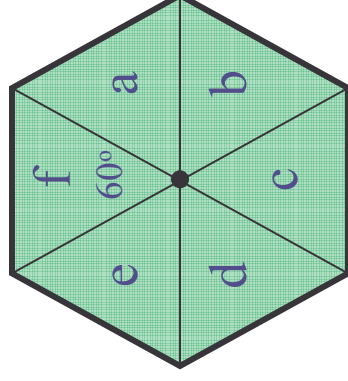
(b). 120° sector



(c). 120° sector (alternate)



(d). 90° sector

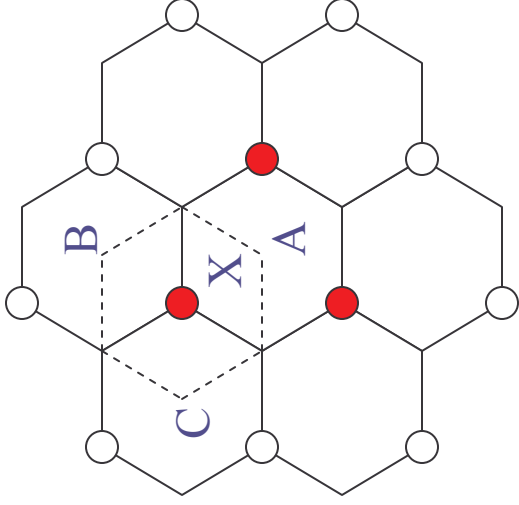


(e). 60° sector

# Cell Sectoring by Antenna Design

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- Placing directional transmitters at corners where three adjacent cells meet

