STAT 37797: Mathematics of Data Science

Matrix concentration inequalities



Cong Ma
University of Chicago, Autumn 2021

Recap: matrix Bernstein inequality

Consider a sequence of independent random matrices $\{oldsymbol{X}_l \in \mathbb{R}^{d_1 imes d_2}\}$

•
$$\mathbb{E}[X_l] = \mathbf{0}$$

•
$$\|\boldsymbol{X}_l\| \leq B$$
 for each l

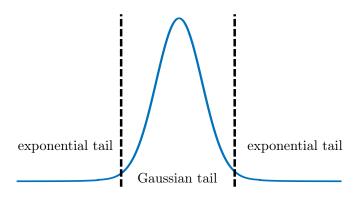
variance statistic:

$$v := \max \left\{ \left\| \mathbb{E} \left[\sum_{l} \boldsymbol{X}_{l} \boldsymbol{X}_{l}^{\top} \right] \right\|, \left\| \mathbb{E} \left[\sum_{l} \boldsymbol{X}_{l}^{\top} \boldsymbol{X}_{l} \right] \right\| \right\}$$

Theorem 3.1 (Matrix Bernstein inequality)

For all
$$\tau \geq 0$$
,
$$\mathbb{P}\left\{\left\|\sum_{l} \boldsymbol{X}_{l}\right\| \geq \tau\right\} \leq (d_{1} + d_{2}) \exp\left(\frac{-\tau^{2}/2}{v + B\tau/3}\right)$$

Recap: matrix Bernstein inequality





This lecture: detailed introduction of matrix Bernstein

An introduction to matrix concentration inequalities

— Joel Tropp '15

Outline

- Matrix theory background
- Matrix Laplace transform method
- Matrix Bernstein inequality
- Application: random features

Matrix theory background

Matrix function

Suppose the eigendecomposition of a symmetric matrix $oldsymbol{A} \in \mathbb{R}^{d imes d}$ is

$$oldsymbol{A} = oldsymbol{U} \left[egin{array}{ccc} \lambda_1 & & & & \ & \ddots & & \ & & \lambda_d \end{array}
ight] oldsymbol{U}^ op$$

Then we can define

$$f(oldsymbol{A}) := oldsymbol{U} \left[egin{array}{ccc} f(\lambda_1) & & & & & & \\ & & \ddots & & & & \\ & & & f(\lambda_d) \end{array}
ight] oldsymbol{U}^ op$$

Examples of matrix functions

• Let $f(a) = c_0 + \sum_{k=1}^{\infty} c_k a^k$, then

$$f(\mathbf{A}) := c_0 \mathbf{I} + \sum_{k=1}^{\infty} c_k \mathbf{A}^k$$

- matrix exponential: $e^{A}:=I+\sum_{k=1}^{\infty}\frac{1}{k!}A^{k}$ (why?)
 monotonicity: if $A \prec H$, then $\operatorname{tr} e^{A} < \operatorname{tr} e^{H}$
- matrix logarithm: $\log(e^A) := A$
 - \circ monotonicity: if $0 \preceq A \preceq H$, then $\log A \preceq \log(H)$

Matrix moments and cumulants

Let X be a random symmetric matrix. Then

• matrix moment generating function (MGF):

$$M_X(\theta) := \mathbb{E}[e^{\theta X}]$$

• matrix cumulant generating function (CGF):

$$\Xi_{\boldsymbol{X}}(\theta) := \log \mathbb{E}[e^{\theta \boldsymbol{X}}]$$

Matrix Laplace transform method

Matrix Laplace transform

A key step for a scalar random variable Y: by Markov's inequality,

$$\mathbb{P}\left\{Y \ge t\right\} \le \inf_{\theta > 0} e^{-\theta t} \,\mathbb{E}\left[e^{\theta Y}\right]$$

This can be generalized to the matrix case

Matrix Laplace transform

Lemma 3.2

Let Y be a random symmetric matrix. For all $t \in \mathbb{R}$,

$$\mathbb{P}\left\{\lambda_{\max}(\boldsymbol{Y}) \geq t\right\} \leq \inf_{\theta > 0} e^{-\theta t} \mathbb{E}\left[\operatorname{tr} e^{\theta \boldsymbol{Y}}\right]$$

ullet can control the extreme eigenvalues of $oldsymbol{Y}$ via the trace of the matrix MGF

