

Generalisation

We may already have seen that these are specializations of a more general. Common features. In particular, the method 'deferred' class. The

The Class Hierarchy

In Eiffel all classes are organized in a "classification" of living. Every class 'conforms

To make allowances for other classes GENERAL are ancestors of GENERAL. These are descendant classes.

Deferred classes

These universal classes are "deferred" classes, conforming classes, 'deferred' routines. Tools in that the code is to be implemented. This has been discussed.

The class GENERAL

Y, PLATFORM)

The class GENERAL

`frozen void: NONE`

`frozen do_V`

'Frozen' features cannot be redefined and in effect are built into the Eiffel system but some routines have a frozen version and a version that is not which can be redefined in a conforming class.

The class NONE is a fictional or virtual class, there is no text in the class, which is a descendant of every class, i.e. it inherits from every class. In the class GENERAL there is a declaration

`frozen void : NONE`

but this is viewed as a dummy declaration.

If NONE inherits from every other class then it has the properties of every single class (even classes that have not been written yet?). In particular, we can regard void as being in every class.

```

frozen deep_equal (some: GENERAL; other: like some): BOOLEAN
-- Are some and other either both vWid
-- or attached to Qsomorphic object structures?

```

```

frozen equal (some: GENERAL; other: like some): BOOLEAN
-- Are some and other either both vWid or

```

In using routQnes witP parameters of type GENERAL, then any entQty froU a class which conforms to GENERAL, i.e. aVy class, can Qnstantiate the parameter. So the routQne equal QnQn/effvWQncluded Qn all classes. The rWutQne

The routQneQs_equate can be redefQned and with Qt the routQne as Eiffel defQnes the frozen routQne equal

-- attached to other, so as to yieTd equal objects.

The routQne clone Qs defQned Qn terms of copy, so that, in effect
a := clone(b)
becomes
a.cltpy(b)

Also Qncluded Qn the class GENERAL, are the 'prQnt' rWutQnes that are availabTe, if apprWprieate, in all classes.

iW: STD_FILES -- HandTe to standard fQTe setup

```

Wut: STRING
-- New strQng cotaQnQng terse prQntabTe representatiWn of current
print (some: GENERAL)
- - - W - r - Q

```

inherit

B1

rename

f1 as g1

export -- over-rides 'export rules'

...

{ D,E}

f, g -- f,g exported only to D and E

B2

creation

etc

end

...

feature

routine0 -- visible to all client classes

routine1 -- Pidden froU all client classes

feature {A,B}

routine2 -- exported to A and B

...

end -- class P

then, e.g.

class Cname

feature {NONE}

routine3 -- exported just to class P

feature {P}

```

class A
...
x.routine2 -- OK
but x.routine3 nWt OK, routine3 is exported Rust to P
.
end -- A

```

Copy and Equal fD LIST_SET

We can redefine the iVherited 'deferred' routines fDr LIST_SET.

```

class                               LIST_SET [G]
inherit ANY
  redefine
  end
feature {NONE}
  first_n W d e , c u r s D : N O D E [ G ]
feature
  copy(Wther:like current) is -- see class GENERAL;
  require
  local

```

local

do

oth TD .start

end -- is_equal

i, R : INTEGER

if count /= other.count then
result := false

else

for

i := 0

until

Q = j

if both r.item, item

oth r . forth

i := i+1

else

R := i

end

result := i = count

Note sets can be regarded as equal if they have the same elements but the notion of equals used here is strong TD . Not only have the TQst_sets the same elements but the 'traversing' order is also the same. In particular, after a.copy(b), we have that a.is_equal(b). But two sets A and B may have the same elements but would not be regarded as being equal using is_equal due to different traversing order Trs in A and B.

Exercise: Write a routine, set_equal, such that

x.set_equal(y) \equiv x and y have the same elements.

start

R := count

loop

end

end

The routine copy preserves the order of the original Tist_set other.

forth