STAT253/317 Lecture 6 Time Reversibility (§4.8)

4.8.1 Backward Markov Chain

If $\{\ldots,X_{n-1},X_n,X_{n+1},\ldots\}$ is a Markov chain, the backward chain $\{\ldots, X_{n+1}, X_n, X_{n-1}, \ldots\}$ is also a Markov chain.

Proof.

$$P(X_m = j \mid X_{m+1} = i, X_{m+2}, X_{m+3}, ...)$$

$$= \frac{P(X_m = j, X_{m+1} = i, X_{m+2}, X_{m+3}, \dots)}{P(X_{m+1} = i, X_{m+2}, X_{m+3}, \dots)}$$

$$P(X_{m+1} = i, X_{m+2}, X_{m+3}, ...)$$

$$P(X_{m+2}, X_{m+3}, ... | X_m = j, X_{m+1} = i)) P(X_m = j, X_{m+1} = i)$$

$$P(X_{m+2}, X_{m+3}, ... | X_{m+1} = i) P(X_{m+1} = i)$$

$$X_{m+1} = i$$
) $(X_m = j, X_m = j)$

$$= \frac{P(X_{m+2}, X_{m+3}, \dots \mid X_{m+1} = i)}{P(X_{m+2}, X_{m+3}, \dots \mid X_{m+1} = i)} P(X_m = j, X_{m+1} = i)$$
 (Markov Property

$$|X_{m+1}=i)$$

$$= \frac{P(X_m = j, X_{m+1} = i)}{P(X_{m+1} = i)} = P(X_m = j \mid X_{m+1} = i)$$

Transition Probabilities of the Backward Markov Chain

Consider a Markov chain $\{X_n : n = 0, 1, 2, ...\}$ with transition probabilities $\{P_{ij}\}$.

Let $\{\pi_i^{(m)} = P(X_m = j)\}$ be the marginal distribution of X_m .

The transition probabilities $\{Q_{ij}^{(m)}\}$ of the backward Markov chain are

$$Q_{ij}^{(m)} = P(X_m = j \mid X_{m+1} = i)$$

$$= \frac{P(X_m = j, X_{m+1} = i)}{P(X_{m+1} = i)}$$

$$= \frac{P(X_m = j)P(X_{m+1} = i \mid X_m = j)}{P(X_{m+1} = i)} = \frac{\pi_j^{(m)}P_{ji}}{\pi_i^{(m+1)}}$$

We can see the backward Markov chain is **NOT** stationary because the transition probabilities $Q_{ii}^{(m)}$ depend on m.

To make the backward Markov chain **stationary**, the forward chain must start with its stationary distribution $\{\pi_i\}$ so that

$$P(X_m = j) = \pi_i$$
 for all m

the transition probabilities $\{Q_{ii}\}$ of the backward Markov chain is

$$Q_{ij} = P(X_m = j \mid X_{m+1} = i)$$

$$= \frac{P(X_m = j, X_{m+1} = i)}{P(X_{m+1} = i)}$$

$$= \frac{P(X_m = j)P(X_{m+1} = i \mid X_m = j)}{P(X_{m+1} = i)} = \frac{\pi_j P_{ji}}{\pi_i}$$

which does not depend on m

Time Reversible Markov Chains & Detailed Balanced Equations

A Markov chain is said to be time reversible iff

$$Q_{ij}=P_{ij},$$

i.e., it behaves exactly the same no matter running forward or backward when in the stationary state.

Because Q_{ij} equals $\pi_j P_{ji}/\pi_i$, a Markov chain is time reversible if and only if its stationary distribution $\{\pi_j\}$ satisfies the equations

$$\pi_i P_{ij} = \pi_j P_{ji}$$
 for all i, j .

This set of equations is called the detailed balanced equation.

Balanced Equations v.s. Detailed Balanced Equations

Recall a distribution π_j for a Markov chain is said to be stationary if and only if it satisfies

$$\pi_j = \sum_{i \in \mathfrak{X}} \pi_i P_{ij}$$
 for all $j \in \mathfrak{X}$.

This set of equations is called the **balanced equations**.

Balanced Equations v.s. Detailed Balanced Equations

Recall a distribution π_j for a Markov chain is said to be stationary if and only if it satisfies

$$\pi_j = \sum_{i \in \mathfrak{X}} \pi_i P_{ij} \quad \text{for all } j \in \mathfrak{X}.$$

This set of equations is called the **balanced equations**.

