

UNIVERSITY OF DUBLIN

TRINITY COLLEGE

Faculty of Engineering and Systems Sciences
Department of Computer Science

B.A.Computer Science
Senior Freshman Examination

Trinity Term 2002

2BA2 - Programming Techniques

Saturday 25th May

Exam Hall

14.00 - 17.00

Dr. Hugh Gibbons

Attempt FOUR questions

(In presenting programs explain clearly the design of the Eiffel code)

Qs 1.

The following routine, `ins_sort`, sorts an array using the algorithm 'sorting by insertion'.

```

ins_sort (a:ARRAY[G]; low,high:INTEGER) is
  require
    a /= void and then low <= high
    a.lower <= low and high <= a.upper;
  local
    k: INTEGER;
  do
    from
      k := low + 1
    until
      k > high
    loop
      search(a, low, k-1, a.item(k));
      insert(a.item(k), index+1, a, low, k-1);
      k := k + 1
    end
  ensure
    sorted: is_ordered(a, low, high)
end ; -- ins_sort

```

The routine call,

```
search(a, low, k-1, a.item(k))
```

searches where to insert the item, `a.item(k)`, into the ordered array segment `a[low .. k-1]` using 'a binary search technique'.

The routine call,

```
insert(a.item(k), index+1, a, low, k-1)
```

inserts the item, `a.item(k)` at position, `index+1`, in the array segment, `a[low .. k-1]`.

i) Present an Eiffel routine

```

Search(a:ARRAY[G]; low,high:INTEGER; x:G) is
  require
    Ordered: Is_Ordered(a,low,high)
  ensure
    --(found → x = a.item(index)) & (¬found → x ∉ a[low..high])

```

that 'binary searches' an array segment `a[low .. high]` for an item `x`.

If there is more than one occurrence of `x`, it returns the position of the 'rightmost' one, i.e. if `a.item(index) = x`, then `a.item(index+1) > x`.

ii) Present an Eiffel routine

```

insert(x:G; i:INTEGER; a:ARRAY[G]; low,high:INTEGER)
  require
    a /= void and low <= high
    a.lower <= low and high < a.upper
    low <= i and i <= high
  ensure
    equal(a.item(i), x)

```

that inserts `x` at position `i` in array, `a`, by moving items to the right.

iii) Given an array attribute, $a:\text{ARRAY}[G]$, present an Eiffel routine

$\text{reverse}(\text{low}, \text{high} : \text{INTEGER})$

that reverses the array segment $a[\text{low} .. \text{high}]$.

If the array, a , was in ascending order then the call

$\text{reverse}(a.\text{lower}, a.\text{upper})$

changes the array, a , into descending order.

Qs. 2

a) Implement boolean function

$\text{is_equal}(a, b : \text{ARRAY}[G]) : \text{BOOLEAN}$

which determines whether or not the arrays, a and b , are equal item by item.

b) Present an Eiffel procedure,

$\text{partition}(L0, R0 : \text{INTEGER}; p : G)$

that will partition an array attribute, $A:\text{ARRAY}[G]$, about a pivot, p , such that after partition we will have

$A[L0 .. R] \leq p \leq A[L .. R0]$

where L and R are integer attributes.

c) Using a procedure for partitioning an array, present a procedure that will Quicksort an array, $A:\text{ARRAY}[G]$.

Qs 3

Present Eiffel routines that will

a) Generate all $n!$ permutations of $\{1..n\}$ in order starting with the 'largest' $n, n-1, \dots, 2, 1$ and ending with the 'smallest' $1, 2, 3, \dots, n$

b) Generate all 2^n subsets of $\{1..n\}$ starting with the set $\{1, 2, \dots, n\}$ and finishing with the empty set, $\{\}$.

Q 4

Assume we are given the classes, LIST_BAG and NODE with the following short forms (repeated items are allowed),

```

class interface LIST_BAG [G]

    add_first(x : G)
    -- add x, in new first node

    add_last(x:G)
    --add x in new last node

    count : INTEGER

    empty : BOOLEAN

    has (x : G) : BOOLEAN

    remove_first_node
    -- remove first node, if any.

    remove_last_node
    -- remove last node, if any

    reverse
    -- reverse the list

    <other routines>

end -- class LIST_BAG

```

```

class interface NODE[G]
    item : G
    next : NODE[G]

    set_item(x : G)
    set_next(n : NODE[G])

end -- NODE

```

Implement, using linked nodes, the routines,

add_first(x : G), add_last(x:G), remove_first_node, remove_last_node
and the routine
reverse, which reverses the linked list.

Use linked list diagrams in explaining the routines.

Qs. 5.

Assume that a Directed Acyclic Graph (DAG), D, is stored as an adjacency list.

Present Eiffel routines that will

- a) Breadth First Traverse the DAG, D.
- b) Topological Sort the DAG, D.

Qs 6

Assume class interfaces for a Binary Search Tree class, BST, and a BIN_NODE class as follows: (Repeated items in BST.)

```
class interface
  BST [G -> COMPARABLE]
  -- Binary Search Tree
feature

    empty -- make empty

    is_empty: BOOLEAN

    add (x: G)
    -- add x to BST;
    -- (repeated items allowed)

    array2bst(a:ARRAY[G])
    -- add all items from array, a, to
    -- BST.

    bst2array : ARRAY[G]
    --puts items in BST into an array
    -- so that the array is ordered.

    root: BIN_NODE [G]

    count: INTEGER -- # items in tree
end -- class BST
```

```
class interface
  BIN_NODE [G]
feature

    value: G;

    left: BIN_NODE [G];

    right: BIN_NODE [G];

    value_set (v: G)

    left_set (n: BIN_NODE [G])

    right_set (n: BIN_NODE [G])

    build (v: G; l, r: BIN_NODE [G])

end -- class BIN_NODE
```

Implement the routines
 add, array2bst, bst2array
 in the class BST

Suggestion:

The routine, add(x : G), may use an auxiliary routine, insert, as follows:

```
    Add(x:G) is
      do
        if root /= void then
          insert(x,root)
        else
          !!root
          root.build(x,void,void)
          count := 1
        end
      end -- Add
```

Implement the routine,
 insert(x : G; t : BIN_NODE[G])
 that inserts an item, x, into a tree with root, t.

The routine, bst2array, may use an auxillary routine

```
inord(t:BIN_NODE[G])
```

as follows:

```
bst2array : ARRAY[G] is
  require
    not is_empty
  do
    !!arr_ord.make(1,size)
    index:= 1
    inord(root)
    result := clone(arr_ord)
  end -- bst2array
```

The array, arr_ord, and integer, index, are hidden attributes of the class. The counter, index and array, arr_ord, may be updated during the execution of inord. The routine, inord, inorder traverses the tree, starting at the root, and adds the items of the tree to the array, arr_ord. Implement the routine, inord.

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