

# Virtual Memory

**Introduction**

Demand  
Replacement  
Allocation

## ■ Reading: OS Concepts: Chapter 10

## ■ Contents

- Introduction
- Demand Paging
- Page Replacement
- Frame Allocation

# What is virtual memory?

## ■ Complete separation of logical and physical memory

- ☞ \_\_\_\_\_
- ☞ \_\_\_\_\_
- ☞ \_\_\_\_\_
- ☞ \_\_\_\_\_

## ■ Virtual memory can be implemented via:

- ☞ Demand paging
- ☞ Demand segmentation

## Demand paging

### ■ Technique



### ■ Advantages & disadvantages



### ■ There are 3 possibilities when a page is referenced

☞ If it is in-memory: \_\_\_\_\_

☞ If it is not in-memory: \_\_\_\_\_

☞ If it is invalid: \_\_\_\_\_

3

## Handling memory references

### ■ So first we must test if the page is legal

### ■ Then check if it is present

☞ Using the valid-invalid bit which



Frame #	valid-invalid bit
	1
	1
	1
	1
	0
⋮	
	0
	0

page table

4

## Handling page faults

- A memory reference causes a page table lookup
- If the page is invalid a **page fault** (interrupt) occurs

\_\_\_\_\_

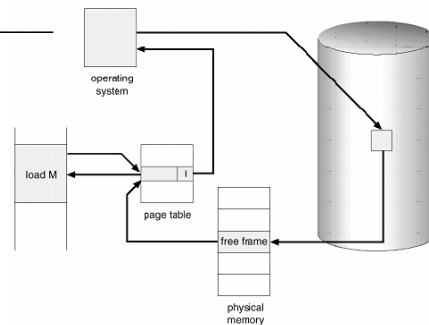
\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_



5

## Performance

- $EAT = (1 - p) * EAT_{\text{paging}} + p * [\text{Page Fault Service Time}]$
- $p$  is the **Page Fault Rate**
  - $p = 0.0$  implies \_\_\_\_\_
  - $p = 1.0$  implies \_\_\_\_\_

■  $EAT_{\text{paging}}$ : \_\_\_\_\_

■ Page Fault Service Time: \_\_\_\_\_

### Example

- $m = 100\text{ns}$
- $\epsilon = 10\text{ns}$
- $\alpha = 0.98$
- $n = 2$
- $p = 0.000001$
- Page Fault Service Time =  $25\text{ms}$

$EAT_{\text{paging}} =$  \_\_\_\_\_

$EAT =$  \_\_\_\_\_

6

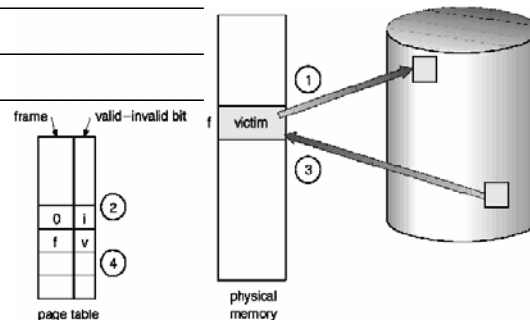
## Page Replacement

- Page replacement is necessary when there is no free memory into which to load a referenced page.

- Basic steps

- Somehow pick a victim page

- 1. \_\_\_\_\_
    - 2. \_\_\_\_\_
    - 3. \_\_\_\_\_
    - 4. \_\_\_\_\_



7

## Page Replacement Issues

- The Page Fault Service Time (PFST) can be reduced if the swapped out page has not been modified

- \_\_\_\_\_

- We also need to consider ownership of the victim

- \_\_\_\_\_

- \_\_\_\_\_

- We now consider possible algorithms for the replacement...

