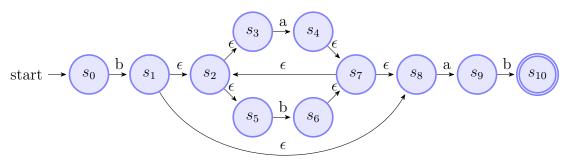
1 给定正则表达式 $b(a|b)^*ab$

1.1 用Thompson算法构造出相应的NFA(15分)

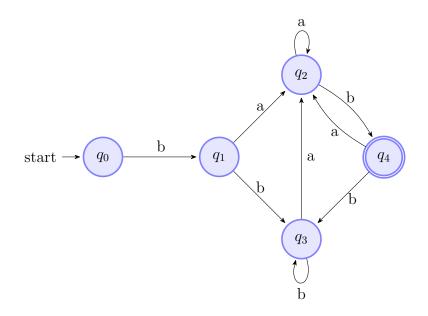


1.2 用子集构造法将NFA转换为DFA,写出状态转换表(10分)

,并画出相应的DFA状态图(10分)

新状态	NFA状态集合	a	b
q_0	$\{s_0\}$		$\{s_1, s_2, s_3, s_5, s_8\}$
q_1	$\{s_1, s_2, s_3, s_5, s_8\}$	$\{s_2, s_3, s_4, s_5, s_7, s_8, s_9\}$	$\{s_2, s_3, s_5, s_6, s_7, s_8\}$
q_2	${s_2, s_3, s_4, s_5, s_7, s_8, s_9}$	${s_2, s_3, s_4, s_5, s_7, s_8, s_9}$	$\{s_2, s_3, s_5, s_6, s_7, s_8, s_{10}\}$
q_3	$\{s_2, s_3, s_5, s_6, s_7, s_8\}$	${s_2, s_3, s_4, s_5, s_7, s_8, s_9}$	$\{s_2, s_3, s_5, s_6, s_7, s_8\}$
q_4	${s_2, s_3, s_5, s_6, s_7, s_8, s_{10}}$	$\{s_2, s_3, s_4, s_5, s_7, s_8, s_9\}$	$\{s_2, s_3, s_5, s_6, s_7, s_8\}$

表 1: 状态转换表



1.3 将得到的DFA图的状态最小化(15分)

 G_1 是接受状态的集合, G_2 是非接受状态的集合,所以

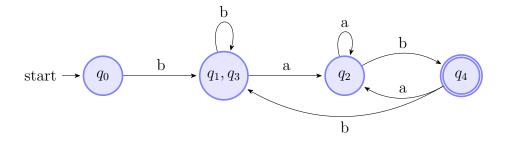
$$G_1 = \{q_4\}, G_2 = \{q_0, q_1, q_2, q_3\}$$

 G_1 无法继续分割,故考察 G_2 可以看出, G_2 中的状态 q_1 和 q_3 可以合并故可以

状态	a	b
q_0		q_1
q_1	q_2	q_3
q_2	q_2	q_4
q_3	q_2	q_3

表 2: 状态转换表 (最小化)

最小化DFA



2 给定文法

$$S \to \text{if } E \text{ then } S \text{ else } S$$

$$S \to \text{if } E \text{ then } S$$

$$S \to \text{other}$$

$$E \to b$$

2.1 计算文法中非终结符的First集(10分)和Follow集(10分)

提取左公因子, 文法改写为

$$S \to \text{if } E \text{ then } S \ S' \mid \text{other}$$

$$S' \to \text{else } S \mid \epsilon$$

$$E \to b$$

计算First集和Follow集

$$First(S) = \{if, other\}$$

 $First(S') = \{else, \epsilon\}$
 $First(E) = \{b\}$
 $Follow(S) = \{else, \$\}$
 $Follow(S') = \{else, \$\}$
 $Follow(E) = \{then\}$

2.2 构造它的预测分析表(共20分,每空4分)

非终结符	if	then	else	other	b	\$
S	$S \to \text{if } E \text{ then } S S'$			$S \to \text{ other}$		
S'			$S' \to \epsilon$			$S' \to \epsilon$
			$S' \to \text{else } S'$			
E					$E \rightarrow b$	

表 3: 预测分析表

2.3 判断其是否为LL(1)文法(共10分,理由5分,结论5分)

文法无左递归, 但是

$$First(S') \cup Follow(S') = \{ \text{else} \}$$

(或答: 预测分析表中出现冲突)

所以文法不是LL(1)文法