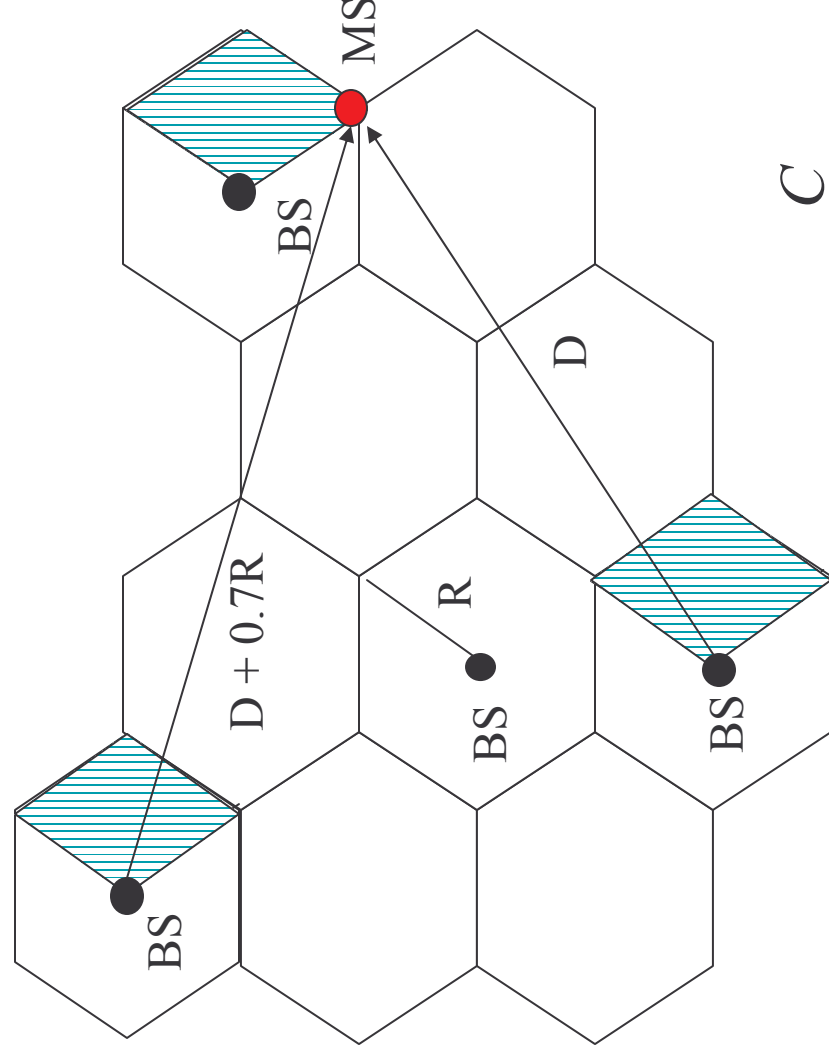




# Worst Case for Forward Channel

## Interference in Three-sectors

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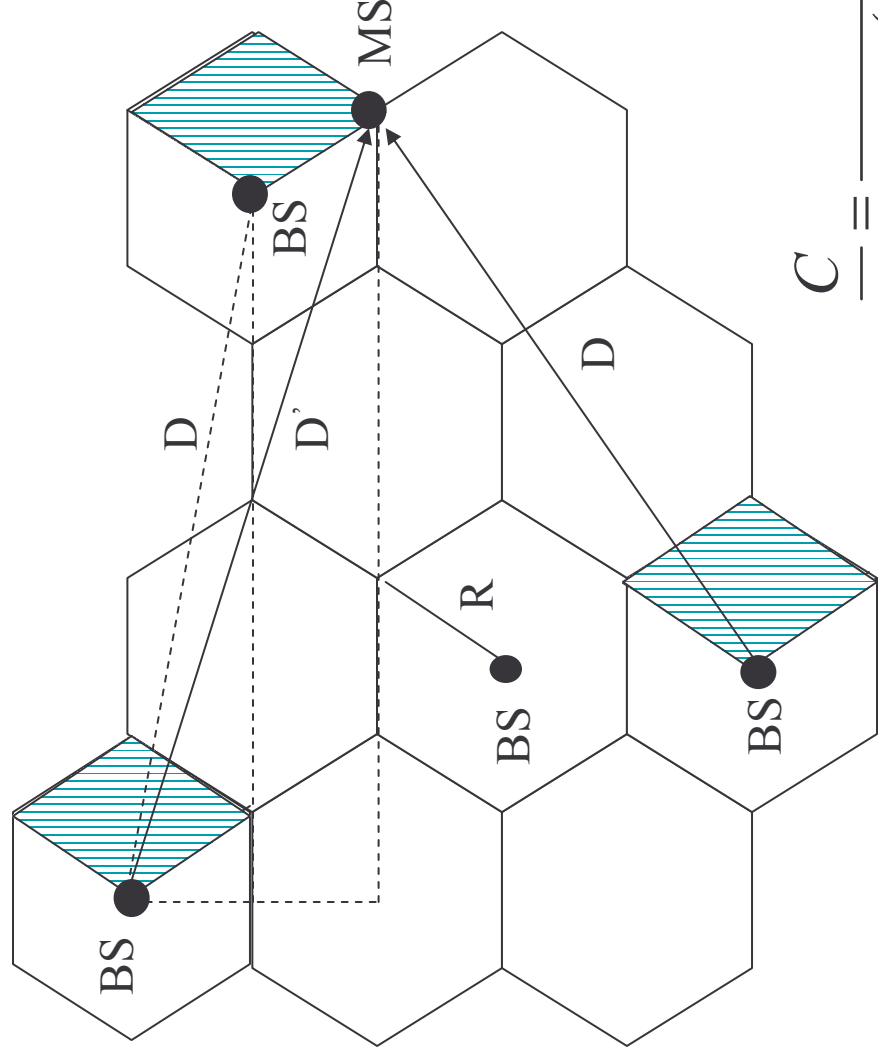
$$\frac{C}{I} = \frac{C}{q^{-\gamma} + (q + 0.7)^{-\gamma}}$$