A solution to the **detailed balanced equations** must also be a solution to the **balanced equations**, because

$$\sum_{i \in \mathfrak{X}} \pi_i P_{ij} = \sum_{i \in \mathfrak{X}} \pi_j P_{ji} = \pi_j \sum_{i \in \mathfrak{X}} P_{ji} = \pi_j \cdot 1 = \pi_j$$

Balanced Equations v.s. Detailed Balanced Equations

Recall a distribution π_j for a Markov chain is said to be stationary if and only if it satisfies

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It is possible that the balanced equations have solutions but the detailed balanced equations do not.

Interpretation of the Balanced Equation

$$\pi_j = \sum_{i \in \mathfrak{X}} \pi_i P_{ij} \quad ext{for all } j \in \mathfrak{X}$$
 $\Leftrightarrow \quad \pi_j (1 - P_{jj}) = \sum_{i \in \mathfrak{X}, i
eq j} \pi_i P_{ij} \quad ext{for all } j \in \mathfrak{X}$

rate of transitions ${f out}$ of state $j={f rate}$ of transitions ${f into}$ state j



How many equations are there in the Balanced euqations? # of balanced equations = # of states

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Interpretation of the Detailed Balanced Equation

Detailed balanced equations are easier to solve but there might not be a solutions

Balanced equations are more difficult to solve but more likely to have a solutions.

$$\pi_i P_{ij} = \pi_j P_{ji}$$

rate of transitions from i to j =rate of transitions from j to i

$$i \longrightarrow j$$
 $j \longrightarrow i$

How many detailed balanced equations are there in total?

Ans: # of equations = number of pairs of states
If there are N states, there are N(N-1)/2 equations

How many unknowns are there? Ans: # of states = N

Example 4.35

Consider a random walk with states 0, 1, ..., M and transition probabilities

$$\begin{aligned} P_{i,i+1} &= \alpha_i = 1 - P_{i,i-1}, & \text{for } i = 1, \dots, M-1, \\ P_{0,1} &= \alpha_0 = 1 - P_{0,0}, \\ P_{M,M} &= \alpha_M = 1 - P_{M,M-1} \\ 0 &\xrightarrow{\alpha_0} 1 \xrightarrow{\alpha_1} 2 \cdots \longrightarrow M-1 \xrightarrow{\alpha_{M-1}} M \overset{\alpha_M}{\circlearrowleft} \\ \xrightarrow{1-\alpha_0} 0 &\xrightarrow{1-\alpha_1} 1 \xrightarrow{\text{elphas.}} M - 1 \xrightarrow{1-\alpha_M} M \end{aligned}$$

Example 4.35 (Cont'd)

The stationary distribution π can be solved via the detailed balanced equation

$$\pi_i P_{i,i-1} = \pi_i (1 - \alpha_i) = \pi_{i-1} P_{i-1,i} = \pi_{i-1} \alpha_{i-1}$$

So

$$\pi_i = \frac{\alpha_{i-1}}{1 - \alpha_i} \pi_{i-1} = \dots = \frac{\alpha_{i-1} \alpha_{i-2} \dots \alpha_0}{(1 - \alpha_i)(1 - \alpha_{i-1}) \dots (1 - \alpha_1)} \pi_0$$

Since $\sum_{i=0}^{M} \pi_{i} = 1$, one can solve π_{0} via

$$\pi_0 \left[1 + \sum_{i=1}^M \frac{\alpha_{i-1}\alpha_{i-2}\dots\alpha_0}{(1-\alpha_i)(1-\alpha_{i-1})\dots(1-\alpha_1)} \right] = 1$$

A Non-Time-Reversible Markov Chain

In Exercise 4.34 (a flea moving around the vertices of a triangle),

the transition probabilities, and the stationary distribution are respectively

where $C = 3 - p_2q_3 - p_3q_1 - p_1q_2$. One can easily verify that

$$\pi_1 P_{12} = \pi_1 p_1 \neq \pi_2 P_{21} = \pi_2 q_2$$

The chain is NOT time reversible.

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Other Non-Time-Reversible Markov Chains

► A Markov chain with **transient states** cannot be time-reversible because then running forward and backward in time will not be equivalent.