Proof of Lemma 3.2

For any $\theta > 0$,

$$\begin{split} \mathbb{P}\left\{\lambda_{\max}(\boldsymbol{Y}) \geq t\right\} &= \mathbb{P}\left\{\mathrm{e}^{\theta\lambda_{\max}(\boldsymbol{Y})} \geq \mathrm{e}^{\theta t}\right\} \\ &\leq \frac{\mathbb{E}[\mathrm{e}^{\theta\lambda_{\max}(\boldsymbol{Y})}]}{\mathrm{e}^{\theta t}} \qquad \text{(Markov's inequality)} \\ &= \frac{\mathbb{E}[\mathrm{e}^{\lambda_{\max}(\theta\boldsymbol{Y})}]}{\mathrm{e}^{\theta t}} \\ &= \frac{\mathbb{E}[\lambda_{\max}(\mathrm{e}^{\theta\boldsymbol{Y}})]}{\mathrm{e}^{\theta t}} \qquad (\mathrm{e}^{\lambda_{\max}(\boldsymbol{Z})} = \lambda_{\max}(\mathrm{e}^{\boldsymbol{Z}})) \\ &\leq \frac{\mathbb{E}[\operatorname{tr}\mathrm{e}^{\theta\boldsymbol{Y}}]}{\mathrm{e}^{\theta t}} \end{split}$$

This completes the proof since it holds for any $\theta > 0$

Issues of the matrix MGF

The Laplace transform method is effective for controlling an independent sum when MGF decomposes

• in the scalar case where $X = X_1 + \cdots + X_n$ with independent $\{X_l\}$:

$$M_X(\theta) = \mathbb{E}[\mathrm{e}^{\theta X_1 + \dots + \theta X_n}] = \mathbb{E}[\mathrm{e}^{\theta X_1}] \cdots \mathbb{E}[\mathrm{e}^{\theta X_n}] = \underbrace{\prod_{l=1}^n M_{X_l}(\theta)}_{\text{look at each } X_l \text{ separately}}$$

Issues in the matrix settings:

$$\mathrm{e}^{m{X}_1+m{X}_2}
eq \mathrm{e}^{m{X}_1}\mathrm{e}^{m{X}_2}$$
 unless $m{X}_1$ and $m{X}_2$ commute
$$\mathrm{tr}\,\mathrm{e}^{m{X}_1+\cdots+m{X}_n}
eq \mathrm{tr}\,\mathrm{e}^{m{X}_1}\mathrm{e}^{m{X}_1}\cdots\mathrm{e}^{m{X}_n}$$

Subadditivity of the matrix CGF

Fortunately, the matrix CGF satisfies certain subadditivity rules, allowing us to decompose independent matrix components

Lemma 3.3

Consider a finite sequence $\{X_l\}_{1 \leq l \leq n}$ of independent random symmetric matrices. Then for any $\theta \in \mathbb{R}$,

$$\underbrace{\mathbb{E}\Big[\operatorname{tr} e^{\theta \sum_{l} X_{l}}\Big]}_{\operatorname{tr} \exp \left(\Xi_{\sum_{l} X_{l}}(\theta)\right)} \leq \underbrace{\operatorname{tr} \exp\left(\sum_{l} \log \mathbb{E}\left[e^{\theta X_{l}}\right]\right)}_{\operatorname{tr} \exp\left(\sum_{l} \Xi_{X_{l}}(\theta)\right)}$$

• this is a deep result — based on Lieb's Theorem!

Lieb's Theorem



Elliott Lieb

Theorem 3.4 (Lieb '73)

Fix a symmetric matrix $oldsymbol{H}$. Then

$$A \mapsto \operatorname{tr} \exp(H + \log A)$$

is concave on positive-semidefinite cone

Lieb's Theorem immediately implies (exercise: Jensen's inequality)

$$\mathbb{E}[\operatorname{tr}\exp(\boldsymbol{H} + \boldsymbol{X})] \le \operatorname{tr}\exp(\boldsymbol{H} + \log \mathbb{E}[e^{\boldsymbol{X}}])$$
(3.1)

