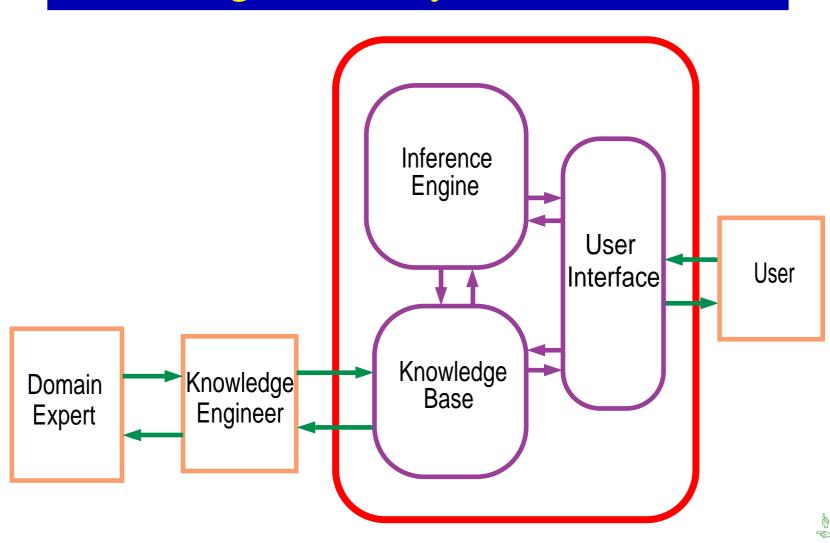
# Knowledge Engineering

#### Overview:

- ➤ How representation and reasoning systems interact with humans.
- Roles of people involved in a RRS.
- ➤ Building RRSs using meta-interpreters.
- ➤ Knowledge-based interaction and debugging tools

## Knowledge-based system architecture



## Roles for people in a KBS

- Software engineers build the inference engine and user interface.
- Knowledge engineers design, build, and debug the knowledge base in consultation with domain experts.
- Domain experts know about the domain, but nothing about particular cases or how the system works.
- Users have problems for the system, know about particular cases, but not about how the system works or the domain.

# Implementing Knowledge-based Systems

To build an interpreter for a language, we need to distinguish

- Base language the language of the RRS being implemented.
- Metalanguage the language used to implement the system.

They could even be the same language!

# Implementing the base language

Let's use the definite clause language as the base language and the metalanguage.

- We need to represent the base-level constructs in the metalanguage.
- We represent base-level terms, atoms, and bodies as meta-level terms.
- ➤ We represent base-level clauses as meta-level facts.
- In the non-ground representation base-level variables are represented as meta-level variables.

#### Representing the base level constructs

- Base-level atom  $p(t_1, \ldots, t_n)$  is represented as the meta-level term  $p(t_1, \ldots, t_n)$ .
- Meta-level term  $oand(e_1, e_2)$  denotes the conjunction of base-level bodies  $e_1$  and  $e_2$ .
- Meta-level constant *true* denotes the object-level empty body.
- The meta-level atom clause(h, b) is true if "h if b" is a clause in the base-level knowledge base.



### Example representation

The base-level clauses

connected\_to(
$$l_1, w_0$$
).  
connected\_to( $w_0, w_1$ )  $\leftarrow$  up( $s_2$ ).  
 $lit(L) \leftarrow light(L) \wedge ok(L) \wedge live(L)$ .

can be represented as the meta-level facts

```
clause(connected\_to(l_1, w_0), true).
clause(connected\_to(w_0, w_1), up(s_2)).
clause(lit(L), oand(light(L), oand(ok(L), live(L)))).
```



# Making the representation pretty

- Use the infix function symbol "&" rather than *oand*.
  - $\rightarrow$  instead of writing  $oand(e_1, e_2)$ , you write  $e_1 \& e_2$ .
- Instead of writing clause(h, b) you can write  $h \Leftarrow b$ , where  $\Leftarrow$  is an infix meta-level predicate symbol.
  - Thus the base-level clause " $h \leftarrow a_1 \wedge \cdots \wedge a_n$ " is represented as the meta-level atom  $h \Leftarrow a_1 \& \cdots \& a_n$ .

### Example representation

#### The base-level clauses

connected\_to(
$$l_1, w_0$$
).  
connected\_to( $w_0, w_1$ )  $\leftarrow$  up( $s_2$ ).  
 $lit(L) \leftarrow light(L) \wedge ok(L) \wedge live(L)$ .

can be represented as the meta-level facts

```
connected\_to(l_1, w_0) \Leftarrow true.
connected\_to(w_0, w_1) \Leftarrow up(s_2).
lit(L) \Leftarrow light(L) \& ok(L) \& live(L).
```