▶ If there exists two states *i* and *j* such that

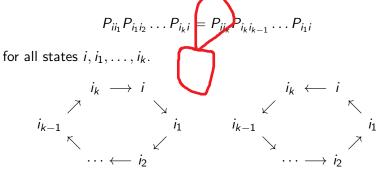
$$P_{ij} > 0$$
 but $P_{jj} = 0$

then the Markov chain cannot be time-reversible because then when running backward in time

$$Q_{ij} = \frac{\pi_j P_{ji}}{\pi_i} = 0 \neq P_{ij}.$$

Theorem 4.2

An ergodic Markov chain for which $P_{ij}=0$ whenever $P_{ji}=0$ is time reversible if and only if starting in state i, any path back to i has the same probability as the reversed path. That is, if



For some ergodic Markov chains that have no cycles, like 1-D random walk, they must be time-reversible

Theorem 4.2 — Proof of Necessity

If a Markov chain is time reversible, we have

$$\pi_i P_{ij} = \pi_j P_{ji}, \quad \pi_k P_{kj} = \pi_j P_{jk}.$$

implying (if $P_{ij}P_{jk} > 0$) that

$$\frac{\pi_i}{\pi_k} = \frac{P_{ji}P_{kj}}{P_{ij}P_{jk}},$$

but $\pi_i P_{ik} = \pi_k P_{ki}$ also implies $\pi_i / \pi_k = P_{ki} / P_{ik}$. Thus

$$P_{ik}P_{kj}P_{ji}=P_{ij}P_{jk}P_{ki}.$$

This proves for the case $i \rightarrow j \rightarrow k \rightarrow i$. The general case for longer cycle can be proved similarly

Theorem 4.2 — Proof of Sufficiency

Consider the cycle $i \rightarrow i_1 \rightarrow i_2 \rightarrow \ldots \rightarrow i_k \rightarrow j \rightarrow i$.

$$P_{ii_1}P_{i_1i_2}\dots P_{i_kj}P_{ji} = P_{ij}P_{ji_k}P_{i_ki_{k-1}}\dots P_{i_1i_j}$$

Summing the preceding over all states i_1, \ldots, i_k yields

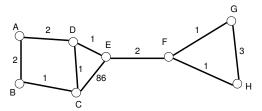
$$P_{ij}^{(k-1)}P_{ji}=P_{ij}P_{ji}^{(k-1)}$$

Letting $k \to \infty$ yields

$$\pi_i P_{ii} = P_{ij} \pi_i$$

which proves the theorem.

Example 4.36 Random Walk on a Weighted Graph (p.241)



A graph = a set of vertices (or nodes) + a set of arcs (or edges) connecting some pairs of vertices. We consider random walk on a connected graph such that

- \triangleright each pair (i, j) of vertices are connected by at most one arc;
- ▶ all arcs are undirected: arc (i,j) = arc(j,i);
- ▶ there is a path consists of arcs connecting any pair of vertices;
- each arc (i, j) is associated with a weight $w_{ii} > 0$
 - \triangleright $w_{ij} = 0$ if there is not arc connecting (i, j)
 - $ightharpoonup w_{ij} = w_{ji}$

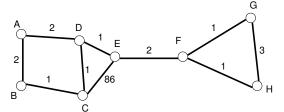
Example 4.36 Random Walk on a Weighted Graph (p.241)

A particle moving from vertices to vertices that if at any time the particle is at node i, then it will next move to node j with probability

$$P_{ij} = \frac{w_{ij}}{\sum_{k} w_{ik}},$$

E.g., in the graph below, there are two arcs from vertices B with weights $w_{BA}=2$ and $W_{BC}=1$ respectively. So,

$$P_{BA} = \frac{w_{BA}}{w_{BA} + w_{BC}} = \frac{2}{2+1} = \frac{2}{3}, \quad P_{BC} = \frac{w_{BC}}{w_{BA} + w_{BC}} = \frac{1}{2+1} = \frac{1}{3}.$$



Random walk on a graph is **irreducible** because the graph is **connected**.

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Random Walk on a Weighted Graph is Time Reversible

Solving the detailed balanced equation:

$$\pi_i P_{ij} = \frac{\pi_i w_{ij}}{\sum_k w_{ik}} = \frac{\pi_j w_{ji}}{\sum_k w_{jk}} = \pi_j P_{ji} \quad \text{for all } i, j$$

or, equivalently, since $w_{ij} = w_{ji}$,

$$\frac{\pi_i}{\sum_k w_{ik}} = \frac{\pi_j}{\sum_k w_{jk}} \quad \text{for all } i, j,$$

