

# UNIVERSITY OF DUBLIN

## TRINITY COLLEGE

Faculty of Engineering and Systems Sciences  
Department of Computer Science

B.A.(Mod.) Computer Science  
Senior Freshman Examination

Trinity Term 1999

### *2BA2 – Programming Techniques*

**Monday 31<sup>st</sup> May**

**Exam Hall**

**14.00 - 17.00**

Dr. Hugh Gibbons

Attempt FOUR questions

(In presenting programs explain clearly the design of the Eiffel code)

## Qs 1

- i) Present an Eiffel function

```
is_ordered(a : ARRAY[G]; L, H : INTEGER) : BOOLEAN
  require
    a /= void and L <= H;
    a.lower <= L and H <= a.upper
  ensure
    -- ( All i | L ≤ i < H : a.item(i) ≤ a.item(i+1))
```

that will check whether an array segment,  $a[L..H]$ , is ordered.

- ii) Present an Eiffel routine

```
Search(a:ARRAY[G]; L,H:INTEGER; x:G) is
  require
    Ordered: Is_Ordered(a,L,H)
  ensure
    --(found → x = a.item(index)) & (¬found → x ∉ a[L..H])
```

that binary searches an array section,  $a[L..H]$ , for an item,  $x$ .

- iii) Assume we have an attribute,  $a:ARRAY[G]$ , in a class.

Present a routine

```
reverse(L,H:INTEGER)
  require
    Ordered: Is_Ordered(a,L,H)
  ensure
    -- ( All i | L ≤ i < H : a.item(i) ≥ a.item(i+1))
```

that reverses the items in array segment,  $a[L..H]$ . Ensure that none of the items in the array are overwritten.

## Qs 2

A Matrix  $M$  has a Saddle Point iff for some position  $(i,j)$ ,  $M(i,j)$  is the minimum of row  $i$  and maximum of column  $j$ .

- a) Show that all saddle points have the same value,  
i.e. if  $M(i,j)$  and  $M(s,t)$  are Saddle Points then  $M(i,j) = M(s,t)$ .
- b) Assume we have the following classes for VECTOR and MATRIX with short forms,

```
class interface
  VECTOR [G -> COMPARABLE]
creation
  make
feature
  arr_copy (a: ARRAY [G])
    -- copy an array, a, into current vector

  item (i: INTEGER): G

  make (n: INTEGER)

  max_index: INTEGER
    -- an index for maximum item

  min_index: INTEGER
    -- an index for minimum item

  put (x: G; i: INTEGER)

  size: INTEGER
end -- class VECTOR
```

```
class interface
  MATRIX [G -> COMPARABLE]
creation
  make
feature
  cols: INTEGER -- # columns

  item (i, j: INTEGER): G

  make (r, c: INTEGER)

  max_row: VECTOR [G]
    -- vector of maximums for each row

  min_row: VECTOR [G]
    -- vector of minimums for each row

  put (x: G; i, j: INTEGER)

  put_row (a: ARRAY [G]; i: INTEGER)

  rows: INTEGER -- # rows

  transpose: like Current
end -- class MATRIX
```

- i) Present an Eiffel routine  
    **one\_saddle** (m: MATRIX [INTEGER])  
    that will find the location of a saddle point in matrix, m.
- ii) Present an Eiffel routine  
    **all\_saddle** (m: MATRIX [INTEGER])  
    that will find the location of all saddle points in matrix, m.

### Qs 3

Assume we have classes, LIST\_BAG and PAIR with short forms

```
class interface LIST_BAG [G]
```

```
  add(x : G)
```

```
  -- Add x, maybe again
```

```
  join(other : LIST_BAG[G])
```

```
  -- join to the end of current
```

```
  count : INTEGER
```

```
  empty : BOOLEAN
```

```
  has (x : G) : BOOLEAN
```

```
  remove (x : G)
```

```
  copy(s: LIST_BAG[G])
```

```
  -- traversal routines
```

```
  item : G    -- item at cursor
```

```
  start  -- set cursor back to start
```

```
  first -- first item in list
```

```
  forth  -- move cursor forward
```

```
  off : Boolean -- Cursor beyond end?
```

```
end -- class LIST_BAG
```

```
class interface PAIR[G]
```

```
  first, second : G
```

```
  set_first(x :G) is
```

```
    set_second(x :G)
```

```
end -- PAIR
```

Present an Eiffel class that will provide a routine for sorting a LIST\_BAG object via the algorithm for quicksort, i.e. provide routines that will quicksort a list. Use linked list diagrams to explain the routines.