Proof of Lemma 3.3

$$\mathbb{E}\left[\operatorname{tr} e^{\theta \sum_{l} \boldsymbol{X}_{l}}\right] = \mathbb{E}\left[\operatorname{tr} \exp\left(\theta \sum_{l=1}^{n-1} \boldsymbol{X}_{l} + \theta \boldsymbol{X}_{n}\right)\right]$$

$$\leq \mathbb{E}\left[\operatorname{tr} \exp\left(\theta \sum_{l=1}^{n-1} \boldsymbol{X}_{l} + \log \mathbb{E}\left[e^{\theta \boldsymbol{X}_{n}}\right]\right)\right] \quad \text{(by (3.1))}$$

$$\leq \mathbb{E}\left[\operatorname{tr} \exp\left(\theta \sum_{l=1}^{n-2} \boldsymbol{X}_{l} + \log \mathbb{E}\left[e^{\theta \boldsymbol{X}_{n-1}}\right] + \log \mathbb{E}\left[e^{\theta \boldsymbol{X}_{n}}\right]\right)\right]$$

$$\leq \cdots$$

$$\leq \operatorname{tr} \exp\left(\sum_{l=1}^{n} \log \mathbb{E}\left[e^{\theta \boldsymbol{X}_{l}}\right]\right)$$

Master bounds

Combining the Laplace transform method with the subadditivity of CGF yields:

Theorem 3.5 (Master bounds for sum of independent matrices)

Consider a finite sequence $\{X_l\}$ of independent random symmetric matrices. Then

$$\mathbb{P}\left\{\lambda_{\max}\left(\sum_{l} \boldsymbol{X}_{l}\right) \geq t\right\} \leq \inf_{\theta > 0} \frac{\operatorname{tr}\exp\left(\sum_{l} \log \mathbb{E}[e^{\theta \boldsymbol{X}_{l}}]\right)}{e^{\theta t}}$$

• this is a general result underlying the proofs of the matrix Bernstein inequality and beyond (e.g. matrix Chernoff)

Matrix Bernstein inequality

Matrix CGF

$$\mathbb{P}\left\{\lambda_{\max}\left(\sum_{l} \boldsymbol{X}_{l}\right) \geq t\right\} \leq \inf_{\theta > 0} \frac{\operatorname{tr}\exp\left(\sum_{l} \log \mathbb{E}\left[e^{\theta \boldsymbol{X}_{l}}\right]\right)}{e^{\theta t}}$$

To invoke the master bound, one needs to $\underbrace{\mathsf{control}}$ the matrix $\widehat{\mathsf{CGF}}$ main step for proving matrix Bernstein

Symmetric case

Consider a sequence of independent random symmetric matrices $\{m{X}_l \in \mathbb{R}^{d imes d}\}$

• $\mathbb{E}[X_l] = \mathbf{0}$

- $\lambda_{\max}(\boldsymbol{X}_l) \leq B$ for each l
- ullet variance statistic: $v := \left\| \mathbb{E}\left[\sum_{l} oldsymbol{X}_{l}^{2}
 ight] \right\|$

Theorem 3.6 (Matrix Bernstein inequality: symmetric case)

For all
$$\tau \geq 0$$
,
$$\mathbb{P}\left\{\lambda_{\max}\left(\sum_{l} \boldsymbol{X}_{l}\right) \geq \tau\right\} \leq d \exp\left(\frac{-\tau^{2}/2}{v + B\tau/3}\right)$$

Bounding matrix CGF

For bounded random matrices, one can control the matrix CGF as follows:

Lemma 3.7

Suppose
$$\mathbb{E}[X] = \mathbf{0}$$
 and $\lambda_{\max}(X) \leq B$. Then for $0 < \theta < 3/B$,
$$\log \mathbb{E}[\mathrm{e}^{\theta X}] \leq \frac{\theta^2/2}{1 - \theta B/3} \mathbb{E}[X^2]$$

Proof of Theorem 3.6

Let $g(\theta) := \frac{\theta^2/2}{1-\theta B/3}$, then it follows from the master bound that

$$\begin{split} \mathbb{P}\left\{\lambda_{\max}\big(\sum_{i}\boldsymbol{X}_{i}\big) \geq t\right\} &\leq \inf_{\theta>0} \frac{\operatorname{tr}\exp\big(\sum_{i=1}^{n}\log\mathbb{E}[\mathrm{e}^{\theta\boldsymbol{X}_{i}}]\big)}{\mathrm{e}^{\theta t}} \\ &\leq \inf_{0<\theta<3/B} \frac{\operatorname{tr}\exp\big(g(\theta)\sum_{i=1}^{n}\mathbb{E}[\boldsymbol{X}_{i}^{2}]\big)}{\mathrm{e}^{\theta t}} \\ &\leq \inf_{0<\theta<3/B} \frac{d\,\exp\big(g(\theta)v\big)}{\mathrm{e}^{\theta t}} \end{split}$$