$$q = D / R$$

# Worst Case for Forward Channel

## Interference in Three-sectors (Cont'd)

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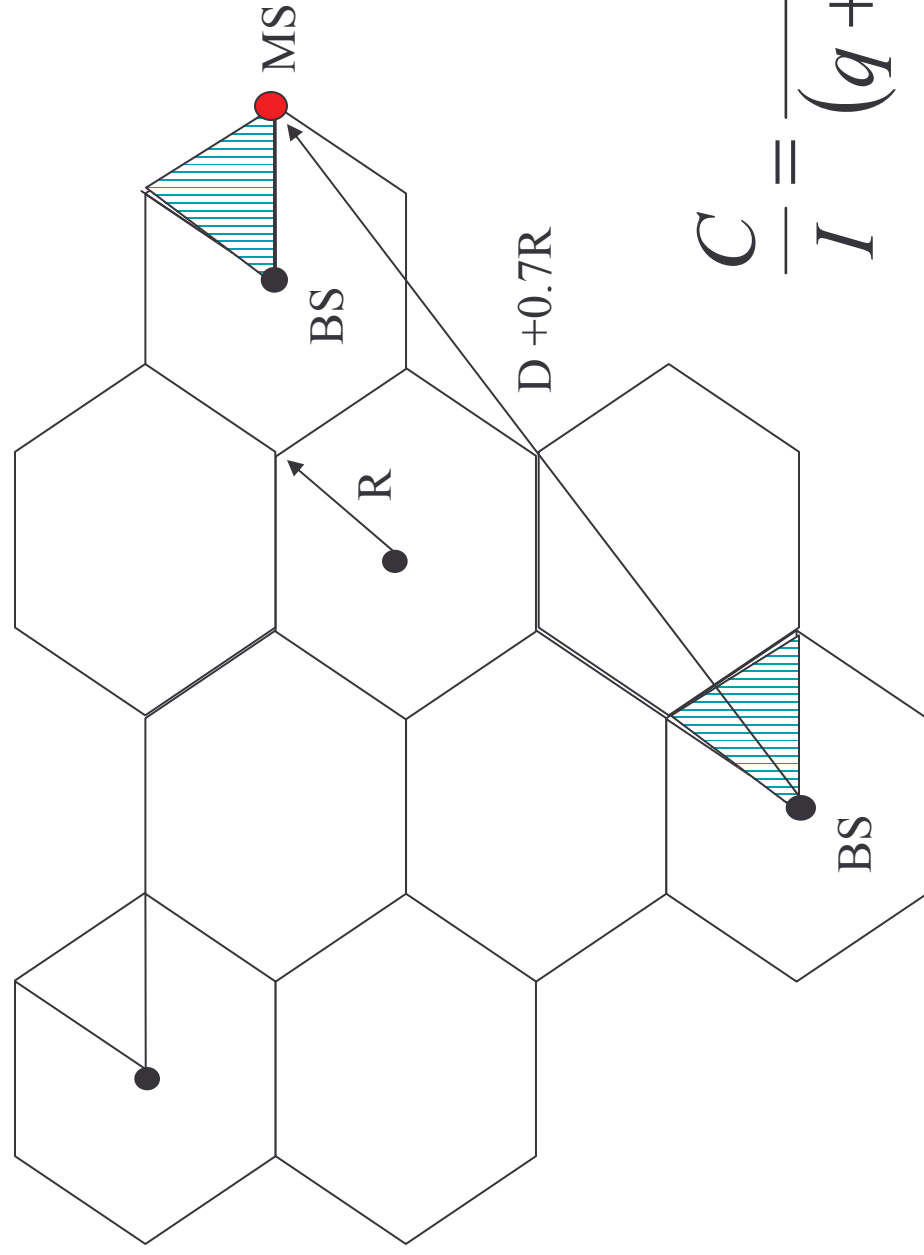
$$\frac{C}{I} = \frac{C}{q^{-\gamma} + (q + 0.7)^{-\gamma}}$$

$$q = D / R$$

# Worst Case for Forward Channel

## Interference in Six-sectors

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$$\frac{C}{I} = \frac{C}{(q + 0.7)^{-\gamma}}$$

$$q = D/R$$