Suggestion:

Provide a routine/function

```
  partition(s : LIST_BAG[G]; pivot : G) : PAIR[LIST_BAG[G]]
```

that will partition a list, s, into a pair of lists and a routine/function

```
  quicksort(s : LIST_BAG[G]) : LIST_BAG[G]
```

that will do the sorting.

#### Qs 4

A derangement of 1..n is a permutation, p, of 1..n such that  $p(i) \neq i$ .

- a) Write out all the derangements of 1..4.
- b) The number of derangements of 1..n can be given by the recursive function
$$\begin{aligned} \text{Der}(1) &= 0 \\ \text{Der}(2) &= 1 \\ \text{Der}(n) &= (n-1)(\text{Der}(n-1) + \text{Der}(n-2)), \text{ for } n > 2. \end{aligned}$$
Present an iterative/non-recursive function that will calculate  $\text{Der}(n)$ .
- c) Present an Eiffel routine that will generate all the derangements of 1..n

#### Qs 5

Assume class interfaces for a Binary Search Tree class, BST, and a BIN\_NODE class as follows: (Note: No repeated items in an object of type BST.)

```
class interface
  BST [G -> COMPARABLE]
  -- Binary Search Tree
feature
  empty -- make current empty
  is_empty: BOOLEAN
  inorder: ARRAY [G]
    require
      not is_empty
  add (x: G) -- add x, if not in current.
  array2bst(a: ARRAY [G])
    -- add all items of array, a, to current.
    -- repeat items are not added.
  inorder : ARRAY [G]
    -- puts items in tree into an array
  remove (x: G)
    require
      not is_empty
  root: BIN_NODE [G]
  count: INTEGER -- # items in tree
end -- class BST
```

```
class interface
  BIN_NODE [G]
feature
  value: G;
  left: BIN_NODE [G];
  right: BIN_NODE [G];
  value_set (v: G)
  left_set (n: BIN_NODE [G])
  right_set (n: BIN_NODE [G])
  build (v: G; l, r: BIN_NODE [G])
end -- class BIN_NODE
```

- a) The routine, `add(x : G)`, which uses an auxiliary routine, `insert`, is as follows:

```
Add(x:G) is
do
    if root /= void then
        insert(x,root)
    else
        !!root
        root.build(x,void,void)
        count := 1
    end
end -- Add
```

Implement the routine,

`insert(x : G; t : BIN_NODE[G])`

that inserts an item, `x`, into a tree with root, `t`, if it is not already there.

- b) Implement the routine,

`array2tree(a:ARRAY[G])`

that adds all the items in the array, `a`, to the binary search tree, with no repetitions.

- c) Implement the auxiliary routine,

`inord(t:BIN_NODE[G])`

which is used in the routine, `inorder`, as follows:

```
inorder : ARRAY[G] is
require
    not is_empty
do
    !!a_ord.make(1,size)
    count_ord := 1
    inord(root)
    result := clone(a_ord)
end -- inorder
```

The array, `a_ord`, and integer, `count_ord`, are hidden attributes of the class.

The routine, `inord`, `inorder` traverses the tree, starting at the root, and adds the items of the tree to the array, `a_ord`.

#### Qs. 6.

Assume that a Directed Acyclic Graph (DAG), `D`, is stored as an adjacency list.

- a) Present an Eiffel routine that will Breadth First Traverse the DAG, `D`.

- b) Present an Eiffel routine that will Topological Sort the DAG, `D`.