Taking $\theta = \frac{t}{v+Bt/3}$ and simplifying the above expression, we establish matrix Bernstein

Proof of Lemma 3.7

Define $f(x) = \frac{e^{\theta x} - 1 - \theta x}{x^2}$, then for any X with $\lambda_{\max}(X) \leq B$:

$$\mathbf{e}^{\theta X} = \mathbf{I} + \theta \mathbf{X} + \left(\mathbf{e}^{\theta X} - \mathbf{I} - \theta \mathbf{X}\right) = \mathbf{I} + \theta \mathbf{X} + \mathbf{X} \cdot f(\mathbf{X}) \cdot \mathbf{X}$$

\(\times \mathbf{I} + \theta \mathbf{X} + f(B) \cdot \mathbf{X}^2

In addition, we note an elementary inequality: for any $0 < \theta < 3/B$,

$$f(B) = \frac{\mathrm{e}^{\theta B} - 1 - \theta B}{B^2} = \frac{1}{B^2} \sum_{k=2}^{\infty} \frac{(\theta B)^k}{k!} \le \frac{\theta^2}{2} \sum_{k=2}^{\infty} \frac{(\theta B)^{k-2}}{3^{k-2}} = \frac{\theta^2/2}{1 - \theta B/3}$$

$$\implies$$
 $e^{\theta X} \leq I + \theta X + \frac{\theta^2/2}{1 - \theta B/3} \cdot X^2$

Since X is zero-mean, one further has

$$\mathbb{E}\left[e^{\theta \boldsymbol{X}}\right] \leq \boldsymbol{I} + \frac{\theta^2/2}{1 - \theta B/3} \mathbb{E}[\boldsymbol{X}^2] \leq \exp\left(\frac{\theta^2/2}{1 - \theta B/3} \mathbb{E}[\boldsymbol{X}^2]\right)$$

Application: random features

Kernel trick

A modern idea in machine learning: replace the inner product by kernel evaluation (i.e. certain similarity measure)

Advantage: work beyond the Euclidean domain via task-specific similarity measures

Similarity measure

Define the similarity measure Φ

- $\Phi(\boldsymbol{x}, \boldsymbol{x}) = 1$
- $|\Phi(\boldsymbol{x}, \boldsymbol{y})| \leq 1$
- $\bullet \ \Phi(\boldsymbol{x}, \boldsymbol{y}) = \Phi(\boldsymbol{y}, \boldsymbol{x})$

Example: angular similarity

$$\Phi(\boldsymbol{x}, \boldsymbol{y}) = \frac{2}{\pi} \arcsin \frac{\langle \boldsymbol{x}, \boldsymbol{y} \rangle}{\|\boldsymbol{x}\|_2 \|\boldsymbol{y}\|_2} = 1 - \frac{2\angle(\boldsymbol{x}, \boldsymbol{y})}{\pi}$$

Kernel matrix

Consider N data points $x_1,\cdots,x_N\in\mathbb{R}^d$. Then the kernel matrix $G\in\mathbb{R}^{N\times N}$ is

$$G_{i,j} = \Phi(\boldsymbol{x}_i, \boldsymbol{x}_j) \qquad 1 \le i, j \le N$$

ullet Kernel Φ is said to be positive semidefinite if $m{G}\succeq m{0}$ for any $\{m{x}_i\}$

Challenge: kernel matrices are usually large

ullet cost of constructing ${m G}$ is $O(dN^2)$

Question: can we approximate G more efficiently?

Random features

Introduce a random variable $oldsymbol{w}$ and a feature map ψ such that

$$\Phi(\boldsymbol{x},\boldsymbol{y}) = \mathbb{E}_{\boldsymbol{w}}[\underbrace{\psi(\boldsymbol{x};\boldsymbol{w})\cdot\psi(\boldsymbol{y};\boldsymbol{w})}_{\text{decouple }\boldsymbol{x} \text{ and } \boldsymbol{y}}]$$