which means $\frac{\pi_i}{\sum_{i} w_{ik}}$ is a constant c for all i, i.e.,

$$\pi_i = c \sum_{i} w_{ik}.$$

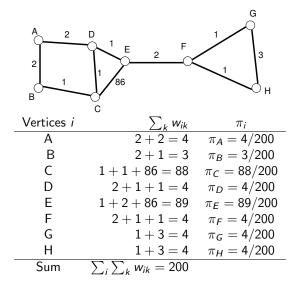
Since $1=\sum_j \pi_j=c\sum_j \sum_k w_{jk},$ we know $c=1/(\sum_j \sum_k w_{jk}),$ and hence

$$\pi_i = rac{\sum_k w_{ik}}{\sum_i \sum_k w_{jk}}$$

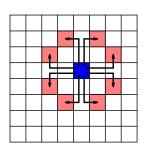
is a solution to the detailed balanced eq. The process is therefore time-reversible.

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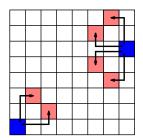
Random Walk on a Weighted Graph

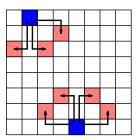


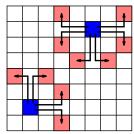
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- ► The Knight moves in an L shape in any direction.
- At the blue square, the Knight can move to any of the 8 red squares.
- From a square near the boundary, the Knight has fewer possible moves as it cannot move out of the Chessboard (see the 3 graphs below.)

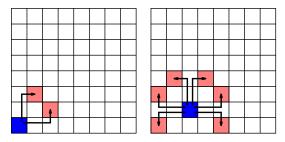






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- ► A Knight moves randomly on an empty chessboard.
- ▶ In each step, it's equally like to take any of its legal moves. E.g., at the corner, it has prob. 1/2 each to move to either of the two red squares, from which it has prob. 1/6 each to move to any of the 6 possible squares.



- ▶ Each move is indep. of the history of moves up to that time.
- ► The position of on knight after *n*th move is a Markov chain where states are the 64 squares on the chessboard.

The Knight's random walk on a Chessboard is also a random walk on weighted graph where

- the vertices are the 64 squares on the chessboard;
- there is an arc between any two squares that Knight can move in 1 step;
- ▶ all the arcs have weight $w_{ij} = 1$.

The transition probability of a random walk on weighted graph from square i to square j is

$$P_{ij} = \frac{w_{ij}}{\sum_k w_{ik}} = \frac{1}{\# \text{ of squares that connected with square } i \text{ with an arc}}$$
$$= \frac{1}{\# \text{ of legal moves from square } i}$$

which is exactly the random walk of the knight.

Using the property of random walks on a graph, the stationary distribution of the Knight's random walk is

$$\pi_i = \frac{\sum_k w_{ik}}{\sum_j \sum_k w_{jk}} = \frac{\# \text{ of legal moves from square } i}{\sum_j (\# \text{ of legal moves from square } j)}$$

The numbers of legal moves from the squares are as follows:

2	3	4	4	4	4	3	2
3	4	6	6	6	6	4	3
4	6	8	8	8	8	6	4
4	6	8	8	8	8	6	4
4	6	8	8	8	8	6	4
4	6	8	8	8	8	6	4
3	4	6	6	6	6	4	3
2	3	4	4	4	4	3	2

The sum of the number of possible moves over all squares is

$$2 \times 4 + 3 \times 8 + 4 \times 20$$

 $+ 6 \times 16 + 8 \times 16 = 336.$

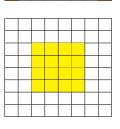
The long run proportion of time that the Knight is in a specific square is simply the counts in the table above divided by 336.

Return Time of a Random Knight

Recall that $1/\pi_i = \mathbb{E}[T_i]$ is the expected time between two visits of the Markov chain to state i.

Starting from one of the four corners, it takes $1/\pi_i=336/2=168$ moves on average for a Knight to return to its initial position.

Starting from the center of the chessboard, it takes $1/\pi_i = 336/8 = 42$ moves on average for a Knight to return to its initial position

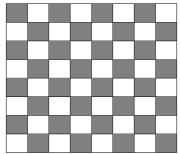


More Questions

- ▶ Is this Markov chain irreducible? That is, can the Knight visit every square from every square?
- What is the period of this Markov chain?

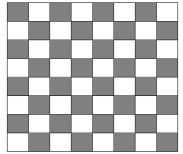
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- ▶ Is this Markov chain irreducible? That is, can the Knight visit every square from every square?
- What is the period of this Markov chain?



Every "L" move can only from a gray square to a white square or a white square to a gray square.