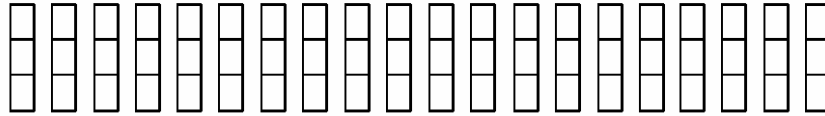
- FIFO
  - Optimal
  - LRU

8

## FIFO Page Replacement

reference string

7 0 1 2 0 3 0 4 2 3 0 3 2 1 2 0 1 7 0 1



page frames

■ Simple concept: \_\_\_\_\_

■ \_\_\_\_\_  
 \_\_\_\_\_

■ Belady's Anomaly: \_\_\_\_\_

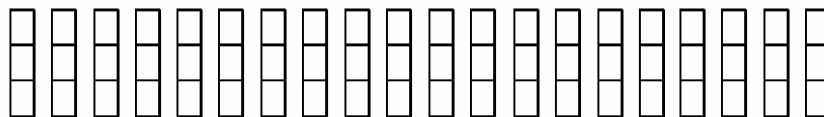
\_\_\_\_\_

9

## Optimal Page Replacement

reference string

7 0 1 2 0 3 0 4 2 3 0 3 2 1 2 0 1 7 0 1



page frames

■ \_\_\_\_\_  
 \_\_\_\_\_

■ \_\_\_\_\_  
 \_\_\_\_\_

10

## LRU Implementations

### ■ Counter implementation

- \_\_\_\_\_
- \_\_\_\_\_
- \_\_\_\_\_

### ■ Stack implementation

- \_\_\_\_\_
- \_\_\_\_\_
- \_\_\_\_\_

## LRU Approximations

### ■ Additional Reference bits algorithm

- \_\_\_\_\_
- \_\_\_\_\_
- \_\_\_\_\_
- \_\_\_\_\_

### ■ Second chance algorithm

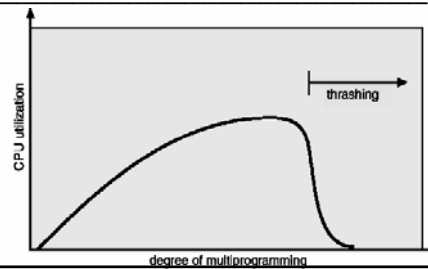
- \_\_\_\_\_
- \_\_\_\_\_
- \_\_\_\_\_

## Thrashing

- Thrashing is when a process is busy (i.e. spends most of its time) swapping pages in and out

■

■

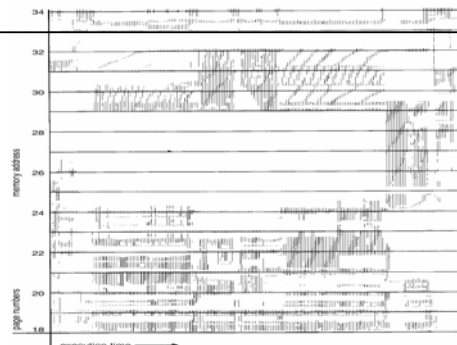


## Locality model

- Paging works because the order of memory accesses are not random.

■

- So why does thrashing occur then?

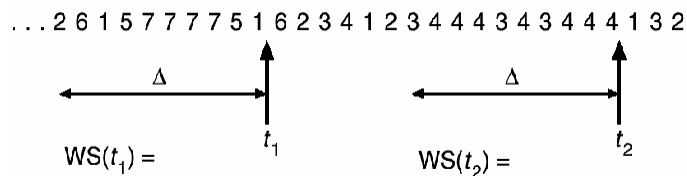


Solving Thrashing 1: **Working Set Model**

- The working set of a process is the set of page references used over some time / no. of instructions  $\Delta$

- If  $\Delta$  is too small: \_\_\_\_\_

- If  $\Delta$  is too large: \_\_\_\_\_



- \_\_\_\_\_
- \_\_\_\_\_

15

**Implementing Working Set**

- It is too complex to implement the working set model as we would need to keep incredibly long lists of previous page references for each process.

- Instead we can approximate it with

- \_\_\_\_\_
  - \_\_\_\_\_

- Example: \_\_\_\_\_

- How accurate is the approximation?

- \_\_\_\_\_
  - \_\_\_\_\_

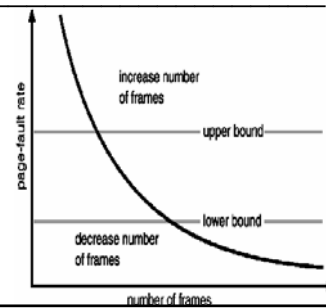
16



## Solving Thrashing 2: Page Fault Frequency Scheme

- This is a simpler solution...
- Establish an 'acceptable' page fault rate

- \_\_\_\_\_
- \_\_\_\_\_



## Programming Considerations

Given Array  $A[1024, 1024]$  of integer, which should you write?

Program 1      **for** (*int*  $i=0$ ; ( $i<1024$ );  $i++$ )  
                   **for** (*int*  $j=0$ ; ( $j<1024$ );  $j++$ )  
                    $A[i,j] = 0$ ;

Program 2      **for** (*int*  $j=0$ ; ( $j<1024$ );  $j++$ )  
                   **for** (*int*  $i=0$ ; ( $i<1024$ );  $i++$ )  
                    $A[i,j] = 0$ ;

## Windows NT

- Uses Demand paging with clustering
  - \_\_\_\_\_
- Processes are assigned Working set minimum & maximum
  - Minimum: \_\_\_\_\_
  - Maximum: \_\_\_\_\_
- When free memory falls below a threshold
  - \_\_\_\_\_

19

## Solaris 2

- Maintains a list of free pages
- Lotsfree \_\_\_\_\_
- Pageout process: \_\_\_\_\_
- Scanrate: \_\_\_\_\_

20