• example (angular similarity)

$$\underline{\Phi(\boldsymbol{x}, \boldsymbol{y}) = 1 - \frac{2\angle(\boldsymbol{x}, \boldsymbol{y})}{\pi} = \mathbb{E}_{\boldsymbol{w}}[\operatorname{sgn}\langle \boldsymbol{x}, \boldsymbol{w} \rangle \cdot \operatorname{sgn}\langle \boldsymbol{y}, \boldsymbol{w} \rangle]}$$
Grothendieck's identity (3.2)

with $oldsymbol{w}$ uniformly drawn from the unit sphere

Random features

Introduce a random variable $oldsymbol{w}$ and a feature map ψ such that

$$\Phi(\boldsymbol{x},\boldsymbol{y}) = \mathbb{E}_{\boldsymbol{w}}[\underbrace{\psi(\boldsymbol{x};\boldsymbol{w})\cdot\psi(\boldsymbol{y};\boldsymbol{w})}_{\text{decouple }\boldsymbol{x} \text{ and } \boldsymbol{y}}]$$

this results in a random feature vector

$$oldsymbol{z} = \left[egin{array}{c} z_1 \ dots \ z_N \end{array}
ight] = \left[egin{array}{c} \psi(oldsymbol{x}_1; oldsymbol{w}) \ dots \ \psi(oldsymbol{x}_N; oldsymbol{w}) \end{array}
ight]$$

 \circ $zz^{ op}$ is an unbiased estimate of G, i.e. $G = \mathbb{E}[zz^{ op}]$

Example

Angular similarity:

$$\begin{split} \Phi(\boldsymbol{x}, \boldsymbol{y}) &= 1 - \frac{2\angle(\boldsymbol{x}, \boldsymbol{y})}{\pi} \\ &= \mathbb{E}_{\boldsymbol{w}} \left[\operatorname{sign} \langle \boldsymbol{x}, \boldsymbol{w} \rangle \operatorname{sign} \langle \boldsymbol{y}, \boldsymbol{w} \rangle \right] \end{split}$$

where $oldsymbol{w}$ is uniformly drawn from the unit sphere

As a result, the random feature map is $\psi({m x},{m w}) = {\sf sign}\langle {m x},{m w}
angle$

Random feature approximation

Generate n independent copies of ${m R} = {m z} {m z}^{ op}$, i.e. $\{{m R}_l\}_{1 \le l \le n}$

Estimator of the kernel matrix G:

$$\hat{\boldsymbol{G}} = \frac{1}{n} \sum_{l=1}^{n} \boldsymbol{R}_{l}$$

Question: how many random features are needed to guarantee accurate estimation?

Statistical guarantees for random feature approximation

Consider the angular similarity example (3.2):

• To begin with,

$$\mathbb{E}[\boldsymbol{R}_{l}^{2}] = \mathbb{E}[\boldsymbol{z}\boldsymbol{z}^{\top}\boldsymbol{z}\boldsymbol{z}^{\top}] = N\mathbb{E}[\boldsymbol{z}\boldsymbol{z}^{\top}] = N\boldsymbol{G}$$

$$\implies v = \left\|\frac{1}{n^{2}}\sum_{l=1}^{n}\mathbb{E}[\boldsymbol{R}_{l}^{2}]\right\| = \frac{N}{n}\|\boldsymbol{G}\|$$

- Next, $\frac{1}{n} \| \mathbf{R} \| = \frac{1}{n} \| \mathbf{z} \|_2^2 = \frac{N}{n} =: B$
- Applying the matrix Bernstein inequality yields: with high prob.

$$\|\hat{\boldsymbol{G}} - \boldsymbol{G}\| \lesssim \sqrt{v \log N} + B \log N \lesssim \sqrt{\frac{N}{n}} \|\boldsymbol{G}\| \log N + \frac{N}{n} \log N$$

$$\lesssim \sqrt{\frac{N}{n}} \|\boldsymbol{G}\| \log N \qquad \qquad \text{(for sufficiently large } n\text{)}$$

Sample complexity

Define the intrinsic dimension of G as

$$\mathsf{intdim}(\boldsymbol{G}) = \frac{\mathrm{tr}\boldsymbol{G}}{\|\boldsymbol{G}\|} = \frac{N}{\|\boldsymbol{G}\|}$$

If $n \gtrsim \varepsilon^{-2}$ intdim(G) $\log N$, then we have

$$\frac{\|\hat{\boldsymbol{G}} - \boldsymbol{G}\|}{\|\boldsymbol{G}\|} \le \varepsilon$$

Reference

- "An introduction to matrix concentration inequalities," J. Tropp, Foundations and Trends in Machine Learning, 2015.
- "Convex trace functions and the Wigner-Yanase-Dyson conjecture," E. Lieb, Advances in Mathematics, 1973.
- "User-friendly tail bounds for sums of random matrices," J. Tropp, Foundations of computational mathematics, 2012.
- "Random features for large-scale kernel machines," A. Rahimi, B. Recht, Neural Information Processing Systems, 2008.